# NOPT042 Constraint programming: Tutorial 11 - Tabling

## Dynamic programming with tabling

The "t" in Picat stands for "tabling": storing and resusing subcomputations, most typically used in dynamic programming (divide & conquer). We have already seen the following classical example of usefulness of tabling:

#### Example: Fibonacci sequence

```
In [1]: %load_ext ipicat
       Picat version 3.7
In [2]: %picat -n fib
         fib(0, F) => F = 0.
         fib(1, F) => F = 1.
         fib(N, F), N > 1 \Rightarrow fib(N - 1, F1), fib(N - 2, F2), F = F1 + F2.
In [3]: %*picat -n fib_tabled
         table
         fib tabled(0, F) \Rightarrow F = 0.
         fib_tabled(1, F) \Rightarrow F = 1.
         fib tabled(N, F), N > 1 => fib tabled(N - 1, F1), fib tabled(N - 2, F2), F = \frac{1}{2}
         Compare the performance:
In [4]: %picat
        main =>
             time(fib_tabled(42, F)),
             println(F),
             time(fib(42, F)),
             println(F).
       CPU time 0.0 seconds.
       267914296
       CPU time 28.318 seconds.
       267914296
```

Example: shortest path

Find the shortest path from source to target in a weighted digraph. Code from the book:

```
table(+,+,-,min)
            sp(X,Y,Path,W) ?=>
              Path = [(X,Y)],
              edge(X,Y,W).
            sp(X,Y,Path,W) =>
              Path = [(X,Z)|Path1],
              edge(X,Z,Wxz),
              sp(Z,Y,Path1,W1),
              W = W \times z + W1.
        Recall that ?=> means a backtrackable rule. Consider the following simple
        instance:
            index (+,-,-)
            edge(a,b,5).
            edge(b,c,3).
            edge(c,a,9).
            source(a).
            target(c).
In [5]: !cat shortest-path/instance2.pi
        !time picat shortest-path/shortest-path.pi instance2.pi
       edge(1, 2, 1).
       edge(1, 4, 8).
       edge(1, 7, 6).
       edge(2, 4, 2).
       edge(3, 2, 14).
       edge(3, 4, 10).
       edge(3, 5, 6).
       edge(3, 6, 19).
       edge(4, 5, 8).
       edge(4, 8, 13).
       edge(5, 8, 12).
       edge(6, 5, 7).
       edge(7, 4, 5).
       edge(8, 6, 4).
       edge(8, 7, 10).
       source(1).
       target(6).
```

```
path = [(1,2),(2,4),(4,8),(8,6)]
     w = 20
            0m0.011s
      real
      user
            0m0.000s
            0m0.010s
      SYS
In [6]: !cat shortest-path/shortest-path.pi
      Adapted from
       spl.pi
       from Constraint Solving and Planning with Picat, Springer
       by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
      main([Filename]) =>
       cl(Filename),
       source(S),
       target(T),
       sp(S,T,Path,W),
       println(path = Path),
       println(w = W).
      table(+,+,-,min)
      sp(X,Y,Path,W) ?=>
       Path = [(X,Y)],
       edge(X,Y,W).
      sp(X,Y,Path,W) =>
       Path = [(X,Z)|Path1],
       edge(X,Z,Wxz),
       sp(Z,Y,Path1,W1),
       W = W \times z + W1.
```

### Table mode declaration

We can tell Picat what to table using a table mode declaration:

```
table(s1,s2,...,sn)
my_predicate(X1,...,Xn) => ...
```

where si is one of the following:

- + : input, the row/column/etc. where to store
- - : output, the value to store
- min or max : objective, only store outputs with smallest/largest value of this
- nt : not tabled, as if this argument was not passed; last coordinate only, you can use this for global data that do not change in the subproblems, or for

arguments functionally dependent (1-1, easily computable) on the + arguments

For example:

```
table(+,+,-,min)
sp(X,Y,Path,W)
```

means for every X and Y store (only) the  $\operatorname{Path}$  with minimum weight W (only rewrite  $\operatorname{Path}$  if its W is smaller).

#### Index declaration

The *index declaration* index (+,-,-) does not change semantics but facilitates faster lookup when unifying e.g. terms edge(a,X,W), see Wikipedia. The + means that the corresponding coordinate is indexed ("an input"), - means not indexed ("an output"). There can be multiple index patterns, e.g. an undirected graph can be given as:

```
index (+,-) (-,+)
edge(a,b).
edge(a,c).
edge(b,c).
edge(c,b).
```

if we want to traverse the edges in both ways. (This example is from the guide.)

```
In [7]: !cat table-mode-example.pi
     /***********************
       table mode.pi
       from Constraint Solving and Planning with Picat, Springer
       by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
     main ?=>
        p(a,Y),
        println("Y" = Y).
     table(+,max)
     index (-,+)
     p(a,2).
     p(a,1).
     p(a,3).
     p(b,3).
     p(b,4).
In [8]: !picat table-mode-example.pi
     Y = 3
```

## Exercise: shortest shortest path

Modify the above example so that among the minimum-weight paths, only one with minimum *length*, meaning number of edges, is chosen.

```
In [9]: !cat shortest-path/instance.pi
        index (+,-,-)
        edge(a,b,5).
        edge(b,c,3).
        edge(c,a,9).
        source(a).
        target(c).
In [10]: !picat shortest-path/shortest-shortest-path.pi instance.pi
        path = [(a,b),(b,c)]
        w = (8,2)
In [11]: !cat shortest-path/instance3.pi
         !picat shortest-path/shortest-shortest-path.pi instance3.pi
        % this instance is unsatisfiable
        edge(2, 4, 2).
        edge(3, 2, 14).
        edge(3, 4, 10).
        edge(3, 5, 6).
        edge(3, 6, 19).
        edge(4, 5, 8).
        edge(4, 8, 13).
        edge(5, 8, 12).
        edge(6, 5, 7).
        edge(7, 4, 5).
        edge(8, 6, 4).
        edge(8, 7, 10).
        source(1).
        target(6).
        *** error(failed,main/1)
        !cat shortest-path/shortest-shortest-path.pi
In [12]:
```

```
Adapted from
 sp2.pi
 from Constraint Solving and Planning with Picat, Springer
 by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
main([Filename]) =>
 cl(Filename),
 source(S),
 target(T),
 ssp(S,T,Path,W),
 println(path = Path),
 println(w = W).
table(+,+,-,min)
ssp(X,Y,Path,WL) ?=>
 Path = [(X,Y)],
 WL = (Wxy, 1),
 edge(X,Y,Wxy).
ssp(X,Y,Path,WL) =>
 Path = [(X,Z)|Path1],
 edge(X,Z,Wxz),
 ssp(Z,Y,Path1,WL1),
 WL1 = (Wzy, Len1),
 WL = (Wxz+Wzy, Len1+1).
```

% The order in `WL = (Weight, Length)` matters, otherwise we would choose mi nimum-weight path among minimum-edges paths.

#### Exercise: edit distance

Find the (length of the) shortest sequence of edit operations that transform Source string to Target string. There are two types of edit operations allowed:

- insert: insert a single character (at any position)
- delete: delete a single character (at any position)

Once you can compute the distance, try also outputing the sequence of operations.

```
In [13]: # this should output 4
!picat edit-distance/edit.pi saturday sunday

dist = 4
[del(2,a),del(2,t),ins(3,n),del(4,r)]

In [14]: !cat edit-distance/edit.pi
```

```
Adapted from
      edit.pi
      from Constraint Solving and Planning with Picat, Springer
      by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
main([Source, Target]) =>
      edit(Source, Target, Distance, Seq, 1),
      writeln(dist=Distance),
      writeln(Seq).
table(+,+,min)
% base
edit([],[],D,Seq, I) =>
      D=0,
      Seq=[].
% match
edit([X|P],[X|T],D,Seq,I) =>
      edit(P,T,D,Seq,I+1).
% insert
edit(P,[X|T],D,Seq,I) ?=>
      edit(P,T,D1,Seq1,I+1),
      Seq=[sins(I,X)|Seq1],
      D=D1+1.
% delete
edit([X|P],T,D,Seq,I) =>
      edit(P,T,D1,Seq1,I),
      Seq=[\$del(I,X)|Seq1],
      D=D1+1.
```

## Exercise: 01-knapsack

Write a dynamic program for the 01-knapsack problem.

```
Code adapted from
 knapsack.pi
 from Constraint Solving and Planning with Picat, Springer
 by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
main([Filename]) =>
   cl(Filename),
   instance(ItemNames, Capacity, Values, Weights),
   Items = [(ItemNames[I], Values[I], Weights[I]) : I in 1..ItemNames.lengt
h],
   knapsack(Items, Capacity, ChosenItems, TotalValue),
   output(ChosenItems, TotalValue).
table(+,+,-,max)
knapsack([], Capacity, ChosenItems, Value) =>
 ChosenItems = [], Value = 0.
knapsack( , Capacity, ChosenItems, Value), Capacity =< 0 =>
 ChosenItems = [], Value = 0.
knapsack([ | RemainingItems], Capacity, ChosenItems, Value) ?=>
 % Don't take the item
 knapsack(RemainingItems, Capacity, ChosenItems, Value).
knapsack([Item@(ItemName, ItemValue, ItemWeight) | RemainingItems], Capacit
y, ChosenItems, Value), Capacity >= ItemWeight =>
 % Take the item
 ChosenItems = [Item | PrevChosenItems],
 knapsack(RemainingItems, Capacity - ItemWeight, PrevChosenItems, PrevValu
e),
 Value = PrevValue + ItemValue.
output(ChosenItems, TotalValue) =>
   println(total=TotalValue),
   foreach(Item in ChosenItems)
     println(Item)
   end.
```