Apache Spark Internals

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Acknowledgments & Sources

- Research papers: https://spark.apache.org/research.html
- Presentations:
- M. Zaharia, "Introduction to Spark Internals" https://www.youtube.com/watch?v=49Hr5xZyTEA
- A. Davidson, "A Deeper Understanding of Spark Internals", https://www.youtube.com/watch?v=dmL0N3qfSc8

Introduction and Motivations

What is Apache Spark

Project goals

- Generality: diverse workloads, operators, job sizes
- Low latency: sub-second
- Fault tolerance: faults are the norm, not the exception
- Simplicity: often comes from generality

Motivations

Software engineering point of view

- Hadoop code base is huge
- Contributions/Extensions to Hadoop are cumbersome
- Java-only hinders wide adoption, but Java support is fundamental

System/Framework point of view

- Unified pipeline
- Simplified data flow
- Faster processing speed

Data abstraction point of view

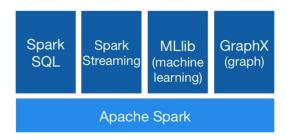
- New fundamental abstraction RDD
- Easy to extend with new operators
- More descriptive computing model

Hadoop: No Unified Vision

General Batching	Specialized systems			
	Streaming	Iterative	Ad-hoc / SQL	Graph
MapReduce	Storm	Mahout	Pig	Giraph
	S4		Hive	
	Samza		Drill	
			Impala	

- Sparse modules
- Diversity of APIs
- Higher operational costs

SPARK: A Unified Pipeline

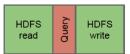


- Spark Streaming (stream processing)
- GraphX (graph processing)
- MLLib (machine learning library)
- Spark SQL (SQL on Spark)

A Simplified Data Flow



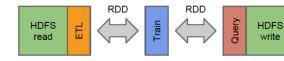


















Hadoop: Bloated Computing Model

```
public class WordCount {
 public static class Map extends Mapper < LongWritable, Text, Text, IntWritable> {
    private final static IntWritable one = new IntWritable(1);
   private Text word = new Text();
    public void map(LongWritable key, Text value, Context context) throws IOException, InterruptedException
        String line = value.toString();
        StringTokenizer tokenizer = new StringTokenizer(line);
        while (tokenizer.hasMoreTokens()) {
            word.set(tokenizer.nextToken());
           context.write(word, one);
 public static class Reduce extends Reducer<Text, IntWritable, Text, IntWritable> {
   public void reduce(Text key, Iterable<IntWritable> values, Context context)
      throws IOException, InterruptedException {
        int sum = 0;
        for (IntWritable val : values) {
            sum += val.get();
        context.write(key, new IntWritable(sum));
 public static void main(String[] args) throws Exception {
    Configuration conf = new Configuration();
        Job job = new Job(conf, "wordcount");
    job.setOutputKeyClass(Text.class);
    job.setOutputValueClass(IntWritable.class);
    job.setMapperClass(Map.class);
    iob.setReducerClass(Reduce.class):
    job.setInputFormatClass(TextInputFormat.class);
    job.setOutputFormatClass(TextOutputFormat.class);
    FileInputFormat.addInputPath(job, new Path(args[0]));
    FileOutputFormat.setOutputPath(job, new Path(args[1]));
    job.waitForCompletion(true);
```

SPARK: Descriptive Computing Model

```
val file = sc.textFile("hdfs://...")

val counts = file.flatMap(line => line.split(" "))
    .map(word => (word,1))
    .reduceByKey(_ + _)

counts.saveAsTextFile("hdfs://...")
```

- Organize computation into multiple stages in a processing pipeline
 - ► **Transformations** apply user code to distributed data in parallel
 - Actions assemble final output of an algorithm, from distributed data

Faster Processing Speed

Let's focus on iterative algorithms

- Spark is faster thanks to the simplified data flow
- We avoid materializing data on HDFS after each iteration

Example: k-means algorithm, 1 iteration

- ▶ HDFS Read
- Map(Assign sample to closest centroid)
- GroupBy(Centroid ID)
- ► NETWORK Shuffle
- Reduce(Compute new centroids)
- ► HDFS Write

Code Base (2012)

Spark core: 16,000 LOC Scheduler: 2500 Operators: 2000 Interpreter: 3300 LOC Block manager: 2700 Networking: 1200 Accumulators: 200 Broadcast: 3500 Mesos backend: Standalone backend: Hadoop I/O: 400 LOC **700 LOC** 1700 LOC

- 2012 (version 0.6.x): 20,000 lines of code
- 2014 (branch-1.0): 50,000 lines of code

Anatomy of a Spark Application

A Very Simple Application Example

```
val sc = new SparkContext("spark://...", "MyJob", home,
    jars)

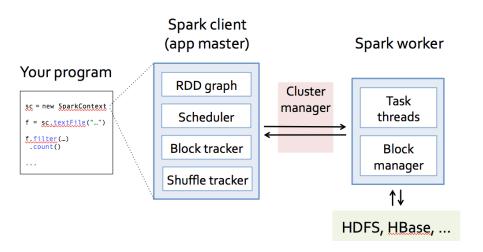
val file = sc.textFile("hdfs://...") // This is an RDD

val errors = file.filter(_.contains("ERROR")) // This is
    an RDD

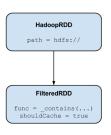
errors.cache()
errors.count() // This is an action
```

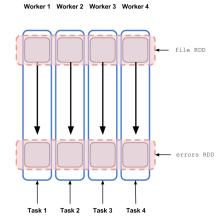
Apache Spark adopts the lazy evaluation model

Spark Components: details



The RDD graph: dataset vs. partition views





Data Locality

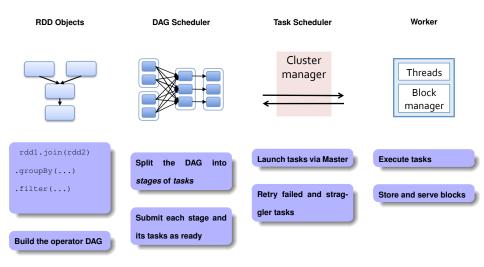
Data locality principle

- Same as for Hadoop MapReduce
- Avoid network I/O, workers should manage local data

Data locality and caching

- First run: data not in cache, so use HadoopRDD's locality prefs (from HDFS)
- Second run: FilteredRDD is in cache, so use its locations
- If something falls out of cache, go back to HDFS

Lifetime of a Job in Spark



In Summary

Our example Application: a jar file

- Creates a SparkContext, which is the core component of the driver
- Creates an input RDD, from a file in HDFS
- Manipulates the input RDD by applying a filter (f: T => Boolean) transformation
- Invokes the action count () on the transformed RDD

The DAG Scheduler

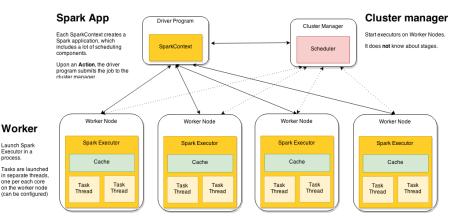
- Gets: RDDs, functions to run on each partition and a listener for results
- Builds Stages of Tasks objects (code + preferred location)
- Submits Tasks to the Task Scheduler as ready
- Resubmits failed Stages

The Task Scheduler

- Launches Tasks on executors
- Relaunches failed Tasks
- Reports to the DAG Scheduler

Spark Deployments

Spark Components: System-level View



Spark Deployment Modes

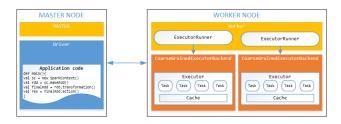
The Spark Framework can adopt several cluster managers

- Local Mode
- Standalone mode
- Apache Mesos
- Hadoop YARN

General "workflow"

- Spark application creates SparkContext, which initializes the DriverProgram
- Registers to the ClusterManager
- Ask resources to allocate Executors
- Schedule Task execution

Worker Nodes and Executors



Worker nodes are machines that run executors

- ▶ Host one or multiple Workers
- ▶ One JVM (= 1 UNIX process) per Worker
- ▶ Each Worker can spawn one or more Executors

Executors run tasks

- Run in child JVM (= 1 UNIX process)
- ▶ Execute one or more task using threads in a ThreadPool

Comparison to Hadoop MapReduce

Hadoop MapReduce

- One Task per UNIX process (JVM), more if JVM reuse
- MultiThreadedMapper, advanced feature to have threads in Map Tasks
- → Short-lived Executor, with one large Task

Spark

- Tasks run in one or more Threads, within a single UNIX process (JVM)
- Executor process statically allocated to worker, even with no threads
- → Long-lived Executor, with many small Tasks

Benefits of the Spark Architecture

Isolation

- Applications are completely isolated
- Task scheduling per application

Low-overhead

- Task setup cost is that of spawning a thread, not a process
- 10-100 times faster
- ► Small tasks → mitigate effects of data skew

Sharing data

- Applications cannot share data in memory natively
- Use an external storage service like Tachyon

Resource allocation

- Static process provisioning for executors, even without active tasks
- Dynamic provisioning under development

Resilient Distributed Datasets

M. Zaharia, M. Chowdhury, T. Das, A. Dave, J. Ma, M. McCauley, M.J. Franklin, S. Shenker, I. Stoica.

Resilient Distributed Datasets: A Fault-Tolerant Abstraction for In-Memory Cluster Computing,

USENIX Symposium on Networked Systems Design and Implementation, 2012

What is an RDD

- RDD are partitioned, locality aware, distributed collections
 - ▶ RDD are immutable

- RDD are data structures that:
 - Either point to a direct data source (e.g. HDFS)
 - Apply some transformations to its parent RDD(s) to generate new data elements

- Computations on RDDs
 - Represented by lazily evaluated lineage DAGs composed by chained RDDs

RDD Abstraction

Overall objective

- Support a wide array of operators (more than just Map and Reduce)
- Allow arbitrary composition of such operators

Simplify scheduling

Avoid to modify the scheduler for each operator

→ The question is: How to capture dependencies in a general way?

RDD Interfaces

- Set of partitions ("splits")
 - Much like in Hadoop MapReduce, each RDD is associated to (input) partitions
- List of dependencies on parent RDDs
 - This is completely new w.r.t. Hadoop MapReduce
- Function to compute a partition given parents
 - ► This is actually the "user-defined code" we referred to when discussing about the Mapper and Reducer classes in Hadoop
- Optional preferred locations
 - ► This is to enforce data locality
- Optional partitioning info (Partitioner)
 - This really helps in some "advanced" scenarios in which you want to pay attention to the behavior of the shuffle mechanism

Hadoop RDD

- partitions = one per HDFS block
- dependencies = none
- compute(partition) = read corresponding block
- preferredLocations(part) = HDFS block location
- partitioner = none

Filtered RDD

- partitions = same as parent RDD
- dependencies = one-to-one on parent
- compute(partition) = compute parent and filter it
- preferredLocations(part) = none (ask parent)
- partitioner = none

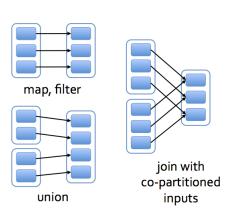
Joined RDD

- partitions = one per reduce task
- dependencies = shuffle on each parent
- compute(partition) = read and join shuffled data
- preferredLocations(part) = none
- partitioner = HashPartitioner(numTask)¹

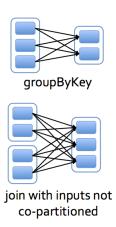
¹Spark knows this data is hashed.

Dependency Types (1)

Narrow dependencies



Wide dependencies



Dependency Types (2)

Narrow dependencies

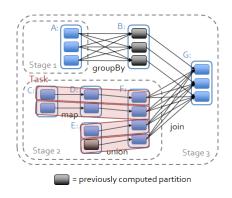
- Each partition of the parent RDD is used by at most one partition of the child RDD
- Task can be executed locally and we don't have to shuffle. (Eg: map, flatMap, filter, sample)

Wide Dependencies

- Multiple child partitions may depend on one partition of the parent RDD
- ► This means we have to shuffle data unless the parents are hash-partitioned (Eg: sortByKey, reduceByKey, groupByKey, cogroupByKey, join, cartesian)

Dependency Types: Optimizations

- Benefits of Lazy evaluation: The DAG Scheduler optimizes Stages and Tasks before submitting them to the Task Scheduler
 - Examples:
 - ★ Piplining narrow dependencies within a Stage
 - ★ Join plan selection based on partitioning
 - * Cache reuse



Operations on RDDs: Transformations

Transformations

- Set of operations on a RDD that define how they should be transformed
- As in relational algebra, the application of a transformation to an RDD yields a new RDD (because RDDs are immutable)
- Transformations are lazily evaluated, which allow for optimizations to take place before execution

• Examples (not exhaustive)

- map(func), flatMap(func), filter(func)
- grouByKey()
- reduceByKey(func), mapValues(func), distinct(), sortByKey(func)
- ▶ join (other), union (other)
- sample()

Operations on RDDs: Actions

Actions

- Apply transformation chains on RDDs, eventually performing some additional operations (e.g., counting)
- Some actions only store data to an external data source (e.g. HDFS), others fetch data from the RDD (and its transformation chain) upon which the action is applied, and convey it to the driver

Examples (not exhaustive)

- reduce(func)
- collect(), first(), take(), foreach(func)
- count(), countByKey()
- saveAsTextFile()

Operations on RDDs: Final Notes

Look at return types!

- Return type: RDD → transformation
- ► Return type: built-in scala/java types such as int, long, List<Object>, Array<Object> → action

Caching is a transformation

Hints to keep RDD in memory after its first evaluation

Transformations depend on RDD "flavor"

- ▶ PairRDD
- ▶ SchemaRDD

Detailed Example: Word Count

Spark Word Count: the driver

Driver and SparkContext

- A SparkContext initializes the application driver, the latter then registers the application to the cluster manager, and gets a list of executors
- ▶ Then, the driver takes full control of the Spark job

Spark Word Count: the code

```
val lines = sc.textFile("input")
val words = lines.flatMap(_.split(" "))
val ones = words.map(w => (w, 1))
val counts = ones.reduceByKey(_ + _)
val result = counts.collectAsMap()
```

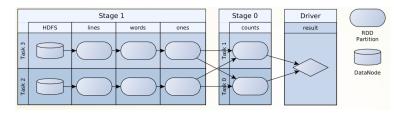
RDD lineage DAG is built on driver side with

- Data source RDD(s)
- Transformation RDD(s), which are created by transformations

Job submission

An action triggers the DAG scheduler to submit a job

Spark Word Count: the DAG



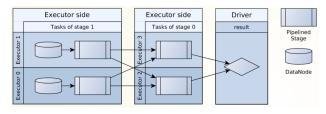
Directed Acyclic Graph

Built from the RDD lineage

DAG scheduler

- Transforms the DAG into stages and turns each partition of a stage into a single task
- Decides what to run

Spark Word Count: the execution plan



Spark Tasks

- Serialized RDD lineage DAG + closures of transformations
- Run by Spark executors

Task scheduling

- The driver side task scheduler launches tasks on executors according to resource and locality constraints
- ▶ The task scheduler decides where to run tasks

Spark Word Count: the Shuffle phase

```
val lines = sc.textFile("input")
val words = lines.flatMap(_.split(" "))
val ones = words.map(w => (w, 1))
val counts = ones.reduceByKey(_ + _)
val result = counts.collectAsMap()
```

reduceByKey transformation

- Induces the shuffle phase
- ► In particular, we have a wide dependency
- Like in Hadoop MapReduce, intermediate <key,value> pairs are stored on the local file system

• Automatic combiners!

 The reduceByKey transformation implements map-side combiners to pre-aggregate data

Advanced Topics

Caching and Storage

Spark's Storage Module

The storage module

- Access (I/O) "external" data sources: HDFS, Local Disk, RAM, remote data access through the network
- Caches RDDs using a variety of "storage levels"

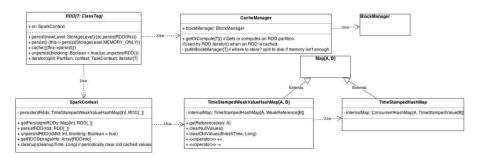
Main components

- The Cache Manager: uses the Block Manager to perform caching
- ► The Block Manager: distributed key/value store

Class Diagram of the Caching Component

RDD Caching flow in Spark-core

Class diagram



How Caching Works

Frequently used RDD can be stored in memory

- Deciding which RDD to cache is an art!
- ▶ One method, one short-cut: persist(), cache()

SparkContext keeps track of cached RDD

- Uses a data-structed called persistentRDD
- Maintains references to cached RDD, and eventually call the garbage collector
- ► Time-stamp based invalidation using
 TimeStampedWeakValueHashMap[A, B]

How Caching Works



```
- rdd.iterator(split:Partition)
-- cacheManager.getOrCompute(rdd, partition)
--- val key = RDDBlockId(rdd.id, partition.index)
--- switch(blockManager.get(key))
        case Some: return
        case None:
                val computedValue = rdd.computeOrReadCheckPoint(partition)
                if isRunningLocally
                        return computedValue
                else
                        cachedValue = putInBlockManager(key, computedValue, storageLevel)
 BlockManager, scala
 // write-once key-value
 // serves both cachedRdds and shuffle data
 // tracks storage level of each block
 // can drop data on disk if RAM mem is low
 // can replicate data across nodes
 def get(blockId){
         trvGetLocal()
         tryGetRemote()
         if no, return None
```

The Block Manager

"Write-once" key-value store

- One node per worker
- No updates, data is immutable

Main tasks

- Serves shuffle data (local or remote connections) and cached RDDs
- Tracks the "Storage Level" (RAM, disk) for each block
- Spills data to disk if memory is insufficient
- Handles data replication, if required

Storage Levels

The Block Manager can hold data in various storage tiers

- org.apache.spark.storage.StorageLevel contains flags to indicate which tier to use
- Manual configuration, in the application
- Deciding the storage level to use for RDDs is not trivial

Available storage tiers

- RAM (default option): if the the RDD doesn't fit in memory, some partitions will not be cached (will be re-computed when needed)
- Tachyon (off java heap): reduces garbage collection overhead, the crash of an executor no longer leads to cached data loss
- Disk

Data format

- Serialized or as Java objects
- Replicated partitions

Resource Allocation: Spark Schedulers

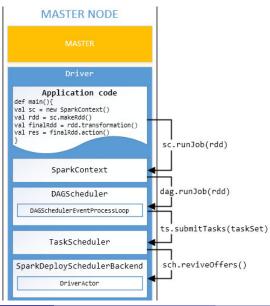
Spark Schedulers

- Two main scheduler components, executed by the driver
 - The DAG scheduler
 - The Task scheduler

Objectives

- Gain a broad understanding of how Spark submits Applications
- Understand how Stages and Tasks are built, and their optimization
- Understand interaction among various other Spark components

Submitting a Spark Application: A Walk Through



Submitting a Spark Application: Details

```
TaskSchedulerimpl.scala
1 - finalRdd.action()
2 -- sc.runJob()
                                                                          27 def submitTasks(taskSet){
3 --- dagScheduler.runJob()
                                                                                  schedulableBuilder.addTaskSetManager(new TaskSetManager(taskSet))
                                                                                  if (!runInLocal && !hasReceivedTask){
4 ---- dagScheduler.submitJob()
5 ---- DAGSchedulerEventProcessLoop.post(JobSubmitted)
                                                                          30
                                                                                          starvationTimer.scheduleAtFixedRate(new TimerTask(){..., STARVA
                                                                          31
                                                                          32
                                                                                  schedulerBackend.reviveOffers()
                                                                          33 }
6 - DAGSchedulerEventProcessLoop.onRecieve(JobSubmitted)
                                                                         34 def sparkDeplovSchedulerBackend.reviveOffers(){
7 -- dagScheduler.handleJobSubmitted()
                                                                                  driverActor | ReviveOffers
8 --- finalStage = new Stage()
                                                                         36 }
9 --- submitStage(finalStage)
10 ---- missingStages = getMissingParentStages(stage).sortById()
11 ---- if missingStages != Nil
                                                                          CoarseGrainedSchedulerBackend.scala
        foreach parent <- missingStages: submitStage(parent)
                                                                          37 def receiveWithLogging(){
13 ---- else
                                                                                  case ReviveOffers => makeOffers(){
        dagScheduler.submitMissingTasks(stage)
                                                                          39
                                                                                          launchTasks(tasks){
                                                                          40
                                                                                                   foreach task <- tasks:
                                                                          41
                                                                                                           executorActor(task.executorId) ! LaunchTask(
        15 def getMissingParentStages(stage){
                                                                          42
        16
                case ShuffleDependency: missing += new Stage()
                                                                          43
        17
                 case NarrowDependency: waiting4Visit.push()
                                                                          44 }
         18 }
         19 def submitMissingTasks(stage){
                 val loc = getPreferredLocs(stage.rdd)
         20
         21
                 if stage.isShuffleMap
         22
                          tasks:Seq[Task] = partitions.foreach(yield new ShuffleMapTask(p, loc))
         23
                 else
         24
                          tasks:Seq[Task] = partitions.foreach(yield new ResultTask())
         25
         26
                 taskScheduler.submitTasks(new TaskSet(tasks))
         27 }
```

The DAG Scheduler

Stage-oriented scheduling

- ► Computes a DAG of stages for each job in the application Lines 10-14, details in Lines 15-27
- Keeps track of which RDD and stage output are materialized
- Determines an optimal schedule, minimizing stages
- Submit stages as sets of Tasks (TaskSets) to the Task scheduler Line 26

Data locality principle

 Uses "preferred location" information (optionally) attached to each RDD

Line 20

 Package this information into Tasks and send it to the Task scheduler

Manages Stage failures

- Failure type: (intermediate) data loss of shuffle output files
- Failed stages will be resubmitted
- NOTE: Task failures are handled by the Task scheduler, which simply resubmit them if they can be computed with no dependency on previous output

The DAG Scheduler: Implementation Details

Implemented as an event queue

- Uses a daemon thread to handle various kinds of events
 Line 6
- ▶ JobSubmitted, JobCancelled, CompletionEvent
- The thread "swipes" the queue, and routes event to the corresponding handlers

• What happens when a job is submitted to the DAGScheduler?

- ▶ JobWaiter object is created
- JobSubmitted event is fired
- ► The daemon thread blocks and wait for a job result Lines 3, 4

The DAG Scheduler: Implementation Details (2)

• Who handles the JobSubmitted event?

► Specific handler called handleJobSubmitted Line 6

Walk-through to the Job Submitted handler

- Create a new job, called ActiveJob
- New job starts with only 1 stage, corresponding to the last stage of the job upon which an action is called
- Use the dependency information to produce additional stages
 - ★ Shuffle Dependency: create a new map stage Line 16
 - Narrow Dependency: pipes them into a single stage qetMissingParentStages

Lines 8-9

More About Stages

What is a DAG

- Directed acyclic graph of stages
- Stage boundaries determined by the shuffle phase
- Stages are run in topological order

Definition of a Stage

- Set of independent tasks
- All tasks of a stage apply the same function
- All tasks of a stage have the same dependency type
- All tasks in a stage belong to a TaskSet

Stage types

- Shuffle Map Stage: stage tasks results are inputs for another stage
- Result Stage: tasks compute the final action that initiated a job (e.g., count (), save (), etc.)

The Task Scheduler

Task oriented scheduling

- Schedules tasks for a single SparkContext
- Submits tasks sets produced by the DAG Scheduler
- Retries failed tasks
- Takes care of stragglers with speculative execution
- Produces events for the DAG Scheduler

Implementation details

- The Task scheduler creates a TaskSetManager to wrap the TaskSet from the DAG scheduler
 Jine 28
- ► The TaskSetManager class operates as follows:
 - ★ Keeps track of each task status
 - Retries failed tasks
 - Imposes data locality using delayed scheduling Lines 29,30
- Message passing implemented using Actors, and precisely using the Akka framework

Running Tasks on Executors



```
Executor, scala
1 def launchTask(serializedTask){
        val tr = new TaskRunner(serializedTask)
        threadPool.execute(tr)
4 }
    TaskRunner, scala
    5 def run(){
            executorBackend.statusUpdate(RUNNING)
            val task = serializer.deserialize(serializedTask)
            val value = task.run()
            val res = new DirectTaskResult(serializer.serialize(value))
            executorBackend.statusUpdate(FINISHED)
    11
            if res.size > akkaFrameSize
    12
                    blockManager.putBytes(taskId, res)
            else
    13
    14
                    return res
    15 }
        16 task.run()
        17 if task is ResultTask
                val (rdd, func) = ser.deserialize(RDD, taskContext)
        19
                return result = func(rdd.iterator(partition))
        20
        21 if task is ShuffleMapTask
                val (rdd, dep) = ser.deserialize(RDD, ShuffleDependency, taskContext)
                shuffleWriter = shuffleManager.getWriter(dep.shuffleHandler, partitionId)
        24
                shuffleWriter.write(rdd.iterator(partition))
        25
                return shuffleWriter.stop().get()
```

Running Tasks on Executors

Executors run two kinds of tasks

- ResultTask: apply the action on the RDD, once it has been computed, alongside all its dependencies Line 19
- ► ShuffleTask: use the Block Manager to store shuffle output using the ShuffleWriter
 Lines 23,24
- ► The ShuffleRead component depends on the type of the RDD, which is determined by the compute function and the transformation applied to it

Data Shuffling

The Spark Shuffle Mechanism

Same concept as for Hadoop MapReduce, involving:

- Storage of "intermediate" results on the local file-system
- Partitioning of "intermediate" data
- Serialization / De-serialization
- Pulling data over the network

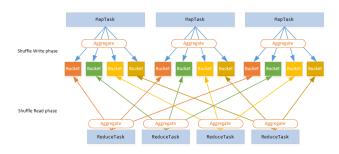
Transformations requiring a shuffle phase

```
▶ groupByKey(), reduceByKey(), sortByKey(), distinct()
```

Various types of Shuffle

- Hash Shuffle
- Consolidate Hash Shuffle
- Sort-based Shuffle

The Spark Shuffle Mechanism: an Illustration

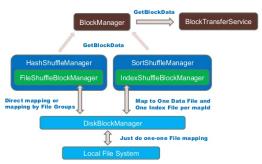


Data Aggregation

- ▶ **Defined on** ShuffleMapTask
- Two methods available:
 - ★ AppendOnlyMap: in-memory hash table combiner
 - ★ ExternalAppendOnlyMap: memory + disk hash table combiner

Batching disk writes to increase throughput

The Spark Shuffle Mechanism: Implementation Details



Pluggable component

- Shuffle Manager: components registered to SparkEnv, configured through SparkConf
- Shuffle Writer: tracks "intermediate data" for the MapOutputTracker
- Shuffle Reader. pull-based mechanism used by the ShuffleRDD
- Shuffle Block Manager: mapping between logical partitioning and the physical layout of data

The Hash Shuffle Mechanism



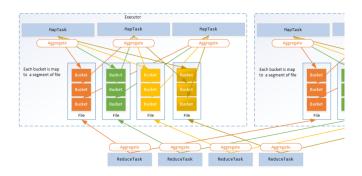
Map Tasks write output to multiple files

- Assume: m map tasks and r reduce tasks
- ► Then: m × r shuffle files as well as in-memory buffers (for batching writes)

• Be careful on storage space requirements!

- Buffer size must not be too big with many tasks
- Buffer size must not be too small, for otherwise throughput decreases

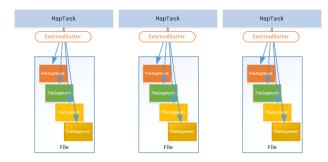
The Consolidate Hash Shuffle Mechanism



Addresses buffer size problems

- Executor view vs. Task view
 - Buckets are consolidated in a single file
- ▶ Hence: $F = C \times r$ files and buffers, where C is the number of Task threads within an Executor

The Sort-based Shuffle Mechanism



Implements the Hadoop Shuffle mechanism

- Single shuffle file, plus an index file to find "buckets"
- Very beneficial for write throughput, as more disk writes can be batched

Sorting mechanism

- Pluggable external sorter
- Degenerates to Hash Shuffle if no sorting is required

Data Transfer: Implementation Details

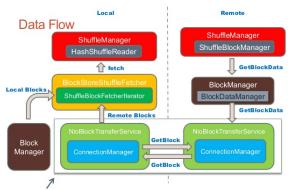
BlockTransfer Service

- General interface for ShuffleFetcher
- Uses BlockDataManager to get local data

Shuffle Client

- Manages and wraps the "client-side", setting up the TransportContext and TransportClient
- Transport Context: manages the transport layer
- Transport Server: streaming server
- Transport Client: fetches consecutive chunks

Data Transfer: an Illustration



Can Switch to different BlockTransferService