# LINEAR PROGRAMMING PROJECT

**GROUP 2** REPORT FOR OPTIMIZATION F24 COURSE

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#### 1 Introduction

Initial problem is formulated as following:

$$\min_{x' \in \mathbb{R}^p} \|Ax' - y'\|_1$$
 s.t.  $0 \le x' \le 1$ 

where  $A \in \mathbb{R}^{m \times p}$  with  $m \geq p$  — message encoding matrix, y' — received encoded (noisy) message, x' — encoded initial message to be find.

#### 2 Notations

Notation	Meaning
$e_i$	unit vector with 1 at index $i$ and all other zeroes
$1_n$	vector of $n$ ones
$I_n$	identity matrix of size $n \times n$
$0_n$	vector of $n$ zeroes
$0_{m \times n}$	zero matrix of size $m \times n$
$x_i (\text{or } (Ax)_i)$	i-th component of vector $x$ (or $Ax$ )

## 3 Q1: Linear problem formulation

Initial problem (eq. (1)) is not linear as cost function  $||Ax'-y'||_1 = \sum_{i=1}^m |(Ax')_i-y_i'|$ , is not linear. However, this objective function is **piecewise linear convex** function. Therefore, each element  $|(Ax')_i-y_i'| = \max((Ax')_i-y_i',y_i'-(Ax')_i)$  can be substituted with new variable z' with the following additional constraints:  $z_i \geq (Ax')_i-y_i'$  and  $z_i \geq y_i'-(Ax')_i$ .

So the following problem is **linear** and is equivalent to the initial one:

$$\min_{x' \in \mathbb{R}^p, z \in \mathbb{R}^m} \sum_{i=1}^m z_i$$
s.t.  $x' \ge 0$ 

$$x' \le 1$$

$$z_i \ge (Ax')_i - y'_i, \ i = 1 \dots m$$

$$z_i \ge y'_i - (Ax')_i, \ i = 1 \dots m$$
(2)

## 4 Q2: Linear problem in standard form

For the easier and more evident deviation of standard form of Equation (2), linear problem will be firstly rewritten in geometric form, and only then — in standard. The obtained linear optimization problem in standard form will be equivalent to initial problem (eq. (1)).

#### 4.1 Geometric form

The equivalent **geometric** form of Equation (2) is

$$\min_{z' \in \mathbb{R}^{p+m}} c^T z'$$
s.t.
$$\begin{pmatrix}
I_p & \vdots & 0_{p \times m} \\
\vdots & \vdots & \vdots & \vdots \\
-I_p & \vdots & 0_{p \times m} \\
\vdots & \vdots & \vdots & \vdots \\
-A & \vdots & I_m \\
A' & b'
\end{pmatrix}$$

$$z' \ge \begin{pmatrix}
0_p \\
-1_p \\
-y' \\
y'
\end{pmatrix},$$

$$(3)$$

where  $c = \sum_{i=p+1}^{p+m} e_i \in \mathbb{R}^{(p+m)}, b' \in \mathbb{R}^{2p+2m}$  and  $A' \in \mathbb{R}^{(2p+2m)\times(p+m)}$ .

The first p components of z' correspond to the components of x', and the next m components correspond to the components of z from Equation (2). Rows and columns of A' representation in Equation (3) are separated in blocks for clarity: the vertical separation is for x' and z correspondingly, and the horizontal separations denote corresponding constraints from Equation (2).

#### 4.2 Standard form

Note that  $z' = (x', z)^T$  from Equation (3) is already non-negative, because  $x' \ge 0$  by problem definition and  $z \ge 0$  by construction<sup>1</sup>. Therefore, to convert Equation (3) to standard form, only introduction of slack variables is needed to get rid of inequality sign. The equivalent **standard** form of Equation (3) is

$$\min_{\tilde{x} \in \mathbb{R}^{2p+3m}} c^T \tilde{x}$$
s.t. 
$$\underbrace{\begin{pmatrix}
-I_p & 0_{p \times m} & -S^{1,p} \\
-A & I_m & -S^{p+1,p+m} \\
A & I_m & -S^{p+m+1,p+2m}
\end{pmatrix}}_{\tilde{x} > 0,} z' = \underbrace{\begin{pmatrix}
-1_p \\
-y' \\
y'
\end{pmatrix}}_{b'}, \tag{4}$$

where  $c=\sum_{i=p+1}^{p+m}e_i\in\mathbb{R}^{(2p+3m)},$   $b'\in\mathbb{R}^{p+2m},$   $A'\in\mathbb{R}^{(p+2m)\times(2p+3m)},$  and  $S^{a,b}$  — slack variable matrix of size  $(b-a+1)\times(p+2m)$  with rows  $S_i^{a,b}=e_{a+i-1}$ , which represents necessary slack variables.

The first p components of  $\tilde{x}$  correspond to the components of x', the next m components correspond to the components of z from Equation (2) and the last (p+2m) components correspond to slack variables s. Rows and columns of A' representation in Equation (4) are again separated in blocks for clarity: the vertical separations are for x', z and s correspondingly, and the horizontal separations are related to corresponding constraints from Equation (3) (except first one, as non-negativity in standard from is separate constraint).

<sup>&</sup>lt;sup>1</sup>Intuitively, z substitutes the absolute value, so is non-negative. Formally, from Equation (2),  $z \ge t$  and  $z \ge -t$  for some t. So if  $t \ge 0$ , then  $z \ge t \ge 0$ , and if  $t \le 0$  then  $z \ge -t \ge 0$