Topics in Algebra, Chapter 1 Solutions

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1.1 - Set Theory

(1a) If $A \subset B$ and $B \subset C$, then every element $a \in A$ is also $a \in B$. Likewise, $B \subset C \Leftrightarrow (a \in B \Rightarrow a \in C)$. Thus, every element $a \in A$ is also in C. Therefore, $A \subset C$.

(1b) If we presuppose that $B \subset A$, this means that $x \in B \Rightarrow x \in A$. In set-builder notation, we know that

$$A \cup B = \{x \mid x \in A \vee x \in B\}$$

Because every element of B is also in A, the set builder is reduced to

$$A \cup B = \{x \mid x \in A\} = A$$

Conversely, if we start by supposing that $A \cup B = A$, this implies, by logical necessity, that

$$(x \in A) \Leftrightarrow (x \in A \vee x \in B)$$

As alluded to above, this is logically transformable into the statement that $x \in B \Rightarrow x \in A$, and thus $B \subset A$.

1.1 - Set Theory

(1c) In the set $B \cup C$, every element satisfies $x \in B \vee x \in C$. Likewise, in $A \cup C$, every element satisfies $x \in A \vee x \in C$. In the case that $x \in C$, the right side of the disjunction is satisfied in both sets. If $x \in B$, then $B \subset A$ implies that $x \in A$, and so both sets are satisfied again. This means that every element in $B \cup C$ is in $A \cup C$, and so $B \cup C \subset A \cup C$.

Similarly, every element of $B \cap C$ satisfies $x \in B \wedge x \in C$ and every element of $A \cap C$ satisfies $x \in A \wedge x \in C$. Because, again, due to $B \subset A$, we have that $x \in B \Rightarrow x \in A$ and thus every element in $B \cap C$ must also be in $A \cap C$ and so $B \cap C \subset A \cap C$.

1.1 - Set Theory

(2a) The commutativity of set intersection ($\cap$) and set union ($\cup$) follow from the analogous commutativity properties of logical operators $\wedge$ and $\vee$, respectively. Because $A \cap B$ consists of the elements which satisfy

$$x \in A \wedge x \in B \quad \equiv \quad x \in B \wedge x \in A$$

and thus it follows that $A \cap B = B \cap A$. The same argument follows for $A \cup B = B \cup A$.

(2b) As with the last problem, the associativity of $\cap$ follows from the associativity of logical $\wedge$. That is, the set $(A \cap B) \cap C$ consists of elements which satisfy,

$$(x \in A \wedge x \in B) \wedge x \in C \quad \equiv \quad x \in A \wedge (x \in B \wedge x \in C)$$

And thus, exactly those same elements satisfy the necessary condition for being elements of $A \cap (B \cap C)$, giving the equality $(A \cap B) \cap C = A \cap (B \cap C)$.

1.1 - Set Theory

(3) Logically speaking, the elements of $A \cup (B \cap C)$ are those which satisfy $x \in A \vee (x \in B \wedge x \in C)$. By the distributivity of $\vee$, it follows that this logical condition is equivalent to the condition,

$$x \in A \vee (x \in B \wedge x \in C) \quad \equiv \quad (x \in A \vee x \in B) \wedge (x \in A \vee x \in C)$$

And thus, the elements in the former set must be exactly those elements in the latter set, giving the equality $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$.

(4a) The elements of $(A \cap B)'$ are those which satisfy $\neg(x \in A \wedge x \in B)$. Applying De Morgan's logical negation rules, we get that this is logically equivalent to $x \notin A \vee x \notin B$. The set whose elements all satisfy this condition is identically $A' \cup B'$.

(4b) As in the last problem, we translate the set $(A \cup B)'$ into the membership condition $\neg(x \in A \vee x \in B)$. De Morgan transforms this into the membership condition $x \notin A \wedge x \notin B$. This membership condition corresponds to the identical set, $A' \cap B'$.

1.1 - Set Theory

(5) There are two cases: (i) A is disjoint from B ($A \cap B = \emptyset$), or (ii) A is not disjoint from B ($A \cap B \neq \emptyset$).

In case (i), we trivially have that $o(A \cap B) = o(\emptyset) = 0$. So, it is trivially the case that every element of A is represented in $A \cup B$ and likewise for B, and all of these elements are distinct. Therefore, $o(A \cup B) = o(A) + o(B)$, which works for our hypothesis given that $o(A \cap B) = 0$.

In case (ii), let's suppose that there are k elements in common between A and B, such that $o(A \cap B) = k$. Because $A \cup B$ will only contain one copy of each of the common elements, we must subtract $o(A \cap B)$ from $o(A) + o(B)$ to get the number of unique elements in $A \cup B$.

1.1 - Set Theory

(6) First, we construct on A a bijective index $i : A \to [n]$. Every subset $S \subset A$ is uniquely identified by an n-tuple of binary variables $(b_1, \ldots, b_n) \in B$, with $b_i \in \{0, 1\}$. It's trivial to see that the set of all binary n-tuples B has 2^n elements. From this set, we have a bijection which generates (and indexes) the set of all subsets of A (hereafter called the *power set* of A, $\mathcal{P}(A)$). That is, we define $\sigma : B \to \mathcal{P}(A)$ via the map

$$\sigma : b \mapsto \bigcup_{b_j = 1} i^{-1}(j)$$

In other words, if a particular subset had $b_j = 1$, this means the element $a \in A$ with $i(a) = j$ is included in the subset. Because there exists a bijective map between the finite sets B and $\mathcal{P}(A)$, we have that they are the same size. Therefore, $\mathcal{P}(A)$ has 2^n elements.

1.1 - Set Theory

(7) Let S be the set of Americans. Let C be the set of Americans that like cheese, and A be the set of Americans that like apples. The proportions given suggest that $|C|/|S| = 0.63$ and that $|A|/|S| = 0.76$. Because both $C \subset S$ and $A \subset S$, we expect $C \cup A \subset S$. This necessitates that

$$|C \cup A| \leq |S|$$
$$|C| + |A| - |C \cap A| \leq |S|$$

Rearranging terms and dividing through by $|S|$, we get a bound on the proportion of Americans which like both cheese and apples:

$$|C \cap A| \geq \frac{|C|}{|S|} + \frac{|A|}{|S|} - 1.0$$
$$|C \cap A| \geq 0.63 + 0.76 - 1.0$$
$$|C \cap A| \geq 0.39$$

In English, no fewer than 39% of Americans like both cheese and apples.

1.1 - Set Theory

(8) Recall that set difference $A - B$ entails the membership condition $x \in A \wedge x \notin B$. This means that the *symmetric difference* $A * B$ entails the membership condition,

$$(x \in A \wedge x \notin B) \vee (x \in B \wedge x \notin A)$$

By double-distributing the disjunction, we get to a CNF representation, and then set-difference:

$$\begin{gathered}(x \in A \vee x \in B) \wedge (x \in A \vee x \notin A) \wedge \\ (x \notin B \vee x \in B) \wedge (x \notin B \vee x \notin A) \\ \Downarrow \\ (x \in A \vee x \in B) \wedge (x \notin B \vee x \notin A) \\ \Downarrow \\ (x \in A \vee x \in B) \wedge \neg(x \in A \wedge x \in B)\end{gathered}$$

This shows logical equivalence in membership between $A * B$ and $(A \cup B) - (A \cap B)$.

1.1 - Set Theory

(9a) For brevity, we argue non-symbolically. Every element of $(A + B)$ is either unique to A or unique to B. Taking the symmetric difference, then, $(A + B) + C$ must result in those elements which are either unique to A, unique to B, unique to C, or common to all three of them.

Likewise, the elements of $(B + C)$ are those elements unique to B and unique to C. Taking the symmetric difference, we see that $A + (B + C)$ is the set of elements which are unique to A, unique to B, unique to C, or common to all three. Therefore, the sets are equal and so $(A + B) + C = A + (B + C)$.

1.1 - Set Theory

(9b) Reducing the symmetric differences and distributing the leftmost conjunction,

$$A \cap ((B - C) \cup (C - B))$$
$$(A \cap (B - C)) \cup (A \cap (C - B))$$

We then use the fact that intersection distributes over set difference:

$$((A \cap B) - (A \cap C)) \cup ((A \cap C) - (A \cap B))$$
$$\Downarrow$$
$$A \cdot B + A \cdot C$$

(9c) It's trivially the case that $A \cdot A = A \cap A = A$.

1.1 - Set Theory

(9d) Every element of A is also (identically) an element of A, so there are no elements unique to either set. Therefore, $A + A = \emptyset$

(9e) Taking the left-symmetric-difference from A of both sides of the equation and applying the associativity from part (a) with the cancellation property of part (d), we get:

$$A + (A + B) = A + (A + C)$$
$$(A + A) + B = (A + A) + C)$$
$$\emptyset + B = \emptyset + C$$

Trivially, the symmetric difference of any set with the empty set is simply the set itself, which gives us $B = C$.

1.1 - Set Theory

(10a) The relation of having a common ancestor is reflexive, symmetric and transitive. Therefore, it is a valid equivalence relation.

(10b) The relation of living within 100 miles of each other is reflexive and symmetric, but not transitive. Therefore, it is not a valid equivalence relation.

(10c) The relation of having the same father is reflexive, symmetric and transitive. Therefore, it is a valid equivalence relation.

(10d) The relation of having the same absolute value is reflexive, symmetric and transitive. Therefore, it is a valid equivalence relation.

(10e) The relation of being strictly greater and strictly lesser than one another is impossible to satisfy. Therefore, it is not a valid equivalence relation.

(10f) The relation of two lines having the same slope in the plane is reflexive, symmetric and transitive. Therefore, it is a valid equivalence relation.

1.1 - Set Theory

(11a) Using only symmetry and transitivity, we do not have a guarantee that there exists a b for a such that $a \sim b$. Thus, this argument does not account for cases that each equivalence class $[a]$ are each only individual elements.

(11b) If we include a property known as *seriality*, which necessitates that every a has a b such that $a \sim b$. Under this assumption, both symmetry and transitivity imply reflexivity.

1.1 - Set Theory

(12) Clearly, $a \sim a$ because 0 is a multiple of n. Likewise, if $a \sim b$, then it means $a - b = pn$ and so $b - a = -pn$, which is also a multiple of n, implying $b \sim a$. Finally, $a \sim b$ and $b \sim c$ mean that $a - b = pn$ and $b - c = qn$. Adding the equations together gives us $a - c = (p + q)n$, so $a \sim c$. Therefore, differing by a multiple of n is a valid equivalence relation. Each of the equivalence classes are defined to be $\mathrm{cl}(i) = \{x \in \mathbb{Z} \mid x \equiv i \pmod{n}\}$. Because every integer must have a remainder in $[n]$ after division by n, we have that $\mathrm{cl} : [n] \to \mathbb{Z}$ is a surjection, and so there are at most n equivalence classes. Because each equivalence class $\mathrm{cl}(i)$ trivially contains i for $0 \leq i < n$, we have that there are at least n equivalence classes. Therefore, there are exactly n equivalence classes.

1.1 - Set Theory

(13) It's clear that being in the same mutually disjoint subset A_α is a valid equivalence relation. This follows from the demonstrable reflexivity, symmetry and transitivity. Moreover, the equivalence classes must be the distinct A_α's because that is how we defined our equivalence.

1.2 - Mappings

As a matter of notation, I invoke $\sqrt{t}$ as meaning the *positive square root* of t, so $\sqrt{t} \geq 0$.

(1a) σ is surjective, but not injective. Every $t \in T$ is mapped to the set $\sigma^{-1}(t) = \{-\sqrt{t}, \sqrt{t}\}$.

(1b) σ is both injective and surjective. Every $t \in T$ is mapped to its pre-image $\sigma^{-1}(t) = \sqrt{t}$.

(1c) σ is injective, but not surjective. Only the perfect squares $t \in T$ have pre-images of the form $\sigma^{-1}(t) = \sqrt{t}$.

(1d) σ is injective, but not surjective. Only the even integers $t \in T$ have pre-images of the form $\sigma^{-1}(t) = t/2$.

1.2 - Mappings

(2) I define the injection $\alpha : S \times T \to T \times S$ as the "swap" map, $\alpha : (s, t) \mapsto (t, s)$. This map is provably an injection because

$$(t, s) = (t', s') \quad \Leftrightarrow \quad (s, t) = (s', t')$$

So, this injection, α, is evidence of the one-to-one correspondence between the sets.

(3a) As in problem (2), the injection $\alpha : (S \times T) \times U \to S \times (T \times U)$ defined by $\alpha : ((s, t), u) \mapsto (s, (t, u))$ evidences a one-to-one correspondence between the sets.

(3b) We define the injection $\alpha : (S \times T) \times U \to S \times T \times U$ via the map $\alpha : ((s, t), u) \mapsto (s, t, u)$, which evidences the one-to-one correspondence between the sets.

1.2 - Mappings

(4a) If there's a one-to-one correspondence between S and T, then there exists an injection $\sigma : S \to T$ satisfying $\sigma(s) = \sigma(s') \Rightarrow s = s'$. This injection induces an inverse injection, $\sigma^{-1} : T \to S$ defined by $\sigma^{-1}(t) = \left\{t \in T \mid \bigcup_{s \in S} \sigma(s) = t\right\}$. The injection σ^{-1} evidences a one-to-one correspondence between T and S.

(4b) Suppose that the injections $\alpha : S \to T$ and $\beta : T \to U$ evidence the one-to-one correspondences between S and T, and T and U respectively. Then we can define the composed injection $\sigma : S \to U$ via the map $\sigma : s \mapsto \beta(\alpha(s))$. Because $\beta(t) = \beta(t') \Rightarrow t = t'$ and $\alpha(s) = \alpha(s') \Rightarrow s = s'$, it follows that $\sigma(s) = \sigma(s') \Rightarrow s = s'$ and thus σ is injective. This shows that there is a one-to-one correspondence between S and U.

1.2 - Mappings

(5) Because the identity automorphism $\iota : s \mapsto s$, it's clear that $\sigma \circ \iota : s \mapsto \sigma(s)$, which is identical to the map $\sigma : s \mapsto \sigma(s)$. An identical argument holds for $\iota \circ \sigma$, and so it holds that $\sigma = \sigma \circ \iota = \iota \circ \sigma$.

(6) Because we know that $|S^*| > |S|$ for any set, it is impossible for any mapping of the $|S|$ elements of S to cover the $|S^*|$ elements of S^*. Therefore, no mapping $S \to S^*$ will ever be surjective.

(7) We can consider constructing an element $\sigma \in A(S)$ as a sequence of choosing unique images $\sigma(s_i)$ for each $s_i \in S$, $1 \leq i \leq n$. s_1 could have any one of the n elements of S as its image. Each subsequent s_i will have one fewer option for its image. This generates a set of $n!$ distinct automorphisms, and this set is $A(S)$.

1.2 - Mappings

(8a) Suppose for the sake of argument that $\sigma : S \to S$ is surjective, but not injective. Specifically, suppose that there exists an $s^* \in S$ with more than one preimage, $\gamma = \sigma^{-1}(s^*)$. Then, it follows that σ is still surjective if its restriction $\bar{\sigma} : S - \gamma \to S - s^*$ is surjective (note that we must remove γ from our domain in order to enforce the constraint that an element of the domain cannot have two different images). However, the domain of $\bar{\sigma}$ is smaller than its codomain, which means that $\bar{\sigma}$ cannot be surjective because no element of the domain can map to two different images. Therefore, neither $\bar{\sigma}$ nor σ itself can be surjective without also being injective.

1.2 - Mappings

(8b) Using the reverse argument from problem (8a), we suppose for the sake of contradiction that σ is not surjective. Because it *is* injective, we have that every element $s \in S$ in the domain is mapped to a unique image $\sigma(s) \in S$. Because we assume it is not surjective, we assume that there is an $s^* \in S$ which does not have a preimage. So, σ maps the $|S|$ elements of S to the $|S| - 1$ elements of $S - \{s^*\}$. By the pigeonhole principle, σ must map two distinct inputs to the same image, and so σ cannot be injective. We have a contradiction, so we prove the result

$$\sigma \text{ is injective} \quad \Leftrightarrow \quad \sigma \text{ is surjective}$$

1.2 - Mappings

(8c) Consider the mapping $\sigma : \mathbb{Z} \to \mathbb{Z}$ via the assignment $n \mapsto \lfloor n/2 \rfloor$. Clearly, σ is surjective because every integer can be doubled to result in one of its preimages. However, it is not injective because n has the multiple preimages $\{2n, 2n+1\}$. So, σ is an infinite counter-example to (8a).

On the other hand, the mapping $\sigma : \mathbb{Z} \to \mathbb{Z}$ via the assignment $n \mapsto 2n$ is clearly injective because every integer can be doubled to result in a unique image. However, it is not surjective because the odd integers have no preimages (they are not the result of any integer being doubled). This evidences an infinite counter-example to (8b).

1.2 - Mappings

(9a) Consider the maps

$$\begin{aligned} \sigma &: [2] \to [2] \qquad & \sigma &: i \mapsto 1 \\ \tau &: [2] \to [1] \qquad & \tau &: 1 \mapsto 1 \end{aligned}$$

Clearly, σ is not surjective because 2 has no preimage. However, $\sigma \circ \tau$ is surjective because 1 does have a preimage. So, the converse of the lemma is false.

(9b) Consider the maps

$$\begin{aligned} \sigma &: [1] \to [2] \qquad & \sigma &: 1 \mapsto 1 \\ \tau &: [2] \to [2] \qquad & \tau &: i \mapsto 1 \end{aligned}$$

Clearly, τ is not injective because both 1 and 2 are mapped to the same image, 1. However, $\sigma \circ \tau$ is injective because 1 has only one preimage. So, the converse of the lemma is false.

1.2 - Mappings

(10) We define the map $\kappa : \mathbb{Z} \to \mathbb{Q}$ via the assignment based on prime factorization of each integer:

$$p_1^{r_1} p_2^{r_2} \cdots p_n^{r_n} \quad \mapsto \quad \frac{p_1^{r_1}}{p_2^{r_2} \cdots p_n^{r_n}}$$

It's clear that the numerator and denominator are coprime, and so the image of each integer is a valid rational. Moreover, every rational has at most one preimage. If two images (p/q) and (p'/q') are equal, then it means that their components $p = p'$ and $q = q'$, and so their preimages are equal, $pq = p'q'$. This shows that κ is an injection and so there is a one-to-one correspondence between the integers and the rationals.

1.2 - Mappings

(11a) Trivially, because σ is a mapping into T, and $\sigma(A) = \sigma_A(A)$, σ_A is thus a mapping $A \to T$.

(11b) If we assume that σ is injective, then we know that every element of $\sigma(S)$ is the image of exactly one $s \in S$. Because $\sigma(A) \subset \sigma(S)$, we have that every element of $\sigma(A)$ is the image of exactly one $a \in A$. Because $\sigma(A) = \sigma_A(A)$, the same property is true of σ_A's domain. Therefore, σ_A is injective.

(11c) If the set of elements, $\bar{T} \subset T$, in T which have more than one preimage is disjoint from $\sigma(A)$ such that $\sigma(A) \cap \bar{T} = \emptyset$, then σ will still have the property of mapping each $a \in A$ to a unique image $\sigma(a) \in \sigma(A)$. Therefore, σ_A can still be injective, even if σ as a whole is not.

1.2 - Mappings

(12) First, we recognize that $\sigma(A) \subset A$, so $\sigma \circ \sigma(A) \subset A$. On the other hand, $\sigma_A(A) = \sigma(A) \subset A$, and likewise for $\sigma_A \circ \sigma_A(A)$. Because of this equality, we have that $\sigma \circ \sigma(A) = \sigma_A \circ \sigma_A(A)$, and so the domain of both $(\sigma \circ \sigma)_A$ and $\sigma_A \circ \sigma_A$ are mapped to the same images, and so the functions are the same.

(13a) Consider the proper subset $5\mathbb{Z} \subset \mathbb{Z}$. The injection $\sigma : x \mapsto 5x$ evidences a one-to-one correspondence between the sets, and so $\mathbb{Z}$ is infinite.

(13b) Consider the proper subset $\mathbb{R}_{>0} \subset \mathbb{R}$. The injection $\sigma : x \mapsto e^x$ evidences a one-to-one correspondence between the sets, and so $\mathbb{R}$ is infinite.

(13c) Because A is infinite, it has a subset $\bar{A} \subset A \subset S$ with which it has a one-to-one correspondence. This means that S also has the same one-to-one correspondence with $\bar{A}$ and so S is also infinite.

1.2 - Mappings

(14) Let "$S \to \mathbb{Z}$" denote that there is a one-to-one correspondence, called α, between S and $\mathbb{Z}$. Because $\mathbb{Z} \to \mathbb{Q}$ via the mapping κ from problem (10) and the transitivity proved in problem (4b), we have $S \to \mathbb{Q}$. Moreover, via the injection $f : \mathbb{Q} \to \mathbb{Z} \times \mathbb{Z}$ with map $p/q \mapsto (p, q)$, we establish $\mathbb{Q} \to \mathbb{Z} \times \mathbb{Z}$. Finally, by inverting and pairng α, we can inject $A^{-1} : \mathbb{Z} \times \mathbb{Z} \to S \times S$ via the assignment $A^{-1} : (p, q) \mapsto (\alpha^{-1}(p), \alpha^{-1}(q))$. This injection completes our transitive chain by implying $S \to S \times S$, as desired.

1.2 - Mappings

(15) First, we show that there exists a surjective map $\sigma : U \to S$. This follows from the existence of surjective maps $\alpha : U \to T$ and $\beta : T \to S$, which compose to give $\sigma = \alpha \circ \beta$.

Second, we consider a hypothetical surjective map $\gamma : S \to U$. If such a map existed, its inputs could first be mapped into T, then mapped surjectively onto U. This violates our assumption that there does not exist a surjective map between T and U. So, there cannot exist a surjective map from $S \to U$. Therefore, $S < U$.

1.2 - Mappings

(16) If we start by supposing that $m < n$, then we can easily find a surjective map $\sigma : T \to S$ by the pigeonhole principle. Moreover, we cannot form a surjective map $\gamma : S \to T$. If we assign every $s \in S$ in the domain to a unique image $t \in T$, then there will always be at least one element $t^* \in T$ with no preimage in S. Therefore, there can be no surjective map from $S \to T$. So, $S < T$.

Conversely, if we start by supposing that $S < T$, then it is evident that there cannot exist a surjective map $S \to T$. If $m \geq n$, then we can easily find a surjective map onto T by the pigeonhole principle. Therefore, this enforces that $m < n$.

Both directions prove that $S < T$ for finite sets S, T equivalently means $m < n$.