# Programming LSQUIC

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#### Presentation outline

- Introduction
- LSQUIC history, features, and architecture
- Using API
  - Objects: engines, connections, streams
  - Sending and receiving packets
  - Instantiation and callbacks
- Bonus section, time permitting

## Example Program: tut.c

- You can follow along
- Snippets of tut.c are used as examples
- Compile the code while I talk about architecture

```
git clone https://github.com/litespeedtech/lsquic-tutorial
# This pulls in LSQUIC and BoringSSL
git submodule update --init --recursive
# Tested on Ubuntu
cmake .
make
```

## What is QUIC (keep it short)

- Why: Ossification
- Began as HTTP/2 over UDP
- Who: First Google, now everyone:
  - LiteSpeed Technologies, Microsoft, Apple, Facebook, Akamai, and others
- General-purpose transport protocol
- Killer feature: HTTP/3 (end of 2020)
- Experimental: datagrams, multipath, DNS-over-QUIC, NETCONF-over-QUIC, and so on
- Future: bright

## Which QUIC do you mean?

- Google QUIC vs IETF QUIC
- gQUIC? iQUIC? We all QUIC for yogurt?
- Google QUIC is on the way out
- In the slides that follow, "QUIC" means "IETF QUIC"

### Introducing LSQUIC

- Began as proprietary, then open-sourced
- Vanilla C
- Minimal dependencies
- Performance
- Robustness
- Flexibility

## History

- 2016. Goal: add Google QUIC support to LiteSpeed Web Server.
- 2017. gQUIC support is shipped (Q035); LSQUIC is released on GitHub (client only).
- 2018. IETF QUIC work begins. LSQUIC is the first functional HTTP/3 server.
- 2019. LSQUIC 2.0 is released on GitHub (including the server bits). HTTP/3 support is shipped.
- 2020. HTTP/3 is an RFC. (One can hope.)

#### Features

- Latest (Draft-27, Draft-28, and Draft-29) IETF QUIC and HTTP/3 support, including
  - ECN, spin bits, path migration, NAT rebinding
  - Push promises, key updates
  - Several experimental extensions:
    - loss bits
    - timestamps
    - delayed ACKs
    - QUIC grease bit
- Google QUIC versions Q043, Q046, and Q050 (what Chrome currently uses)
- Many, many knobs to play with

### Architecture

- Bring your own event loop
- Bring your own networking
- Bring your own TLS context
- Scalable connection management

## Objects

- Engine
- Connection
- Stream

## Engine

- Client mode or server mode
- If need both, instantiate two
- HTTP mode

#### HTTP mode

- Single library for both QUIC and HTTP/3
- Hide HTTP logic: control streams, header compression, framing
- Identical interface for gQUIC and HTTP/3
- Historical or strategic?
- Optimization: write-through

### Connection

- Created, managed, and destroyed by engine
- Client initiates connection; object created before handshake is successful.
- Server: by the time user code gets the connection object, the handshake has already been completed
- Can have many streams during lifetime

#### Stream

- Belongs to a connection
- Bidirectional
- Usually corresponds to request/response exchange depending on the application protocol
- API tries to mimic socket
  - But one can only take it so far

#### Include files

A single include file, contains all the necessary LSQUIC declarations:

```
#include "lsquic.h"
```

It pulls in auxiliary lsquic\_types.h

## Library initialization

Before the first engine object is instantiated, the library must be initialized.

This will initialize the crypto library, gQUIC server certificate cache, and, depending on the platform, monotonic timers.

If you plan to instantiate engines only in a single mode, client or server, you can omit the appropriate flag.

### Introducing tut.c

- Program to illustrate LSQUIC API use
- Contains both client and server code
- Echo service: client sends a line of text to server, the server returns the line, reversed
- Several examples that follow are excerpts from tut.c

### Running tutorial program

- Peruse online help: use the -h flag
- Running client or server: the server takes -c and -k arguments

```
# Server:
sh$ ./tut -c mycert-cert.pem -k mycert-key.pem ::0 12345
# Client:
sh$ ./tut ::1 12345 -L debug -f client.log
Hello!
!olleH
^D
```

## Engine constructor

- Server or client
- HTTP mode

### Specifying engine callbacks

Pass pointer to struct lsquic\_engine\_api

```
/* Minimal configuration */
struct lsquic_engine_api engine_api = {
    .ea_packets_out = send_packets_out,
    .ea_packets_out_ctx = sender_ctx,
    .ea_stream_if = &stream_callbacks,
    .ea_stream_if_ctx = &some_context,
    .ea_get_ssl_ctx = get_ssl_ctx, /* Server only */
};
```

## Excerpt from tut.c

#### Packets in

- Single function to feed packets to engine instance
- Specify: datagram, peer and local addresses, ECN value
- peer\_ctx is associated with peer address: it is passed to send packets function.

```
/* 0: processed by real connection
 * 1: h.a.n.d.l.e.d.
 * -1: error: invalid arguments, malloc failure
 */
int
lsquic engine packet in (lsquic engine t *,
    const unsigned char *udp payload, size t sz,
    const struct sockaddr *sa local,
    const struct sockaddr *sa peer,
    void *peer ctx, int ecn);
```

## Why specify local address

- Becomes source address on outgoing packets
  - Important in multihomed configuration
- Path change detection
  - QUIC sends special frames to validate path

#### Packets out

- Callback
- Called during connection processing (explicit call by user)

```
/* Returns number of packets successfully sent out or -1 on error.
 *
 * If not all packets could be sent, call
 * lsquic engine send unsent packets() when can send again.
 */
typedef int (*lsquic packets out f)(
    void
                                   *packets out ctx.
    const struct lsquic_out_spec
                                  *out_spec,
    unsigned
                                   n packets_out
);
```

#### When an error occurs

- errno is examined
- EAGAIN (or EWOULDBLOCK) means retry later; engine enters the "can't send packets" mode
- Call lsquic\_engine\_send\_unsent\_packets() when sending is possible again
- EMSGSIZE means the packet is too large. This happens when MTU probes are sent. The engine retries sending without the offending packet.
- Other errno values cause immediate termination of corresponding connection.

### Outgoing packet specification

• Why iovec? • UDP datagram can contain more than one QUIC packet • Packet coalescing struct lsquic\_out\_spec struct iovec \*iov: size t iovlen; const struct sockaddr \*local sa: const struct sockaddr \*dest sa: void \*peer ctx; ecn: /\* 0 - 3: see RFC 3168 \*/ int

};

## Packets out example

```
static int
my packets out (void *ctx, const struct lsquic out spec *specs,
                                            unsigned n specs) {
    struct msghdr msg; memset(&msg, 0, sizeof(msg));
    unsigned n;
    for (n = 0; n < n \text{ specs}; ++n) {
        msg.msg_name
                     = (void *) specs[n].dest_sa;
        msg.msg_namelen = sizeof(struct sockaddr in);
        msg.msg_iov = specs[n].iov;
        msg.msg_iovlen = specs[n].iovlen;
        if (sendmsg((int) specs[n].peer_ctx, &msg, 0) < 0)</pre>
            break:
    return (int) n;
```

### When to process connections

- Connections are either tickable immediately or at some future point
- Future point may be a retransmission, idle, or some other alarm
- Only exception when idle timeout is disabled (on by default)
- Engine knows when to process connections next

```
/* Returns true if there are connections to be processed, in
  * which case `diff' is set to microseconds from current time.
  */
int
lsquic_engine_earliest_adv_tick (lsquic_engine_t *, int *diff);
```

## Example with event loop

```
/* Abbreviated, see full version in tut.c */
void tut process conns (struct tut *tut) {
    ev tstamp timeout;
    int diff:
    ev timer stop();
    lsquic engine process conns(engine);
    if (lsquic engine earliest adv tick(engine, &diff) {
        if (diff > 0)
            timeout = (ev_tstamp) diff / 1000000; /* To seconds */
        else
            timeout = 0.:
        ev timer init(timeout)
        ev timer start();
```

#### Tickable connection

- There are incoming packets
- A stream is both readable by the user code and the user code wants to read from it
- A stream is both writeable by the user code and the user code wants to write to it
- User has written to stream outside of on\_write() callbacks (that is allowed) and now there are packets ready to be sent
- A control frame needs to be sent out
- A stream needs to be serviced or created

## Required engine callbacks

- Callback to send packets
- Connection and stream callbacks
  - on\_new\_conn(), on\_read(), and so on
- Callback to get default TLS context (server only)

### Optional callbacks

- Certificate lookup by SNI (server only)
- Outgoing packet memory allocation
- Connection ID lifecycle: new, live, and old CIDs
- Shared memory hash
- HTTP header set processing

### Stream and connection callbacks

- Specified in struct lsquic\_stream\_if
- Mandatory callbacks:
  - on\_new\_conn() new connection is created
  - on\_conn\_closed()
  - on\_new\_stream()
  - on read()
  - on write()
  - on close()
- Optional callbacks:
  - on\_goaway\_received()
  - on\_new\_token() (client only)
  - on\_new\_coken() (chefit only
  - on\_hsk\_done() (client only)
  - on\_sess\_resume\_info() (client only)

#### On new connection

- Server: handshake successful; client: object created
- Chance to create custom per-connection context

```
/* Return pointer to per-connection context. OK to return NULL. */
static lsquic conn ctx t *
my on new conn (void *ea stream if ctx, lsquic conn t *conn)
    struct some_context *ctx = ea_stream_if_ctx;
    struct my_conn_ctx *my_ctx = my_ctx_new(ctx);
    if (ctx->is client)
        /* Need a stream to send request */
        lsquic conn make stream(conn):
    return (void *) my_ctx;
```

#### On new stream

- Depending on situation, register interest in reading or writing
  - Or just read or write
- Chance to create per-stream context

```
/* Return pointer to per-connection context. OK to return NULL. */
static lsquic stream ctx t *
my on new stream (void *ea stream if ctx, lsquic stream t *stream) {
    struct some context *ctx = ea stream if ctx:
    /* Associate some data with this stream: */
    struct my_stream_ctx *stream_ctx
                  = mv stream ctx new(ea stream if ctx);
    stream ctx->stream = stream;
    if (ctx->is_client)
        lsquic_stream_wantwrite(stream, 1);
    return (void *) stream ctx;
```

#### On read

• Read data – or collect error

```
static void
my_on_read (lsquic_stream_t *stream, lsquic_stream_ctx_t *h) {
    struct my stream ctx *my stream ctx = (void *) h;
    unsigned char buf[BUFSZ];
    ssize t nr = lsquic stream read(stream, buf, sizeof(buf));
    /* Do something with the data.... */
    if (nr == 0) /* EOF */ {
        lsquic stream shutdown(stream, 0);
        lsquic stream wantwrite(stream, 1): /* Want to reply */
```

#### On write

• If called, you should be able to write *some* bytes

#### On stream close

• After this, stream will be destroyed, drop all pointers to it

```
/* Made-up example */
static void
my_on_close (lsquic_stream_t *stream, lsquic_stream_ctx_t *h) {
    lsquic_conn_t *conn = lsquic_stream_conn(stream);
    struct my_conn_ctx *my_ctx = lsquic_conn_get_ctx(conn);
    if (!has_more_reqs_to_send(my_ctx)) /* For example */
        lsquic_conn_close(conn);
    free(h);
}
```

#### On stream close in tut.c

#### On connection close

• After this, connection will be destroyed, drop all pointers to it

```
static void
my on conn closed (lsquic conn t *conn) {
    struct my_conn_ctx *my_ctx = lsquic_conn_get_ctx(conn);
    struct some context *ctx = my ctx->some context;
    --ctx->n_conns;
    if (0 == ctx->n conn && (ctx->flags & CLOSING))
        exit event loop(ctx):
    free(my_ctx);
```

#### On connection close in tut.c

- This deletes last event from the event loop the event loop exits and the program terminates
- Note that the same thing happens if handshake fails

```
static void
tut_client_on_conn_closed (struct lsquic_conn *conn)
{
   struct tut *const tut = (void *) lsquic_conn_get_ctx(conn);
   LOG("client connection closed -- stop reading from socket");
   ev_io_stop(tut->tut_loop, &tut->tut_sock_w);
}
```

#### More about streams

- When writing to a stream, data is placed directly into packets
  - Except when it isn't
- Writes smaller than packet size are buffered
- Inside on\_read and on\_write callbacks, reading and writing succeeds
- Unless stream is reset, in which case reading from stream returns -1
  - This is done so that user can collect error

#### More stream functions

```
/* Flush any buffered data. This triggers packetizing even a single
 * byte into a separate frame.
 */
int
lsquic stream flush (lsquic stream t *);
/* Possible values for how are 0, 1, and 2. See shutdown(2). */
int
lsquic stream shutdown (lsquic stream t *, int how);
int
lsquic stream close (lsquic stream t *);
```

#### Stream return values and error codes

- Reading and writing interface modeled on read(2) and write(2)
- Including the use of errno
  - If no data to read, error is EWOULDBLOCK
  - If stream is closed, error is EBADF
  - If stream is reset, error is ECONNRESET
  - After this, error codes fit only if you squint very hard
    - EINVAL argument to shutdown is not 0, 1, or 2
    - EILSEQ cannot send HTTP payload before headers
    - EBADMSG sending HTTP headers is not allowed (several reasons)
- When lsquic\_stream\_read() returns 0, it means EOF
- lsquic\_stream\_write() returns 0 when flow control or congestion control limit is reached.

### More ways to read and write

- Scatter/gather
   Similar to readv(2) and writev(2)
   ssize\_t
   lsquic\_stream\_readv (lsquic\_stream\_t \*, const struct iovec \*,
- ssize\_t

```
lsquic_stream_writev (lsquic_stream_t *, const struct iovec *,
```

int iovcnt):

# Read using a callback

- Use for zero-copy stream processing
- lsquic\_stream\_read() and lsquic\_stream\_readv() are just wrappers
- Callback returns number of bytes processed
  - Pointer to user-supplied context;
  - Pointer to the data;
  - Data size (can be zero); and
  - Indicator whether the FIN follows the data.
- If callback returns 0 or value smaller than len, reading stops

```
ssize_t
lsquic_stream_readf (lsquic_stream_t *,
    size_t (*readf)(void *ctx, const unsigned char *, size_t len, int fin),
    void *ctx);
```

### Stream read: copy data

- The classic way to read
- Easy, but it copies data

```
static void /* v0 is the default. On command line: -b 0 */
tut client on read v0 (lsquic stream t *stream, lsquic stream ctx t *h)
 struct tut *tut = (struct tut *) h;
 unsigned char buf[3];
 ssize t nread = lsquic stream read(stream, buf, sizeof(buf));
 if (nread > 0)
      fwrite(buf, 1, nread, stdout);
     fflush(stdout):
/* --- 8< --- snin --- 8< --- */
```

#### Stream read take 2: use callback

```
static void /* On command line: -b 1 */
tut_client_on_read_v1 (lsquic_stream_t *stream, lsquic_stream_ctx t *h)
  struct tut *tut = (struct tut *) h:
  size t nread = lsquic stream readf(stream, tut client readf v1, NULL);
  if (nread == 0)
      LOG("read to end-of-stream: close and read from stdin again");
      lsquic stream shutdown(stream, 0);
      ev io start(tut->tut loop, &tut->tut u.c.stdin w);
/* --- 8< --- snip --- 8< --- */
```

#### Stream read take 2: the callback itself

- len may be arbitrary (and larger than packet size)
- Return value smaller than len to stop callback

```
static size t
tut client readf v1 (void *ctx, const unsigned char *data,
                                                   size t len, int fin)
    if (len)
        fwrite(data, 1, len, stdout):
        fflush(stdout):
    return len;
```

#### Stream read take 3: use FIN

• Save one on read() call struct client read v2 ctx { struct tut \*tut; lsquic stream t \*stream; }; static void /\* On command line: -b 2 \*/ tut client on read v2 (lsquic stream t \*stream, lsquic stream ctx t \*h) { struct tut \*tut = (struct tut \*) h: struct client\_read\_v2\_ctx v2ctx = { tut, stream, }; ssize t nread = lsquic stream readf(stream, tut client readf v2, &v2ctx): if (nread < 0)/\* ERROR \*/

#### Stream read take 3: the callback itself

```
static size_t
tut client readf v2 (void *ctx, const unsigned char *data,
                                               size t len, int fin) {
  struct client read v2 ctx *v2ctx = ctx;
  if (len)
    fwrite(data, 1, len, stdout);
  if (fin)
    fflush(stdout):
    LOG("read to end-of-stream: close and read from stdin again");
    lsquic stream shutdown(v2ctx->stream, 0);
    ev io start(v2ctx->tut->tut loop, &v2ctx->tut->tut u.c.stdin w):
  }
  return len;
```

#### Stream write take 1

```
static void /* v0 is the default. On command line: -w 0 */
tut server on write v0 (lsquic stream t *stream, lsquic stream ctx t *h)
  struct tut_server_stream_ctx *const tssc = (void *) h;
  ssize t nw = lsquic stream write(stream.
      tssc->tssc buf + tssc->tssc off. tssc->tssc sz - tssc->tssc off);
 if (nw > 0)
      tssc->tssc off += nw;
      if (tssc->tssc off == tssc->tssc sz)
          lsquic_stream_close(stream);
/* --- 8< --- snip --- 8< --- */
```

## Write using callbacks

```
struct lsquic reader {
  /* Return number of bytes written to buf */
  size t (*lsqr read) (void *lsqr ctx, void *buf, size t count);
 /* Return number of bytes remaining in the reader. */
  size t (*lsqr size) (void *lsqr ctx);
 void *lsgr ctx:
};
/* Return umber of butes written or −1 on error. */
ssize t
lsquic stream writef (lsquic stream t *, struct lsquic reader *);
```

#### Stream write take 2

- Useful when reading from external data source, such as file descriptor
- Write data directly into stream frame

```
static void /* On command line: -w 1 */
tut server on write v1 (lsquic stream t *stream, lsquic stream ctx t *h)
    struct tut_server_stream_ctx *const tssc = (void *) h;
    struct lsquic reader reader = { tssc read, tssc size, tssc, };
    ssize t nw = lsquic stream writef(stream, &reader);
    if (nw > 0 && tssc->tssc off == tssc->tssc sz)
        lsquic stream close(stream):
/* --- 8< --- snip --- 8< --- */
```

#### Reader size callback

• Return number of bytes remaining

```
static size_t
tssc_size (void *ctx)
{
   struct tut_server_stream_ctx *tssc = ctx;
   return tssc->tssc_sz - tssc->tssc_off;
}
```

#### Reader read callback

- count is calculated using tssc\_size()
- If larger than amount of remaining data, may indicate truncation

```
static size t
tssc read (void *ctx, void *buf, size t count)
  struct tut server stream ctx *tssc = ctx;
  if (count > tssc->tssc sz - tssc->tssc off)
    count = tssc->tssc sz - tssc->tssc off:
  memcpy(buf, tssc->tssc buf + tssc->tssc off, count);
  tssc->tssc off += count:
  return count;
```

### Client: making connection

```
lsquic conn t *
lsquic engine connect (lsquic engine t *,
      enum lsquic version, /* Set to N LSQVER for default */
      const struct sockaddr *local_sa,
      const struct sockaddr *peer sa,
     void *peer ctx.
      lsquic conn_ctx_t *conn_ctx,
      const char *hostname. /* Used for SNI */
     unsigned short base plpmtu, /* 0 means default */
      const unsigned char *sess resume, size t sess resume len,
      const unsigned char *token, size t token sz);
```

# Excerpt from tut.c

```
tut.tut u.c.conn = lsquic engine connect(
    tut.tut engine, N LSQVER,
    (struct sockaddr *) &tut.tut local sas, &addr.sa,
    (void *) (uintptr_t) tut.tut_sock_fd, /* Peer ctx */
    NULL, NULL, O, NULL, O, NULL, O);
if (!tut.tut u.c.conn)
    LOG("cannot create connection"):
    exit(EXIT FAILURE):
tut process conns(&tut);
```

## Specifying QUIC version

- There is no version negotiation version in QUIC... yet
- When passed to lsquic\_engine\_connect, N\_LSQVER means "let the engine pick the version"
  - The engine picks the highest it supports, so that's a good way to go

#### Server: additional callbacks

• SSL context and certificate callbacks

```
typedef struct ssl_ctx_st * (*lsquic_lookup_cert_f)(
    void *lsquic cert lookup ctx, const struct sockaddr *local,
    const char *sni):
struct lsquic engine api {
 lsquic lookup cert f ea lookup cert;
 biov
                        *ea cert lu ctx;
 struct ssl ctx st * (*ea get ssl ctx)(void *peer ctx);
 /* (Other members of the struct are not shown) */
};
```

# Engine settings

- Engine API's ea\_settings may be pointed to settings.
- struct lsquic\_engine\_settings.
- There are many settings (over 50).
- Do use lsquic\_engine\_init\_settings().

### Settings helper functions

```
/* Initialize `settings' to default values */
void
lsquic engine init settings (struct lsquic engine settings *,
 /* Bitmask of LSENG SERVER and LSENG HTTP */
                             unsigned lsquic engine flags);
/* Check settings for errors, return 0 on success, -1 on failure. */
int
lsquic engine check settings (const struct lsquic engine settings *,
                              unsigned lsquic engine flags,
                              /* Optional, can be NULL: */
                              char *err buf, size t err buf sz);
```

### Settings example in tut.c 1/2

```
case 'o': /* For example: -o version=h3-27 -o cc algo=2 */
 if (!settings initialized) {
   lsquic engine init settings(&settings,
                   cert_file || key_file ? LSENG_SERVER : 0);
   settings initialized = 1:
 /* ... */
 else if (0 == strncmp(optarg, "cc_algo=", val - optarg))
   settings.es cc algo = atoi(val);
```

# Settings example in tut.c 2/2

```
/* Check settings */
if (0 != lsquic engine_check_settings(&settings,
                tut.tut flags & TUT SERVER ? LSENG SERVER : 0,
                errbuf, sizeof(errbuf)))
  LOG("invalid settings: %s", errbuf);
  exit(EXIT FAILURE);
/* ... */
eapi.ea settings = &settings;
```

## Logging mechanism

• By default, log messages are thrown away

```
struct lsquic logger if {
 int (*log buf)(void *logger ctx, const char *buf, size t len);
}:
enum lsquic logger timestamp style { LLTS NONE, LLTS HHMMSSMS,
   LLTS YYYYMMDD HHMMSSMS, LLTS CHROMELIKE, LLTS HHMMSSUS,
   LLTS YYYYMMDD HHMMSSUS, N LLTS };
void lsquic logger init(const struct lsquic logger if *,
   void *logger ctx, enum lsquic logger timestamp style);
```

### Logging levels and modules

- Eight log levels
  - Debug, info, notice, warning, error, alert, emerg, crit.
  - Only first five are used; warning is the default
- Many modules
  - Event, engine, stream, connection, bbr, and many more
  - Refer to documentation

```
/* Set log level for all modules */
int lsquic_set_log_level (const char *log_level);
/* Set log level per module "event=debug" */
int lsquic_logger_lopt (const char *optarg);
```

# Logging in tut.c 1/2

```
• Check out -f, -1, and -L flags
• E.g. -f log.file (goes to stderr by default)
• E.g. -1 event=debug, stream=info or -L debug
static int
tut log buf (void *ctx, const char *buf, size t len) {
  FILE *out = ctx:
  fwrite(buf, 1, len, out);
  fflush(out);
  return 0:
static const struct lsquic logger if logger if = { tut log buf, };
lsquic logger init(&logger if, s log fh, LLTS HHMMSSUS);
```

# Logging in tut.c 2/2

```
case 'l': /* e.g. -l event=debug, cubic=info */
 if (0 != lsquic logger lopt(optarg)) {
     fprintf(stderr, "error processing -l option\n");
     exit(EXIT FAILURE);
 break:
case 'L': /* e.g. -L debug */
 if (0 != lsquic set log level(optarg)) {
     fprintf(stderr, "error processing -L option\n");
     exit(EXIT FAILURE);
 break;
```

#### Tools: Wireshark

- Wireshark supports IETF QUIC
- Need version 3.3 for Internet-Draft 29 support
- Export TLS secrets using ea\_keylog\_if
- Example: -G option in tut:

```
if (key_log_dir) {
  eapi.ea_keylog_if = &keylog_if;
  eapi.ea_keylog_ctx = (void *) key_log_dir;
}
```

# LSQUIC API for key logging

```
/* Secrets are logged per connection. Interface to open file (handle),
 * log lines, and close file.
struct lsquic keylog if {
    void * (*kli open) (void *keylog ctx, lsquic conn t *);
    void (*kli_log_line) (void *handle, const char *line);
    void (*kli_close) (void *handle);
}:
struct lsquic_engine_api {
  /* --- 8< --- snip --- 8< --- */
  const struct lsquic keylog if
                                      *ea keylog if;
 void
                                      *ea kevlog ctx:
};
```

# Key logging in tut.c

```
static struct lsquic keylog if keylog if = {
  .kli open = kevlog open,
  .kli_log_line = keylog_log_line,
  .kli_close = keylog_close,
};
static void
keylog log line (void *handle, const char *line)
 fputs(line, handle);
 fputs("\n", handle);
  fflush(handle);
```

#### Wireshark screenshot

```
No.
         Time
                                    Destination
                                               Protocol Length Info
                       Source
      16 2.570135
                       127.0.0.1
                                   127.0.0.1
                                               OUTC
                                                         83 Protected Pavload (KPO), DCID=e77ce754b2d32f6c, PKN: 7, ACK
      17 2.924101
                       127.0.0.1
                                   127.0.0.1
                                               OUIC
                                                         78 Protected Pavload (KP0), DCID=e77ce754b2d32f6c, PKN: 8, STREAM(0)
      18 2.925751
                       127.0.0.1
                                   127.0.0.1
                                               OUTC
                                                         90 Protected Pavload (KPO), DCID=be3ac0381d9049e8, PKN: 11, ACK, STREAM(...
      19 2.927297
                       127.0.0.1
                                   127.0.0.1
                                               OUTC
                                                         81 Protected Payload (KP0), DCID=e77ce754b2d32f6c, PKN: 9, ACK
                                                         97 Protected Pavload (KP0), DCID=be3ac0381d9049e8, PKN: 12, NCI
      20 2.942133
                       127.0.0.1
                                   127.0.0.1
                                               OUIC
                                                         81 Protected Payload (VPA) DCTD-e77ce754b2d32f6c DVN: 10 ACV
      21 2 044500
                       127 0 0 1 127 0 0 1
                                              OUTC
   V ACK
```

```
Frame Type: ACK (0x0000000000000000)
     Largest Acknowledged: 8
     ACK Delav: 46
     ACK Range Count: 0
     First ACK Range: 3

✓ TIME STAMP

     Frame Type: TIME STAMP (0x000000000000002f5)
     Time Stamp: 365556

✓ STREAM id=0 fin=1 off=0 len=7 uni=0

  ✓ Frame Type: STREAM (0x0000000000000000b)
        ..... 1 = Fin: True
        .... ..1. = Len(gth): True
        .... .0.. = Off(set): False
     Stream ID: 0
     Length: 7
     Stream Data: 216f6c6c65480a
```

```
0000 02 08 2e 00 03 42 f5 80 05 93 f4 0b 00 07 21 6f
```

#### Connection ID

```
#define MAX CID LEN 20
typedef struct lsquic cid
   uint fast8 t
                   len;
   union {
       uint8 t buf[MAX CID LEN];
       uint64 t id:
                   u cid:
#define idbuf u_cid.buf
} lsquic cid t;
#define LSQUIC_CIDS_EQ(a, b) ((a)->len == 8 ? \
    (b)->len == 8 && (a)->u cid.id == (b)->u cid.id : \
    (a)->len == (b)->len && 0 == memcmp((a)->idbuf, (b)->idbuf, (a)->len)
```

#### Get this-and-that API

```
const lsquic cid t *
lsquic conn id (const lsquic conn t *);
lsquic conn t *
lsquic stream conn (const lsquic stream t *);
lsquic engine t *
lsquic_conn_get_engine (lsquic_conn_t *);
int
lsquic conn get sockaddr (lsquic conn t *,
      const struct sockaddr **local, const struct sockaddr **peer);
```

### Stream priorities

- 1 through 256, where 1 is the highest priority
- Controls dispatch of on\_read() and on\_write() callbacks
- Not HTTP/3 priorities (those are coming later)

```
/* Set stream priority. Valid priority values are 1 through 256,
  * inclusive. Lower value means higher priority.
  */
int
lsquic_stream_set_priority (lsquic_stream_t *, unsigned priority);
/* Return current priority of the stream */
unsigned
lsquic stream priority (const lsquic stream t *);
```

### But wait, there is more!

- LSQUIC has more features that we did not cover
- Shared memory callbacks to support persistent crypto contexts and multi-process setups
- Callbacks to handle memory allocation for outgoing packets
- Callbacks to observe CID life cycle: new, live, and dead. Useful in multi-process setups
- Server certificate verification callbacks
- And more: auto-tuning; r/w callback no-progress detection; r/w callback dispatch: once vs loop; different congestion controls; optimistic ACK attack detection; stream fragmentation, reassembly, and commitment attack migitation
- Refer to documentation and to more involved example programs in the LSQUIC distribution
- https://lsquic.readthedocs.io/

Bonus Section #1

Linux Wishlist

#### GSO and ECN

- Rather, GSO or ECN
- Cannot have both

```
/* Have kernel split `packet_sz` bytes into packets... */
cmsg->cmsg level = SOL UDP;
cmsg->cmsg type = UDP SEGMENT;
cmsg->cmsg len = CMSG LEN(sizeof(uint16 t));
*((uint16_t *) CMSG_DATA(cmsg)) = packet_sz;
/* ... but how can one specify ECN for each? */
cmsg->cmsg level = IPPROTO IP:
cmsg->cmsg type = IP TOS;
cmsg->cmsg len = CMSG LEN(sizeof(tos));
memcpy(CMSG DATA(cmsg), &tos, sizeof(tos));
```

## DPLPMTUD: Suppressing EMSGSIZE

ip(7) has this to say:

It is possible to implement RFC 4821 MTU probing with SOCK\_DGRAM or

SOCK\_RAW sockets by setting a value of IP\_PMTUDISC\_PROBE (available since Linux 2.6.22). This is also particularly useful for diagnostic tools such as

tracepath(8) that wish to deliberately send probe packets larger than the observed

```
/* Even with this setting, sendmsg(2) returns -1 with EMSGSIZE
  * if datagram larger than local interface's MTU
  */
on = IP_PMTUDISC_PROBE;
s = setsockopt(fd, IPPROTO_IP, IP_MTU_DISCOVER, &on, sizeof(on));
```

• Painful: have to handle EMSGSIZE specially and retry sendmmsg(2)

Path MTU.

Bonus Section #2

HTTP/3

# HTTP/3 differences

- ALPN is required
  - This step is handled by the library
  - QUIC version I-D 29 corresponds to ALPN "h3-29" and so on
  - QUIC version 1 will correspond to "h3"
- SNI is required
  - Pass it to lsquic\_engine\_connect()
- Send headers before sending payload
  - Use lsquic\_stream\_send\_headers()
- Optional: header callbacks via ea\_hsi\_if (HSI: header set interface)
  - $\bullet$  If not specified, LSQUIC will pretend you are reading HTTP/1.x-like stream.
- Server must use ea\_lookup\_cert callback

# h3cli.c: a simple HTTP/3 client

- Specify hostname, port number (or service name), and path
- Default method is GET
- Websites to try:
  - www.litespeedtech.com
  - www.google.com
  - www.facebook.com

./h3cli www.litespeedtech.com 443 / -M HEAD

# h3cli.c: use LSENG\_HTTP flag

#### h3cli.c: connect

```
h3cli.h3cli_conn = lsquic_engine_connect(
   h3cli.h3cli_engine, N_LSQVER,
   (struct sockaddr *) &h3cli.h3cli_local_sas, &addr.sa,
   (void *) (uintptr_t) h3cli.h3cli_sock_fd, /* Peer ctx */
   NULL,
   h3cli.h3cli_hostname, /* <=== This becomes SNI in ClientHello */
   0, NULL, 0, NULL, 0);</pre>
```

### h3cli.c: send requests

```
static void h3cli client on write (struct lsquic stream *stream,
                                          lsquic stream ctx t *h) {
 struct header_buf hbuf;
 struct lsxpack header harray[5];
  struct lsquic_http_headers headers = { 5, harray, };
 h3cli_set_header(&harray[0], &hbuf, V(":method"), V("GET")):
 h3cli set header(&harray[1], &hbuf, V(":scheme"), V("https"));
 /* --- 8< --- snip --- 8< --- */
 if (0 == lsquic stream send headers(stream, &headers, 0))
      lsquic stream shutdown(stream, 1);
      lsquic_stream_wantread(stream, 1);
```

# HTTP/3: configuration options

- What's a few more tunable parameters between friends?
- Push promises
- Compression
  - QPACK choices: table size; allow blocked streams?
- Priorities (coming as an extension)
- LSQUIC aims for sane defaults
- Change to fit your needs and performance goals; YMMV