

Course: Modern Cryptography

CCA Security

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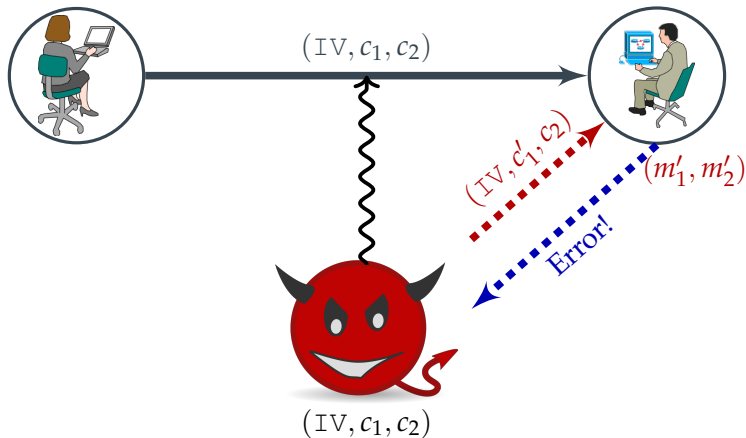
Padding-Oracle Attack¹

m_1	$m_20\times0606\dots06$
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(m_1, m_2)

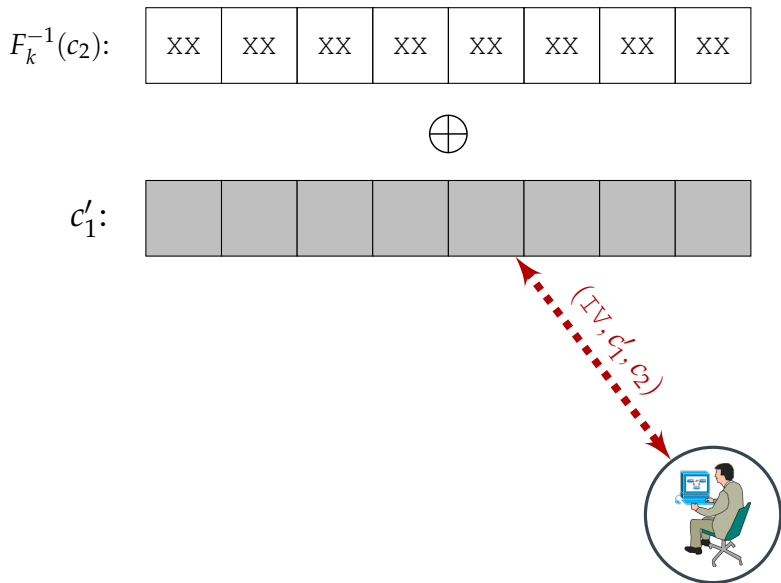
$$m'_2 = F_k^{-1}(c_2) \oplus c'_1$$

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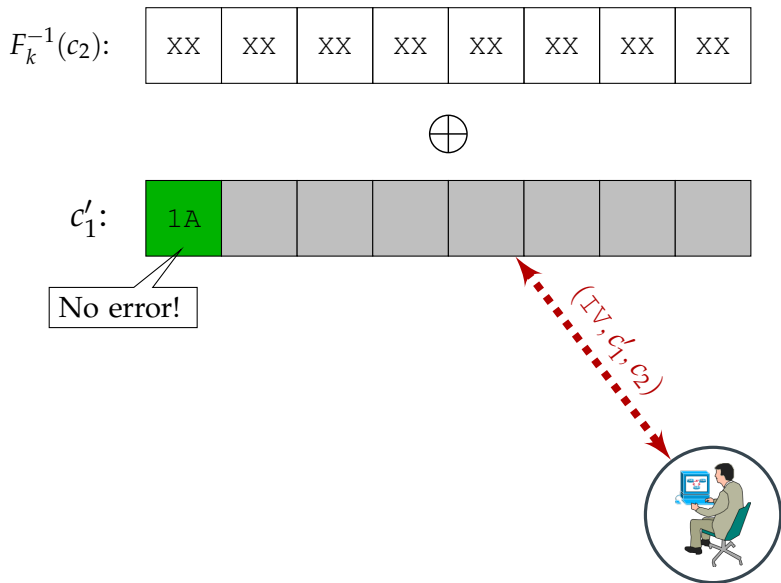


¹Thanks to Prof. Katz for the example.

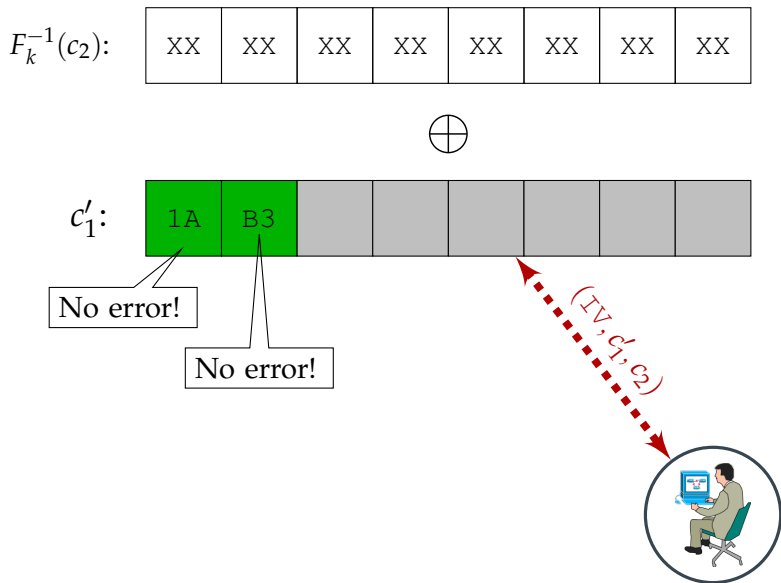
Padding-Oracle Attack..



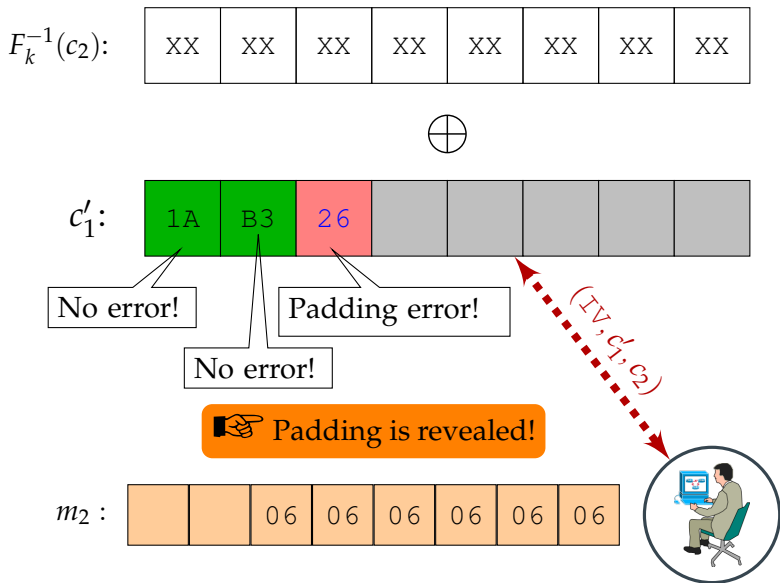
Padding-Oracle Attack..



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Padding-Oracle Attack..

$F_k^{-1}(c_2)$:

		4C	16	6A	1D	3D	8A
		\oplus	\oplus	\oplus	\oplus	\oplus	\oplus
		06	06	06	06	06	06



c_1 :

3A	AB	4C	16	6A	1D	3D	8A
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m_2 :

		06	06	06	06	06	06
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Padding-Oracle Attack..

$F_k^{-1}(c_2):$

		4C	16	6A	1D	3D	8A
		\oplus	\oplus	\oplus	\oplus	\oplus	\oplus
		06	06	06	06	06	06



$c'_1:$

XX	XX	4D	17	6B	1C	3C	8B
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$8A \oplus 06 \oplus 07$



$m'_2:$

		07	07	07	07	07	07
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Padding-Oracle Attack..

$F_k^{-1}(c_2)$:

		4C	16	6A	1D	3D	8A
		\oplus	\oplus	\oplus	\oplus	\oplus	\oplus
		06	06	06	06	06	06



👉 We query the padding oracle for all possibilities for second byte of c'_1 . No error reveals the second byte of $F_k^{-1}(c_2)$ and hence of m_2 .

c'_1 :

XX	XX	4D	17	6B	1C	3C	8B
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👉 Similarly we can obtain the first byte of m_2 as well.

m'_2 :

		07	07	07	07	07	07
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