

# NOPT042 Constraint programming: Tutorial 3 – Improving your model

## What was in Lecture 3

Look-back search algorithms (tree search)

- backtracking (what order of variables and values?)
- backjumping (to the source of conflict)
- graph-directed backjumping (driven by constraint network, several jumps in sequence)
- Gaschnig backjumping (considers violated constraints, only one jump, to highest-level)
- conflict-driven: combines both, carry the source of conflict through backjumps
- dynamic backtracking: change order of variables (don't rework easy parts outside of conflict)
- backmarking (avoid repeated constraint checks: remember results of tests)

## Tips & Tricks

### Debugging

- Try very small instances where you understand the solution set.
- Unit-test constraints one by one.
- Start with a model that is as simple as possible.
- For advanced debugging options you can use `debug`, see "How to Use the Debugger" from the Picat Guide.

### How to import a file

Use `cl(instance)` to compile (to bytecode `instance.qi`) & load the file `instance.pi` (anywhere in `$PATH`).

```
main =>
    cl(instance),
    puzzle(Vars),
    solve(Vars).
```

See also Modules ( `import` ).

# Improving the model

It is a good practice to first create a baseline model, and then try to improve. Ways to create more efficient model include:

- **global constraints:** e.g. `all_different`
- **symmetry breaking:** if there is a symmetry in the search space, e.g. in variables or in values, we can fix one element of the orbit and thus limit the part of the space that needs to be explored
- choosing the best solver for your model (or the best model for your solver)
- choosing a good solver configuration (e.g. search strategy---see the next tutorial)

## Example: Map coloring

Create a model to color the map of Australian states and territories 7 with four colors (cf. The 4-color Theorem). (We exclude the Australian Capital Territory, the Jervis Bay Territory, and the external territories. Map coloring is a special case of [graph coloring](#), see [this map](#)).



Let's use the following decision variables:

```
Territories = [WA, NT, SA, Q, NSW, V, T]
```

```
In [1]: !picat map-coloring/map-coloring-baseline
```

```
[1,2,3,1,2,1,1]
```

```
In [2]: !cat map-coloring/map-coloring-baseline.pi
```

```

import cp.

color_map(Territories) =>
    % variables
    Territories = [WA, NT, SA, Q, NSW, V, T],
    Territories :: 1..4,

    % constraints
    WA #!= NT,
    WA #!= SA,
    NT #!= SA,
    NT #!= Q,
    SA #!= Q,
    SA #!= NSW,
    SA #!= V,
    Q  #!= NSW,
    V  #!= NSW.

main =>
    color_map(Territories),
    solve(Territories),
    println(Territories).

```

In [ ]:

How can we improve the model?

In [3]: `!picat map-coloring/map-coloring-improved`

```

Western Australia is red.
Northern Territory is green.
South Australia is blue.
Queensland is red.
New South Wales is green.
Victoria is red.
Tasmania is red.

```

In [4]: `!cat map-coloring/map-coloring-improved.pi`

```

import cp.

color_map(Territories) =>
    % variables
    Territories = [WA, NT, SA, Q, NSW, V, T],
    Territories :: 1..8,

    % constraints
    Edges = [
        {WA,NT},
        {WA,SA},
        {NT,SA},
        {NT,Q},
        {SA,Q},
        {SA,NSW},
        {SA,V},
        {Q ,NSW},
        {V ,NSW}
    ],
    foreach(E in Edges)
        E[1] #!= E[2]
    end.

% symmetry breaking constraints
precolor(Territories) =>
    WA #= 1,
    NT #= 2,
    SA #= 3.

% better output than `println(Territories)` (we could also use a map, i.e. a dictionary)
output(Territories) =>
    Color_names = ["red", "green", "blue", "yellow"],
    Territory_names = ["Western Australia", "Northern Territory", "South Australia",
"Queensland", "New South Wales", "Victoria", "Tasmania"],
    foreach(I in 1..Territories.length)
        writef("%s is %s.\n", Territory_names[I], Color_names[Territories[I]])
    end.

main =>
    color_map(Territories),
    % precolor(Territories),
    solve(Territories),
    output(Territories).

```

What else is wrong with this model? We always want to separate the model from the data. (See the exercise Graph-coloring below.)

## Choosing a solver

Picat provides the following four solvers (implemented as modules):

- cp

- sat
- smt
- mip

What are the differences?

## Example: Balanced diet (optimization)

This is (one of?) the first optimization problem for which Linear programming was used. Given a list of food items together with their nutritional values and prices, the goal is to choose a balanced diet---one that contains required minimal amounts of nutrients---while minimizing total price.

Note how we pass options to the solver: `solve($[min(XSum)],Xs)` The \$ sign tells the solver to interpret the following as a term, rather than evaluating it as a function.

We will use the `mip` solver. It requires an external MIP solver. Here we use the Computational Infrastructure for Operations Research (COIN-OR)'s Cbc (branch and cut).

First, we need to install the Cbc package.

```
sudo apt-get install coinor-cbc coinor-libcbc-dev
```

Or without root privileges:

```
wget https://raw.githubusercontent.com/coin-or/coinbrew/master/coinbrew
chmod u+x coinbrew
./coinbrew fetch Cbc@master
./coinbrew build Cbc
export PATH=$PATH:~/coinbrew/dist/bin
```

In [5]: `!picat balanced-diet/balanced-diet`

```

'Done'
GLPSOL--GLPK LP/MIP Solver 5.0
Parameter(s) specified in the command line:
  --cpxlp -o __tmp.sol __tmp.lp
Reading problem data from '__tmp.lp'...
23 rows, 9 columns, 41 non-zeros
9 integer variables, none of which are binary
47 lines were read
GLPK Integer Optimizer 5.0
23 rows, 9 columns, 41 non-zeros
9 integer variables, none of which are binary
Preprocessing...
5 rows, 9 columns, 23 non-zeros
9 integer variables, none of which are binary
Scaling...
  A: min|aij| = 1.000e+00  max|aij| = 5.000e+02  ratio = 5.000e+02
  GM: min|aij| = 5.946e-01  max|aij| = 1.682e+00  ratio = 2.828e+00
  EQ: min|aij| = 3.536e-01  max|aij| = 1.000e+00  ratio = 2.828e+00
  2N: min|aij| = 2.500e-01  max|aij| = 1.562e+00  ratio = 6.250e+00
Constructing initial basis...
Size of triangular part is 5
Solving LP relaxation...
GLPK Simplex Optimizer 5.0
5 rows, 9 columns, 23 non-zeros
    0: obj = 0.000000000e+00 inf = 1.512e+02 (4)
    4: obj = 1.266666667e+02 inf = 0.000e+00 (0)
*    5: obj = 9.000000000e+01 inf = 0.000e+00 (0)
OPTIMAL LP SOLUTION FOUND
Integer optimization begins...
Long-step dual simplex will be used
+    5: mip = not found yet >= -inf (1; 0)
+    5: >>>> 9.000000000e+01 >= 9.000000000e+01 0.0% (1; 0)
+    5: mip = 9.000000000e+01 >= tree is empty 0.0% (0; 1)
INTEGER OPTIMAL SOLUTION FOUND
Time used: 0.0 secs
Memory used: 0.1 Mb (62497 bytes)
Writing MIP solution to '__tmp.sol'...
{0,3,1,0}

```

In [6]: !cat balanced-diet/balanced-diet.pi

```
% adapted from Constraint Solving and Planning with Picat, Springer  
% by Neng-Fa Zhou, Hakan Kjellerstrand, and Jonathan Fruhman
```

```
import mip.
```

```
main =>  
  data(Prices,Limits,Nutrients),  
  Len = length(Prices),  
  Xs = new_array(Len),  
  Xs :: 0..10,  
  foreach (I in 1..Nutrients.length)  
    scalar_product(Nutrients[I],Xs,#>=,Limits[I])  
  end,  
  scalar_product(Prices,Xs,XSum),  
  solve($[min(XSum)],Xs),  
  writeln(Xs).
```

```
% plain scalar product
```

```
scalar_product(A,Xs,Product) =>  
  Product #= sum([A[I]*Xs[I] : I in 1..A.length]).
```

```
scalar_product(A,Xs,Rel,Product) =>  
  scalar_product(A,Xs,P),  
  call(Rel,P,Product).
```

```
data(Prices,Limits,Nutrition) =>  
  % prices in cents for each product  
  Prices = {50,20,30,80},  
  % required amount for each nutrition type  
  Limits = {500,6,10,8},
```

```
% nutrition for each product
```

```
Nutrition =  
  {{400,200,150,500}, % calories  
   { 3,  2,  0,  0}, % chocolate  
   { 2,  2,  4,  4}, % sugar  
   { 2,  4,  1,  5}}. % fat
```

```
In [7]: !cat __tmp.lp
```



```

Minimize
  obj: X0
Subject To
  -X0 >= -1800
  X0 >= 0
  -X1 >= -10
  X1 >= 0
  -2 X1 - 2 X2 - 4 X3 - 4 X4 + X6 = 0
  -X2 >= -10
  X2 >= 0
  -2 X2 - 3 X1 + X7 = 0
  -20 X2 - 30 X3 - 50 X1 - 80 X4 + X0 = 0
  -X3 >= -10
  X3 >= 0
  -150 X3 - 200 X2 - 400 X1 - 500 X4 + X8 = 0
  -X3 - 2 X1 - 4 X2 - 5 X4 + X5 = 0
  -X4 >= -10
  X4 >= 0
  -X5 >= -120
  X5 >= 8
  -X6 >= -120
  X6 >= 10
  -X7 >= -50
  X7 >= 6
  -X8 >= -12500
  X8 >= 500
Bounds
  0 <= X0 <= 1800
  0 <= X1 <= 10
  0 <= X2 <= 10
  0 <= X3 <= 10
  0 <= X4 <= 10
  8 <= X5 <= 120
  10 <= X6 <= 120
  6 <= X7 <= 50
  500 <= X8 <= 12500
Integers
  X0
  X1
  X2
  X3
  X4
  X5
  X6
  X7
  X8
End

```

In [8]: !cat \_\_tmp.sol

Problem:  
 Rows: 23  
 Columns: 9 (9 integer, 0 binary)  
 Non-zeros: 41  
 Status: INTEGER OPTIMAL  
 Objective: obj = 90 (MINimum)

No.	Row name	Activity	Lower bound	Upper bound
1	r.4	-90	-1800	
2	r.5	90	0	
3	r.6	0	-10	
4	r.7	0	0	
5	r.8	0	0	=
6	r.9	-3	-10	
7	r.10	3	0	
8	r.11	0	0	=
9	r.12	0	0	=
10	r.13	-1	-10	
11	r.14	1	0	
12	r.15	0	0	=
13	r.16	0	0	=
14	r.17	0	-10	
15	r.18	0	0	
16	r.19	-13	-120	
17	r.20	13	8	
18	r.21	-10	-120	
19	r.22	10	10	
20	r.23	-6	-50	
21	r.24	6	6	
22	r.25	-750	-12500	
23	r.26	750	500	

No.	Column name	Activity	Lower bound	Upper bound
1	X0 *	90	0	1800
2	X1 *	0	0	10
3	X2 *	3	0	10
4	X3 *	1	0	10
5	X4 *	0	0	10
6	X6 *	10	10	120
7	X7 *	6	6	50
8	X8 *	750	500	12500
9	X5 *	13	8	120

Integer feasibility conditions:

KKT.PE: max.abs.err = 0.00e+00 on row 0  
 max.rel.err = 0.00e+00 on row 0  
 High quality

KKT.PB: max.abs.err = 0.00e+00 on row 0  
 max.rel.err = 0.00e+00 on row 0  
 High quality

End of output

# Exercises

## Exercise: Coins grid

Place coins on a  $31 \times 31$  board such that each row and each column contain exactly 14 coins, minimize the sum of quadratic horizontal distances of all coins from the main diagonal. (Source: Tony Hurlimann, "A coin puzzle - SVOR-contest 2007")

## Exercise: Sudoku

A traditional constraint satisfaction example: solve an  $n \times n$  sudoku puzzle. Try the following simple instance (from [the book](#)):

```
Board = {  
    {4, _, _, _},  
    {_, 1, _, _},  
    {_, _, _, 1},  
    {_, _, _, 2}  
}.
```

## Exercise: Graph-coloring

1. Write a program that solves the (directed) graph 3-coloring problem with a given number of colors and a given graph. The graph is given by a list of edges, each edge is a 2-element array. We assume that vertices of the graph are  $1, \dots, n$  where  $n$  is the maximum number appearing in the list.
2. Generalize your program to graph  $k$ -coloring where  $k$  is a positive integer given on the input.
3. Modify your program to accept the incidence matrix (a 2D array) instead of the list of edges.
4. Add the flag `-n` to output the minimum number of colors (the chromatic number) of a given graph. For example:

```
picat graph-coloring "[{1,2},{2,3},{3,4},{4,1}]"  
picat graph-coloring "[{1,2},{2,3},{3,1}]" 4  
picat graph-coloring "{{0,1,1},{1,0,1},{1,1,0}}" 4  
picat graph-coloring -n "[{1,2},{2,3},{3,4},{4,1}]"
```