CS515 - Project 1: Implementation, design, and results.

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#### 0.1 Implementation

The server submitted here implements a basic event driven architecture. The server is structured in the following way:

```
set up socket
while True do
   select()
   for i = 0 - > max\_file\_descriptors do
       if i set to read then
          if i == server\_socket then
             accept new connection
             create Responder for this connection
          else
             receive on i
             if receive_is_terminating then
                 close(i)
              else
                 send header
                 add Responder to the write queue
             end if
          end if
       end if
       if i set to write then
          read Responders next chunk of data from disk
          send(i)
       end if
   end for
end while
```

The main structure to allow this to work is a class I've called *Responder*. These class will contain all the information the server will require to interact with a client. As clients connect to the server Responders are created accordingly and these Responders will allow the server to switch between connections and send smaller chunks of data to each client. One difficult choice was whether to read in the entire file from the disk in the beginning or to read in chunks at a time as they were needed. It seems that for large files this method will be much more feasible. However, with small files this approach might introduce significant overhead. There are several files involved the server:

#### myhttpd.[ch]:

This is the main server code. Contains code for the algorithm described above.

#### file\_handler.[ch]:

As the name suggests this code handles all interactions with files. This includes checking for any errors in the request related to the file (e.g. file not found). In addition this will check the type of file and get its length for use in the header. Of course the code to read a chunk of a file is here as well. This method will also handle making sure we don't read more of the file than what exists.

#### request\_handler.[ch]:

The main purpose of this class is to create the header. This works by the server calling  $get\_header()$ . From here various checks are made to make sure the request can be fulfilled. If the request can be filled the server will proceed otherwise the reply will consist of the error code, and a page letting the user know what went wrong.

#### Responder.[ch]:

This class is a container to help the server when sending to multiple clients at the same time. Here we track the current offset into the file (so we know where to read from for the next send), the id of the Responder (i.e. the socket number its using), whether or not it's currently in the sending state, a pointer to the file\_handler attached to the file its sending, and a time that corresponds to the last activity from this Responder (used for HTTP/1.1).

The main design decision for 1.1 was to store records in a min-heap, sorted on time value. The heap will contain a pair (time, Responder id). When a Responder takes an action their time is updated. Then at the end of each main loop for the server the heap is examined. The first time to do is check the top value of the heap. If this value hasn't timed out yet no one has timed out. However, if this has timed out an arbitrary number of other elements in the heap could have timed out as well. So, we keep popping until we come across a value that hasn't timed out. One other thing that must happen before a Responder has been considered timed out is, when popping values off the heap, we need to check if this value is the same value recorded inside the Responder. If the Responder has been active recently it will have a newer time, so this must be updated.

The client application is very straightforward. We simply spawn a number of threads and each makes a connection to the server. Each thread then requests a file, receives the response and disconnects.

#### 0.2 Results

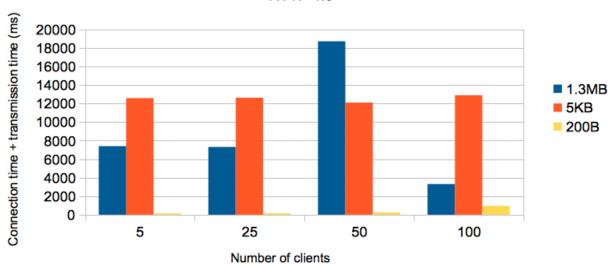
For the experimentation portion of this project I create a custom topology inside of mininet. This topology simply contained 1 switch with all the hosts connected to the switch. All tests were performed with links with a bandwidth of 10Mb/s and a delay of 6ms. All the tests were performed again on links of the same bandwidth but a delay value of 24ms. The initial delay was used because this produced the same RTT I discovered between my home server and a machine on campus. The second value was used to simulate a congested network, operating at 1/4 the link speed of a 'normal' network. The organization of the tests was as follows:

- All tests were performed 5 times
- There were 3 files to be requested from the server of varying sizes (1.3MB, 5KB, and 200B)
- Each test used a client which would run 5 threads
- The number of hosts were incremented from each test (5,25,50,100)

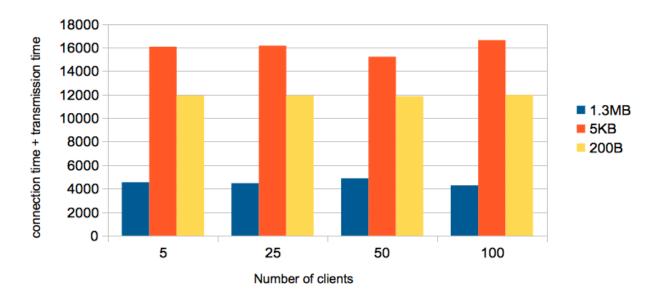
Several charts follow illustrated some of the times discovered in the tests.

### Connection Time + Transmission Time vs File Size

**HTTP 1.0** 

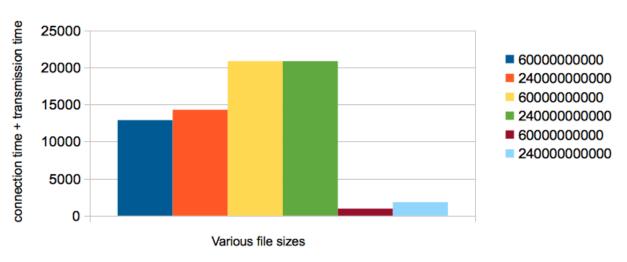


### Connection Time + Transmission Time vs File Size



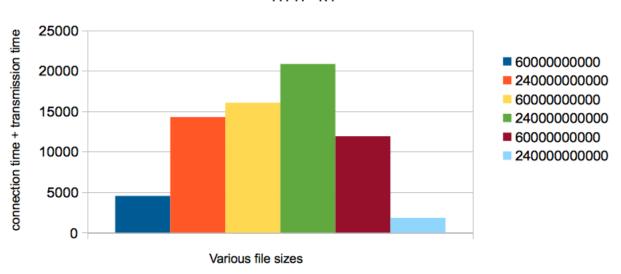
### Transmission+Connection Times vs. Delay Bandwitch Product

**HTTP 1.0** 



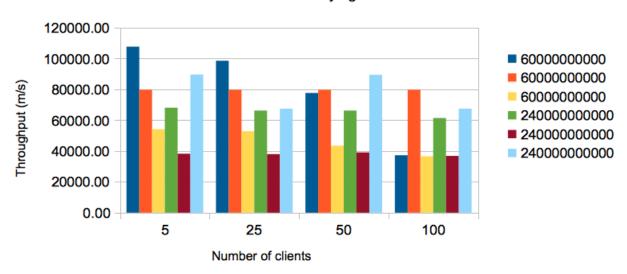
# Transmission and connection time vs. delay-bandwidth product

**HTTP 1.1** 



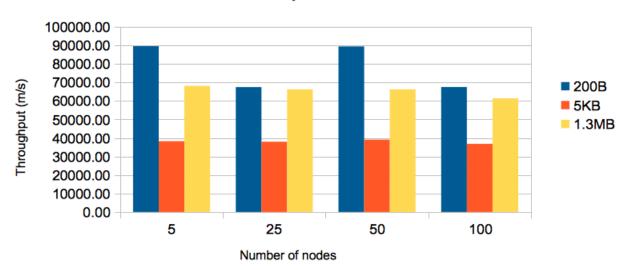
## Delay-Bandwitdth vs. Throughput

For various file sizes and varying amounts of clients



### Throughput vs File Size

### Delay of 24ms



# Throughput vs. File Size

### Delay of 6ms

