

2.1

- A* under consistent and admissible heuristic will output the minimum cost path.
- BFS outputs the path that has the lowest number of edges regardless of edge weights
- BFS outputs N0-N5.
- This means A* should also return N0-N5.
- This means edge weight N0-N5 should be lesser than the second lowest cost path (N0-N1-N2-N5).
- Edge weight(N0-N5) < 1 + 2 + 3
- Edge weight(N0-N5) in (0, 6)

Correct option: E

2.2

Given $h(N6) = 7$

Consistency implies

$$h(N5) \leq c(N5-N6) + h(N6)$$

$$h(N5) \leq 17$$

$$h(N4) \leq c(N4-N5) + h(N5)$$

$$h(N4) \leq 10 + h(N5)$$

$$h(N4) \leq 27$$

Given h is non-negative

$$h(N4) \geq 0$$

$h(N4)$ is in $[0, 27]$

Does admissibility hold through the interval?

$h^*(N4) = 30$ so admissibility holds throughout the interval $[0, 27]$

Can consistency hold throughout the interval $[0, 27]$?

TLDR; yes

Between any two neighboring nodes p and q

$$h(p) \leq c(p-q) + h(q)$$

$$h(p) \leq 10 + h(q) \quad \text{--- eq1}$$

And also, in the other direction

$$h(q) \leq c(p-q) + h(p)$$

$$h(q) \leq 10 + h(p) \quad \text{--- eq2}$$

CASE 1: $h(N4)$ in $[0, 17]$

Lets us say $h(N4) = x$, where x in $[0, 17]$

We design a heuristic h such that

- $h(N1) = h(N2) = h(N4) = h(N5) = h(S) = x$
- $h(N3) = h(N6) = 7$

- $h(T) = 0$

Eq1 and eq2 hold for all pairs of neighbors with this designed h.

CASE 2: $h(N4)$ in $[17, 27]$

Let us say $h(N4) = x$, where x in $[17, 27]$

We design a heuristic h such that

- $h(N1) = h(N4) = h(S) = x$
- $h(N2) = h(N5) = 17$
- $h(N3) = h(N6) = 7$
- $h(T) = 0$

Eq1 and eq2 hold for all pairs of neighbors with this designed h.

Correct option: A

2.3

```
visited[:] = false; //all unvisited nodes
dfs(s, visited, [])

dfs(node, visited, current_path):
    if node is visited:
        return
    visited[node] = true

    if node is target:
        print current_path
        exit

    for nbr in all_neighbors(node): //any neighbor could be expanded based
on tie breaking
        if nbr is not visited:
            current_path.append(nbr)
            dfs(nbr, visited, current_path)
            current_path.pop_back(nbr)
```

Based on the above code and considering there can be various tie breaking strategies. Any path from source to target can be formed by some tie-breaking scheme. For example, Starting with S

Then unvisited N1 neighbor of S is visited

Then unvisited N2 neighbor of N1 is visited

Then unvisited N4 neighbor of N2 is visited

Finally, T target node is visited for the first time => print the path S-N1-N2-N4-T (option D)

Correct options: B, C, D, E