# NOPT042 Constraint programming: Tutorial 11 - Tabling

#### What was in Lecture 8?

#### Combining search and inference

- search (complete, slow) + consistency (incomplete, fast)
- integrate other solving techniques (e.g. integer programming)
- · look-back: maintain consistency among already instantiated vars
- forward-checking: (partial/full) look-ahead, preventing failure better than checking

#### Variable ordering

- important!
- FAIL FIRST (smaller domain first)
- harder first: most constrained variables, more constraints to past variables

#### Value ordering

- SUCEED FIRST (prefer values belonging to a solution)
- prefer value with more supports (from AC-4)
- prefer value leading to less domain reduction (compute singleton consistency)
- problem-driven heuristics are better

#### Branching strategies

- enumeration (X#=0 or X#=1 or ... or X#=N-1)
- step labeling (X#=i or X#!=i)
- bisection/domain-splitting (X#<=i or X#>i)

### Cycle cutset

- acyclic constraint network can be solved by backtrack-free AC
- make it acyclic labeling variables on cycles
- cycle cutset = set of vertices whose removal splits all cycles
- heuristics to find: order vertices by degrees, while G cyclic remove first V
- MAC Extended (MACE): combine AC with cycle cutset

# Dynamic programming with tabling

The "t" in Picat stands for "tabling": storing and resusing subcomputations, most typically used in dynamic programming (divide & conquer). We have already seen the following classical example of usefulness of tabling:

## Example: Fibonacci sequence

```
In [1]: %load_ext ipicat
       Picat version 3.7
In [2]: | %%picat -n fib
        fib(0, F) \Rightarrow F = 0.
        fib(1, F) => F = 1.
        fib(N, F), N > 1 \Rightarrow fib(N - 1, F1), fib(N - 2, F2), F = F1 + F2.
table
        fib_tabled(0, F) \Rightarrow F = 0.
         fib_tabled(1, F) \Rightarrow F = 1.
         fib_{tabled}(N, F), N > 1 \Rightarrow fib_{tabled}(N - 1, F1), fib_{tabled}(N - 2, F2), F = F1 + F
         Compare the performance:
In [4]: %%picat
        main =>
            time(fib_tabled(42, F)),
             println(F),
             time(fib(42, F)),
             println(F).
       CPU time 0.0 seconds.
       267914296
       CPU time 30.419 seconds.
       267914296
```

## Example: shortest path

Find the shortest path from source to target in a weighted digraph. Code from the book:

```
table(+,+,-,min)

sp(X,Y,Path,W) ?=>
  Path = [(X,Y)],
  edge(X,Y,W).

sp(X,Y,Path,W) =>
```

```
Path = [(X,Z)|Path1],
              edge(X,Z,Wxz),
              sp(Z,Y,Path1,W1),
              W = Wxz+W1.
        Recall that ?=> means a backtrackable rule. Consider the following simple
        instance:
            index (+,-,-)
            edge(a,b,5).
            edge(b,c,3).
            edge(c,a,9).
            source(a).
            target(c).
In [5]: !cat shortest-path/instance2.pi
       edge(1, 2, 1).
       edge(1, 4, 8).
       edge(1, 7, 6).
       edge(2, 4, 2).
       edge(3, 2, 14).
       edge(3, 4, 10).
       edge(3, 5, 6).
       edge(3, 6, 19).
       edge(4, 5, 8).
       edge(4, 8, 13).
       edge(5, 8, 12).
       edge(6, 5, 7).
       edge(7, 4, 5).
       edge(8, 6, 4).
       edge(8, 7, 10).
       source(1).
       target(6).
In [6]: !picat shortest-path/shortest-path instance2
       path = [(1,2),(2,4),(4,8),(8,6)]
       w = 20
```

In [7]: !cat shortest-path/shortest-path.pi

```
% Adapted from Constraint Solving and Planning with Picat, Springer
% by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
main([Filename]) =>
 cl(Filename),
  source(S),
  target(T),
  sp(S,T,Path,W),
  println(path = Path),
  println(w = W).
table(+,+,-,min)
sp(X,Y,Path,W) ?=>
 Path = [(X,Y)],
  edge(X,Y,W).
sp(X,Y,Path,W) =>
  Path = [(X,Z)|Path1],
  edge(X,Z,Wxz),
  sp(Z,Y,Path1,W1),
  W = Wxz+W1.
```

## Table mode declaration

We can tell Picat what to table using a table mode declaration:

```
table(s1,s2,...,sn)
my_predicate(X1,...,Xn) => ...
```

where si is one of the following:

- + : input, the row/column/etc. where to store
- - : output, the value to store
- min or max: objective, only store outputs with smallest/largest value of this
- nt: not tabled, as if this argument was not passed; last coordinate only, you
  can use this for global data that do not change in the subproblems, or for
  arguments functionally dependent (1-1, easily computable) on the +
  arguments

For example:

```
table(+,+,-,min)
sp(X,Y,Path,W)
```

means for every X and Y store (only) the  ${\bf Path}$  with minimum weight W (only rewrite  ${\bf Path}$  if its W is smaller).

## Index declaration

The *index declaration* index (+,-,-) does not change semantics but facilitates faster lookup when unifying e.g. terms edge(a,X,W), see Wikipedia. The + means that the corresponding coordinate is indexed ("an input"), - means not indexed ("an output"). There can be multiple index patterns, e.g. an undirected graph can be given as:

```
index (+,-) (-,+)
edge(a,b).
edge(a,c).
edge(b,c).
edge(c,b).
```

if we want to traverse the edges in both ways. (This example is from the guide.)

```
In [8]: !cat table-mode-example.pi

% Adapted from Constraint Solving and Planning with Picat, Springer
% by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
main ?=>
        p(a,Y),
        println("Y" = Y).

table(+,max)
   index (-,+)
   p(a,2).
   p(a,1).
   p(a,3).
   p(b,3).
   p(b,4).

In [9]: !picat table-mode-example
Y = 3
```

# Exercise: shortest shortest path

Modify the above example so that among the minimum-weight paths, only one with minimum *length*, meaning number of edges, is chosen.

```
In [10]: !cat shortest-path/instance.pi
```

```
index (+,-,-)
        edge(a,b,5).
        edge(b,c,3).
        edge(c,a,9).
        source(a).
        target(c).
In [11]: !picat shortest-path/shortest-shortest-path instance
        path = [(a,b),(b,c)]
        W = (8,2)
In [12]: !cat shortest-path/instance3.pi
         !picat shortest-path/shortest-shortest-path instance3
        % this instance is unsatisfiable
        edge(2, 4, 2).
        edge(3, 2, 14).
        edge(3, 4, 10).
        edge(3, 5, 6).
        edge(3, 6, 19).
        edge(4, 5, 8).
        edge(4, 8, 13).
        edge(5, 8, 12).
        edge(6, 5, 7).
        edge(7, 4, 5).
        edge(8, 6, 4).
        edge(8, 7, 10).
        source(1).
        target(6).
        *** error(failed,main/1)
In [13]: !cat shortest-path/shortest-shortest-path.pi
```

```
% Adapted from Constraint Solving and Planning with Picat, Springer
% by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
main([Filename]) =>
 cl(Filename),
  source(S),
  target(T),
  ssp(S,T,Path,W),
  println(path = Path),
  println(w = W).
table(+,+,-,min)
ssp(X,Y,Path,WL) ?=>
  Path = [(X,Y)],
  WL = (Wxy, 1),
  edge(X,Y,Wxy).
ssp(X,Y,Path,WL) =>
  Path = [(X,Z)|Path1],
  edge(X,Z,Wxz),
  ssp(Z,Y,Path1,WL1),
  WL1 = (Wzy, Len1),
  WL = (Wxz+Wzy, Len1+1).
```

% The order in `WL = (Weight, Length)` matters, otherwise we would choose minimum-we
ight path among minimum-edges paths.

## Exercise: edit distance

Find the (length of the) shortest sequence of edit operations that transform

Source string to Target string. There are two types of edit operations allowed:

- insert: insert a single character (at any position)
- delete: delete a single character (at any position)

Once you can compute the distance, try also outputing the sequence of operations.

```
In [14]: # this should output 4
!picat edit-distance/edit.pi saturday sunday

dist = 4
  [del(2,a),del(2,t),ins(3,n),del(4,r)]

In [15]: !cat edit-distance/edit.pi
```

```
% Adapted from Constraint Solving and Planning with Picat, Springer
% by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
main([Source, Target]) =>
        edit(Source, Target, Distance, Seq, 1),
        writeln(dist=Distance),
        writeln(Seq).
table(+,+,min)
% base
edit([],[],D,Seq, I) =>
        D=0,
        Seq=[].
% match
edit([X|P],[X|T],D,Seq,I) =>
        edit(P,T,D,Seq,I+1).
% insert
edit(P,[X|T],D,Seq,I) ?=>
        edit(P,T,D1,Seq1,I+1),
        Seq=[\sin(I,X)|Seq1],
        D=D1+1.
% delete
edit([X|P],T,D,Seq,I) =>
        edit(P,T,D1,Seq1,I),
        Seq=[$del(I,X)|Seq1],
        D=D1+1.
```

## Exercise: 01-knapsack

Write a dynamic program for the 01-knapsack problem.

```
In [16]: !cat knapsack/instance.pi
    instance(ItemNames, Capacity, Values, Weights) =>
        ItemNames = {"tv", "desktop", "laptop", "tablet", "vase", "bottle", "jacket"},
        Capacity = 23,
        Values = {500,350,230,115,180,75,125},
        Weights = {15,11,5,1,7,3,4}.

In [17]: !picat knapsack/knapsack instance
    total = 845
    (tv,500,15)
    (laptop,230,5)
    (tablet,115,1)

In [18]: !cat knapsack/knapsack
```

cat: knapsack/knapsack: No such file or directory