

Problem 1

$$1) A \rightarrow B, B \rightarrow C, A \rightarrow C$$

1. Transitivity - If $A \rightarrow B$ and $B \rightarrow C$, then $A \rightarrow C$

- Given $X \rightarrow Y$ on a schema R , for any valid instance r conforming to R , and any two records t_1 and t_2 from r , by definition we know that $\pi_A(t_1) = \pi_A(t_2)$, $\pi_B(t_1) = \pi_B(t_2)$ and $\pi_Z(t_1) = \pi_Z(t_2)$.

We also know that $\pi_A(t_1) = \pi_A(t_2)$ implies $\pi_B(t_1) = \pi_B(t_2)$, $\pi_B(t_1) = \pi_B(t_2)$ implies $\pi_C(t_1) = \pi_C(t_2)$.

By the rule of Transitivity we know $\pi_A(t_1) = \pi_A(t_2)$ implies $\pi_C(t_1) = \pi_C(t_2)$

From the above we can concur that: Given $A \rightarrow B$ and $B \rightarrow C$, then $A \rightarrow C$.

2)

$$A \rightarrow B, C \rightarrow D, \text{ then } AC \rightarrow BD$$

With Augmentation we can multiply $A \rightarrow B$ by C and get $AC \rightarrow BC$.

since $C \rightarrow D$ we get $AC \rightarrow BD$.

3)

$$AB \rightarrow C \text{ implies we do not have } A \rightarrow C \text{ NOR } B \rightarrow C$$

By Reflexivity $AB \rightarrow ABC$

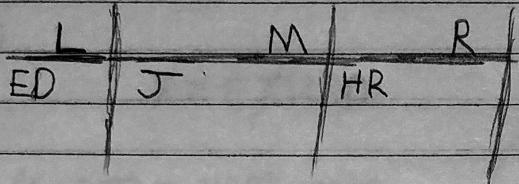
By Augmentation we have $ABC \rightarrow C$

No rules bring us to the implication that $A \rightarrow C$ and $B \rightarrow C$, therefore we do not have $A \rightarrow C$ NOR $B \rightarrow C$.

Problem 2

key

$$W = EHJRD \quad F = \{EJ \rightarrow H, EJD \rightarrow R, E \rightarrow J\}$$



$$E^+ = E \rightarrow J, E \rightarrow EJ, EJ \rightarrow EJH$$

*superkey

$$D^+ = D \rightarrow D$$

$$EJD \rightarrow REJDH$$

$$EJ^+ = EJ \rightarrow EJH$$

$$DJ^+ = DJ \rightarrow DJ$$

$$ED^+ = ED \rightarrow EJHDR$$

*Key

1) False, because given an employee, a job and hours are implied, meaning each employee has one job and a set amount of hours. $E \rightarrow EJH, EJ \rightarrow H$
 $E \rightarrow J$

2) False, as before, an employee implies a job and hours associated with that job. $E \rightarrow J$

3) False, an employee is given a rating each day they work. $EJD \rightarrow R$

4) True, a job never implies an employee, only employees imply jobs.

5) ED is a key and EJD is a superkey

6) Not in BCNF form, not all FD's are keys or superkeys.

7)

$$X = E \cancel{J} H R D$$

$$X = E J \cancel{H} R D$$

$R_1(EH)$
 $R_2(EJ)$
 $R_3(ERD)$

$$(X - EJ) = HRD^+ = HRD$$

$$(X - EH) = JRD^+ = JRD$$

$$(X - ER) = JHD^+ = JHD$$

$$(X - ED) = JHR^+ = JHR$$

$$(X - JH) = ERD^+ = EJ \cancel{H} RD$$

$$(X - HRD) = E J H$$

$$X = ERD$$

$$Y = E J R D$$

$$(Y - EJ) = RD^+ = RD$$

$$(Y - ER) = JD^+ = JD$$

$$(Y - ED) = JR^+ = JR$$

$$(Y - JR) = ED^+ = E J H R D$$

$$(Y - JD) = ER^+ = E J H R$$

$$(Y - RD) = EJ^+ = E J H$$

$$(Y - JRD) = E^+ = E \cancel{J} H$$

* can't do anymore

It is dependency preserving because everything is derivable between the tables due to E.

$E H J \cancel{E} R D$

<u>E J H R D</u>				
1	2	3	4	5
6	7	8	9	10
11	12	13	14	15

<u>E H</u>	<u>E J</u>
1 3	1 2
6 8	6 7
11 13	11 12

<u>E R D</u>	<u>E H J E R D</u>
1 4 5	1 3 2 1 4 5
1 3 2 1 4 5	1 3 2 6 9 10
1 3 2 1 4 5	1 3 2 1 4 5
6 9 10	6 8 7 1 4 5
6 8 7 1 4 5	6 8 7 6 9 10
6 8 7 1 4 5	6 8 7 1 4 5
1 1 1 4 1 5	1 1 1 3 1 2 1 4 1 5

<u>E H J E</u>	<u>E H E J</u>	<u>E J E</u>	<u>E</u>
1 3 1 2	1 1 1 3 1 2	1 1 1 2	1
6 8 1 2	1 1 1 3 6 7	6 8 7	
6 8 6 7	1 1 1 3 1 1 1 2	1 1 1 3 1 2	
6 8 1 1 1 2			

<u>E H J</u>
1 3 2

<u>E H J R D</u>
1 3 2 4 5
6 8 7 9 10
1 1 1 3 1 2 1 4 1 5

Lossless.

8)

$$F = \{JD \rightarrow H\}$$

We only need to worry about Job and Day implying hours, everything else doesn't matter.

Problem 3

1)

$$R \left\{ \begin{array}{l} R_1 = R - Z \\ R_2 = WZ \end{array} \right. \quad W \cap Z = \emptyset$$

$$R_1 \cap R_2 = (R - Z) \cap WZ \\ = W$$

* $R - Z$ leaves only W and doing an intersection
 $W \cap WZ$ leaves only W , making it lossless.

2)

$$W \rightarrow Z = W'A \rightarrow CZ' = W'C \rightarrow Z' \\ = W \rightarrow Z'$$

This gives us $W \cap Z' = \emptyset$

so your new decomposition would be:

$$R_1 = R - Z'$$

$$R_2 = WZ'$$

Problem 4

Hourly_Emps (ssn, name, lg+, rating, wage_per_hr,
hrs_per_wk)

$$T = SNL RWH$$

$$F = \{ S \rightarrow SNL RWH, R \rightarrow W \}$$

1)

Not in BCNF because not all FD's are superkeys or keys

In 3NF because in $R \rightarrow W$; W is part of the key

2)

$R_1(SNL RH)$ $R_2(RW)$

lossless!

In BCNF form

because each
R has an
FD in them.

	S	N	L	R	W	H	R	W
1	2	3	4	5	6		4	5
7	8	9	10	11	12		10	11
13	14	15	16	17	18		16	17

S	N	L	R	H	R	W	S	N	L	R	H	R	W
1	2	3	4	6	4	5	1	2	3	4	6	5	
+ 2	3	4	6	10	11		7	8	9	10	12	11	
- 1	2	3	4	6	16	17	- 13	14	15	16	18	17	
- 7	8	9	10	12	4	5							
- 7	8	9	10	12	10	11							
- 7	8	9	10	12	16	17							
- 13	14	15	16	18	4	5							
- 13	14	15	16	18	10	11							
- 13	14	15	16	18	16	17							

QUESTION

3)

The decomposition of T into R_1 and R_2 is dependency preserving because we can derive W_2 from R_1 to R_2 and everything else is also derivable just from R_1 .

$$HW_1 \cup W_2 = T$$

$$\{W \leftarrow R \mid HW_1 \cup W_2 \subseteq R\} = T$$

(1)

Problem 5

$$R = ABC$$

$$F = \{A \rightarrow B, B \rightarrow C, C \rightarrow A\}$$

1) $R_1(AC)$ $R_2(BC)$

<u>ABC</u>	<u>AC</u>	<u>BC</u>
1 2 3	1 3	2 3
4 5 6	4 6	5 6
7 8 9	7 9	8 9

Not Dependency Preserving because we can't check $A \rightarrow B$ without a join.

<u>ACBC</u>	<u>ACB</u>
1 3 2 3	1 3 2
1 3 5 6	4 6 5
1 3 8 9	7 9 8
4 6 2 3	
4 6 5 6	
4 6 8 9	
7 9 2 3	
7 9 5 6	
7 9 8 9	

lossless!

2) $F = \{A \rightarrow B, C \rightarrow B\}$ $R_1(AC)$ $R_2(BC)$

<u>ABC</u>	<u>AC</u>	<u>BC</u>
1 2 3	1 3	2 3
4 5 6	4 6	5 6
7 8 9	7 9	8 9

Not dependency preserving because we can't check $A \rightarrow B$ without a join.

<u>ACBC</u>
1 3 2 3
1 3 5 6
4 6 1 3
4 6 5 6
4 6 8 9
7 9 1 3
7 9 5 6
7 9 8 9
1 3 8 9

Lossless!