Recap: basic data structures Data Structures and Algorithms for Computational Linguistics III ISCL-BA-07 - Some basis data etrust Çağrı Çöltekin ccoltekin@sfs.uni-tuebingen.de - Lists - Stacks - Oueses Winter Semester 2021-2022 Abstract data types and data structures Arrays An array is simply a contiguous sequence of objects with the same size . An abstract data type (ADT), or abstract data structure, is an object with · Arrays are very close to how computers sto well-defined operations. For example a stack supports push () and pop () data in their memory a[1] operations · Arrays can also be multi-dimensional. For An abstract data structure can be implemented using different data structure can be implemented using a linked list, or an array example, matrices can be represented with 2-dimensional arrays · Sometimes names, usage is confusingly similar · Arrays support fast access to their elements through indexing On the downside, resizing and inserting values in arbitrary locations are expensive a = [3, 6, 8, 9, 3, 0] Arrays Lists Main operations for list ADT are - append - prepend - head (and tail) Lists are typically im ited using linked lists (but array-based lists are a = [3, 6, 8, 9, 3, 0] a[0] # 3 · No built-in array data also common) structure in Python · Python lists are array-based a[-1] . Lists are indevable a[1:4] # [6, 8, 9] For proper/faster arrays, use the numpy library a2d = [[3, 6, 8], [9, 3, 0]] a2d[0][1] # 6 head 3 6 8 9 3 0 head 3 6 8 9 3 0 Stacks Oueues . A queue is a first-in-first (FIFO) out data structure A stack is a last-in-first (LIFO) out data stru Two basic op Two basic operations: - enqueue - dequeu - push - pop Queues can be implemented using linked lists (or maybe arrays) . Stacks can be implemented using linked lists (or arrays) enqueue (3) 0 0 Other common ADT Studying algorithms . In this course we will study a series of important algorithms, including Strings are often implemented based on character arrays
 Maps or dictionaries are similar to arrays and lists, but allow indexing with (almost) arbitrary data types Sorting
Pattern matching
Graph traversal For any algorithm we design/use, there are a number of desirable properties
Correctness an algorithm should dow that it is supposed to do
Robustness an algorithm should downer(by) handle all possible inputs it may receive
Efficiency an algorithm should be simple as possible inputs it may receive
Simplicity an algorithm should be a simple as possible - Maps are generally implemented using hashing (later in this course) Sets implement the mathematical (finite) sets: a collection unique elements without order Trees are used in n nany algorithms we discuss later (we will revisit trees as data structures) * We will briefly touch upon a few of these issues with a simple case study A simple problem: searching a sequence for a value Linear search: take 2 def linear_search(seq, val):
answer = None
for i in range(len(seq)):
if seq[i] == val:
answer = i linear_search(seq, val):
for i in range(len(seq)):
 if seq[i] == val:
 return i return None Can we do even better? Is this a good algorithm? Can we improve it?

Linear search: take 3 Binary search def linear_search(seq, val):
n = len(seq) - 1 : def binary_search(seq, val):
2 left, right = 0, lem(seq)
3 while left = right:
4 mid = (left + right) // 2
5 if seq[mid] = val:
7 return mid
7 if seq[mid] > val:
7 right = mid - 1 n = lem(seq) - i
leat = seq[n]
seq[n] = val
i = 0
while seq[i] != val:
i + i
seq[n] = last
if i < n or seq[n] == val:
return i
else:
return Mone • Is this better? · Any disadvantages? the sequence is sorted. . Can we do even better? else: left = mid + 1 Binary search A note on recursion . Some problems are much easier to solve recursively. idef binary_search_recursive(seq, val, left=None, right=None):
 if left is None: Recursion is also a mathematical concept, properties of recursive algorith are often easier to prove if left is None:

left = 0

if right is None:

right = len(seq)

if left > right:

return None

mid = (left + right) // 2

if seq[mid] == val: • Reminder: - You have to define one or more base cases (e.g., if left > right for binary search)

- Each recursive step should approach the base case (e.g., should run on a smaller portion of the data) ir seq[mid] == val:
 return mid
if seq[mid] > val:
 return binary_search_recursive(seq, val, left, mid - :)
else: \ast We will see quite a few recursive algorithms, it is time for getting used to if you are not return binary_search_recursive(seq, val, mid + 1, right) Exercise: write a recursive function for linear search. Summary An interesting (but not-extremely-relevant) anecdote This lecture is a review of some basic data structure and algorithm How hard can binary search could be? We will assume you know these concepts, please revise your earlier knowledge if needed . It was first suggested in a lecture in 1946 (by John Mauchly) * First fix to this version was suggested in 1960 (by Derrick Henry Lehmer Next: * Another, fix/improvement over this was published on 1962 (by Hermann Analysis of algorithms (Reading: textbook (goodrich2013) chapter 3)
 A few common patterns for designing (efficient) algorithms . In 2006, a bug in Java's binary search implementation was di Acknowledgments, credits, references + Some of the slides are based on the previous year's course by Corina Dima