

1.11:

$g_1(n) = n^2$  for even  $n \geq 0$ ,  $g_2(n) = n$  for  $0 \leq n \leq 100$ ,  $g_1(n)/g_2(n) = n^2/n = n \rightarrow \infty$  as  $n \rightarrow \infty$   
So,  $g_1(n)$  is  $\Omega(g_2(n))$  for even  $n$  in  $0 \leq n \leq 100$   
 $g_2(n) = n^3$  for  $n > 100$ ,  $g_1(n)/g_2(n) = n^2/n^3 = 1/n \rightarrow 0$  as  $n \rightarrow \infty$   
So,  $g_1(n)$  is  $O(g_2(n))$  for even  $n$  in  $n > 100$   
 $g_3(n) = n^{2.5}$ ,  $g_1(n)/g_3(n) = n^2/n^{2.5} = 1/n^{0.5} \rightarrow 0$  as  $n \rightarrow \infty$   
So,  $g_1(n)$  is  $O(g_3(n))$  for even  $n \geq 0$   
 $g_1(n) = n^3$  for odd  $n \geq 1$ ,  $g_2(n) = n$  for  $0 \leq n \leq 100$ ,  $g_1(n)/g_2(n) = n^3/n = n^2 \rightarrow \infty$  as  $n \rightarrow \infty$   
So,  $g_1(n)$  is  $\Omega(g_2(n))$  for odd  $n$  in  $0 \leq n \leq 100$   
 $g_2(n) = n^3$  for  $n > 100$ ,  $g_1(n)/g_2(n) = n^3/n^3 = 1 \rightarrow 1$  as  $n \rightarrow \infty$   
So,  $g_1(n)$  is  $O(g_2(n))$  for odd  $n$  in  $n > 100$   
 $g_3(n) = n^{2.5}$ ,  $g_1(n)/g_3(n) = n^3/n^{2.5} = n^{0.5} \rightarrow \infty$  as  $n \rightarrow \infty$   
So,  $g_1(n)$  is  $\Omega(g_3(n))$  for odd  $n \geq 1$   
 $g_2(n) = n$  for  $0 \leq n \leq 100$ ,  $g_1(n) = n^2$  for even  $n \geq 0$ ,  $g_2(n)/g_1(n) = n/n^2 = 1/n \rightarrow 0$  as  $n \rightarrow \infty$   
So,  $g_2(n)$  is  $O(g_1(n))$  for even  $n$  in  $0 \leq n \leq 100$   
 $g_1(n) = n^3$  for odd  $n \geq 1$ ,  $g_2(n)/g_1(n) = n/n^3 = 1/n^2 \rightarrow 0$  as  $n \rightarrow \infty$   
So,  $g_2(n)$  is  $O(g_1(n))$  for odd  $n$  in  $0 \leq n \leq 100$   
 $g_3(n) = n^{2.5}$ ,  $g_2(n)/g_3(n) = n/n^{2.5} = 1/n^{1.5} \rightarrow 0$  as  $n \rightarrow \infty$   
So,  $g_2(n)$  is  $O(g_3(n))$  for  $0 \leq n \leq 100$   
 $g_2(n) = n^3$  for  $n > 100$ ,  $g_1(n) = n^2$  for even  $n \geq 0$ ,  $g_2(n)/g_1(n) = n^3/n^2 = n \rightarrow \infty$  as  $n \rightarrow \infty$   
So,  $g_2(n)$  is  $\Omega(g_1(n))$  for even  $n$  in  $0 \leq n \leq 100$   
 $g_1(n) = n^3$  for odd  $n \geq 1$ ,  $g_2(n)/g_1(n) = n^3/n^3 = 1 \rightarrow 1$  as  $n \rightarrow \infty$   
So,  $g_2(n)$  is  $O(g_1(n))$  for odd  $n$  in  $0 \leq n \leq 100$   
 $g_3(n) = n^{2.5}$ ,  $g_2(n)/g_3(n) = n^3/n^{2.5} = n^{0.5} \rightarrow \infty$  as  $n \rightarrow \infty$   
So,  $g_2(n)$  is  $\Omega(g_3(n))$  for  $n > 100$

$g_3(n)$  is neither  $O(g_1(n))$  nor  $\Omega(g_1(n))$

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1.13:

a.  $17/1 \rightarrow 17$ , So 17 is  $O(1)$

b.  $n(n-1)/2 / n^2 = (n^2-n) / (2n^2) = 1/2 - 1/(2n) \rightarrow 1/2$ , So  $n(n-1)/2$  is  $O(n^2)$

c.  $\max(n^3, 10n^2) / n^3 = \max(1, 10/n) \rightarrow \max(1, 0) \rightarrow 1$ , So  $\max(n^3, 10n^2)$  is  $O(n^3)$

d. If  $k > 0$ ,  $\sum_{i=1}^n i^k = 1^k + 2^k + \dots + n^k \rightarrow \infty$ ,  $\sum_{i=1}^n i^k / n^{k+1} \rightarrow 0$ , So  $\sum_{i=1}^n i^k$  is  $\Omega(n^{k+1})$

If  $k < 0$ ,  $\sum_{i=1}^n i^k = 1^k + 2^k + \dots + n^k \rightarrow \text{some constant} > 1^k$ ,  $\sum_{i=1}^n i^k / n^{k+1} \rightarrow 0$ , So  $\sum_{i=1}^n i^k$  is  $O(n^{k+1})$

e.  $k$  polynomial  $P(n) = a_k n^k + a_{k-1} n^{k-1} + \dots + a_0$

If  $k > 0$ ,  $(a_k n^k + a_{k-1} n^{k-1} + \dots + a_0) / n^k = a_k + a_{k-1}/n + a_{k-2}/n^2 \dots + a_0/n^k \rightarrow a_k$  as  $n \rightarrow \infty$ , So  $P(n)$  is  $O(n^k)$

If  $k < 0$ ,  $(a_k n^k + a_{k-1} n^{k-1} + \dots + a_0) / n^k = a_k + a_{k-1}n + a_{k-2}n^2 \dots + a_0 n^k \rightarrow \infty$  as  $n \rightarrow \infty$ , So  $P(n)$  is  $\Omega(n^k)$

1.14:

a is true. Since  $T_1(n)$  is  $\Omega(f(n))$ ,  $T_1(n) \geq c_1 * f(n)$  for  $n > n_0$  and  $T_2(n)$  is  $\Omega(g(n))$ ,  $T_2(n) \geq c_2 * g(n)$ .

$$T_1(n) + T_2(n) \geq c_1 * f(n) + c_2 * g(n) \rightarrow c * (f(n) + g(n))$$

Since  $f(n) + g(n) = \max(f(n), g(n))$  by the book.

Therefore  $T_1(n) + T_2(n) \geq c * (f(n) + g(n))$  is equivalent to  $T_1(n) + T_2(n)$  is  $\Omega(\max(f(n), g(n)))$ .

b is false. Because either  $T_1(n)$ ,  $T_2(n)$  or even both of them are in the range  $(0,1)$ , then the multiplication will be  $T_1(n) * T_2(n) \leq f(n) * g(n)$ , so it is not always true.

1.17:

$n = 2, 4, 8, 16 \dots$

$n = 2$	count = 0
$n = 4$	count = 1
$n = 8$	count = 2
$n = 16$ $\cdot$ $\cdot$ $\cdot$ $n = 2^{\text{count} + 1}$	count = 3 $\cdot$ $\cdot$ $\cdot$ count = $\log_2 n - 1$