

1) $L_1 \subseteq L_2$

a) $L_1 = \{a^n b^n \mid n \geq 1\}$ $L_2 = a^* b^*$

We can see for this language, everything in L_1 can be represented by L_2 . However, L_1 is not irregular. So therefore $L_1 \subseteq L_2$ and L_1 is not regular does NOT imply L_2 is irregular

b) $L_1 \subseteq L_2$ if L_2 is not regular, then L_1 is not regular

Let $L_2 = \{a^n b^n \mid n \geq 1\}$ and $L_1 = ab$

$L_1 \subseteq L_2$ in this case, but L_1 is regular and L_2 is not. Therefore, $L_1 \subseteq L_2$ when L_2 is non-regular does NOT imply L_1 is non-regular

c) If L_1 and L_2 are not regular, then $L_1 \cup L_2$ is non-regular.

Let $L_1 = \{a^n b^n \mid n \leq m\}$ & $L_2 = \{a^n b^n \mid n > m\}$

Then $L_1 \cup L_2 = a^* b^*$ which is regular.

d) If L_1 and L_2 are not regular, then $L_1 \cap L_2$ is non-regular

$L_1 = \{a^n b^n \mid n \leq m\}$ $L_2 = \{a^n b^n \mid n \geq m\}$

Then $L_1 \cap L_2 = \{a^m b^m\}$ which is regular.

Thus, that statement is FALSE

e) If L_1 is regular and L_2 is not regular, then $L_1 \cup L_2$ is not regular

Let $L_1 = a^* b^*$ and $L_2 = \{a^n b^n \mid n \geq 1\}$ then $L_1 \cup L_2$?

then $L_1 \cup L_2 = a^* b^*$ which is regular

Thus, this statement is FALSE

2) a) $L = \{0^i 1^j 0^k \mid k > i + j\}$ Assume L is regular

1) select w , $w = 0^p 1^p 0^{2p}$ $|w| > M \checkmark w \in L$

$N = \text{PUMPING LENGTH}$

$w = xyz$ $|y| \geq 1$ $|xy| \leq N$

$w = 0110000$

$i=1$
 $\begin{array}{c} 0110000 \\ \times \quad y \quad z \end{array}$

$i=2$
 $\begin{array}{c} 01110000 \\ \times \quad y \quad z \end{array}$ but $i=1, j=3$ and $k=4$, $i+j=4$. so $k \nless i+j$.

so L is not regular

b) $\{0^i 1^j \mid j=i \text{ or } j=2i\} = L$ Assume L is regular

Let $w = 0^N 1^{2N}$ so $|w| \geq N$ and $w \in L$

$$w = x y z$$

$$i=1$$

$$w' = \begin{array}{|c|c|c|} \hline 0 & 1 & 1 \\ \hline x & y & z \\ \hline \end{array}$$

$$i=2$$

$$w'' = \begin{array}{|c|c|c|} \hline 0 & 1 & 1 \\ \hline x & y & z \\ \hline \end{array}$$

But now $j=3i$, which is not accepted by L .

Therefore, L is not regular.

c) $L = \{x \in \{0,1\}^* \mid x \text{ is a palindrome}\}$ Assume L is regular

Let $w = 0^N 11 0^N$ so $|w| \geq N$ and $w \in L$

$$w = x y z \quad |xy| \leq N \quad |y| \geq 1$$

$$i=1$$

$$w = x y z \quad \text{because } |xy| \leq N, \text{ must all be } 0\text{'s}$$

$$\begin{array}{|c|c|c|c|c|} \hline 0 & 0 & 1 & 1 & 0 & 0 \\ \hline x & y & & & z \\ \hline \end{array}$$

$$i=2$$

$$w = x y y z$$

$$\begin{array}{|c|c|c|c|c|c|} \hline 0 & 0 & 0 & 1 & 1 & 0 & 0 \\ \hline x & y & & & & z \\ \hline \end{array} \quad \text{but } w \text{ is not a palindrome here.}$$

Thus, L is not a regular language.

d) $L = \{0^{2^n} \mid n \geq 0\}$ Assume L is Regular

$$w = 0^{2^N} \quad w \in L \quad |w| = 2^N \quad |w| > N$$

$$w = x y z \quad |xy| \leq N \quad |y| \geq 1$$

$$w' = x y y z$$

$$|w'| = |w| + |y| = 2^N + |y|$$

$$2^N < |w'| \leq 2^N + N$$

$|w|$ is supposed to be of size 2^N , yet we see by pumping we can get $2^N + N$.

Therefore, L is not regular.