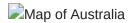
# NOPT042 Constraint programming: Tutorial 3 – Improving your model

It is a good practice to first create a baseline model, and then try to improve. Ways to create more efficient model include:

- gobal constraints: e.g. all different
- **symmetry breaking**: if there is a symmetry in the search space, e.g. in variables or in values, we can fix one element of the orbit and thus limit the part of the space that needs to be explored
- choosing the best solver for your model (or the best model for your solver)
- choosing a good solver configuration (e.g. search strategy---see the next tutorial)

#### Example: Map coloring

Create a model to color the map of Australian states and territories 7 with four colors (cf. The 4-color Theorem). (We exclude the Australian Capital Territory, the Jervis Bay Territory, and the external territories. Map coloring is a special case of graph coloring, see this map.



Let's use the following decision variables:

solve(Territories),
println(Territories).

```
Territories = [WA, NT, SA, Q, NSW, V, T]
In [1]:
        !picat map-coloring/map-coloring-baseline.pi
         [1,2,3,1,2,1,1]
In [2]:
        !cat map-coloring/map-coloring-baseline.pi
        import cp.
        color map(Territories) =>
            % variables
            Territories = [WA, NT, SA, Q, NSW, V, T],
            Territories :: 1..4,
            % constraints
            WA \# != NT,
            WA \# != SA,
            NT #!= SA,
            NT #!= Q,
            SA \#!= Q,
            SA \#! = NSW,
            SA \#!= V,
             Q \#! = NSW,
             V \# != NSW.
        main =>
            color map(Territories),
```

exercise Graph-coloring below.)

```
In [3]: !picat map-coloring/map-coloring-improved.pi
         Western Australia is red.
         Northern Territory is green.
         South Australia is blue.
         Queensland is red.
         New South Wales is green.
         Victoria is red.
         Tasmania is red.
In [4]: !cat map-coloring/map-coloring-improved.pi
         import cp.
         color map(Territories) =>
             % variables
             Territories = [WA, NT, SA, Q, NSW, V, T],
             Territories :: 1..8,
             % constraints
             Edges = [
                 {WA,NT},
                 {WA,SA},
                 {NT,SA},
                 {NT,Q},
                 {SA,Q},
                 {SA,NSW},
                 {SA,V},
                 {Q ,NSW},
                 {V ,NSW}
             ],
             foreach(E in Edges)
                 E[1] #!= E[2]
             end.
         % symmetry breaking constraints
         precolor(Territories) =>
             WA \#=1,
             NT #= 2.
             SA \#= 3.
         % better output than `println(Territories)` (we could also use a map, i.e. a dictionary)
         output(Territories) =>
             Color_names = ["red", "green", "blue", "yellow"],
Territory_names = ["Western Australia", "Northern Territory", "South Australia", "Qu
         eensland", "New South Wales", "Victoria", "Tasmania"],
             foreach(I in 1..Territories.length)
                 writef("%s is %s.\n", Territory names[I], Color names[Territories[I]])
             end.
         main =>
             color map(Territories),
             % precolor(Territories),
             solve(Territories),
             output(Territories).
         What else is wrong with this model? We always want to separate the model from the data. (See the
```

# Choosing a solver

Picat provides the following four solvers (implemented as modules):

- cp
- sat
- smt
- mip

What are the differences?

# Example: Balanced diet (optimization)

This is (one of?) the first optimization problem for which Linear programming was used. Given a list of food items together with their nutritional values and prices, the goal is to choose a balanced diet---one that contains required minimal amounts of nutrients---while minimizing total price.

Note how we pass options to the solver: solve(\$[min(XSum)],Xs) The \$ sign tells the solver to interpret the following as a term, rather than evaluating it as a function.

We will use the mip solver. It requires an external MIP solver. Here we use the Computational Infrastructure for Operations Research (COIN-OR)'s Cbc (branch and cut). (How does branch and cut work? See NOPT059 Large-scale optimization: Exact methods.)

First, we need to install the Cbc package.

```
sudo apt-get install coinor-cbc coinor-libcbc-dev
```

Or without root privileges:

```
cd ~
git clone https://www.github.com/coin-or/coinbrew
cd coinbrew
chmod u+x coinbrew
./coinbrew build Cbc@2.10.8 --no-prompt
export PATH=$PATH:~/coinbrew/dist/bin
```

Welcome to the CBC MILP Solver

Version: 2.10.3

Build Date: Mar 24 2020

command line - cbc \_\_tmp.lp solve solu \_\_tmp.sol (default strategy 1)

Continuous objective value is 90 - 0.01 seconds

Cgl0004I processed model has 4 rows, 4 columns (4 integer (0 of which binary)) and 14 elements

Cutoff increment increased from 1e-05 to 9.9999

Cbc0012I Integer solution of 90 found by DiveCoefficient after 0 iterations and 0 nodes (0.01 seconds)

Cbc0001I Search completed - best objective 90, took 0 iterations and 0 nodes (0.01 seconds)

Cbc0035I Maximum depth 0, 0 variables fixed on reduced cost

Cuts at root node changed objective from 90 to 90

Probing was tried 0 times and created 0 cuts of which 0 were active after adding rounds of cuts (0.000 seconds)

Gomory was tried 0 times and created 0 cuts of which 0 were active after adding rounds o f cuts (0.000 seconds)

Knapsack was tried 0 times and created 0 cuts of which 0 were active after adding rounds of cuts (0.000 seconds)

Clique was tried 0 times and created 0 cuts of which 0 were active after adding rounds o f cuts (0.000 seconds)

MixedIntegerRounding2 was tried 0 times and created 0 cuts of which 0 were active after adding rounds of cuts (0.000 seconds)

FlowCover was tried 0 times and created 0 cuts of which 0 were active after adding round s of cuts (0.000 seconds)

TwoMirCuts was tried 0 times and created 0 cuts of which 0 were active after adding roun ds of cuts (0.000 seconds)

ZeroHalf was tried 0 times and created 0 cuts of which 0 were active after adding rounds of cuts (0.000 seconds)

#### Result - Optimal solution found

Objective value: 90.00000000

Enumerated nodes: 0
Total iterations: 0
Time (CPU seconds): 0.01
Time (Wallclock seconds): 0.02

Total time (CPU seconds): 0.02 (Wallclock seconds): 0.03

{0,3,1,0}

In [6]: !cat balanced-diet/balanced-diet.pi

```
from Constraint Solving and Planning with Picat, Springer
 by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
import mip.
main =>
 data(Prices, Limits, Nutrients),
 Len = length(Prices),
 Xs = new array(Len),
 Xs :: 0..10,
 foreach (I in 1..Nutrients.length)
   scalar product(Nutrients[I], Xs, #>=, Limits[I])
 end,
 scalar product(Prices, Xs, XSum),
 solve($[min(XSum)],Xs),
 writeln(Xs).
% plain scalar product
scalar product(A,Xs,Product) =>
 Product \#= sum([A[I]*Xs[I] : I in 1..A.length]).
scalar product(A,Xs,Rel,Product) =>
 scalar_product(A,Xs,P),
 call(Rel,P,Product).
data(Prices,Limits,Nutrition) =>
 % prices in cents for each product
 Prices = \{50, 20, 30, 80\},\
 % required amount for each nutrition type
 Limits = \{500, 6, 10, 8\},
 % nutrition for each product
 Nutrition =
   {{400,200,150,500}, % calories
    { 3, 2, 0, 0}, % chocolate
    { 2, 2, 4, 4}, % sugar
    { 2, 4, 1, 5}}. % fat
```

# **Exercises**

# Exercise: Coins grid

Place coins on a  $31 \times 31$  board such that each row and each column contain exactly 14 coins, minimize the sum of quadratic horizontal distances of all coins from the main diagonal. (Source: Tony Hurlimann, "A coin puzzle - SVOR-contest 2007")

# Exercise: Sudoku

A traditional constraint satisfaction example: solve an  $n \times n$  sudoku puzzle. Try the following simple instance (from the book:

```
Instance = {
     {4, _, _, _},
     {_, 1, _, _},
     {_, _, _, 1},
```

```
{_, _, _, 2}
}.
```

#### Exercise: Minesweeper

Identify the positions of all mines in a given board. A traditional constraint satisfaction example: solve an  $n \times n$  sudoku puzzle. Try the following instance (from the book:

# Exercise: Magic square

Arrange numbers  $1, 2, \ldots, n^2$  in a square such that every row, every column, and the two main diagonals all sum to the same quantity. How many magic squares are there for a given n? Allow also for a partially filled instance.

# Exercise: Graph-coloring

- 1. Write a program that solves the (directed) graph 3-coloring problem with a given number of colors and a given graph. The graph is given by a list of edges, each edge is a 2-element array. We assume that vertices of the graph are  $1, \ldots, n$  where n is the maximum number appearing in the list.
- 2. Generalize your program to graph k-coloring where k is a positive integer given on the input.
- 3. Modify your program to accept the incidence matrix (a 2D array) instead of the list of edges.
- 4. Add the flag n to output the minimum number of colors (the chromatic number) of a given graph.

For example:

```
picat graph-coloring.pi "[{1,2},{2,3},{3,4},{4,1}]"
picat graph-coloring.pi "[{1,2},{2,3},{3,1}]" 4
picat graph-coloring.pi "{{0,1,1},{1,0,1},{1,1,0}}" 4
picat graph-coloring.pi -n "[{1,2},{2,3},{3,4},{4,1}]"
```

#### Homework: boardomino

In the *boardomino* puzzle, the goal is to cover an  $n \times n$  chess board with some fields missing by domino tiles, if it is possible. The input is given by a positive integer n (the size of the board) and a list of pairs (2-element arrays) of missing fields. If a covering exists, output "yes" and some reasonable representation of it. Else, the output should contain the word "failed" (e.g. the standard message of the cp solver).

Here are some sample instances (the first one is unsatisfiable):

```
picat boardomino.pi 4 "[{1,1},{4,4}]"
picat boardomino.pi 8 "[{1,1},{1,2}]"
picat boardomino.pi 8 "[{1,1},{8,8}]"
```

Try different models and solvers and choose the best option based on the performance on the unsatisfiable instances of the form  $\begin{bmatrix} n & \{1,1\},\{n,n\} \end{bmatrix}$ . (Your program should not be much slower than the best one.)

NOTE: If you use the mip solver, all tests will fail (as the external MIP solver won't be found).