

1.11:

$g_1(n) = n^2$ for even $n \geq 0$ , $g_2(n) = n$ for $0 \leq n \leq 100$ , $g_3(n) = n^{2.5}$ ,	$g_1(n)/g_2(n) = n^2/n = n \rightarrow \infty$ as $n \rightarrow \infty$ So, $g_1(n)$ is $\Omega(g_2(n))$ for even $n$ in $0 \leq n \leq 100$
$g_2(n) = n^3$ for $n > 100$ ,	$g_1(n)/g_2(n) = n^2/n^3 = 1/n \rightarrow 0$ as $n \rightarrow \infty$ So, $g_1(n)$ is $O(g_2(n))$ for even $n$ in $n > 100$
$g_3(n) = n^{2.5}$ ,	$g_1(n)/g_3(n) = n^2/n^{2.5} = 1/n^{0.5} \rightarrow 0$ as $n \rightarrow \infty$ So, $g_1(n)$ is $O(g_3(n))$ for even $n \geq 0$
$g_1(n) = n^3$ for odd $n \geq 1$ , $g_2(n) = n$ for $0 \leq n \leq 100$ ,	$g_1(n)/g_2(n) = n^3/n = n^2 \rightarrow \infty$ as $n \rightarrow \infty$ So, $g_1(n)$ is $\Omega(g_2(n))$ for odd $n$ in $0 \leq n \leq 100$
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$g_2(n) = n$ for $0 \leq n \leq 100$ , $g_1(n) = n^2$ for even $n \geq 0$ ,	$g_2(n)/g_1(n) = n/n^2 = 1/n \rightarrow 0$ as $n \rightarrow \infty$ So, $g_2(n)$ is $O(g_1(n))$ for even $n$ in $0 \leq n \leq 100$
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$g_3(n) = n^{2.5}$ ,	$g_2(n)/g_3(n) = n^3/n^{2.5} = n^{0.5} \rightarrow \infty$ as $n \rightarrow \infty$ So, $g_2(n)$ is $\Omega(g_3(n))$ for $n > 100$
$g_3(n)$ is neither $O(g_1(n))$ nor $\Omega(g_1(n))$	
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1.13:

- a.  $17/1 \rightarrow 17$ , So 17 is  $O(1)$
- b.  $n(n-1)/2 / n^2 = (n^2-n) / (2n^2) = \frac{1}{2} - \frac{1}{2n} \rightarrow \frac{1}{2}$ , So  $n(n-1)/2$  is  $O(n^2)$
- c.  $\max(n^3, 10n^2) / n^3 = \max(1, 10/n) \rightarrow \max(1, 0) \rightarrow 1$ , So  $\max(n^3, 10n^2)$  is  $O(n^3)$
- d. If  $k > 0$ ,  $\sum_{i=1}^n i^k = 1^k + 2^k + \dots + n^k \rightarrow \infty$ ,  $\sum_{i=1}^n i^k / n^{k+1} \rightarrow \infty$ , So  $\sum_{i=1}^n i^k$  is  $\Omega(n^{k+1})$   
If  $k < 0$ ,  $\sum_{i=1}^n i^k = 1^k + 2^k + \dots + n^k \rightarrow$  some constant  $> 1^k$ ,  $\sum_{i=1}^n i^k / n^{k+1} \rightarrow$  constant, So  $\sum_{i=1}^n i^k$  is  $O(n^{k+1})$
- e. k polynomial  $P(n) = a_k n^k + a_{k-1} n^{k-1} + \dots + a_0$   
If  $k > 0$ ,  $(a_k n^k + a_{k-1} n^{k-1} + \dots + a_0) / n^k = a_k + a_{k-1}/n + a_{k-2}/n^2 \dots + a_0/n^k \rightarrow a_k$  as  $n \rightarrow \infty$ , So  $P(n)$  is  $O(n^k)$   
If  $k < 0$ ,  $(a_k n^k + a_{k-1} n^{k-1} + \dots + a_0) / n^k = a_k + a_{k-1} n + a_{k-2} n^2 \dots + a_0 n^k \rightarrow \infty$  as  $n \rightarrow \infty$ , So  $P(n)$  is  $\Omega(n^k)$

1.14:

a is true. Since  $T_1(n)$  is  $\Omega(f(n))$ ,  $T_1(n) \geq c_1 * f(n)$  for  $n > n_0$  and  $T_2(n)$  is  $\Omega(g(n))$ ,  $T_2(n) \geq c_2 * g(n)$ .  
 $T_1(n) + T_2(n) \geq c_1 * f(n) + c_2 * g(n) \rightarrow c * (f(n) + g(n))$

Since  $f(n) + g(n) = \max(f(n), g(n))$  by the book.

Therefore  $T_1(n) + T_2(n) \geq c * (f(n) + g(n))$  is equivalent to  $T_1(n) + T_2(n)$  is  $\Omega(\max(f(n), g(n)))$ .

b is false. Because either  $T_1(n)$ ,  $T_2(n)$  or even both of them are in the range  $(0, 1)$ , then the multiplication will be  $T_1(n) * T_2(n) \leq f(n) * g(n)$ , so it is not always true.

1.17:

$n = 2, 4, 8, 16\dots$

$n = 2$	count = 0
$n = 4$	count = 1
$n = 8$	count = 2
$n = 16$	count = 3
.	.
.	.
$n = 2^{\text{count} + 1}$	count = $\log_2 n - 1$