NOPT042 Constraint programming: Tutorial 9 - Implicit constraints

What was in Lecture 6

Higher levels of consistencies

- arc consistency: vertex extends to an edge
- path consistency: edge extends to a triangle
- **k-consistency**: any consistent assignment of (k-1) different variables can be extended to one additional variable
- **strong k-consistency**: j-konsistency for all $j \le k$
- NC is 1-consistency=strong1-consistency, AC is (strong) 2-consistency, PC is (strong) 3-consistency, NC+AC+PC = strong path consistency
- example: n-vertex graphs, neither n-consistency nor k-consistency for all $n \leq k$ is enough
- **backtrack-free search**: "for some order of variables we can find a value for each variable compatible with the values of already assigned variables" (k backward edges needs (k+1)-consistency)
- graph width w: minimal width (max number of backward edges) among all node orders
- **Theorem**: strongly (w+1)-consistent => there is an order of variables giving a backtrack-free solution
- directional k-consistency: in some order of vars, any consistent assignment of k-1 variables can be extended to any kth variable coming after
- **adaptive consistency**: ensure directional i-consistency where i depends on node width
- [strong] (i,j)-consistency: extend any consistent assignment on i variables to j additional variables
- **inverse consistency** is (1,k)-consistency: look for support of a value in other variables, neighborhood inverse consistency
- **singleton consistency**: assign value to a variable, then test consistency
- nonbinary constraints, generalized arc consistency (GAC), AC-3 adapted for GAC.
- bounds consistency: GAC only for boundary values of the domains (often used in practice)

Exercise: Seesaw

Adam, Boris, and Cecil want to sit on a 10-feet long seesaw such that they are at least 2 feet apart and the seesaw is balanced. Adam weighs 36 lbs, Boris 32 lbs, and Cecil 16 lbs. Write a general model. (You can assume that the length is even, the distance is integer, and that they can only sit at integer points.)

(Problem from Marriott & Stuckey "Programming with Constraints", page 257. Instance from R. Barták's tutorial.)

Possible decision variables?

- Position on the seesaw for each person.
- Distances between persons, position of the first person, and order of persons.
- Person or empty for each position on the seesaw.

Global constraints? Symmetry breaking? Multiple modeling? Search strategies?

```
In [2]: !ls seesaw
        !cat seesaw/instancel.pi
       instance1.pi instance2.pi instance3.pi instance4.pi instance5.pi
       % sample instance
       instance(NumPeople, Length, Distance, Weights) =>
           NumPeople = 3,
           Length = 10,
           Distance = 2,
           Weights = [36, 32, 16].
In [3]: !time picat seesaw/.solution/seesaw1.pi seesaw/instance4.pi
        !time picat seesaw/.solution/seesaw4.pi seesaw/instance4.pi
       [-16, -15, -14, -13, -12, -11, -10, -9, -8, -7, -6, -5, -4, -3, -2, -1, 0, 1, 2, 6, 8, 7, 9, 14, 15,
       10,16,11,12,13]
       real
               0m25.314s
       user
               0m25.260s
               0m0.031s
       SVS
       [-8,0,1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,-7,-6,-5,-4,-3,-9,-10,-13,-12,-
       11,-16,-14]
       real
               0m0.013s
       user
               0m0.004s
               0m0.008s
       SYS
```

Exercise: Golomb's ruler

A Golomb's ruler is an imaginary ruler with n marks such that the distance between every two marks is different. Find the shortest possible ruler for a given

(The solution for N=28 was announced on Nov 23, 2022! The length is 585.)

- What length are you able to solve in reasonable time?
- Add suitable implicit constraints. (We will discuss this in class.)

Redundant (implicit) constraints

Redundant constraints do not restrict the solution set but rather express properties of a solution from a different viewpoint. This can lead to

- · faster domain reduction,
- a significant boost in propagation,
- improved communication between variables.

We have already seen one example last week in the Magic sequence problem: adding the scalar product constraint.

Implicit constraints based on the following:

$$dist[i, j] = dist[i, i + 1] + dist[i + 1, i + 2] + \ldots + dist[j - 1, j]$$

Now estimate distances by 1, sum from i to j:

CPU time 22.974 seconds. Backtracks: 1224484

[0,1,4,13,28,33,47,54,64,70,72]

length = 72

```
\label{eq:foreach} \begin{array}{lll} \text{foreach}(I \text{ in } 1..N\text{-}1, \text{ J in } I\text{+}1..N) \\ & \text{Distances}[I,J\text{-}I] \text{ $\#>$=$ } (J\text{-}I)^*(J\text{-}I\text{+}1) \text{ div } 2, \\ & \text{Distances}[I,J\text{-}I] \text{ $\#<$=$ Length - } (N\text{-}J\text{+}I\text{-}1)^*(N\text{-}J\text{+}I) \text{ div } 2 \\ \text{end} \end{array}
```