

# Discrete Mathematics

## Tutorial 2

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# Induction

We need to prove

$$P(n), \text{ for any } n \geq 1$$

Initial step  $n = 1$

Prove  $P(1)$

Inductive hypothesis

Suppose that

$$P(1), P(2), \dots, P(k)$$

Inductive step

Prove  $P(k + 1)$

# Induction

## Fibonacci numbers

$$f_0 = 0, \quad f_1 = 1, \quad f_n = f_{n-1} + f_{n-2}$$

0 1 1 2 3 5 8 13 21 ...

$$f_n = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^n - \left(\frac{1-\sqrt{5}}{2}\right)^n}{\sqrt{5}}$$

# Induction

## Fibonacci numbers

$$f_0 = 0, \quad f_1 = 1, \quad f_n = f_{n-1} + f_{n-2}$$

Initial step. 1)  $n = 0$

$$f_0 = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^0 - \left(\frac{1-\sqrt{5}}{2}\right)^0}{\sqrt{5}} = \frac{1-1}{\sqrt{5}} = 0$$

# Induction

## Fibonacci numbers

$$f_0 = 0, \quad f_1 = 1, \quad f_n = f_{n-1} + f_{n-2}$$

Initial step. 2)  $n = 1$

$$\begin{aligned} f_1 &= \frac{\left(\frac{1+\sqrt{5}}{2}\right)^1 - \left(\frac{1-\sqrt{5}}{2}\right)^1}{\sqrt{5}} = \\ &= \frac{1 + \sqrt{5} - (1 - \sqrt{5})}{2\sqrt{5}} = \frac{2\sqrt{5}}{2\sqrt{5}} = 1 \end{aligned}$$

# Induction

## Fibonacci numbers

$$f_0 = 0, \quad f_1 = 1, \quad f_n = f_{n-1} + f_{n-2}$$

Induction hypothesis. Suppose that there is  $k$  such that, for any  $i \leq k$ ,

$$f_i = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^i - \left(\frac{1-\sqrt{5}}{2}\right)^i}{\sqrt{5}}$$

# Induction

## Fibonacci numbers

$$f_0 = 0, \quad f_1 = 1, \quad f_n = f_{n-1} + f_{n-2}$$

Inductive step. We need to prove

$$f_{k+1} = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^{k+1} - \left(\frac{1-\sqrt{5}}{2}\right)^{k+1}}{\sqrt{5}}$$

# Induction

## Fibonacci numbers

$$f_0 = 0, \quad f_1 = 1, \quad f_n = f_{n-1} + f_{n-2}$$

$$f_{k+1} = f_k + f_{k-1}$$

$$f_k = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^k - \left(\frac{1-\sqrt{5}}{2}\right)^k}{\sqrt{5}}$$

$$f_{k-1} = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^{k-1} - \left(\frac{1-\sqrt{5}}{2}\right)^{k-1}}{\sqrt{5}}$$



# Induction

$$\begin{aligned}f_{k+1} &= \frac{\left(\frac{1+\sqrt{5}}{2}\right)^k - \left(\frac{1-\sqrt{5}}{2}\right)^k}{\sqrt{5}} + \frac{\left(\frac{1+\sqrt{5}}{2}\right)^{k-1} - \left(\frac{1-\sqrt{5}}{2}\right)^{k-1}}{\sqrt{5}} = \\&= \frac{\left(\frac{1+\sqrt{5}}{2} + 1\right) \left(\frac{1+\sqrt{5}}{2}\right)^{k-1} - \left(\frac{1-\sqrt{5}}{2} + 1\right) \left(\frac{1-\sqrt{5}}{2}\right)^{k-1}}{\sqrt{5}} = \\&= \frac{\left(\frac{3+\sqrt{5}}{2}\right) \left(\frac{1+\sqrt{5}}{2}\right)^{k-1} - \left(\frac{3-\sqrt{5}}{2}\right) \left(\frac{1-\sqrt{5}}{2}\right)^{k-1}}{\sqrt{5}} =\end{aligned}$$

# Induction

$$\begin{aligned}f_{k+1} &= \frac{\left(\frac{3+\sqrt{5}}{2}\right) \left(\frac{1+\sqrt{5}}{2}\right)^{k-1} - \left(\frac{3-\sqrt{5}}{2}\right) \left(\frac{1-\sqrt{5}}{2}\right)^{k-1}}{\sqrt{5}} = \\&= \frac{\left(\frac{1+\sqrt{5}}{2}\right)^2 \left(\frac{1+\sqrt{5}}{2}\right)^{k-1} - \left(\frac{1-\sqrt{5}}{2}\right)^2 \left(\frac{1-\sqrt{5}}{2}\right)^{k-1}}{\sqrt{5}} = \\&= \frac{\left(\frac{1+\sqrt{5}}{2}\right)^{k+1} - \left(\frac{1-\sqrt{5}}{2}\right)^{k+1}}{\sqrt{5}}\end{aligned}$$

$$\left(\frac{1+\sqrt{5}}{2}\right)^2 = \frac{1+2\sqrt{5}+5}{4} = \frac{3+\sqrt{5}}{2} \quad \left(\frac{1-\sqrt{5}}{2}\right)^2 = \frac{1-2\sqrt{5}+5}{4} = \frac{3-\sqrt{5}}{2}$$

# Logic

## Truth Table

| $P$ | $\neg P$ |
|-----|----------|
|-----|----------|

|   |   |
|---|---|
| 0 | 1 |
| 1 | 0 |

| $P_1$ | $P_2$ | $P_1 \& P_2$ | $P_1 \vee P_2$ | $P_1 \rightarrow P_2$ | $P_1 \leftrightarrow P_2$ |
|-------|-------|--------------|----------------|-----------------------|---------------------------|
| 0     | 0     | 0            | 0              | 1                     | 1                         |
| 0     | 1     | 0            | 1              | 1                     | 0                         |
| 1     | 0     | 0            | 1              | 0                     | 0                         |
| 1     | 1     | 1            | 1              | 1                     | 1                         |

## Logic

| $P_1$ | $P_2$ | $P_1 \& P_2$ | $P_1 \vee P_2$ | $P_1 \rightarrow P_2$ | $P_1 \leftrightarrow P_2$ |
|-------|-------|--------------|----------------|-----------------------|---------------------------|
| 0     | 0     | 0            | 0              | 1                     | 1                         |
| 0     | 1     | 0            | 1              | 1                     | 0                         |
| 1     | 0     | 0            | 1              | 0                     | 0                         |
| 1     | 1     | 1            | 1              | 1                     | 1                         |

$$0_{10} = 00_2$$

$$1_{10} = 01_2$$

$$2_{10} = 10_2$$

$$3_{10} = 11_2$$

## Logic

| $P_1$ | $P_2$ | $P_1 \& P_2$ | $P_1 \vee P_2$ | $P_1 \rightarrow P_2$ | $P_1 \leftrightarrow P_2$ |
|-------|-------|--------------|----------------|-----------------------|---------------------------|
| 0     | 0     | 0            | 0              | 1                     | 1                         |
| 0     | 1     | 0            | 1              | 1                     | 0                         |
| 1     | 0     | 0            | 1              | 0                     | 0                         |
| 1     | 1     | 1            | 1              | 1                     | 1                         |

$$(\neg a \rightarrow b) \leftrightarrow (a \& b)$$

| $a$ | $b$ | $\neg a$ | $\neg a \rightarrow b$ | $a \& b$ | $(\neg a \rightarrow b) \leftrightarrow (a \& b)$ |
|-----|-----|----------|------------------------|----------|---|
| 0   | 0   | 1        | 0                      | 0        | 1   |
| 0   | 1   | 1        | 1                      | 0        | 0   |
| 1   | 0   | 0        | 1                      | 0        | 0   |
| 1   | 1   | 0        | 1                      | 1        | 1   |

## Logic

$$(\neg a \rightarrow b) \leftrightarrow (\neg b \& c)$$

| $a$ | $b$ | $c$ | $\neg a$ | $\neg b$ | $(\neg a \rightarrow b)$ | $\neg b \& c$ | $(\neg a \rightarrow b) \leftrightarrow (\neg b \& c)$ |
|-----|-----|-----|----------|----------|--------------------------|---------------|--|
| 0   | 0   | 0   | 1        | 1        | 0                        | 0             | 1  |
| 0   | 0   | 1   | 1        | 1        | 0                        | 1             | 0  |
| 0   | 1   | 0   | 1        | 0        | 1                        | 0             | 0  |
| 0   | 1   | 1   | 1        | 0        | 1                        | 0             | 0  |
| 1   | 0   | 0   | 0        | 1        | 1                        | 0             | 0  |
| 1   | 0   | 1   | 0        | 1        | 1                        | 1             | 1  |
| 1   | 1   | 0   | 0        | 0        | 1                        | 0             | 0  |
| 1   | 1   | 1   | 0        | 0        | 1                        | 0             | 0  |

# Logic

## Truth Table

$$0_{10} = 000_2$$

$$1_{10} = 001_2$$

$$2_{10} = 010_2$$

$$3_{10} = 011_2$$

$$4_{10} = 100_2$$

$$5_{10} = 101_2$$

$$6_{10} = 110_2$$

$$7_{10} = 111_2$$

For  $n$  variables, the truth table contains  $2^n$  rows.

# Logic

Let  $n = 100$ .

How long does it take to compute on a computer?



# Logic

Let  $n = 100$ .

How long does it take to compute on a computer?

1 GHz = 1.000.000.000 Hertz =  $10^9$  elem.operations/sec.

10 GHz = 10.000.000.000 Hertz =  $10^{10}$  elem.operations/sec.

# Logic

Let  $n = 100$ .

How long does it take to compute on a computer?

$$1 \text{ GHz} = 1.000.000.000 \text{ Hertz} = 10^9 \text{ elem.operations/sec.}$$

$$10 \text{ GHz} = 10.000.000.000 \text{ Hertz} = 10^{10} \text{ elem.operations/sec.}$$

$$2^{100} = (2^{10})^{10} = (1024)^{10} > (1000)^{10} = (10^3)^{10} = 10^{30}$$

# Logic

$$n = 100$$

$$2^{100} > 10^{30}$$

$$\text{Seconds} = \frac{10^{30}}{10^{10}} = 10^{30-10} = 10^{20}$$

$$60 \times 60 \times 24 = 86400 \text{ seconds in each day} < 100000 = 10^5$$

$$\text{Days} = \frac{10^{20}}{10^5} = 10^{20-5} = 10^{15}$$

# Logic

$$n = 100$$

$$\text{Days} = \frac{10^{20}}{10^5} = 10^{20-5} = 10^{15}$$

$$365 < 500$$

$$\text{Years} = \frac{10^{15}}{500} = 2 \times \frac{10^{15}}{10^3} = 2 \times 10^{15-3} = 2 \times 10^{12} =$$

$$2.000.000.000.000 \text{ years}$$

## Logic

$$A \leftrightarrow B = (A \rightarrow B) \& (B \rightarrow A)$$

$$A \rightarrow B = \neg A \vee B$$

$$\begin{aligned} A \leftrightarrow B &= (A \rightarrow B) \& (B \rightarrow A) = \\ &= (\neg A \vee B) \& (\neg B \vee A) = (\neg A \vee B) \& (A \vee \neg B) \end{aligned}$$

## Logic

$$A \& (B \vee C) = (A \& B) \vee (A \& C)$$

$$A \& \neg A = 0$$

$$\begin{aligned} A \leftrightarrow B &= (\neg A \vee B) \& (A \vee \neg B) = \\ &= (\neg A \& A) \vee (\neg A \& \neg B) \vee (B \& A) \vee (B \& \neg B) = \\ &= (A \& B) \vee (\neg A \& \neg B) \end{aligned}$$

# Logic

$$A \leftrightarrow B = (\neg A \vee B) \& (A \vee \neg B) = (A \& B) \vee (\neg A \& \neg B)$$

$$A \rightarrow B = \neg A \vee B$$

$$\neg(\neg a) = a$$

## Example

$$\begin{aligned}(\neg a \rightarrow b) \leftrightarrow (\neg b \& c) &= (\neg(\neg a) \vee b) \leftrightarrow (\neg b \& c) = \\ &= (a \vee b) \leftrightarrow (\neg b \& c) =\end{aligned}$$

# Logic

$$A \leftrightarrow B = (\neg A \vee B) \& (A \vee \neg B) = (A \& B) \vee (\neg A \& \neg B)$$

$$\neg(A \vee B) = \neg A \& \neg B \quad \neg(A \& B) = \neg A \vee \neg B$$

## Example

$$\begin{aligned} (\neg a \rightarrow b) \leftrightarrow (\neg b \& c) &= (a \vee b) \leftrightarrow (\neg b \& c) = \\ &= [(a \vee b) \& (\neg b \& c)] \vee [\neg(a \vee b) \& \neg(\neg b \& c)] = \\ &= [(a \vee b) \& (\neg b \& c)] \vee [(\neg a \& \neg b) \& (b \vee \neg c)] = \\ &= (a \& \neg b \& c) \vee (b \& \neg b \& c) \vee (\neg a \& \neg b \& b) \vee (\neg a \& \neg b \& c) = \\ &= (a \& \neg b \& c) \vee (\neg a \& \neg b \& c) \end{aligned}$$



# DNF/CNF

## Algorithm

1. Dispose of  $\rightarrow$  and  $\leftrightarrow$ 
  - $A \leftrightarrow B = (\neg A \vee B) \& (A \vee \neg B) = (A \& B) \vee (\neg A \& \neg B)$
  - $A \rightarrow B = \neg A \vee B$
2. Use De Morgan's laws
  - $\neg(A \vee B) = \neg A \& \neg B$
  - $\neg(A \& B) = \neg A \vee \neg B$
3. Use Distributivity laws
  - $A \& (B \vee C) = (A \& B) \vee (A \& C)$  (for DNF)
  - $A \vee (B \& C) = (A \vee B) \& (A \vee C)$  (for CNF)

| $a$ | $b$ | $F$ |
|-----|-----|-----|
| 0   | 0   | 0   |
| 0   | 1   | 1   |
| 1   | 0   | 1   |
| 1   | 1   | 0   |

| $a$ | $b$ | $F$ |
|-----|-----|-----|
| 0   | 0   | 0   |
| 0   | 1   | 1   |
| 1   | 0   | 1   |
| 1   | 1   | 0   |

$$\text{DNF: } F = (\neg a \& b) \vee (a \& \neg b)$$

| $a$ | $b$ | $F$ |
|-----|-----|-----|
| 0   | 0   | 0   |
| 0   | 1   | 1   |
| 1   | 0   | 1   |
| 1   | 1   | 0   |

$$\text{DNF: } F = (\neg a \& b) \vee (a \& \neg b)$$

$$\text{CNF: } F = (a \vee b) \& (\neg a \vee \neg b)$$

Thank you for your attention!