

# Pset 3

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- (c) Collaborators:
- (d) I have followed the academic integrity and collaboration policy
- (e) Hours:

## 1 Checking Connectivity

(a) **function** REDFIND( $G, u, v$ )  
    **for all**  $v \in V$  **do**  
        visited[ $v$ ]  $\leftarrow$  **False**  
     $Q \leftarrow$  MAKEQUEUE()  $\triangleright$  standard FIFO queue  
     $Q$ .PUSH( $u$ )  
    **while**  $Q$  not empty **do**  
         $q \leftarrow$   $Q$ .POP()  
        **if**  $q == v$  **then return** **True**  
        visited[ $q$ ]  $\leftarrow$  **True**  
        **for all**  $(q, v) \in E$  **do**  
            **if** COLOR( $q, v$ ) == **RED** and not visited[ $v$ ] **then**  
                 $Q$ .PUSH( $v$ )  
    **return** **False**

**Correctness**

Blah blah blah

**Running Time**

```

(b)  function DFS( $G, u, v$ )
      for all  $w \in V$  do
        visited[ $w$ ]  $\leftarrow$  False
      visited[ $u$ ]  $\leftarrow$  True
      for all  $(u, w) \in E$  do
        if COLOR( $u, w$ ) == Red then
          if EXPLORE( $G, w, v, \text{Red}$ ) then return True
        else if EXPLORE( $G, w, v, \text{Blue}$ ) then return True
      return False

function EXPLORE( $G, w, v, \text{edgeColor}$ )
  visited[ $w$ ]  $\leftarrow$  True
  if  $w == v$  then return True
  for all  $(w, u) \in E$  do
    if edgeColor == Blue then
      if COLOR( $w, u$ ) == Blue and not visited[ $u$ ] then
        if EXPLORE( $G, u, v, \text{Blue}$ ) then return True
    else ▷ previous edge is still on red path
      if not visited[ $u$ ] then
        if EXPLORE( $G, u, v, \text{COLOR}(w, u)$ ) then return True
  return False

```

## 2 Road Trip

```

(a) function ISREACHABLE( $G, s, t$ )
    for all  $v \in V$  do
        visited[ $v$ ]  $\leftarrow$  False
     $Q \leftarrow \text{MAKEQUEUE}()$  ▷ standard FIFO queue
     $Q.\text{PUSH}(s)$ 
    while  $Q$  not empty do
         $q \leftarrow Q.\text{POP}()$ 
        if  $q == t$  then return True
        visited[ $q$ ]  $\leftarrow$  True
        for all  $(q, v) \in E$  do
            if  $L \geq \ell(v, q)$  and not visited[ $v$ ] then
                 $Q.\text{PUSH}(v)$ 
    return False

(b) function LOWESTGAS( $G, \ell, s, t$ )
    for all  $v \in V$  do
        maxL[ $v$ ]  $\leftarrow \infty$ 
    maxL[ $s$ ]  $\leftarrow 0$ 
     $P \leftarrow \text{MAKEQUEUE}(V)$  ▷ priority queue with maxL values as keys
    while  $P$  not empty do
         $v \leftarrow \text{EXTRACTMIN}(P)$ 
        for all  $(v, w) \in E$  do
            if maxL[ $w$ ]  $>$  maxL[ $v$ ] and maxL[ $w$ ]  $>$   $\ell(v, w)$  then
                maxL[ $w$ ]  $\leftarrow \text{MAX}(\text{maxL}[v], \ell(v, w))$ 
                 $\text{CHANGEKEY}(P, w)$ 
    return maxL[ $t$ ]

```

### 3 Counting Shortest Paths

Struct Definition

```
struct {
    int dist, numPaths;
} PathStruct;

function NUMSHORTEST( $G, \ell, s, t$ )
    for all  $v \in V$  do
        paths[v].dist  $\leftarrow \infty$ 
        paths[v].numPaths  $\leftarrow 0$ 
    paths[s].dist  $\leftarrow 0$ 
    paths[s].numPaths  $\leftarrow 1$ 
     $P \leftarrow \text{MAKEQUEUE}(V)$   $\triangleright$  priority queue with dist values as keys
    while  $P$  not empty do
         $v \leftarrow \text{EXTRACTMIN}(P)$ 
        for all  $(v, w) \in E$  do
            if paths[w].dist  $>$  paths[v].dist +  $\ell(v, w)$  then
                paths[w].dist  $\leftarrow$  paths[v].dist +  $\ell(v, w)$ 
                paths[w].numPaths  $\leftarrow$  paths[v].numPaths
            else if paths[w].dist == paths[v].dist +  $\ell(v, w)$  then
                paths[w].numPaths  $\leftarrow$  paths[w].numPaths + paths[v].numPaths
    return paths[t].numPaths
```

## 4 Spanning Tree with Leaves

We implement Kruskal's algorithm two times.

```
function LEAFSPANNINGTREE( $G, U, w$ )  
  for all  $v \in V$  do  
    MAKESET( $v$ )  
   $F \leftarrow \{ \}$   
  sort edges  $E$  by increasing weight  
   $E' \leftarrow \{ \{u, v\} \in E \mid u \notin U \wedge v \notin U \}$   
   $E'' \leftarrow \{ \{u, v\} \in E \mid u \in U \oplus v \in U \}$   
  for all  $\{u, v\} \in E'$  do  
    if FIND( $u$ )  $\neq$  FIND( $v$ ) then  
       $F \leftarrow F \cup \{ \{u, v\} \}$   
      UNION( $\{u, v\}$ )  
  for all  $\{u, v\} \in E''$  do  
    if FIND( $u$ )  $\neq$  FIND( $v$ ) then  
       $F \leftarrow F \cup \{ \{u, v\} \}$   
      UNION( $\{u, v\}$ )
```

## 5 Perfect Matching in a Tree

```

function CHECKPERFECTMATCHING( $G$ )
     $Q \leftarrow \text{MAKEQUEUE}()$                                  $\triangleright$  standard FIFO queue
     $P \leftarrow \text{MAKEQUEUE}()$                                  $\triangleright$  standard FIFO queue
    for all  $v \in V$  do
         $\text{isPaired}[v] \leftarrow \text{False}$ 
        if degree of  $v == 1$  then
             $Q.\text{PUSH}(v)$ 
    while  $Q$  not empty and  $P$  not empty do
        while  $Q$  not empty do
             $q \leftarrow Q.\text{POP}()$ 
            if  $\text{isPaired}[q]$  then
                continue
             $r \leftarrow \text{None}$ 
            for all  $\{q, v\} \in E$  do                                 $\triangleright$  at most 1 unpaired neighbor exists
                if not  $\text{isPaired}[v]$  then
                     $r \leftarrow v$ 
                break
            if  $r == \text{None}$  then return False
             $\text{isPaired}[q] \leftarrow \text{True}$ 
             $\text{isPaired}[r] \leftarrow \text{True}$ 
             $r\text{NumUnpairedNeighbors} \leftarrow 0$ 
             $r\text{UnpairedNeighbor} \leftarrow \text{None}$ 
            for all  $\{r, v\} \in E$  do
                 $r\text{NumUnpairedNeighbors} \leftarrow r\text{NumUnpairedNeighbors} + 1$ 
                 $r\text{UnpairedNeighbor} \leftarrow v$ 
            if  $r\text{NumUnpairedNeighbors} == 1$  then
                 $Q.\text{PUSH}(r\text{UnpairedNeighbor})$ 
            else
                 $P.\text{PUSH}(r)$ 
        while  $P$  not empty do
             $p \leftarrow P.\text{POP}()$ 
            if  $\text{isPaired}[p]$  then
                continue
             $\text{numUnpairedNeighbors} \leftarrow 0$ 
            for all  $\{p, v\} \in E$  do
                if not  $\text{isPaired}[v]$  then
                     $\text{numUnpairedNeighbors} \leftarrow \text{numUnpairedNeighbors} + 1$ 
                if  $\text{numUnpairedNeighbors} > 1$  then
                    skip to next  $p$  in queue
             $Q.\text{PUSH}(p)$                                  $\triangleright$  node only has 1 unpaired neighbor
    return True

```