
Copyright © 1987 by Borland International, INC.

Jede Vervielfältigung dieses Handbuchs sowie des Turbo Modula-2 Softwareprogrammes wird strafrechtlich verfolgt. Die Rechte an der Dokumentation und die Rechte an dem Softwareprogramm Turbo Modula-2 liegen bei Borland International.

Der rechtmäßige Erwerb der Programmdisketten und des Handbuchs erlaubt die Nutzung der Programme analog der Benutzung eines Buches. Entsprechend der Unmöglichkeit, daß ein Buch zugleich an verschiedenen Orten von mehreren Personen gelesen wird, darf das Softwareprogramm Turbo Modula-2 nicht gleichzeitig von verschiedenen Personen an verschiedenen Orten und auf verschiedenen Geräten benutzt werden. Diskettenkopien dürfen lediglich zum Zweck der Datensicherung angefertigt werden.

Sie (der Lizenznehmer von Turbo Modula-2) dürfen eigene Programme, die Sie mit dem Turbo Modula-2 Sprachcompiler geschrieben und compiliert haben, ohne Einschränkung oder Verpflichtung zu zusätzlichen Lizenzgebühren verwenden, weitergeben oder verkaufen. Sie sind ebenfalls nicht verpflichtet anzugeben, daß Ihre Programme mit dem Turbo Modula-2 Compiler entwickelt wurden, oder daß Sie Quellcode aus Toolboxen der Borland Sprachprodukte enthalten.

Die Beispielprogramme von Turbo Modula-2, die auf den Originaldisketten mitgeliefert werden, demonstrieren die verschiedenen Eigenschaften von Turbo Modula-2. Sie dürfen diese Beispielprogramme editieren oder für den eigenen Gebrauch modifizieren. Sie dürfen diese Programme jedoch nicht als Ganzes oder in Teilen in compilierter Form oder als Quelltext weitergeben oder verkaufen. Sie dürfen die Routinen aus den Beispielprogrammen in eigenen Anwenderprogrammen verwenden, solange die resultierenden Programme nicht einem Beispielprogramm oder einem Teil eines Beispielprogrammes in Erscheinung oder Funktion gleichen.

Einschränkung der Gewährleistung

Es wird keine Garantie für die Richtigkeit des Inhaltes dieses Handbuchs übernommen. Da sich Fehler, trotz aller Bemühungen, nie vollständig vermeiden lassen, sind wir für Hinweise jederzeit dankbar.

Herausgeber:

Heimsoeth Software GmbH & Co. Produktions- und Vertriebs-KG
Fraunhoferstr. 13, Postfach 14 02 80
8000 München 5 Telefon: 089/2 60 94 67, Telex: 52 12 637 mcm d
1. Auflage August 1987

MODULA-2

TABLE OF CONTENTS

Introduction 17

Structure of This Manual 17

Typography 18

Acknowledgements 19

Chapter 1. Getting Started 21

Before Use 21

Files on Disk 21

Preparing a Work Disk 23

Running Turbo Modula-2 26

Chapter 2. A Brief Tour of Turbo Modula-2 29

The Menu System 29

Chapter 3. Language Elements 37

Character Set 37

Vocabulary 38

Numbers 38

Whole Numbers 38

Real Numbers 40

Strings and Characters 40

Strings 40

Characters 41

(Delimiters and Comments 41

Delimiters 41

Comments 42

Operators 42

Reserved Words 43

Identifiers 44

Standard Identifiers 44

User-Defined Identifiers 45

Library Identifiers 45

Chapter 4. Expressions 49*Properties of an Expression 49**Operands 50**Operators 50**Operator Precedence 51***Chapter 5. Data Structure 53***Data Objects 53**Elementary Data Types 54**Declarations 56**Constant Declarations 57**Type Declarations 58**Variable Declarations 58**User-Defined Unstructured Types 59**Enumeration Types 59**Subrange Types 60**Pointer Types 61**Structured Types 62**Array Types 62**Record Types 64**Variant Records 66**Set Types 67**Procedure Types 68***Chapter 6. Statements 71***Assignment Statements 72**WITH Statements 73**Conditional Statements 75**IF Statements 75**CASE Statements 76**Repetitive Statements 77**FOR Statements 77**WHILE Statements 79**REPEAT Statements 79**LOOP Statements 80**EXIT Statements 80**Procedural Statements 81**Procedure Calls 82**Procedure Declarations 82**Parameters 84*

Open Array Parameters 86

Function Procedures 89

RETURN Statements 90

Nested Procedures 91

Scope of Visibility 92

FORWARD Statements 93

Standard Procedures 94

Chapter 7. Modules 97

The Main Module 97

Library Modules 100

Definition and Implementation Modules 101

Compiled Modules 104

Opaque Export 105

Local Modules 106

Scope and Local Modules 107

Chapter 8. Low-Level Facilities 111

Type-Transfer Functions 113

Type Transfer and Type Conversion 114

Low-Level Types and the Pseudomodule System 115

Untyped Parameters 116

Absolute Addresses 117

Couroutines and Interrupts 118

Couroutines 118

Interrupts 119

Chapter 9. Turbo Modula-2 Extensions 121

Input and Output Extensions 121

String Extensions 122

3-dimensional Open Arrays 122

Exception-Handling Extensions 122

Syntax and Semantics of Exception Handling 123

Declaration of Exceptions 124

Raising Exceptions 124

Exception Handlers 125

Exception Propagation 127

Chapter 10. System Operations 131*File Management Utilities 132**Options 135**Avoiding the Menus 137**The Turbo Editor 138**Operating the Editor 139**Editing Commands 139**Cursor Movement Commands 141**Extended Movement Commands 143**Insert and Delete Commands 144**Block Commands 145**Find and Replace Commands 147**Miscellaneous Editing Commands 149**The Librarian 151**Searching Libraries 151**The Compiler 152**Operating the Compiler 153**The Listing 154**Error Correction 157**Running Out of Memory 159**When the Compiler Runs Out of Memory 159**When Your Program Runs Out of Memory 160**Symbol Files 162**Compiler Options and Switches 163**The Linker 167**Linking with Overlays 169**Linking the Linker 173**Version Control 174**Utilities 175**Linking Microsoft Relocatable Files 175**Profile 178***Chapter 11. The Standard Library 181***Overview of Input and Output Modules 182**Overview of Utility Modules 183**Overview of System and Low-Level Modules 183**Details of the Module Library 184**Input and Output 184**Streams 186**Logical Devices 187*

<i>The Texts Module</i>	188
<i>Standard Text Streams</i>	188
<i>Stopping the Program During Input and Output</i>	192
<i>Opening, Creating, and Closing a Text</i>	193
<i>Renaming, Deleting, and Other File Operations</i>	195
<i>Reading and Writing</i>	195
<i>READ and WRITE Statements</i>	201
<i>User-Defined I/O Drivers</i>	205
<i>The InOut Module</i>	207
<i>The Files Module</i>	210
<i>Errors During File Handling</i>	210
<i>Operations on Entire File</i>	214
<i>Code Processing</i>	216
<i>Files with Elements of Mixed Types</i>	222
<i>The Terminal Module</i>	224
<i>The ComLine Module</i>	229
<i>Redirection of Input and Output</i>	229
<i>Utility Modules</i>	233
<i>The MathLib and LongMath Modules</i>	233
<i>The Strings Modules</i>	237
<i>The Convert Module</i>	241
<i>The Doubles Module</i>	242
<i>System-Dependent Modules</i>	243
<i>The Processes Module</i>	243
<i>The Pseudomodule SYSTEM</i>	248
<i>Low-Level Access to Data</i>	248
<i>Array of Word</i>	248
<i>Coroutines and Interrupts</i>	251
<i>Z80-Specific Procedures</i>	253
<i>Interface to CP/M</i>	253
<i>Assembler Interface</i>	254
<i>Modules in Memory Management</i>	258
<i>The STORAGE Module</i>	260
<i>Standard Procedures Dependent on Storage</i>	260
<i>Heap Pointer</i>	260
<i>Dynamic Variable Errors</i>	263
<i>The Loader Module</i>	264
Chapter 12. Turbo Modula-2 Reference Directory	271
<i>ABS</i>	273
<i>ADDRESS</i>	273
<i>ADR</i>	275

ALLOCATE 276
AND 277
Append 278
Arctan 279
ArgumentError 280
ARRAY 281
available 284
Awaited 285
BADOVERLAY 286
BDOS 287
BEGIN 288
BIOS 288
BITSET 290
BOOLEAN 291
BusyRead 292
BYTE 294
Call 295
CAP 296
CAPS 296
CARD 297
CARDINAL 298
CardToStr 299
CASE 300
CHAR 303
CHR 304
clearEol 305
ClearScreen 306
ClearToEOL 307
Close 308
CloseInput 309
CloseOutput 310
CloseText 311
CODE 312
Col 313
ComLine 314
commandLine 317
ConnectDriver 318
console 319
CONST 320
Convert 321

Copy 322
Cos 323
Create 324
CreateText 325
DeadLock 326
DEALLOCATE 327
DEC 328
DEFINITION 330
Delete Files 332
Delete Strings 333
DeleteLine 334
(viceError 335
DiskFull 336
DISPOSE 337
DIV 339
Done InOut 339
Done Texts 340
DOUBLE 341
Doubles 342
DoubleToStr 343
END 343
EndError 344
Entier 345
EOF 346
EOL 347
EOLN 348
EOT 349
EXCEPTION 350
EXCL 351
EXIT 352
('tScreen 353
Lxp 353
EXPORT 354
FALSE 356
FILE 356
Files 357
FileSize 359
FILL 360
firstDrive 360
FLOAT 361

Flush 362
FOR 363
FORWARD 364
FREEMEM 365
GetName 366
GotoXY 367
HALT 368
haltOnControlC 368
HIGH 370
Highlight 371
highlightNormal 372
HRESULT 373
IF 374
IMPLEMENTATION 375
IMPORT 378
INC 379
INCL 382
Init 382
InitScreen 383
inName 384
InOut 385
IORESULT 386
INP 387
input 388
insert 389
InsertDelete 389
InsertLine 391
INT 392
INTEGER 393
IntToStr 394
IOTRANSFER 395
legal 396
Length 397
LoadError 397
LongMath 398
LONGREAL 399
Ln 401
Loader 402
LONG 403
LONGINT 404

LongMath 405
LongToStr 406
LOOP 407
MARK 408
MathLib 409
MAX 410
MIN 411
MOD 412
MODULE 413
NEW 414
NEWPROCESS 416
(extPos 417
NIL 419
Normal 419
NoTrailer 420
numCols 421
numRows 422
ODD 423
Open 423
OpenInput 425
OpenOutput 426
OpenText 427
OpSet 428
OR 429
ORD 430
OUT 431
outName 432
OUTOFMEMORY 433
output 434
OVERFLOW 435
(^INTER 436
ros 438
PROC 439
PROCEDURE declaration 440
PROCEDURE type 441
PROCESS 443
Processes 445
progName 446
PromptFor 446
QUALIFIED 448

RAISE 449
Random 450
Randomize 451
READ 452
ReadAgain (Terminal) 453
ReadAgain (Texts) 454
ReadByte 455
ReadBytes 456
ReadCard 458
ReadChar (Terminal) 459
ReadChar (Texts) 459
ReadDouble 460
ReadInt 461
ReadLine (Texts) 462
ReadLine (Terminal) 463
READLN 464
ReadLong 465
ReadReal 466
ReadRec 466
ReadString 468
ReadWord 468
REAL 470
REALOVERFLOW 472
RealToStr 473
RECORD 475
RedirectInput 479
RedirectOutput 480
RELEASE 481
Rename 482
REPEAT 484
ResetSys 484
RETURN 485
SEND 486
SET 487
SetCol 489
SetPos 490
SIGNAL 491
Sin 492
SIZE 493
SpecialOps 494

Sqrt 495
StartProcess 496
StatusError 497
STORAGE 498
String 499
StringError 500
Strings 501
StrToCard 502
Str.ToDouble 503
StrToInt 504
StrToLong 505
(ToReal 506
SYSTEM 506
termCH 508
TEXT 510
TextDriver 511
TextFile 512
TextNotOpen 513
Texts 514
TooLarge 516
TooManyTexts 517
TRANSFER 518
TRUE 519
TRUNC 520
TSIZE 521
TYPE 522
UseError 523
VAL 524
VAR 525
WAIT 526
(FILE 527
WITH 528
WORD 529
WRITE 531
WriteByte 532
WriteBytes 533
WriteCard 533
WriteChar (Terminal) 534
WriteChar (Texts) 535
WriteDouble 536

WriteHex 537
WriteInt 538
WRITELN 538
WriteLn (Terminal) 540
WriteLn (Texts) 541
WriteLn (InOut) 541
WriteLong 542
WriteOct 543
WriteReal 544
WriteRec 545
WriteString (Terminal) 545
WriteString (Texts) 546
WriteWord 547

Appendices

Appendix A. Turbo Modula-2 and Turbo Pascal 549

What's the Difference 550

Vocabulary 551

Identifier Names 551

Characters 551

Numbers 551

Strings 552

Set Constants 553

Comments 553

Declarations 554

Constant Declarations 554

Type Declarations 555

Arrays 555

Records 556

Procedure Types 557

Variable declarations 559

Procedure Declarations 559

Open Array Parameters 559

Untyped Parameters 560

Function Procedures 560

Expressions 561

Set Operators 562

Statements 564

Procedure Calls 564

Looping Statements 564
CASE Statements 565
WITH Statements 566
RETURN Statements 566
Standard Procedures in Turbo Modula-2 566

Appendix B. Installation Procedures 571

Installing M2 571
Screen Installation 572
Manual Installation 573
Using the Turbo Pascal TINST.DIA File 574
Entering Terminal Codes 574
Terminal Properties 575
Installation of Editing Commands 577
Compiler Installation (Miscellaneous) 581

Appendix C. Summary of Compiler Directives 583

Appendix D. Error Diagnosis 585

Format of a Runtime Error Message 585
Errors Detected by the Interpreter 586
Errors Detected by Support Modules 589
Exceptions Issued by Module Files 589
Exceptions Issued by Module Loader 589
The Calling Chain 590
Finding Runtime Errors 592
Compiler Error Messages 592

Appendix E. BNF Syntax for Turbo Modula-2 597

Introduction

This book is a reference manual for the Turbo Modula-2 system, implemented for the CP/M-80 operating system running on Z80 computers. Although there are many examples throughout the book, this is not a tutorial for Modula-2 programming; a basic knowledge of Modula-2 or Pascal is assumed.

The Modula-2 language was designed in 1980 by Niklaus Wirth, who also created Pascal. While Modula-2 provides the powerful data and statement structures of Pascal, it also incorporates a modular structure as well as basic facilities for multiprogramming applications. Turbo Modula-2 is Borland's practical implementation of Modula-2. It follows closely the definition of standard Modula-2 as defined by N. Wirth in his book, *Programming in Modula-2* (3rd Ed. New York: Springer Verlag, 1984). However, it differs from Modula-2 in two main areas: It has an easy-to-use I/O library and optional extensions.

Because of the wide range of possible external devices, and in order to make the language truly machine-independent, the strict definition of Modula-2 does not include any input/output or low-level facilities. Turbo Modula-2, however, provides an extensive and flexible library of these facilities, which for the most part follow Wirth's suggestions.

Turbo Modula-2 also includes several constructs not mentioned in the original definition; for example, general-purpose *READ* and *WRITE* statements, string comparison and assignment, multidimensional open arrays and exception handling.

Structure of This Manual

This manual is divided into four main areas: an introduction, the module library, a look-up section, and appendices.

Chapter 1 will get you started using the Turbo Modula-2 system, walking you through such operations as copying your distribution disk and running your first program.

Chapters 2 through 7 provide a thorough description of Turbo Modula-2 programming, with a general discussion of the Modula-2 language elements and a more detailed explanation of data, statement, and program structures, as well as local and library modules. Chapters 5 and 6 serve as your guide to writing simple Modula-2 programs.

Chapters 8 and 9 cover the more advanced topics in Turbo Modula-2, including system-specific functions and extensions and low-level facilities.

Chapter 10 covers system operations, including discussion of the editor, the compiler, the linker, the librarian, and other operational features.

Chapter 11 contains the library modules, which consist of 14 modules containing over 100 procedures. You may want to read the short description of each library module now, and go back to study the specifics when you need to use a particular module.

Chapter 12 comprises Turbo Modula-2's extensive alphabetical look-up section, which contains entries of reserved words standard identifiers, and library identifiers.

The five appendices provide information on the operation of Modula-2, while also presenting some comparisons between Turbo Pascal and Turbo Modula-2. Appendix A provides the language comparisons. Appendix B details the instructions for installing Turbo Modula-2 on your system. Appendix C summarizes the compiler directives discussed in Chapter 9. Appendix D lists error messages and their definitions/diagnoses. Appendix E provides the BNF syntax diagrams of Turbo Modula-2.

Typography

The body of this manual is printed in a normal typeface. Special typefaces are used for the following purposes:

Alternate Alternate characters are used to illustrate program examples and screen displays.

Italics *Italics* are used to emphasize certain concepts and first-mentioned terms.

Boldface Boldface is used to mark reserved words in text as well as in programming examples.

Acknowledgements

Several programs, languages, and operating systems are referenced in this manual; the following lists them and their respective companies.

- Turbo Pascal is a registered trademark of Borland International.
- WordStar is a registered trademark of MicroPro International.
- CP/M is a registered trademark of Digital Research Inc.
- Microsoft is a registered trademark of Microsoft Corp.

Chapter 1

Getting Started

Before Use

When you receive your Turbo Modula-2 disk, complete and mail in the license agreement at the front of this manual. This agreement allows you to make as many copies as you need for your personal use and backup purposes only.

For your own protection, make a backup copy of the distribution disk with your file-copy or disk-copy program before you start using Turbo Modula-2. Make certain all files have transferred successfully, then store the original disk in a safe place. If anything happens to the backup copy, you can make a new backup copy from the original.

Files on Disk

The files you have just copied from the distribution disk to your backup disk are described here.

System Files

M2.COM	The Turbo Modula-2 system file contains the M-code interpreter, the Overlay Manager, and the runtime system. Entering the command M2 on your terminal will load this file and get Turbo Modula-2 up and running.
M2.OVR	Overlay file for M2.COM.
SHELL.MCD	The Turbo Modula-2 menu shell.
COMPILE.MCD	The Turbo Modula-2 M-code compiler.
GENZ80.MCD	The Turbo Modula-2 compiler's optional second-pass used for generating native Z80 code.

ERRMSG.S.OVR	Text file containing error messages.
LIBRARY.MCD	A library manager that combines several compiled modules into a single .LIB file.
SYSLIB.LIB	The standard modules supplied with the Turbo Modula-2 system.
LINK.MCD	The Turbo Modula-2 static overlay linker that is used to produce stand-alone (.COM) files.

Installation Files

INSTM2.COM	This installation program allows you to install the Turbo Modula-2 compiler (M2.COM), as well as other .COM programs produced by the linker.
INSTM2.OVR	Installation program overlay file.
INSTM2.DTA	Installation data file compatible with Turbo Pascal's TINST.DTA file (see Appendix B, »Installation Procedures«, for more details).

Example Files

*.MOD	Sample Modula-2 programs.
-------	---------------------------

Other Files

READ.ME	If present, this file contains the latest corrections or suggestions on the use of the system.
---------	--

Utility Files

REL.MCD	A Turbo Modula-2 utility that converts Microsoft REL (relocatable) object files into the .MCD format used by Turbo Modula-2.
---------	--

PROFILE.MCD A Turbo Modula-2 utility that counts instructions in the various procedures of an M-code program. This utility can help you improve the efficiency of your Turbo Modula-2 programs.

Preparing a Work Disk

The most effective way to use Turbo Modula-2 is to have the system files on your boot disk, allowing you to use other drives for programs and data. To make your working disk, follow these steps:

1. Prepare a system boot disk, which will become your work disk; also include a file-copy program such as PIP. For more information on this step, consult your operating system's user manual.
2. Now boot your system using the work disk, and place the Turbo Modula-2 backup disk (the one you made after reading the first paragraph of this chapter) into a free disk drive.
3. Copy the following files from your **Turbo Modula-2 backup disk to the work disk**:

M2.OVR
M2.COM
SHELL.MCD
INSTM2.COM
INSTM2.MSG
INSTM2.DTA

Now run the installation program by typing INSTM2 and pressing [RET] at the system prompt. This message will appear:

Install program ([RETURN] for M2.COM):

Press [RET] again and the following screen will appear:

Modula-M2 system installation menu.
Choose installation item from the following:

[S]creen installation		[C]ommand installation
[M]iscellaneous		[Q]uit

Enter S, C, M, or Q:

5. Press **s** for Screen installation and pick the appropriate terminal type. For more details on installing your system, see Appendix B.

Note: For Turbo Pascal owners with custom terminals, you may substitute your Turbo Pascal terminal data file for INSTM2.DTA by copying the Pascal data file to your work disk and renaming it from TINST.DTA to INSTM2.DTA.

6. Now press **M** for Miscellaneous. Tell the system which drives to search in for library and work files (generally, all your drives will be included).

Note: Do not try to install the keyboard commands until you become more familiar with the editor.

7. Now select **Q** for Quit and delete the installation files from your work disk (INSTM2.*). Then delete the file-copy program we had you install in Step 1. Your bootable work disk should contain an installed copy of M2.COM, - M2.OVR, and SHELL.MCD.
8. With your work disk in drive A and your backup disk in drive B, type M2 at the system prompt. The Turbo Modula-2 main menu will appear.

Selected drive: A

Work file:

Edit	Compile	Run	eXecute
Link	Options	Quit	liBrarian
Dir	Filecopy	Kill	reName Type

>

- Note: Do not attempt to execute any menu items (except Filecopy) until all necessary files have been copied to the work disk.
9. To finish preparing your work disk, you must copy some files using Turbo Moudula-2. To begin, press **F** for Filecopy at the main menu. Then at the "Copy from:" prompt, type in the name of the file you want to copy; for example:

Copy from: B:*.MCD [RET]

Press **RET** to get the "Copy to:" prompt:

Copy to: A: [RET]

Or do it in one step, like so:

Copy from: B:ERRMSG.S.OVR A: [RET]

or

Copy from: B:SYSLIB.LIB A: [RET]

Copy the following files using any of the preceding procedures:

COMPILE.MCD
SYSLIB.LIB
ERRMSG.S.OVR

And optionally, if there is room left on your disk, copy these files as well:

GENZ80.MCD
LIBRARY.MCD
LINK.MCD
REL.MCD
PROFILE.MCD

(The files M2.COM, M2.OVR, and SHELL.MCD should already be on your work disk.)

At this point, it would be wise to make a backup of your work disk to avoid having to reinstall the system. Now you're ready to try out the Turbo Modula-2 compiler.

Running Turbo Modula-2

Now that you've prepared a work disk and logged onto the drive containing it, you are ready to load Turbo Modula-2 into memory. After the system prompt, type M2, press **RET**, and the following message will appear (but your terminal will be listed):

Turbo Modula-2 System

Version 1.00
CP/M-80, Z80

Copyright (C) 1984, 1985, 1986

Borland International

Terminal: No Terminal Selected

This will be followed by the main menu:

Selected drive: A

Work file:

Edit Compile Run execute

Link Options Out Library

Dir	Filecopy	Kill	reName	Type
-----	----------	------	--------	------

2

To run a program, place the backup disk containing the sample programs (*.MOD) in the B drive. At the main menu, press **C**. You will see this message:

Workfile name:

Type Hello and press **[RET]**. Turbo Modula-2 will now compile the sample program. The compilation terminates with the following message:

Compiled bytes: 32
M-code file A00:HELLO.MCD produced.

Now press **[R]** and then **[RET]** to run the program. The screen will look like this:

() Run MCD-file: A00:HELLO

Hello World!

>

After the prompt (>) appears in the main menu, quit Turbo Modula-2 by pressing **[Q]**; this will return you to the operating system.

Chapter 2

A Brief Tour of Turbo Modula-2

For those of you who want to learn Turbo Modula-2 quickly, or are already acquainted with some aspects of it, this chapter gets you into the thick of things in a hurry. Only a minimum of explanation is given here; to use Turbo Modula-2 to its full potential, you should refer to Chapter 10. This chapter will take you through the menu system and briefly explain how to use each feature.

The Menu System

To start, place your work disk in the logged drive (to make a work disk, refer to Chapter 1), then type M2 and press [RET] to bring up the following screen:

Turbo Modula-2 System	Version 1.00
	CP/M-80, Z80
Copyright(C) 1984,1985, 1986	Borland International

Terminal: No terminal selected

This screen will be quickly followed by the main menu:

Selected drive: A

(...ork file:

Edit	Compile	Run	eXecute
Link	Options	Quit	liBrarian
Dir	Filecopy	Kill	reName
Type			
>			

The main menu contains all of the major functions performed by Turbo Modula-2. The following section describes each element of the main menu and displays a sample screen where appropriate. (Note: The italicized items in the example indicate that you must type in your own data.) To select each item, you press the letter highlighted on your screen.

Selected drive. Use this to change the default drive; press [**s**] and type the letter of the drive that is to become the new default drive. You can also use this function to reset the system or to log the disk; this is handy when you need to insert a new disk to copy files to.

> S

New drive: a:

Work file. Sets a default file name for other menu commands to use, such as Edit, Compile, Link, and so on. You can override this default by pressing the space bar before selecting the desired menu item (see the section «Avoiding the Menus» in Chapter 10).

> W

Workfile name: Myfile

Edit. Invokes the Turbo Editor, a WordStar-like editor. This editor is similar to the Turbo Pascal editor; however, it is a «one-pass» virtual editor that limits file sizes to the disk space available rather than internal memory.

> E

Edit file: COO:MYFILE.MOD

Compile. Turbo Modula-2 is an incremental compiler, which means it saves its state when a compile-time error is found. You can then enter the editor, correct the error and exit, at which point compilation continues at the closest block to where the error was corrected.

> C

Compile file: COO:MYFILE.MOD

Run. Runs compiled code without going through the linking step. The Run command does dynamic linking while it is reading in the support modules to run a program.

> R

Run MCD-file: C00:MYFILE

eXecute. Provides a way to run most external programs, such as STAT or even the Turbo Modula-2 installation program.

() eXecute COM-file: INSTM2.COM

Link. This command serves two purposes: It links separately compiled modules into a stand-alone .COM file and links specified modules so that they will load quickly when linked dynamically during the Run command. In addition, the linker provides the facility to generate overlays.

> L

Link main module: C00:MYFILE

Options. This command invokes the following Options menu:

compiler options:

List	(OFF)	Native	(OFF)	eXtensions	(ON)
Test	(OFF)	Overflow	(OFF)	Upper=lower	(OFF)

Path to search: SYSLIB

{ and run-time error

Save current selection Quit

>

The compiler options are global; that is, they influence the entire compilation unless overridden with internal switches. You can toggle the compiler options by pressing the key for the capital (highlighted) letter in each option (for example, L for List). The following describes the function of each compiler option:

- List determines if source output is displayed on the monitor during compilation.
- Native determines which type of code is produced, M-code or Native.
- eXtensions tells the compiler to issue warning messages if a program is using any of the Turbo Modula-2 extensions.
- Test determines if the compiler generates test code for array bounds and subrange checking.
- Overflow determines if the compiler generates code to check for integer overflow.
- Upper=lower tells the compiler whether or not to be case-sensitive.
- Path to search tells the system which library files to search when looking for external modules.

>P

New search path: SYSLIB MYLIB

- Find run-time error helps you find the runtime error position.

>F

Module name: C00:MYFILE.MOD

Enter PC: 23

- Save current selection allows you to save the options and search path you find most comfortable..
- Quit returns you to the main menu.

Now we'll get back to describing the main menu options.

Quit. Returns you to the operating system.

librarian. This command prompts you for a library **name** and then displays the librarian menu, as shown in the following:

>B

Select library: SYSLIB

Selected library: A00:SYSLIB.LIB

Dir Include Copy

(1 Compress Quit

The library keeps many precompiled library modules in one file. The compiled versions of the definition module and the implementation module are contained in the library as .SYM files and .MCD files, respectively.

Notice that the librarian menu has a different prompt, distinguishing the following library management utilities from the main menu's file-management routines. Again, initiate each option by pressing the capital (or highlighted) letter of its name.

- Selected library* prompts for a name prior to displaying the librarian menu. Press **s** to change the selected library without returning to the main menu.

*S

Select library:

- Dir* lists the .SYM and .MCD files, that are stored in the selected library. In addition, it lists the size of each file and the cumulative size of all the files.
*D

Directory of library: COO.SYSLIB

1: COMLINE	.MCD	0.5K	.SYM	0.5K
2: CONVERT	.MCD	2.0K	.SYM	1.0K
3: DOUBLES	.MCD	3.5K	.SYM	1.0K
4: FILES	.MCD	3.0K	.SYM	1.5K
5: INOUT	.MCD	1.0K	.SYM	1.0K
6: LOADER	.MCD	2.5K	.SYM	1.0K
7: LONGMATH	.MCD	2.0K	.SYM	0.5K
8: MATHLIB	.MCD	1.5K	.SYM	0.5K
9: PROCESSE	.MCD	0.5K	.SYM	0.5K
10: STRINGS	.MCD	0.5K	.SYM	1.0K
11: TERMINAL	.MCD	1.0K	.SYM	1.0K
12: TEXTS	.MCD	2.0K	.SYM	2.0K
Total size:		39.5K	Unused:	8.0K

- Include* takes external files (either .SYM files or .MCD) and places them in the library.

*I

Include file:

- Copy* copies files from the selected library to either another library or to a stand-alone .SYM or .MCD file.

*C

Copy module:

- Kill* erases a specified library module.

*K

Kill module:

- cOmpress* eliminates unnecessary space in the library files. Since this command can take some time, we recommend using it only on stable (debugged) library files.

*0

Compressing library, please wait.

*

- Quit** returns you to the main menu.

The remainder of the main menu options described here are file-management commands.

Dir. The directory command allows you to display a full or partial directory of any disk. It accepts drive name, user areas, and wildcard file names (such as *.MOD or **) to build the directory listing. When a directory listing is displayed, it is shown with a number before each file name. These numbers may be used in subsequent file-management commands, such as **Filecopy** and **Kill**.

>D

Directory mask:

1: COMPILE .MCD	4: LIBRARY .MCD	7: M2 .OVR	10: SHELL .MCD
2: ERRMSG.S .OVR	5: LINK .MCD	8: MYFILE .BAK	11: SYSLIB .LIB
3: GENZ80 .MCD	6: M2 .COM	9: MYFILE .MOD	

Bytes Remaining on A: 96K

- Filecopy.** This command will accept a drive name, a user number, and a wildcard file name as the source file. In addition, the source file may be a list or range of numbers that reference the last directory command.

>F

Copy from : 1 3-5
Copy to : b:

- Kill.** This command allows you to delete disk files without leaving the shell. It also accepts drive names, user numbers, and wildcard file names, as well as a list or range of numbers referencing the last directory command.

```
>K  
Kill file: *.bak  
Deleting D00:MYFILE.BAK
```

- reName.** This command accepts one new file name (with optional drive and user area) for the old name and one file specification for the new name.

```
>N
```

```
Rename from : myfile.mod  
Rename to   : myfile.def
```

- Type** This command displays the specified file on the screen. You can pause output by pressing **[Ctrl][S]** and terminate output by pressing **[Ctrl][C]**.

```
>T
```

```
Type file: myfile.def
```

```
DEFINITIONMODULE MyFile;  
  PROCEDURE MyProc;  
END MyFile.
```

Chapter 3

Language Elements

Language elements are the building blocks that form a program and are the fundamental units recognized by the compiler. Each element is used to build a different level of abstraction. For example, a digit is an element that can be combined with other digits to form a number that is an element at a different level of abstraction. Thus, we can continue abstracting until a level is reached where an entire program is represented by one word, such as *MyProgram*.

This chapter describes the elements used to assemble meaningful program statements in Turbo Modula-2. In the following section, we will look at the first level of abstraction used by Turbo Modula-2; namely, the characters that comprise symbols.

Character Set

The Turbo Modula-2 character set includes all characters that are legal in Modula-2 declarations, expressions, and statements. This set is made up of alphabetic, numeric, and special characters.

Alphabetic (uppercase and lowercase)

A to *Z* and *a* to *z*

Numeric

0 1 2 3 4 5 6 7 8 9

Special Characters

~ + - * / = ^ < > () [] { } . , : ; & # ' " !

Not all characters are available on all terminals; thus Turbo Modula-2 recognizes the following synonyms:

(.	for	[Left parentheses and period for left index bracket
.)	for]	Period and right parentheses for right index bracket
(:	for	{	Left parentheses and colon for left brace
):	for	}	Colon and right parentheses for right brace
!	for		Exclamation mark for vertical bar

Using these available symbols, it is possible to create a great number of Modula-2 programs.

Vocabulary

The Modula-2 compiler recognizes certain groups of characters as symbols themselves, including numbers, characters, strings, delimiters, operators, and reserved words and identifiers. These symbols (or characters) comprise the second level of abstraction, providing you with the means to form sentences and thus programs in Modula-2.

The Turbo Modula-2 vocabulary can be divided into five classes: numbers, strings and characters, delimiters and comments, operators, and reserved words and identifiers (user-defined, standard, and library).

Numbers

There are two types of numbers defined in Modula-2: whole numbers and real numbers. Turbo Modula-2 recognizes three subtypes of whole numbers and two subtypes of real numbers.

Whole Numbers

Turbo's three types of whole numbers may be further defined as two signed numbers and one unsigned number. The two signed whole numbers are single precision and double precision, which are INTEGER and LONGINT, respectively. The unsigned whole number is single precision and is called a CARDINAL.

Following are the ranges for whole numbers:

Type	Range
INTEGER	(single-precision, signed whole number)
LONGINT	(double-precision, signed whole numbers)
CARDINAL	(single-precision, unsigned whole numbers)

Single-precision integers can be formed in any of these three bases: decimal, octal, or hexadecimal (hex). Decimal numbers consist of the digits 0 to 9. Octal numbers comprise the digits 0 to 7, followed by the letter *B*. Hex numbers consist of the digits 0 to 9 and the letters *A* to *F*. The first character in a hexadecimal number must be a digit and the hexadecimal number must end with the letter *H*. The following are examples of legal and illegal single-precision whole numbers:

Legal (base)

Single	Precision	Illegal	Single Precision
1986	(Decimal)	12.34	Decimal point illegal.
-10	(Decimal)	FH	First character must be a digit.
10B	(Octal)	08B	8 is not a legal octal digit.
0FFFFFH	(Hex)	0F	Requires <i>H</i> for hexadecimal.
1AH	(Hex)	0GH	G is not a legal hexadecimal digit.
62345	(Decimal)	-60000	Too small for single-precision integer.

Double-precision integers, or long integers, can only be formed in decimal notation, followed by the letter *L*. Following are legal and illegal long integers:

Legal Double Precision	Illegal Double Precision
123L	123
456L	45.6
120392237L	0FFFFFFFH

Real Numbers

Single-precision and double-precision numbers are the two types of real numbers recognized by Turbo Modula-2. They are called REAL and LONGREAL, respectively.

Single-precision real numbers can be accurate up to 6 decimal points, while double-precision reals can be accurate up to 14 decimal points. The ranges of REAL and LONGREAL are as follows:

Type	Range		
REAL	-6.80565E+38	to	6.80565E+38
LONGREAL	-3.5953862697246D+308	to	3.5953862697246D+308

As shown in the previous ranges, a real number contains a sequence of numeric characters, containing a decimal point with an optional *scale factor*. The scale factor is specified by the letter *E* or *D* (*E* for single precision and *D* for double precision) and an integer, beginning with an optional plus (+) or minus (-) sign. The following are samples of legal and illegal real numbers:

Legal Real Numbers		Illegal Real Numbers	
12.34	Single precision	1 453	Blank space illegal
0.1E3	Equals 100.0	12	Needs decimal point
3.6E-5	Equals 0.000036	3.6E-99	Exponent too large
12.34D0	Double precision	1,423.0	Comma is illegal
0.1D3	Equals 100.0	12d-10	<i>D</i> and <i>E</i> must be uppercase
3.6D-5	Equals 0.000036	12.0D-999	Exponent too small

Strings and Characters

Strings

A *string* is a portion of text that can be handled as a single unit, such as a message to be written on the screen. Strings provide a means of manipulating text within a program. Constant strings are formed by enclosing a sequence of printable characters in either single or double quotation marks. The opening and closing quotation marks must be of the same kind, and that kind cannot occur within the string. Here are some examples of legal and illegal constant strings:

Legal Constant Strings

'My name is:'
 'He said, "What is the time?"'
 "12 34"

Illegal Constant Strings

"Hello'
 "She said, "The
 time is 4 p.m."

Quotes must match
 Opening quote type
 cannot occur within
 the string

Characters

A *character* is a single, printable ASCII symbol enclosed in matching quotes (single or double), or an octal integer (up to 377) followed by the letter C. In theory, characters occupy 1 byte of storage, and thus may take on any value between 0 and 255 decimal (or 377 octal). In practice, characters may occupy an entire word (2 bytes) in memory. The only time a character actually takes 1 byte is when it is declared as part of an array (arrays pack characters together). Any other declaration involving a character always results in the character occupying 2 bytes. Following are legal and illegal characters:

Legal Characters

"A"
 "*"
 15C
 377C

Illegal Characters

The letter A
 Asterisk
 Carriage return
 Largest character

"AB"
 "Q"
 8C
 777C

Only one character allowed
 Must have matching quotes
 8 is not an octal digit
 Octal number too large for
 character

When a string contains only one character enclosed in quotes, the string is considered a special case and is compatible with a character. This only applies to constant strings like "A" or "*", not character variables.

Delimiters and Comments**Delimiters**

A *delimiter* is one or more characters that separate other syntactic entities. For example:

Blanks separate identifiers and reserved words.

The vertical bar separates CASE, exception, and variant record statements.

,	The comma separates items in parameter lists and declarations.
;	The semicolon separates program statements and declarations.
* , + , - , /	The mathematical symbols (or operators) also serve as delimiters between operands.
Comments	A comment may appear anywhere one of the preceding delimiters is allowed. In addition, a comment may serve as a substitute for a blank space between symbols.

Comments

Comments contain descriptive or explanatory text about a program. Since, for the most part, they are ignored by the compiler (except for compiler switches, see Chapter 10), they can contain any characters enclosed by a set of parentheses and asterisks.

Unlike Pascal, which allows either the symbol sets (* and *) or { and } to delimit comments, Modula-2 uses only (* and *) to delimit comments and uses { and } exclusively to delimit sets. Comments may be nested to any depth (limited only by memory space). They may occur anywhere a delimiter is allowed, and are treated as blanks. The following is a sample of legal and illegal comments:

Legal Comments	Illegal Comments	
(* This is a comment *)	{This isn't a comment}	Braces not allowed
(*Comments can be nested (* like this *) *)	(* Hi (* there *)	Missing second closing comment

Operators

Operators in Modula-2 are made up of either special symbols or reserved words. In the case of reserved words, they are always printed in uppercase letters (see "Compiler Switches," in Chapter 10) and cannot be used as identifiers. The following lists the special symbol and function of each operator:

+	Addition and set union
-	Subtraction and set difference
*	Multiplication and set intersection
/	Division and symmetric set difference
&	Logical AND
~	Logical NOT
^	Dereferencing
:=	Assignment
=	Equality
# <>	Inequality
<	Less than
>	Greater than
<=	Less than or equal
>=	Greater than or equal
()	Parentheses
[]	Index brackets
{}	Set braces
, ; .. : “ ”	Punctuation

In addition to these special symbols, the following reserved words are also operators:

AND	Logical AND
OR	Logical OR
NOT	Logical NOT
IN	Set membership
DIV	Integer division
MOD	Integer modulus

Operators are closely related to the operands they work on. In most cases, an operator will only work with certain kinds of operands; for example, to divide real numbers you must use the real division operator: 2.8/2.9. In addition, the result of an operation could be of a different type from the original operands; for example, comparing two integers results in a Boolean value: 1 = 2. (Chapter 4 discusses operators and operands in more depth.)

Reserved Words

Reserved words in Modula-2 are symbols with a predefined and static meaning;

they cannot be used for any other purpose. They are always written in uppercase letters.

AND	ELSIF	LOOP	REPEAT
ARRAY	END	MOD	RETURN
BEGIN	EXIT	MODULE	SET
BY	EXPORT	NOT	THEN
CASE	FOR	OF	TO
CONST	FROM	OR	TYPE
DEFINITION	IF	POINTER	UNTIL
DIV	IMPLEMENTATION	PROCEDURE	VAR
DO	IMPORT	QUALIFIED	WHILE
ELSE	IN	RECORD	WITH

The following Turbo Modula-2 reserved words provide useful extensions to the Modula-2 standard:

EXCEPTION FORWARD RAISE

In this manual (and other Borland manuals), all reserved words are in boldface type; however, they will not appear on your screen this way.

The preceding symbols and reserved words will be discussed in the next few chapters. For an immediate discussion of reserved words, refer to Chapter 12, "Turbo Modula-2 Reference Directory."

Identifiers

Identifiers are unique names given to constants, types, variables, procedures, and modules. Identifiers are sequences of alphabetic and numeric characters; the first character must always be a letter. Note that Turbo Modula-2 distinguishes between uppercase and lowercase letters; thus, unlike in Pascal, in Turbo Modula-2 *Var1* and *VARI* are considered two unique identifiers.

Standard Identifiers

Turbo Modula-2 has a number of predefined identifiers for special purposes. Standard identifiers are "visible" and available in all modules without explicitly importing them. The standard identifiers can be redefined, but doing so will cost you the function offered by the predefined identifier. Also, a redefined standard

identifier only affects the module where it is redefined, not other imported modules. (Again, note that the use of uppercase and lowercase is significant.)

ABS	DISPOSE	INT	ODD
BITSET	DOUBLE	INTEGER	ORD
BOOLEAN	EXCL	LONG	PROC
CAP	FALSE	LONGINT	REAL
CARD	FLOAT	LONGREAL	SIZE
CARDINAL	HALT	MAX	TRUE
CHAR	HIGH	MIN	TRUNC
CHR	INC	NEW	VAL
DEC	INCL	NIL	

The following standard procedures provide extensions to the Modula-2 standard:

READ READLN WRITE WRITELN

User-Defined Identifiers

Users may declare their own identifiers; however, there are two rules to follow: Identifiers must begin with an alphabetic character and may consist only of alphabetic and numeric characters; spaces, underscores, and other special characters are not allowed. The following is a sampling of legal and illegal user-defined identifiers:

Legal Identifiers	Illegal Identifiers
MyName	My Name
Tr612	Public-Transit
hello	lst
Modula2	LAST__PRIME
	Blank space illegal
	Hyphen illegal
	First character must be a letter
	Underscore illegal

Library Identifiers

Turbo Modula-2 contains an extensive library (in SYSLIB.LIB) with a number of predefined modules. The modules and their identifiers can be redefined, but doing

so will cause you to lose the function offered by them, or at least make using them more awkward.

In one sense, library identifiers are no different than user-defined identifiers since you can create libraries with new identifiers, thus adding to the list of available library identifiers. However, there are two differences.

One is that Turbo Modula-2 provides a set of standard identifiers that are necessary for programming in Modula-2. We presume that some of these standard identifiers are in every implementation of Modula-2, while we have added others for machine-specific reasons. The second difference is that Turbo Modula-2 "comprehends" the functions of system-dependent library modules, such as *SYSTEM* and *STORAGE*. Thus, when an item is used from the machine-dependent *SYSTEM* module, the compiler already knows about it and has no need to look at additional symbol files to determine usage. Table 3-1 is a list of library identifiers.

Table 3-1

ADDRESS	EndError	Open	SpecialOps
ADR	Entier	OpenInput	Sqrt
ALLOCATE	EOF	OpenOutput	StartProcess
Append	EOL	OpenText	StatusError
Arctan	EOLN	OpSet	STORAGE
Argument Error	EOT	OUT	String
available	ExitScreen	outName	StringError
Awaited	Exp	OUTOFMEMORY	Strings
BDOS	FILE	output	StrToCard
BIOS	Files	OVERFLOW	Str.ToDouble
BusyRead	FileSize	Pos	StrToInt
BYTE	FILL	PROCESS	StrToLong
Call	firstDrive	Processes	StrToReal
CAPS	Flush	progName	SYSTEM
CardToStr	FREEMEM	PromptFor	termCH
clearEol	GetName	Random	Terminal
ClearScreen	GotoXY	Randomize	TEXT
ClearToEOL	haltOnControlC	READ	TextDriver
Close	Highlight	ReadAgain	TextFile
CloseInput	highlightNormal	ReadByte	TextNotOpen
CloseOutput	HI_RESULT	ReadBytes	Texts
CloseText	Init	ReadCard	TooLarge
CODE	InitScreen	ReadChar	TooManyTexts

Col	inName	ReadDouble	TRANSFER
ComLine	InOut	ReadInt	TSIZE
commandLine	INP	ReadLine	UseError
ConnectDriver	Insert	ReadLn	WAIT
console	insertDelete	READLN	WORD
Convert	InsertLine	ReadLong	Write
Copy	IntToStr	REALOVERFLOW	WRITEBYTE
Cos	ICRESULT	ReadReal	WRITEBYTES
Create	IOTRANSFER	ReadRec	WriteCard
CreateText	legal	ReadString	WriteChar
DeadLock	Length	ReadWord	WriteDouble
DEALLOCATE	Ln	RealToStr	WriteHex
Delete	Loader	RedirectInput	WriteInt
DeleteLine	LoadError	RedirectOutput	WriteLn
DeviceError	LongMath	RELEASE	WRITELN
DiskFull	LongToStr	Rename	WriteLong
DISPOSE	MARK	ResetSys	WriteOct
Done	MathLib	SEND	WriteReal
Doubles	NEW	SetCol	WriteRec
DoubleToStr	NEWPROCESS	SetPos	WriteString
	NextPos	SIGNAL	WriteWord
	Normal	Sin	
	NoTrailer	SIZE	
	numCols		
	numRows		

For more detailed information regarding library modules, refer to Chapter 11, "The Standard Library." For information regarding the preceding library identifiers, refer to Chapter 12, "Turbo Modula-2 Reference Directory."

Chapter 4

Expressions

An *expression* is a sequence of language elements that combine to form a temporary data object with a possibly different value and type than either original element. An expression consists of operators and operands that perform certain operations when an expression is "evaluated."

In the previous chapter, we looked at the character set and vocabulary used to build Modula-2 programs. Here we'll take the next step by showing you how to form expressions with language elements, while also providing the foundation for our discussion of abstract data types in the next chapter.

Properties of an Expression

To explain expressions, let's look at a simple example from mathematics:

$$1 + 2$$

1. The expression consists of **operands and an operator**: The two operands are 1 and 2, and the operator is +.
2. Expressions have a result value that can be assigned or used in further expression evaluation. The preceding example has a result value of 3.
3. addition, the expression is considered to have a type that is determined by the result. In this example, the result type is CARDINAL.

Expressions are evaluated by applying each operator to its operands. In general, expressions are evaluated from left to right. Operators that take precedence over others are executed first, despite the position of the operator within the expression (see the later section, "Operator Precedence").

Operands

In Turbo Modula-2, the operands of an expression can be any *elementary type* (such as INTEGER or REAL), including function procedures that return an elementary type (such as standard function procedures ABS, CHR, or SIZE). However, operands within a particular expression must be of the same type; for instance, in the previous example both operands are integers. The following shows an illegal expression, one that mixes integers and reals:

3.4 + 30

The purpose of this type-checking is to make explicit the types of operands involved in the expression. Unlike Pascal and some other languages, Modula-2 in general has no implicit type conversions; however, there are ways to override this strict type-checking. The explicit conversion in the following example makes the preceding example legal:

3.4 + FLOAT(30)

The result of this expression is 33.4 and is of type REAL. Note that it is possible for the result type to be different from its operands. Consider the following:

1 = 2

Here is an expression containing integer operands and one operator. When the operator is applied to the operands, the resulting value is FALSE and the result's type is BOOLEAN.

Operators

The most distinguishing feature of an operator is the type (or types) it works on. There are four classes of operator:

Arithmetic	Performs normal mathematical operations. Operations are performed on numbers, such as INTEGER and REAL.
Relational	Does a comparison of like items to obtain a BOOLEAN result of TRUE or FALSE. The items compared can be of any elementary type; for example, INTEGER, POINTER, string, and so on.

Logical	Performs the combination and the negation of BOOLEAN expressions. For example, the expression, NOT Raining AND (Today=Friday) combines the negation of <i>Raining</i> and the truth of <i>(Today=Friday)</i> .
Set	Performs bit operations on operands of type BITSET or logical set operations on user-defined SET types; for example, you may wish to mask the high bit of bytes in a WordStar document file, or you may want to include the enumerated value <i>Red</i> in a user-defined set of colors.

Since the selection of operators is dependent on the type they're operating on, we'll defer further description of each operator until its type is introduced in Chapter 5.

Operator Precedence

When more than one operator appears in an expression, its meaning or result is dependent on the order of evaluation. Operations of highest precedence are resolved first; operations of equal precedence proceed from left to right. Operations within parentheses are of the highest priority, and thus are evaluated first. *Operator precedence* is defined as follows:

1st (Highest)	priority:	NOT, ~
2nd	priority:	* , /, DIV, MOD, AND, &
3rd	priority:	+, -, OR
4th (Lowest)	priority:	=, #, <, >, >=, <>, IN

As an example, let's look at a simple mathematical expression:

(

1.0 + 2.0 * 3.0

Evaluation begins on the left with $1.0 + 2.0$; however, this subexpression is not fully evaluated until the priority level of the next operator is determined. The previous table, indicates that multiplication has a higher priority than addition; thus the subexpression $2.0 * 3.0$ is evaluated first, with its product added to 1.0 for a result of 7.0 (type REAL).

Parentheses can alter the interpretation of operator precedence. Whenever parentheses are found in an expression, the contents within the parentheses are evaluated before being combined with any other item in the expression.

If the prior example is changed to include parentheses, as shown in the following:

$$(1.0 + 2.0) * 3.0$$

then evaluation will start with the contents of the parentheses, and the sum will be multiplied by 3.0. This time the result will be 9.0 (type REAL).

Boolean expressions are evaluated almost the same as other expressions. However, in Boolean expressions, it is not always necessary to evaluate the entire expression in order to determine that expression's result. This is called *short-circuit evaluation*. For example, consider the following Boolean expression:

$$(2+4=4) \text{ AND } (3+4=7)$$

To begin evaluation, we would first resolve the contents of the first parenthesis ($2+4=4$); its result is FALSE, which is saved. The AND operator is next, because both operands to an AND must be TRUE for the entire expression to be TRUE, we need only evaluate the contents of the second set of parentheses if the first operand is TRUE. In this case, the first operand of the AND is FALSE so the second operand to AND is not evaluated (the expression is short-circuited).

This order of evaluation is useful for certain algorithms that need to access an element of an array only if the index is not a number that would cause an array bounds error. Thus, assuming you understand arrays, the following expression would *not* cause a bounds error if I is 0 and the array a is undefined when equal to 0:

$$(I \neq 0) \text{ AND } (a[I] = 100)$$

Other Boolean operators, such as OR, are evaluated in a similar manner. That is, if the first operand of an OR expression is TRUE, the second operand is not evaluated.

Chapter 5

Data Structure

Modula-2 is often described as a language that encourages structured programming. It permits a large program to be broken down into smaller, more manageable sections that can be separately compiled into object code. Within a program, objects are structured into logical units that make them easier to manipulate.

Modula-2 has, as you might expect, many components that join to form a modern block-structured language. As Wirth has pointed out, programs are a combination of algorithmic structures and data structures. Modula-2 has well-defined constructs for expressing both structures. In this chapter, we will look at object manipulation in programs and the data structures used to describe these objects.

Data Objects

Data objects are the information a program processes. This information can be numbers, characters, or anything the program problem requires. Every data object in your Modula-2 program has the following properties:

- | | |
|------------|--|
| Identifier | A unique name distinguishing one data item from another. An identifier is given to most variables, any defined types, many constants, and program parts. Identifiers (except for standard identifiers) must be defined in a declaration. |
| Value | Each data item has a value. This value may be constant and unchangeable throughout the program or the data item may take on different values at different times during the course of the program. |
| Type | Each data item has a type that determines its use. The <i>type</i> of a data item dictates the values it can be assigned and the operations that can be performed on it. |

As an example, suppose your program includes the variable *Count*, which can take any whole number as a value. The variable *Count* has the following properties:

- Count* is the *identifier*, or unique name, of this data item.
- The *value* of the variable *Count* is determined at runtime by an assignment and may be changed at any point.
- The *type* of *Count* is INTEGER, which means that it may receive only the values in the range -32K to 32K and that only integer math operations may be performed on it.

There are two categories of types: *unstructured* and *structured*. First, we will look at the elementary data types in Modula-2, which are predefined unstructured types. Then we will see how declarations define a data item's type, identifier, and value. Next, we will discuss user-defined unstructured types. And lastly, for more complicated objects, we will examine the building blocks Modula-2 provides to define larger structures--structured types.

Elementary Data Types

In the previous chapter, we described some of the objects that might be manipulated in a program (numbers, strings, and so on). In addition, we showed that an object, such as a number, may have different representations called types. For simple items, Modula-2 has predefined names that represent these types; for example, INTEGER, LONGINT, REAL, and so forth. These are called elementary data types.

The following eight types are predefined in Modula-2. They are considered basic types and are always available. Generally, variables are declared as one of these types; however, you can also build new types from these basic types.

CHAR type: Range [0C to 377C]. Takes a single character as a value. A CHAR constant is denoted by enclosing a printable ASCII character in single or double quotation marks, or by specifying the octal value of the ASCII character followed by a C (as in the range given here). Sample CHAR values include the following:

"A" "12" "#" " " 1C 377C

BOOLEAN type: Range [FALSE to TRUE]. Assumes one of two logical truth states denoted by the standard identifiers TRUE and FALSE. Use BOOLEAN-type variables when you expect a yes/no- or on/off-type answer. For example, "Is it raining?" has a yes/no answer that could be represented as

Raining = TRUE

or

Raining = FALSE

CARDINAL type: Range [0 to 65535]. Use CARDINAL-type variables anytime values are limited to positive whole numbers or zero, such as for a person's age or address, or lengths and distances. The type CARDINAL is used more frequently than the type INTEGER, and a CARDINAL's range is also double that of an INTEGER in the positive direction. Sample CARDINAL values:

8088 2132 0 64000

INTEGER type: Range [-32768 to 32767]. Use the INTEGER type when values are expected to drop below zero (for instance, in temperature ranges). In practice, this type is not used often since most values in programs are non-negative. Examples of INTEGER values:

8088 -2132 0

BITSET type: Range [0..15] (bit-wise). BITSET is a predefined set type (see "Set Types" later in this chapter), whose primary use is in performing bit-wise operations (such as masking) on word-length variables. BITSETS (and sets in general) has several predefined operations. Operators provided are *union*(+), *intersection*(-), *difference*(*), and *symmetrical difference*(/). In addition, there are standard procedures (*INCL* and *EXCL*) that are used to include and exclude any element from a set variable, and a reserved word (IN) to test whether a particular element is a member of a set element.

Sets are designated by enclosing the members in curly brackets, like this: {3,4,6}. You can visualize the type BITSET by comparing the set notation to its binary representation:

{ }	= 0000000000000000
{4,5,6,7,12,13,14,15}	= 1111000011110000
{0,1,2,3,4,5,6,7,8,9,10,11,12,13,14,15}	= 1111111111111111

In the first example, all the bits are clear which means there are no members in the set. In the second example, bits 4 through 7 and bits 12 through 15 are set (the right-most bit is bit 0). In the last example, all bits are set.

REAL type: Range [-6.80565E+38 to 6.80565E+38]. Use REAL numbers to represent objects with a fractional part or to represent numbers that are very large or very small, including money, the distance to the sun (93,000,000 miles), or measurements of temperature and time. Real numbers have 6 digits of precision; for example:

3.0E+8 98.6 0.0 2.3E-11

Turbo Modula-2 supports two types of double-precision numbers, LONGINT and LONGREAL. Double precision means the number is stored in twice the storage space as its single-precision counterpart, resulting in greater accuracy since the number is represented with more bits. Compare the ranges and number of significant digits in the single- and double-precision numbers.

LONGINT type: Range [-2147483648 to 2147483647]. Double-precision integers are used when a whole number greater than 65535 or less than -32768 is needed. For example, the size of a file may be expressed with a LONGINT as 100000L bytes. Other examples of long integers include the following:

0L 1234567889L -20000000L

LONGREAL type: Range [-3.5953862697246D+308 to 3.5953862697246D+308]. These are double-precision reals. Variables of this type have a precision of 14 digits; for example, multiplication is accurate to 14 decimal points. Use LONGREALs when numbers must be very large or require great accuracy.

Declarations

In general, Modula-2 programs have a declaration part and a program part. The declaration part must contain all constant, type, and data declarations necessary

to describe the data objects used in the program part. Data items must be defined before they are used.

Constant Declarations

Constants benefit programs in two ways: (1) They place a descriptive name where otherwise obscure numbers may be. (2) They simplify future modifications by limiting certain changes to those constant definitions that have a global effect.

A *constant declaration* associates an identifier with a value and a type. Constant declarations start with the reserved word CONST and are followed by any number of declarations terminated by semicolons. Each declaration statement contains an identifier followed by an equal sign, followed by a constant expression. A constant expression is an expression containing only constants.

In the following declaration, the identifier *Code* represents the value *A* and *Rate* represents the value 1.20:

```
CONST
  Code = "A";
  Rate = 1.20;
  Amount = 350.0 * Rate;
```

These values determine the type of the constant, which for *Code* is CHAR and for *Rate* is REAL. In the third expression, the identifier *Amount* represents the value obtained when the constant expression is evaluated. Notice that the identifier *Rate* must be a constant that has already been defined.

In some cases, the constant's type is not obvious; for example:

```
CONST
  N = 100;
```

Is the constant *N* of type INTEGER or of type CARDINAL? In this case, the answer is both. However, if *N* had been declared as

```
CONST
  N = 60000;
```

then its type would be only CARDINAL, since the value is out of the range for integers.

Type Declarations

Each data item has a type defining the possible values it can assume and the operations that can be performed on it. In addition, a item type implies its structure and how much storage the data item occupies. The type can be one of the predefined types (such as REAL or CARDINAL) or can be a user-defined type.

A *type declaration* begins with the reserved word TYPE and is followed by any number of declarations terminated by semicolons. Each declaration contains an identifier, followed by an equal sign, followed by a type identifier or statement. For example:

```
TYPE
  Ages = CARDINAL;
  Time = INTEGER;
```

Here two new types are introduced into the program. These types can now be used to declare data items with descriptive type names. This method has the same advantage as that of naming constants: By using a descriptive name for a defined type that is to be used several times in the variable declaration, the declaration becomes more readable and easier to modify. For instance, in the previous example, you can change all variables declared as type *Time* from INTEGER to REAL by simply changing the word INTEGER to REAL.

Variable Declarations

Two things are required to describe a variable: a unique identifier and a type. The value of the variable is left undefined so it may assume different values at runtime. Modula-2 has no provision for initialization of variables at load time; all variables must be initialized at runtime.

The *variable declaration* defines the variable's type and identifier and optionally the location in memory where the variable will exist. The declaration begins with the reserved word VAR and is followed by any number of variable declarations terminated by semicolons. Each declaration starts with an identifier list, optionally followed by an absolute location, then by a colon, and lastly by the type

of the variable. The absolute location is specified by a left square bracket ([), followed by an address, and a right square bracket (]). For example:

```
VAR  
Count: CARDINAL;  
SwitchOn: BOOLEAN;  
Screen[0C000H] : ScreenType;
```

This example defines three variables: a CARDINAL called *Count*, a BOOLEAN called *SwitchOn*, and an absolute variable called *Screen* residing at memory location 0C000H (the location must be a constant). This latter variable is of type *ScreenType*. These three variables are completely defined and ready to use in a program (assuming the type *ScreenType* has been defined previously).

User-Defined Unstructured Types

We have already seen some predefined unstructured types; however, there are also user-defined unstructured types, which can be used to define variables that can hold one value at a time.

Users can define three unstructured types: *enumeration*, *subrange*, and *pointer*. In general, user-defined types may be declared directly in the variable declaration, or they may be defined as types in the type definition section.

Enumeration Types

Often, predefined types such as INTEGER or REAL are not sufficiently descriptive for certain applications. For instance, suppose you are trying to describe the four points of a compass. One solution is to decide that North is equal to 1, East is 2, and so on. However, when you are reading and writing the code, you must always remember which number stands for the compass point you need. Also, the variables you are working with are declared as CARDINAL or INTEGER, which tells you nothing about how the variables should be used. Enumeration types (with a user-defined range) provide an elegant solution to this problem. First, you define a new type:

```
TYPE  
CompassPoints = (North, East, South, West);
```

Then, you may declare variables of this type, like so:

VAR

Direction: CompassPoints;

When you use the variable, you simply assign or test for the values you have defined for that type:

```
Direction := North; (* Assign the variable the value North *)
Direction >= East; (* Compare the variable to the value East *)
```

Enumeration types can also be declared directly in the variable section. Once an enumeration type is defined, the names of its values are also declared and cannot be used in another enumeration type. Here the type *Weather* can assume the values *Clear*, *Rain*, *Wind*, or *Snow*. For example,

VAR

```
Weather : (Clear,Rain,Wind,Snow);
Water   : (Clear,Murky,Opaque);
Color   : (Violet,Orange,Green,Black);
```

The use of the value *Clear* in defining *Water* would produce a compiler error because the value *Clear* has already been defined as a value assigned to *Weather*. If the second *Clear* were changed to *Transparent*, then the compiler would accept the declaration.

Enumeration types have an implied order by the way they are declared. In the first example, *Clear* is the first value of the type *Weather*, with an ordinal value of 0. The value *Snow* is greater than *Clear*, *Rain*, or *Wind*. Variables of enumeration types can be compared to find their relative rank. Thus a statement like *Clear < Snow* is TRUE and *Rain > Wind* is FALSE.

Subrange Types

The *subrange type* includes certain variables that can be included in one of the previously mentioned types but which never takes values outside of a restricted range. For example, today's date is of the type INTEGER, but is always in the range of 1 to 31. We can state this by declaring it a subrange of the CARDINAL type. You may define subranges of INTEGER, CARDINAL, or enumeration types. Sample subrange types include the following:

TYPE

```
Today      = [1..31];      (* CARDINAL type with range 1 to 31 *)
Unusual    = [Rain..Snow]; (* Subrange of Weather
```

To specify a subrange of positive integers, you may prefix the subrange with the word INTEGER.

VAR

```
Cents: INTEGER[0..100];
```

However, in general, the compiler will know which type is specified. If the range contains a negative bound, then it assumes INTEGER; otherwise, it assumes CARDINAL.

Pointer Types

Pointers point to another variable that can be of any type, including another pointer. Pointer operations include comparison and assignment, plus a special *dereferencing operator* that allows you to access the object pointed to by the pointer. Pointers are used mainly for pointing to other structured data types such as records or arrays (see "Structured Types" later in this chapter). These pointers are used to build dynamic structures such as linked lists or trees. Sample pointer types include the following:

TYPE

```
CharPtr      = POINTER TO CHAR;
IntPtr       = POINTER TO INTEGER;
CardPtr      = POINTER TO CARDINAL;
RealPtr      = POINTER TO REAL;
SetPtr       = POINTER TO BITSET;
LgPtr        = POINTER TO LONGINT;
BoolPtr      = POINTER TO BOOLEAN;
LRealPtr     = POINTER TO LONGREAL;
PtrPtr       = POINTER TO POINTER TO CHAR;
Person       = POINTER TO PERSONRECORD
```

The value of a pointer is an address in memory. Thus, you may assign a pointer value a 16-bit address. Assuming the following variable declarations:

```
VAR  
cp1, cp2: CharPtr;
```

you may make the following assignments:

```
cp1 := ADDRESS(39283); (* Assign an absolute address to cp1 *)  
cp2 := cp1; (* Make cp2 point to where cp1 points *)
```

Notice that the ADDRESS operator converts the CARDINAL number to a generic pointer type, which can be assigned to any pointer. (See ADDRESS in Chapter 12, "Turbo Modula-2 Reference Directory," or in the module *SYSTEM*.)

Dereferencing is used to access the item pointed to by a pointer. The operator used is the caret (^). To dereference a pointer, you simply append the caret to the end of the pointer's identifier. For example, to assign a character to the location pointed to by the previous declaration, you would use the following statement:

```
cp1^ := 'A';
```

Structured Types

A structured type is a composite of more than one data object. A structured type is analogous to a group of compartments, where each compartment may hold another structured or unstructured type.

There are two types of structures in Modula-2: *arrays* and *records*. Arrays are structured types that contain a predefined number of variables of the *same* type. In contrast, a record contains a fixed number of *different* types.

Both structured types can be accessed as a whole, or the parts of **the structure** can be accessed individually.

Array Types

Arrays have a fixed number of elements of an identical type. The type of each individual element is called the *base type* of the array. Each element of the array is accessed via an index. Arrays are declared as follows:

```
VAR  
IntegerList: ARRAY[1..5] OF INTEGER;
```

IntegerList is an array of five integers. The integer elements can be accessed as *IntegerList[1]*, *IntegerList[2]*, and so forth, where the number in the square bracket is the index.

Arrays can be very descriptive when used in combination with enumeration and subrange types. For example, we could use arrays to record the hours of sunshine in each day of the week by using the following declarations:

TYPE

```
Hours = [1..24];
Days = (Sunday, Monday, Tuesday, Wednesday, Thursday, Friday,
        Saturday);
DaysOfTheWeek = [Sunday .. Saturday];
WeeklySunshine = ARRAY DaysOfTheWeek OF Hours;
```

VAR

```
Sunshine: WeeklySunshine;
```

Then to store the fact that *Monday* was gloomy and *Friday* was sunny, we would use the following:

```
Sunshine[Monday] := 1;
Sunshine[Friday] := 12;
```

Notice that each element can be accessed by an identifying **index**. The second item in the *Sunshine* array is accessed with the value *Monday*.

It is also possible to treat the array as a unit:

```
(VAR
  ThisWeek,LastWeek: WeeklySunshine;
```

With this declaration, and assuming the array *LastWeek* already has some values, we can make this assignment:

```
ThisWeek := LastWeek;      (* More of the same weather *)
```

Multidimensional arrays (such as matrixes) are also possible. Given our previous definition of *Weather*, we can build a structure that will store the weather conditions for each hour of every day for a week. The declaration is as follows:

VAR

WeeksWeather: **ARRAY Days OF ARRAY Hours OF Weather;**

An alternative and identical declaration is

VAR

WeeksWeather: **ARRAY Days,Hours OF Weather;**

We can access the elements in several ways, two of which are shown here:

WeeksWeather[Sunday], [12] := Clear;

assigns *Clear* weather to the 12th hour on *Sunday*; and

WeeksWeather[Monday] := WeeksWeather[Sunday];

assigns *Sunday*'s weather to *Monday*.

Record Types

While an array must consist of elements of identical type, a record may consist of elements or *fields* of different types. Each field can be accessed by a unique name. Records are used to group different types of information into a single data type. Use record variables to handle such information as the attributes of a person for example, his or her name, age, address, phone number, or favorite color. This is easily done with the following record declaration:

TYPE

People = RECORD

Name : ARRAY[1..30] OF CHAR;

Age : CARDINAL;

Address: ARRAY[1..40] OF CHAR;

Phone : ARRAY[1..14] OF CHAR;

PreferredColor: Color;

END;

This is a record containing five fields: *Name*, *Age*, *Address*, *Phone* and *PreferredColor*. Each is a different type: three different arrays (*Name*, *Address*, *Phone*), a CARDINAL, and the user-defined type *Color*. Each field in the record can be accessed with the name of the variable, followed by a period and the na-

of the field to be accessed. For example, if the variable *Friend* is defined as type *People*, like so:

VAR

Friend : *People*;

then we can access each field as follows:

```
Friend.Name      := 'Jon';
Friend.Age       := 23;
Friend.Address   := '234 Anywhere Drive';
Friend.Phone     := '333-3456';
Friend.PreferredColor := Blue;
```

We can also nest records two different ways. We can include the nested record directly in the record definition.

TYPE

```
People = RECORD
    Name : RECORD
        First,
        Middle,
        Last: ARRAY[1..15] OF CHAR;
    END;
    Age : CARDINAL;
END;
```

Or we can declare the type separately and nest it using the new type's name:

TYPE

```
Names = RECORD
    First,
    Middle,
    Last: ARRAY[1..15] OF CHAR;
END;
```

```
People = RECORD
    Name : Names;
    Age : CARDINAL;
END;
```

```

VAR
  Person: People;
BEGIN
  Person.Name.First[1]  = 'L'; (* first initial *)
  Person.Name.Middle[1] = 'J'; (* middle initial *)
  Person.Name.Last[1]   = 'G'; (* last initial *)
END

```

Notice that we can now access the initials of a name by selecting the Name field with each of the nested fields (*First*, *Middle*, *Last*), and then index into the first character of each array.

Variant Records

A record type can also have one or more dynamic variant parts. The variant is selected by the value of a *tag field*.

```

Employee = RECORD
  Name: ARRAY[1..20] OFCHAR;
  Age: CARDINAL;
  CASE OwnHome: BOOLEAN OF
    TRUE: Payment      : CARDINAL;
    OwnedSince     : Date
    |
    FALSE: Rent: CARDINAL
  END
END;

```

In the preceding record, there are two fixed fields, *Name* and *Age*, and a variant part consisting of one or two fields, depending on the value of the tag field *OwnHome*. If *OwnHome* is TRUE, then the fields *Payment* and *OwnedSince* are presumed to exist and the *Rent* field should not be accessed. If *OwnHome* is FALSE, then only the *Rent* field should be accessed.

Note: It is possible to access any of the fields regardless of the value of the tag field; it is up to the programmer to ensure that variant fields are only accessed when the tag field has the correct value.

Sometimes you may not want to specify a tag field. In this case, the tag fie

can be omitted, producing a record with fields that overlap one another. The colon and type must be present to indicate that no tag field was chosen.

```
TYPE
HiLo = RECORD
  CASE : BOOLEAN OF
    TRUE: byte: ARRAY[0..1] OFCHAR;
    |
    FALSE: All: CARDINAL;
  END;
END;
```

VAR

MyWord : HiLo;

The variable *MyWord* of type *HiLo* can be accessed as follows:

```
MyWord.byte[1]:= 'A'; (* Assign character 'A' to the low byte *)
MyWord.All      := 65; (* Assign the CARDINAL 65 to the entire word *)
```

Set Types

A set type defines a collection of related members. The members may be of any scalar type. Within a particular set, all members are of the same type and must have an ordinal value between 0 and 15. For example,

```
TYPE
Color = (Red, Green, Blue);
Digits = SET OF [0..9];
Colors = SET OF Color;
```

The set *Digits* has members that are a subrange of type CARDINAL. The members of *Digits* have ordinal values ranging from 0 to 9. *Colors* is a set consisting of members of an enumerated type. The members have ordinal values ranging from 0 to 2.

The value of a set variable is the collection of members that it contains at any point in time. Operations on sets include assignment and comparison, as well as some special set operations. If we define set variables like this:

VAR

```
ColSet1,ColSet2: Colors;
```

we can make the following statements:

```
ColSet1 := Colors[Red,Blue];
ColSet2 := Colors[Red,Green];
```

Thus, a variable of type *Colors* may take as a value any combination of the three members; for example, {Red,Blue}, {Blue}, {Green,Blue}, and { }, which is the empty set.

Notice that the type of the set must prefix the constant set. If no prefix is present, the type of the set is assumed to be BITSET. The same operations applicable to BITSET are valid for user-defined sets (see BITSET in Chapter 12). Thus, we have the following:

ColSet1 + ColSet2	= Colors[Red,Green,Blue]
ColSet1 - ColSet2	= Colors[Blue]
ColSet1 * ColSet2	= Colors[Red]
ColSet1 / ColSet2	= Colors[Green,Blue]
ColSet1 - ColSet1	= Colors[]
ColSet1 / Colors[Green]	= Colors[Red,Green,Blue]

Procedure Types

Procedure types are an advanced language feature included here for completeness. Thus, you may wish first to read more about procedures and parameters (in Chapter 6), and then come back to procedure types.

Procedures can be assigned to a variable in order to be passed as arguments to other procedures. A procedure-type declaration specifies the number and type of its parameters and an optional function result. The following is a sample procedure type:

TYPE

```
Proc1 = PROCEDURE (CARDINAL, VAR INTEGER);
```

This is a procedure type with two arguments. The first argument is a val-

parameter of type CARDINAL and the second is a variable parameter of type INTEGER. For example, suppose *P* is a procedure variable of type *Proc1*; declared as

VAR

P: *Proc1*;

The procedure variable *P* can take two arguments: one a CARDINAL, the other an INTEGER variable, and only in that order. The variable *P* could then be passed as an argument to another procedure.

A procedure variable may only take on the value of a globally defined procedure (see "Procedure Declarations" in Chapter 6) with an identical parameter list. For example, the variable *P* may take the value of procedures declared as follows:

```
PROCEDURE Test(:CARDINAL; VAR i:INTEGER);  
BEGIN
```

```
END Test;
```

or

```
PROCEDURE Fees(accountNo:CARDINAL; VAR Balance:INTEGER);  
BEGIN
```

```
END Fees;
```

The assignment operation is the only operation allowed on procedure variables.

```
P := Test;  
P := Fees;
```

Assuming that *Fees* is assigned to *P* and that *CurrentBalance* is declared as an INTEGER variable, the procedure *Fees* can be called as follows:

```
P(121,CurrentBalance);
```

The most useful aspect of procedure variables is passing them as parameters to algorithms accessing complicated data structures (such as tree traversal procedures). The benefit is that many different operations can be performed on each element of the data structure, but only one data-access routine need be written to do so.

Chapter 6

Statements

In the previous chapter we discussed how to define data with constant, variable, and type declarations. Now we'll provide you with the remaining item you'll need to write your programs: the statement.

Statements provide a means to handle data and define the flow of control. Modula-2 program statements can be divided into four classes: *assignment*, *conditional*, *repetitive*, and *procedural* (which includes procedure calls and a special flow of control statements). Each statement type will be examined in detail, but first we'll talk about programs in general in Modula-2.

In Modula-2, programs are called modules. Every main program begins with the reserved word MODULE, followed by the name of the program and a semicolon. The module statement is followed by a series of declaration statements that include the data declarations discussed in the last chapter (namely constant, type, and variable declarations), and possibly external and procedure declarations (discussed in the following chapter and at the end of this chapter, respectively).

The following sample program outlines the general form of a Modula-2 program:

```
MODULE Skeleton;
(* External Declarations (Next Chapter) *)

CONST
  (* Constant Declarations (Last Chapter) *)
TYPE
  (* Type Declarations (Last Chapter) *)
VAR
  (* Variable Declarations (Last Chapter) *)

(* Procedure Declarations (This Chapter) *)
```

```
BEGIN
(* Program statements (This Chapter) *)
END Skeleton.
```

The program body follows the declaration section, and contains statements separated by semicolons. The statements are within a block enclosed by the reserved words BEGIN and END, followed by the name of the program and a period. Modula-2 statements can be divided into the following classes:

Assignment	Assignments are made to variables, resulting in a change in the value that the variable holds. The source of the assignment may be any expression that has the same resulting type as the destination. A <i>simple assignment</i> may look like this: $X := 4$.
Conditional	This class groups statements together so that they are executed only if a specified condition is TRUE; for example, »If it's raining, then....« This group includes the IF and CASE statements.
Repetitive	Looping statements group one or more statements that are to be executed a number of times. For example, a loop may »Switch a light on and off ten times.« This group includes the FOR, WHILE, REPEAT, and LOOP statements.
Procedural	These statements call and control subroutines. They include procedure calls and the RETURN statement, which affects the flow of control in procedures.

Assignment Statements

A simple assignment statement replaces the current value of a variable with the result of an expression. The expression may be as simple as a constant or another variable, or it may be a complicated mathematical expression containing constants, variables, and functions. For example:

```
MODULE Assignment;
CONST
Pi = 3.14159;
```

```
VAR
  X,Y: CARDINAL;
  CircleArea,Radius,Z : REAL;

BEGIN
  (* simple *)
  X := 2;
  Y := X + 1;
  Z           := 3.23;

  (* complex *)
  CircleArea := 2.0 * Pi * Radius;
  Z := (CircleArea*(45.2/29.3)+Z*(Z+2.0))/Pi;
END Assignment;
```

The result of the expression on the right must be assignment compatible with the variable on the left. In the first two examples, the expressions result in type CARDINAL; thus *X* and *Y* must be of type CARDINAL. The second set of expressions have REAL results; thus *CircleArea* and *Z* are of type REAL.

Note that *:=* is the *assignment operator* and should be read as »becomes« (or gets), as in *I* becomes *I+1*, or *I* gets *I+2*.

WITH Statements

The WITH statement is used to alter the scope of identifiers, and is often used with assignment statements (though it is not itself an assignment statement). WITH is used with record variables, eliminating the need to name the record identifier for each field accessed within the WITH block. The WITH statement takes the following form:

```
WITH record identifier DO
  Statement sequence
END
```

This statement improves code readability when you are assigning values to fields of record variables, because the code reflects the immediate action of assignment to individual fields without the clutter of additional identifiers. For example:

```

MODULE WithStatement;
TYPE
  Months = (Jan, Feb, Mar, Apr, May, Jun, July, Aug, Sep, Oct, Nov, Dec);
  Date   = RECORD
    Day:[1..31];
    Month:Months;
  Year:CARDINAL
  END;
  Person= RECORD
    Name      : ARRAY[0..30] OFCHAR;
    Age       : CARDINAL;
    Height    : REAL;
    BirthDay : Date;
  END;
VAR
  Friend: ARRAY [1..2] OF Person;

BEGIN
  Friend[1].Name := 'Rodney';
  Friend[1].Age := 28;
  Friend[1].Height := 5.5;
  Friend[1].BirthDay.Day:= 3;
  Friend[1].BirthDay.Month:= Dec;
  Friend[1].BirthDay.Year:= 1958;
  WITH Friend[2] DO
    Name := 'Judith';
    Age := 32;
    Height := 5.4;
    WITH BirthDay DO
      Day:= 27;
      Month:= Oct;
      Year:= 1954;
    END
  END
ENDWithStatement.

```

The first record, *Friend* [1], receives values by simple assignment. Note that for every field accessed the name of the record is present. The second record is accessed using the **WITH** statement. Notice that the structure of the code reflects the structure of the data. The change in scope is accentuated by the indentation used to code the **WITH** statement.

Conditional Statements

IF Statements

The IF statement executes a different *statement sequence* depending on the result of a Boolean expression. The IF statement has the form

```
IF BOOLEAN expression THEN
  Statement sequence
{ELSIF BOOLEAN expression THEN
  Statement sequence}
[ELSE
  Statement sequence]
END
```

The ELSIF (else if) and ELSE parts are optional. If the Boolean expression following the IF is TRUE, the first statement sequence is executed and control continues with the first statement after the END. If the first expression is FALSE and the ELSIF part is present, then its Boolean expression is evaluated. If it is TRUE, then its statement sequence is executed; otherwise, the next ELSIF is evaluated. There may be any number of ELSIF parts after the IF THEN, including none. If all expressions are FALSE, then the ELSE part, if present, is executed. For example:

```
MODULE IfThenElsifElse;

VAR
  SwitchOn,PowerOn: BOOLEAN;
  Status : (OK, Danger, Emergency);

BEGIN
  (* Statements that set Boolean flags *)
  IF SwitchOn AND PowerOn THEN
    Status := OK
  ELSIF NOT SwitchOn THEN
    Status := Danger
  ELSE
    Status := Emergency
  END;
  (* Other statements *)
END IfThenElsifElse.
```

After the Boolean flags have been set, the expression following the IF is evaluated. If it is TRUE, then *Status* gets *OK* and execution continues after the END; otherwise, the ELSIF expression is evaluated. If it is TRUE, then *Status* gets *Danger*. However, if the ELSIF expression is FALSE, the ELSE is executed and *Status* is set to *Emergency*. After either the IF part, the ELSIF part, or the ELSE part is executed, control continues with the first statement after the END.

CASE Statements

The CASE statement executes a different statement sequence according to the value of an expression. The result of the expression is compared against value in a constant list (a list of constants separated by commas) until a match is found or there are no constants left. A CASE statement may have several constant lists each associated with a statement sequence and separated by vertical bars.

The elements of the constant lists must be of the same type as the expression following the CASE. The CASE statement takes the form

```
CASE expression OF
    Constant List : statement sequence
    |
    Constant List : statement sequence
    .
    .
    .
ELSE
    statement sequence
END
```

where a particular statement sequence is executed if the expression results a value contained in the constant list. The optional ELSE part is executed if none of the constants match the expression. If none of the constants match the expression and there is no ELSE part, the behavior of the code is dependent on the setting of the TEST switch when the source was compiled (see »Compiler Options and Switches,« in Chapter 10). If TEST is ON, then the exception CaseSelectError is generated. If TEST is OFF, then execution continues with the instruction after the END that corresponds to the CASE. For example:

```
MODULE Cases;
  VAR
    Temp : (Freezing, Cold, Cool, Perfect, Warm, Hot, Searing);
    HeaterOn,
    AirConditionerON : BOOLEAN;
  BEGIN
    (* Statements which set Temp *)
    CASE Temp OF
      Freezing..Cool : HeaterOn := TRUE
    |
      Perfect, Warm, Hot : HeaterOn := FALSE
    ELSE(* must be Searing *)
      HeaterOn := FALSE;
      AirConditionerON := TRUE;
    END
  END Cases.
```

This example executes one of three statement sequences depending on the value of *Temp*. If *Temp* has a value between *Freezing* and *Cool*, then the heater is turned on. If the value is between *Perfect* and *Hot*, then the heater is turned off. However, *Temp* may take on a value that is in neither constant list. In that case, the *ELSE* part of the *CASE* statement is executed and the heater is turned off and the air conditioner is turned on.

Note: The *CASE* statement is for situations where the lists of constants have adjacent ordinal values. When conditional execution is dependent on nonadjacent values, the *IF THEN ELSIF* statement should be used; otherwise, nonadjacent values used in *CASE* statements will generate a relatively large amount of code.

Repetitive Statements

FOR Statements

The *FOR* statement encloses a sequence of statements that is **repeated** a fixed number of times. The *FOR* statement takes the form

```
FORVar:= First TO Last BY Step DO
  statement sequence
END
```

where *Var* is a variable called the *control variable*, *First* is the initial value of *Var*, *Last* is the final value of *Var*, and *Step* is the increment added to *Var* for each iteration. *First*, *Last*, and *Var* can be any ordinal type, but they must all be of the same type.

The statement sequence is repeatedly executed and the control variable incremented until the control variable equals or exceeds the value of *Last*. For example:

```
MODULE ForLoop;
CONST
  Increment = 0.1;
VAR
  I:CARDINAL;
  Result:REAL;
BEGIN
  FOR I:= 0 TO 10 BY 2 DO
    Result:= Result + Increment
  END
END ForLoop.
```

This example repeats the statement sequence six times, with *I* equal to values 0, 2, 4, 6, 8, and 10. If the optional step value is omitted, then it is assumed to be 1. The step can also be negative; for example:

```
MODULE ForLoopCountDown
VAR
  Time,Add:INTEGER;
BEGIN
  FORTime:= 3 TO -3 BY -1 DO
    Add:= 2 * Time
  END
END ForLoopCountDown.
```

Note that the control variable can take part in the statement sequence but cannot be altered by it.

WHILE Statements

The WHILE statement repeats a statement sequence until a Boolean expression yields a FALSE result. The WHILE statement takes the form

```
WHILE Boolean expression DO  
    statement sequence  
END
```

where the statement sequence is repeatedly executed when the Boolean expression is TRUE. For example:

```
MODULE WhileLoop;  
VAR  
    Control:CARDINAL;  
BEGIN  
    Control:= 4;  
    WHILE Control>0 DO  
        Control:= Control - 1  
    END  
END WhileLoop.
```

This example repeats the statement sequence until *Control* is equal to 0.

Note that the Boolean expression is evaluated before the statement sequence; therefore, it is possible that the statement sequence will not be executed. In the preceding example, this would be true if *Control* was set to 0 instead of 4. Also note that the values that make up the control expression must be affected by some part of the loop, otherwise the loop will never end.

REPEAT Statements

The REPEAT statement replicates a statement sequence until a Boolean expression yields TRUE. The REPEAT statement differs from the WHILE statement in that the expression is evaluated *after* the statement sequence; therefore, the statement sequence is always executed at least once. The REPEAT statement has the form

```
REPEAT
    statement sequence
UNTIL Boolean expression
```

where the statement sequence is repeatedly executed until the expression is TRUE. For example:

```
MODULE RepeatLoop;
VAR
    N: REAL;
BEGIN
    N := 1.0;
    REPEAT
        N := 2.0 * N;
    UNTIL N > 100.0;
END RepeatLoop.
```

This loop repeats until the value of *N* is greater than 100.

LOOP Statements

The LOOP statement repeats a statement sequence until terminated by an EXIT statement. The LOOP statement has the form

```
LOOP
    statement sequence
END
```

This statement is generally used when the termination condition can only be determined in the middle of the loop. Since Modula-2 has no GOTO statement, the EXIT statement is used to terminate the loop and continue control at the statement after the END.

EXIT Statements

The EXIT statement specifies termination of a LOOP statement. A common example of this occurs in sequential processing, where a data object is obtained and then processed. In this instance, a problem can occur if the program can only test if the object was obtained successfully *after* the statement to obtain it. Thus, there must be a test directly after the Obtain statement to see if it was successful. If

it was not, no further processing can be done and control must be passed around the processing statements and go to the statement after the loop. The following example simulates this situation with arrays:

```
MODULE LoopLoop;
TYPE
  DataObject = RECORD
    Name : ARRAY[0..30] OFCHAR;
    Age  : CARDINAL
  END;
VAR
  Record: ARRAY[1..10] OFDataObject;
  WorkRecord : DataObject;
BEGIN
  I := 0;
  Record[10].Name := 'LastRecord';
  LOOP
    I := I + 1;
    WorkRecord := Record[I];
    IF WorkRecord.Name = 'LastRecord' THEN EXIT END;
    (* statements to process the record *)
  END
END LoopLoop.
```

This example presents a more common problem with sequential file processing: The flag in the *Name* field represents the end-of-file condition; thus, the records are continually read and processed until the end-of-file condition is TRUE. At this point the EXIT statement is executed and control continued after the loop's END.

More than one EXIT can be present in a loop. However, if there are too many EXIT statements, it is difficult to understand the meaning of the loop and to determine when and where it will terminate.

Procedural Statements

Procedures are named subroutines and subroutines are pieces of code declared elsewhere and invoked by name. In the following, we will look at how a procedure is called and how it is defined.

Procedure Calls

A *procedure call* is an identifier, optionally followed by a parameter list. The identifier represents a sequence of statements defined elsewhere. When a procedure call is encountered in a statement sequence, the flow of control is transferred to the statements that that identifier represents. A procedure call may have parameters specified by listing variables or constants, which are separated by commas and enclosed in parentheses. These are called *actual parameters*. The following example shows three procedure calls in the body of the module:

```
MODULE ProcedureCall;
  (* Statements that define the procedure Read *)
  VAR
    Number,
    Count      : CARDINAL;
    First      : INTEGER;
  BEGIN
    Read(Number); Read(Count); INC(First);
  END ProcedureCall.
```

The first *Read* statement has one actual parameter, *Number*. This parameter can be different each time the procedure is invoked, as with the second procedure call of *Read*, where *Count* is the actual parameter.

Note that there must be a declaration statement defining *Read*. However, this is not so with the third procedure call in the example, (*INC(First);*). This is because it is a standard procedure, predefined in Modula-2 and callable from anywhere without explicit declarations. (There is a list of standard procedures at the end of this chapter.)

Procedure Declarations

Procedure declarations allow the programmer to define a statement sequence once and use it many times. Procedures allow a series of actions to be abstracted into a single meaningful name. For example, it could take hundreds of lines of code to implement the following procedure calls:

```
ReadChar;
AccessFile;
WriteScreen;
```

However, these three statements abstract the lines of code into names that describe the action the code performs. (This lends itself to the concept of information hiding: The details of each procedure are not relevant to the immediate code; they are hidden in the declarations of each procedure.)

As mentioned at the beginning of this chapter, the declaration section of a module may contain procedure declarations. Each procedure declaration starts with the reserved word PROCEDURE and is followed by the procedure's identifier name, an optional parameter list, an optional function result type, and a semicolon. Next is a declaration section that contains constant, type, variable, and possibly additional procedure declarations. The following example shows a skeleton procedure declaration:

```
PROCEDURE Skeleton((* parameter declarations *));
CONST
  (* Constant Declarations *)
TYPE
  (* Type Declarations *)
VAR
  (* Variable Declarations *)

  (* Procedure Declarations *)

BEGIN
  (* Program Statements *)
END Skeleton;
```

Notice the similarities between this procedure skeleton and the module skeleton at the beginning of this chapter. Both have a declaration section and a body, which is convenient when you need to convert a stand-alone program to a subroutine. The following example shows how procedures are declared and how they can be called repeatedly:

```
MODULE ProcedureDeclaration;
VAR
  x: CARDINAL;

PROCEDURE ExampleProcedure;
BEGIN
  x := x + 1;
```

```
END ExampleProcedure;

BEGIN
  x := 0;
  ExampleProcedure;
  ExampleProcedure;
  ExampleProcedure;
END ProcedureDeclaration.
```

The module *ProcedureDeclaration* has defined one procedure called *ExampleProcedure*. The body of the module calls it three times, incrementing *x* by 1 each time the procedure is invoked.

Parameters

Procedure declarations may have parameters that are passed to the procedure in a procedure call. The parameters are specified in the procedure statement after the procedure identifier, and enclosed in parentheses. This is called a *parameter declaration*.

The parameter declaration defines the names and types of each parameter, which are its *formal parameters*. For example:

```
PROCEDURE foo(a,b,c: INTEGER; d: CARDINAL; e: REAL);
```

Notice that parameters of the same type may be grouped together, as in variable declarations. Each group of parameters is separated by semicolons. In addition to name and type, the declaration determines the kind of parameter that is being passed.

There are two kinds of parameters: *variable parameters* and *value parameters*. Variable parameters link actual and formal identifiers, and are identified by prefixing a group of parameters with the reserved word VAR. Value parameters pass a value as an initial condition to a formal identifier. Since VAR does not appear in the preceding declaration, all of the parameters are value parameters.

To explain the difference between variable and value parameters, let's look at the procedure call. When a procedure call passes a parameter to a subroutine, the program may or may not want the variable it passes to be changed. If it does not, the formal parameter should be declared as a value parameter. But if the call is

expecting the procedure to do some work on the variable and return the new value, then the formal parameter should be a variable parameter.

When a value parameter is passed, a copy of the value is made. Any reference to that parameter will affect only the copy and not the original; thus, constants as well as variables may be passed as value parameters. For example:

```
MODULE ValParameters;

PROCEDURE Power(VAR A:REAL;I:CARDINAL);
VAR
  Temp: REAL;
BEGIN
  Temp:= 1.0;
  WHILE I > 0 DO
    Temp:= Temp * A;
    I:= I-1
  END;
  A := Temp;
END Power;

VAR
  x : REAL;
  e : CARDINAL;
BEGIN
  x := 10.0; e := 2;
  Power(x,e);      (* Raise 10 to the 2nd power *)
  (* Now x = 100.0 *)
  Power(x,3);      (* Raise 100 to the 3rd power *)
  (* Now x = 1000000.0 *)
END ValParameters.
```

(The reserved word VAR in the procedure heading declares *A* to be a variable parameter; the cardinal *I* is a value parameter.)

In the first call, the value parameter is the variable *e* set to the value 2. Notice that the value of *I* is changed. This will not be reflected in the value of *e*; it will remain the same. In the second call, the constant 3 is passed as a value parameter. A copy of this value is made within the procedure, thus it may be treated as a variable within the procedure.

When a variable parameter is passed, the address of the variable is passed. Thus, any reference to that parameter is to the original copy. For this reason, constants cannot be passed as variable parameters.

The following example shows the use of variable parameters:

```
MODULE VarParameters;

PROCEDURE Swap(VAR I,J: INTEGER);
VAR
  Temp: INTEGER;
BEGIN
  Temp:= I;
  I:= J;
  J:= Temp;
END Swap;

VAR
  a,b,c: INTEGER;
BEGIN
  a := 1; b := -1; c := 0;
  Swap(a,b);          (* Now a = -1 and b = 1 *)
  Swap(c,a);          (* Now c = -1 and a = 0 *)
END VarParameters.
```

The integers *I* and *J* are formal parameters. The formal parameters are used inside a procedure in place of the external actual parameters (*a*, *b*, *c*) in the example.

In the first call to *Swap*, *a* and *b* are aliases for *I* and *J*, respectively, which allows the procedure to change the values that *a* and *b* contain. The procedure interchanges the values of *I* and *J*, as well as the values of the external variables *a* and *b*. The second call then interchanges *c* and *a*, whereby *c* ends up with *b*'s original value, *a* gets the value of *c*, and *b* holds the original value of *a*.

Open Array Parameters

Modula-2 has a special feature that allows an array of unspecified size to be passed as a parameter to procedures. These are called *open array parameters*. They are specified in the parameter declaration list in the following form:

ARRAY OF T

where *T* is the type of the item in the array. The lower bound of the array is always 0, and the upper bound is determined by a call to the standard procedure *HIGH*. *HIGH* returns the highest legal index value based on the declaration of the array that was passed, but which has been adjusted so that the lower bound is zero. For example, consider this module:

```
MODULE OpenArrays;
TYPE
Employee = RECORD
    Name : ARRAY[0..20] OF CHAR;
    Age  : CARDINAL;
    Wage : REAL;
END;
VAR
AllEmployees : ARRAY[1..20] OF Employee;
CurrentEmployees : ARRAY[0..9] OF Employee;

PROCEDURE PrintEmployeeRecords(VAR e: ARRAY OF Employee);
VAR
I: CARDINAL;
BEGIN
FOR I := 0 TO HIGH(e) DO
    WITH e[I] DO
        WRITELN(Name,Age:10,Wage:5:2);
    END
END
END PrintEmployeeRecords;

VAR
I: CARDINAL;
BEGIN
FOR I := 1 TO 20 DO
    WITH AllEmployees[I] DO
        Name := 'Hilary';
        Age := 11;
        Wage := 3.35;
    END
END;
```

```
FOR I := 0 TO 9 DO
  WITH CurrentEmployees[I] DO
    Name := 'Nate';
    Age := 2;
    Wage := 3.35;
  END
END;
PrintEmployeeRecords(AllEmployees);
PrintEmployeeRecords(CurrentEmployees);
END OpenArrays.
```

The procedure *PrintEmployeeRecords* will accept and correctly print the two Employee arrays even though they have different bounds. Note that the function *HIGH* gives the highest legally accessible item in the array.

Although *HIGH* provides the largest index, you may still need another way to determine how full the array is. For instance, it is often useful to define procedures that have open arrays of characters, because you need to pass strings of differing length to the same algorithm. However, the character arrays may not be completely filled when passed to the open array; thus, *HIGH* points to the end of the array rather than the end of the string. Modula-2 appends a null character (0C) if a constant string is passed to a procedure with an open array. Thus, the procedure can determine the actual length by scanning for the null character.

With open arrays of other types that may not be completely filled, you must provide some mechanism for the procedure to determine the last valid item in the array. This may be a flag record or an additional parameter that gives the count.

There is an even more powerful version of an open array that allows you to pass any type to a procedure; for more details, refer to Chapter 8. Turbo Modula-2 allows multidimensional open-array parameters, specified as

M: ARRAY OF ARRAY OF T

The high bound of the first dimension is *HIGH(M)* and the high bound of the second dimension is *HIGH(M[0])*. Thus you may find the high bound of the *n*th dimension by passing *HIGH* the array with *n-1* subscripts.

Function Procedures

As mentioned earlier, a procedure may return a result. This is called a *function procedure*. A call to a function procedure is either an expression or part of an expression. The types of values they return must be defined in the procedure statement, as follows:

```
PROCEDURE Foo( ): INTEGER;
```

This procedure heading defines a function that returns a value of type INTEGER. Note that even though there are no parameters, the empty parameter list must be present. This is true for both formal and actual parameter lists. For example, if the preceding function is used to assign a value to an integer variable, then it will look like this:

```
IntVar := Foo( );
```

With the empty parameter list, it is clear whether an identifier occurring in an expression is a variable or a function procedure. This is not a problem with normal parameterless procedures, because they cannot occur in expressions.

As an example, we will change the *Power* procedure we used earlier to a function:

```
MODULE FunctionProcedure;

PROCEDURE Power(A:REAL; I:CARDINAL):REAL;
VAR
  Temp: REAL;
BEGIN
  { op:= 1.0;
  WHILE I>0 DO
    Temp:= Temp * A;
    I:= I-1;
  END;
  RETURN Temp;
END Power;
```

```

VAR
  x: REAL;
  e: CARDINAL;
BEGIN
  x := 10.0; e := 2;
  x := Power(x,e); (* Raise 10 to the 2nd power *)
  (* Now x = 100.0 *)
  x := Power(x,3); (* Raise 100 to the 3rd power *)
  (* Now x = 1000000.0 *)
END FunctionProcedure.

```

The *FunctionProcedure* example has the same outcome as the procedure *Pow* example, but the methods are different. The parameter declaration for *A* has been changed from a variable parameter to a value parameter. Thus, the actual parameter is not changed by the body of the function, but by the assignment statement in the body of the module. Another difference is that the function *Power* has an explicit RETURN statement to define the result and terminate the procedure.

In general, it is a good practice to have only value parameters in a function procedure. When a function's variable parameter or global variable is changed with the course of the function call, it is called a side effect.

RETURN Statements

The value returned by a function procedure is explicitly named by a RETURN statement. This statement takes the form

RETURN value;

where *value* is an expression with the same resulting type as the function procedure. Thus, in an earlier example, the statement

RETURN Temp;

specifies the result of the function. The RETURN statement may also be used in normal procedures with no statement after RETURN. In a normal procedure the RETURN statement has the same effect as if the procedure had come to its final END statement. This permits procedures to be exited at any point. As we

the EXIT statement previously described, too many RETURN statements in a procedure make it difficult to understand.

Examples of RETURN statements:

```
RETURN ;          (* normal procedure *)
RETURN 26;        (* function procedure returning an Integer *)
```

Nested Procedures

As shown in the skeleton procedure, the declaration section can also include additional procedure declarations called *nested procedures*, or *local procedures*. (The procedures described earlier are global procedures.) The following is an example of a nested procedure:

```
MODULE NestedProcedures;

PROCEDURE LevelOne;

PROCEDURE LevelTwo;

PROCEDURE LevelThree;
BEGIN
END LevelThree;

BEGIN (* LevelTwo procedure body *)
END LevelTwo;

PROCEDURE AnotherLevelTwo;
BEGIN
END AnotherLevelTwo;

BEGIN (* LevelOne procedure body *)
END LevelOne;

BEGIN (* Module body *)
END NestedProcedures.
```

This example defines four procedures: one is global (procedure *LevelOne*), and the remaining three are nested. Two are considered local to the procedure

LevelOne; they are *LevelTwo* and *AnotherLevelTwo*. The procedure *LevelTwo* has a local procedure called *LevelThree*. At certain points in this program, only certain procedure identifiers are available to be called. This is because procedures restrict the visibility of identifiers declared within them. This is referred to as the *scope* of an identifier.

Scope of Visibility

Here we present the rules for the *scope of visibility* of an identifier for a procedure. The rules for modules are slightly different and are presented in the next chapter.

The procedure in which an identifier is declared defines the scope of visibility for that identifier. For procedures, the rules are as follows:

- An identifier exists while the declaring procedure is active.
- An identifier is visible within the declaring procedure after the identifier is declared and it is also visible within nested procedures unless redeclare a nested procedure. The new identifier exists until the redeclaring procedure is inactive, then the identifier reverts to the original declaration condition.
- All identifiers that are visible outside of the procedure are also visible inside the procedure, unless redeclared in that procedure.

For example consider the following module:

```
MODULE ProcedureScope;  
  
PROCEDURE A;  
VAR  
  i,j: CARDINAL;  
  
PROCEDURE B;  
VAR  
  j: REAL;  
BEGIN          (* i and j (the REAL) are visible here *)  
END B;
```

```
PROCEDURE C;  
BEGIN (* i and j (the CARDINAL) are visible here *)  
END C;  
  
BEGIN (* i and j (the CARDINAL) are visible here *)  
END A;  
  
BEGIN  
END ProcedureScope.
```

An identifier declared in *A* is visible in *A*, *B*, *C*, and all the procedures nested within *B* and *C*. On the other hand, an identifier declared in procedure *B* is visible only in *B* and any of *B*'s nested procedures, but not in *A* or *C*.

Notice that the identifier *j* is declared in both *A* and *B*. The *A* version of *j* (the CARDINAL) is visible throughout *A*, *C*, and the procedures nested in *C*. In procedure *B*, the *A* version is invisible, being replaced by the *B* version of *j* (the REAL). The *B* version is therefore visible in *B* and all the procedures nested in *B*.

FORWARD Statements

The original definition of Modula-2 specified that identifiers declared at the same level have the same scope. Thus a procedure declared textually at the top of a program could call a procedure that was declared later in the text. This was possible because the first Modula-2 compilers were multipass compilers (in most cases four passes) that could resolve an undeclared reference on a subsequent pass.

Turbo Modula-2 and Wirth's latest compilers are one-pass, recursive-decent compilers, which means there must be additional syntax to allow a single pass. This is done with the FORWARD statement. Its syntax and usage is similar to that of Pascal's. To use the FORWARD statement, simply copy the entire procedure heading, including the parameter list and the optional function result, to the location where it is needed. Then append the FORWARD statement and a semicolon. You need not change the procedure at all. The syntax is as follows:

```
PROCEDURE identifier ( parameter list ) :result type ; FORWARD ;
```

This statement is required for algorithms involving mutual recursion, where one procedure calls another, which then calls the first. For example:

```
MODULE MutualRecursion;

PROCEDURE One; FORWARD ;
PROCEDURE Two;
BEGIN
  WRITELN('Procedure Two'); READLN;
  One;
END Two;

BEGIN
  WRITELN('Procedure Two'); READLN;
  One;
END Two;

PROCEDURE One;
BEGIN
  WRITELN('Procedure One'); READLN;
  Two;
END One;

BEGIN
  One;
END MutualRecursion.
```

Thus, the FORWARD statement is a way of extending the scope of a procedure's identifier, making it available over a larger program area.

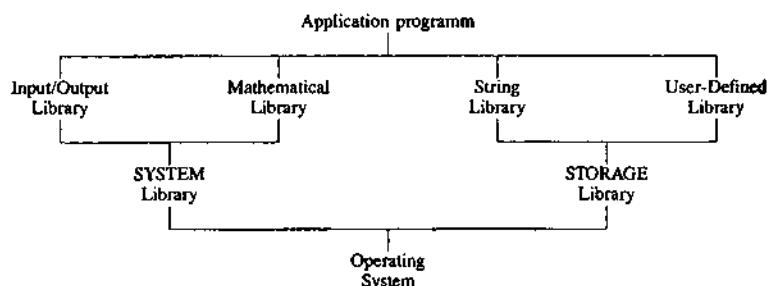
Standard Procedures

Modula-2 includes several predefined procedures that are visible anywhere within a program; in other words, the identifiers do not require importing into a module. (To see examples of these procedures, turn to the respective entries in Chapter 1.) The following is a list of all standard procedures in Modula-2:

ABS(X)	Absolute value of variable X
CAP(Ch)	Uppercase of letters Ch
CARD(X)	Conversion of X to CARDINAL
CHR(X)	The character with ordinal number X
DEC(X)	Replace X by its predecessor

DEC(X,N)	Replace X by its Nth predecessor
EXCL(S,I)	Exclude element I from set S
FLOAT(X)	Conversion of X to REAL
HALT	Halt program execution
HIGH(A)	Upper index of array A
INC(X)	Replace X by its successor ($X := X + 1$)
INC(X,N)	Replace X by its Nth successor ($X := X + N$)
INCL(S,I)	Include element I in set S
FLOAT(X)	Conversion of X to REAL
HALT	Halt program execution
HIGH(A)	Upper index of array A
INC(X)	Replace X by its successor ($X := X + 1$)
INC(X,N)	Replace X by its Nth successor ($X := X + N$)
INCL(S,I)	Include element I in set S
INT(X)	Conversion to INTEGER
LONG(X)	Conversion to LONGINT
MAX(T)	Largest element of type T
MIN(T)	Smallest element of type T
ODD(X)	Returns Boolean TRUE if X is odd
ORD(X)	Ordinal number of X in its type
SIZE(X)	Returns the size in bytes of variable X
SIZE(T)	Returns the size in bytes of type T
TRUNC(X)	Truncate real X to CARDINAL
VAL(T,X)	Value with ordinal number X and type T

Figure 7-1: Hierarchy of Modula-2 Program



The application program is the main module, requiring the services of the *Input/Output*, *Mathematical*, *String*, and *User-defined Library* modules. These modules in turn may require the services of the so-called, »low- level« module (*SYSTEM* and *STORAGE* Libraries), which in turn require the services of the underlying operating system.

Each library module is a self-contained unit that is compiled separately. Users can build their own library modules and use them as if they were provided by the Modula-2 system. When a main module uses library procedures, the facilities of that library are linked to the main program. In general, this process is the same for both predefined and user-defined modules.

As in any hierarchy, the application program may bypass the intermediate modules to use the services of the low-level modules.

When a module requires the services of a library module, it must explicitly state this as the first item in its declaration section. This is done with the **IMPORT** declaration, which can take the following forms:

FROM <module name> IMPORT <identifier list> ;

or

IMPORT <identifier list>;

The first form names a library module after the reserved word **FROM**, then names the identifiers it requires from that module after the reserved word **IMPORT**. The identifier list is terminated with a semicolon.

Chapter 7

Modules

In Modula-2, the module is a syntactical construct that serves to encapsulate certain parts of a program, letting the outside world see only identifiers that are explicitly made visible. Up till now, all the sample modules we have seen have been stand-alone programs, or main modules (which are analogous to entire Pascal programs). Typically, though, Modula-2 programs are made up of many modules, where each support module contains a small amount of code to do its part.

In the Turbo Modula-2 system, the smallest modules have less than 100 lines, the most complex ones nearly 1,500. (The upper limit also mirrors the memory limitations of the computer we're working on [Z80].) How large should a module be? There is no general rule. The task is to break the problem into natural pieces and to have little interconnection between the parts. The amount of interconnection is mirrored by the size of import and export lists; the smaller the lists, the better.

In this chapter, we'll examine the types of modules that can be used in Modula-2.

The Main Module

The main module defines a program. When a main module is running, it usually requires the services of library modules and the underlying operating system. Since these may be viewed as disjoint units (only minimal connection), we may look at a Modula-2 program as hierarchically structured, as shown in Figure 7-1.

The second form imports all identifiers made available by the module specified after the reserved word **IMPORT** and followed by a semicolon. However, unlike the first form, use of imported identifiers requires qualification. This means that any identifiers used from this module must be prefixed with the module name, as follows:

ModuleName.IdentifierName

Typically, this form is used when identifiers imported from different modules have the same name. Thus, the module name specifies which identifier is intended. The following example shows the use of the two forms of the **IMPORT** statement:

```
MODULE ImportLists;
FROM Terminal IMPORT WriteChar, WriteString, WriteLn, ReadChar;
IMPORT Texts;

VAR
  ch: CHAR;
  num: INTEGER;

BEGIN
  WriteString('Enter a character: ');
  ReadChar(ch); WriteLn;
  WriteString('The character you entered was " ');
  WriteChar(ch);
  WriteString(' " . '); WriteLn;

  Texts.WriteLine(Texts.output);
  Texts.WriteString(Texts.output,'Enter an integer and')
  Texts.WriteString(Texts.output,'press return: ');
  Texts.ReadInt(Texts.input,num);
  Texts.WriteString(Texts.output,'The number you entered was ');
  Texts.WriteInt(Texts.output,num,2);
  Texts.WriteLine(Texts.output);
END ImportLists.
```

The facilities of the two library modules *Terminal* and *Texts* are used in the program *ImportLists*. Individual identifiers are imported from the module *Terminal*: *WriteString*, *WriteChar*, *ReadChar*, and *WriteLn*. All facilities from the module

Texts are made available, though they must be qualified. The identifiers used from *Texts* are *output*, *input*, *WriteLn*, *WriteString*, *ReadInt*, and *WriteInt*. Notice that some procedures are available in both modules, which is why the module *Texts* was imported as a whole and then its identifiers qualified. (For more details on the operations of *Texts* and *Terminal*, refer to Chapter 11, »The Standard Library,« or the specific procedures in Chapter 12.)

Library Modules

Library modules have two advantages over stand-alone modules: *information hiding* and *module decoupling*.

In Modula-2, library modules hide the operation and details from the module user; information needed to use the library facilities is separate from the implementation details.

In the preceding example, we needed to know the types and number of parameters of the *WriteString* procedure in order to use it in the module *Terminal*. The declaration

```
PROCEDURE WriteString(s: ARRAY OF CHAR);  
(* Write String to Screen *)
```

tells us that the procedure *WriteString* takes an array of characters and prints them to the screen. This information is contained in a *definition module* (explained in the next section). Thus by looking at the definition module, we can determine how to use its facilities.

In addition to information hiding, library modules allow the details of operations to be changed without affecting modules that depend on its definition. This is called *module decoupling*. This feature of Modula-2 aids in the development of large programs by minimizing recompilation after code changes. Thus, if a library module is made more efficient, users (or »clients«) of that module will only notice the increase in efficiency and will not need to recompile (unless producing a .COM file) or change their source code. A client need only recompile or change source code when the actual definition (and therefore meaning) of a module changes.

Definition and Implementation Modules

Modula-2 formalizes the separation of a module's definition and its implementation with these two modules: the definition module and the implementation module. Every library module consists of one definition module and one implementation module.

The definition module contains all the declarations of the library module that are to be made public, which may include constants, types, variables, and procedures. These declarations do not contain code, but instead contain the interface needed to use the procedures they export, along with comments that describe how the exported identifiers should be used. Definition modules take the following form, beginning with the reserved words **DEFINITION MODULE**:

```
DEFINITION MODULE <module name> ;  
  
< Import Sections >  
< Export Sections >  
  
END <module name> .
```

where Import Sections declare any external identifiers needed to define exported objects and procedures. The Export Sections include the constant, type, variable, and procedure declarations that define identifiers exported by this library module.

Note: Definition modules may not contain an **EXPORT** list (see «Local Modules,» later in this chapter). All identifiers appearing in the definition module are automatically exported.

The following definition module is an example of a library module that provides random-number facilities:

```
DEFINITION MODULE RandomNumbers;  
(* This module provides a reproducible set of random numbers. *)  
  
PROCEDURE Randomize(NewSeed: Cardinal);  
(* Randomize sets the internal seed to the specified value, thus  
allowing a random sequence to be reproduced when the same value is  
passed to Randomize.
```

```
*)
```

```
PROCEDURE Random( ): REAL;
(* The Random function returns a random real between 0 and 1 *)
```

```
END RandomNumbers.
```

This module provides information about the use of the random-number facilities its implementation provides. By reading this Modula-2 text, clients may incorporate random numbers into their programs without writing their own random-number generator or understanding how it works. The corresponding implementation module hides all the details from the client.

The implementation module holds the code that performs the actions stated in the definition module. Objects declared in the definition module are visible in the implementation module. Thus, constants, types, and variables should not be redeclared in the implementation module; doing so will cause a duplicate definition error during compilation. Also, the procedure heading declared in the definition module must be identical to the procedure heading in the implementation module or the compiler will flag the implementation's procedure in error as »Two different declarations of same procedure.«

Implementation modules must define all procedures mentioned in the definition module or the compiler will flag unresolved identifiers.

Implementation modules can be considered identical to main modules with two exceptions. The first is the form of declaration, as shown in the following:

```
IMPLEMENTATION MODULE <module name> ;
(* Import Declarations *)
(* Data Declarations *)
(* Procedure and Module Declarations *)

BEGIN
(* Initialization Code *)
end <module name> .
```

Note that the module name must be the same one used in the definition module.

The second exception is the optional initialization code, which follows the reserved word **BEGIN**. This code is used primarily to set up variables before any procedure is called or any variable is accessed. Thus, you are guaranteed that this initialization code will be executed before the client modules use any function provided by the library module.

If the initialization code is not present, then the reserved word **BEGIN** is optional. Of course, the reserved word **END**, followed by the module name and a period, must be present.

The following module implements the previously defined module *RandomNumbers*:

```
IMPLEMENTATION MODULE RandomNumbers;
CONST
  multiplier = 100;
  modulus = 257;
VAR
  Seed: CARDINAL;

PROCEDURE Randomize(NewSeed: CARDINAL);
BEGIN
  Seed := NewSeed MOD modulus;
END Randomize;

PROCEDURE RANDOM( ): REAL;
BEGIN
  Seed := (Seed * multiplier) MOD modulus;
  RETURN FLOAT(Seed) / FLOAT(modulus);
END Random;
BEGIN
  Randomize(1986)
END RandomNumbers.
```

To complete the examples presented in this section, we can define a main module that imports the facilities of *RandomNumbers*:

```
MODULE PrintRandomNumbers;
FROM RandomNumbers IMPORT Randomize,Random;
FROM Texts IMPORT output, WriteReal, WriteLn;
```

```
VAR
  Count: CARDINAL;

BEGIN
  FOR Count := 1 TO 10 DO
    WriteReal(output,Random( ),5,3); WriteLn(output);
  END
END PrintRandomNumbers.
```

As mentioned earlier, most Modula-2 programs are made up of several modules. The preceding module depends on two library modules, each made up of two modules (a definition and an implementation module). We have defined the two modules required for random numbers, and thus have the source to both. However, the module *Texts* is a standard library module provided by Turbo Modula-2. This means the source code is not available, though the definition module is listed in the next chapter. So how does the program know the interface and code? This information is stored in compiled form; there is one file for the definition module and one for the implementation module.

Compiled Modules

When a definition module is compiled, the result is stored in a symbol file with the extension .SYM. This file is basically a symbol table containing the names and types of exported identifiers. When the compiler is resolving the import list of a module, it looks in this file to obtain information about the imported identifiers. This information is then inserted into the working symbol table of the module being compiled.

Thus, the compiler can tell if an externally declared identifier is being used correctly, at least in a syntactic sense. This concept is referred to as type-checking across compilation units. To compile a module that imports identifiers, the compiler must have access to the identifier's corresponding .SYM file. For example if there is an import from a module named *Texts*, then there must be a file named *Texts.SYM*.

When an implementation module is compiled, its result is stored in a special code file with the extension .MCD. This file contains the code that is linked in when a program is either loaded or linked.

For convenience, Turbo Modula-2 provides a library manager that groups the

.SYM and .MCD files of various library modules into one file with the extension .LIB. For example, all of the standard library modules provided with Turbo Modula-2 are in the library file called SYSLIB.LIB. (For more information on the library manager, refer to Chapter 10, »System Operations.«).

Opaque Export

Sometimes library modules must ensure that variables of an **exported type** are only manipulated by procedures of the library module itself.

Consider a module that exports a file type. One approach is to export the whole structure of this record, making all fields accessible to user programs. If, however, a program updates such descriptor records, chaos may ensue (in the file example, valuable information on disk may be destroyed). In such cases, we would like to hide the fields of the record.

To avoid such problems, Modula-2 offers *opaque export*. This means a name of a type is exported without giving its structure (the actual structure is given in the corresponding implementation module). Thus, user programs may declare variables of that type, but the only operations applied to them will be the provided library procedures. The following example may help clarify this:

```
DEFINITION MODULE Files;
TYPE
  FILE;

PROCEDURE Open(VAR f: FILE; filename: ARRAY OF CHAR);
(* Other declarations *)

END Files.

MODULE OpaqueExport;
FROM Files IMPORT FILE, Open;
```

```

VAR
  f : FILE;
BEGIN
  Open(f, "A:TESTDATA.DAT")
END OpaqueExport.

```

Since the definition module *Files* does not provide information about the type *FILE* the user program cannot apply any operations to it, except for those explicitly provided by *Files* (which in this case is *Open*).

When compiling the example *OpaqueExport*, the compiler must know the space requirements of the variables of opaque types, which in this case is *FILE*. However, it can only inspect the definition module. Therefore, to be able to find out the space needs of opaque-type variables, a restriction must be imposed on opaque export: Only pointer types can be exported in opaque mode.

Local Modules

The remaining module type in Modula-2 is a *local module*, one that is nested within either a main module or an implementation module. Local modules serve to hide details of some task or object from the surrounding environment. Since they are nested, they cannot be separately compiled. Additionally, they have special scope rules.

Local modules may be declared anywhere a procedure declaration is permitted. Thus, local modules may appear in the declaration section of either a module (main, implementation, or local) or a procedure. In general, modules and their contents are static, meaning they exist throughout the duration of the program that surrounds them. There is one exception: When a module is declared local to a procedure, then the module exists only while the enclosing procedure is active.

A local module declaration is similar to a main module and takes the following form:

```

MODULE <module name> ;
(* Import Declarations *)
(* Export Declarations *)
(* Data Declarations *)
(* Procedure and Module Declarations *)

```

BEGIN

```
(* Initialization Code *)  
END <module name> ;
```

Its export declaration is what makes it different from other modules. These declarations cause the listed identifiers to become visible in the surrounding environment. Both import and export declarations cause the scope of identifiers to be altered in both the external and internal environment of the local module.

Scope and Local Modules

Unlike procedure walls, local module walls are not only opaque from the outside but also from the inside. Identifiers that are to be used within a local module must be imported (via an import declaration), and identifiers declared inside a local module that must be visible outside must be exported (via an export declaration). The rules for importing and exporting are completely symmetrical:

- An identifier exists while the declaring module is active. The declaring module is active while its surrounding environment is active.
- An identifier is only visible in a local module if declared by the module or imported from the surrounding scope.
- An identifier declared in a local module is only visible **within the module**, unless explicitly exported.

The form of import declaration for local modules is identical to that for other kinds of modules. The major difference is that any identifier in the surrounding scope may be imported; thus, you may find import lists grouping many different identifiers from many different sources within the same list.

The export declaration that makes identifiers visible outside the local module has the following form:

```
EXPORT <identifier list> ;
```

where the identifier list may contain constants, types, variables, and procedures. The identifier list is prefixed with the reserved word EXPORT and terminated with a semicolon.

An export declaration introduces an identifier to the next highest scope. For example, an export declaration in a local module would introduce an identifier to the surrounding environment. And an import declaration would introduce an identifier into the local module from the next highest scope.

This modular separation hides many details from the programmer and protects the abstraction of the data structure, which is defined in the local module.

The following stack example demonstrates this:

```
MODULE LocalModules;
FROM Terminal IMPORT WriteString, WriteLn;
TYPE
  StackElement = CARDINAL;

MODULE AbstractStack;
IMPORT StackElement, WriteString, WriteLn;
EXPORT Push, Pop;

CONST
  StackSize = 100;
VAR
  SP : [0..StackSize];    (* The stack pointer *)
  Stack : ARRAY [0..StackSize-1] OF StackElement;

PROCEDURE Push(item : StackElement);
BEGIN
  IF SP < StackSize THEN
    Stack[SP] := item;
    INC(SP);
  ELSE
    WriteString('ERROR: Stack overflow.');?>
    WriteLn; HALT;
  END
END Push;

PROCEDURE Pop(VAR item: StackElement);
BEGIN
  IF SP > 0 THEN
    DEC(SP);
    item := Stack[SP]
```

```
ELSE
  WriteString('ERROR: Stack underflow'); WriteLn; HALT;
END
END Pop;

BEGIN (* initialization *)
  SP := 0;
END AbstractStack;

VAR
  i: CARDINAL;
(  WIN (* main program for LocalModules *)
  Push(2);
  Pop(i);
  IF i = 2 THEN WriteString('The stack worked!') END
END LocalModules.
```

Since the stack pointer variable *SP* is not exported, it is not visible outside the module. Note that all objects declared outside of the module must be imported if they are used. The only exception to this rule are standard identifiers, such as *INTEGER*, *BOOLEAN*, *ORD*, and so forth. Notice the use of the standard procedure *HALT* to stop execution when a stack error occurs.

There is a further facility to restrict the number of visible identifiers: *qualified export*. If we had written

```
MODULE AbstractStack;
IMPORT ...
EXPORT QUALIFIED Push, Pop;
```

Then the procedures *Push* and *Pop* would have to be denoted by *AbstractStack.Push* and *AbstractStack.Pop*, respectively, in the surrounding scope.

Since it is sometimes difficult to avoid name clashes of different modules' identifiers, qualified export allows you to clarify ambiguous names. Note that identifiers exported from library modules are always exported in a qualified mode, but the qualification is usually overridden with the **FROM** statement.

Local modules may also override the qualified statement by importing with the **FROM** clause. This allows local modules to share data without cluttering the next

highest scope with unnecessary identifier names. The following example demonstrates **QUALIFIED** exports:

```
MODULE t;

MODULE one;
EXPORT QUALIFIED a;

MODULE two;
EXPORT QUALIFIED a;
VAR
  a: CARDINAL;
BEGIN
  a := 4;
END two;

MODULE three;
FROM two IMPORT a;
BEGIN
  a := 5;
END three;

BEGIN
  WRITELN('one',two.a);
END one;

BEGIN
  WRITELN(one.a);
END t.
```

Note: In the main body of module one, the identifier *a* must be qualified; however, it is not qualified in the body of module three.

Chapter 8

Low-Level Facilities

Modula-2 is a strongly typed language, an aspect that greatly contributes to its programming safety, yet one that is often too restrictive for system programming. Low-level facilities are the avenue to system programming in Modula-2. They include type-transfer functions, special types of the pseudomodule **SYSTEM**, absolute addresses, and coroutines and interrupts. It is helpful to restrict their use to small sections of code grouped into one module, so that when the program is moved to another system only the one module need be changed.

Before we begin discussing low-level facilities, though, we should consider some of the decisions that were made in implementing Turbo Modula-2. These decisions have a direct impact on certain low-level facilities, such as type-transfer functions and **SYSTEM** types. The general philosophy at Borland is small and fast; thus we have restrictions like word alignment of variables and benefits such as register variables.

Register variables are used whenever possible. Programmers can hand-optimize procedures by declaring important variables in the first four words of local storage. Only simple variables of unstructured types are chosen as register variables. If there are not four words of storage declared in a procedure, then the procedure's parameters are placed in registers. Consider the following procedure:

```
PROCEDURE RegVars(x,y: CARDINAL);
  VAR
    i,j,amount: CARDINAL;
  BEGIN
END RegVars;
```

In this procedure, three words of local storage (*i,j,amount*) and one word of parameters (*x*) would be allocated to the registers as follows:

<i>x</i>	sits in BC
<i>amount</i>	sits in DE
<i>j</i>	sits in BC'
<i>i</i>	sits in DE'

A consequence of using register variables is that while a variable sits in a register that variable will have no address. We impose the following restriction:
Simple variables of unstructured type are not accepted by ADR.

Of course, these variables do not have to be kept in registers during the entire procedure. For example, they may be placed in memory if a procedure is called. However, if a procedure has no procedure calls, then the register variables will not be moved to memory, and the procedure can be very efficient and fast.

Word alignment provides both advantages and disadvantages. Word alignment means placing all variables on word boundaries, and the variables that require less than a word of storage are allocated a full word anyway. Thus, a character that needs only 1 byte of storage is allocated one word.

The reason for allocating variables in a word-aligned fashion is that it makes M-code and the M-code interpreter highly efficient and small. Since the M-code interpreter is part of the runtime system (always present), it makes sense to use M-code for non-time-critical operations. This is because M-code takes up approximately one-third the space of native (machine) code, allowing more code to be fit into Z80 memory space. Turbo Modula-2 allows a single program to be made up of both M-code and native code modules so that very large and efficient programs can be written.

The choice to use an efficient M-code interpreter, however, does impose certain restrictions on accessing memory. To keep the interpreter efficient, it may only access entire words at a time. Thus variables are word aligned, which leads to the following restrictions:

- Byte-sized array elements cannot be substituted for VAR parameters.
- Byte-sized array elements cannot be arguments for INC or DEC.

Notice these restrictions are only for byte-sized array elements. This is because byte-sized array elements are not word aligned, they are packed. This is apparent when contrasting arrays and records. A record variable that contains two character fields will always occupy two words of storage. An array that is declared to contain two character elements is always packed into 2 bytes (one word).

When considering certain low-level facilities, keep in mind that variables are usually word aligned. Word alignment impacts transfer functions that involve characters and bytes and the SYSTEM type *BYTE*.

Type-Transfer Functions

Modula-2 offers facilities to relax strict type-checking, forsaking some portability for the ability to perform system programming.

For this purpose, type identifiers are used as function names. The type-transfer function allows programs to explicitly override the type-checking of the compiler. Type-transfer functions are not meant to perform any computation or manipulation. They merely change the interpretation of the bit pattern contained in their argument.

(Since the machine representations of types are entirely implementation dependent, so are the correspondences between representations. Modula-2 does not, for example, define which bit pattern corresponds to the integer value 15. Therefore, the use of type-transfer functions makes a program nonportable.

Type-transfer functions can only be applied to simple types. Moreover, type transfers are only allowed between types that occupy the same amount of storage. Since different systems may have types of different sizes, the applicability of type-transfer functions depends upon the implementation.

In Turbo Modula-2 implementation on Z80 computers, simple types have three sizes: REAL and LONGINT occupy 4 bytes, LONGREAL occupies 8 bytes, and the rest (including pointers, set, characters, and subranges) occupy 2 bytes. Because of the nature of the M-code interpreter, all data objects in Turbo Modula-2 are word aligned. This means that a single character declaration allocates one word, not 1 byte. Further, a record structure consisting of two characters occupies 4 bytes. Only arrays are packed so that characters are in an adjacent byte. Thus, an ARRAY[0..1] OF CHAR occupies only one word.

(Type-transfer functions are illustrated in the following module:

```
MODULE TypeTransfer;
VAR
  ch : CHAR;
  c : CARDINAL;
  i : INTEGER;
  b : BOOLEAN;
  s : BITSET;
```

```

r : REAL;
l : LONGINT;

BEGIN
  ch:= CHAR(c);
  s := BITSET(i)*BITSET(b); (* bitwise AND operator *)
  r := REAL(l);
END TypeTransfer.

```

This module shows the use of type-transfer functions. The first line in the body allows the CARDINAL *c* to be assigned to the CHAR *ch*. The second line transfers two different types into type BITSET, performs a set operation, and then assigns the resulting set to the variable *s*. The last line transfers the LONGINT *l* to a REAL. No physical conversion takes place in any transfer; the bits of the variables remain the same, but are viewed in a different way.

Type Transfer and Type Conversion

There is a difference to note between type-transfer functions and type conversions. While conversions try to preserve meaning, changing the bit patterns if necessary, type-transfer functions leave the bit pattern unchanged but interpret it differently. For example, in the type conversion call

`LONG(1.0) = 1L`

the real value is truncated and converted to type LONGINT. But in the type-transfer call

`LONGINT(1.0) = 1065353216`

the bit pattern standing for the real value 1.0 is interpreted as a LONGINT number.

In addition, type-transfer functions perform no checking. The statement

`b := BOOLEAN(3)`

where *b* is declared as a BOOLEAN, will compile and execute without any error messages. Execution may, of course, lead to somewhat strange results.

In some cases it may appear as if there is no difference between a type-transfer

function and a type-conversion function. Consider the following code:

```
PROCEDURE TransferAndConversion(i: INTEGER);
VAR
  c : CARDINAL;
BEGIN
  c := CARDINAL(i); (* Type transfer *)
  c := CARD(i);      (* Type conversion *)
END TransferAndConversion;
```

- As long as the value of *i* is positive, the two statements will have the same meaning and effect. However, if the value of *i* is -1, the first statement will execute fine, but the value that *c* receives may not be what you think (65535). And if *i* is -1, the second statement will cause a runtime error because the conversion is impossible using *CARD*--you cannot convert a negative number to a *CARDINAL* value.

If you intend to look at the same bits in a different way, then you need the low-level facilities of type-transfer functions.

Low-Level Types and the Pseudomodule System

Besides type-transfer functions, there are the special types *BYTE*, *WORD*, and *ADDRESS* exported by the pseudomodule *SYSTEM* (so-called because it does not really reside in a library module).

Type *WORD* is compatible with all types that occupy one machine word of storage. Since objects of type *WORD* are considered to have no specific interpretation, no operations (except assignment) may be applied to them. Type-transfer functions must be used to indicate the desired interpretation.

There is also the type *BYTE*, which is compatible with 8-bit types such as *CHAR*. Like *WORD*, the assignment operation is the only one that may be applied to variables of type *BYTE*. Note that whether a variable of type *BYTE* takes only 1 byte or 2 bytes of memory is dependent on the declaration used (see the beginning of this chapter).

The type *ADDRESS* is compatible with all pointer types and has the declaration

```
TYPE  
  ADDRESS = POINTER TO WORD
```

In contrast to type *WORD*, arithmetic operators may be applied to operands of type *ADDRESS*; such operands will behave like CARDINAL operands. This feature is especially useful for storage management algorithms (also called pointer arithmetic). The resulting type from pointer arithmetic is *ADDRESS*.

All knowledge about the pseudomodule *SYSTEM* is built into the compiler. Although there is no actual implementation module that corresponds to the definition module of *SYSTEM*, you may use *SYSTEM* as if it were defined as the definition module presented in Chapter 11, «The Standard Library.»

Untyped Parameters

Modula-2 allows a procedure to declare a parameter that will accept any type passed to it. This is called an *untyped parameter*. Untyped parameters have the following type declaration form:

```
PROCEDURE Foo(chunk: ARRAY OF WORD);
```

Notice that the untyped parameter declaration is similar to an open array parameter declaration. The difference is that it uses the low-level type *WORD* exported from the pseudomodule *SYSTEM*.

Untyped parameters are functionally different from open array parameters in that they accept any type of variable as an actual parameter, not just arrays of one type. You may pass either a single character or a large record structure to a procedure with an untyped parameter. For example:

```
MODULE UntypedParameters;  
FROM SYSTEM IMPORT WORD;  
  
PROCEDURE HowBigIs(chunk: ARRAY OF WORD);  
BEGIN  
  WRITELN('Size is ',SIZE(chunk):3,' High is ',HIGH(chunk):3);  
END HowBigIs;
```

```

VAR
  c : CARDINAL;
  r : REAL;
  s : ARRAY [0..20]OF CHAR;
  t : RECORD
    a,b,c,d: LONGREAL;
  END ;
BEGIN
  WRITE('CARDINAL      : '); HowBigIs(c);
  WRITE('REAL          : '); HowBigIs(r);
  WRITE('String literal  : '); HowBigIs('This is a string');
  s := 'This is a string';
  WRITE('String variable : '); HowBigIs(s);
  WRITE('Record variable : '); HowBigIs(t);
END UntypedParameters.

```

The following are the results of this program:

```

CARDINAL      : Size is   2 High is   0
REAL          : Size is   4 High is   1
String literal : Size is  16 High is   7
String variable: Size is  22 High is  10
Record variable: Size is  32 High is  15

```

The standard function *SIZE* gives you the number of bytes that were passed, and the standard function *HIGH* tells you the index of the last valid word in the open array.

Absolute Addresses

Modula-2 has a facility to explicitly specify the address of a variable and thus override the space allocation scheme used by the Modula-2 system. This is especially useful for memory-mapped I/O devices. In the example that follows, a memory-mapped video screen is declared to reside at address 0C000H:

```

VAR
  screen[0C000H] : ARRAY [0..31],[0..127] OF CHAR;

```

The desired address is specified in brackets after the variable identifier; in every other respect, this is a normal variable declaration. The programmer must ensure

correct and consistent use of such variables. Declaring the previous screen variable on a computer with a different screen would lead to chaos.

A useful application of this facility is to define an array that stretches across all memory. Using the low-level types *WORD* and *BYTE*, we can define an array from the start to finish of memory, like so:

```
VAR
Mem [1]:RECORD
  CASE : BOOLEAN OF
    TRUE: b : ARRAY [1..65535] OF BYTE |
    FALSE: w: ARRAY [1..32767] OF WORD |
  END
END ;
```

Memory locations can be accessed like this:

```
Mem.b[10] := 0C;
Mem.w[CARDINAL(ADR(SomeDataRecordOrArray))+SomeOffset] := 39201;
```

Note that you cannot define an absolute variable at location 0. However, the second byte of a word defined at location 0FFFF resides at location 0.

Coroutines and Interrupts

The pseudomodule SYSTEM provides certain mechanisms to implement *coroutines* and *interrupt handlers*, in a high-level manner. These facilities are machine-specific but are presumed to be in all implementations of Modula-2. The actual synchronization and scheduling of processes must be done by supporting modules. As an example, the standard library module *Processes* offers one possible implementation of coroutines; there are many other ways.

Coroutines

Coroutines can be created by calling the procedure *NEWPROCESS* exported by SYSTEM. Switching between coroutines must be done explicitly by calling the procedure *TRANSFER*. Both the source and destination coroutines have to be identified in the *TRANSFER* statement.

Programs using coroutines can be thought of as consisting of several programs, each with their own program counters. Each program corresponds to a coroutine: One is always active, the others are sleeping or frozen. *TRANSFER* freezes the currently active coroutine and activates another one. Unlike procedures, when a coroutine is reactivated it will always continue processing from the point where it was frozen. In contrast, when a procedure is activated, processing always starts at the beginning.

Note that more than one coroutine may use the same procedure. However, each coroutine must have its own *PROCESS* variable and its own work space. This allows several processes to share the same code.

Interrupts

Usually, an interrupt is an unscheduled jump to some special code triggered by a hardware condition, although, software may also trigger interrupts. The code that receives control *after* an interrupt is called an *interrupt handler*. It services the interrupt and then returns control to the interrupted program.

In Modula-2, interrupts are considered coroutine transfers and interrupt handlers are coroutines. When an interrupt occurs, the currently executing code is suspended, and an unscheduled transfer to the coroutine waiting to service that interrupt takes place. When the interrupt handler is finished, control is returned to the suspended code by using a special transfer statement, *IOTRANSFER*, exported from the pseudomodule *SYSTEM*. This procedure not only acts as a coroutine transfer but also initializes (or re-initializes) the interrupt vector for the next interrupt.

Of course, the coroutine waiting for an interrupt must somehow notify the system which interrupt it wants. For this purpose, interrupt vectors are used. (The Z80 CPU uses this scheme when it is in interrupt mode 2.) Note that it is up to the programmer to ensure that the processor is in the correct interrupt mode and that the interrupt jump table is in the correct location.

Note that the CP/M operating system is not reentrant, making it difficult to use interrupts. To guarantee proper functioning, be certain the operating system is not active when user interrupts are possible.

Chapter 9

Turbo Modula-2 Extensions

Turbo Modula-2 offers several extensions to Wirth's definition that are not covered by standard Modula-2, such as general-purpose *READ* and *WRITE* statements, string comparison and assignment, multidimensional open arrays, and exception handling. These extensions can be suppressed with the embedded compiler option (* -X- *), or from the Options menu. If this option is turned off, the compiler will flag all extensions with a warning message.

Input and Output Extensions

Although the various input/output modules provide procedures for reading and writing data to external devices, Turbo Modula-2 defines four new statements that allow a simplified approach to input/output: *READ*, *WRITE*, *READLN*, and *WRITELN*.

The four statements provide you with a quick way to do output. The following program shows some of the many possible uses for these extensions.

```
MODULE READWRITE;
VAR aNumber: CARDINAL;
    aCharacter: CHAR;

BEGIN
    WRITELN('The READ and WRITE statements',' take any number ',
            'of arguments');
    WRITE('Enter a number and a character: ');
    READ(aNumber,aCharacter);
    WRITE('You may even mix types. The number is ',aNumber);
    WRITELN(' The character is ',aCharacter);
    WRITELN('Number may be formatted, as in Pascal ');
    WRITELN('The number 3 in a field of length of 4: " ',3:4,' " ');
    WRITELN('A real number ',34.556:10:2);
END READWRITE.
```

READ and *WRITE* extensions have essentially the same function as the Pascal equivalents: They eliminate the need to use Modula-2's precise library procedures.

During compilation, the *READ* and *WRITE* statements are translated into calls to the appropriate input/output procedures from the library module *Texts*. Using these statements replaces the need to import the specific procedure from the *Texts* module.

For a more thorough discussion of these statements, refer to the module *Texts* in Chapter 11, »The Standard Library.«

String Extensions

Standard Modula-2 allows the assignment of arrays only if both sides are of the same type, and it completely forbids comparisons of arrays. In Turbo Modula-2, assignments and comparisons of strings are allowed. A string is any variable whose type is an array with elements of type CHAR. The starting and ending bounds do not matter; strings are assigned and compared as if both strings involved start at the same lower bound.

The end of a string is denoted either by the end of the array (if the array is completely filled) or by the null character 0C. Any two strings can be assigned to each other, even if they are defined as different lengths or start at different bounds.

Additionally, any two strings can be compared using the relational operators < >, #, < , > , <=, or >=.

Multidimensional Open Arrays

Turbo Modula-2 allows open array parameters to be of any dimension. This extension is invaluable for programs that use matrixes. The standard procedure *HIGH* will return the highest bound for each dimension. The next higher bound is found by passing *HIGH* the variable with »[0]« appended (see Chapter 6, »Open Array Parameters«).

Error-Handling Extensions

Program errors can be divided into the following three types:

Compiler-time errors	Syntax errors discovered by the compiler (for example, a missing semicolon or a misspelling)
I/O errors	Problems occurring during input/output operations (for example, file not on disk or file not open)
Runtime errors	Errors occurring while executing a program (for example, division by zero or integer type passing out of range)

Compiler errors are corrected before the program can be run. However, I/O and runtime errors are only apparent while executing a program.

Should an error occur while running a program, the computer has two options: (1) to write an error message on the screen and halt the program, or (2) to send an error signal to the program.

The first option is fine for simple programs, but could have disastrous results in real-time applications. The second method requires that the running program test some kind of flag after each operation capable of producing an error. This can clutter up the program's logic and reduce efficiency, particularly if the tested condition occurs only in rare circumstances.

As an example, consider a disk-write operation. Testing for a full disk after each disk write is clearly inconvenient. On the other hand, halting the program unconditionally prevents the running program from reacting to this error.

Pascal solves this problem by using a compiler option that determines the program's reaction to I/O errors. However, this is not possible in Modula-2 since disk writes are performed in a library module not within the control of the compiler. It is evident that a system consisting of several largely independent modules needs some way to signal error conditions.

Exceptions serve this purpose well. We do suggest, however, that you use exceptions only if there is no other appropriate way to handle errors. In general, their use should be confined to the signaling of errors across module boundaries. (There are a number of predefined exceptions in Appendix D.)

Syntax and Semantics of Exception Handling

An exception consists of three parts: the *exception declaration*, which defines the

exception identifier and uses the reserved word **EXCEPTION**; the *exception handler*, which is the programmed response to the error condition; and the **R4** statement, which calls the exception.

An advantage of exception handling is that handlers are present statically in program text, near the location they are needed. In comparison to other error handling methods, exception handlers are much clearer. They also follow the nested structure of the Modula-2 language itself.

Declaration of Exceptions

An exception declaration contains the reserved word **EXCEPTION**, followed by a list of identifiers. Exception declarations have the following form:

```
EXCEPTION <identifier list> ;
```

And the following example is taken from the module *Files*:

```
EXCEPTION
  EndError, StatusError, UseError, DeviceError, DiskFull;
```

All of the usual scope rules of Modula-2 apply here. Exception identifiers can be exported and imported like normal Modula-2 identifiers.

Raising Exceptions

Exceptions are raised when the program detects an error condition; for example, when the module *Files* has detected that the disk is full. Raising an exception will transfer control to an exception handler that is provided by either the user or the system.

A program may raise an exception with the reserved word **RAISE**, followed by the exception identifier and optionally by a string. The **RAISE** statement takes the form

```
RAISE <exception identifier> , <string> ;
```

This next example is taken from the standard module *MathLib*.

```
IF x < 0.0 THEN  
  RAISE ArgumentError, 'Negative argument for Sqrt';  
END ;
```

When an exception is raised, the system looks in the current procedure for a matching exception handler. If none is found, the calling procedure is examined, then the caller of that procedure, and so on, until a matching exception handler is found. This handler is then executed, and the procedure containing the handler is exited. If no handler is found, the system prints the exception identifier's name and the optional message string.

There is an alternate form of the RAISE statement that is only allowed within exception handlers. It is the reserved word RAISE by itself, as shown in the following:

```
RAISE ;
```

This has the effect of passing an exception through one handler and on to the next in the calling chain.

Exception Handlers

Exception handlers are written at the end of procedures and modules to handle exceptions issued by a RAISE statement. The syntax is similar to the familiar CASE statement. Exception handlers have the following form:

```
EXCEPTION  
  <vertical bar>  
  <exception identifier list> : <statement sequence>  
  <vertical bar>  
  <exception identifier list> : <statement sequence>  
  . . .  
ELSE  
  <statement sequence>  
END <procedure identifier> ;
```

Exception handlers consist of the reserved word EXCEPTION, followed by any number of exception cases. Each exception case is separated by a vertical bar, and each consists of a list of exception identifiers followed by a colon and a statement

sequence. In addition, an ELSE part is executed if the raised exception is not any of the exception identifier lists.

In the following example, pretend you are running a program-controlled lab experiment and want to guarantee that the experiment turns off in all cases.

```
MODULE LaserExperiment;
IMPORT MathLib;

MODULE GuardedMath;
FROM MathLib IMPORT Exp;
EXPORT ExpG, UndefinedExp;
EXCEPTION UndefinedExp;          (* Exception Declaration *)

PROCEDURE ExpG(X,Y:REAL):REAL;

BEGIN
  IF (1.E-10 < X) AND (X < 1.E10) THEN RETURN Y/X
  ELSE RAISE UndefinedExp      (* Signal an error *)
  END
END ExpG;

END GuardedMath;

(* Procedure Declarations *)

BEGIN
  SetUpExperiment;
  PerformExperiment;
  (* Normal termination *)
  TurnOffLaser;
  ShutDownLab;
  PrintOutResults;

EXCEPTION
  Undefined Exp: WRITELN("WARNING: Bad Math");
  TurnOffLaser;
  ShutDownLab;
  DumpProgramVars;
```

```
ELSE
    TurnOffLaser;
    ShutDownLab;
END LaserExperiment.
```

There is one explicit exception raised in a special exponentiation routine. The exception handler in the main body of code has two clauses: one to catch the exception explicitly raised in *ExpG*, the other to catch all other possibilities.

If no exception handler is present for a given exception condition, the RAISE statement writes the exception identifier and an optional message to the screen, and halts the program. For example:

```
RAISE ZeroDivision;
```

would react to an error by printing the following message:

```
ZeroDivision in module Division
Press "C" for calling chain >
```

If a more comprehensive error message is desired, an additional string can be included in the RAISE statement. The string can be a literal or any variable whose type is an array with elements of type CHAR. The string will be printed below the exception name, like so:

```
DiskFull in module MYPROG
While processing file OUTDATA
Press "C" for calling chain >
```

A program can respond to an exception by appending a handler to the end of the program unit where an exception may be raised. Such a program unit can be a procedure or the main program. When an exception is raised, normal execution is suspended and the corresponding exception handler is invoked; that is, the handler is executed instead of the rest of the program unit.

Exception Propagation

A scheme requiring an exception handler in the program unit where the exception is raised is really not very useful. After all, you need exceptions to signal your

program that an error has occurred in a called library module. There must be some way to propagate exceptions. If the procedure that raises an exception does not contain a handler for it, the procedure is aborted completely and its caller is searched for a handler.

If the calling procedure also does not contain a handler, the search is continued in the procedure that called the caller, and so on. This continues until a handler for the raised exception is found, or until the main program (or a coroutine) is reached. Note that the order in which procedures are searched for exception handlers duplicates the order in which they are displayed in the calling chain.

If a handler for a raised exception is found in the calling chain, it is invoked and thus replaces the rest of the procedure containing the handler. If all statements in the handler are executed, the procedure returns. Note that a program cannot be restarted at the point where the exception is raised. This restriction makes the exception-handling mechanism relatively safe. The restriction is not severe, since an exception can easily be converted into a flag that indicates an unsuccessful operation.

A handler for an exception can of course raise an exception itself. In this case the exception is always propagated to the next procedure in the calling sequence thus preventing an infinite loop where every time an exception is handled, another is raised. To propagate the handler exception to higher-level procedures, a handler can raise the handled exception again by using the short form of the RAISE statement. If the following statement is found in some exception handler, the handled exception is propagated to the calling procedure (or module):

EXCEPTION

```
EndError,DeviceError: Close(infile); RAISE  
END ReadFile;
```

If a library module should signal an exception, it must first be declared with an exception declaration. An exception declaration can be given anywhere in a program where constant, variable, or type declarations would be legal. And like other identifiers, exceptions can be exported and imported.

Many of the library modules export exception identifiers that can be incorporated into programming error handlers. For example, the following program reads data input from the keyboard and writes it to disk. The exception DiskFu

is declared in the module *Files* and raised when the disk is full (*Texts* is implemented with *Files*).

```
MODULE DataWrite;
FROM Files IMPORT DiskFull;
FROM Texts IMPORT TEXT,OpenText,CloseText;
FROM Comline IMPORT PromptFor;
VAR Value: CARDINAL;
DataFile: TEXT;
Str: ARRAY [0..30] OF CHAR;
BEGIN
Str:= "Input value (99 to stop): ";
(* Request filename and open file *)
PromptFor("Data file name: ",DataName);
OpenText(DataFile,DataName);
(* Request data from keyboard and write to file *)
WRITELN(Str);
READ(Value);
WHILE Value # 99 DO
  WRITE(DataFile,Value);
  WRITELN(Str);
  READ(Value2);
END ;
(* Close completed file *)
CloseText(DataFile);
(*If disk full, die *)
EXCEPTION
DiskFull: WRITELN("DISKFULL - YOU JUST LOST ALL YOUR INPUT");
END DataWrite.
```


Chapter 10

System Operations

Turbo Modula-2 is a menu-driven package similar to Turbo Pascal. However, Turbo Modula-2 is a more complex system, containing more menu options and operating differently than Turbo Pascal.

In this chapter, we'll take a look at the basic system operations of Turbo Modula-2. As shown in Chapter 2, the Turbo Modula-2 main menu looks like this:

Selected drive: A

Work file:

Edit	Compile	Run	eXecute
Link	Options	Quit	liBrarian
Dir	Filecopy	Kill	reName Type >

The first two items on the main menu (Selected drive and Work file) are default values used by the other menu items.

Selected drive

Allows you to log a drive as the default disk drive. This command has the side effect of resetting the drive that is specified, so the BDOS read-only error is avoided. (This command is the same as Turbo Pascal's *L* command.)

Work file

Work file is a file name or partial file name that can be prompted with certain menu commands, such as Compile, Link, Run, Edit, and Find runtime error. It is essentially a default user response in which the user can either specify a whole file name (like A:MYMOD) or a partial string (like B:). These default responses can be overridden; see the section in this chapter, »Avoiding the Menus.«

The Run and eXecute commands are used to execute runnable files from within the shell. The remaining menu items are used for file management. (Note: Each item is initiated by pressing the item's capital letter.)

File-Management Utilities

The following five file-management commands provide most of the file facilities you will need when writing Modula-2 programs. Whenever these commands accept a drive specifier, you may also specify a user area. It takes the form

DU:filename.ext

where *D* stands for the drive letter (*A* through *P*) and *U* stands for a user area (0 through 32). Hard disk users will find this facility invaluable for isolating different projects in different user areas.

Dir

Lets you display a directory. You can mask certain file names by specifying an ambiguous file name at the mask prompt. In addition, the Dir command lists the remaining space on the disk from which the directory is read.

You can use question marks (?) and asterisks (*) as »wild cards« in the mask. A question mark represents an arbitrary letter (including none) and an asterisk represents an arbitrary string. Here are some examples of allowed directory masks:

- | | |
|-----------|--|
| A: | Lists the names of all files on disk drive A: |
| B:*.SYM | Lists the names of all files having extension .SYM on drive B: |
| A:TEXTS.* | Lists the names of all files with the file name <i>TEXTS</i> , regardless of extension |
| T????* | Lists the names of all files on the logged drive whose file name begins with a <i>T</i> and has five letters or less |
| A10:** | Lists the names of all files on drive A: in user area 10 |

If you enter nothing after the prompt and press **Enter** then the names of all files on the currently logged disk drive are displayed.

The files are listed by number, which you can use to refer to individual files or groups of files in the Filecopy and Kill commands. Groups of files are specified as a range, for example, 4-9 or 1-20. If there are more file names than will fit on the screen, you will receive the following prompt:

"C" to continue >

Press **C** to continue the directory listings; pressing any other key will stop it.

Filecopy

Allows you to copy a file to another file. The source can be specified either as a complete or ambiguous file name, or as file numbers; the two methods cannot be mixed. The files can be referred to by number only if a Dir command has been performed. As an example, if copying files by numbers:

>D

```
Directory mask: a:  
1: PIPES      .DEF    4: PIPES      .SYM    7: PIPETEST   .MOD    10:T2      .MOD  
2: PIPES      .MCD    5: PIPETEST   .BAK    8: T         .MOD    11:T3      .MOD  
3: PIPES      .MOD    6: PIPETEST   .MCD    9: T1        .MOD    12:TYPETRAN .MOD  
Bytes Remaining on A: 145K
```

>F

Copy from: 1 2 8-11

Copy to : c:

```
Copying A00:PIPES      .DEF  
Copying A00:PIPES      .MCD  
Copying A00:T          .MOD  
Copying A00:T1         .MOD  
Copying A00:T2         .MOD  
Copying A00:T3         .MOD
```

>

Kill

Allows you to delete a file. The files may be referred to by their number if a Dir command has been performed. When Killing by number, you are prompted for

each deletion. If you specify only a mask (ambiguous file name), then you are not prompted at all. For instance, look at this example of Killing by numbers:

>D

Directory mask: B:

1: SAMPLE .DEF
2: SAMPLE .MCD
3: SAMPLE .MOD
4: SAMPLE .SYM

Bytes Remaining on B: 75K

>K

Kill file: 2 4

Delete B00:SAMPLE.MCD (Y/N)? Y

Delete B00:SAMPLE.SYM (Y/N)? N

Note that since CP/M allows numbers as file names, the Kill command distinguishes between names and file numbers by insisting that a space be placed before file names that consist of numbers (this also applies to Filecopy).

reName

Allows you to change the name of a file. You may only change one name at a time. A warning is given if you overwrite another file. For example:

>N

Rename from: sample.mod

Rename to : sample.mud

Type

Allows you to display the contents of a file on the screen without leaving Turbo Modula-2 or using its editor. (This works the same as CP/M's TYPE utility.) Press **[CTRL] [S]** to pause and **[CTRL] [C]** to stop. For example:

> T

Type file: sample.mod

MODULE sample;

END sample.

Options

Provides you with a collection of infrequently used but occasionally useful options. If you press **[o]** (for Options), the following submenu will appear:

compiler options:

List (ON) Native (OFF) eXtensions (ON)
Test (OFF) Overflow (OFF) Upper=lower (OFF)

Path to search: SYSLIB

Find run-time error

Save current selection Quit

>

The first six items on the option menu are compiler switches, which affect the way source code is compiled. These switches are global in that they affect the whole file; however, they are overridden if the source being compiled has embedded switches.

Each of the six switches has an embedded counterpart. A brief description of each is given here. (For a complete description of the following compiler switches and how to use embedded compiler switches, refer to the section in this chapter,

»Compiler Options and Switches.« A brief description is also given in Appendix C.)

- List toggles compiler listing on and off.
- Native toggles native code generation on and off.
- eXtensions allows or disallows Turbo Modula-2 extensions.
- Test toggles bounds-checking by assignment-compatible assignments assignments (for more information see »Compiler Options and Switches,« later in this chapter).
- Overflow toggles INTEGER overflow-checking.
- Upper=lower toggles case sensitivity on and off in identifiers.

Path to search

This entry holds the names of the libraries files that will be searched to find separately compiled imported modules. Library files are created with the librarian (see the section, »The Librarian,« later in this chapter). This option contains the entry SYSLIB, which refers to the file SYSLIB.LIB that contains the standard library modules.

Find runtime error

This utility helps you find bugs that occur at runtime. By accepting a module name and a PC number, it can determine where a program has stopped. The module name and PC number are obtained from the calling chain that is displayed when the error occurs. (For details on the calling chain, refer to the section Appendix E, »The Calling Chain.«)

Save current selection

This option selects the current switch settings and library path as the new default setting. The current selections are made permanent by writing them to the M2.COM file; thus, this file must always be available.

Quit

This selection returns you to the main menu.

Avoiding the Menus

A menu-driven shell makes life a lot easier for beginners, but for experienced users it can sometimes be tiresome. Therefore, we have included a way to circumvent the menu scheme. If you answer the menu prompt with the space bar, you are allowed to enter the full selection sequence, optionally followed by additional arguments on the same line. After the selected command is executed, you are again placed in the menu where the blank was entered. The following examples demonstrate this:

> C B:MYPROG

Compiles the file B:MYPROG.MOD at once. It is equivalent to this dialogue:

> W

Work File: B:MYPROG

> C

Compile file: B:MYPROG

This next example shows how prompts are avoided:

> D B:

lists the directory of drive B, as the following dialogue would

> D

Directory mask: B:

This method only works on the main menu; you cannot, for example, toggle the native code generation option with the string > ON.

The Turbo Editor

The built-in editor is a screen editor specifically designed for creating program text. If you are familiar with Turbo Pascal or WordStar, you'll need little instruction in the use of the editor. The Turbo editor includes some extensions, plus you can install your own commands on top of the WordStar commands (described in Appendix B, »Installation Procedures«) and the WordStar commands will remain usable.

Using the Turbo editor is simple: After you have defined a work file and pressed **E** (for Edit), the menu will disappear and the editor is activated. If the work file exists on the logged drive, it is loaded and the first page of text displayed. If it is a new file, the screen is blank, apart from the status line at the top.

To terminate a line, press the **Enter** key. When you have filled the screen with text, the top line will scroll off the screen, out of view. You may page back and forth in your text with the editing commands described later in this section.

First, let's take a look at the information the status line provides at the top of the screen.

X:FILENAME.TYP. Shows the drive, user area, and name of the file being edited.

Line n. Shows the line number that contains the cursor, counting from the start of the file.

Col n. Shows the column number that contains the cursor, counting from the left side of the screen.

Char n. Shows the character number that contains the cursor, counting from the beginning of the file.

Insert. Indicates that characters entered on the keyboard are inserted at the cursor position. Existing text in front of the cursor is pushed to the right. The insert mode on/off command, **Ctrl V**, switches this message to *Overwrite*, which means text entered on the keyboard overwrites characters under the cursor, instead of being inserted.

Indent. Indicates that auto-indentation is active. It may be switched off with the auto-indent on/off command, **Ctrl Q I**, in which case this space on the status line is blank.

Operating The Editor

This editor is a full-screen editor, which means you can move the cursor anywhere on the screen and begin writing. This is done by using a special group of control characters: pressing the **Ctrl** key while simultaneously pressing any of the keys, **A**, **S**, **D**, **F**, **E**, **R**, **X**, or **C**.

The characters are arranged on the keyboard in a manner that logically indicates their use. For example, in the following display:

E
S D
X

pressing **Ctrl E** will move the cursor up, **Ctrl X** moves it down, **Ctrl S** moves it to the left, and **Ctrl D** moves it to the right. If your keyboard has repeating key capability, you may hold down the **Ctrl** key and one of these four keys to move the cursor rapidly across the screen.

Editing Commands

The editor accepts and uses many editing commands that move the cursor, page through the text, find and replace text strings, and so on. These commands can be grouped into the following categories:

- Cursor movement commands
- Extended movement commands
- Insert and delete commands
- Block commands
- Find and replace commands
- Miscellaneous commands

Each group contains logically related commands that are described in the following sections. (A summary of the commands is provided in Table 10-1.) The following descriptions consist of a command definition, followed by the default keystrokes used to activate the command. If you would like to redefine the commands, refer to »Installation of Editing Commands« in Appendix B.

Table 10-1 Summary of Turbo Editor Commands

Cursor Movement Commands

Character left	Ctrl	S
Character right	Ctrl	D
Word left	Ctrl	A
Word right	Ctrl	F
Line up	Ctrl	E
Line down	Ctrl	X
Scroll up	Ctrl	W
Scroll down	Ctrl	Z
Page up	Ctrl	R
Page down	Ctrl	C

Extended Movement Commands

Left on line	Ctrl	Q	S
Right on line	Ctrl	Q	D
Top of window	Ctrl	Q	E
Bottom of window	Ctrl	Q	X
To top of file	Ctrl	Q	R
To end of file	Ctrl	Q	C
To beginning of block	Ctrl	Q	B
To end of block	Ctrl	Q	K
To last cursor position	Ctrl	Q	P

Insert and Delete Commands

Insert mode on/off	Ctrl	V
Delete left character	Del	
Delete character under cursor	Ctrl	G
Delete right word	Ctrl	T
Insert line	Ctrl	N
Delete line	Ctrl	Y
Delete to end of line	Ctrl	Q
Delete line up to cursor position	Ctrl	H

Cursor Movement Commands

Block Commands

Mark block begin

Ctrl K B

Mark block end

Ctrl K K

Hide block

Ctrl K H

Copy block

Ctrl K C

Move block

Ctrl K V

Delete block

Ctrl K Y

Read block from disk

Ctrl K R

{ Write block to disk

Ctrl K W

Find and Replace Commands

Find

Ctrl Q F

Find and replace

Ctrl Q A

Repeat last find

Ctrl L

Miscellaneous Editing Commands

Delete file on disk

Ctrl K J

Save file, exit

Ctrl K D

Save, edit

Ctrl K

Quit, no save

CTRL K Q

Tab

CTRL I

Auto-indent on/off

CTRL Q I

Control character prefix

CTRL P

Abort operation

CTRL U

Cursor Movement Commands

Character left

CTRL S

Moves the cursor one character to the left nondestructively (without affecting any characters). When at the start of the line, the cursor will move to the end of the previous line.

Character right

CTRL D

Moves the cursor one character to the right nondestructively (without affecting

any characters). When the last character on the line is reached, the cursor will move to the first character on the next line.

Word left

[CTRL]

A

Moves the cursor to the beginning of the word to the left.

Word right

[CTRL]

F

Moves the cursor to the beginning of the word to the right.

Line up

[CTRL]

E

Moves the cursor to the preceding line. If there is no character in the current column, the cursor is moved to the end of the line.

Line down

[CTRL]

X

Moves the cursor to the proceeding line. If there is no character in the current column, the cursor is moved to the end of the line.

Scroll up

[CTRL]

W

Scrolls the file up one line toward the beginning of the file (the entire screen scrolls down).

Scroll down

[CTRL]

Z

Scrolls the file down one line toward the end of the file (the entire screen scroll up).

Page up

[CTRL]

R

Moves the cursor one page up.

Page down

[CTRL]

C

Moves the cursor one page down.

Extended Movement Commands

The editor provides commands to quickly move to either end of a line, to the beginning and end of the text, and to the previous cursor position. These commands require two control characters to be entered simultaneously: press **[CTRL] [Q]** and then one of the control characters, **S**, **D**, **E**, **X**, **R**, or **C**. Their keyboard arrangement repeats the pattern previously shown.

	E	R
S	D	
X	C	

Left on line

[CTRL]	Q	S
---------------	----------	----------

Moves the cursor to the far left of the screen (column 1).

Right on line

[CTRL]	Q	D
---------------	----------	----------

Moves the cursor to the far right (the end) of the current line.

Top of window

[CTRL]	Q	E
---------------	----------	----------

Moves the cursor to the top of the screen.

Bottom of window

[CTRL]	Q	X
---------------	----------	----------

Moves the cursor to the bottom of the screen.

Top of file

[CTRL]	Q	R
---------------	----------	----------

Moves to the first character in the file unless the top of the file has been paged out, in which case the cursor is moved to the first character in the buffer. To move to the start of a file that has had the beginning paged out, you must use **[CTRL] [K] [S]**, which also saves any changes that were made. If you don't want to save your changes, you must quit (**K** **Q**) and restart the editor.

End of file

[CTRL]	Q	C
---------------	----------	----------

Moves to the last character in the file.

Beginning of block

CTRL	Q	B
-------------	----------	----------

Moves the cursor to the block-begin marker set with **CTRL** **K** **B**.

End of block

CTRL	Q	K
-------------	----------	----------

Moves the cursor to the block-end marker set with **CTRL** **K** **K**. (Block commands only work if a block is fully marked.)

Last cursor position

CTRL	Q	P
-------------	----------	----------

Moves to the last position of the cursor. This command is particularly useful after a find or find/replace operation has been executed and you'd like to return to the last position before its execution.

Insert and Delete Commands**Insert mode on/off**

CTRL	V
-------------	----------

When you write text, you may choose between two entry modes: *Insert* and *Overwrite*. The current mode is indicated in the status line. The **CTRL** **V** command allows you to switch between these modes.

Insert mode inserts characters at the cursor position; existing text to the right of the cursor will move to the right as you write new text.

Overwrite mode may be chosen if you wish to replace old text with new text. Characters entered will overwrite any characters under the cursor.

Delete left character

DEL

Deletes the character to the left of the cursor. This can also be used to remove line breaks.

Delete character under cursor

CTRL	U
-------------	----------

Deletes the character under the cursor.

Delete right word**CTRL** **T**

Deletes the word or part of the word to the right of the cursor. It may also be used to remove line breaks.

Insert line**CTRL** **N**

Inserts a new line at the cursor position.

Delete line**CTRL** **Y**

Deletes the line containing the cursor. No provision exists to restore a deleted line, so use this cautiously.

Delete to end of line**CTRL** **Q** **Y**

Deletes all text from the cursor position to the end of the line.

Block Commands

All block commands are extended commands (that is, the standard command definition consists of two characters). You can use them to move, delete, or copy whole chunks of text, and to perform certain file operations.

A block of text (which you determine) is marked by placing a block-begin marker at the first character and a block-end marker after the last character of the desired text.

The marked block may be copied, moved, deleted, or even written to a disk file. There are also commands to read or write an external disk file into or out of the editor as a block.

Mark block begin**CTRL** **K** **B**

Marks the beginning of a block. The marker itself is not visible on the screen, and the block only becomes visibly marked if the block-end marker is set. (This only applies if your terminal has highlight capabilities.)

Mark block end**CTRL** **K** **X**

Marks the end of a block. As stated previously, the marker itself is not visible on the screen; it is only visible if the block-begin marker is also set.

Hide block**CTRL** **K** **H**

Causes the visual marking of a block to be switched off and removes block markers. After this command, no block markers are visible and all block commands are invalid.

Copy block**CTRL** **K** **C**

Copies a previously marked block to the cursor position. The original block is left unchanged, and the markers are placed around the new copy of the block.

Move block**CTRL** **K** **V**

Moves a previously marked block from its original position to the cursor position. The markers remain around the block in its new position. If no block is marked, nothing happens.

Delete block**CTRL** **K** **Y**

Deletes the previously marked block. No provision exists to restore a deleted block, so use this cautiously.

Read block from disk**CTRL** **K** **R**

Reads a file into the current text at the cursor position, exactly as if it were a block that was moved or copied. The text read in from the disk is marked as a block.

When using this command, you are prompted to give the name of the file to be read (after which you must press **Enter**). The file specified may be any legal file name; no default extension is supplied.

Write block to disk**CTRL** **K** **W**

Writes a previously marked block to a file. The block is left unchanged, and the markers remain in place. When this command is issued, you are prompted for the

name of the file to write to. If the file specified already exists, you are asked if you want to continue before the existing file is overwritten. If no block is marked, you will get an error message.

The file specified may be any legal file name; no default extension is supplied.

Find and Replace Commands

These commands allow you to search for a string of characters or replace one string with another, which is useful when correcting a repetitive error.

Find

CTRL **Q** **F**

Lets you search for any string of up to 30 characters. When you enter this command, the status line is cleared and you are prompted for a search string. Enter the string you desire and terminate by pressing **J**.

Search strings may be edited with the Character left, Character right, Word left, and Word right commands. To recall the previous search string, press **CTRL** **L** when prompted for the new string. A search operation may be aborted with the Abort command, **CTRL** **U**.

The search string may contain any characters, even control characters. Enter control characters into the search string by using a **CTRL** **P** prefix. For example, to enter a Control-A, hold down the **CTRL** key and press **P**, then press **A**. Thus you may include a line break in a search string by including Control-M/Control-J (press **CTRL** **P** **M** and **CTRL** **P** **J**, respectively).

When the search string is specified, you are requested for search options. The following options are available:

- U** Ignores uppercase/lowercase. Uppercase and lowercase alphabetic characters are regarded as equal.
- W** Searches for whole words only; skips matching patterns that are embedded in other words.
- B** Searches backward.

digits Finds the *n*th occurrence of a search string.

For example:

BU Searches backward and ignores uppercase/lowercase. 'Block' will match both 'blockhead' and 'BLOCKADE'.

W Searches for whole words only. The search string 'term' will only match the word 'term', not the word 'terminal'.

B5 Searches backward and finds the fifth occurrence.

Terminate the list of options (if any) by pressing **Enter**; the search will then begin. If the text contains a target matching the search string, the cursor is positioned at the end of the target, unless you're searching backward, in which case the cursor is positioned at the start of the target.

The search operation may be repeated by the Repeat-Last-Find command, **[CTRL] L**.

Find and replace

CTRL **Q** **A**

Lets you search for any string of up to 30 characters and replace it with any other string of up to 30 characters.

When you enter this command, the status line is cleared and you are prompted for a search string. Enter the string you require and terminate by pressing **Enter**. Again, however, control characters entered into the search string must use a **[CTRL] P** prefix. (See the previous example in the »Find« section.)

As mentioned in the »Find« section, search strings may be edited with the Character left, Character right, Word left, and Word right commands. (Refer to the »Find« section for more information.)

When the search string is specified, you are asked to enter the string that will replace the search string. You can enter up to 30 characters; control-character entry and editing is performed as stated previously. If you press **Enter**, the target string will not be replaced but will be deleted.

Finally you are prompted for options. The find-and-replace options include the following:

- G Global search and replace; searches and replaces in the entire text, starting at the current cursor position.
- N Replaces every item without asking *Replace (Y/N)* for each occurrence of the search string.
- U Ignores uppercase/lowercase; uppercase and lowercase alphabetic characters are considered equal.
- W Searches and replaces whole words only; skips matching patterns that are embedded in other words.

For example:

GWU Find and replace whole words in the remaining text. Ignore uppercase/lowercase.

Terminate the list of options (if any) by pressing **[Enter]**; the search-and-replace operation will begin. If found (and the N option is not specified), the cursor is positioned at the end of the target and you are asked

Replace (Y/N)?

You may abort the search-and-replace operation at this point with the Abort command, **[CTRL] [U]**.

The search-and-replace operation may be repeated by the Repeat-Last-Find command, **[CTRL] [L]**.

Repeat last find

[CTRL] [L]

Repeats the latest find-and-replace operation exactly as if all information had been reentered.

Miscellaneous Editing Commands

This section lists the commands that do not fall into any of the earlier categories.

Save and quit

CTRL	K	D
-------------	----------	----------

Saves the text file and also gives the original file (if any) the last name .BAK. You are then returned to the Turbo Modula-2 shell.

Tab

TAB	CTRL	I
------------	-------------	----------

Unlike Turbo Pascal, tab positions are fixed to multiples of eight; however, Modula-2 contains the same Autotab feature present in Turbo Pascal.

Indent on/off

CTRL	Q	I
-------------	----------	----------

The indent feature provides automatic indentation. When active, the indentation of the current line is repeated on each following line. To change this indentation, use the space bar and **—** keys to select the new column.

When indent is active, the message INDENT is displayed in the status line; when passive, the message is removed. Autotab is active by default.

Save and edit

CTRL	K	S
-------------	----------	----------

Saves the file and remains in the editor. Press **CTRL** **Q** **P** to return to where you issued the command.

Quit/no save

CTRL	K	Q
-------------	----------	----------

Terminates the editor. If changes were made to the text, you are prompted as to whether or not you want to abandon the edited file. Press **Y** for yes; otherwise, press any other key.

Delete file

CTRL	K	J
-------------	----------	----------

Allows you to delete files from within the editor. This is handy if you run out of disk space after entering lots of text you want to save.

Abort operation

CTRL	U
-------------	----------

Pressing **CTRL** **U** will let you abort any command whenever it pauses for input.

The Librarian

A typical Modula-2 development system uses a large number of files. There are four files for each library file: two that define it and two that make it usable. It is typical in Modula-2 to write and debug a library routine once and then use it many times. Thus it is the second two files that must be kept around. These are the .SYM and .MCD files, which hold the compiled definition and the compiled code, respectively.

The liBrarian makes the storage of these files easier by allowing you to keep many .SYM and .MCD files in a CP/M file called a library file. Thus, if you keep 10 library modules for your general use, you can store the compiled versions in one file instead of 20.

So, the compiler, the program loader, and the linker all need to know which libraries should be searched for files. This information is given in the search path. As mentioned earlier, the search path can be changed in the Options submenu, which initially consists only of the entry SYSLIB.

Searching Libraries

After a library has been created, the modules included in the library will be invisible to the system until the name of the new library is entered in the »Path to search« selection of the Options submenu.

Libraries are manipulated via the liBrarian, which is selected by pressing **B** in the main menu. The liBrarian first prompts the user to select which library to work on:

```
>B  
(Select library:
```

If you enter MYLIB.LIB, then the following menu will appear:

Selected library: A:MYLIB.LIB

Dir	Include	Copy
Kill	c0mpress	quit
*		

You may enter your selection from this menu after the * prompt. The choice have the following effects:

Include	A CP/M file is included in the specified library. Only the CP/M extensions .MCD or .SYM are allowed.
Dir	Displays all the modules already included in the used library; includes size information.
Copy	Moves files out of the library, either by copying them into another library or to a CP/M file. If the destination name has a drive code, it is considered to be a CP/M file name without a drive code, the name specifies the destination library.
Kill	Terminates a module's existence inside the used library.
cOmpress	Copies all files in a library in such a way that there are no unused »holes.« This command is used quite rarely since it takes a lot of time.
Select	Selects another library.
Quit	Returns the user to the main menu.

The Compiler

The Turbo Modula-2 compiler translates a Modula-2 program (the *source*) into a sequence of instructions (the M-code) that form the code of a virtual machine. They are optimized for compactness, to make possible execution of large programs in the limited address space of 8-bit computers.

The Z80 processor cannot execute M-code directly; instead, every instruction is executed by a special program, the M-code Interpreter. In comparison to directly interpreted language like BASIC, M-code operates much faster. In most applications, it approaches machine speed but needs less memory. For those cases where speed is critical, there is a native code option that tells the compiler to pro-

duce Z80 machine code. (For more on this topic, see the section, »Compiler Options and Switches,« later in this chapter.)

The translated program is called an *object program*. The source can only be read by the programmer, while the instructions of an object program can only be executed by the computer. The translation proceeds in a single pass through the source program; in other words, Turbo Modula-2 uses a one-pass compiler for generating M-code. When native code is desired, an additional pass is executed.

The Turbo Modula-2 compiler serves another purpose: If several modules are designed to work together, the programmer must first prepare one or more definition modules. A definition module contains information about parts of the corresponding implementation module that can be used by other modules. The compiler translates the definition module's information into a symbol file. These files do not contain M-code instruction; they merely reference parts of the corresponding implementation module. Whether a program's text constitutes a definition module or not is decided by the reserved word **DEFINITION** in the header line of the module. (Definition modules usually have the extension .DEF in their file name.)

Operating the Compiler

Before a source program is compiled, its program text is written using the editor and then stored on disk. You will have usually provided Turbo Modula-2 with a work-file name for this purpose. The compiler is started from the Turbo Modula-2 shell by pressing **C**. Assuming a work-file name of DEMO and a default disk drive of B, Turbo Modula-2 will respond by prompting you with the message: Compile file:B:DEMO. You can confirm this by pressing **J**.

Source programs usually have the extension .MOD. This extension can be omitted in the file name, and the compiler will append it automatically.

Note: If a source file has no extension at all, the file name cannot be entered »as is,« since the compiler will then look for a file with the extension .MOD. This problem is solved by appending a single period (.) to the file name. Also note that the Turbo Modula-2 editor assumes the default extension .MOD.

Once the name of the source is entered, the compiler starts execution. It continuously reads the source from disk and, if the List option is on, displays the translated portion of the source on the screen. When compilation is finished, the

object file is written to the disk where the source resides. You can override this convention by including a drive code (letter plus a colon) on the command line somewhere after the name of the source. For example, Compile file:B:DEMO A translates B:DEMO.MOD into A:DEMO.MCD.

The generated file now assumes the same name as the source file. If it is an object module, it is given the extension .MCD; if it is a symbol file, it assumes the type *SYM*.

The Listing

During translation the compiler can produce a listing of the compiled program text. The listing can be switched on and off in the Options submenu. When switched on, it is usually displayed continuously on the screen, constantly updating you about how far the compiler has proceeded in the source text. The listing can be redirected to a disk file or to any of CP/M's logical output devices, such as the printer, by including a redirection argument *n* the command line after the **Compile** command. Remember that an argument for output redirection consists of the > symbol, followed by the name of the desired output medium. For example:

Compile file:B:POWERS >B:DEMO.LST

creates a file B:DEMO.LST and sends the listing to this file.

A program listing could like this:

```

| (* $U- *)
| MODULE Powers;
| FROM Texts IMPORT Done,input;
| FROM SYSTEM IMPORT ADR,MOVE,WORD;
|
| PROCEDURE Power(x: REAL; i: INTEGER): REAL;
| VAR t: REAL;
| BEGIN
0|   t:=1.0'
9|   WHILE i > 0 DO
12|     t := * x;
21|     i := i - 1;
25|   END ;
27|   RETURN t

```

```
27| END Power;
---- Size = 31
|
| PROCEDURE GetNum(s: ARRAY OF CHAR; VAR n: ARRAY OF WORD);
| VAR rec: RECORD
|     CASE : BOOLEAN OF
|         | TRUE   : r: REAL;
|     | FALSE      : c: INTEGER;
|     END
| END ;
|
| BEGIN
| WITH rec DO
11| REPEAT
11|   WRITE(s);
16|   IF SIZE(n) = SIZE(REAL) THEN
22|     READLN(r); MOVE(ADR(r),ADR(n),SIZE(r))
33|   ELSE
38|     READLN(c); MOVE(ADR(c),ADR(n),SIZE(c))
47|   END
50| UNTIL Done(input)
50| END
57| END GetNum;
--- Size = 59
|
| VAR
|   x: REAL; n: INTEGER;
| BEGIN
0| LOOP
1|   WRITELN;
2|   GetNum("X= ",x);
3|   GetNum("N= ",n);
28|   IF n < 0 THEN EXIT END ;
8888.
34|   WRITELN(x :8:4, " to the power of ",n :0,
4|   = ",Power(x,n):8:4);
91| END
91| END Powers.
--- Size = 94
|
```

End of source reached.

Compiled bytes: 208

M-code file D00:POWERS.MCD produced.

Besides displaying the program text, this listing offers other information. After each procedure and each module, the line »Size =« indicates the length in bytes of the M-code that was generated. The line »Compiled bytes:« indicates the total length of the M-code generated by the compiler.

The left margin of the source program text is marked by a column of vertical bars. The numbers at the left of the bars are only given on lines that actually generate M-code. They indicate the amount (in bytes) of M-code that is generated between the start of the procedure or module and the current line. These numbers form a coordinate system (procedure name and offset) that allows you to find the location of an error should one occur during execution. Notice that this offset is only correct with M-code. If the native code generation switch has been toggled on, these numbers are meaningless.

If a program's execution fails, the calling chain specifies the offset of the point of error from the enclosing procedure or module. This offset corresponds (in the case of M-code) directly to the numbers in the compiler listing, thereby indicating the error's location in the source. You can demonstrate this effect by entering and compiling the previous program example. If you then run the program and present input that is too large, the program will fail.

In the case of native code, the offset from the calling chain of a runtime error corresponds to the set of numbers given during the second native code pass of the compiler. The numbers will look like this:

End of source reached.

NAME	START	len
POWER	4	102
GETNUM	110	240
POWERS	354	313

Compiled bytes: 675

Native-code file D00.POwers.MCD produced

s, a runtime error with offset = 53 refers to the second line:

POWER 4 102

since 53 is within the first 106 bytes of the file. Since the procedure power is 102 bytes in length, you know the error occurred about halfway through the procedure. An easier method to finding runtime errors is to use the Find runtime error utility provided in the compiler Options menu (described earlier).

The primary use for the numbers in the listing is to determine the size of procedure and modules. This is helpful when using overlays. Normally you will want to keep the listing option off; using the List option reduces compiler speed by about 30 percent.

Error Correction

The Turbo Modula-2 compiler provides for an interactive error-correction scheme. If an error is detected during the compiler run (for example, a syntax error or a type mismatch), the compiler stops and displays the type of error and the point at which it occurred. You are given the option of loading the source immediately into the editor for correction or stopping the compiler. When correction is completed, the compiler takes over again, resuming its work near where you corrected the error (usually at the start of the same line). This is all done automatically, without having to use the Edit and Compile commands everytime an error occurs. The correction cycle is explained further in the following example.

Assume that the ninth line (labeled by the byte offset 12) of the program listing in the preceding section reads $t = t * x$ instead of $t := t * x$. This line is now illegal, since assignment must be expressed by $:=$ instead of $=$. The compiler will translate up to the first illegal line, displaying all processed lines in the listing. When it detects an error, it lists the currently processed line in full before giving a message and stopping. With the List option on, it would look like this:

```
| (* $U- *)
| MODULE Powers;
|
| PROCEDURE Power(x: REAL; i: INTEGER): REAL;
|   VAR t: REAL;
| BEGIN
```

```

2l  t:=1.0'
9l  WHILE i > 0 DO
12l    t = t * x;
          ^ Error in Syntax:
" := " expected, but " = " found
E(dit, Q(uit >

```

The caret (^) points to the last symbol read before detection of the error. You can abort the compiler run by pressing **[Q]**, or you can correct the error pressing **[E]**. If you select **[E]**, the editor is started for correction, your program is loaded automatically, and the cursor is positioned at the place where the error was detected. In the example, the following line would be in the center of your screen:

```
t = t * x;
```

The cursor (denoted by the ^) is already placed on the faulty “ = ”. You can insert the missing “ : ” at once.

The editing commands available are those in the Turbo Modula-2 editor. Of course, you can also »walk around« in the source, fixing other pieces of text. To make space for a text buffer, the compiler must store some of its internal data on disk before editing takes place. The amount needed varies between 0 and 15K depending on the length of your program. To be on the safe side, reserve a working space of about 15K on your source disk, plus the size of a .BAK file if there is one yet.

Once correction is completed, you can leave the editor. To save your corrections, press **[CTRL] [K] [D]**. As in normal editing, the old version is still present, but now has the extension .BAK. The compiler takes over again and translation is resumed near the initial change in your program. (Pressing **[CTRL] [K] [Q]** will save the old version of your source file and the old .BAK file; all changes you have made will be lost and compilation will be aborted.)

If, after saving your changes and continuing on with compilation, there are still more errors in your program, the correction cycle will be repeated.

Running Out of Memory

Following are two instances of memory shortage and their respective causes and solutions.

When the Compiler Runs Out of Memory

The compiler uses the free memory space for two purposes: for its own runtime stack and the heap (that is, the symbol table). Note that the size of the compiled code is inconsequential since it is written to disk as soon as the compiler's internal buffer gets full. In most cases, it's the symbol table that gets too big, which means too many identifiers are visible at a certain point in time. In order to circumvent space problems, you must reduce the number of identifiers. Following are several approaches:

- First check whether your program imports objects that it doesn't use. Note that unqualified import (for example, **IMPORT Texts**) causes all objects exported by the imported module to be put into the symbol table. The space it occupies equals about the length of the module's .SYM file. It is therefore more economic to import only those identifiers that are actually used; for example,
FROM Texts IMPORT TEXT, OpenText, CloseText;
- Second, try to structure your program. After compiling a procedure, the compiler no longer needs to know about the objects local to the procedure (except for the parameters). Therefore, it can throw away all the information concerning local objects and reuse the space for other purposes. After compiling a local module, the compiler keeps only the exported identifiers. Declare objects as global only as necessary and also declare them as local as possible. Use local modules to hide information that is unimportant to the rest of the program. The number of exported identifiers should be as small as possible. This not only reduces symbol table space, it also makes for a better program.
- As a last resort, you may eliminate symbolic constants and enumerated types, replacing their occurrences by their values. For readability and maintenance, the names should be placed in comments. For example:

```
VAR
  color : (*red,green,blue)*) CARDINAL

BEGIN
  FOR color := 0(*red*) TO 2(*blue*) DO
```

```
...
END
END
```

Note that changing *red* to *0(*red*)* in the whole program can be achieved most efficiently via the editor's global replacement command (**CTRL** **Q** **A**).

When Your Program Runs Out of Memory

When this happens, your program probably has too much data and code to fit into memory. You might use overlays if your program has too much code. When you have compiled to M-code, data is usually more critical than code. Following are some common pitfalls that can be avoided and some techniques to make better use of memory:

- Be sure to close all your files as soon as possible. Closing a file frees the space used for the file buffer and the file descriptor. When you reopen a file without first closing it, you allocate a new buffer and a new file descriptor without first giving back the old one. Thus, the memory occupied by it is never reused. If you do that within a loop, you will quickly run out of memory.
- Passing arrays as value parameters is not bad practice. It ensures that the parameter is not changed by the called procedure. However, it forces the compiler to make a copy of the entire array. If the array is large, this also eats up your memory space. So if you are having space troubles, pass large arrays as VAR parameters instead. This also applies to large records.
- In contrast to other languages (for example, the CP/M version of Turbo Pascal), the data space of a procedure is not allocated statically, but is created only when you enter the procedure. Therefore, it might be advantageous to declare large structures local to a procedure. Here is an example:

```
MODULE example;
VAR a, b : ARRAY [1..10000] OF CHAR;

PROCEDURE p;
BEGIN
  (* statements involving a *)
END p;
```

```
PROCEDURE q;
BEGIN
  (* statements involving b *)
END q;

BEGIN
  p; q;
END EXAMPLE.
```

This module requires 20 Kbytes of data space to run (10 Kbytes for each of the arrays). This can be reduced to 10 Kbytes:

```
MODULE example;

PROCEDURE p;
VAR a : ARRAY [1..10000] OF CHAR;
BEGIN
  (* statements involving a *)
END p;

PROCEDURE q;
VAR b : ARRAY [1..10000] OF CHAR;
BEGIN
  (* statements involving b *)
END q;

BEGIN
  p; q;
END example.
```

In the second example, the arrays *a* and *b* have been declared local to the procedures *p* and *q*. Thus, *a* and *b* never exist at the same time. As soon as *p* is called, 10 Kbytes of memory are used for *a*. When *p* is finished, the space is given back and used for *b* when *q* is called. Thus, you never need more than 10 Kbytes of memory for the data.

- Sometimes you don't know exactly how large an array should be, and using a large one to be safe would take up too much memory. In this situation, you might use a pointer to an array and allocate only as much as you need:

```

MODULE example;
FROM STORAGE IMPORT ALLOCATE;
VAR a, b: POINTER TO ARRAY [1..10000] OF CHAR;
...
BEGIN
(* Let us assume the program knows at this point that a *)
(* and b need only be 2000 characters long *)
    ALLOCATE(a,2000*SIZE(CHAR)); ALLOCATE(b,2000*SIZE(CHAR));
END example.

```

This technique is not without danger, however. The compiler does not know that the array is actually shorter than the declaration says. When the array is accessed, it will therefore check index expressions against the declared bounds and not against the actual bounds. Thus, in the preceding example you can access *a* ^ [8000] without receiving a runtime error. Instead, you will access memory that does not belong to *a*. You may thus destroy other variables, your program, or the operating system, and make your program crash in very interesting ways--use this technique with caution.

Symbol Files

Quite often a module imports objects from other modules. It then contains one or more **IMPORT** statements following the header line. An **IMPORT** statement could look like this:

```
FROM MathLib IMPORT Sin,Cos, Exp;
```

The compiler must »know« how the imported procedures *Sin*, *Cos*, and *Exp* are defined in *MathLib*; for example, in order to check for type mismatches. This is achieved by compiling the definition part of the imported module into a symbol file (a file with the extension .SYM). This file must be on the disk either by itself or in a library file (a file with the extension .LIB, which contains many .SYMs and their corresponding .MCD files).

The example presented expects the file SYSLIB.LIB to be on your work disk. If the importing module is compiled, the compiler must access the symbol files of all definition modules from which objects are imported. In the example, the compiler would search for a file called MATHLIB.SYM, in order to »learn« about

the definitions of *Sin*, *Cos*, and *Exp*. If not found, then the compiler would check the library path and search for any files found there.

All symbol files needed during a compiler run must be online. Specifications concerning the drive where one particular symbol file resides need not be given. This is because the compiler automatically searches all drives that you indicated when you installed the compiler with INSTM2 and searches all libraries in the search path for a needed symbol file.

The fact that all needed symbol files must be online can cause problems when the source files of one large programming project are distributed over several disks. The recommended scheme in this case is to combine all the .SYM files and their .MCD files into a common library. This is done with the Turbo Modula-2 Librarian (described in the previous section).

Compiler Options and Switches

The compiler accepts some options that affect the set of programs that are recognized as valid, the generated object code, and the output of a compiler listing. An option can be set or reset in the source program itself. You can either set options globally by using the switches on the Option menu or set options locally by embedding switches in the code.

This is done by including a dollar sign (\$), an uppercase letter, and either the symbol + or - anywhere between two comment brackets; for example, (* \$T + *). As you would expect, + turns on the option denoted by the uppercase letter; using - turns it off. An option remains set until the opposite option is given.

Options are set to either on or off by default. The default settings are displayed in the Options submenu of the Turbo Modula-2 shell, where they may also be changed for the current session and, if desired, saved for future sessions by using the *S* option.

Default settings are save directly in the M2.COM file. Thus, the M2.COM file must be online. Remember the switches on the option menu are global and may be overridden with embedded switches. All available options appear on the menu as follows:

L ist (OFF) N ative (OFF) eX tensions (OFF)
T est (OFF) O verflow (OFF) U pper-lower (OFF)

Each of these switches has the same effect as their embedded counterparts, but these may be overridden. One exception is the Native switch, which is only effective from this menu. The following describes each switch and presents its embedded code:

\$L The List Option

- \$L+ generates a listing.
- \$L- generates none.

If this option is on, the compiler will emit a listing. The listing appears on the screen unless it is redirected, which is done by entering »> listfile« after the »Compile file« message. For example:

```
>C  
Compile file: B:MYFILE >MYFILE:LST
```

If \$L is turned off, the listing is suppressed; however, the names of procedures and modules will be displayed when encountered.

\$O The Overflow Check Option

- \$O+ checks for INTEGER and CARDINAL overflow.
- \$O- does not check for overflow.

If this option is on, the generated object program checks for an overflow when certain operations are performed. An overflow occurs when the result of an operation is too large (or too small) to be represented in the computer's memory. For example:

```
1.  VAR a,b: INTEGER; c: CARDINAL
2.  BEGIN
3.    a := 20000;
4.    b:=a + 15000;
5.    c:=10;
6.    WHILE c >= 0DO
7.      c := c - 1
8.    END ;
```

In both the fourth and seventh line an overflow will occur. In the fourth line the sum evaluates to 35000, while 32767 is the largest representable INTEGER. In the sixth line the program loops. The CARDINAL *c* takes values from 10 down to 0, then *c* is decremented to -1. Since -1 is not a value that can be represented by a CARDINAL, an overflow occurs in line seven. Overflows are signaled by the runtime error message shown here:

```
OVERFLOW in (Name of module)
Press "C" for Calling Chain >
```

If \$O is turned off, no error message is given on overflow; instead, the result of the operation is computed *modulo* 2^{16} . This method is sometimes desirable for computations; for example, in a random-number generator. Furthermore, suppression of overflow checks can speed up a program considerably. The reader should be warned, however, against some unexpected results, particularly in the case of CARDINAL arithmetic. In the previous example, the variable *c* would be decremented from 0 to 65535; hence the condition *c* $>= 0$ would always be fulfilled and the program would loop infinitely!

Overflow checks can be suppressed only for operations involving CARDINAL and INTEGER addition and subtraction and CARDINAL multiplication. Other operations, such as INTEGER multiplication and all operations dealing with REALs and LONGINTs, should always be checked for overflow or have exception handlers provided to trap errors.

\$T The Test Option

- \$T+ checks array indices and subrange variables.
- \$T- does not check them.

If this option is on, code for several tests is inserted into the generated object program. The tests include checks for subrange variables or array indices lying outside their admissible bounds; functions returning no result; and CASE statements with an ELSE clause, where none of the other alternatives apply. If it is turned off, none of this test code is generated. Like overflows, turning off the test option can speed up a program and make it more compact. As a rule, however, the resulting safety loss does not make up for this advantage until a program is completely debugged.

\$U The Upper=Lower Case Option

- \$U+ does not distinguish between uppercase and lowercase letters.
- \$U- does distinguish between them.

If this option is on, the Turbo Modula-2 compiler treats an uppercase letter exactly like its lowercase equivalent. If it is off, an uppercase letter and its lowercase equivalent are considered to be different characters. For example:

```
VAR n,N: CARDINAL;  
BEGIN
```

```
  if n < N THEN ....
```

If this program fragment is compiled with \$U turned on, the VAR declaration would be illegal because *n* and *N* would denote the same identifier.

Compilation with \$U turned off would cause no problems in this line. However, the third line of the program fragment would then be illegal, since *if* would be considered an identifier rather than the reserved word IF. The Modula-2 standard prescribes that lowercase and uppercase letters are different; that is, that \$U should be turned off. Programmers conscious of the standard can reset this option permanently in the Options menu by pressing **[S]**.

\$X The Extension Option

- \$X+ allows Turbo Modula-2 extensions.
- \$X- does not allow them.

Turbo Modula-2 supports several programming constructs not included in the standard, including the following:

- Pascal's *READ* and *WRITE* statements can be used for readable input and output.
- Exceptions similar to Ada's are supported.
- Two arrays with elements of type CHAR (strings) can be assigned even if they are not of the same type.

- Any two arrays of CHAR can be compared.
- Open arrays can be multidimensional.

These extensions are supported only if \$X is turned on. If this option is turned off, nonstandard constructs are flagged by the compiler as errors.

If a programmer intends to use only standard Modula-2, he can turn off \$X permanently in the options menu and save it.

\$N The Native Code Generation Option

- N ON* generates native code for the Z80.
- N OFF* generates M-code.

Turbo Modula-2 can also generate native Z80 machine code that will execute faster than the M-code normally generated. In particular, floating-point math operations and I/O will probably not benefit from this option. The generated native code will be about three times as large as the equivalent M-code, and may run more than ten times faster.

Since this option applies only to compilation units, it must be specified from the Options menu; using (* \$N+ *) in the text will have no effect.

The Linker

Turbo Modula-2's linker combines several separately compiled modules in a single file. A typical large program system consists of a main module and several other support modules. Before such a program is executed, the main module and all needed support modules are loaded into main memory.

How these modules are loaded depends on the type of output the linker produces. There are two output formats produced by the linker: (1) an .MCD file that contains object code and (2) a .COM file that contains executable code.

The purpose of linking many object files into one object .MCD file is to speed up loading time by eliminating disk searching and the time it takes to open more than one file.

If a program's support modules exist as separate files, they are automatically brought into main memory. Every module involved in the program must be found on disk before it can be read in. The Turbo Modula-2 system optimizes the search for the necessary files, achieving relatively short loading times. Nevertheless, some gain in speed can be attained if all modules, or even only part of them, are contained in a single file.

The second function of the linker is to produce an executable .COM file by combining the main module, all support modules, and a runtime system.

The linker takes several object files with the extension .MCD and combines them into a single file. Its output consists of either another .MCD file or a .COM file. While .MCD files are run from inside the Turbo Modula-2 environment, .COM files are executed directly from the operating system. For example, if the file DEMO exists with both extensions .COM and .MCD, you would use the execute command to run DEMO.COM and the Run command to execute DEMO.MCD. In addition, you could run the DEMO.COM file from your operating system prompt like any other program.

You start the linker by pressing **L** at the main menu. You are then prompted for the program's main module and the name of the linker's output, like so:

Link main module:

Output file:

If any of these arguments do not contain an extension, the default extension .MCD is supplied. If the output file name has an extension, it must be either .COM or .MCD. You must specify the extension .COM if you wish to produce an executable file.

After that, you are asked:

Include all needed modules (Y/N)?

Normally, you would press **Y** (for yes), in which case the linker does the rest of the work for you. It searches for all the modules needed on all the drives that are present in the search path. These modules are combined into one file, which is then written to disk.

If you have several versions of the same object file on different disks and want to pick out one of these to be included in the linked file, you would press **N** (for no) and be prompted like this:

Include modules:

You can then enter all modules you wish to include on one line. The .MCD extension can be omitted; it will be supplied as a default. If you want to include more modules than would fit on a single line, terminate the line by pressing the space bar and a **+**. The »Include modules:« prompt is then repeated:

(.include modules: TEXTS FILES B:MATHLIB +
Include modules: B:WINDOWS B:GRAPHICS

The linker then displays a list of all linked modules (and produces a linked object file), all modules that are referred to but not linked, and all modules that are missing. An executable file is produced only if all modules are accounted for and unresolved.

Note that unresolved references are perfectly legal if the output file is of type .MCD. The loader will bring all modules that are referred to but not linked automatically into memory during a Run command.

It would be a waste of memory to link one of the modules *Tests*, *Files*, *ComLine*, *Convert*, *Loader*, or *Doubles* into a .MCD file, since they already reside in memory. If you let the linker search library modules automatically, all these modules are left out. When linking a .COM file, the situation changes. When a .COM file is started directly from CP/M, the flexible Turbo Modula-2 loader is not available. Therefore, all modules the main program imports from must be included in the linked file.

(The included files usually consist of a single object module. It is possible, however, that they are linked and include several modules. In that case, all modules are transferred to the new linked file.

Linking with Overlays

The linker can produce .COM files with overlays. The scheme is more flexible than the one employed in Turbo Pascal, but a bit more difficult. Every overlay consists of a full module; you need not change the module's source code to make it an overlay.

To link a file with overlays, the user must answer the initial question whether to include all needed modules with »No.« The user must then specify all modules that should be linked, but not overlaid. These modules must be entered on a single line or on several lines terminated by characters. If a module is not preceded by a drive code, its .MCD file is first searched in all libraries of the search path, then on all drives of the system. A name with a drive specification will cause the module to be searched only on the specified drive. After all these modules are processed, the user is asked whether overlays should be used. Typing *Y* will cause the following prompt to appear:

Include overlay module:

Map onto linked module:

First, the name of the module to be overlaid is entered. Next, the module whose space is to be used when it is loaded is entered. Now both modules will be linked to occupy the same memory segment. The prompts will be repeated until the user answers the first question with a single carriage return.

Note: The second module must already have been included in the linked file, otherwise its space will not be available to the overlay module at runtime. Also, more than one overlay module may be mapped onto the same linked module.

As an example, suppose you want to link the module *MAIN* with three overlay modules (*OVER1*, *OVER2*, and *OVER3*.) and one resident module (*SUPPORT*) into the file *MYPROG.COM*. The dialogue with the linker would look like this:

>L

Link main module : MAIN
Output file : MYPROG.COM

Include all needed modules (Y/N)? N

Include modules : SUPPORT OVER1

Linking MAIN
Linking SUPPORT
Linking OVER1

Use overlays (Y/N)? Y

```

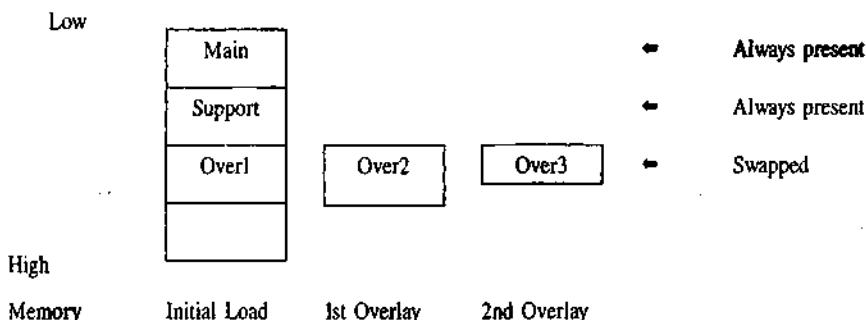
Include overlay module      : OVER2
Map onto linked module     : OVER1
Linking OVER2

Include overlay module      : OVER3
Map onto linked module     : OVER1
Linking OVER3

Include overlay module      : (RETURN)

```

Figure 10-1. Diagram of Linking Overlay Files



Note that the linker accepts input redirection. Complicated schemes for linking overlay files can be edited once in a file and are then fed to the linker as input. A file to accomplish the task in Figure 10-1 could look like this:

```

SUPPORT      OVER1      (modules to be included)
OVER2        OVER1      (first overlay)
OVER3        OVER1      (second overlay)

```

If the file had the name **LINKPROG** and you started the linker with

> L

```
Link main module : MAIN <LINKPROG
Output file      : MYPROG.COM
```

you would only have to answer two (Y/N) questions; entering all modules to be included or used as overlays would be unnecessary. Of course, input such as this:

>L

```
Link main module : MAIN MYPROG.COM <LINKPROG
```

or even

```
L MAIN MYPROG.COM <LINKPROG
```

is also legal.

A few restrictions must be made for overlay modules. First, an overlay module may not contain an initialization part (if so, the linker will issue an error message). Second, two overlay modules residing in the same segment should not import from each other. If the import concerns only constants, types, and variables, this is harmless. However, calling a procedure in another overlay module in the same segment will lead to trouble. In this case the linker issues only a warning message.

A program with overlays will be linked into two files with the extensions .COM and .OVR. The .OVR files contain all overlays needed for the program. They must be online when the program is started.

Every time a procedure in an overlay module is called, the module is brought into memory (provided it is not already there), overwriting whatever was in the segment previously. The same thing happens if a procedure returns to its caller and the caller has moved out of memory in the meantime. The calls to procedures in overlay modules do not need to follow a tree-like structure; they can proceed circularly. By specifying the mapping scheme, you can minimize overlay disk feeds and not bother with address computations, which is done by the linker.

The linker writes all overlay modules to disk as they are processed. Therefore, the length of all overlays is not restricted by memory size (the .OVR file can have a length of up to 256 Kbytes). However, the linker buffers modules that are not

overlays in main memory. Therefore, the size of the .COM file is restricted by the memory available to the linker.

Linking the Linker

If you experience an OUTOFMEMORY error during linking, you can increase the available memory by making the linker smaller. The linker consists of the file LINK.MCD, which is loaded on top of the Turbo Modula-2 system. The Turbo system contains several modules that are unnecessary for the linker's operation and that may be taking up valuable memory space (the editor, for example). These modules can be removed by linking the linker itself. To do so, simply give the following commands:

```
>L
```

```
Link main module : LINK  
Output file      : LINK.COM
```

and answer the question

```
Link all needed modules (Y/N)?
```

with . The result is a file LINK.COM. If the linker is now started from CP/M with the command A > LINK, the memory available for its operation is increased. If this is still not sufficient, you can use overlays for the linker, producing the files LINK.COM and LINK.OVR. You can do this by editing a file called LINKLINK or something similar that consists of the lines

```
TEXTS TERMINAL COMLINE LOADER
```

```
{ FILES LOADER  
CONVERT LOADER
```

Next, start the linker with this file as input, and answer »No« to whether to include all modules. The modules *Files* and *Convert*, which are needed by the linker, are both mapped as overlays onto the segment of the module *Loader*. They no longer use main memory for themselves. Therefore, the space needed by these modules is free to link even longer files.

Notice that the standard, but less efficient, »overlay scheme« that uses the procedure *Call* in *Loader* is also available.

Version Control

To prevent a program from running erroneous versions of separately compiled modules, a version number is given to each definition module. There are several ways to generate version numbers. Turbo Modula-2 determines a version number by calculating a checksum of the output produced when compiling a definition module. The number is kept in the symbol (.SYM) files as well as code (.MCD) files that are imported from each library module.

This method of generating numbers is handy since recompiling an unchanged module does not produce a new version number and requires unnecessary recompilation of dependent modules. With this process, you can even add comments to the definition module and then recompile it, which will have no effect on the version number generated.

Version control is the act of checking version numbers for consistency. There are three distinct points where version control is performed: compile time, link time, and runtime.

The compiler checks .SYM files for version numbers. Version conflicts occur when a library module's definition and implementation modules both import from another library module at different times (that is, different definitions, or version numbers, of the imported module are used by the definition and implementation modules). This situation is resolved by recompiling the module that was first compiled of the two importing modules.

If a library definition and implementation module both import from the same library module, then the compiler checks to make sure the definition (version number) has not changed between compiling the definition module and compiling the implementation module.

The linker and loader check .MCD files for version numbers. Version conflicts occur when trying to link or load a module that imports from a library module that has had its definition and implementation modules changed (and recompiled). The importing module must be recompiled to resolve the conflict.

Conflicts can also occur when a module imports from two library modules that in turn import from another library module at different times (that is, the lowest level module being imported from has changed its definition). To resolve this conflict, you must recompile the first compiled of the two library implementation modules.

The linker compares the version numbers of all the import modules. Thus, if more than one module imports a particular library module, all the version numbers for that module must match. In that case, the version numbers are found in the importing .MCD files.

Once the loader does dynamic linking, it can also find the same kind of version conflicts the linker finds (in case linking was not done). The loader can also detect version conflicts when loading overlays, which occurs when an overlay module imports from its caller and the caller's definition subsequently changes. This conflict is resolved by recompiling the overlay module.

Utilities

The following utility programs will add to your collection of programming tools. The first utility converts preexisting object files to a form usable by the Turbo Modula-2 system. The second tool allows you to evaluate the performance of your programs.

Linking Microsoft Relocatable Files

In some applications, it is desirable to use existing libraries with Modula-2. On Z80-CP/M systems, these libraries exist mostly in Microsoft relocatable format. Object files in this format are generated by such translators as the Microsoft 6800 assembler M80, the Microsoft FORTRAN compiler F80, Pascal MT/+, and others.

Turbo Modula-2 offers a utility to link programs in the Microsoft .REL format with Modula-2 programs. Only library procedures can be called from Modula-2, but constants and variables cannot be accessed. Nor can, say, a FORTRAN subroutine call a Modula-2 procedure.

Linkage is accomplished by a special program named REL. It takes as input a compiled definition module (with the extension .SYM) and one or more

relocatable files (with the default extension .REL). Its output is a normal Modula-2 .MCD file.

In the Modula-2 definition module, all procedures to be called from Modula-2 must be declared with their parameter lists. The names of the procedures must appear as entry points in the relocatable files. Parameters are normally passed on the stack. The definition module could look like this:

```
DEFINITION MODULE Z80Stuff;

FROM SYSTEM IMPORT ADDRESS, BYTE;

PROCEDURE MoveL(source,dest : ADDRESS; l : CARDINAL);

<other procedure declarations>

END Z80Stuff.
```

The input file (called, for example, Z80STUFF.MAC) to the M80 macro assembler consists of the following:

```
.Z80
PUBLIC MOVEL      <more identifiers>
MOVEL:POP        HL;    return address
              POP    BC;    length
              POP    DE;    source
              EX     (SP),HL; exchange return address and destination
              EX     DE,HL;  exchange source and destination
              LD     A,B
              OR     A,C
              RET    Z
              LDIR
              RET
<more code>
```

You can then translate these files by invoking the Turbo Modula-2 compiler at M80. Next, REL is executed by entering REL after the run-file prompt; for example:

Run MCD-file: REL Z80STUFF Z80 STUFF

The first parameter given is the name of the definition module. No extension may be given; the extension .SYM is always used. The second parameter given is the name of the relocatable file. If no extension is given, .REL is assumed.

REL produces an M-code file with the name of the definition module and the extension .MCD. This file may be used normally by a program, for example:

```
MODULE demo;

FROM Z80Stuff IMPORT MoveL;

BEGIN
  s := ...; d := ...; l := ...;
  MoveL(s,d,l);
END demo.
```

The Microsoft FORTRAN convention for passing parameters, however, is different from the scheme explained previously. The difference between these two schemes when linking (for example, FORTRAN routines or libraries) is that parameters are not passed on the stack but in registers. REL offers the command line switch »F« (/FORTRAN) to accommodate this. The FORTRAN switch may be given anywhere on the command line--its effect is global. The FORTRAN parameter-passing scheme is used for all REL files on the command line.

REL also has a search switch (»S«) that permits the user to search libraries and link only those modules that are actually needed. Unlike the FORTRAN switch, the search switch has only a local effect. It is given after a file name and causes only that file to be searched.

Here is a FORTRAN example:

```
FUNCTION IDIV(I,J)
IDIV = I/J
RETURN
END
```

And the Modula-2 definition module might look like this:

```
DEFINITION MODULE divide;

PROCEDURE idiv( i, j : INTEGER ) : INTEGER;

END divide;
```

After compiling both programs, link them by entering the following:

Run MCD-file: REL DIVIDE IDIV /F FORLIB /S

(We call our FORTRAN source file IDIV.FOR.) Once again, you can use function procedure *IDIV* any way that it's allowed in Modula-2.

Warnings: Parameters of structured types are always passed as VAR parameters to FORTRAN, even if a value parameter is specified in the definition module (such parameters are not copied).

FORTRAN routines are not reentrant, nor is the scheme used by REL to pass parameters. Since FORTRAN routines cannot currently call Modula-2 procedures, there are no problems with recursive procedures. Programs using interrupts, however, may have trouble.

Profile

Profile is a utility included with Turbo Modula-2 that allows you to estimate the time spent in various procedures of your program.

If the profiler is invoked with the command

> R

Run MCD-file:PROFILE

the main menu prompt > will reappear on the screen as if nothing happened. The profiler actually loads a second version of the shell. You can then run the program to be profiled by giving another Run command. When this program has executed, the profiler outputs a detailed list of all procedures and modules and the percentage of time spent in them.

The profiler can only monitor M-code procedures; native code procedures will not appear in its output. In fact, the profiler simply counts M-code operations.

Since all M-code instructions do not execute in the same amount of time, the measurement is slightly inaccurate; but it should be sufficient to find out which modules and procedures are time-critical.

With this information you can now proceed to compile modules into native code or to write extremely time-critical procedures in assembly.

Since output of the profiler is quite long, output redirection is normally used. For example:

Run MCD-file: PROFILE > MYPROG.PRO.

(The following example monitors MYPROG.MCD performance using PROFILE and then routes the report to file MYPROG.PRO.

>R

Run MCD-file: PROFILE > MYPROG.PRO

>R

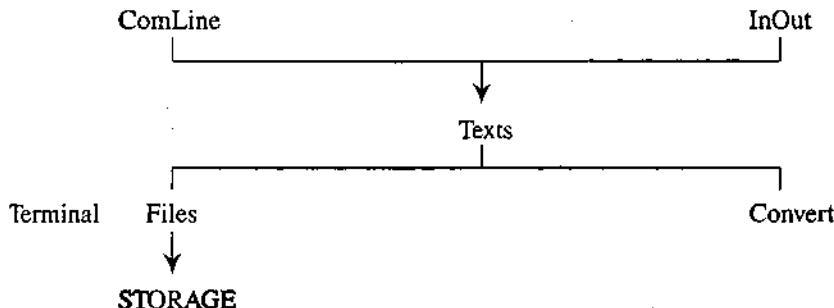
Run MCD-file: MYPROG.PRO

Chapter 11

The Standard Library

Some constructs needed for programming are not covered by the language Modula-2. Instead, these operations are performed by a set of predefined modules that cover input and output, mathematical functions, string handling, and storage and process management. A good part of the modules deal with input and output, and thus form a hierarchy (see Figure 11-1).

Figure 11-1. Library Module Hierarchy



The arrows in Figure 11-1 denote dependencies on other modules. In general, modules at the top of this drawing can be regarded as high-level modules, and those at the bottom as low-level modules. The module *Terminal* is an independent low-level module; that is, it does not depend on other modules and no module depends on it. The module *STORAGE* does not deal with input and output, but manages heap memory for control blocks and buffers.

The modules *MathLib*, *LongMath*, *Strings*, and *Doubles* comprise what is called a utility library. They provide facilities for manipulating objects.

The low-level modules *Loader*, *SYSTEM*, *STORAGE*, and *Processes* are system-dependent. *Loader* and *Processes* are higher level in that the programs that use them can become portable with minor changes. The modules *SYSTEM* and *STORAGE* constitute a special case: They enclose system-dependent primitives

needed for low-level programming. The compiler translates all operations imported from them directly into code. *SYSTEM* and *STORAGE* have neither a definition nor an implementation module (nor compiled counterparts), since technically they are not library modules. Because the compiler knows about them in advance, the use of the facilities from these two modules forces the programmer to explicitly specify (in an import list) the use of system-dependent, low-level constructs. Figure 11-2 displays the interdependencies of these modules.

Figure 11-2. Interdependencies of System-Dependent, Low-Level Modules



We can see that the Modula-2 library can be divided into three sections: input/output modules, utility modules, and system and low-level modules. In the next section we will briefly discuss the function of each library module within these three module types; a detailed description of each module will follow.

Overview of Input and Output Modules

Most programs involve input and output of data; for example, entering characters at the keyboard or writing program output to a disk file. Although these processes are different, Turbo Modula-2 treats them in a similar manner.

ComLine Module. Gives a program access to command-line parameters. The standard texts *input* and *output* can be redirected from the standard devices with this module.

Files Module. Procedures providing low-level access to disk files, including random access as well as sequential access. This module is used by the higher-level modules *InOut* and *Texts* to interface with the operating system.

InOut Module. Wirth's standard procedures for high-level input and output.

Terminal Module. Procedures to read the keyboard and write to the console with

highlighting and cursor control. Includes a procedure to clear the screen and one to check for characters in the keyboard buffer or console.

Texts Module. Procedures to read and write text (character streams) to and from external files and devices. Includes provisions for installable input/output device drivers.

Overview of Utility Modules

Utility modules provide string-handling procedures, mathematical functions, and conversion from numbers to strings and back again.

Convert Module. Contains procedures to convert strings into numbers and numbers into strings; handles CARDINAL, INTEGER, and REAL numbers.

Doubles Module. Contains procedures to convert strings into double-precision numbers and double-precision numbers into strings; handles LONGINT and LONGREAL numbers.

LongMath Module. Contains several mathematical functions for double-precision numbers, including the common transcendental functions and the natural power functions.

MathLib Module. Contains several mathematical functions for single-precision numbers, including the common transcendental functions and the natural power functions; this module also includes a random-number generator.

Strings Module. Contains procedures for string handling, including deletion, insertion, and copying of substrings.

Overview of System and Low-Level Modules

These modules provide an interface to the operating system, memory, and ports. They allow the programmer to develop system routines for memory management, multiprogramming, and other low-level constructs.

Processes Module. Enables the formulation of loosely coupled processes that are

implemented as coroutines. You can use this module to simulate concurrent programming.

STORAGE Module. Handles memory management for allocation and deallocation of dynamic variables; also keeps track of available memory.

SYSTEM Module. Contains system-dependent procedures that work with addresses, ports, operating systems and assembly language calls, and facilities for low-level process control.

Loader Module. Handles loading and execution of compiled object modules, enabling the splitting of large programs into smaller parts; the less commonly used parts are loaded from a disk as overlays. This allows the program to manage its own overlays rather than having it done by the runtime system, which is how the overlays produced by the linker are handled.

Details of the Module Library

Rather than defining specific input/output and system-specific statements in the language, these operations are provided in a hierarchy of modules. The modular structure has the effect (and advantage) of hiding computer-specific details, while also providing a means of including new machine-dependent routines when required. What follows is a full description of the services offered by the predefined modules.

Input and Output

Input and output remain one of the biggest problems in high-level language design. This is not surprising, since these operations have to deal with a large number of peripheral devices whose characteristics cannot be contained in a few simple abstractions. In every input/output system, simplicity and brevity of details are in conflict with efficient and flexible use of peripheral devices.

In Modula-2, this dilemma has been circumvented by not including any input/output primitives in the language at all. Instead, all communication with external devices is provided by a set of modules supplied with the system. And since these primitives are not part of the language (or the compiler), they can be freely supplemented according to the user's needs.

However, such a library of input/output functions has an intrinsic disadvantage: It leads to a large set of different primitives that must be remembered by the user. This is because Modula-2 does not allow procedure names to be overloaded (which is when more than one procedure has the same name, but is differentiated by the types of its parameters). Instead, there must be a separate routine for processing each type, even if each routine is essentially the same. For example, you must have one routine to write a character to the screen, another to write an INTEGER, another to write a REAL, and so on.

In addition, Modula-2 does not allow procedures to have a variable number of parameters. Thus, when you wish to write three strings to the screen, you must call *WriteString* three separate times rather than make one call that passes three parameters.

On the other hand, the advantage of having one call for one parameter for each type means you need not guess at the type of the parameter or search out its definition--it is made explicit directly in the code.

In our experience, it is often more difficult to remember the correct usage of the various input/output operations than to deal with the syntax of Modula-2. This is clearly an undesirable situation, particularly for the novice. In Turbo Modula-2, we opted for the best of both worlds: We have included all input/output modules postulated by the standard and allowed the use of general-purpose read and write statements as well, namely *READ*, *READLN*, *WRITE*, and *WRITELN*.

The form of these read and write statements is similar to that of procedures, but they accept varying numbers of parameters, where each may be a different type. This type of procedure is seen with other standard procedures such as *INC* and *DEC*, which accept any scalar type and an optional count parameter.

(*ID*, *READLN*, *WRITE*, and *WRITELN* are not declared in any standard module; instead, they must be thought of as extensions to the language Modula-2. The compiler translates these statements into corresponding procedure calls to the module *Texts*. This occurs transparently since you can supply your own drivers to the *Texts* module. You lose nothing in terms of efficiency or flexibility, but gain a lot in simplicity of input and output.

The input and output operations form a hierarchy: At its top are the general *READ* and *WRITE* statements, and at its bottom the underlying operating system. Between the extremes there is a sequence of modules incorporating primitives

with varying degrees of abstraction. Before we delve into the explanation of these modules, let's work out some of the characteristics of most input and output operations.

Input and output usually consists of the reading and writing of a sequence of data items. The process is often strictly sequential; for many peripheral devices, this is the only way to do input or output. A line printer, for example, prints one character after the other, and input from the keyboard consists of a sequence of characters as well. Even if some peripheral device allows random access to data (for example, to a disk drive or a video screen where it is possible to write characters at any position), the sequential operation is often the most convenient. A structure that consists of a sequence of data items is called a *stream*. (This is a purely abstract notion and is not part of the Modula-2 language.)

Streams

As you type at the keyboard, picture a stream of data flowing into the computer from the keyboard. The data output to the disk file is another stream disappearing from the program and being stored on the disk.

A stream, as discussed here, is a sequence of data that is either input from or output to an external medium. The external medium may be either a disk file or a logical device (the keyboard, video screen, or printer); both forms are treated in the same way.

In general, the predefined input/output procedures allow you to link either an input or an output stream to a medium, pass data down the stream, and then sever the connection upon completion.

The following lists the main characteristics of a stream:

- The number of elements of a stream, called the *length* of a stream, can vary.
- Only one element of a stream is visible at any one time. Reading starts at the beginning of a stream. After a read operation, the next element in the stream becomes visible.
- A stream can be modified only by appending elements at its end. Appending an element is called *writing*.

- A stream has a mode that can be either *read* or *write*. The mode is specified before any reading or writing is carried out and cannot be changed thereafter.

Notice that none of these characteristics specify the type of the elements of a stream. They can be of any type; however, it is useful to make a broad classification. We generally distinguish between legible and illegible input and output. A legible stream consists of elements of type CHAR, which is also called a *text*. Texts provide the means of communication between the user and the computer, where illegible streams store data for further processing by the computer.

There are several input/output modules that differ in both flexibility and type of stream handled. For example, the module *Terminal* deals only with the screen and keyboard text streams, while the module *Files* can handle common sequential I/O or random access files that are not streams at all.

A special stream type is the *text stream*, which contains only objects of type CHAR; for example, a stream of characters input from the keyboard or output to the screen or a printer. Text streams are a special case because they are formatted with EOT and EOL characters. (An EOT character marks the end of the text stream, while an EOL character terminates a line.)

Logical Devices

We have just discussed where data originates from and how a program may abstract it; where the data ends up is also relevant. In general, operating systems permit you to write to either a disk file or to *logical devices*. A logical device acts as a code word, specifying to the operating system which path leads to which device, whether it is a printer, a terminal, or another device. Which logical devices you have available will depend on your operating system and your hardware. For example, the CP/M operating system recognizes the following code names:

- {
 - CON: The video screen for output, the keyboard for input
 - LST: The printer (output only)
 - RDR: An additional input interface
 - PUN: An additional output interface

Note: Your configuration determines which devices can be accessed by RDR: a by PUN:.

The Texts Module

Texts provides the means to input and output a legible stream to and from peripheral devices. The module *Texts* uses the following characters to format streams:

- EOL (=36C, 30 decimal) denotes the end of a line.
- EOT (=32C, 26 decimal) denotes the end of the text.

Note that EOT is not defined as a constant in *Texts*.

Peripheral devices usually employ either a single carriage return (15C, decimal) or a carriage-return/line-feed (12C, 10 decimal) pair to denote line endings. The conversion from these control codes to the EOL character and reverse is done automatically by the module *Texts*.

At the center of this module is the type *TEXT*. It describes any legible, sequential input and output. This type is declared as a subrange:

```
TYPE TEXT: [1..16];
```

Variables of type *TEXT* serve as indices to various internal tables of this module. The tables themselves are not exported; they remain hidden to guarantee the integrity of the system.

Standard Text Streams

Texts exports three standard predefined texts that are declared as

```
input, output, console: TEXT;
```

The frequently used text streams are offered as predefined standard texts already linked to commonly used mediums. These standard texts do not require the usual linking and severing operations mentioned earlier. For example:

- Input* is the primary input stream usually assigned to the keyboard, but can be redirected to another medium, such as a card reader or other input device.

- Output* is the primary output stream to the video console, but can be redirected to another medium, such as a printer or a file.
- Console* is the output stream to the video console; it cannot be redirected. It is used to print error messages to the screen when the normal output stream is redirected elsewhere.

Texts Module Specification

```

DEFINITION MODULE Texts;
FROM Files IMPORT FILE;
TYPE TEXT = [1..16];

VAR input,output,console: TEXT;

PROCEDURE ReadChar      (t: TEXT;      VAR ch: CHAR);
PROCEDURE ReadString    (t: TEXT;      VAR s:  ARRAY OF CHAR);
PROCEDURE ReadInt       (t: TEXT;      VAR i:  INTEGER);
PROCEDURE ReadCard      (t: TEXT;      VAR c:  CARDINAL);
PROCEDURE ReadLong      (t: TEXT;      VAR l:  LONGINT);
PROCEDURE ReadReal      (t: TEXT;      VAR r:  REAL);
PROCEDURE ReadLn         (t: TEXT);

PROCEDURE WriteChar     (t: TEXT;      ch:      CHAR);
PROCEDURE WriteString   (t: TEXT;      s:       ARRAY OF CHAR);
PROCEDURE WriteInt      (t: TEXT;      i:       INTEGER; n: CARDINAL);
PROCEDURE WriteCard     (t: TEXT;      c, n:   CARDINAL);
PROCEDURE WriteLong     (t: TEXT;      l:       LONGINT; n: CARDINAL);
PROCEDURE WriteReal     (t: TEXT;      r:       REAL; n: CARDINAL;
                        digits: INTEGER);

PROCEDURE WriteLn        (t: TEXT);

PROCEDURE ReadLine       (t: TEXT;      VAR s:  ARRAY OF CHAR);
PROCEDURE ReadAgain     (t: TEXT);

PROCEDURE Done           (t: TEXT): BOOLEAN;
PROCEDURE EOLN           (t: TEXT): BOOLEAN;
PROCEDURE EOT            (t: TEXT): BOOLEAN;
PROCEDURE Col             (t: TEXT): CARDINAL;
PROCEDURE SetCol          (t: TEXT;      column: CARDINAL);
PROCEDURE TextFile        (t: TEXT): FILE;
PROCEDURE OpenText        (VAR t: TEXT; name: ARRAY OF CHAR):
                           BOOLEAN;
PROCEDURE CreateText     (VAR t: TEXT; name: ARRAY OF CHAR);
PROCEDURE CloseText       (VAR t: TEXT);

```

```
CONST EOL=36C;

TYPE TextDriver = PROCEDURE(TEXT, VAR CHAR);

PROCEDURE ConnectDriver(VAR t: TEXT; p: TextDriver);

PROCEDURE Init; (* used only by system *)

VAR haltOnControlC : BOOLEAN; (* TRUE by default *)

EXCEPTION TextNotOpen, TooManyTexts;

END Texts.
```

The following example shows the use of the *Texts* module. It reads ten from the keyboard and writes them to a disk file.

```
MODULE DataSave;
FROM Texts IMPORT TEXT, input, Done, CreateText, CloseText,
          ReadCard, WriteCard;
VAR
  DataFile: TEXT;
  Value,Count: CARDINAL;
BEGIN
  (* Create disk file and link text stream *)
  CreateText(DataFile, 'DATA.DTA');
  (* Read 10 values, save on disk *)
  FOR Count:= 1 TO 10 DO
    ReadCard(Input,Value);
    IF Done(Input) THEN
      WriteCard(DataFile,Value,16)
    END
  END; (* Close disk file *)
  CloseText(DataFile)
END DataSave.
```

Note that since *input* is a standard type *TEXT* declared in module *Texts*, it does not require opening or closing.

Stopping the Program During Input and Output

The following variable lets the programmer determine when a program can be interrupted while using the *Texts* module:

```
VAR haltOnControlC: BOOLEAN;
```

The variable *haltOnControlC* is exported by *Texts* and can be specified by the user. If *haltOnControlC* is TRUE (the default), then the program will halt on either of the following conditions:

- If you enter a **[CTRL]** **C** as the first character of an input line;
- If you enter a **[CTRL]** **S** to stop the output to the screen, followed by a **[CTRL]** **C**. (Note that *WriteLn* is the only output statement that checks for keyboard input; thus output to the screen can only be stopped if it is used.)

If *haltOnControlC* is set to FALSE by the programmer, then the user may only interrupt by program control.

Opening, Creating, and Closing a Text

Before data can be written or read, a connection must be established between a variable of type *TEXT* and some external medium, such as a disk file or a line printer. This is accomplished with the three procedures that follow.

OpenText (VAR *t*: TEXT; Name: ARRAY OF CHAR): BOOLEAN;

Associates an internal text value assigned to *t* with an existing external file identified by the string *Name*. The variable *t* can be used only for input. The parameter *Name* must be the name of a disk file or the code name for one of the operating system's logical input devices, such as CON: or RDR:. In the case of a disk file, a new internal file is created and connected to the existing external file using the procedure *Open* in module *Files*. When an input device is specified, the variable *t* is linked to the external device. In either case, *OpenText* returns TRUE if the text or device is opened successfully; FALSE if the file or device is not found.

CreateText(VAR *t*: TEXT; Name: ARRAY OF CHAR);

Establishes a new external text specified by *Name* and connects it to the text variable *t*. The text *t* can be used only for writing. Usually, the parameter *Name* identifies a new disk file, which is then created by the procedure *Create* in module *Files*. If the specified name denotes one of the logical output device names, such as CON: or LST:, then the text *t* is connected to that particular device.

After input or output has been completed on a text, the association with an external medium can be severed by use of the procedure *CloseText*. It is always recommended to close a text when processing is complete; otherwise, output files may not be updated and input files will not have control blocks disposed of properly.

CloseText(VAR *t*: TEXT);

Severs the connection between the given internal text *t* and its associated external text. If the external medium is a disk file, this file is closed using *Close* in module *Files*. When *t* is used for output, the EOT character is appended to the text before it is closed.

A special convention is used for the standard texts *input*, *output*, and *console*. These texts are always open and connected by default to the terminal (imagine them opened with the name CON:). Thus, when reading from the keyboard or writing to the video screen, *input*, *output*, or *console* do not need to be opened or created. When *OpenText* is applied to *input*, the standard input text is redirected to the external medium identified by the parameter *Name*. Likewise, *CreateText(output,name)* redirects standard output to the medium given by *name*. Closing *input* or *output* will establish the default connection again. Following are some examples.

OpenText(t,"CON:");

Establishes *t* as an input text with data from the keyboard.

OpenText(input,"RDR:");

Opens the standard text *input* with a different device. This time the standard text *input* is linked to the RDR: device instead of the keyboard; that is, *input* is *redirected*.

OpenText(InText,"B:INDATA.DAT");

Opens existing file INDATA.DAT on drive B for input. Subsequent input from text *InText* is taken from this external file.

CreateText(Printer,"LST:");

Connects the internal output text *Printer* with the external printer device.

CreateText(OutText,"OUTDATA.DAT");

Creates a new output file on the currently logged drive. Output statements specifying *OutText* write to this file.

CreateText(output,"PROTOCOL");

Redirects the standard output text. Standard output is no longer displayed on the screen; it is written to a file PROTOCOL that is created by this call.

CloseText(OutText);

Closes the output file B:OUTDATA.DAT, thereby establishing this file in the directory of drive B.

CloseText(input)

Returns the standard *input* from the RDR: device to the keyboard.

CloseText(output)

Closes file PROTOCOL and returns output to the video screen.

Renaming, Deleting, and Other File Operations

In addition to the previously described functions, you may need to rename, delete, or perform other operations on text files. This can be done via the module *Files* described on page #. All of the services offered in *Files* can be performed on *Texts* by using the function *TextFile*.

TextFile(VAR t: TEXT): FILE;

Returns the value of *FILE* variable associated with the *TEXT* *t*. If this text has not yet been opened, the value NIL is returned. Here are some examples.

Rename(TextFile(t),'NEWNAME.DAT')

Renames the file associated with the text *t* to 'NEWNAME.DAT,' and also closes it.

Size := FileSize (TextFile(t))

Puts the size of *t* in bytes in the variable *Size* of type LONGINT.

Flush(TextFile(t))

Flushes the *t* file buffer.

Reading and Writing

Once a text has been opened, data can be read from it. Reading generally involves two operations: (l) the transmission of one or several characters from the external

medium into the computer's memory, and (2) an interpretation of these characters. If, for example, a variable of type INTEGER is read, the decimal representation on the external device must be translated into the binary form used internally by the computer. The following set of procedures are applicable when inputting formatted data.

ReadChar(t: TEXT; VAR ch: CHAR);

Reads a single character from text *t* into *ch*.

ReadString(t: TEXT; VAR s: ARRAY OF CHAR);

Reads a string; that is, a sequence of characters not containing blanks or control characters. Any leading blanks are skipped. The string *s* receives as many legal characters as will fit into it. Input is terminated by a blank, an EOL character, or any other control code (a character with a decimal value less than 32). Note that strings cannot be separated by commas.

The next procedures can be used for numeric input.

ReadInt(t: TEXT; VAR i: INTEGER);

Reads a string and converts it to an integer. Again, leading blanks are ignored. Valid input is any sequence of digits 0 to 9, possibly preceded by an arithmetic sign (+ or -). The integer must be in the range *MIN(INTEGER)* to *MAX(INTEGER)*.

ReadCard(t: TEXT; VAR c: CARDINAL);

Reads a string and converts it to CARDINAL. Leading blanks are ignored and no sign is allowed. The CARDINAL must be in the range from 0 to *MAX(CARDINAL)*.

ReadReal (t: TEXT; VAR r: REAL);

Reads a string and converts it to a real number. Leading blanks are ignored. The string may contain a decimal point, a mantissa, and an exponent, but these are optional.

ReadLong(t: TEXT; VAR l: LONGINT);

Reads a string and converts it to a number of type **LONGINT**. Leading blanks are ignored.

Note: To read and write double-precision real numbers, use the module *Doubles* described later in this chapter.

Note that when *t* is the standard text *input*, all characters read from the keyboard are echoed on the screen. Line editing is possible until a **[RET]** is entered. If you wish to either input a character from the keyboard without echo or to enter a carriage return, use the procedure *ReadChar* provided in the module *Terminal*.

Done(t:TEXT):BOOLEAN;

Monitors correct execution of numeric input requests. When applied after a numerical read operation, *Done* returns a Boolean value that indicates whether the input is valid.

Thus, when a number is read, *Done* is FALSE if the input is syntactically illegal (for example, ABC) or too large to be converted (for example, 1000000) when a cardinal is requested.

Done is unaffected by all other operations. Note the following examples:

Declarations:

```
VAR ch: CHAR; i: INTEGER; c: CARDINAL; r: REAL;
    s: ARRAY [0..20]OF CHAR;
    t: TEXT;
```

(
Operation: ReadInt(input,i);
Done is TRUE: "0", "123", "-10".
Done is FALSE: "ABC", "1.0", <EOL>, "60000" (too large).
Valid Range: -32768 to +32767

Operation: ReadCard(input,c);
Done is TRUE: "0", "123", "60000",
Done is FALSE: "ABC", "1.0", "100000" (too large), "-1."
Valid Range: 0 to 65535

Operation: ReadReal(t,r);
 Done is TRUE: "0", "-3.", "3.14", "10.00E+10", "10E10".
 Done is FALSE: "E0", <EOL>, "10E40" (too big)
 Valid Range: -6.80564E38 to +6.80564E38

Operation: ReadLong(t,l);
 Done is TRUE: "0", "-10000000", "78"
 Done is FALSE: "ABC", "1.0", "10000000000" (too big)
 Valid Range: -2,147,483,624 to +2,147,483,623

ReadLine(t: TEXT; VAR Line: ARRAY OF CHAR);

Reads a whole line of text into the string variable *Line*. This is similar to *ReadString*, but no leading blanks are skipped and input is terminated only by an EOL or EOT character.

ReadLn(t: TEXT);

Advances the read position past the EOL character of the last line read. All characters on the current line, including the next EOL, are skipped. The next read position starts a new line. This procedure is useful to catch trailing blanks on a line or for skipping over unwanted input.

ReadAgain(t: TEXT);

Causes the last character read from text *t* to be returned upon the next read operation. For example:

```
Read(t,ch); ReadAgain(t);
IF (ch>="0") & (ch<="9") THEN ReadCard(t,x)
ELSE
  ReadString(t,s)
END ;
```

EOLN(t: TEXT): BOOLEAN;

Returns TRUE if the last character read was an EOL (36C) character.

EOT(*t*: TEXT): BOOLEAN;

Returns TRUE if the last character read was an EOT (32C) character.

Once a text has been created, data can be written to it. Again, writing involves two operations: (1) the translation of the internal representation of data in the computer's memory to one or more characters, and (2) the transmission of these characters to the external medium. Following are procedures to write data.

WriteChar(*t*: TEXT; *ch*: CHAR);

(Writes a single character to the text *t*.

WriteString(*t*: TEXT; *s*: ARRAY OF CHAR);

All characters in the given string *s* are written, starting from the first element (with index 0) up to the first null character (0C) or the end of *s*.

WriteInt(*t*: TEXT; *i*: INTEGER; FieldWidth: CARDINAL);

Write the integer *i* to text *t*. The given number appears right-justified in a field of *FieldWidth* characters. If *FieldWidth* is too small, more space is allocated.

WriteCard(*t*: TEXT; *c*,FieldWidth: CARDINAL);

Write the CARDINAL *c* to text *t*. The given number appears right-justified in a field of *FieldWidth* characters. If *FieldWidth* is too small, more space is allocated.

WriteReal(*t*: TEXT; *r*: REAL; FieldWidth: CARDINAL; Digits:INTEGER);

(Writes real number *r*, using at least *FieldWidth* characters; *r* is right-justified in the *FieldWidth* character field.

The parameter *Digits* controls the number of digits used in the mantissa of *r*. If it is positive, *r* is written in fixed-point format with *Digits* characters after the decimal point. If it is zero, the decimal point is omitted. A negative value of *Digits* indicates scientific notation; that is, the decimal point follows the most significant digit and an exponent is present. In this case, the number of digits in the mantissa is equal to *ABS(Digits)*.

WriteLong(t: TEXT; l: LONGINT; FieldWidth: CARDINAL);

Writes long integer *l*, right-justified in a *FieldWidth* character field. If more than *FieldWidth* columns are needed to print *l*, more space is allocated.

WriteLn(t: TEXT);

Terminates output line. *WriteLn(t)* is equivalent to *WriteChar(t,EOL)*. The EOL character is translated to a carriage-return/line-feed pair.

SetCol(t: TEXT; Col: CARDINAL);

Advances the write position of text *t* to column *Col* in the current output line. If *Col* is not greater than the current column number of *t*, nothing happens. Columns are numbered from 0; that is, the left-most column has number 0.

Col(t: TEXT): CARDINAL;

Returns the current column position of text *t*. After a *WriteLn*, *Col* returns 0; it increases by 1 with every character written.

The following are examples of the preceding procedures:

Operation	Output
<code>WriteInt(output,1025,5);</code>	" 1025 "
<code>WriteInt(t,1025,7);</code>	" 1025 "
<code>WriteReal(output,12.28,10,);</code>	" 12.3 "
<code>WriteReal(output,12.28,5,0);</code>	" 12 "
<code>WriteReal(output,12.18,12,-5);</code>	" 1.21800E+01 "
<code>WriteLong(output,-1000L*1000L,10);</code>	" -1000000 "
<code>SetCol(output,Col(output) DIV 10 + 1) * 10);</code>	Tabs to the next multiple of 10

The following example sums up input line by line. (This is recommended for very portable code; for an easier method, see the *READ* and *WRITE* statements that follow.)

```

MODULE SumUp;
FROM Texts IMPORT input, output, EOLN, EOT, ReadReal, Done,
              WriteChar, WriteReal, WriteLn, WriteString;

VAR
  item,sum: REAL;
BEGIN
  REPEAT
    WriteString(output, "enter real numbers > ");
    sum:=0.0;
    ( REPEAT
      ReadReal(input,item);
      IF Done(input) THEN sum:=sum+item END ;
      UNTIL EOLN(input) OR EOT(input);
      IF NOT EOT(input) THEN
        WriteString(output, "Sum = " );
        WriteReal(output, sum,12,-5 );
        WriteLn(output);
      END
    UNTIL EOT(input)
  END SumUp.

```

READ and WRITE Statements

The number of procedures in the preceding section may seem overwhelming; however, the *READ* and *WRITE* statements are much easier to use. The compiler translates them into appropriate calls to procedures in the module *Texts*, which causes an implicit import from that module. *READ*, *READLN*, *WRITE*, and *WRITELN* are predeclared identifiers, not reserved words. Note that *READ* and *WRITE* statements are an extension to Modula-2 and their use may not be portable to other Modula-2 implementations. Their full definition follows:

- Let t denote a variable of type *TEXT* and let v_1, v_2, \dots, v_n denote variables of a readable type; that is, of one of the types *INTEGER*, *CARDINAL*, *CHAR*, any subrange of them, *REAL*, *LONGINT*, or any *ARRAY* with *CHAR* elements.
- Let p, p_1, \dots, p_n denote parameters of one of the forms e_1 , $e_1 : e_2$, or $e_1 : e_2 : e_3$, where e_1 is an expression of a readable type, e_2 is a cardinal expression, and e_3 is an integer expression.

READ(t,v1,...,vn)

Causes v_1, \dots, v_n to be read from text t . The read items must be separated by blanks or line ends. The text variable t may be omitted.

READ(v1,...vn) Causes v_1, \dots, v_n to be read from the standard text *input*.

READLN(t,v1,...,vn)

Works exactly like *READ*, except that the remainder of the last line read is skipped.

WRITE(t,pl,...,pn)

Causes p_1, \dots, p_n to be written to text t . If the text variable is omitted, output goes to the standard text *output*, for example.

WRITE(pl,...,pn)

Characters and strings are written without preceding or subsequent blanks (except for blanks in the string itself).

Integers and cardinals are written right-justified in a field of six characters, which can be overridden with a field specifier.

Reals are displayed in scientific notation, right-justified in a 12-character field. However, this may be overridden with a field specifier.

The field width of numbers can be changed by appending a colon and the desired width to the expression. The representation of real numbers can be changed to fixed-point by appending a second colon and the desired width of the fractional part (also see points 6 and 7 that follow).

WRITELN(t,pl,...,pn)

Works exactly like *WRITE*, except that the line is terminated when all parameters are written. The argument list may be empty; if so, only the EOL character is written.

The translation process from the *READ* and *WRITE* statements to calls of procedures in module *Texts* is as follows:

1. *READ(v1,...,vn)* is translated to *READ(Input,v1,...,vn)*.
WRITE(p1,...,pn) is translated to *WRITE(Output,p1,...,pn)*.

The same equivalence holds for *READLN* and *WRITELN*.

2. *READLN(t,v1,...,vn)* is translated to *READ(t,v1,...,vn);*
ReadLn(t).

- (*"WRITELN(t,p1,...,pn)* is translated to *WRITE(t,p1,...,pn);*
WriteLn(t).

3. *READ(t,v1,...,vn)* is translated to
READ(t,v1);...;READ(t,vn).
WRITE(t,p1,...,pn) is translated to
WRITE(t,p1);...;WRITE(t,pn).

4. *READ(t,v)* is translated to

<i>ReadChar(t,v)</i>	if v is of type CHAR.
<i>ReadString(t,v)</i>	if v is of type ARRAY OF CHAR.
<i>ReadInt(t,v)</i>	if v is of type INTEGER.
<i>ReadCard(t,v)</i>	if v is of type CARDINAL.
<i>ReadReal(t,v)</i>	if v is of type REAL.
<i>ReadLong(t,v)</i>	if v is of type LONGINT.
<i>ReadDouble(t,v)</i>	if v is of type LONGREAL.

5. *WRITE(t,p)* is translated to

- (*WriteChar(t,p)* if p is of type CHAR.
WriteString(t,p) if p is of type ARRAY OF CHAR.
WriteInt(t,p,6) if p is of type INTEGER.
WriteCard(t,p,6) if p is of type CARDINAL.
WriteReal(t,p,12,-5) if p is of type REAL.
WriteLong(t,p,12) if p is of type LONGINT.
WriteDoubles(t,p,22,-14) if p is of type LONGREAL.

Note that the standard width for the representation of a number is 6 in the case

of integers and cardinals and 12 in the case of **reals**; **reals** are normally written in scientific notation with a 5-digit mantissa.

Also note that you cannot replace the procedure call *ReadLine* with a *READLN* statement. *READLN(line)* would only read the first string of a line, stopping at the first blank encountered, then skipping to the end of the line without reading anything. On the other hand, *ReadLine(Input, Line1)* reads in the full line from the start to the next EOL character.

6. If x is of type **INTEGER**, **CARDINAL**, or **LONGINT** and n is a cardinal expression, then *WRITE(t,x:n)* is translated to *WriteInt(t,x,n)*, *WriteCard(t,x,n)*, or *WriteLong(t,x,n)*, respectively; that is, the default width specification is overridden.
7. If x is of type **REAL** or **LONGREAL**, then *WRITE(t,x:n)* is translated to *WriteReal(t,x,n,-5)*. Again, the default width specification is overridden.

Furthermore, if i is an integer, *WRITE(t,x:n:i)* is translated to *WriteReal(t,x,n,i)*.

In this case, the form of notation (scientific or fixed-point) is indicated by the programmer. As in the previous explanation of procedure *WriteReal*, positive values of m indicate fixed-point notation with i digits behind the decimal point, whereas negative values stand for scientific notation.

The following example presents the same example from the earlier »Reading and Writing« section, except this time we make use of Turbo Modula-2's *READ* and *WRITE* extensions. The module becomes shorter and easier to read, even though it is translated into (nearly) the same code as the previous example. The *READ* and *WRITE* extensions do not in themselves shorten code generated by the compiler, rather they help to write shorter and clearer source programs.

```
MODULE SumUp;
FROM Texts IMPORT input, EOLN, EOT, Done;

VAR
  item,sum: REAL;
BEGIN
  REPEAT
    WRITE( "enter real numbers> " );
    READ( item );
    sum := sum + item;
  UNTIL item = 0.0;
  WRITELN( sum );
  Done();

```

```
sum:=0.0;  
REPEAT  
  READ(item);  
  IF Done(input) THEN sum:=sum+item END ;  
  UNTIL EOLN(input) OR EOT(input);  
  IF NOT EOT(input) THEN WRITELN("Sum =",sum:12:-5) END  
UNTIL EOT(input)  
END SumUp.
```

User-Defined I/O Drivers

These provide communication with disk files and the logical devices of the operating system. If you wish to communicate with some device not falling into these categories, you can install your own I/O drivers and procedure for input and output. Such a procedure must match the following procedure type:

```
TYPE  
  TextDriver = PROCEDURE(TEXT, VAR CHAR);
```

TextDriver is a template for the many possible I/O drivers. All of these procedures must accept a *TEXT* as their first parameter. Their second argument must be a *VAR* parameter of type *CHAR*.

If a text is used for output, the I/O driver sends its second argument; while a driver used for input deposits a character that was read in this parameter. Driver procedures must provide for the correct conversion of the EOL character into end-of-line character sequences used by the external device. The first parameter (of the subrange type *TEXT*) may serve as an index into a table that contains further information about the text on which the I/O takes place. In most situations it can be ignored.

To install a user-defined I/O driver, use the following procedure:

```
ConnectDriver(VAR t: TEXT; driver: TextDriver);
```

After a call to this procedure, I/O on text *t* will be handled via the procedure that was substituted for the *driver* argument. This procedure has the effect of opening text *t*; no further *OpenText* or *CreateText* need be given.

The following is an example of installing a driver for an IEEE interface. We have presented the sketch of a local module, exporting the two texts *InIEEE* and

InIEEE, communicating with an IEEE interface. It is assumed that line ends are denoted by single carriage-return characters.

```
IMPLEMENTATION MODULE IEEE;
FROM Texts IMPORT TEXT,EOL,ConnectDriver;
EXPORT InIEEE,OutIEEE;
CONST CR=15C;
VAR
  InIEEE,OutIEEE: TEXT;
PROCEDURE Put(t: TEXT; VAR ch: CHAR); (* Output driver *)
BEGIN
  IF ch = EOL THEN ch:=CR END ;
  (* code to write ch to the IEEE interface *)
END Put;

PROCEDURE Get(t: Text; VAR ch: CHAR); (* Input driver *)
BEGIN
  (* code to read ch from IEEE interface *)
  IF ch = CR THEN ch := EOL END ;
END Get;

BEGIN (* Initialization: Install Get and Put *)
  ConnectDriver(inIEEE, Get);
  ConnectDriver(OutIEEE ,Put)
END IEEE;
```

The texts *InIEEE* and *OutIEEE* may now be used to input or output characters, strings, or numbers. Statements like the following are now possible:

```
READ(InIEEE,x,y,s);
WRITELN(OutIEEE,"Size of Sample: ",SampleSize);
```

The InOut Module

InOut serves about the same purpose as the module *Texts*: It provides communication to external mediums without any assumptions about the nature of peripheral devices. The module *InOut* is postulated to be present in every Modula-2 system. It is included in the Turbo Modula-2 system for the sake of compatibility, even if abstract input and output is generally accomplished more conveniently by using the *READ* and *WRITE* extensions to Turbo Modula-2 in conjunction with the module *Texts*.

The operations provided by *InOut* can be thought of as a subset of those provided by *Texts*. *InOut* is more restrictive, however. Only operations on the standard text streams *input* and *output* are allowed. These standard streams are hidden by *InOut*, but can be accessed directly from *Texts*. As a consequence, procedures provided by *InOut* do not require a parameter to indicate which text stream is used. Read operations always use *input*, while write operations work on *output*. Correct execution of input procedures is monitored in the global variable *Done*, and the global variable *termCH* always contains the last character read.

In addition to the procedures provided by *Texts*, *InOut* exports two more output procedures.

WriteOct (*c,n*: CARDINAL);

Writes cardinal *c* in octal representation (range 0 to 177777 octal).

WriteHex (*c,n*: CARDINAL);

Writes cardinal *c* in hexadecimal representation (range 0 to FFFF hex).

Our implementation of *InOut* reads and writes real numbers as well as CARDINALs and INTEGERs. The user should keep in mind that other implementations might choose to separate these tasks. Procedures to read and write REALs would most likely be imported from the modules *RealInOut* or *RealIO*.

Note that when using the procedures *OpenInput* or *OpenOutput*, no file name is provided by the program. Instead, the name is requested from the user at the terminal, which limits the usefulness of these procedures.

InOut Module Specification

```
DEFINITION MODULE InOut;
CONST EOL=36C;
VAR Done: BOOLEAN;
    termCH: CHAR;

PROCEDURE OpenInput (defext: ARRAY OF CHAR);
(* Requests a file name at the terminal and connects input with that
   file. The file name may indicate one of the logical devices of
   CP/M. If the file name does not contain an extension itself, the
   default extension defext is appended.
*)

PROCEDURE OpenOutput (defext: ARRAY OF CHAR);
(* Same as above but for output *)

PROCEDURE CloseInput;
(* Closes input file, returns input to terminal. *)

PROCEDURE CloseOutput;
(* Closes output file, returns output to terminal. *)

PROCEDURE Read (VAR ch: CHAR);
(* Done := Not past end of input *)

PROCEDURE ReadString (VAR ss: ARRAY OF CHAR);
(* Reads string as in Texts. Done := the returned string is not empty
*)

PROCEDURE ReadInt (VAR x: INTEGER);
(* Reads integer as in Texts. Done := integer was read *)

PROCEDURE ReadCard (VAR x: CARDINAL);
(* Done := cardinal was read *)

PROCEDURE Write(ch: CHAR);
PROCEDURE WriteLn;
PROCEDURE WriteString(s: ARRAY OF CHAR);
PROCEDURE WriteInt(x: INTEGER; n: CARDINAL);
```

```
PROCEDURE WriteCard(x,n: CARDINAL);
PROCEDURE WriteHex(x,n: CARDINAL);
PROCEDURE WriteOct(x,n: CARDINAL);

PROCEDURE ReadReal (VAR x: REAL);
PROCEDURE WriteReal(x: REAL; n,digits: CARDINAL);
END InOut.
```

The Files Module

The next lower level of the module hierarchy connecting *InOut* and *Texts* with the computer's disk-operating system embodies the concept of a file. There are two ways to look at a file. The first is as an external object: *A file describes a collection of data stored on a magnetic disk.* Often the data is of the same type. The second is as an internal object: *A file represents an internal data structure that describes the external data and how it is accessed.* The two notions should not be confused. The type *FILE* as exported by the module *Files* describes access to data items, not the data itself. In some operating systems this is referred to as a *file control block*.

Files offers the possibility to access data collected in a disk file. The elements of a file may be of any type and can be legible or illegible. No interpretation is performed; elements in a file are represented in computer internal binary form. As an example, writing the integer 279 to a *TEXT* consists of the output of the three numerals 2, 7 and 9, or in binary form:

```
| 00110010 | 00110111 | 00111001 |
```

When writing the same number as an integer to a file, its internal representation is written without prior conversion into ASCII digits, like so:

```
| 00000001 | 00011001 |
```

Physically, a file is a sequence of bytes on disk that can be read or written. The read or write position can be set to any byte of the file, thus allowing data to be accessed randomly.

It is often convenient to think of files at a higher level of abstraction. At this level, files are ordered collections of variables that are usually all of the same type. The type of the file elements is arbitrary--it could be CHAR, a record, or any other type.

Errors During File Handling

Files signals error conditions by raising one of five exceptions. If such an excep-

tion is raised, the executing program is stopped and a message describing the cause of the error is displayed on the screen. Programs may recover from such error conditions if the programmer provides an exception handler for the exception raised (refer to Chapter 9). *Files* defines the following exceptions:

StatusError Indicates request of an illegal file-processing operation, such as reading from an unopened file or opening a file twice.

Files Module Specification

```

DEFINITION MODULE Files;
FROM SYSTEM IMPORT BYTE, WORD, ADDRESS;
TYPE FILE;

PROCEDURE Open          (VAR f: FILE; name: ARRAY OF CHAR):
                        BOOLEAN;
PROCEDURE Create        (VAR f: FILE; name: ARRAY OF CHAR);
PROCEDURE Close         (VAR f: FILE);
PROCEDURE Delete        (VAR f: FILE);
PROCEDURE Rename        (VAR f: FILE; name: ARRAY OF CHAR);

PROCEDURE GetName       (f: FILE; VAR name: ARRAY OF CHAR);

PROCEDURE FileSize      (f: FILE): LONGINT;
PROCEDURE EOF            (f: FILE): BOOLEAN;

PROCEDURE ReadByte       (f: FILE; VAR ch: BYTE);
PROCEDURE ReadWord       (f: FILE; VAR w: WORD);
PROCEDURE ReadRec        (f: FILE; VAR rec: ARRAY OF WORD);
PROCEDURE ReadBytes      (f: FILE; buf: ADDRESS; nbytes: CARDINAL):
                        CARDINAL;

PROCEDURE WriteByte      (f: FILE; ch: BYTE);
PROCEDURE WriteWord      (f: FILE; w: WORD);
PROCEDURE WriteRec        (f: FILE; VAR rec: ARRAY OF WORD);
PROCEDURE WriteBytes     (f: FILE; buf: ADDRESS; nbytes: CARDINAL);

PROCEDURE Flush           (f: FILE);
(* Flushes the file's internal buffer to disk. Is used to detect
 DiskFulls at once *)

PROCEDURE NextPos        (f: FILE): LONGINT;
PROCEDURE SetPos          (f: FILE; pos: LONGINT);

PROCEDURE NoTrailer      (f: FILE);

PROCEDURE ResetSys        ( );

```

```
EXCEPTION EndError, StatusError, UseError, DeviceError, DiskFull;  
END Files.
```

EndError	Indicates attempt to read past the end of a file.
UseError	Indicates file creation, renaming, or deletion on a write-protected disk, like the CP/M-message BDOS ERROR ON (drivename): R/O, which is suppressed in this case. Note that changed disks are always write-protected before a system reset. Also note that CP/M will not recognize a physically write-protected disk as R/O when reset; thus, <i>UseError</i> is not reliable for physically write-protected disks.
DeviceError	Indicates data is not readable, presumably because of a bad disk sector. This exception is also raised if a <i>Close</i> operation is unsuccessful.
DiskFull	Indicates that data cannot be written due to a full disk, or files or extents cannot be created due to a disk-directory overflow.

Operations on Entire Files

Before any processing can be carried out, the internal file must be connected with an external (disk) file. After the connection is made, the file is then said to be open. New files are opened with *Create* and existing files are opened with *Open*:

```
Open  (VAR f: FILE; name: ARRAY OF CHAR): BOOLEAN;
Create (VAR f: FILE; name: ARRAY OF CHAR);
```

These procedures connect an internal file *f* with an external file. *Open* searches the directory for an existing file with the given name, returning FALSE if no such file is found. *Create* always creates a new empty file; if a file with the same name already exists, it is deleted.

The following procedure performs a reset of the disk system. It is useful to prevent *UseErrors* caused by changed disks.

ResetSys;

After a call to *ResetSys*, it is possible to write to drives where disks have been swapped, but all open output files will be lost.

After processing has been completed on a file, the connection may be severed by the *Close* procedure.

Close(VAR f: File);

Disconnects *f* with its associated external file. If output has been sent to *f*, closing is mandatory; otherwise, all data of *f* will be lost. Since closing a file recovers the memory used for control blocks and data buffering, it is advisable to close a file in any case, even if it has only been used for input.

The following procedures are used to delete or rename a file:

Delete (VAR f: FILE);
Rename (VAR f: FILE; name: ARRAY OF CHAR);

Both procedures assume that *f* is open. *Delete* deletes the directory entry that corresponds to *f*, while *Rename* renames the entry to its given name. Both procedures have the side effect of closing the internal file *f*.

The external name of file *f* is returned by the following:

GetName(f: FILE; VAR name: ARRAY OF CHAR);

The name is returned in a standard format: The first character is the drive code, which can assume the values *A* through *P*. The second and third characters are the user area of the file. The fourth character is always a colon. The file name is next, consisting of up to 8 characters, with its end marked by a period. If the file name contains an extension, it is appended. To hold a returned file name, the *name* parameter should be declared with at least 16 characters.

For example, assuming the procedure calls

Open(f, "B4:INDATA.DAT");
Create(out, "OUTDATA");

then the results will be *GetName(f, fname)*, which will yield »B04:IN-
DATA.DAT,« and *GetName(out, outname)*, which will yield »A00:OUTDATA,«

The following procedure returns the size of an external file:

FileSize (f: FILE): LONGINT;

This returns the exact number of bytes in the file. Since files can be larger than

65535, the maximal CARDINAL value, the result is of type LONGINT. (since its maximum value is in the trillions, it should be adequate for some time)

Unfortunately, CP/M has no means to determine the exact size of a file in bytes. To correct this restriction, Turbo Modula-2 uses a special convention: The byte of the last CP/M record (128 bytes) in a file indicates the number of defined bytes in the record. This last byte has an offset value of 128; that is, if it is no bytes in the record are defined. A value of 255 indicates 127 defined bytes. A value below 128 means that the last record is completely filled. Using this byte and the BDOS function that returns the number of 128-byte records in a file, the file's exact size can be computed.

This scheme imposes almost no restriction on the file formats that can be represented. In particular, reading ASCII files generated by WordStar or other programs is possible without encountering problems. If you run into problems transferring files between Turbo Modula-2 programs and other software, you can use this procedure:

```
NoTrailer (f: FILE);
```

Note that the file *f* must be already open, thus allowing *NoTrailer* to cause the last byte of the last CP/M record of *f* to be interpreted as a data byte instead of a length byte. Consequently, a subsequent call of *FileSize(f)* will always return a multiple of 128. If the length of *f* is changed by write operations, no length will be appended to *f* when it is closed.

File Processing

We can think of a file as a sequence of bytes stored on disk. The number of bytes in a file is called the *length* of the file. A file's elements are all of the same type and each can consist of several bytes. However, only one element can be accessed at a time.

The position number of the accessible element is called the *current position* of the file. The current position can point to any element, or it may denote the end of the file; that is, it can range from 0 up to the *FileSize(f)*. Initially, if no positioning has been carried out, the current position will be 0. Reading advances the current position past the read data. Likewise, writing advances the current position past the modified data.

Once a file is open (an *Open* or *Create* operation applied to it), reading

writing becomes possible. The following operations are available for file processing and can be carried out only on open files. If one of these operations is applied to a file that is not open, the exception *StatusError* is raised. The file operations that follow read data from the given file *f* and differ in the kind of data that is read:

```
ReadByte  (f: FILE; VAR ch: BYTE);  
ReadWord  (f: FILE; VAR w: WORD);  
ReadRec   (f: FILE; VAR rec: ARRAY OF WORD);
```

ReadByte reads a single byte and assigns it to the parameter *ch*. *BYTE* is a special type, imported from the pseudomodule *SYSTEM*. It matches every variable occupying 1 byte of memory. This includes variables of the type *CHAR*, *BOOLEAN*, subrange types with bounds in the range 0 to 255, and enumeration types with, at most, 256 elements. If the file elements are one of these types, *ReadByte* should be used for reading.

ReadWord reads 2 bytes (a word) and assigns them to the parameter *w*. The type *WORD* is also imported from the pseudomodule *SYSTEM*. It matches any variable occupying 2 bytes of memory. Examples are variables of type *INTEGER*, *CARDINAL*, *BITSET*, sets, and all pointers. If the file elements are one of these types, *ReadWord* should be used for reading.

ReadRec reads everything not covered by *ReadByte* and *ReadWord*. The type *ARRAY OF WORD* matches variables of any type. This works since the storage of every element is rounded up to a word offset. As many words as are needed to fill out the given parameter are read. Employing this procedure for reading is appropriate for file elements occupying more than one word; in particular, for records, arrays, long integers, and reals.

In all three procedures, the current position is advanced by as many bytes as are read. A subsequent read operation will then access the next element in the file. If the current read position is higher than the end position, the exception *EndError* is raised.

The following file operations write data to the given file *f* and differ in the kind of data that is written:

```
WriteByte (f: FILE; ch: BYTE);
WriteWord (f: FILE; w: WORD);
WriteRec (f: FILE; VAR rec: ARRAY OF WORD);
```

These procedures give the value specified in the second parameter to the element at the current position in the given file *f*. If the current position is not higher than the end position, the corresponding file element is modified. If the current position points to the end of the file, data is appended and the length of the file is incremented.

While *WriteByte* is appropriate for file elements occupying 1 byte, *WriteWord* is used for word-sized elements and *WriteRec* applies to elements of a larger size (for example, records, arrays, or reals).

If writing is impossible, the exception *DiskFull* is raised. There are also two lower level procedures for reading and writing data.

ReadBytes (f: FILE; buf: ADDRESS; nbytes: CARDINAL): CARDINAL;

Reads as many bytes as are specified in the *nbytes* parameter into consecutive addresses, starting at the address denoted by *buf*. If less than *nbytes* bytes remain in the file *f*, only the number of remaining bytes are read. Hence, a call to this procedure will never cause the exception *EndError* to be raised. The number of bytes read is returned as a function result.

WriteBytes (f: FILE; buf: ADDRESS; nbytes: CARDINAL);

Writes as many bytes as are specified by the given *nbytes*. Data is written from consecutive addresses, starting at the address *buf*.

EOF(f: FILE): BOOLEAN;

Returns TRUE if the end of file *f* is reached; returns FALSE if there are still bytes in *f* to be read.

Note that this procedure provides a limited form of *lookahead*. You know that the end of a file has been reached before undefined data has been read. This works even if the file is empty, a fact that distinguishes a *FILE* from a *TEXT*. With *TEXT* you have to read the EOT character to know that the *TEXT* has been completely read. The appropriate form to process a *FILE* is a WHILE loop:

```
IF Open( f, fname ) THEN
  WHILE NOT EOF(f) DO
    ReadRec(f, data);
    Process(data);
  END;
  Close(f);
END
```

SetPos(f: FILE; pos: LONGINT);

Allows data to be accessed randomly. Sets the current position of the file *f* to *pos*. This parameter denotes the byte position where subsequent file processing will take place. The next read operation will then access the element at position *pos*. The next write operation will modify the element at this position. The parameter *pos* can point to any element in the file or it can point to the end of the file; its legal range is 0 up to and including *FileSize(f)*. If *pos* is larger than *FileSize(f)*, the exception *EndError* is raised.

If all file elements are of the same type *T*, and you know the ordinal number *n* of the selected element in the sequence (counting again from 0), you can compute its byte position as

BytePosition := n * SIZE(*T*)

Thus, a call like *SetPos(f, LONG(n)*LONG(SIZE(T)))* would set the current file position correctly. Note that *SetPos* requires a parameter of type LONGINT. Since type coercion is not performed in Modula-2, the construct *SetPos(f,0)* is illegal; use *SetPos(f, LONG(0))* or *SetPos(f, 0L)* instead.

The current file position can be examined by calling

NextPos (f: FILE): LONGINT;

which will return the current position of the given file *f*.

To demonstrate how files work we will present two examples. The first example shows a series of operations on a file (represented by boxes) and how these procedures affect the file and the file pointer (the up arrow). The declarations VAR f: FILE; c: CHAR; are used in the procedure calls in Example 1.

Example 1

1. Open(f, "INDATA");

```
| 01001100 | 01100101 | 01101100 | 01101100 |
```

 ^

Steps 2 through 4 show how to modify a file element: First read it in, then set the current position back by one and write the modified data.

2. ReadByte(f,c);

```
| 01001100 | 01100101 | 01101100 | 01101100 |
```

 ^

c now has the value 01001100 binary (>L<).

3. SetPos(f,NextPos(f)-1L);

```
| 01001100 | 01100101 | 01101100 | 01101100 |
```

 ^

4. INC(c); WriteByte(f,c);

```
| 01001101 | 01100101 | 01101100 | 01101100 |
```

 ^

Steps 5 and 6 show how to append data at the end of a random access file: Set the position to the length of the file and write the data to be appended.

5. `SetPos(f,Size(f));`

| 01001101 | 01100101 | 01101100 | 01101100 |

^

6. `(writeByte(f,"y");`

| 01001101 | 01100101 | 01101100 | 01101100 | 01111001 |

^

Example 2 is a practical implementation of a file copy utility that shows the typical usage of the *Files* module. Its operation is similar to the *Filecopy* command in the Turbo Modula-2 shell. Data is first written into a temporary file with the extension ---. If the transfer succeeds, the temporary file is renamed to whatever the user indicates. This scheme avoids an accidental loss of data.

Among others, two procedures of the module *FileExtensions* are imported, *StripExt* and *AppendExt*. (This module is not in the standard library. It is used as a programming example in the section on the module *Strings*.) In short, *StripExt* strips the extension from a file name and *AppendExt* appends an extension. The procedure *PromptFor*, imported from the standard module *Comline*, reads .ext from the command line or, if none is found, prompts the user for input.

Example 2

```
MODULE FileCopy;
FROM SYSTEM IMPORT ADR;
FROM FileExtensions IMPORT StripExt,AppendExt;
(* See the discussion of the Strings Module for a listing of the
   FileExtensions routines.
*)
```

```

FROM ComLine IMPORT PromptFor;
FROM Files IMPORT FILE,Open,Create,Close,Rename,Delete,ReadBytes,
              WriteBytes,DiskFull;

CONST BufferSize = 20000;
VAR
  inFile,outFile      : FILE ;
  inName,outName,tempName : ARRAY [0..20] OF CHAR;
  buffer             : ARRAY [1..BufferSize] OF CHAR;
  fetched            : CARDINAL;

BEGIN
  PromptFor("Copy file: ", inName);
  IF inName[0] # OC THEN (* no empty string entered *)
    PromptFor("Copy to : ", outName);
    IF outName[0] # OC THEN
      IF Open(inFile, inName) THEN
        tempName:=outName;
        StripExt(tempName);
        AppendExt(tempName,"$$$");
        Create(outFile, tempName);
        REPEAT
          fetched := ReadBytes(inFile,ADR(buffer),BufferSize);
          WriteBytes(outFile, ADR(buffer), fetched);
        UNTIL fetched < BufferSize;
        Rename(outFile, outName);
        Close(inFile);
      ELSE WRITELN(inName, " not found.") END
    END
  END
EXCEPTION
  DiskFull: WRITELN("DISK FULL"); Delete(outFile)
END FileCopy.

```

Files with Elements of Mixed Types

Up to now, we have only encountered files with elements all of the same type. However, sometimes it is convenient to include elements of varying types in one file; for example, a file can have a header as its first element, containing various information about a file's data. The actual data would follow the header in a different format. You can process such a file the same way you would a file with

uniform elements. The procedures *ReadByte*, *ReadWord*, *ReadRec*, *ReadBytes*, *WriteByte*, *WriteWord*, *WriteRec*, and *WriteBytes* are functional. Of course, the procedure chosen must match the type of accessed element. For example, *ReadByte* would be used to read a character, while *ReadRec* would serve to input a record.

The Terminal Module

The module *Terminal* provides communication over the operator's terminal (the keyboard and the video screen), without the abstraction of a text stream. The following procedures show how to implement basic input/output routines.

ReadChar(VAR ch: CHAR);

Reads a single character from the keyboard. The character is not echoed on the screen, which distinguishes the module *Terminal* from *Texts* or *InOut*. Line ends are denoted by carriage returns (15C); no conversion to the EOL character is performed.

BusyRead(VAR ch: CHAR);

Tests to see if a character has been typed. A typed character is returned in the given *ch*. If no key is pressed, *ch* is set to 0C.

ReadAgain;

Causes the last character read to be returned again upon the next call of *ReadChar* or any read statement from *Texts* or *Terminal*.

ReadLine(VAR s: ARRAY OF CHAR);

Reads a line (terminated by a carriage return) into a string.

WriteChar(ch: CHAR);

Writes a single character; no conversion from EOL to carriage-return/line-feed pairs is performed.

WriteLn;

Terminates a line.

WriteString(s: ARRAY OF CHAR);

Writes a string to the terminal.

Terminal Module Specification

```
DEFINITION MODULE Terminal;

PROCEDURE ReadChar (VAR ch: CHAR);
PROCEDURE BusyRead (VAR ch: CHAR);
PROCEDURE ReadAgain;
PROCEDURE ReadLine (VAR CHAR);      s:          ARRAY           OF

PROCEDURE WriteChar (ch: CHAR);
PROCEDURE WriteLn;
PROCEDURE WriteString (s: ARRAY OF CHAR);

TYPE
  SpecialOps (clearEol, insertDelete, highlightnormal);
  OpSet      = SET OF SpecialOps;

VAR available      : OpSet;
  numRows, numCols : CARDINAL;

PROCEDURE ClearScreen;
PROCEDURE GotoXY(x,y: CARDINAL);

PROCEDURE ClearToEOL;
PROCEDURE InsertLine;
PROCEDURE DeleteLine;

PROCEDURE Highlight;
PROCEDURE Normal;

PROCEDURE InitScreen;
PROCEDURE ExitScreen;

END Terminal.
```

Terminals may require a different set of control codes to do certain screen functions, such as cursor positioning or clearing the screen. Thus the following procedures of this module work only if the Turbo Modula-2 system has been properly installed for the terminal running the system. Note: It is possible to install stand-alone Turbo Modula-2 programs as well as the system (see Appendix B »Installation Procedures»).

Some of the screen functions available in the module *Terminal* are described here.

ClearScreen ();

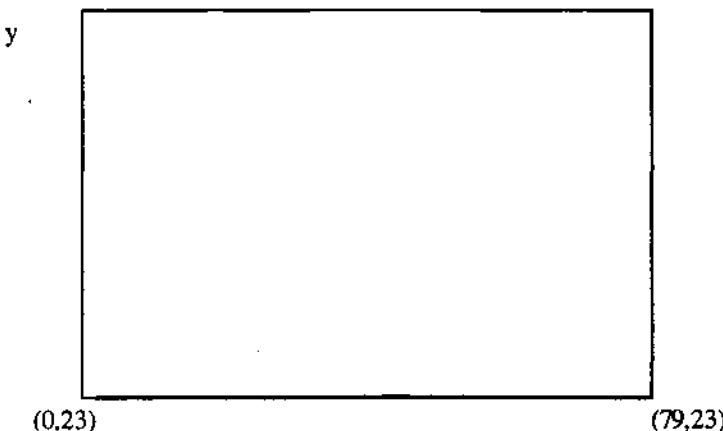
Clears the screen and homes the cursor (places it in the upper left corner of the screen).

GotoXY(x,y: CARDINAL);

Places the cursor on column *x*, row *y* of the screen. Columns and rows are numbered from 0. The upper left corner is at (0,0). A typical display has 80 columns and 24 lines, as shown in Figure 11-3.

Figure 11-3. Typical Display Using GotoXY.

(0,0) x (79,0)



Inside the module *Terminal* are two variables, *numRows* and *numCols*, set by

the terminal initialization to give the screen height and width. They are defined as follows:

VAR

```
numRows, numCols: CARDINAL;
```

where for common terminals, *numRows* = 24 and *numCol* = 80.

You can highlight a piece of text with the procedures *Highlight* and *Normal*.

(*Highlight* highlights output to the screen by underlining it, using inverse video, changing its brightness. This effect is dependent on the characteristics of your terminal and how the system is installed. *Normal* returns you to normal display mode.

The remaining procedures, *ClearToEOL*, *DeleteLine*, and *InsertLine*, are used mainly for text editing.

ClearToEOL deletes all characters from the cursor position to the end of the cursor line. *DeleteLine* deletes the cursor line completely and moves all following lines up one to close the gap. *InsertLine* inserts a blank line at the current cursor position.

After a call to any of these three functions, the cursor position is undefined. The cursor has to be placed explicitly with a call to *GotoXY*.

Since these functions are not available on all terminals, there is an enumeration type (*SpecialsOps*) and a set variable (*OpSet*) exported by *Terminal* that allow you to determine which functions were installed. They are shown in the following:

(**TYPE**

```
SpecialOps = (clearEol, insertDelete, highlightNormal);  
OpSet      = SET OF SpecialOps;
```

VAR

```
available : OpSet;
```

The variable *available* can be used as shown in the following examples. To clear to the end of a line, first determine if the procedure *ClearToEOL* is installed on the terminal.

The ComLine Module

1. This module provides access to the command line and handles the redirection of input and output. The command line consists of any arguments that follow the command that began the program. *ComLine* will work in both CP/M and Turbo Modula-2 modes. The command line is simply declared in *ComLine* as a text.

VAR

```
commandLine: TEXT;
```

Once imported, it can be read like any other text; for instance, *READ(commandLine,fileName)* works perfectly. Of course, this text is quite short, consisting of only a single line terminated by one EOT (32C) character. The text does not need to be opened explicitly; *ComLine* does it on initialization.

You may have noticed that all Turbo Modula-2 system programs accept arguments on the command line. If no argument is given there, the user is prompted to enter it. This is accomplished by a call to the following procedure:

```
PromptFor (prompt: ARRAY OF CHAR; VAR s: ARRAY OF CHAR);
```

This returns an input string in the result parameter *s*. *PromptFor* first tries to perform a *ReadString* from the command line. If no further data is found there, it displays the prompt string and performs a *ReadString* on *s* from the terminal. This method is user friendly, since it allows for interactive use and batch processing with command line arguments.

Redirection of Input and Output

The command line can contain redirection arguments: The symbol <, followed by a file name, redirects the standard text *input*; while the symbol >, followed by a file name, redirects *output*. Since many interactive programs will not work with I/O redirection, redirection of input or output is not automatic; instead, there are the procedures *RedirectInput* and *RedirectOutput*.

If an input redirection argument is given on the command line, *RedirectInput* connects the standard text *input* with the file or device specified by that argument. Similarly, *RedirectOutput* connects the standard text *output* with some device or file, provided an output redirection argument is included in the command line.

If the command line is read, any redirection arguments are ignored to avoid interference with normal command line arguments.

Note: Redirected files are not automatically closed when a program terminates. Therefore, if output goes to a disk file, you must close the text (and the associated file) explicitly with the *CloseText* procedure from *Texts*. *CloseText(output)* called before program termination will cause the redirected output file to be correctly closed.

ComLine Module Specification

```
DEFINITION MODULE ComLine;
FROM Texts IMPORT TEXT ;

VAR
  commandLine      : TEXT ;
  inName,outName   : ARRAY [0..19] OF CHAR;
  progName         : ARRAY [0..7] OF CHAR;

PROCEDURE RedirectInput;
PROCEDURE RedirectOutput;

PROCEDURE PromptFor (prompt:ARRAY OF CHAR; VAR s:ARRAY OF CHAR);

END ComLine
```

The names following the symbols < or > on the command line are saved in the two variables *inName* and *outName*. These are declared ARRAY [0..19] OF CHAR;. If no redirection arguments are given, both *inName* and *outName* are set to »CON:«.

The last item exported by *ComLine* is the variable *ProgName*. It is declared as ARRAY [0..7] OF CHAR;.

This contains the name of the executing program (without drive code or extension). If a program is called from CP/M, this string is empty.

As an example, for *Comline* the following program CO copies the standard input to the standard output text. Using I/O redirection, it can be used for a large number of purposes, displayed in the following:

CO <B:TEXTFILE	Displays the contents of TEXTFILE
CO <B:TEXTFILE>LST:	Prints out TEXTFILE
CO >B:TEXTFILE	Enters characters typed at the keyboard into TEXTFILE until a ^ Z is read
CO (no arguments)	Guess

```
MODULE CO;
FROM ComLine IMPORT RedirectInput,RedirectOutput;
FROM Texts IMPORT EOT,input,output,CloseText;
VAR
  ch: CHAR;
BEGIN
  RedirectInput; RedirectOutput;
  LOOP
    READ (ch);
    IF EOT(input) THEN EXIT END ;
    WRITE (ch);
  END ;
  CloseText(input); CloseText(output)
END CO.
```

The module *ComLine* has an unexported exception called *NoInputFile*. This ex-

ception is not usually trapped by user programs. If necessary it can be trapped with the **ELSE** clause in an exception handler, like so:

```
MODULE TrapUnexportedException;
FROM ComLine IMPORT RedirectInput, inName;
BEGIN
  WRITELN("inName is ",inName);
  RedirectInput;
EXCEPTION
  ELSE
    WRITELN('Must be non-existent file');
END TrapUnexportedException.
```

The exception is only raised if the name found in *inName* is not found. Thus, if »JUNK« is not a valid file, the following invocation of the program causes the error

Run MCD-file: MOO:UNEXPEXC <junk

```
inName is JUNK
Must be nonexistent file
```

Utility Modules

This section describes library modules that provide common mathematical and string-handling functions, and conversions to and from strings and numbers.

The **MathLib** and **LongMath** Modules

MathLib provides common mathematical functions, while *LongMath* provides the same functions for double-precision variables, replacing every declared REAL with a LONGREAL. Following are the functions provided by *MathLib*.

sqrt(x: REAL): REAL;

Computes the square root of the (positive) argument *x*. If *x* is negative, the exception ArgumentError is raised.

Exp(x: REAL): REAL;

Computes the natural exponent of the argument *x*. To prevent overflow from occurring, *x* must be smaller than 87.4; otherwise, the exception ArgumentError is raised.

Ln(x: REAL): REAL;

Returns the natural logarithm of the (positive) argument *x*. The exception ArgumentError is raised on negative *x*.

Sin(x: REAL): REAL;

Cos(x: REAL): REAL;

Returns sine and cosine of the argument *x*. The argument must be given in radians (radians = degrees * pi/180).

ArcTan(x: REAL): REAL;

Returns the arc tangent of argument *x*. The result is given in radians and lies within the interval of -pi/2 to pi/2.

```
DEFINITION MODULE MathLib;

PROCEDURE Sqrt      (x: REAL): REAL;
PROCEDURE Exp       (x: REAL): REAL;
PROCEDURE Ln        (x: REAL): REAL;
PROCEDURE Sin       (x: REAL): REAL;
PROCEDURE Cos       (x: REAL): REAL;
PROCEDURE ArcTan    (x: REAL): REAL;
PROCEDURE Entier    (x: REAL): INTEGER;

PROCEDURE Randomize(n: CARDINAL);
PROCEDURE Random ( ): REAL;

EXCEPTION ArgumentError;
(* Is raised if an argument is outside its legal range *)

END MathLib.
```

LongMath Module Specification

```
DEFINITION MODULE LongMath;

PROCEDURE Sqrt      (x: LONGREAL): LONGREAL;
PROCEDURE Exp       (x: LONGREAL): LONGREAL;
PROCEDURE Ln        (x: LONGREAL): LONGREAL;
PROCEDURE Sin       (x: LONGREAL): LONGREAL;
PROCEDURE Cos       (x: LONGREAL): LONGREAL;
PROCEDURE ArcTan    (x: LONGREAL): LONGREAL;
PROCEDURE Entier    (x: LONGREAL): LONGINT;

EXCEPTION ArgumentError;
(* Is raised if an argument is outside its legal range *)

END LongMath.
```

Entier(x: REAL): INTEGER;

Returns the integer part of a real number rounded toward negative infinity. For example:

if $x \geq 0$, then $\text{entier}(x) = \text{INT}(x)$
 if $x \leq 0$, then $\text{entier}(x) = \text{INT}(x-1)$

Thus,

$\text{Entier}(2.7) = 2$
 $\text{Entier}(-2.7) = -3$
 $\text{Entier}(-1.1) = -2$
 $\text{Entier}(1.1) = 1$

Other mathematical functions can be easily constructed by combining the preceding functions. For example:

$$\begin{array}{ll} \text{Tan}(x) & = \text{Sin}(x) / \text{Cos}(x) \\ x \text{ raised to the } y & = \text{Exp}(\text{Ln}(x)^*y) \end{array}$$

The last transformation is recommended only when y is not a whole number. For integer y , the following algorithm is more efficient:

```
MODULE PowerTest;

PROCEDURE power(x: REAL; n: INTEGER): REAL;
VAR
  i: CARDINAL;
  z: REAL;
BEGIN
  i := ABS(n); z := 1.0;
  WHILE i>0 DO
    (* z * x**i = x0**ABS(n) *)
    IF ODD(i) THEN z := z * x END ;
    x := x * x; i := i DIV 2
  END ;
  IF n>=0 THEN RETURN z
  ELSE RETURN 1.0/z
END
```

```
END power;

BEGIN
  WRITELN(power(23.9,2));
END PowerTest.
```

In addition to the previous functions, *MathLib* (but not *LongMath*) contains a random-number generator. There are two user calls to the random-number generator: *Randomize* and *Random*.

Random(): REAL;

Returns the next element of a pseudo-random-number sequence. Numbers are distributed evenly in the interval from 0.0 to 1.0.

Randomize(n: CARDINAL);

Reseeds the generator. A different random-number sequence corresponds to every possible argument *n*. Any random sequence may be repeated by providing *Randomize* with the same value of *n*.

Randomize takes one argument, a long (32 bit) integer, and sets the seed for the random-number generator equal to that number. *Random* takes no arguments, and returns a floating-point number *x* in the range $0 <= x < 1$. The number may be single precision or double precision.

Since there is a default value for the seed, most users need only call *Random*. *Randomize* is provided so that users can replicate simulations, since *Random* is a deterministic function of the seed. (In computer science jargon, it is a pseudo-random-number generator and not a true random-number generator.)

You should not call *Randomize* too frequently, because if *Randomize* is called for each call to *Random* then the numbers will not be much more random than the arguments to *Randomize*.

The least-significant bits in the real number returned by *Random* may not be particularly random.

The random-number generator maintains a 32-bit seed, which it treats as an unsigned integer. Each call to *Random* changes the seed by the formula

`seed := 134775813 * seed + 1 mod 4294967296`

and then returns the real number `seed / 4294967296`.

For example, if we call *Randomize* with the argument 1, then the seed is set to 1 and the next call to *Random* changes the seed to *134775814* and returns

`134775814 / 4294967296 + 0.03137994...`

No matter what the seed is, the random-number generator will pass through 4,294,967,296 numbers before it repeats.

The default value of the seed is 860098850.

The Strings Modules

A string in Modula-2 is an array with elements of type CHAR. The current contents of a string do not have to fill out the array completely; the end of the string is marked by a null character (0C). For convenience, there is the predefined type

TYPE

`String = ARRAY [0..80] OF CHAR;`

The module *Strings* offers a number of functions to manipulate strings, concatenate them, extract or delete substrings, and so on. Since their parameters are open arrays, they will work with any array of characters, not only with the predefined type *String*. Keep in mind, however, that formal, open array parameters map an actual array parameter into an index range starting at zero. Therefore, all index parameters passed to the following procedures are expected to have a range starting at zero. Thus, we recommend setting the lower bound to zero when you define your own string type. The exported string procedures follow.

Length (VAR str: ARRAY OF CHAR): CARDINAL;

Returns the length of the contents of the given array. The returned length does not need to coincide with the size of the *str*. If the array *a* is filled completely, *HIGH(a)+1* is returned; otherwise, *Length* will yield the index of the first null character encountered.

For example, let *Alfa* be declared as **ARRAY [0..100] OF CHAR**. Thus,

Alfa := ""; **WRITE(Length(*Alfa*))** will give 0
Alfa := "ABC"; **WRITE(Length(*Alfa*))** will give 3

Pos (substr,str: ARRAY OF CHAR): CARDINAL;

Pos scans for the first occurrence of the given string *substr* in string *str*. The procedure returns the index of the substring's first element. If *substr* is not contained in *str*, *HIGH(str)+1* is returned. Hence, this procedure can give back any value between 0 (if *substr* is a prefix of *str*) and *HIGH(str)+1* (if *substr* is not found).

Insert (substr: ARRAY OF CHAR; VAR str: ARRAY OF CHAR; inx: CARDINAL);

Inserts the given string *substr* into the given *str* at the index position *inx*. As always, indices are assumed to start at zero. If the array *str* is not large enough to hold the augmented string, the exception **StringError** is raised. Note that no error is raised if the *substr* is inserted at an index greater than the length of the string but still within the size of the string.

Strings Module Specification

```

DEFINITION MODULE Strings;
TYPE
  String = ARRAY [0..80] OF CHAR;

PROCEDURE Length (VAR str: ARRAY OF CHAR): CARDINAL;
PROCEDURE Pos      (substr,str: ARRAY OF CHAR): CARDINAL;
PROCEDURE Insert (substr: ARRAY OF CHAR; VAR str: ARRAY OF CHAR;
                  inx: CARDINAL);
PROCEDURE Delete (VAR str: ARRAY OF CHAR; inx,len: CARDINAL);
PROCEDURE Append (substr: ARRAY OF CHAR; VAR str: ARRAY OF CHAR);
PROCEDURE Copy    (VAR str: ARRAY OF CHAR; inx,len: CARDINAL;
                  VAR result: ARRAY OF CHAR);
PROCEDURE CAPS   (VAR str: ARRAY OF CHAR);

EXCEPTION StringError;

END Strings.

```

Delete (VAR str: ARRAY OF CHAR; inx,len: CARDINAL);

Deletes *len* characters in the given string *str*, starting at the element with index *inx*.

Append (substr: ARRAY OF CHAR; VAR str: ARRAY OF CHAR);

Appends the given string *substr* to the given *str*. If the array *str* is not large enough to hold the augmented string, the exception *StringError* is raised.

(Copy(VAR str: ARRAY OF CHAR; inx,len: CARDINAL; VAR result: ARRAY OF CHAR);

Assigns a substring of *str*, starting at position *inx*, consisting of *len* characters to string *result*.

CAPS (VAR str: ARRAY OF CHAR);

Changes all the characters in *str* to uppercase.

The following utility module demonstrates use of the *Strings* module. When dealing with files, it is often necessary to supply default extensions or to use several files with the same name but different extensions. These tasks can be simplified by using a library module like the following one:

```
DEFINITION MODULE FileExtensions;
TYPE
  Extension = ARRAY [0..2] OF CHAR;
PROCEDURE AppendExt (VAR str: ARRAY OF CHAR; ext: Extension);
PROCEDURE StripExt VAR str: ARRAY OF CHAR);
PROCEDURE GetExt    (str: ARRAY OF CHAR; VAR ext: Extension);

END FileExtensions.
```

The procedure *AppendExt* supplies a default extension if the given *str* does not contain an extension itself. *StripExt* removes any extension in *str*, while *GetExt* returns the extension of *str* or the empty string if *str* does not have an extension. Using the services exported by *Strings*, these procedures can be implemented as follows:

```
IMPLEMENTATION MODULE FileExtensions;
FROM Strings IMPORT Length, Pos, Append, Delete;

PROCEDURE AppendExt(VAR str: ARRAY OF CHAR; ext: Extension);
VAR len: CARDINAL;
BEGIN
  IF Pos(".", str) = HIGH(str)+1 THEN (* no "." in str *)
    len := Length(str);
    str[len] := ".";
    (* Strings go from 0 to Length(str)-1 *)
    Append(ext, str)
  END
END AppendExt;

PROCEDURE StripExt(VAR str: ARRAY OF CHAR);
VAR
  period: CARDINAL;
BEGIN
  period := Pos(".", str);
  IF period # HIGH(str)+1 THEN
    Delete(str, period, Length(str)-period)
  END
END StripExt;

PROCEDURE GetExt(str: ARRAY OF CHAR; VAR ext: Extension);
VAR
  period,elen: CARDINAL;
BEGIN
  period := Pos(".",str);
  elen := HIGH(str)-period;
  IF elen > 3 THEN elen := 3 END ;
  Copy(str, period+1, elen, ext)
END GetExt;

END FileExtensions.
```

The Convert Module

Convert exports procedures to convert strings into standard Modula-2 numeric types and back again. (*Doubles* provides the same services for double-precision real numbers.)

The first four procedures of the *Convert* module convert strings to INTEGERs, CARDINALs, LONGINTs, and REALs, respectively. The result is placed into the second parameter. The Boolean function result indicates whether a conversion is possible. It is TRUE if the string represents a legal number; otherwise, it is FALSE.

The second four procedures convert numbers to strings. The number is placed into the string starting at the right-most index. The left end of the string is padded with blanks. If the number does not fit into the given string, the exception TooLarge is raised. The *digits* parameter of procedure *RealToStr* has the same meaning as its equivalent in the procedure *WriteReal* (see the module *Texts*): It controls the size of the mantissa and the form of the string result (scientific or fixed point). For example, given the declaration

```
VAR  
  s: ARRAY [0..11] OF CHAR;
```

a call of *RealToStr*(2.5, *s*, 2); gives »2.50« and *RealToStr*(2.5, *s*, -4) gives »2.5000E+00«.

Convert Module Specification

```
TRCMD3 DEFINITION MODULE Convert;
```

```
PROCEDURE StrToInt (VAR s: ARRAY OF CHAR; VAR i: INTEGER) : BOOLEAN;  
PROCEDURE StrToCard (VAR s: ARRAY OF CHAR; VAR c: CARDINAL) : BOOLEAN;  
PROCEDURE StrToLong (VAR s: ARRAY OF CHAR; VAR l: LONGINT) : BOOLEAN;  
PROCEDURE StrToReal (VAR s: ARRAY OF CHAR; VAR r: REAL) : BOOLEAN;  
  
PROCEDURE IntToStr (i: INTEGER; VAR s: ARRAY OF CHAR);  
PROCEDURE CardToStr (c: CARDINAL; VAR s: ARRAY OF CHAR);  
PROCEDURE LongToStr (l: LONGINT; VAR s: ARRAY OF CHAR);  
PROCEDURE RealToStr (r: REAL; VAR s: ARRAY OF CHAR; digits: INTEGER);  
EXCEPTION TooLarge;  
END Convert.
```

The Doubles Module

Doubles exports procedures that provide support for double-precision real numbers (*LONGREALS*), similar to those provided for single-precision real numbers (*REALS*) in *Texts* and *Convert*.

As shown in the definition module, the variable *legal* monitors input from *ReadDouble* in a similar fashion to *Done* and *ReadReal* in *Texts*. The other routines are identical to their single-precision *REAL* counterparts.

Doubles Module Specification

```
DEFINITION MODULE Doubles;  
FROM Texts IMPORT TEXT;  
  
VAR legal: BOOLEAN;  
  
PROCEDURE ReadDouble (t: TEXT; VAR d: LONGREAL );  
PROCEDURE WriteDouble (t: TEXT; d: LONGREAL; n: CARDINAL;  
                      m: INTEGER);  
  
PROCEDURE StrToDouble (VAR s: ARRAY OF CHAR;  
                      VAR d: LONGREAL ): BOOLEAN;  
PROCEDURE DoubleToStr (d: LONGREAL;VAR s: ARRAY OF CHAR;  
                      digits: INTEGER);  
  
END Doubles.
```

System-Dependent Modules

The four modules described here provide utilities of a low-level nature. *Processes* and *Loader* are heavily dependent on the underlying system. They are coded in a high-level manner because *SYSTEM* and *STORAGE* provide a standard interface to machine-dependent operations such as memory allocation and process control.

The Processes Module

This module provides the synchronization between loosely coupled or largely independent processes. A process consists of a piece of program and a work space (its data. It is realized by the standard Modula-2 concept of a coroutine using the primitives *PROCESS*, *NEWPROCESS*, and *TRANSFER* in the pseudomodule *SYSTEM*.

A process takes the form of a parameterless procedure, which must be declared at the global level. Processes are not invoked by a procedure call, however. Instead, flow of control between processes is achieved via variables of type *SIGNAL*. Only two operations are applicable to signals: *SEND* or *Awaited*. When a process waits for a signal, its execution is paused; that is, control transfers to some other process. The waiting process is resumed when the awaited signal is sent by the currently executing process. More than one process can wait for the same signal. In that case, waiting processes are inserted into a queue, with the first process in the queue resuming once it has been sent the required signal. Sending an unanticipated signal has no effect.

The services exported by *Processes* are described next.

StartProcess(P: PROC; n: CARDINAL);

(
 A concurrent process with program *P* and a workspace of size *n* (in bytes). A minimum workspace consists of about 50 to 100 words. After a call to this procedure, control is transferred to the newly created process. For more information on this procedure, see the procedure *NEWPROCESS* in the *SYSTEM* module and Wirth's book.

Processes Module Specification

```
DEFINITION MODULE Processes;  
TYPE SIGNAL;  
  
PROCEDURE StartProcessP: PROC; n: CARDINAL);  
PROCEDURE SEND      (VAR s: SIGNAL);  
PROCEDURE WAIT      (VAR s: SIGNAL);  
PROCEDURE Awaited    (s: SIGNAL): BOOLEAN;  
PROCEDURE Init       (VAR s: SIGNAL);  
  
EXCEPTION DeadLock;  
  
END Processes.
```

SEND(VAR s: SIGNAL);

If there is no process waiting for *s*, *SEND* has no effect. If one is waiting, it is resumed. If several processes are waiting for the same signal, they are inserted into a queue. The first process in the queue is resumed when the call to *SEND(s)* occurs.

WAIT(VAR s: SIGNAL);

Causes the currently executing process to be suspended until it receives the signal (Processes ready for execution are ordered in a queue. After a call to *WAIT*, the next process in the queue is resumed. If no process is ready for execution (that is, all processes are waiting), the exception *DeadLock* is raised.

Awaited(s: SIGNAL): BOOLEAN;

Returns TRUE if at least one process is waiting for *s*; otherwise, it returns FALSE.

Init(VAR s: SIGNAL);

Initialization of a *SIGNAL*; this is mandatory before a *SIGNAL* is used.

DeadLock;

This exception is raised if all current processes are waiting for a signal; it indicates an error in your signal-sending logic. A procedure constituting the program of a process cannot be called, nor can it return like a normal procedure. Its instructions are usually enclosed in a **LOOP** statement without a corresponding **EXIT**. To enable transfer of control to some other process, such a procedure must contain at least one *SEND* or *WAIT* statement.

A state can arise where all processes are waiting for some signal. In this case, no process can execute and subsequently the system will come to a halt. This situation is known as a *deadlock*, or *deadly embrace*. When it is detected by *Processes*, the exception *DeadLock* is raised.

For example, when polling a number of external devices; assume the computer is connected to a number of external devices that can be serviced independently from one another. The computer will ask each device: Do you require service?

When a service request is found, it sends out the message: Device No. *n* needs service. If there is a process ready to serve this device, it then takes control of the processor. When it is finished, this process starts the polling again by putting itself into a wait state for the next request. If the service is time-intensive, it could just as well enable polling repeatedly in the middle of its operation. This is achieved by waiting for the signal RESUME, which transfers control to a further location through the polling routine. Only when no further service is needed is the signal RESUME sent, causing one of the processes waiting for it to be resumed. A sketch of a module that services devices by polling follows:

```
MODULE Poll;
FROM Processes IMPORT SIGNAL, SEND, WAIT, Awaited, Init,
StartProcess;
VAR
Service1,...,ServiceN, Resume : SIGNAL;

PROCEDURE Device1;
BEGIN
LOOP
WAIT(Service1);
(* "Service device #1"*)
IF (*"Service is complicated") THEN
(* Check if there are other devices waiting for service *)
WAIT(Resume)
END ;
(* "Resume service of device #1"*)
END
END Device1;

PROCEDURE Device2;
(* Body of Device2 similar to Device1 *)
END Device2;

.

.

.

PROCEDURE DeviceN;
(* Body of DeviceN similar to Device1 *)
END DeviceN;
```

```
BEGIN (* Main program, containing the polling loop *)
  Init(Service1); ... Init(ServiceN);
  Init(DeviceCheck); StartProcess(Device1,100);
  StartProcess(Device2,500);

  .
  .

  StartProcess(DeviceN,200);
LOOP
  IF (*"ServiceRequest of Device #1"*) & Awaited(Service1) THEN
    SEND(Service1)
  ELSIF (*"ServiceRequest of Device #2"*) & Awaited(Service2) THEN
    SEND(Service2)

  .
  .

ELSIF (*"ServiceRequest of Device #N"*) & Awaited(ServiceN) THEN
  SEND(ServiceN)
ELSE
  (* If there are any pending services, resume one of them *)
  SEND(Resume)
END
END
END Poll.
```

Note that the polling process encapsulated in the main program does not contain any *WAIT* statements. It is always available for execution, excluding the possibility of a deadlock. Every other process in this module pauses only by executing a *WAIT* statement; therefore, the only process ready for execution after issuing a wait is the polling process--a guarantee that polling occurs regularly.

A more rudimentary polling scheme can be implemented without the notion of processes and signals. In this scheme, the polling loop calls the service procedures directly when a service request is detected. However, in order to express a feature like the resume operation conveniently, processes are essential.

The Pseudomodule SYSTEM

This module contains types and procedures for low-level programming. An implementation part for this does not exist; rather, the compiler translates all exported operations directly into code, which is why *SYSTEM* is called a *pseudomodule*. Everything in it must be imported to become available, thereby making the use of low-level constructs explicit. The following sections describe the definition part of *SYSTEM*.

Low-Level Access to Data

The type *WORD* represents a machine word (2 bytes, 16 bits). Every type occupying a word or less is assignment-compatible with *WORD*. In particular, a formal parameter of type *WORD* accepts any actual parameter occupying one machine word. No operations (except assignment) can be performed on variables of type *WORD*. However, a variable of type *WORD* can be the argument of a type-transfer function.

The type *BYTE* represents a byte. It can be considered a subrange of type *WORD*. A *BYTE* can be assigned any value with a decimal equivalent in the range of 0 to 255.

Array of Word

If a formal parameter is declared as an **ARRAY OF WORD**, its corresponding actual parameter can be of any type. The upper bound of the open array is adjusted to match the size of the actual parameter; thus, no restrictions are imposed on a procedure's arguments. The most common examples of these procedures are *ReadRec* and *WriteRec* found in the library module *Files*.

Pseudomodule SYSTEM Specification**DEFINITION MODULE SYSTEM;****TYPE**

WORD; BYTE; ADDRESS; PROCESS;

VAR

IRESULT, HRESULT: CARDINAL

PROCEDURE ADR(VAR v: AnyType): ADDRESS;**PROCEDURE TSIZE(AnyType): CARDINAL;****PROCEDURE TRANSFER(VAR source, dest: PROCESS);****PROCEDURE IOTRANSFER(VAR source, dest: PROCESS; n: CARDINAL);****PROCEDURE NEWPROCESS(p: PROC; a: ADDRESS; n: CARDINAL;****VAR q: PROCESS);****PROCEDURE BIOS(n: CARDINAL; w: WORD): CARDINAL;****PROCEDURE BDOS(n: CARDINAL; w: WORD): CARDINAL;****PROCEDURE CODE(AnyStringLiteral);****PROCEDURE MOVE(source,dest: ADDRESS; len: CARDINAL);****PROCEDURE FILL(adr: ADDRESS; len: CARDINAL; val: BYTE);****PROCEDURE INP(port: WORD): CARDINAL;****PROCEDURE OUT(port: WORD; data: CARDINAL);****EXCEPTION OVERFLOW, REALOVERFLOW, OUTOFMEMORY, BADOVERLAY;****END SYSTEM.**

ADDRESS

ADDRESS is a type that represents memory locations. It can be thought of as a *POINTER TO WORD*; thus, dereferencing an *ADDRESS* will yield a *WORD*. Addresses can be computed; it is possible to add and subtract *CARDINALS* from addresses with the result being of type *ADDRESS*. *ADDRESS* is compatible with a *CARDINAL* and with every pointer type. In particular, if a formal parameter is of type *ADDRESS*, the corresponding actual parameter may be of any pointer type.

ADR

The procedure *ADR* returns the address of its argument. It is declared as

```
ADR(VAR x: "AnyStructuredType"): ADDRESS;
```

Note that since Turbo Modula-2 implements register variables in local procedures (see Chapter 8, «Low-Level Facilities»), the *ADR* function will not accept simple variables of unstructured types. This is because a scalar variable in a procedure may reside in a register that has no address.

TSIZE

The procedure *TSIZE* accepts the name of a type as its argument. It returns the minimum size of any variable of the argument type. Thus, when variable *x* is declared as *VAR x: T*, then *SIZE(x) = TSIZE(T)*.

TSIZE has an alternate syntax similar to the *NEW* procedure: If *T* is a record type whose last field is a variant, a tag value of this variant can be indicated behind the type. In this case, *TSIZE* returns the amount of memory used by a variable created with a call to *NEW* with the same tag-value indication. If the last variant contains another variant as its last field, a tag value for this variant may be given, too. Thus, *TSIZE* is informally declared as

```
TSIZE("AnyType")
```

or

```
TSIZE("AnyType", TagValue, TagValue, ...)
```

Coroutines and Interrupts

Modula-2 includes the concept of a coroutine, which consists of instructions and a workspace separate from the main program. A coroutine can represent a process or detail how statements will manipulate data. (The main program itself is a process, since it consists of statements acting on data.) Coroutines lead to the introduction of loosely coupled processes. At any one time, a single-processor computer can execute only one process. The pseudomodule *SYSTEM* exports the type *PROCESS* to be used to reference a coroutine.

Although the body of a coroutine looks like that of a parameterless procedure, its execution is more like that of a main program. A coroutine cannot be called in a call statement or exited by a **RETURN** statement—coroutines should never reach the final end. Since coroutines are never exited by a **RETURN** or **END** statement, their procedure body is usually enclosed by a **LOOP** statement without a corresponding **EXIT**.

Control to and from a coroutine is done via the *TRANSFER* statement. When a coroutine transfers control away from itself, its state is preserved until it is reactivated by a transfer from another coroutine. When a coroutine's state is preserved, all variables retain their values and modules remain initialized. When reactivated, the coroutine continues with the statement directly after the last *TRANSFER*.

Note that the program bodies of coroutines and processes (as described in the section, »The Processes Module«) are remarkably similar. In fact, processes are modeled by coroutines. The module *Processes* implements *StartProcess*, *SEND*, and *WAIT* with the primitives that follow.

NEWPROCESS(p: PROC; a: ADDRESS; n: CARDINAL; VAR q: PROCESS);

Creates a coroutine and initializes the given process *q* with procedure body *p* and a workspace of *n* words starting at address *a*. This workspace must be large enough to contain *p*'s local variables, local stack, and local heap. A subsequent *TRANSFER(r,q)* will transfer control to the first statement of procedure *p*. *PROC* is a standard procedure type representing a parameterless procedure. It is exported from the *SYSTEM* module as

```
TYPE  
PROC = PROCEDURE( );
```

TRANSFER(VAR source, dest: PROCESS);

Transfers control to the coroutine *dest*. The currently active coroutine is assigned to *source* and suspended. A subsequent *TRANSFER* to *source* will resume the suspended coroutine with the next statement after the transfer.

Coroutines are used to express interrupt handling, such that a coroutine waiting for an interrupt will transfer control to some other coroutine (the main program) to do something in the meantime. When the interrupt occurs, control is transferred back to the interrupt-handling coroutine by an unscheduled *TRANSFER* statement.

It is as if a transfer instruction is explicitly inserted into the executing coroutine at the point of interrupt. Since we do not know the state of execution of the interrupted routine, no explicit *TRANSFER* can be used. Thus, a special transfer instruction is provided to allow unscheduled transfers.

IOTRANSFER(VAR source, dest: PROCESS; n: CARDINAL);

Performs a normal transfer from *source* to *dest*. When the interrupt specified by the given *n* occurs, control is transferred back to the next statement after *IOTRANSFER*. Any available interrupt vector can become *n*.

Interrupts must sometimes be disabled to avoid disturbing critical program parts. This is done by placing critical procedures in a module that is given a priority. Such a module is called a *monitor*. The priority (a numeric value) is indicated in square brackets after the module name; for example:

```
MODULE CriticalStuff[1];
```

IMPLEMENTATION MODULE Processes[1];

When a procedure in a monitor is executed, only interrupts of a higher priority than the monitor are enabled. However, in Turbo Modula-2 there are no interrupts with priorities since the Z80 hardware does not usually recognize them. All interrupts are disabled when a piece of a monitor is executed; that is, any number in brackets after the module name shuts out all (maskable) interrupts.

Warning: CP/M is not designed for interrupts; in particular, it is not reentrant. Therefore, when a CP/M operation is suspended by an interrupt, the operation

(and possibly others as well) cannot be invoked by the interrupting routine. (Since interrupts are only of limited use when dealing with CP/M, those of you interested in a more detailed discussion can refer to Wirth's book.)

Z80-Specific Procedures

The following procedures allow the programmer to take direct advantage of certain Z80 machine instructions. *MOVE* and *FILL* are very fast and their use is recommended anywhere speed is important. *INP* and *OUT* are for accessing the Z80 I/O ports. The native code generator produces highly efficient inline code for *INP* and *OUT*.

MOVE(source,destin: ADDRESS);

Moves blocks of data in memory. Its operation is also appropriate for blocks of overlapping memory: If *source* < *destin*, then the upper addresses are moved first; and if *source* >= *destin*, then the lower addresses are moved first.

FILL(adr:ADDRESS; nbytes: CARDINAL; val: BYTE);

Fills the block of memory starting at *adr*, which is *nbytes* large, with the byte value *val*.

INP(port: WORD): CARDINAL;

Receives (inputs) a value from an I/O port.

OUT(port: WORD; data: CARDINAL;);

Outputs a value to an I/O port.

interface to CP/M

The next two procedures allow you to directly access the operations of CP/M.

BDOS(n,w: WORD);

Results in a call to memory location 5, the CP/M entry location. The function number is specified in the given *n*; *w* denotes an additional parameter. If some *BDOS* operation returns a result in the accumulator *A*, it can be found in the variable *IORESULT*. If an operation returns a 16-bit result in register *HL*, it can be retrieved from the variable *HLRESULT*.

IORESULT and *HLRESULT* are declared as CARDINALS:

```
VAR HLRESULT, IORESULT: CARDINAL;
```

The action performed by BDOS is shown in the following:

- Load register C with *n*.
- Load register DE with *w*.
- CALL 5.
- Assign the contents of accumulator A to *IORESULT*.
- Assign the contents of register HL to *HLRESULT*.

```
BIOS(n,w: WORD): CARDINAL;
```

Results in a call to the BIOS jump vector. The given *n* indicates which entry is called: 0=WARMBOOT, 1=CONSTAT, 2=CONIN, and so on. Its exact operation is as follows:

- Load A with the given *n*.
- Load BC with the given *w*.
- CALL *WARMBOOT+n*3*.
- Assign the contents of accumulator A to *IORESULT*.
- Assign the contents of register HL to *HLRESULT*.

For example, to read a character directly from the console, do the following:

```
BDOS(1,0); ch:=CHR(IORESULT);
```

To do the same using a call to the BIOS vector:

```
BIOS(2,0); ch:=CHR(IORESULT);
```

Assembler Interface

Assembly routines can be included in a Turbo Modula-2 program by using

```
CODE("Any string literal");
```

A call to *CODE* can be inserted behind a procedure heading, replacing the body of the procedure. Its body may not contain a declaration part, BEGIN statement, or any other statements. The final END must be present. Compilation of a *CODE*

statement causes the .COM file whose name is given as an argument to be inserted into the object program. For example:

```
PROCEDURE MOVE(source,dest: ADDRESS; n: CARDINAL);  
  CODE( "MOVE" )  
END MOVE;
```

establishes *MOVE* as an assembly routine. During compilation, the file *MOVE.COM* is read from disk and inserted into the object program. A call to *MOVE* will execute the assembly routine *MOVE.COM*. The file extension *.COM* is supplied automatically.

Assembly routines can have an arbitrary number of parameters and can return results; that is, they can be used as function procedures. It is essential to declare function procedures with the correct result type. Assembly routines that return a result when none is declared can crash the system.

Since where an object program will be loaded is unknown at compile time, included assembly routines must be fully relocatable. Only relative jumps may be used, and calls to subroutines or absolute access to data within the routine are not allowed. Furthermore, to determine the exact size of a routine, it must contain its length in the first 2 bytes. These bytes are stripped off by the compiler; they are not executed.

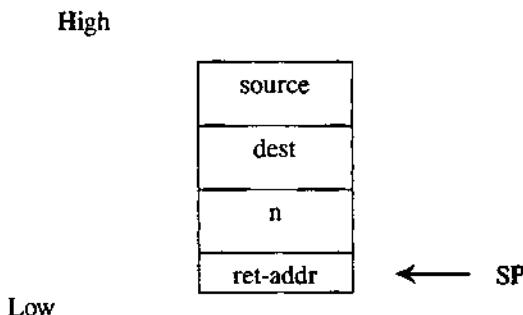
An assembler routine can expect the following environment:

- () Before an assembly procedure is called, its parameters are pushed onto the stack exactly as they are declared in the procedure heading. Value parameters of unstructured types are loaded immediately. VAR parameters and parameters of an **ARRAY** or **RECORD** type are loaded by address. The return address is at the top of the stack.
- () Upon entry, register HL points to the beginning of the routine's object code. Thus, data embedded in the program text can be retrieved using HL plus an index.
- () If an assembly routine returns a result, it is pushed onto the stack. In this case, be sure that all parameters have been removed from the stack and the return address is saved.

- The stack itself is 64 bytes deep. If more space is needed, it must be relocated and restored upon return. All registers can be destroyed.

To continue with the *MOVE* example, we will show how the stack is set up after a call to *MOVE*, but before any of *MOVE*'s instructions have executed. *MOVE*'s parameter *source* was pushed first, followed by *dest*, and then the count *n*. The return address is on the top of the stack and the current PC value is in HL.

For example, the parameters for the procedure *MOVE* are situated on the stack as follows:



The assembly code for *MOVE* pops all parameters off the stack, replacing the return address. It then checks the count word and if it is greater than zero, it performs the move and returns.

A possible implementation of *MOVE* is shown in the following:

```

ORG 100H           ;unimportant, must be relocatable

DEFW MOVEEND-MOVE ;indicate size of object code

MOVE POP HL        ;save return address
      POP BC        ;load BC with count n
      POP DE        ;load DE with destination
      EX  (SP),HL    ;load HL with source and
                      ;put return address back on stack
      LD   A,B        ;if BC=0 then return
      OR   A,C
      RET

```

LDIR ;Move

RET

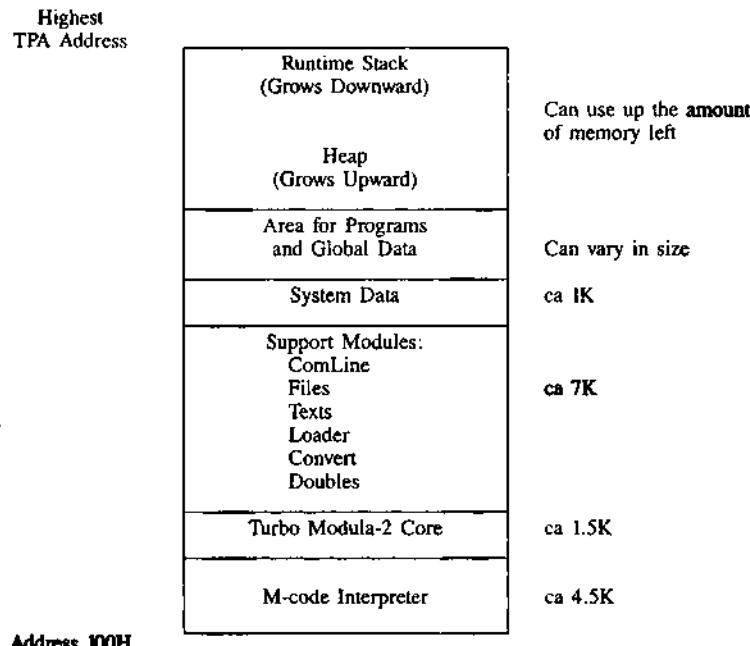
MOVEEND **END**

Modules in Memory Management

This section introduces the modules *STORAGE* and *Loader*. *Storage* allocates and deallocates dynamic data, while *Loader* loads and executes overlay modules. Since both modules affect the way data and programs are held in memory, let's first take a look at the memory organization of Turbo Modula-2.

The memory available for Turbo Modula-2 is the *transient program area* (TPA) of CP/M. The TPA starts at address 100 hex and extends to the start of the BD part of CP/M. How Turbo Modula-2 uses this area is shown in Figure 11-5.

Figure 11-5. Turbo Module-2's Use of TPA



At the bottom of the TPA is the M-code Interpreter. All other parts of the Turbo Modula-2 system and all application programs written within it and not compiled with the native code option exist in an intermediate code called M-code. This is a variant of the instruction set of Wirth's Modula computer, Lilith. The code

optimized for space conservation. Every M-code instruction is interpreted by this part of the system.

The object code of the Turbo Modula-2 core resides directly above the interpreter. It constitutes a runtime supervisor that has the main tasks of setting up the rest of the system and diagnosing and handling exceptions (errors).

There are a few support modules used by virtually every program: *Texts*, *Files*, *Loader*, *Convert*, and *ComLine*. A program's loading time can be improved if these modules are already present in memory. Thus all of these modules are linked together in the file M2.COM, and are loaded into memory when the system is started.

The resident support modules occupy about 7K of main memory. If an application does not need these modules and memory space is at a premium, the program can be transformed into a .COM file using the linker (see Chapter 10 for more about the linker's operation). The linker will include into the .COM file only those modules that are needed. This can result in a memory savings of up to 6K.

There is space for other programs and their data above the area occupied by M2.COM. Such programs include the system file SHELL.MCD, which interprets the user's commands and the compiler COMPILE.MCD as well as application programs. The object modules together with their global data reside at the bottom of this area.

The remaining space is partitioned dynamically between the runtime stack for local data and the heap. The heap holds all variables that were created dynamically by the *NEW* and *ALLOCATE* procedures. The stack is used for all local variables, procedure parameters, and sometimes to evaluate complicated arithmetic or logical expressions. A program is out of memory when the runtime stack and the heap meet; in that case, the exception OUTOFMEMORY is raised.

The STORAGE Module

This module handles allocation and deallocation of dynamically created variables. All such variables reside in the heap area. Like *SYSTEM*, *STORAGE* is a pseudomodule. It has neither an object file (*STORAGE.MCD*) nor a symbol file (*STORAGE.SYM*). Instead, all the following services offered by this module are built into Turbo Modula-2. The following procedures are provided by the module *STORAGE*.

ALLOCATE(VAR a: ADDRESS; size: CARDINAL);

Allocates a space of *size* bytes on the heap and makes the pointer *a* point to the beginning of the allocated area.

Note: Due to the allocation scheme employed by *STORAGE*, a minimum of 64 bytes is always allocated, even if the *size* parameter specifies a smaller value.

DEALLOCATE(VAR a: ADDRESS; size: CARDINAL);

Releases the area referred to by the given pointer *a*. The released area is assumed to have the given *size*. The *size* must be equal to the size indicated when the area was created. The deallocated area is available for future allocation of other data. The pointer *a* is set to nil.

Standard Procedures Dependent on Storage

NEW(a: POINTER TO »AnyType«);

NEW is translated into a call to *ALLOCATE* of the form *ALLOCATE(a, TSIZE(»AnyType«))*. *ALLOCATE* must be imported if the standard procedure *NEW* is used.

DISPOSE(a: POINTER TO »AnyType«);

DISPOSE is translated into a call to *DEALLOCATE* of the form *DEALLOCATE(a, TSIZE(»AnyType«))*. *DEALLOCATE* must be imported if a program makes use of the standard procedure *DISPOSE*.

Heap Pointer

The top of the heap is marked by the heap pointer, which is inaccessible to a

plication programs. The heap can contain usable space because of the deallocation of variables. If a variable residing in the middle of the heap is disposed, the released area is appended to a *free list*. There is one exception to this rule: If the variable residing at the high end of the heap is disposed, its area is not appended to the free list; instead, the heap pointer is decremented and the heap is compressed. When a variable is allocated, the free list is searched to find an area large enough to hold the required data. If the search is unsuccessful, the area is allocated at the top of the heap; that is, the heap pointer is incremented. Thus, the heap status consists of two elements: (1) the heap pointer, marking its top end, and (2) the free list, containing information about holes available for future allocation.

STORAGE Module Specification

```
DEFINITION MODULE STORAGE;  
FROM SYSTEM IMPORT ADDRESS;
```

```
PROCEDURE ALLOCATE (VAR a: ADDRESS; size: CARDINAL);  
PROCEDURE DEALLOCATE (VAR a: ADDRESS; size: CARDINAL);
```

```
PROCEDURE MARK (VAR a: ADDRESS);  
PROCEDURE RELEASE (VAR a: ADDRESS);
```

```
PROCEDURE FREEMEM ( ): CARDINAL;
```

```
END STORAGE.
```

A user may want to handle the heap in a stack-like fashion; for example, create a number of variables and then throw them away all at once. The following procedures allow for this.

MARK(VAR a: ADDRESS);

Creates a new heap segment. The new segment is directly above the old heap pointer. Its free list is initially empty. Until a call to *RELEASE(a)*, all allocations will allocate space in the new segment. The given address *a* is made to point to the bottom of the segment; that is, it equals the old heap pointer.

RELEASE(VAR a: ADDRESS);

Releases the heap segment created by a call to *MARK(a)*. The heap pointer and free list of the old segment are restored. The heap pointer is reset to the given address *a*. This address must have been set by a previous call to *MARK(a)*. The address *a* has the value nil after a call to *RELEASE(a)*.

In a newly created heap segment, variables can be allocated and disposed of in the normal way. While allocation is always safe, disposing presents some problems. The user must make sure that only variables residing in the currently active heap segment are disposed. Disposing of a variable that resides in a segment buried below the one last created by a call to *MARK* will result in a *BadHeap* exception. Similarly, disposing of a variable in an already-released segment causes the same exception to be raised.

The following instruction sequence is legal:

```
MARK(a);
NEW(p);
NEW(q);
DISPOSE(p);
RELEASE(a)
```

And the following sequences result in the exception *BadHeap*:

```
NEW(p);
MARK(a);
DISPOSE(p);    (Disposing a variable in some buried
                segment is not allowed)

MARK(a);
NEW(p);
RELEASE(a);
DISPOSE(p);    (p is already released)
```

Note: Creating or opening a file allocates space on the heap. When the file or text is closed, this area is disposed again. As a result, the preceding precautions must also be taken when a file or text is closed. For example, the following is illegal:

```
Open(f, "INDATA");
MARK(a);
Close(f)
```

The user can interrogate the space that remains for heap allocation with the next procedure.

FREEMEM(): CARDINAL;

Returns the number of bytes between the current value of the heap pointer and the current top of the runtime stack. The size of available holes contained in the free list is not taken into account. For this reason, *FREEMEM* gives somewhat pessimistic information about the space available for heap allocation. The exact number of available bytes is returned only if procedures *ALLOCATE* and *DEALLOCATE* are used in a stack-like manner.

Dynamic Variable Errors

While dynamic variables are a very powerful concept, they must be used with some care. The following are common sources of errors and how the module *STORAGE* reacts to them:

- A variable is disposed when there are still other references to it. In this case, the value(s) of the disposed variable is undefined. Worse, another variable may have been allocated in the same area. Assignment to the disposed variable would result in overwriting other data. Assignment could also ruin the free-list management, which would result in raising the exception BadHeap.
- The heap becomes too big. If the heap pointer overruns the runtime stack during an allocation request, no space is allocated and the exception OUTOFMEMORY is raised. Since OUTOFMEMORY is exported by SYSTEM, it can be handled by an exception handler in the application program. Note that it is possible to generate another OUTOFMEMORY exception in the exception handler by calling a procedure with many or large parameters (for example, *WriteString* or other I/O statements). Thus, exception handlers for OUTOFMEMORY must have enough memory to operate correctly.
- A variable not residing in the active heap segment is disposed. This error can only occur if MARK, RELEASE, and DISPOSE are not executed in the correct order. This problem will usually cause the exception *BadHeap* to be raised some time in the future.

The Loader Module

The *Loader* module handles the allocation and execution of object modules and provides a standard method to activate *overlays*. These services may be necessary if an application requires more memory than is currently available.

Note: The Turbo Modula-2 system linker provides a more flexible and convenient method for defining overlays, which does not require any code modifications (see the section, "The Linker," in Chapter 10). You should consider using it before using the procedures described here.

If you get an OUTOFMEMORY error during program execution, your program needed more memory than was available. Overlays provide one possible solution to memory problems, they work as follows: A large program usually consists of some number of largely independent pieces, not all of which have to be physically in memory at the same time. The Loader Module provides a method to dynamically load only those modules that are required. Once their execution is terminated, they disappear, freeing the memory space they occupied. Other program modules can be loaded in their place, and executed. Thus the memory re-

quired to run a large program system can be significantly less than the sum of all modules used by the program. *Loader* provides the *Call* procedure to dynamically load a module.

Call (modName: **ARRAY OF CHAR**);

Loads the object module specified by modName into main memory and starts its execution. When execution is complete, the memory space occupied by the object code and data is released again. Modules executed by *Call* involves memory management and disk operations in addition to normal procedure execution time.

As an example, *Call* ("OVER") loads and begins execution of the object file OVER.MCD. The called module can import and access all objects of the calling program. Direct communication in the reverse direction is not possible, because the called module and its data area disappear after execution.

Note: Direct communication between a called module and its caller is not possible. This can be solved by the use of objects (message buffers) imported by the called procedure from the caller. Such objects can be freely modified by the called procedure.

Variables created by an overlay with the *NEW* procedure are preserved after termination of the overlay.

Call searches for the specified module on all logged disk drives (as specified in the .MCD search path during installation). The variable *FirstDrive* can be used to indicate which drive should be searched first for any overlay modules. This can be used to improve the loading time for external overlays.

VAR

firstDrive: [0..15]

Indicates the first drive to be accessed when searching for the next overlay. The encoding is 0 for drive A, 1 for drive B, and so on. An application may set *firstDrive* to provide the drive where the next execution of *Call* should start its search. This is useful if many overlays are used and the program does not know which drive the overlay files will be on until runtime.

If the loading of a module is unsuccessful, the exception *LoadError* is raised.

(The reason for an unsuccessful load is displayed in an additional message.) LoadError can occur because of the following:

- The object file is not found.
- There is not enough space to hold the loaded object modules.
- A version conflict is detected. Turbo Modula-2 maintains a system-wide version-control scheme. If objects of a definition module are imported, the importing module remembers the version of this definition module. If the definition module is later changed and recompiled, all client modules have to be recompiled as well. If this is not done, a version conflict results when the client module is loaded. Note that if a definition module is recompiled but not changed, its version is preserved.

The following example shows communication between overlay modules. The main module is declared as an implementation module so that it may export a communication buffer to the overlay modules. Overlays are controlled via this record.

```
DEFINITION MODULE Main;
TYPE
  Commands = ( SayHello, CallOver, none );
VAR
  OverlayCommBuffer: RECORD
    Command: Commands;
    StrBuf: ARRAY [0..80] OF CHAR;
  END ;
END Main.
```

The main program imports the *Call* procedure from the *Loader* module as shown in the following:

```
IMPLEMENTATION MODULE Main;
FROM Loader IMPORT Call;
FROM Terminal IMPORT BusyRead;

PROCEDURE Pause;
VAR ch: CHAR;
```

```
BEGIN
  REPEAT BusyRead(ch) UNTIL ch#0C;
END Pause;
PROCEDURE DisplayOverString;
BEGIN
  WRITE('The string communication buffer ');
  IF OverlayCommBuffer.StrBuf #' ' THEN
    WRITELN('contains ',OverlayCommBuffer.StrBuf,'')
  ELSE
    WRITELN('is empty.')
  END ;
END DisplayOverString;

BEGIN
  WRITELN('Main program about to call overlay one.');
  OverlayCommBuffer.Command := SayHello;
  OverlayCommBuffer.StrBuf := '';
  Call('OVER1');
  WRITELN('Return from call to overlay 1');
  DisplayOverString;
  WRITELN('Now call overlay one again with a null request.');
  WRITELN('Will it be reloaded, press a key to see . . .');
  Pause;
  OverlayCommBuffer.Command := none;
  OverlayCommBuffer.StrBuf := '';
  Call("OVER1"); WRITELN('Returned from second call to overlay 1',
    'did the drive light go on ?');
  DisplayOverString;
  Pause;
  WRITELN('Call overlay two, which will call overlay one.');
  OverlayCommBuffer.Command := CallOver;
  OverlayCommBuffer.StrBuf := '';
  Call("OVER2");
  WRITELN('Return from call to overlay 2',
    'the drive light did go on, right!');
  DisplayOverString;
  WRITELN('That's all folks!');
  Pause;
END Main.
```

The overlay modules are defined as follows; notice that they also import the *Call* procedure from the *Loader Module*.

```

MODULE Over1;
FROM Loader IMPORT Call;
FROM Main IMPORT OverlayCommBuffer,Commands;

BEGIN
  CASE OverlayCommBuffer.Command OF
    SayHello: WRITELN('Hello there. This is overlay one running');
    WRITELN('Message left in the string area. Goodbye!');
    OverlayCommBuffer.StrBuf :=
      'Message from overlay 1 passed to caller';
    |
    CallOver: WRITELN('Overlay one about to call overlay two. ');
    WRITELN('Say hello number two . . .');
    OverlayCommBuffer.Command := SayHello;
    OverlayCommBuffer.StrBuf := '';
    Call('OVER2');
    IF OverlayCommBuffer.StrBuf # '' THEN
      WRITELN('Overlay two passed me this string: ',
        OverlayCommBuffer.StrBuf );
    END ;
    |
  ELSE
    WRITELN('Overlay two : no command passed');
  END ;
END Over1.

MODULE Over2;
FROM Loader IMPORT Call;
FROM Main IMPORT OverlayCommBuffer,Commands;

BEGIN
  CASE OverlayCommBuffer.Command OF
    SayHello: WRITELN('Hello there.');
    WRITELN('This is overlay two running');
    WRITELN('Message left in the string area.');
    WRITELN('Goodbye!');
    OverlayCommBuffer.StrBuf :=

```

```
'Message from overlay 2 passed to caller';
|
CallOver: WRITELN('Overlay two about to call overlay one. ');
WRITELN('Say hello number one . . .');
OverlayCommBuffer.Command := SayHello;
OverlayCommBuffer.StrBuf := '';
Call('OVER1');
IF OverlayCommBuffer.StrBuf # '' THEN
  WRITELN('Overlay one passed overlay two
           this string: ',
          OverlayCommBuffer.StrBuf );
END ;
|
ELSE
  WRITELN('Overlay two : no command passed');
END ;
END Over2.
```

Loader Module Specification

```
DEFINITION MODULE Loader;
VAR
  FirstDrive: [0..15];

PROCEDURE Call(modName: ARRAY OF CHAR);

EXCEPTION LoadError;
END Loader.
```

Chapter 12

Turbo Modula-2 Reference Directory

This chapter provides a complete alphabetical directory to Turbo Modula-2's standard identifiers, extensions, library procedures and modules, and reserved words.

Each entry contains a definition, usage and any restrictions, general comments, an example, a declaration, and a cross reference(s) to similar items where appropriate.

The following notation is used throughout this look-up section to describe operators, statements, types, procedures, and/or modules.

- () Where appropriate, module names are enclosed in parentheses after an entry.
- [] Items enclosed in square brackets are optional; the items can be used only once or not at all.
- { } Items enclosed in curly brackets can be repeated any number of times, including none.
- "[" , "]" , "{ " , " }" The use of square or curly brackets within quotation marks means the bracket is part of the syntax and not just notational.
- < > Items enclosed in angle brackets indicate that another syntactic construct is defined elsewhere.
- UPPERCASE Words written entirely in uppercase characters generally have a special meaning. All reserved words and standard identifiers are written in uppercase. All identifiers exported from the pseudomodules *SYSTEM* and *STORAGE* are also uppercase. Other library identifiers that are uppercase are usually special; for example, the

file types *TEXT* and *FILE* exported from the modules
Texts and *Files* are uppercase.

Each entry in the Reference Directory contains a Class descriptor. This is a quick reference to describe which of the general categories listed below the entry belongs to. In some cases, an entry belongs to more than one general category, and is listed appropriately. The Class descriptor can take on the following values:

E = Turbo Modula-2 Extension

L = Library

R = Reserved Word

S = Standard Identifier

ABS standard function S

Description	ABS returns the absolute value of the expression X.
Usage	$Y := \text{ABS}(X)$
	Argument X is of type REAL, INTEGER, LONGINT, or LONGREAL.
	Result Y must be the same type as X.
Comments	The absolute value is the positive value of a number; thus, if the function argument X is -12, the result is 12. A positive value is returned unchanged.
Example	<pre>CONST Mass = 120.0; VAR XPar,AbsX:INTEGER Acceleration,AbsForce:REAL BEGIN AbsX := ABS(XPar); AbsForce := ABS(Mass * Acceleration);</pre>

ADDRESS (SYSTEM) L

Description	ADDRESS is a type that represents memory addresses.
Declaration	TYPE ADDRESS = POINTER TO WORD ;
Usage	VAR a:ADDRESS ADDRESS must be imported from the pseudomodule SYSTEM.

Comments	<p>ADDRESS can be considered as POINTER TO WORD; thus, dereferencing an ADDRESS will yield a WORD.</p> <p>Addresses can be computed, and it is possible to add and subtract cardinals from addresses. The result of such a computation is again of type ADDRESS.</p> <p>ADDRESS is compatible with CARDINAL and with every pointer type. In particular, if a formal parameter is of type ADDRESS, the corresponding actual parameter may be of any pointer type.</p> <p>This is a low-level concept; its use will reduce the portability of the code that employs it.</p>
See Also	ADR SYSTEM WORD
Example	<pre> MODULE WhereAreXandY; FROM SYSTEM IMPORT ADDRESS, ADR; FROM InOut IMPORT WriteHex; TYPE atype = RECORD A,B: REAL END ; VAR adrx, adry : ADDRESS; x, y : atype; BEGIN adrx := ADR(x); adry := ADR(y); WRITE('x is at !'); WriteHex(adrx); WRITELN; WRITE('y is at !'); WriteHex(adry); WRITELN; END WhereAreXandY;</pre>

ADR procedure (SYSTEM) L

Description *ADR* returns the address of a structured variable.

Declaration **PROCEDURE ADR(VAR v: AnyType): ADDRESS;**

Usage *y* := *ADR(x)*;

ADR must be imported from the pseudomodule *SYSTEM*.

x may be a variable of any structured type and *y* must be a pointer or a cardinal variable.

Comments *ADR* yields a value of type *ADDRESS*, which is a 2-byte value indicating the address of the argument variable.

This is a low-level routine; its use will reduce the portability of the code that employs it.

See Also **ADDRESS**
SYSTEM

Example

```

MODULE WhereAreXandY;
FROM SYSTEM IMPORT ADDRESS, ADR;
FROM InOut IMPORT WriteHex;
TYPE
    atype = RECORD
        A,B: REAL
    END ;
VAR
    adrx, adry : ADDRESS;
    x, y      : atype;
BEGIN
    adrx := ADR(x);
    adry := ADR(y);
    WRITE('x is at '); WriteHex(adrx); WRITELN;
    WRITE('y is at '); WriteHex(adry); WRITELN;
END WhereAreXandY;

```

ALLOCATE procedure (STORAGE) L

Description	<i>ALLOCATE</i> allocates an area from the heap.
Declaration	PROCEDURE ALLOCATE (VAR a: ADDRESS; size: CARDINAL);
Usage	ALLOCATE(a, size)

ALLOCATE must be imported from the pseudomodule **STORAGE**.

a must be compatible with type **ADDRESS** (any pointer).

size must be of type **CARDINAL**.

Comments A space of *size* bytes is allocated on the heap, and the pointer *a* is made to point to the beginning of the allocated area. If the memory requested is not available the exception **OUTOFMEMORY** is raised. You may check available memory with **FREEMEM**.

A safer and more comfortable way to allocate dynamic variables is by using the procedure **NEW**, although you must still import **ALLOCATE**.

The counterpart routine to release this memory is called **DEALLOCATE**.

Note that the size of an **ADDRESS** is only 2 bytes; thus **ADDRESS** is compatible with **WORD**.

See Also **DEALLOCATE**
DISPOSE
FREEMEM
OUTOFMEMORY
NEW
STORAGE

Example Allocate 100 bytes off of the heap:

```
VAR  
  work: ADDRESS;  
BEGIN  
  ALLOCATE(work, 100);
```

AND R

Description AND is a binary logical operator resulting in a Boolean expression.

Usage IF Raining AND NoUmbrella THEN Run ELSE Walk END ;

AND is a reserved word.

Raining and *NoUmbrella* are Boolean expressions.

Comments AND is used to concatenate two logical expressions, where both subexpressions must be true for the entire expression to be true.

Note that the second subexpression in an AND expression is only evaluated if the first subexpression is true. This is called short-circuit evaluation.

See Also OR
 NOT

Example WHILE statement using an AND to test bounds condition prior to indexing an array (short-circuit evaluation):

```
VAR  
  s: ARRAY [1..maxitems] OF Info;  
BEGIN  
  WHILE i#0 AND s[i] > amount DO  
    (* process s[i] *)  
  END ;  
END ;
```

APPEND procedure (STRINGS) L

Description *Append* appends a substring to a string.

Declaration **PROCEDURE** Append(**substr: ARRAY OF CHAR;**
VAR str: ARRAY OF CHAR);

Usage **Append(substr, str)**

Append must be imported from the library module *Strings*.

str and *substr* can be arrays of character with any bounds.

Comments A string's »end« is marked by a zero byte; thus *Append* will tack on *substr* at the first zero byte in *str*.

If *str* has not been declared to be long enough to hold the two together, the exception *StringError* is raised.

See Also **Delete**
Insert
StringError
Strings

Example Write out 'Hello world !':

```
MODULE Main;
FROM Strings IMPORT Append;
VAR
  str, substr: ARRAY [0..80] OF CHAR;
BEGIN
  str    := 'Hello';
  substr := ' world !';
  Append(substr, str);
  WRITELN(str);
END main.
```

Arctan function (MATHLIB, LONGMATH) L

Description *Arctan* returns the arc tangent of *X*, with the result expressed in radians.

Declaration **PROCEDURE** Arctan (*x*: REAL): REAL;

and

PROCEDURE Arctan (*x*: LONGREAL): LONGREAL;

Usage *Y*:= Arctan(*X*)

Arctan must be imported from the library module *MathLib* for single-precision reals and from *LongMath* for double-precision reals.

Argument *X* and the target *Y* of the result must be the same type, either both REAL or both LONGREAL.

Comments The result is within the range -Pi/2 to Pi/2 radians. If you want the result expressed in degrees instead of radians, use the following expression:

Degrees:= Arctan(*X*) * 180.0/3.141592

The inverse of the *Arctan* function is *tangent*, which you can calculate with the *Sin* and *Cos* functions:

Tan:= SIn(*X*)/Cos(*X*).

See Also Mathlib
LongMath
Sin
Cos

Example **MODULE** ArcTanTest;
FROM MathLib **IMPORT** Arctan;
CONST
 RadsToDegs = 180.0/3.141592;

```
VAR
  Opposite, Adjacent, Angle: REAL;

BEGIN
  Opposite := 1.0;
  Adjacent := 1.0;
  Angle := Arctan(Opposite/Adjacent) * RadsToDegs;
  WRITELN('Angle=',Angle,' and it better be
  45 degrees !!');
END ArcTanTest.
```

ArgumentError exception (MATHLIB) L/E

Description	ArgumentError is raised by the module <i>MathLib</i> if an invalid parameter is passed to one of the procedures exported by <i>MathLib</i> .
Declaration	EXCEPTION ArgumentError;
Usage	EXCEPTION ArgumentError: WRITELN('Argument error');
	ArgumentError must be imported from the module <i>MathLib</i> .
Comments	This exception is raised if the parameter to a procedure is out of the valid range for that function.
See Also	MathLib LongMath

Example

```
MODULE MathError;
FROM MathLib IMPORT Exp, ArgumentError;
VAR x: REAL;
BEGIN
  x := Exp(99.0);
EXCEPTION
  ArgumentError:
    WRITELN('argument to Exp must be less than 87.4');
END MathError.
```

ARRAY standard type R

Description

An array consists of a fixed number of components all of the same type, the *base type*. Each component can be individually accessed by indices whose values are of the *index type*. Arrays can also be treated as a whole when being assigned or passed as parameters.

Usage

```
a = ARRAY IndexType OF ComponentType
```

The index type can be BOOLEAN, CHAR, subrange, or enumeration types. The *ComponentType* can be of any type.

Multidimensional arrays are expressed with additional index types separated with commas.

Comments

You can use an array to handle several related items of identical type; for example, an array of soups of the day could be

```
Day = (Sun,Mon,Tue,Wed,Thu,Fri,Sat);
Soup = (Minestrone, CreamOfAsparagus,
        VegetableBeef, SplitPea, ChickenNoodle,
        Chowder, Bouillabaisse);
```

```
SoupOfDay = ARRAY Day OF Soup;
```

Each component in an array has an index with a value defined by the index type. Thus, in the preceding example, a component is specified with an index of type *Day*; for example, *SoupOfDay[Mon]*, *SoupOfDay[Sat]*, and so on.

A component can take values defined by the component type. In the example, an element of *SoupOfDay* can take any value from *Minestrone* to *Bouillabaisse*.

A multidimensional array is simply an array with an array as a component type. For example:

```
a2=ARRAY IndexType1 OF ARRAY IndexType2 OF CARDINAL
```

This can also be rewritten as

```
a2 = ARRAY IndexType1, IndexType2 OF CARDINAL
```

The components are accessed as either

```
MultiDim[Index1][Index2]
```

or

```
MultiDim[Index1, Index2]
```

No operators other than simple assignment are directly applicable to arrays, but any operator applicable to the component type can operate on a component. Thus, if *Numbers* is an array of REAL components, then

```
Numbers[X]:= Numbers[Y] * Numbers[Z]
```

Simple assignment is possible for identically structured arrays. For example, if *First* and *Second* are arrays of the same type, the following assignment makes each component of *First* equal to the corresponding component in *Second*:

```
First:= Second
```

Example

```
MODULE Arrays;
TYPE
  Alphabet = ['A'..'Z'];
  Days = (Sun,Mon,Tue,Wed,Thu,Fri,Sat);
  Shift = RECORD
    Day : [0..8];
    Swing: [0..8];
    Night: [0..8]
  END ;
VAR
  HoursSun : ARRAY Days OF REAL;
  Character : ARRAY [1..10] OF Alphabet;
  CharType : ARRAY CHAR OF (Letter,Digit);
  WorkRecord: ARRAY Day,[1..100] OF Shift;
BEGIN
  HoursSun[Tue] := 12.6;
  HoursSun[Mon] := HoursSun[Sat] + HoursSun[Sun];
  Character[3]:= "R";
  Character[1]:= Character[3];
  WorkRecord[Sun,1] := 0;
END Arrays.
```

available variable (TERMINAL) L

Description	<i>available</i> holds the currently installed terminal operation, which may not be available on all terminals.
Declaration	TYPE SpecialOps = (clearEol, insertDelete, highlight-normal); OpSet : SET OF SpecialOps; VAR available : OpSet;
Usage	IF clearEol IN available THEN ClearToEOL END ;
	<i>available</i> is a predefined variable declared as type <i>OpSet</i> , which must be imported from the module <i>Terminal</i> .
Comments	This variable is useful for writing terminal-independent programs. The value of <i>available</i> is set by the installation program INSTM2.
See Also	SpecialOps OpSet Terminal

Example

```

MODULE TermStuff;
FROM Terminal IMPORT available, SpecialOps,
          ClearToEOL,
          InsertLine, numRows, numCols;

VAR
  I, CurRow: CARDINAL;
  ScreenLine: ARRAY [0..23] OF ARRAY [0..79] OF
    CHAR;

BEGIN
  IF insertDelete IN available THEN InsertLine
  ELSE
    FOR I := 1 TO numCols DO WRITE(' ') END; WRITELN;
    FOR I := CurRow TO numRows DO
      WRITELN(ScreenLine[I])
    END;
  END;
END TermStuff.
```

Awaited procedure (PROCESSES) L

Description	<i>Awaited</i> tests to see if any processes are waiting for a certain signal.
Declaration	PROCEDURE Awaited(s: SIGNAL): BOOLEAN;
Usage	flag := Awaited(siggy)
	<i>Awaited</i> must be imported from the library module <i>Processes</i> . <i>flag</i> must be of type BOOLEAN and <i>siggy</i> must be of type SIGNAL imported from <i>Processes</i> .
Comments	<i>Awaited</i> only makes sense when used with the related routines <i>WAIT</i> and <i>SEND</i> .
	<i>Awaited</i> will not start up the waiting processes, it only lets the program know if processes are waiting for a certain signal. Presumably, if <i>Awaited</i> returns TRUE , then the signal will be sent with the <i>SEND</i> command.

See Also	Init Processes SEND StartProcess WAIT
Example	IF Awaited(Printer) THEN SEND(PrinterAvailable) END

BADOVERLAY exception (SYSTEM) L/E

Description	BADOVERLAY is raised by the module <i>SYSTEM</i> when an overlay file cannot be read.
Declaration	EXCEPTION BADOVERLAY;
Usage	EXCEPTION BADOVERLAY: WRITELN('Bad overlay'); END
	BADOVERLAY must be imported from the module <i>SYSTEM</i> .
Comments	This exception occurs only when the linker is used to produce overlays.
See Also	SYSTEM OUTOFMEMORY OVERFLOW REALOVERFLOW

Example

```

BadOverlays;
FROM Overlayed IMPORT Proc1;
FROM SYSTEM IMPORT BADOVERLAY;
BEGIN
  Proc1;
EXCEPTION
  BADOVERLAY:
    WRITELN('Could not load overlayed module.');
END BadOverlays.
```

()OS procedure (SYSTEM) L

Description BDOS invokes a BDOS (CP/M operating system) function.

Declaration PROCEDURE BDOS(n, w: WORD);

Usage BDOS(n,w)

BDOS must be imported from the pseudomodule *SYSTEM*.

n and *w* must be compatible with type *WORD* (which can be an integer, character, enumeration, and so on).

Comments The BDOS function *n* will be invoked; *w* is a parameter to the operating system, inserted in register DE.

If a function returns a result in the accumulator *A*, it can be extracted out of the variable *IORESULT*.

If a result is returned in the register HL, it can be extracted from *HLRESULT*.

They are declared in *SYSTEM* as

VAR

HLRESULT, IORESULT: CARDINAL.

See Also BIOS
 SYSTEM
 IORESULT
 HLRESULT
 Also see the BDOS section of your CP/M manual.

Example Read a character from the console:

```
BDOS(1,0); ch:=CHR(IORESULT);
```

BEGIN R

Description BEGIN is a key word that divides the declaration part of a module or procedure from the statement part of the same.

Usage BEGIN

Comments A BEGIN is used to mark the beginning of the code part of a procedure or module.

See Also END

Example MODULE Begin;
 BEGIN
 IF TRUE THEN
 WRITE('Only ');
 WRITELN('one BEGIN')
 END ;
 END Begin.

BIOS procedure (SYSTEM) L

Description BIOS results in a call to the specified BIOS jump vector.

Declaration PROCEDURE BIOS(n, w: WORD);

Usage	<code>BIOS(n,w)</code>
	BIOS must be imported from the pseudomodule SYSTEM.
	<i>n</i> and <i>w</i> must be compatible with type WORD (which can be an integer, character, enumeration, and so on).
Comments	The given <i>n</i> indicates which vector is called: 0=WARMBOOT, 1=CONSTAT, 2=CONIN, and so on. If a function returns a result in the accumulator <i>A</i> , it can be extracted out of the variable <i>IORESULT</i> . If a result is returned in the register HL, it can be extracted from <i>HLRESULT</i> . <i>HLRESULT</i> and <i>IORESULT</i> are declared in SYSTEM as cardinals.
See Also	<code>BDOS</code> <code>SYSTEM</code> <code>HLRESULT</code> <code>IORESULT</code>
Example	Read a character directly from the console: <code>BIOS(2,0); ch := CHR(IORESULT);</code>

BITSET standard type S

Description BITSET is a variable type that can assume as a value a set of up to 16 integers with values between 0 and 15.

Usage As a constant set:

```
[BITSET] "{" [element {,element}] "}"
```

element may be a constant, subrange, or variable of type CARDINAL with a value between 0 and 15. The standard identifier BITSET is optional in constant sets.

The standard identifier is used mainly in variable declarations, as in the following:

VAR

```
mask: BITSET;
```

Comments BITSET is a standard predefined set type. All the operators you can apply to other set types can be used on BITSET-type variables (see »Operators« in Chapter 3).

You denote BITSET constants in the same way as other set types:

```
{0..9,12,14} [ ]      BITSET[6,8,11].
```

A BITSET value requires 2 bytes for storage in memory.

Example

```

MODULE BitSets;
CONST
  Mask = {5..7,12};
VAR
  Bits: BITSET;

BEGIN
  Bits:= {1,3,6..13};
  Bits:= Mask * Bits;    (* AND the bits *)
  (* Bits is now {6,7,12} *)

  Excl(Bits,7);          (* Exclude bit 7 *)
  (* Bits is now {6,12} *)
END BitSets.
```

BOOLEAN standard type S

Description BOOLEAN is a variable type that can assume the logical truth values FALSE and TRUE.

Comments You can use BOOLEAN variables and expressions when you want to calculate a statement that will have the answer TRUE or FALSE; for example, the statement »Today is Tuesday and it is raining« gives a BOOLEAN result.

There are two groups of operators you can apply to obtain BOOLEAN expressions: relational operators and logical operators.

Relational operators compare different variables, constants, and so on, and give a BOOLEAN result. For example, the expression:

(Altitude > 2000.0)

compares the REAL variable *Altitude* with the value 2000.0 to get a TRUE or FALSE result. Thus, relational operators can compare any items of similar type.

On the other hand, logical operators are only applicable to BOOLEAN variables and expressions. For example:

`SwitchOn AND (Altitude > 2000.0)`

has a logical AND operator and is TRUE only if both parts of the expression are TRUE.

A BOOLEAN value requires 1 byte for storage in arrays and 2 bytes as stand-alone variables or fields in a record.

Example

```
MODULE Booleans;
VAR
  PowerAvailable, SwitchOn, PowerOn: BOOLEAN;
  Temperature:REAL;
BEGIN
  PowerAvailable:= SwitchOn AND PowerOn;
  IF PowerAvailable AND (Temperature > 90.0) THEN
    WRITELN('Turn on air conditioner ');
  END
END Booleans.
```

BusyRead procedure (Terminal) L

Description *BusyRead* tests to see if a character has been typed at the keyboard.

Declaration `PROCEDURE BusyRead(VAR ch: CHAR);`

Usage `BusyRead(ch)`

BusyRead must be imported from the system library *Terminal*.

ch must be of type CHAR.

Comments *BusyRead*, together with *ReadAgain*, provides much the same service as Turbo Pascal's *KeyPressed*; however, *BusyRead* and *ReadAgain* are more flexible.

If a character is typed at the keyboard, *BusyRead* will return it in *ch*. If nothing is typed, *ch* is returned as 0C (a zero byte). For example:

```
PROCEDURE KeyPressed( ): BOOLEAN;
VAR ch: CHAR;
BEGIN
  BusyRead(ch);
  IF ch#0C THEN
    ReadAgain; RETURN TRUE
  ELSE
    RETURN FALSE;
  END
END KeyPressed;
```

The procedure emulates the *KeyPressed* function. To use the character again in another portion of the code we can use *ReadAgain*.

Notice that using the function *KeyPressed* requires two steps: (1) determine if there is a character and (2) if there is, read the character. This can be a problem in Modula-2 because of its interrupt handlers.

Suppose a program tests (with *KeyPressed*) whether a character is at the keyboard, and before it can read the character it is interrupted. If the interrupting routine also reads the keyboard, then when control returns to the interrupted program the character will not be available. This problem is solved if the procedure that reads a character simply returns a flag when no characters are available.

See Also *ReadAgain*
 Terminal

Example Display the ordinal value of a key that is pressed:

```
MODULE BusyReads;
FROM Terminal IMPORT BusyRead;
VAR ch: CHAR;
BEGIN
  REPEAT
    WRITELN('Press a key (''Q'' to quit): ');
    REPEAT
      BusyRead(ch);
    UNTIL ch # 0C;
    WRITELN('The ordinal value is ',ORD(ch));
  UNTIL CAP(ch)='Q'
END BusyReads.
```

BYTE type (SYSTEM) L

Description *BYTE* is a low-level type used to represent the scalar types that take up 1 byte of memory.

Usage **VAR**
 b :BYTE;

BYTE must be imported from the pseudomodule *SYSTEM*.

Comments *BYTE* is declared as a subrange of *WORD* and is therefore assignment-compatible with all scalar types.

BYTE is compatible with *BOOLEAN*, *CHAR*, and enumeration types that have 256 elements or less.

BYTE is compatible with *WORD*.

Only assignment operations are allowed on variables of type *BYTE*. A stand-alone variable of type *BYTE* takes 2 bytes of storage (also fields in a record). A *BYTE* in an array is packed into 1 byte of storage.

Example **PROCEDURE** p1(b: BYTE);

Call procedure (Loader) L

Description *Call* allocates and executes an overlay in the form of a module.

Declaration **PROCEDURE** Call(modName: **ARRAY OF CHAR**);

Usage Call(modname)

Call must be imported from the library module *Loader*.

modname must be the character array containing the name of the overlay file.

Comments If *modname* does not have an extension, .MCD will be appended.

The object module specified by *modname* will be loaded into memory and begin executing. In this sense, it is not unlike a procedure call, but it is much slower due to disk access. Also, when execution is completed, the memory needed for object code and data is released again.

Communication between the calling program and the overlaid module is only possible through variables declared in the calling program and imported by the called overlay.

See Also firstDrive

Loader

Also see **linker described in Chapter 10**.

Example Load and execute the module *OVER.MCD*:

Call("OVER")

CAP standard function S

Description *CAP* returns the capital of an alphabetic character.

Usage *Ch* := CAP(*Ch*)

Argument *Ch* (and result) is of type CHAR.

Comments You can use this function if you want to change lowercase letters to uppercase letters. This is useful for testing keyboard input.

Note that if the argument has no uppercase equivalent, the argument is returned unchanged. In other words, spaces and punctuation characters will remain unaltered.

Example **MODULE CAPs;**
VAR ch: CHAR;
BEGIN
LOOP
BusyRead(ch);
IF CAP(ch) = 'Q' THEN EXIT END
END
END CAPs.

CAPS procedure (Strings) S

Description *CAPS* makes an entire string variable uppercase.

Declaration **PROCEDURE CAPS(VAR str: ARRAY OF CHAR);**

Usage **CAPS(str)**

CAPS must be imported from the library module *Strings*.

str must be of type ARRAY OF CHAR.

Comments	<i>CAPS</i> saves you from typing the following code anywhere you don't want to distinguish between uppercase and lowercase:
	<pre> FOR i := 0 TO Length(str)-1 DO str[i] := CAP(str[i]); END ; </pre>
See Also	CAP Length
Example	Capitalize the string »Hello«:
	<pre> str := "Hello"; CAPS(str); (* str now has the value "HELLO" *) </pre>

CARD standard function S

Description	<i>CARD</i> converts a numeric argument to type CARDINAL.
Usage	$Y := \text{CARD}(X)$
	Argument <i>X</i> can be of INTEGER, CARDINAL, REAL, LONGINT, or LONGREAL.
	The argument must be in the range 0 to 65535.
	Result <i>Y</i> is of type CARDINAL.
Comments	 <ul style="list-style-type: none"> When you use this function with a REAL argument, the value is truncated (the fractional part removed); thus, value 12.35 will become 12. <i>CARD</i> is not applicable to enumeration or BOOLEAN- and CHAR- type arguments. Use <i>ORD</i> with these types. If the argument cannot be converted (that is, to negative), the runtime error BoundsError is raised.

Example

```
MODULE CARDConversion;
VAR
  r: REAL;
  c: CARDINAL;
BEGIN
  r := 3.14;
  c := CARD(r);
  (* c now has the value 3 *)
END CARDConversion.
```

CARDINAL standard type S**Description**

CARDINAL is a standard type with variables that can assume the values between 0 and 65535. In addition, CARDINAL may be used as a type-transfer function.

Usage

VAR c: CARDINAL;

or

c := CARDINAL(~1)

Comments

You can use CARDINAL variables whenever you know that possible values are limited to positive whole numbers. For example, if you are counting the number of people in a room you know the count can only be a positive whole number. If your program suddenly produces a negative count, an error message will be written.

The CARD standard function will convert other simple data types into CARDINAL. The use of CARDINAL as a type transfer function should be considered a low-level facility that may not be portable.

A CARDINAL value requires 2 bytes for storage.

Example

```

MODULE CARDINALS;
CONST
  SearchValue = 3.65;
  N = 20;
VAR
  DataArray: ARRAY [0..N] OF REAL;
  I, Count : CARDINAL;
BEGIN
  FOR I:= 0 TO N DO
    IF (DataArray[I] = SearchValue) THEN
      Count:= Count + 1
    END
  END
END CARDINALS.
```

CardToStr procedure (Convert) L

Description

CardToStr converts a CARDINAL to a string.

Declaration

```

PROCEDURE CardToStr(c: CARDINAL;
                      VAR s: ARRAY OF CHAR);
```

Usage

`CardToStr(card, string)`

CardToStr must be imported from the library module *Convert*.

card must be of type CARDINAL.

string can be any array of characters.

Comments

The number is placed into the string starting at the right-most index. The left end of the string is padded out with blanks. If the number does not fit into the given string, the exception *TooLarge* is raised.

See Also

Convert
IntToStr
RealToStr
StrToCard

Example

```
MODULE xcts;
FROM Convert IMPORT CardToStr;
VAR
  card: CARDINAL;
  string: ARRAY [0..4] OF CHAR;
BEGIN
  card := 12;
  CardToStr(card, string);
  WRITE('"', string '"'); (* " 12" is output *)
END xcts.
```

CASE statement R

Description CASE selects a different statement sequence depending on the value resulting from an expression.

Usage

```
CASE Selector OF
  Case {; Case}
  [ELSE Statement sequence]
END
```

Selector is a variable or expression of type CARDINAL, INTEGER, BOOLEAN, CHAR, enumeration, or subrange.

Case = [Case label: Statement sequence]

Statement sequence = Statement {; Statement}

Each case label can be either a single constant value or a range of constant values of the same type as the selector, separated by commas.

The ELSE part is executed if it is present and there are no matching values in the CASE label lists.

Comments

The selector, whether it is an expression or a variable, gives a value that causes the selection and execution of the statement sequence containing that value in its label. A value can appear in only one label.

Example

```

CASE X + Y OF
| 0..5 : Z := 0; Y:= R * X
| 6..8 : Z := X + Y - 5
| 9..10: Z := 8; Y:= R * X
END ;

CASE Day OF
Sun:           Message:= 'Closed all day';
                OpenHours:= 0
Mon..Wed,Fri: Message:= 'Normal hours';
                OpenHours:= 8
Thu,Sat:       Message:= 'Half day';
                OpenHours:= 4
END

```

You can define an empty case if necessary. For example, suppose in the last example you have no statement sequence for the value *Sun*:

```

CASE Day OF
Sun:           |
Mon..Wed,Fri: Message:= 'Normal hours';
                OpenHours:= 8
Thu,Sat:       Message:= 'Half day';
                OpenHours:= 4
END

```

Notice that the case of *Sun* explicitly states that no action is performed. When you have code to execute for every value except those in the CASE lists, then you can use an ELSE:

```

CASE Day OF
| Sat : WRITELN('Go for a picnic');
| Sun : WRITELN('Go shopping');
ELSE
    WRITELN('Go to work')
END ;

```

CHAR standard type S

Description CHAR is a standard type with variables that can assume a character as a value. In addition, CHAR may be used as a type-transfer function.

Usage

```
VAR  
  ch: CHAR;  
AND  
  ch := CHAR(27);
```

Comments You can use CHAR-type variables and constants wherever you want to handle text.

A CHAR variable can take any ASCII character as a value (which are ordered by their ASCII value); for example, the ordinal of character *A* is equal to the ASCII value 65, while the character *B* has an ordinal of 66.

Constants of type CHAR are denoted by a character enclosed within single or double quotation marks, or are octal numbers within the range 0 to 377 followed by a C; for example, »9«, 'D', 0C.

You cannot apply arithmetic operators directly on CHAR-type variables, but you can convert the value to CARDINAL and vice versa. To convert a character to a CARDINAL type, you use the standard function *ORD*. To convert a CARDINAL to a CHAR, use the standard function *CHR*.

Turbo Modula-2 permits the comparison and assignment of strings of characters, and the module *Strings* provides several procedures for string handling.

A CHAR variable occupies 2 bytes in memory unless it is declared as an element of an array, in which case it occupies 1 byte in memory.

See Also **BYTE**
 CHR
 ORD
 VAL

Example **MODULE Characters;**
 ch: CHAR;
 BEGIN
 LOOP
 READ(ch);
 IF CAP(ch) = 'Q' **THEN EXIT END**
 IF ch < ' ' **THEN** WRITE('.')
 ELSIF ch > CHR(127) **THEN** WRITE('!!')
 ELSIF (ch >= 'a') **AND** (ch <= 'f') **THEN**
 WRITE(CAP(ch))
 ELSE WRITE(ch)
 END
 END
 END Characters.

CHR standard function S

Description *CHR* returns the character with ordinal *X*.

Usage Ch := CHR(X)

Argument *X* is of type INTEGER or CARDINAL.

Result *Ch* is of type CHAR.

The argument must be in the range 0 to 255.

Comments You can use the *CHR* standard function to convert from an ASCII value in the range 0 to 255 to the corresponding character.

Note: The characters *A* to *Z* have ASCII values 65 to 90; characters *a* to *z* have ASCII values 97 to 122; numerals 0 to 9 have ASCII values 48 to 57.

Converting a character to ASCII is performed by the standard function *ORD*.

(See Also CHAR
ORD

Example

```
PROCEDURE NumToChar(N: CARDINAL): CHAR;
BEGIN
  IF N < 10 THEN RETURN CHR(N + 48)
  ELSE RETURN CHR(48)
  END
END NumToChar;
```

ClearEol enumerated value (terminal) L

Description *clearEol* is the first value of the enumerated type *Terminal.SpecialOps*.

Declaration SpecialOps = (clearEol, insertDelete, highlightNormal);

Usage IF clearEol IN available THEN ClearToEOL END ;

SpecialOps must be imported from the library module *Terminal*.

Comments	If this value is a member of the set <i>Terminal.available</i> , then the program may use the procedure <i>ClearToEOL</i> to clear to the end of the line.
	To use this identifier, include <i>SpecialOps</i> in your import list. The identifiers of each of <i>SpecialOps</i> 'values become visible automatically.
See Also	<i>ClearToEOL</i> <i>available</i> <i>insertDelete</i> <i>highlightNormal</i> <i>Terminal</i>
Example	Clear the right half of the screen:

```

MODULE CheckAvailableOperations;
FROM Terminal IMPORT
    available,SpecialOps,ClearEOL,
    GotoXY;
VAR row: CARDINAL
BEGIN
  FOR row := 0 TO 24 DO
    GotoXY(40,row);
    IF ClearEol IN available THEN
      ELSE ClearToEOL;
      ManualClearEOL;
    END ;
  END ;
  GotoXY(0,0);
END CheckAvailableOperations.

```

ClearScreen procedure (Terminal) L

Description	<i>ClearScreen</i> clears the screen and homes the cursor.
Declaration	PROCEDURE ClearScreen;

Usage	ClearScreen
	<i>ClearScreen</i> must be imported from the library module <i>Terminal</i> .
Comments	This function will work only if the Turbo Modula-2 system has been installed on your terminal.
See Also	ClearToEOL GotoXY Highlight Terminal
Example	Clear the screen: MODULE Clear; FROM Terminal IMPORT ClearScreen; BEGIN ClearScreen; END Clear.

ClearToEOL procedure (Terminal) L

Description	<i>ClearToEOL</i> clears the screen from the current cursor position until the end of the line.
Declaration	PROCEDURE ClearToEOL;
Usage	ClearToEOL
	<i>ClearToEOL</i> must be imported from the library module <i>Terminal</i> .
Comments	The cursor position is undefined after a call to this routine; reposition it with a call to <i>GotoXY</i> .

See Also	ClearScreen GotoXY Highlight Terminal
Example	Clear the right half of the screen:

```
MODULE ClearHalfScreen;
FROM Terminal IMPORT ClearEOL, GotoXY;
VAR row: CARDINAL
BEGIN
  FOR row := 0 TO 24 DO
    GotoXY(40, row);
    ClearToEOL;
  END ;
  GotoXY(0, 0);
END ClearHalfScreen.
```

Close procedure (Files) L

Description	<i>Close</i> disconnects a <i>FILE</i> variable from an external disk file and updates the disk directory.
Declaration	PROCEDURE Close (VAR f: FILE);
Usage	Close(f);
<i>Close</i> must be imported from the module <i>Files</i> .	
	<i>f</i> must be of type <i>FILE</i> imported from <i>Files</i> .
Comments	If output has been sent to <i>f</i> , closing is mandatory; otherwise, the file size in the directory and the trailer byte will not be updated and <i>f</i> 's data could be lost. Since closing a file frees up the memory used for data buffering, we advise closing a file in any case, even if it has been used only for input.

See Also	Files Open NoTrailer
Example	Perform minimal file handling (open and close a <i>FILE</i>):

```
MODULE OpenAndCloseFile;
FROM Files IMPORT FILE, Open, Close;
VAR oldFile : FILE;
BEGIN
  IF Open(oldFile,'FILE.DAT') THEN
    Close(oldFile)
  END
END OpenAndCloseFile.
```

CloseInput procedure (InOut) L

Description	<i>CloseInput</i> closes an input file and returns input to the terminal.
Declaration	PROCEDURE <i>CloseInput</i> ;
Usage	<i>CloseInput</i>
	<i>CloseInput</i> must be imported from the library module <i>InOut</i> .
Comments	These routines are provided to maintain compatibility with other Modula-2 implementations (the <i>Texts</i> module is easier to work with).
	<i>CloseInput</i> is equivalent to the call <i>CloseText(input)</i> , using the <i>CloseText</i> procedure from <i>Texts</i> .
See Also	OpenInput CloseOutput InOut

Example

```

MODULE Input;
FROM InOut IMPORT
  OpenInput, CloseInput, ReadString,
  WriteString, WriteLn;
VAR
  s: ARRAY [0..255] OF CHAR;
BEGIN
  (* Prompt user for filename *)
  (* If no extention given, use MOD *)
  OpenInput("MOD");
  ReadString(s);      (* Read string from file *)
  WriteString(s);     (* Display it *)
  WriteLn;
  CloseInput;         (* Sever the link to the file *)
  ReadString(s);      (* Input comes from console *)
  WriteString(s);     (* Display it *)
  WriteLn;
END Input.
```

CloseOutput procedure (InOut) L

Description	<i>CloseOutput</i> closes an output file and returns output to the terminal.
Declaration	PROCEDURE CloseOutput;
Usage	<i>CloseOutput</i>
	<i>CloseOutput</i> must be imported from the library module <i>InOut</i> .
Comments	These routines are provided to maintain compatibility with other Modula-2 implementations (the <i>Texts</i> module is easier to work with).
	<i>CloseOutput</i> results in the call <i>CloseText(output)</i> , using the <i>CloseText</i> procedure in <i>Texts</i> .

See Also **OpenOutput**
 OpenInput
 Inout

Example **MODULE Output;**
 FROM InOut IMPORT
 OpenOutput, CloseOutput, ReadString,
 WriteString, WriteLn;
VAR
 s: ARRAY [0..255] OF CHAR;
BEGIN
 WriteString('Enter output file and data:');
 WriteLn;
 OpenOutput("MOD");
 ReadString(s); (* Read data from console *)
 (* Place it in the file *)
 WriteString(s); WriteLn;
 CloseOutput; (* Sever the link to the file *)
 WriteString('The file has been closed.');
 WriteLn;
END Output.

CloseText procedure (*Texts*) L

Description *CloseText* disconnects a *TEXT* file variable from an external disk file.

Declaration **PROCEDURE CloseText(**VAR** t: TEXT);**

Usage *CloseText(t)*

CloseText must be imported from the module *Texts*.

t must be of type *TEXT* imported from *Texts*.

Comments If output has been sent to *t*, closing is mandatory; otherwise, the file size in the directory and the trailer byte will not be updated and *t*'s data could be lost.

Since closing a file frees up the memory used for data buffering, we advise closing a file in any case, even if it has been used only for input.

If the external medium is a disk file, the file is closed internally using the procedure *Close* from the module *Files*.

Closing *input* or *output* re-establishes the default connections (to CON:).

See Also OpenText
TEXT
Texts
NoTrailer

Example Perform minimal text-file handling (open and close a *TEXT*):

```
MODULE OpenAndCloseText;
FROM Texts IMPORT TEXT, OpenText, CloseText;
VAR oldFile : TEXT;
BEGIN
  IF OpenText(oldFile,'T.TXT') THEN
    CloseText( oldFile )
  END
END OpenAndCloseText.
```

CODE procedure (SYSTEM) L

Description *CODE* allows the user to include an assembler routine.

Declaration PROCEDURE CODE(s: ARRAY OF CHAR);

Usage `CODE("Any string literal");`

CODE must be imported from the pseudomodule **SYSTEM**.

The string must be a literal file name.

Comments A call to *CODE* must be inserted after a procedure heading; thus replacing the body of the procedure. No declaration part, BEGIN, or any other statements may be present in its body; however, the final END must be included.

Compilation of a *CODE* statement causes the >.COM<-File whose name is given as argument to be inserted into the object program. The length of the file must be in the first 2 bytes of the file.

A *CODE* procedure may use all Z80 registers without restoring them.

See Also **SYSTEM**

Example Include MOVE.COM written in assembler:

```
PROCEDURE MOVE(source: ADDRESS; n: CARDINAL);
  CODE( "MOVE" )           (* MOVE.COM is inserted *)
  END MOVE;
```

Col function (Texts)

Description *Col* returns the current column position of a text file.

Declaration `PROCEDURE Col(t: TEXT): CARDINAL;`

Usage	<code>currentCol := Col(t)</code>
	<i>Col</i> must be imported from the library module <i>Texts</i> .
	<i>currentCol</i> must be of type CARDINAL.
	<i>t</i> must be of type TEXT imported from <i>Texts</i> .
Comments	After a <i>WriteLn</i> or a <i>CreateText</i> , the column position is zero; it increases by one with every written character.
See Also	<i>SetCol</i> <i>Texts</i>
Example	Start a new line if the current column position of TEXT <i>t</i> is greater than 79. <code>IF Col(t) > 79 THEN WriteLn(t) END</code>

ComLine module L

Description	<i>ComLine</i> enables command line processing (arguments and redirection).
Declaration	DEFINITION MODULE ComLine; FROM Texts IMPORT Text;
	VAR commandLine: Text; inName, outName: ARRAY [0..19] OF CHAR; progName: ARRAY [0..7] OF CHAR;
	PROCEDURE RedirectInput; (* Input redirection as ordered on command line. *)
	PROCEDURE RedirectOutput; (* Output redirection as ordered on command line.*)

```
PROCEDURE PromptFor(Prompt: ARRAY OF CHAR;
                     VAR S: ARRAY OF CHAR);
(* Prompts for command line. *)
END ComLine.
```

The *TEXT* variable *commandLine* contains any arguments written after the command that started the program. It can be read like a normal *TEXT*.

inName contains the input medium, which is normally CON:.

outName contains the output medium, which is normally CON:.

progName contains the program name.

Comments

The command line consists of any characters or arguments that follow the command that started the program. The contents of the command line are accessed through the *TEXT* *commandLine*.

The command line can contain *redirection arguments*:

>*FileName* Redirects the standard text *Output* to *FileName*. The new medium *FileName* is stored in *outName*.

<*FileName* Redirects the standard text *Input* to *FileName*. The new medium *FileName* is stored in *inName*.

The redirection symbols and file names do not appear in the command line.

Note that if there is no redirection argument, *inName* and *outName* contain the name of the standard medium CON:.

See Also

Texts

Also refer to each identifier in the definition module.

Example

If you start a program by typing

MyProg Name

the *TEXT commandLine* will contain the word *Name*. You would access this like so:

```
READ(commandLine,InputFilename);
```

If you write a program to compare two disk files, you will want to specify the file to be compared on the command line.

By using *ComLine* procedures you can pass the file names directly to the program. Your program will include the following lines:

```
MODULE Compare;
FROM Files IMPORT FILE, Open,Close;
FROM ComLine IMPORT
  commandLine, RedirectInput, RedirectOutput;
VAR
  Filename1, Filename2: ARRAY [0..14] OF CHAR;
  f1,f2: FILE;
BEGIN
  RedirectInput; RedirectOutput;
  READ(commandLine,Filename2);
  IF Open(f2,Filename2) THEN
    READ(commandLine,Filename1);
    IF Open(f1,Filename1) THEN
      READLN(options);
      (* compare the files *)
      (* Write differences to standard output *)
    END;
  END ;
  Close(f1);
  Close(f2);
END Compare.
```

For example, you can start a program and define a command line by typing

```
Compare A:testfile B:testfile <B:Compare.Ops  
>A:Compared.dif
```

This line invokes the compare program with two files as arguments. Any input requested from the console is redirected from the keyboard to the file COMPAREOPS. The output of the program is redirected from the screen to the file COMPARED.DIF.

commandLine variable (ComLine) L

Description *commandLine* is Predefined text that contains arguments passed on the command line when a program is invoked.

Declaration **VAR**
 commandLine: TEXT;

TEXT is imported from the module *Texts*.

Usage READ(*commandLine*,*filename*);

filename is a string variable

Comments Treat *commandLine* like any *TEXT* variable. It contains anything typed on the command line after the program name, up to 80 characters long.

See Also PromptFor
 ComLine

Example

```
MODULE OffTheCommandLine;
FROM ComLine IMPORT commandLine;
VAR
  x,y: REAL;
BEGIN
  READ(commandLine,x,y);
  WRITELN(x:5:2,'*',y:5:2,' = ',x*y:5:2);
END OffTheCommandLine.
```

ConnectDriver procedure (Texts) L

Description	<i>ConnectDriver</i> allows the user to install her own input and output drivers.
Declaration	TYPE TextDriver = PROCEDUREE(TEXT, VAR CHAR); PROCEDURE ConnectDriver(VAR t: TEXT; p: TextDriver)
Usage	ConnectDriver(t, driver)
	<i>t</i> must be of type <i>TEXT</i> imported from the library module <i>Texts</i> .
	<i>Driver</i> must be compatible with the type <i>TextDriver</i> .
Comments	Allows users to install their own read and write character routines, which are then used by all I/O statements in the standard module <i>Texts</i> , including the extensions <i>READ</i> , <i>WRITELN</i> , and <i>WRITELN</i> .
See Also	Texts

Example Connect the user-defined driver to the text *t*:

```

MODULE AlwaysUpCase;
FROM Texts IMPORT TEXT, output, ConnectDriver;

PROCEDURE WriteUpcase(t: TEXT; VAR ch: CHAR);
BEGIN
  Write(output, CAP(ch));
END WriteUpcase;

VAR t: TEXT;
BEGIN
  ConnectDriver(t, WriteUpcase);
  WRITELN(t, 'This will be displayed all uppercase');
END AlwaysUpCase.
```

console (Texts) L

Description An auxiliary output *TEXT* attached to the console.

Declaration **VAR** console: TEXT;

Usage WRITELN(console, 'output to the console');

console must be imported from the library module *Texts* and can be used anywhere an output *TEXT* can be used.

Comments *console* is always open and, by default, is connected to the terminal. It is as if the call *CreateText(console, >CON:«)* has taken place before the execution of the main module.

console is useful if the *output* text stream has been redirected but you still want certain messages (prompts, errors, and so on) to go to the terminal. The *console* cannot be redirected.

See Also *output*
input
Texts

Example Write out an error message by passing redirection:

```
MODULE consoleOutput;
FROM Texts IMPORT console;
VAR recnum: CARDINAL;
BEGIN
  (* processing loop *)
  WRITE(console, 'Error at record number: ', recnum);
END consoleOutput.
```

CONST declaration R

Description The reserved word CONST precedes the declaration of constants.

Usage **CONST**
 Pi = 3.14159265;
 TwoPi = 2.0*Pi;

Comments Modula-2 allows the use of constants in the definition of further constants, like the *Pi* in *TwoPi*.

Example **MODULE** Constants;
CONST
 nThings = 100;
 LastThing = nThings-1;
TYPE
 ThingRange = [0..LastThing];
VAR
 Things : **ARRAY** ThingRange **OF** CHAR;
 which : [0..lastthing];
 thingy : CARDINAL;
BEGIN
 FOR thingy := 0 **TO** lastthing **DO**
 Process(things[which])
 END
END Constants.

Convert module L

Description *Convert* is a standard module that you can use when converting between numeric types and strings.

Declaration **DEFINITION MODULE** Convert;

```

PROCEDURE StrToInt      (VAR s: ARRAY OF CHAR;
                           VAR i: INTEGER) : BOOLEAN;
PROCEDURE StrToCard     (VAR s: ARRAY OF CHAR;
                           VAR c: CARDINAL) : BOOLEAN;
PROCEDURE StrToLong     (VAR s: ARRAY OF CHAR;
                           VAR l: LONGINT) : BOOLEAN;
PROCEDURE StrToReal     (VAR s: ARRAY OF CHAR;
                           VAR r: REAL) : BOOLEAN;
```



```

PROCEDURE IntToStr (i: INTEGER
                     VAR s: ARRAY OF CHAR);
PROCEDURE CardToStr (c: CARDINAL
                     VAR s: ARRAY OF CHAR);
PROCEDURE LongToStr (l: LONGINT
                     VAR s: ARRAY OF CHAR);
PROCEDURE RealToStr (r: REAL
                     VAR s: ARRAY OF CHAR;
                     digits: INTEGER);
```



```

EXCEPTION TooLarge;
```

```

END Convert.
```

Usage See individual identifiers

Comments The first four procedures convert strings to numeric types; the next four reverse the conversions.

The exception *TooLarge* is raised in the last four procedures if the number is too large to fit in the target string. (*TooLarge* is only raised by the second four procedures.)

See Also The module *Doubles* provides the same functions for double-precision variables (**LONGREAL**).

Copy procedure (Strings) L

Description	<i>Copy</i> copies a substring starting at a given position with given length of a string.
Declaration	PROCEDURE Copy(VAR str: ARRAY OF CHAR ; inx, len: CARDINAL ; VAR result: ARRAY OF CHAR);
Usage	Copy (string, from, length, resultstring) <i>Copy</i> must be imported from the library module Strings . <i>string</i> and <i>result</i> must be character array variables. <i>from</i> and <i>length</i> must be of type CARDINAL .
Comments	A string is terminated by a zero byte. If a zero byte is countered before <i>length</i> characters have been processed, copy will be prematurely interrupted.
See Also	Concat Delete Insert Strings
Example	Copy 'holes' into <i>holestring</i> : Copy ('Donut holes', 6, 5, holestring)

Cos function (MathLib, LongMath) L

Description *Cos* returns the cosine of *X*, where *X* is expressed in radians.

Declarations **PROCEDURE** Cos(*x*: REAL): REAL;
 PROCEDURE Cos(*x*: LONGREAL): LONGREAL;

Usage *Y*:= Cos(*X*)

Argument *X* and result *Y* must be of the same type--either both REAL or both LONGREAL.

Comments If you want the argument expressed in degrees instead of radians, use the following expression:

Result:= Cos(*X* * 3.141592/180.0)

If you need to calculate the inverse of cosine, you can use the *Arctan* function in the following expression:

ArcCos:= Arctan(*X*/Sqrt(1.0 - *X***X*)) - 3.141592/2.0

Example **MODULE** Cosine;
 FROM MathLib **IMPORT** Cos;
 CONST
 DegsToRads = 3.141592/180.0;
 VAR
 Degs,Radians,CosAngle:REAL
 BEGIN
 Radians:= Degs * DegsToRads;
 CosAngle:= Cos(Radians);
 END Cosine.

Create procedure (Files) L

Description *Create* creates a file on disk with a given name.

Declaration **PROCEDURE Create(VAR f: FILE; name: ARRAY OF CHAR);**

Usage *Create(f, filename)*

Create must be imported from the library module *Files*.

f must be declared as *FILE*.

FILE must be imported from the library module *Files*.

filename must be declared as **ARRAY OF CHAR**.

Comments *filename* must be a legal CP/M filename, and there must be room on the directory and the disk for a new file.

f can only be used for output after being opened with *Create*; use *Open* for reading a file.

If a file is already present, it will be overwritten, so use caution.

After writing to this file, close it to update the file size in the directory and the file trailer byte. Failure to do this will probably result in the loss of all written information.

See Also **Close**
 Files
 Open
 NoTrailer

Example Create an empty file on disk:

```
MODULE CreateAnEmptyFile;
FROM Files IMPORT FILE, Create, Close;
VAR
  newFile :FILE;
BEGIN
  Create(newFile, 'FILE.DAT');
  Close(newFile);
END CreateAnEmptyFile.
```

CreateText procedure (*Texts*) L

Description *CreateText* associates a *TEXT* with an external medium and opens it for writing.

Declaration PROCEDURE CreateText(VAR
t: TEXT; name: ARRAY OF CHAR);

Usage CreateText(*t*, *fname*)

CreateText must be imported from library module *Texts*.

t must be of type *TEXT* imported from *Texts*.

fname must be of type *ARRAY OF CHAR*.

Comments *fname* must be a legal CP/M file name or device, such as »MYFILE.DAT« or »CON:«.

t can only be used for output after being opened with *CreateText*; use *OpenText* for reading a *TEXT* file.

If a file is already present it will be overwritten, so use caution.

After writing to this file, close it to update the file size in the directory and the file trailer byte. Failure to do this will probably result in the loss of all written information.

See Also	CloseText OpenText Texts NoTrailer
Example	Create an empty text file on disk:
	<pre>MODULE CreateAnEmptyText; FROM Texts IMPORT TEXT, CreateText, CloseText; VAR newFile: TEXT; BEGIN CreateText(newFile, 'EMPTY.TXT'); CloseText(newFile); END CreateAnEmptyText.</pre>

DeadLock exception (Processes) L/E

Description	DeadLock is raised in the module <i>Processes</i> when all processes are waiting for a signal.
Declaration	EXCEPTION DeadLock;
Usage	EXCEPTION DeadLock: WRITELN('Deadly embrace: all processes terminated.');
Comments	DeadLock must be imported from the module <i>Processes</i> . When a dead lock occurs, your signal handling logic is failing. You can detect this problem by using this exception.
See Also	Processes SYSTEM

Example **MODULE** Signals;
 FROM Processes **IMPORT** Init, SIGNAL, WAIT, DeadLock;
 VAR s: SIGNAL;
 BEGIN
 Init(s); WAIT(s);
 EXCEPTION
 DeadLock: HALT;
 END Signals.

DEALLOCATE procedure (STORAGE) L

Description *DEALLOCATE* deallocates a block of memory on the heap that had been previously associated with a pointer variable.

Declaration **PROCEDURE** DEALLOCATE(**VAR**
 a: ADDRESS; *size*: CARDINAL);

Usage **DEALLOCATE(***a*, *size*)

DEALLOCATE must be imported from the library module **STORAGE**.

a must be compatible with type **ADDRESS** (any pointer).

size must be of type **CARDINAL**.

Comment You can only deallocate portions of memory that you have previously allocated. The variables *a* and *size* must be the same location and amount as when they were allocated.

The pointer *a* will be set to **NIL** and the deallocated area will be available for future allocation.

See Also ALLOCATE
 NEW
 DISPOSE
 STORAGE
 MARK
 RELEASE
 FREEMEM

Example Allocate and deallocate 100 bytes:

```
MODULE HeapMem;  
FROM STORAGE IMPORT ALLOCATE, DEALLOCATE;  
VAR  
  a: POINTER TO ARRAY [0..99] CHAR;  
BEGIN  
  ALLOCATE(a,100);  
  a^ := 'This string is on the heap';  
  DEALLOCATE(a, 100);  
  (* Now the string is gone. *)  
END HeapMem.
```

DEC standard procedure S

Description *DEC* returns a predecessor of argument *X*.

Usage DEC(*X*)

or

DEC(*X,N*)

Argument *X* may be of type INTEGER, CARDINAL, CHAR, BOOLEAN, or a user-defined scalar type.

The second form returns the *N*th predecessor of *X*.

Comments If you subtract 1 from a CARDINAL value, the result is the value prior to the original value. With DEC, you can do exactly the same thing with any ordinal type.

When decrementing a value, you may reach its *MIN* value; for example, the *MIN(CARDINAL)* is 0. You may choose whether or not this situation will generate an error by using the overflow compiler switch. This can be done globally or locally. Thus with overflow checking turned off, as in

```
VAR
  i: CARDINAL;
BEGIN
  i := MIN(CARDINAL);
  (*$0-*)
  DEC(i);
  (*$0+*)
  WRITELN(1);
END
```

the *i* would wraparound to *MAX(CARDINAL)* and no error will be generated. If overflow checking has been turned on, then the exception OVERFLOW will be raised. This can be trapped as follows:

```
MODULE dec;
FROM SYSTEM IMPORT OVERFLOW;
VAR
  i: CARDINAL;
BEGIN
  (*$0+*)
  i := MIN(CARDINAL);
  DEC(i);
EXCEPTION
  OVERFLOW: WRITELN('Trapped overflow error');
END dec.
```

See Also

INC

Example For example, suppose you have an enumeration variable:

```
VAR Energy: (Kinetic,Potential,Heat);
```

You can assign any of the three values to the variable *Energy*. Now suppose you want to decrement the value of *Energy* to a lower value; for example, to the value preceding the current value:

```
DEC(Energy)
```

You have performed a subtraction operation on the enumeration variable.

The procedure can operate on any ordinal type.

```
MODULE Decrement;
VAR
  i: INTEGER;
  state: BOOLEAN;
BEGIN
  i:= 0;
  DEC(i,2);
  (* now i= -2*)
  state:= TRUE;
  DEC(state);
  (* now state is FALSE*)
END Decrement.
```

DEFINITION declaration R

Description **DEFINITION** specifies to the compiler that a compilation unit is the definition part of a library module.

Usage **DEFINITION MODULE** myMod

Comments	<p>The definition module is the »visible« part of a library module that can be imported into other modules. In the definition module <i>myMod</i> you will find all the declarations necessary to use the facilities of <i>myMod</i>.</p> <p>Every definition module will have a corresponding implementation module that contains the hidden parts of that module. It will contain all executable code, as well as those definitions relevant only to the internal operation of the module.</p> <p>All identifiers mentioned in a definition module are exported. The identifiers are exported in qualified mode; although, in most cases, they will be dequalified using the FROM clause in the module that imports them.</p> <p>There is no code in a definition module. The compiler simply produces a symbol table file that can later be used to do type-checking.</p> <p>Definition modules usually reside in files with the extension .DEF. Compilation results in the creation of a file with the extension .SYM.</p> <p>The definition module cannot have the compilation unit in the same file; thus definition modules and implementation modules must be in separate files. The resulting .SYM and .MCD files can be combined, however, into a .LIB file using the Turbo Modula-2 librarian.</p>
See Also	IMPLEMENTATION MODULE

Example A simple user-definition module in file 'SIMPLE.DEF':

```
DEFINITION MODULE SimpleMath;  
  
PROCEDURE ADD(y,z :INTEGER) : INTEGER;  
  
END SimpleMath.
```

The corresponding implementation module in file 'SIMPLE.MOD':

```
IMPLEMENTATION MODULE SimpleMath;  
  
PROCEDURE ADD(y,z :INTEGER) : INTEGER;  
BEGIN  
    RETURN (y+z);  
END Add;  
  
END SimpleMath.
```

DELETE procedure (Files) L

Description *Delete* erases a file from disk.

Declaration **PROCEDURE** Delete(**VAR** f: FILE);

Usage *Delete(f)*

Delete must be imported from the library module *Files*.

f must be an open file of type *FILE* imported from *Files*.

Comments The file must have already been successfully opened.

The file is closed when it is deleted, thus its buffers are freed.

See Also

- Close
- Files
- Open
- Rename

Example Delete the file 'SCRATCH.SCR' from disk:

```
MODULE deleteFile;
FROM Files IMPORT FILE, Open, Delete;
VAR f: FILE;
BEGIN
  IF Open(f, 'SCRATCH.DAT') THEN Delete(f) END
END deleteFile.
```

DELETE procedure (Strings) L

Description *Delete* deletes a substring from a string.

Declaration PROCEDURE Delete(VAR str: ARRAY OF CHAR;
inx,len: CARDINAL);

Usage *Delete(str, from, length)*

Delete must be imported from the library module *Strings*.

from and *length* must be of type CARDINAL.

str must be a character array variable.

(**Comments** The counting starts from 0, so the user's strings should be declared starting from 0.

See Also

- Copy
- Insert
- Pos
- Strings

Example *Delete(»Lover«, 0, 1)* will yield the string *»over«*.

DeleteLine procedure (*Terminal*) L

Description *DeleteLine* deletes the cursor line completely and to fill the gap moves up all following lines by one.

Declaration **PROCEDURE** DeleteLine;

Usage DeleteLine;

DeleteLine must be imported from the library module *Terminal*.

Comments The current cursor line will be deleted. (This can be set by the procedure *GotoXY* in *Terminal*.) Since the cursor position on a screen is undefined after certain operations, the only way to determine the current cursor position is by keeping track of it yourself.

This function will work only if the Turbo Modula-2 system has been installed on your terminal.

A program can determine if the option has been installed by examining the set variable *Available* exported by *Terminal*. If *insertDelete* IN *Available* is TRUE, then these functions are available as installed on the current terminal.

See *GotoXY* for a map of the screen coordinates.

See Also *GotoXY*
 InsertLine
 Terminal

Example Delete the **twelfth line** from the top:

```
GotoXY(0,11);  
DeleteLine;
```

DeviceError exception (Files) L/E

Description DeviceError is raised by the *Files* module when a write operation on a file is impossible because the disk has a bad sector.

Declaration **EXCEPTION**
 DeviceError;

Usage **EXCEPTION**
 DeviceError: WRITELN('File error');
END

DeviceError must be imported from the module *Files*.

Comments Use the DeviceError exception to prevent your program from crashing when the program encounters a bad disk.

See Also EndError
StatusError
UseError
DiskFull

Example

```
MODULE WriteFile;
FROM Files IMPORT
    FILE, Create, WriteWord, Close, EndError,
    StatusError, UseError, DeviceError;
VAR f: FILE; c: CARDINAL;
BEGIN
    Create(f,'testfile.dat');
    WriteWord(f,c);
    Close(f)
EXCEPTION
    EndError: WRITELN('End of file reached.');
    Close(f)
    | StatusError: WRITELN('Error file not opened');
    | UseError : WRITELN('Disk not logged in.');
    | DeviceError: WRITELN('Error writing disk:
        bad sector');
END WriteFile.
```

DiskFull exception (Files) L/E

Description	DiskFull is raised by the <i>Files</i> module when a write or close operation on a file is impossible because the disk or directory is full.
Declaration	EXCEPTION DiskFull;
Usage	EXCEPTION DiskFull: WRITELN('File error'); END
	DiskFull must be imported from the module <i>Files</i> .
Comments	Use the DiskFull exception to prevent your program from crashing when the program encounters a full disk.
See Also	ErrorHandler StatusError UseError DeviceError

Example

```

MODULE WriteFile;
FROM Files IMPORT
    FILE, Create, WriteWord, Close, EndError,
    StatusError, UseError, DeviceError, DiskFull;
VAR f: FILE; c: CARDINAL;
BEGIN
    Create(f,'testfile.dat');
    WriteWord(f,c);
    Close(f)
EXCEPTION
    | EndError: WRITELN('End of file reached.');
    | Close(f)
    | StatusError: WRITELN('Error file not opened');
    | UseError : WRITELN('Disk not logged in.');
    | DeviceError: WRITELN('Error writing disk: bad
      sector');
    | DiskFull: WRITELN('This disk is FULL, buddy!');
END WriteFile.
```

DISPOSE standard procedure (*STORAGE*) S

Description *DISPOSE* deallocates a dynamic variable from the heap.

Usage *DISPOSE(p)*

or

DISPOSE(p, tag1, tag2, ... tagn)

To use this procedure, *DEALLOCATE* must be imported from the library module *STORAGE*, or it must be defined in some other way.

p must be a pointer to a variable.

tag1, ..., tagn are relevant if the type of the variable is a tagged variant record.

Comments	<p><i>DISPOSE</i> is essentially an abbreviated <i>DEALLOCATE</i> command. If, for example, <i>p</i> is of type <i>POINTER TO LONGREAL</i>, then <i>DISPOSE(p)</i> will be equivalent to <i>DEALLOCATE(p, TSIZE(LONGREAL))</i>.</p> <p><i>p</i> must point to a variable that has been previously allocated using either <i>ALLOCATE</i> or <i>NEW</i>.</p> <p>Care must be taken when using <i>DISPOSE</i> and <i>DEALLOCATE</i> together with <i>MARK</i> and <i>RELEASE</i>.</p> <p>When variant records are allocated using a tag field, you should use the same tag when disposing of that record.</p>
See Also	<p>DEALLOCATE MARK NEW RELEASE STORAGE</p>
Example	<p>Allocate a variable and use it:</p> <pre>MODULE Dispose; VAR p1, p2, p3: POINTER TO REAL; BEGIN NEW(p1); NEW(p2); NEW(p3); (* use them *) DISPOSE(p2); DISPOSE(p3); DISPOSE(p1) END Dispose.</pre>

DIV standard operator R

Description DIV is an integer division operator; it is a binary operator that works on operands of the same type. The operands may be both CARDINAL, both INTEGER, or both LONGINT.

Usage `i := i DIV j;`

where *i* and *j* are of the same type.

The result of the expression is the same type as the operands.

Comments Integer division is useful when you do not care about the fractional parts of a quotient.

See Also MOD

Example

```
MODULE IntegerDivision;
VAR
  i,j : INTEGER;
  c,d : CARDINAL;
  l,m : LONGINT;
BEGIN
  i := 1; c := 1;
  i := 30 DIV i;
  c := d DIV c;
  l := m DIV 344212L;
END IntegerDivision.
```

Done procedure (InOut) L

Description Done checks for successful completion of input.

Declaration VAR Done: BOOLEAN;

Usage IF Done THEN DoSomething END ;

Done must be imported from the library module InOut.

Comments *Done* allows a program to check for valid input before proceeding. Each of the input procedures in *InOut* sets *Done* as follows:

Read	Done := Not past end of input
ReadString	Done := The returned string is not empty
ReadInt	Done := Integer was read
ReadCard	Done := Cardinal was read

See Also *Texts*

Example Get a number from the user:

```
MODULE ReadInteger;
FROM InOut IMPORT Done, ReadInt;
BEGIN
  REPEAT
    WRITE('Gimme a number>');
    ReadInt(i)
  UNTIL Done;
END ReadInteger.
```

Done procedure (*Texts*) L

Description *Done* checks for successful completion of numeric input.

Declaration PROCEDURE Done(*t*: TEXT): BOOLEAN;

Usage *okay* := Done(*t*)

Done must be imported from the library module *Texts*.

t must be of type *TEXT* imported from the library module *Texts*.

Comments	<i>Done</i> allows the program to check for valid numeric input before proceeding.
	<i>Done</i> is a function returning TRUE if the numeric input is both syntactically correct and the input number is in the valid range for the type being read; otherwise FALSE is returned.
See Also	Texts
Example	Get a number from the user:

```
MODULE DoneInput;
FROM Texts IMPORT Done;
BEGIN
  REPEAT
    WRITE('Gimme a number>');
    READ(num);
  UNTIL Done(input);
END DoneInput.
```

DOUBLE standard function S

Description	<i>DOUBLE</i> converts the argument to type LONGREAL.
Usage	<i>Y</i> := DOUBLE(<i>X</i>)
	Argument <i>X</i> is of type CARDINAL, INTEGER, LONGINT, or REAL.
	The result <i>Y</i> is of type LONGREAL.

Comments Do not confuse this function with the LONGREAL type-transfer function, which transfers data without converting it.

See Also FLOAT
LONGREAL
TRUNC

Example

```
MODULE Doubles;
VAR
  Converted: LONGREAL;
  r: REAL;
BEGIN
  r := 234.234;
  Converted := DOUBLE(r);
END Doubles.
```

Doubles L

Description *Doubles* is a library module that provides string handling and I/O support for double-precision variables.

Declaration

```
DEFINITION MODULE Double
FROM Texts IMPORT TEXT;

VAR
  legal :BOOLEAN

PROCEDURE ReadDouble (t: TEXT; VAR d: LONGREAL);
PROCEDURE WriteDouble (t: TEXT; d: LONGREAL;
width: CARDINAL; digits: INTEGER);

PROCEDURE StrToDouble(VAR s: ARRAY OF CHAR;
d: LONGREAL);
PROCEDURE DoubleToStr (s: ARRAY OF CHAR;
VAR d: LONGREAL);

END Doubles.
```

Comments This module provides conversions for LONGREALs, which are similar to the ones that *Convert* and *Texts* provide for the standard types in Modula-2.

DoubleToStr procedure (Doubles) L

Description *DoubleToStr* converts a double-precision variable to a string variable.

Declaration **PROCEDURE DoubleToStr**
 (*d*: LONGREAL; **VAR** *s*: **ARRAY OF CHAR**;
 digits: INTEGER);

Usage **DoubleToStr(***d*, *s*, *digits*)

DoubleToStr must be imported from the library module *Doubles*.

s must be a character array variable.

d must be of type LONGREAL.

See Also **Doubles**
 StrToDouble

Example Convert a double-precision number into a string:

```
d := 2.99D10;  
DoubleToStr(d, str, 15);
```

END R

Description END marks the end of the scope of a control structure, procedure, or module.

Usage **END**

or

END *pname*

Comments	In the case of a control structure (IF, WITH, WHILE, or CASE), the END stands alone. In the case of a module or procedure, the END statement is followed by the name of the module or procedure it terminates. A semicolon is never needed before an END.
Example	<pre>MODULE endexample; BEGIN (* Main *) IF appropriate THEN DO WITH SomeRecord DO (* code *) END END END endexample.</pre>

EndError exception (Files) L/E

Description	EndError is raised by the <i>Files</i> module when an attempt is made to read past the end of the file.
Declaration	EXCEPTION EndError;
Usage	EXCEPTION EndError: WRITELN('END OF FILE'); END
	EndError must be imported from the module <i>Files</i> .
Comments	Use the EndError exception to prevent your program from crashing from user input.
See Also	StatusError UseError DeviceError DiskFull

Example

```

MODULE ReadFile;
FROM Files IMPORT FILE, Open, ReadWord,
Close, EndError;
VAR f: FILE; c: CARDINAL;
BEGIN
  IF Open(f,'testfile.dat') THEN
    LOOP
      ReadWord(f,c);
    END
  END ;
EXCEPTION
  EndError: WRITELN('End of file reached.');?>
END ReadFile.
```

Entier function (MathLib, LongMath) L

Description *Entier* returns the integer part of a real number rounded toward negative infinity.

Declaration

```

PROCEDURE Entier(x: REAL): INTEGER;
PROCEDURE Entier(x: LONGREAL): LONGINT;
```

Usage $Y := \text{Entier}(X)$

Argument X is REAL or LONGREAL.

Result Y is INTEGER or LONGINT.

Comments The Turbo Modula-2 predefined function *INT* can convert any numeric type to INTEGER; for example:

```

Entier(-3.5) = -4
INT(-3.5) = -3
```

Example

```

MODULE EntierTest;
IMPORT MathLib;
IMPORT LongMath;
VAR
  Round: INTEGER;
  Number: REAL;
  LongRound: LONGINT;
  LongNumber: LONGREAL;
BEGIN
  Round:= MathLib.Entier(Number + 0.5);
  LongRound := LongMath.Entier(LongNumber + 0.5);
END EntierTest.
```

EOF function (Files) L

Description	<i>EOF</i> returns true if the file pointer is at the end of the file.
Declaration	PROCEDURE EOF(<i>f</i> : FILE): BOOLEAN;
Usage	<i>endOF</i> := EOF(<i>f</i>)
	<i>EOF</i> must be imported from the library module <i>File</i> .
	<i>f</i> must be of type <i>FILE</i> imported from <i>Files</i> .
Comments	This is the normal way to detect the end of defined data in a file.
See Also	Close Files Open

Example Open a file and process it:

```
MODULE EndOfFile;
FROM Files IMPORT FILE, Open, ReadRec, EOF;
VAR
  f: FILE;
  data : ARRAY [0..99] OF REAL;
BEGIN
  IF Opn(f, "Numbers.dat") THEN
    WHILE NOT EOF(f) DO
      ReadRec(f,data);
      (* Code to process data *)
    END
  END
END EndOfFile.
```

file.

EOL constant (*Texts*) L

Description	<i>EOL</i> holds the value of the end-of-line character, 36C or 30 decimal.
Declaration	CONST EOL = 36C;
Usage	ch := EOL; <i>ch</i> must be defined as a CHAR.
Comments	<i>EOL</i> must be imported from the module <i>Texts</i> . The procedures in the module <i>Texts</i> convert the carriage-return/line-feed sequence to and from the <i>EOL</i> character. When input, the carriage-return/line-feed sequence is converted to an <i>EOL</i> character; and when output, the <i>EOL</i> character is converted to a carriage-return/line-feed pair. Thus in a program-processing text file, the end of line is marked with the <i>EOL</i> character.

See Also Texts

Example

```
MODULE EndOfTheLine;
FROM Texts IMPORT EOL;
VAR ch: CHAR;
BEGIN
REPEAT
  READ(ch)
UNTIL ch=EOL;
END EndOfTheLine.
```

EOLN procedure (Texts) L

Description *EOLN* returns TRUE if the *TEXT* file is at the end of a line.

Declaration

```
PROCEDURE EOLN(t: TEXT): BOOLEAN;
```

Usage

```
endOL := EOLN (t)
```

EOLN must be imported from the library module *Texts*.

t must be of type *TEXT* imported from *Texts*.

endOL must be of type BOOLEAN.

Comments

EOLN returns TRUE if the last character read from *t* is an *EOL* character.

See Also

EOT
Texts

Example Read all characters until the end of the line is encountered:

```
REPEAT
  ReadChar(t,ch);
UNTIL EOLN(t);
```

EOT procedure (Texts) L

Description *EOT* returns TRUE if the *TEXT* file is at its end.

Declaration **PROCEDURE EOT(t: TEXT): BOOLEAN;**

Usage *EndOfText* := *EOT* (*t*)

EOT must be imported from the library module *Texts*.

t must be of type *TEXT* imported from *Texts*.

EndOfText must be of type BOOLEAN.

Comments *EOT* returns TRUE if the last character read from *t* is an *EOT* character (for example, Control-Z or 32C).

Notice there is no *EOT* character defined in Turbo Modula-2; it is reserved for the function described here. Note that sometimes the 32C character is referred to as *EOT*. Also note that the EOT character on other systems may not be 32C.

See Also **EOF**
 EOLN
 Texts

Example Read all characters until the end of the *TEXT* file *t* is encountered:

```
REPEAT
  ReadChar(t, ch);
UNTIL EOT(t)
```

EXCEPTION R

Description **EXCEPTION** is used to declare exception identifiers and declare code to trap exceptions when they occur.

Usage Exception Declarations

```
MODULE deviceHandler;
EXCEPTION ex1, ex2, ex3; (* Exception declaration *)
```

```
PROCEDURE dev1;
BEGIN
(* code with RAISEs *)
END dev1;
```

```
END devicehandler;
```

Exception Handlers

```
MODULE deviceMonitor;
IMPORT ex1,ex2,ex3;
BEGIN
dev1; (* Invoke a procedure with RAISEs *)
EXCEPTION
| ex1:   (* ex1  code *)
| ex2:   (* ex2  code *)
| ex3:   (* ex3  code *)
ELSE    (* other code *)
END deviceMonitor.
```

The first module declares exceptions and code that will raise the exceptions in error situations. The second module calls a procedure that may raise an error; therefore, it declares exception handlers to catch and process the errors.

Comments Exceptions are particularly useful in the event of a program crash; for instance, if you need to turn off a laser, or switch out of graphics mode. It is possible to do this with the Turbo Modula-2 system, because the runtime system can take over (in most cases) when something catastrophic happens to your logic. The **ELSE** clause in the exception handler catches all possible exceptions whether imported or not.

See Also RAISE

Example Declare *PowerFailureImminent* as an exception:

```
EXCEPTION PowerFailureImminent
```

Then raise it:

```
IF ( (* some hint *) ) THEN RAISE
```

```
PowerFailureImminent;
```

Then handle it:

```
EXCEPTION
```

```
PowerFailureImminent: WRITE('Help !');  
END someMainModule;
```

EXCL standard procedure S

Description EXCL excludes element *I* from the set *S*.

Usage EXCL(*S*, *I*)

Argument *S* is of set type.

Argument *I* is of the set's base type.

Comments This standard procedure allows you to remove an element from a set. This has the effect of turning off a bit in the word that represents the set.

See Also BITSET
 INCL
 SET

Example **MODULE** Exclude;
TYPE
 Day = (Sun,Mon,Tues,Wed,Thu,Fri,Sat);

VAR
 Work: **SET OF** Day;
BEGIN
 Work:= Day[Mon..Fri];
(* Set Work contains five elements Mon to Fri *)

 EXCL(Work,Fri);
(* Set Work contains four elements, Wed excluded
END Exclude.

EXIT statement R

Description EXIT enables you to exit from a LOOP statement.

Usage **EXIT**

Comments EXIT jumps to the statement following the **END** of the enclosing LOOP statement.

See Also LOOP

Example LOOP
 IF DoneProcessing **THEN EXIT END** ;
END ;
HALT

ExitScreen procedure (Terminal) L

Description *ExitScreen* sends a terminal reset string after a series of screen operations, if needed.

Declaration **PROCEDURE** ExitScreen;

Usage ExitScreen;

ExitScreen must be imported from the library module *Terminal*.

Comments This function will work only if the Turbo Modula-2 system has been installed on your terminal.

) Many terminals don't need or have a reset string. If this is the case with your terminal, the function will do nothing.

Example Send a terminal reset string:

```
BEGIN
  (* Lots of screen operations *)
  ExitScreen
END
```

Exp function (MathLib, LongMath) L

Description Exp returns the natural exponential of X.

Declaration **PROCEDURE** Exp(x: REAL): REAL;
PROCEDURE Exp(x: LONGREAL): LONGREAL;

Usage Y := Exp(X)

Argument X and result Y are both REAL or both LONGREAL.

Comments	Be certain that argument <i>X</i> is less than 87.4 for single-precision numbers and less than 710.47 for double-precision numbers; otherwise, the <i>ArgumentError</i> exception is raised.
	The inverse of the exponential function is the natural logarithmic function <i>Ln</i> .
	If you need to calculate <i>X</i> to the power of <i>Y</i> , you can use the <i>Exp</i> and <i>Ln</i> functions in the following expression:
	<pre>XPowerY:= Exp(Ln(X) * Y)</pre>

See Also **LongMath**
MathLib

Example

```
MODULE Exp;
FROM MathLib IMPORT Exp;
VAR
  Argument,Result:REAL
BEGIN
  IF Argument < 87.4 THEN
    Result:= Exp(Argument)
  END ;
END Exp.
```

EXPORT R

Description	EXPORT specifies which identifiers inside a module are to be visible in the scope surrounding the module.
Usage	EXPORT var1,var2, type1,type2, proc1,proc2

Comments	<p>IMPORT and EXPORT lists must appear immediately after the MODULE statement and before any constant, variable, type, or procedure declarations.</p> <p>Names identifying constants, variables, types, exceptions, and procedures may be exported. Modules may not be exported.</p> <p>Everything in a definition module is exported; thus no export list is required.</p> <p>Nothing in an implementation module may be exported, unless it has been previously exported by its definition module.</p> <p>Nothing may be exported by a main module since there is nothing to export to.</p> <p>Consequently, export lists really only make sense in local modules or in modules that are embedded inside of a main module, procedure, or an implementation module.</p>
Example	<p><i>local's beans</i> is visible to <i>main</i>, but its <i>rice</i> is not:</p> <pre>MODULE main; MODULE local; EXPORT beans; VAR rice,beans: BOOLEAN; BEGIN (* code to set rice and beans *) END local; (* Any code here cannot see rice, only beans *) END main.</pre>

FALSE standard value S

Description	FALSE denotes the Boolean state of falsity.
Usage	Finished := FALSE; <i>Finished</i> is a variable of type BOOLEAN.
Comments	The ordinal value of FALSE is 0; thus, <i>WRITE(CARDINAL(FALSE))</i> will print a 0. In contrast, the ordinal value of TRUE is 1; thus, truth is greater than falsity.
See Also	BOOLEAN TRUE
Example	MODULE Falsity; VAR b: BOOLEAN; BEGIN REPEAT b := FALSE UNTIL b; (* repeats forever *) END Falsity.

FILE type (Files) L

Description	<i>FILE</i> is an opaque type representing a low-level file. Variables of type <i>FILE</i> may represent illegible streams and random access files.
Declaration	TYPE FILE;
Usage	VAR f: FILE;
Comments	<i>FILE</i> must be imported from the module <i>Files</i> . Use <i>FILE</i> when you are working with files that are not strictly ASCII characters or if you do not want to interpret the contents of the file.

See Also	Files TEXT Texts
Example	<pre>MODULE FileExists; FROM Files IMPORT FILE, Open, Close; VAR f: FILE; BEGIN IF Open(f,'testfile.dat') THEN Close(f); WRITELN('File exists.'); ELSE WRITELN('File does not exist.'); END ; END FileExists.</pre>

Files module	L
--------------	---

Description	<i>Files</i> facilitates the handling of disk files.
Declaration	<pre>DEFINITION MODULE Files; FROM SYSTEM IMPORT BYTE WORD, ADDRESS; TYPE FILE; PROCEDURE Open (VAR f: FILE; name: ARRAY OF CHAR): BOOLEAN; PROCEDURE Create (VAR f: FILE; name: ARRAY OF CHAR); PROCEDURE Close (VAR f: FILE); PROCEDURE Delete (VAR f: FILE); PROCEDURE Rename (VAR f: FILE; name: ARRAY OF CHAR); PROCEDURE GetName (f: FILE; VAR name: ARRAY OF CHAR);</pre>

```
PROCEDURE FileSize    (f: FILE): LONGINT;
PROCEDURE EOF         (f: FILE): BOOLEAN;
PROCEDURE ReadByte    (f: FILE; VAR b: BYTE);
PROCEDURE ReadWord    (f: FILE; VAR w: WORD);
PROCEDURE ReadRec     (f: FILE;
                      VAR rec: ARRAY OF WORD);
PROCEDURE ReadBytes   (f: FILE; buf: ADDRESS;
                       nbytes: CARDINAL): CARDINAL;

PROCEDURE WriteByte   (f: FILE; b: BYTE);
PROCEDURE WriteWord   (f: FILE; w: WORD);
PROCEDURE WriteRec    (f: FILE;
                      VAR rec: ARRAY OF WORD);
PROCEDURE WriteBytes  (f: FILE; buf: ADDRESS;
                       nbytes: CARDINAL);

PROCEDURE Flush        (f: FILE);
PROCEDURE NextPos     (f: FILE): LONGINT;
PROCEDURE SetPos       (f: FILE; pos: LONGINT);

PROCEDURE NoTrailer   (f: FILE);

PROCEDURE ResetSys( );

EXCEPTION EndError, StatusError, UseError,
         DeviceError, DiskFull;

END Files.
```

Comments This library module performs low-level disk access.

See Also

Texts

InOut

Also see individual identifiers declared in the module *Files*.

FileSize procedure (Files) L

Description *FileSize* gets the exact size of the file in bytes.

Declaration PROCEDURE FileSize(f: FILE): LONGINT;

Usage nbytes := FileSize(f)

FileSize must be imported from the library module *Files*.

nbytes must be of type LONGINT.

f must be of type *FILE* imported from *Files*.

Comments Turbo Modula-2 uses the last byte of the last record (all CP/M files have 128 bytes per record) to store the number of bytes actually used in this record. This number is added to the number of bytes in the previous records (*nrecs-l**128) to determine the actual file size in bytes.

This processing of the last byte can be turned off by a call to *NoTrailer*, but then *FileSize* will only return sizes in 128-byte multiples.

See Also FILE

Files

NoTrailer

Example To see just how big *widgetFile* is, do the following:

```
MODULE FileSizes;
FROM Files IMPORT FILE, Open;
VAR
  nbytes: LONGINT;
  widgetFile: FILE;
BEGIN
  IF Open(widgetFile, "widgets.dat") THEN
    nbytes := FileSize(widgetFile)
  END
END FileSizes.
```

FILL procedure (SYSTEM) L

Description	<i>FILL</i> fills a block of memory with a given byte.
Declaration	PROCEDURE FILL(<i>adr</i> : ADDRESS; <i>len</i> : CARDINAL; <i>val</i> : BYTE);
Usage	FILL (<i>start</i> , <i>nbytes</i> , <i>ch</i>)
	<i>FILL</i> must be imported from the library module SYSTEM .
	<i>start</i> must be of type ADDRESS imported from SYSTEM .
	<i>nbytes</i> must be of type CARDINAL .
	<i>ch</i> must be of type BYTE imported from SYSTEM .
Comments	The block of memory between <i>start</i> and the address <i>start</i> + <i>nbytes</i> - 1 is filled with the byte value <i>ch</i> .
See Also	MOVE SYSTEM
Example	Fill the array <i>Sieve</i> with the value TRUE:
	VAR Sieve: ARRAY [0..8191] OF BOOLEAN; BEGIN ... FILL (ADR (Sieve), 8192, TRUE);

firstDrive variable (Loader) L

Description	<i>firstDrive</i> contains the first disk drive to be searched when looking for overlay files.
Declaration	VAR <i>firstDrive</i> : [0..15];

Usage	<code>firstDrive := currentDrive;</code>
	<i>firstDrive</i> must be imported from the module <i>Loader</i> .
	<i>currentDrive</i> is a CARDINAL in the range 0 to 15.
Comments	Once the location of the overlay files has been found you may speed up load time by ensuring that <i>firstDrive</i> is set to the correct drive number.
See Also	<i>Loader</i>
Example	<pre>MODULE LoadOverlay; FROM Loader IMPORT firstDrive, Call; BEGIN firstDrive := 13; (* Search ram disk (M:) first *) Call('overlay.fil') END LoadOverlay.</pre>

FLOAT standard function S

Description	<i>FLOAT</i> converts argument <i>X</i> to REAL.
Usage	<code>Y := FLOAT(X)</code>
	Argument <i>X</i> can be of type INTEGER, CARDINAL, LONGINT, or REAL.
	The argument must be in the REAL range -1E38 to 1E38.
	Result <i>Y</i> is of type REAL.
Comments	Be certain that the argument is in the REAL range -6.80564E38 to +6.80564E38; otherwise, an overflow error will occur, causing the exception OVERFLOW to be raised.

Example

```

VAR Count:CARDINAL
      Time:REAL
BEGIN
      Time:= Time + 0.1 * FLOAT(Count);
END

```

Flush procedure (Files) L

Description *Flush* flushes the file's internal buffer to disk.

Declaration **PROCEDURE** Flush(*f*: FILE);

Usage *Flush(f)*

Flush must be imported from the library module *Files*.

f must be of type *FILE* imported from *Files*.

Comments In order to speed up the input and output to disk files, most computers maintain a file buffer for each open file. In a write operation this buffer is not normally »flushed« until the buffer is full. The *Flush* command forces this process.

This function can be used to detect write errors (like DiskFull) sooner than they would otherwise appear. (They do not usually appear until the internal file buffer is full and the Turbo Modula-2 system orders a flush.)

See Also *Files*

Example Force a DiskFull error at the record where it occurs:

```

FOR i := 1 TO nrecs DO
  WriteRec(f, widget[i]);
  Flush(f);
END ;

```

FOR statement R

Description FOR repeats the execution of a sequence of instructions a specified number of times.

Usage

```
FOR ControlVar := StartExp TO EndExpr [BY Step] DO
  <Statement> {; <Statement> }
END
```

ControlVar is of type CARDINAL, INTEGER, CHAR, enumeration, or subrange. It must not be imported, nor be a procedure parameter or a structured variable's component.

StartExp and *EndExp* are single values or expressions of the same type as the control variable.

Step is a constant expression of type CARDINAL or INTEGER.

Comments The FOR statement repeatedly executes the statement sequence for a progression of control variable values. It begins with the value of the start expression and increments by the value *Step* until the control variable is equal to or greater than the value given by the end expression. If you do not specify a step value, it is assumed to be 1.

As defined in the syntax section, the start and end values can be the results of expressions.

```
FOR Y:= X TO X + N BY Ramp DO
```

A FOR statement cannot be exited before its normal termination. If you want to do this, you can use the REPEAT or WHILE statements instead.

Example

```
FOR X:= 0 TO 10 BY 2 DO
  (* Some statement sequence *)
END
```

This example repeats the statement sequence six times, with *X* equal to the values 0, 2, 4, 6, 8, and 10. You can use *X* in the statement sequence, but you cannot alter its value. For example, you can do the following:

```
FOR Count:= 5 TO 3 BY -1 DO
  T:= 2 * Count + 4
END
```

This example is repeated three times for descending values of *Count* equal to -5, -4, and -3.

FORWARD statement R

Description **FORWARD** allows a procedure to be referenced before it is defined.

Usage PROCEDURE *foo*(): CARDINAL; **FORWARD** ;

foo is a procedure or a function heading followed by the reserved word **FORWARD**.

The body of *foo* is given after a normal procedure heading for *foo*.

Comments The **FORWARD** statement allows Turbo Modula-2 to implement mutual recursion in a one-pass compiler environment. This statement was not defined in the original definition, but is needed for a one-pass implementation of the compiler.

Example

```
MODULE Recurring;
PROCEDURE one; FORWARD ;
PROCEDURE two;
BEGIN
  one; (* Call of procedure before it is defined *)
END two;

PROCEDURE one;
BEGIN
  two;
END one;

BEGIN
  one;
END Recurring.
```

FREEMEM procedure (STORAGE) L

Description *FREEMEM* finds out how much contiguous dynamic memory (heap space) is available.

Declaration **PROCEDURE FREEMEM()**: CARDINAL;

Usage *freeman* := FREEMEM()

FREEMEM must be imported from the pseudomodule **STORAGE**.

freeman must be of type **CARDINAL**.

Comments This function returns the number of bytes between the heap pointer and the current top of the runtime stack. The sizes of holes in the allocated heap is not taken into account.

See Also **ALLOCATE**
DEALLOCATE
STORAGE
OUTOFMEMORY

Example WRITELN('There are ',FREEMEM(),' bytes free on the heap');
 ALLOCATE(p, 2000);
 WRITELN('and now there are 2000 less:',FREEMEM());

GetName procedure (Files) L

Description *GetName* retrieves the name of an open file.

Declaration PROCEDURE GetName(f: FILE; VAR name:
 ARRAY OF CHAR);

Usage GetName(f, fname)

GetName must be imported from the library module *Files*.

f must be of type *FILE* imported from *Files*.

fname must be a character array variable.

Comments This procedure allows library modules to determine the name of a file when only the file variable has been passed.

See Also Create
 FILE
 Files
 Open

Example Get the name of file *f*:

GetName(f, fname)

GotoXY procedure (Terminal) L

Description *GotoXY* positions the cursor on the screen.

Declaration **PROCEDURE GotoXY(x,y: CARDINAL);**

Usage *GotoXY*(*col*, *row*)

GotoXY must be imported from the library module *Terminal*.

col and *row* must be of type CARDINAL.

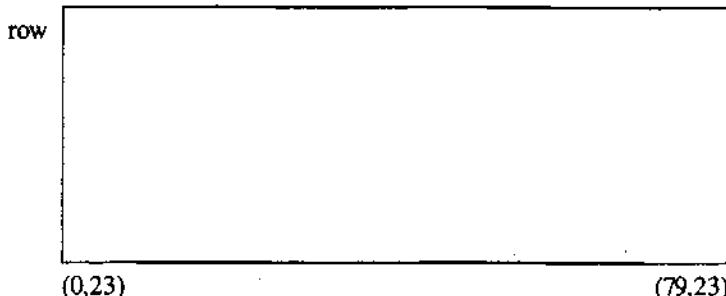
Comments This function will only work if Turbo Modula-2 has been properly installed on your system.

0, 0 is the upper left-hand corner

79,23 is the lower right-hand corner

(0,0) *col*

(79,0)



See Also [ClearScreen](#)
[ClearToEOL](#)
[DeleteLine](#)
[InsertLine](#)
[Highlight](#)
[Normal](#)
[Terminal](#)

Example Put an 'X' in the upper right-hand corner:

```
GotoXY( 79,0 );WriteChar( 'X' );
```

HALT standard procedure S

Description *HALT* stops program execution.

Usage **HALT**

Comments The *HALT* procedure terminates execution of the program and returns you to Turbo Modula-2 or the operating system, depending on where the program execution originated from.

HALT is a standard Modula-2 procedure; a nonstandard way to stop a program and deliver an error message to the terminal is available with the *RAISE* command.

See Also **RAISE**

Example **VAR** ErrorCondition:BOOLEAN
 BEGIN
 IF ErrorCondition **THEN**
 HALT
 END ;

haltOnControlC variable (Texts) L

Description *haltOnControlC* holds the current mode of operation of Control-C during console input and output.

Declaration **VAR**
 haltOnControlC: BOOLEAN;

Usage	<code>haltOnControlC := FALSE;</code>
	<i>haltOnControlC</i> is a predefined variable that must be imported from the module <i>Terminal</i> .
Comments	This variable is useful for preventing a user from stopping a program.
	The default value of <i>haltOnControlC</i> is TRUE, which means that the user can press Ctrl C during an <i>input</i> statement or a <i>WRITELN</i> statement to halt the program.
	Note that the only output procedure that checks for the Control-C character is the <i>WriteLn</i> procedure. Thus, programs that do not use a <i>WriteLn</i> (or <i>WRITELN</i>) statement cannot be interrupted.
See Also	Texts
Example	<pre>MODULE HaltTheProgram; FROM Texts IMPORT haltOnControlC; VAR I: CARDINAL; BEGIN FOR I := 1 TO 20 DO WRITELN('Press control C to stop the program ') END ; haltOnControlC := FALSE; FOR I := 1 TO 20 DO WRITELN('You cannot stop me now . . . ') END ; END HaltTheProgram.</pre>

HIGH standard function S

Description To *HIGH* returns the high-index bound of array *A*.

Usage $Y := \text{HIGH}(A)$

Argument *A* is an array of any index and component type.

Result *Y* is of the same type as the index type.

Comments This function allows you to find the upper index of any array for example, an array defined as

Move: **ARRAY** (Left,Right,Stop) **OF** INTEGER;

has an upper index *Stop*.

Note that for multidimensional arrays, the upper index of the first dimension is returned. For example, if you have defined an array as

MultiDim: **ARRAY** [10..20],(Left,Right,Stop) **OF**
CHAR

the *HIGH* function will return the result 20, which is the high index of the first dimension.

You can also find the other dimensions as well; for example if you have defined

MultiDim: **ARRAY** [0..10] **OF** **ARRAY** [0..20] **OF** WORD;

then $\text{HIGH}(\text{MultiDim}) = 10$ and $\text{HIGH}(\text{MultiDim}[0]) = 20$

HIGH is very useful for finding the upper bound of an open array parameter. The lower bound is always zero and the upper bound is $\text{HIGH}(a)$ where *a* is an open array parameter.

See Also

ARRAY

Example

```

MODULE High;
TYPE
  Weather = (Clear,Rain,Wind,Snow);
  Day     = (Mon,Tue,Wed,Thu,Fri,Sat,Sun);
  Hour    = [1..24]
VAR
  Forecast: ARRAY Day,Hour OF Weather;
  HighForecast: Day;
  HighHour: Hour
BEGIN
  HighForecast:= HIGH(Forecast);
  (* HighForecast is value Sun *)
  HighHour:= HIGH(Forecast[Sun]);
  (* HighHour is 23 *)
END High.
```

Highlight procedure (Terminal) L

Description *Highlight* turns on highlighting (brightens text) for terminal output.

Declaration **PROCEDURE** Highlight;

Usage *Highlight*

Highlight must be imported from the library module *Terminal*.

Comments *Highlight* is only possible on systems with screens capable of displaying text in at least two of the following modes: brightness, inverse video, or underlining.

Highlight will work on your system only after the Turbo Modula-2 system has been properly installed.

Stand-alone programs can see if the terminal they are running on has this capability by checking *OpSet* for the element *highlightNormal*.

See Also	Normal Terminal OpSet highlightNormal
-----------------	--

Example	Highlight the words » <i>This text is high</i> «:
----------------	---

```
Highlight;
WriteLn( 'This text is high ' );
Normal;
WriteLn( 'And this text is normal ' )
```

highlightNormal enumerated value (Terminal) L

Description *highlightNormal* is the third value of the enumerated type *Terminal.SpecialOps*.

Declaration *SpecialOps* = (*clearEol*, *insertDelete*, *highlightNormal*);

Usage IF *highlightNormal* IN *available* THEN *Highlight* END;

SpecialOps must be imported from the library module *Terminal*.

Comments If this value is a member of the set *Terminal.available*, then the program may use the procedure *Highlight* to enhance text written to the screen.

To use this identifier, include *SpecialOps* in your import list. The identifiers of each of *SpecialOps*' values become visible automatically.

See Also	Highlight <i>available</i> <i>insertDelete</i> <i>clearEol</i> <i>Terminal</i>
-----------------	--

Example Write some enhanced text:

```
MODULE CheckAvailableOperations;
FROM Terminal IMPORT
available,SpecialOps,Highlight,Normal;
FROM Strings IMPORT CAPS;
VAR s: ARRAY [0..20] OF CHAR;
BEGIN
  IF highlightNormal IN available THEN
    Highlight;
    WRITELN(s);
    Normal;
  ELSE
    CAPS(s)
    WRITELN(s);
  END ;
END CheckAvailableOperations.
```

HRESULT variable (SYSTEM) L

Description *HRESULT* holds the contents of the *HL* register after a BDOS or BIOS call.

Declaration **VAR**
 HRESULT: CARDINAL;

Usage *c* := *HRESULT*;

HRESULT must be imported from the pseudomodule *SYSTEM*.

c must be of type CARDINAL;

Comments This variable is used to access the register after operating system calls.

See Also	BDOS BIOS HRESULT SYSTEM
Example	<pre>MODULE OperatingSystemCalls; FROM SYSTEM IMPORT BDOS, BIOS, IORESULT, HRESULT; VAR ReadOnlyVector: CARDINAL; BEGIN BDOS(1,0); WRITE('The character read is ',CHR(IORESULT)); BDOS(29,0); ReadOnlyVector := HRESULT; END OperatingSystemCalls.</pre>

IF statement	R
Description	IF executes a sequence of statements that depend on the result of a BOOLEAN expression.
Usage	<pre>IF BooleanExpression THEN StatementSequence { ELSIF BooleanExpression THEN StatementSequence } [ELSE StatementSequence] END</pre> <p style="text-align: center;">StatementSequence = Statement {; Statement}</p>
Comments	<p>Each Boolean expression is evaluated until one results in TRUE. The associated statement sequence is then executed and the statement ends. If none of the IF or ELSIF expressions yield a TRUE result, the optional ELSE statement sequence is executed.</p> <p>Note that Boolean expressions are only evaluated until one is TRUE.</p>

Example

```
VAR
  XAxis,YAxis: REAL;
BEGIN
  IF XAxis < 10.0 THEN
    YAxis:= 2.0 * XAxis
  ELSIF XAxis < 20.0 THEN
    YAxis:= XAxis + 10.0
  ELSE
    YAxis:= 30.0
  END ;
```

IMPLEMENTATION declaration R

Description **IMPLEMENTATION** marks the begining of the implementation code of a module.

Usage **IMPLEMENTATION MODULE** moduleName;
 <declarations>
 BEGIN
 <body>
 END moduleName.

Comments

An implementation module always has a definition module, and together they form a library module from which objects can be imported. The definition module contains the declarations of all the exported objects. The implementation module contains the actual code that defines the objects and procedures that are exported.

An implementation module is a separate compilation unit. In general, it is possible to make a change to an implementation module and recompile it without changing the client modules. Other .MCD files that import objects from the changed module will incorporate those changes without having to be recompiled. This is possible because Turbo Modula-2 programs are linked at runtime when run from the menu shell.

An implementation module is not normally a stand-alone program (though it can be); it is usually executed by importing its code into a so-called »main module«.

The main body of code in an implementation module is executed at load time. This is referred to as the initialization section. It is used to set up variables before any client uses that library's services.

See Also

DEFINITION
Linker
MODULE

Example

Here's a definition module in the file COMPUTE.DEF:

```
DEFINITION MODULE Compute;
VAR
  PI: REAL;

  PROCEDURE Add( Y,Z :INTEGER):INTEGER;
  PROCEDURE Sub( Y,Z :INTEGER):INTEGER
END Compute.
```

And here's the corresponding implementation module in file COMPUTE.MOD:

```
IMPLEMENTATION MODULE Compute;

PROCEDURE Add( Y,Z :INTEGER):INTEGER;
  RETURN Y + Z
END Add;

PROCEDURE Sub( Y,Z :INTEGER): INTEGER;
  RETURN Y - Z
END Sub;

BEGIN
  PI := 3.14159;
END Compute.
```

IMPORT declaration R

Description **IMPORT** specifies which identifiers a module is able to »see« from another module or the surrounding environment.

Usage **IMPORT** identifiers;

or

FROM moduleName IMPORT identifiers;

The first form of **IMPORT** imports identifiers as a whole. If the identifier is a module, then that module's identifiers must be qualified.

The second form de-qualifies identifiers as they are imported.

Comments **IMPORT** lists must come immediately after the **MODULE** declaration and before any constant, variable, type, or procedure declarations.

In local modules, the **IMPORT** declarations must come before the **EXPORT** declarations.

Qualified identifiers can be accessed by prefixing the identifier with the module name, followed by a «.», as **ModuleName.IdentifierName**.

Note that you need not name the module the same as the file name. However, the compiler, the linker, and the loader use the first eight characters of the name given in the **MODULE** heading when searching for symbol files, library code file and linked overlay files. Thus, we recommend that you always use the first eight letters of the module name for the actual file name.

See Also **EXPORT**
 QUALIFIED

Example Input the *Append* procedure from *Strings* and use it in the following:

```
MODULE SayHello;
FROM Strings IMPORT Append;
VAR
  s : ARRAY [0..20] OF CHAR;
BEGIN
  s := 'Hello';
  Append(' Folks', s);
  WRITELN(s);
END SayHello.
```

Import all of *Strings* and do the same thing you did in the previous example:

```
MODULE SayHello;
IMPORT Strings;
VAR
  s : ARRAY [0..20] OF CHAR;
BEGIN
  s := 'Hello';
  Strings.Append( ' Folks', s); (* Qualified *)
  WRITELN(s);
END SayHello.
```

INC standard procedure S

Description *INC* returns a successor of argument *X*.

Usage *INC(X [,N])*

Argument *X* may be of type INTEGER, CARDINAL, CHAR, BOOLEAN or enumeration.

Optional argument *N* must be a CARDINAL; it is the increment that will be added to *X*.

Comments

If you add 1 to a CARDINAL value, the result is the value immediately after the original value. With *INC*, you can do the exact same thing with any ordinal type. For example, suppose you have an enumeration variable

```
VAR Energy: (Kinetic,Potential,Heat);
```

You can assign any of the three named values to the variable *Energy*. To increment the value of *Energy* to a higher value, use *INC(Energy)*.

Now, you have performed an addition operation on the enumeration variable. You can increment by more than one value as shown in the following examples.

The reverse operation selecting the previous value is performed by the standard procedure *DEC*.

When incrementing a value, you may reach its *MAX* value; for example, the *MAX(CARDINAL)* is 65535. You may choose whether or not this situation generates an error by using the overflow compiler switch. This can be done globally or locally. Thus, with overflow checking turned off, as in

```
VAR
  i: CARDINAL;
BEGIN
  i := MAX(CARDINAL);
  (*$O-*)
  INC(i);
  (*$O+*)
  WRITELN(i);
END
```

the *i* would wraparound to *MIN(CARDINAL)*, which is zero, and no error will be generated. If overflow checking has been turned on, then the exception OVERFLOW will be raised. This can be trapped as follows:

```
MODULE inc;
FROM SYSTEM IMPORT OVERFLOW;
VAR
  i: CARDINAL;
BEGIN
  (*$0+*)
  i := MAX(CARDINAL);
  INC(i);
EXCEPTION
  OVERFLOW: WRITELN('Trapped overflow error');
END inc.
```

See Also

DEC

Example

```
TYPE Color = (Red,Orange,Yellow,Green,Blue,
              Indigo,Violet);
VAR Rainbow: Color;
    Increment:CARDINAL

BEGIN
  Rainbow:= Red;
  INC(Rainbow);
  (* Rainbow has the value Orange *)

  INC(Rainbow,2);
  (* Rainbow has the value Green *)
  Increment:= 1;
  INC(Rainbow,Increment);
  (* Rainbow has the value Blue *)
```

INCL standard procedure S

Description	INCL includes element I into the set S.
Usage	INCL(S,I)
	Argument S is of set type.
	Argument I is of the set's base type.
Comments	This standard procedure allows you to include an element into a set. It has the effect of turning on a bit that represents the element in the word representing the set.
See Also	EXCL
Example	<pre>TYPE Day = (Mon,Tues,Wed,Thu,Fri,Sat,Sun); VAR Work: SET OF Day; BEGIN Work:= Day{Mon..Fri}; (* Set Work contains five elements Mon to Fri *) INCL(Work,Sun); (* Set Work contains six elements, Sun included *)</pre>

Init procedure (Processes) L

Description	<i>Init</i> initializes a SIGNAL variable.
Declaration	PROCEDURE Init(VAR s: SIGNAL);
Usage	<i>Init(s)</i>
	<i>Init</i> must be imported from the library module <i>Processes</i> . <i>s</i> must be of type SIGNAL imported from <i>Processes</i> .

Comments This is part of a module that gives a standardized set of coroutine facilities, as well as a rudimentary form of interprocess communications upon which to base scheduling.

The three routines *Init*, *SEND*, and *WAIT* work together in the following way: *Init* initializes a queue belonging to a given signal. Every time a process calls the procedure *WAIT* with this signal, it will be entered in the queue. When a process calls the procedure *SEND* with this signal, the first process in the queue will be resumed and will be purged from the queue. A fourth routine, *Awaited*, returns TRUE if the queue is not empty.

The initialization of a *SIGNAL* by *Init* is mandatory.

See Also *Awaited*
Processes
SEND
SIGNAL
StartProcess
Wait

Example Initialize the signal *DonePrinting*:

```
INIT(DonePrinting)
```

InitScreen procedure (Terminal) L

Description *InitScreen* sends an initialization string to the terminal.

Declaration **PROCEDURE** InitScreen;

Usage *InitScreen*

InitScreen must be imported from the library module *Terminal*.

Comments	<i>InitScreen</i> will send the initialization string specified by the user upon installation of the Turbo Modula-2 system. Many terminals don't need an initialization string, in which case nothing happens.
See Also	<i>EditScreen</i> <i>Terminal</i>
Example	Initialize the screen: <code>InitScreen;</code>

inName variable (ComLine) L

Description	<i>inName</i> is a string variable that contains a string found after the input redirection symbol (<) on the command line.
Declaration	VAR <i>inName</i> : ARRAY [0..19] OF CHAR ;
Usage	<code>WRITE('Input redirected from ',<i>inName</i>);</code>
Comments	The redirection symbol and <i>inName</i> will not appear in the <i>commandLine</i> . <i>inName</i> contains the standard output string CON: if redirection argument has been found on the command line.
See Also	<i>ComLine</i> <i>inName</i> <i>RedirectOutput</i>

Example **MODULE** RedirectInput;
 FROM ComLine **IMPORT** inName, RedirectInput;
 FROM Texts **IMPORT** console;
 BEGIN
 RedirectInput;
 IF inName < > "CON:" **THEN**
 WRITE(console,'Input redirected from ',inName);
 END
 END RedirectInput.

InOut module L

Description *InOut* performs input and output.

Usage **DEFINITION MODULE** InOut;

 CONST EOL=36C;
 VAR Done:BOOLEAN
 termCH: CHAR;

 PROCEDURE OpenInput(defext: **ARRAY OF** CHAR);
 PROCEDURE OpenOutput(defext: **ARRAY OF** CHAR);
 PROCEDURE CloseInput;
 PROCEDURE CloseOutput;

 PROCEDURE Read(**VAR** ch: CHAR);
 PROCEDURE ReadString(**VAR** s: **ARRAY OF** CHAR);
 PROCEDURE ReadInt(**VAR** x: INTEGER);
 PROCEDURE ReadCard(**VAR** x: CARDINAL);

 PROCEDURE Write(ch: CHAR);
 PROCEDURE WriteLn;
 PROCEDURE WriteString(s: **ARRAY OF** CHAR);
 PROCEDURE WriteInt(x: INTEGER; n: CARDINAL);
 PROCEDURE WriteCard(x,n: CARDINAL);
 PROCEDURE WriteHex(x,n: CARDINAL);
 PROCEDURE WriteOct(x,n: CARDINAL);

```
PROCEDURE ReadReal(VAR x: REAL);
PROCEDURE WriteReal(x: REAL; n,digits: CARDINAL);
```

END InOut.

Comments

This is the original input and output module defined by Wirth. It is postulated to be present in every implementation of Modula-2 and is included here to maintain compatibility with other implementations.

The module *InOut* is usually implemented using the module *Texts* with the *input* or *output* text streams.

WriteHex and *WriteOct* can sometimes prove useful for writing addresses.

IORESULT variable (SYSTEM) L

Description IORESULT holds the contents of the *A* register after a BDOS or BIOS call.

Declaration **VAR**
 IORESULT: CARDINAL;

Usage c := IORESULT;

IORESULT must be imported from the pseudomodule *SYSTEM*.

c must be of type CARDINAL;

Comments This variable is used to access the register after operating system calls.

See Also BDOS
 BIOS
 HLRESULT
 SYSTEM

Example **MODULE** OperatingSystemCalls;
 FROM SYSTEM IMPORT BDOS, BIOS, IORESULT, HLRESULT;
 VAR ReadOnlyVector: CARDINAL;
 BEGIN
 BDOS(1,0); WRITE('The character read is
 ',CHR(IORESULT));
 BDOS(29,0);
 ReadOnlyVector := HLRESULT;
 END OperatingSystemCalls.

INP procedure (SYSTEM) L

Description *INP* reads a byte from a given I/O port.

Declaration **PROCEDURE** INP(*port*: WORD): CARDINAL;

Usage *inByte* := INP(*port*)

INP must be imported from the library module **SYSTEM**.

inByte must be of type INTEGER or CARDINAL.

port can be any scalar type compatible with WORD.

Comments The port, whose number is given by *port*, is read and the value is placed in *inByte*.

See Also OUT
 SYSTEM

Example Get a byte by doing the following:

KbdStatus := INP(KbdStatusPort);

input (Texts) L

Description *Input* is the standard *TEXT* variable that can be redefined.

Declaration **VAR** input: TEXT;

Usage *input*

input must be imported from the library module *Texts*.

input is used anywhere an input text is needed.

Comments *input* is declared in *Texts* as a variable of type *TEXT* and read-only.

input is always open and is connected to the terminal by default. It operates as if the call

`OpenText(input, CON:)`

has taken place before the execution of every main module. The call

`CloseText(input)`

reestablishes this connection after *input* has been redirected.

READ and **READLN** always use *input* as their default text stream (users may use *input* without knowing it). On occasions when the user wants to explicitly manipulate the *input* text string, he will have to import *input* from *Texts*.

See Also

input

console

Texts

Example

Read aNumber from *input*:

`ReadInt(input, aNumber)`

Insert procedure (Strings) L

Description *Insert* inserts a substring into a string.

Declaration PROCEDURE Insert(substr: ARRAY OF CHAR;
 VAR str: ARRAY OF CHAR; inx: CARDINAL);

Usage Insert(substr, str, at)

substr and *str* must be of type **ARRAY OF CHAR**. *substr* may be a literal, but *str* must be a variable.

at is the index of the destination byte and it must be of type **CARDINAL**.

Comments *Insert* actually inserts; the old data in *str* is not overwritten but is moved over in the string.

See Also Append
 Copy
 Delete
 Length
 Pos
 Strings

Example str := 'This is Modula-2';
 substr := 'Turbo ';
 Insert(substr, str, 8)
 (* str now is 'This is Turbo Modula-2' *)

insertDelete enumerated value (Terminal) L

Description *insertDelete* is the second value of the enumerated type *Terminal.SpecialOps*.

Declaration SpecialOps = (clearEol, insertDelete, highlightNormal);

Usage	IF insertDelete IN available THEN InsertLine END ;
	<i>SpecialOps</i> must be imported from the library module <i>Terminal</i> .
Comments	If this value is a member of the set <i>Terminal.available</i> , then the program may use the procedures <i>InsertLine</i> and <i>DeleteLine</i> to manipulate text on the screen.
	To use this identifier, include <i>SpecialOps</i> in your import list. The identifiers of each of <i>SpecialOps</i> ' values become visible automatically.
See Also	InsertLine DeleteLine available highlightNormal clearEol Terminal

Example Insert and delete a line on the screen:

```
MODULE CheckAvailableOperations;
FROM Terminal IMPORT
  available,SpecialOps,
  InsertLine,DeleteLine,GotoXY;
VAR
  s: ARRAY [0..20] OF CHAR;
BEGIN
  IF insertDelete IN available THEN
    GotoXY(0,10);
    InsertLine;
    WRITELN(s);
    GotoXY(0,10);
    DeleteLine;
  ELSE
    GotoXY(0,10);
    (* Clear the line and move everything down *)
    WRITELN(s);
    GotoXY(0,10);
    (* Move everything back up one line *)
  END ;
END CheckAvailableOperations.
```

InsertLine procedure (Terminal) L

Description *InsertLine* inserts a blank line onto the screen at the current cursor position.

Declaration **PROCEDURE InsertLine;**

Usage **InsertLine**

InsertLine must be imported from the library module *Terminal*.

Comments	This operation is not available on all systems. To determine if the operation is on the current system, inspect the variable <i>Available</i> imported from <i>Terminal</i> . If <i>insertDelete</i> is an element of the set <i>Available</i> , then insertion and deletion of screen lines can be performed.
See Also	DeleteLine GotoXY Terminal
Example	Insert the text » <i>Here I am</i> « at the sixth line from the top of the screen. <code>GotoXY(0,5) InsertLine; GotoXY(0,5); WRITE('Here I am');</code>

INT standard function S

Description	<i>INT</i> converts the argument to type INTEGER.
Usage	<code>Y := INT(X)</code>
	Argument <i>X</i> can be of type INTEGER, CARDINAL, LONGINT, REAL, and LONGREAL.
	The argument must be in the INTEGER range -32768 to 32767.
	Result <i>Y</i> is type INTEGER.
Comments	When you use this function with a REAL argument, the value is truncated (the fractional part removed); thus, a value 12.35 becomes 12.

Example

```
MODULE ConvertToInteger;
VAR
  r: REAL;
  i: INTEGER;
BEGIN
  x := 39.489;
  i := INT(x) * 10;
END ConvertToInteger.
```

INTEGER standard type S

Description INTEGER is a standard type with variables that can assume whole values between -32768 and 32767. In addition, INTEGER can be used as a type-transfer function.

Usage Variable declarations:

```
VAR
  i,j: INTEGER;
```

Type transfer:

```
i := INTEGER(65535)
```

where *i* is an integer. This statement assigns -1 to the integer *i*.

Comments You can use INTEGER variables whenever you know that the possible values are limited to whole numbers. Note that when the possibilities include only positive whole numbers, you are better off using CARDINAL types. In practice, you rarely need to use INTEGER variables.

You may use any arithmetic operator in INTEGER expressions, and INTEGER variables may take part in relational expressions.

The *INT* standard function will convert other numeric data types into INTEGER.

An INTEGER value requires 2 bytes for storage. The bytes are stored as two's complement, with the least-significant byte stored first.

INTEGER is compatible with WORD. Since BYTE is a subrange of WORD, INTEGER is compatible with BYTE as well. However, the exception BoundsError can occur if negative integer is assigned to a byte (or a positive integer greater than 255).

See Also

CARDINAL
LONGINT
WORD

Example

```
MODULE Integer;
VAR
  Deviation: INTEGER;
  Angle: REAL;
BEGIN
  FOR Deviation:=-3 TO 3 DO
    Angle:= 10.0 * Float(Deviation);
  END
END Integer.
```

IntToStr procedure (Convert) L

Description *IntToStr* converts an integer variable to a string.

Declaration PROCEDURE IntToStr(*i*: INTEGER; VAR *s*: ARRAY OF CHAR);

Usage *IntToStr(i, s)*

IntToStr must be imported from the library module *Conv*.
i must be of type INTEGER.

s must be of type ARRAY OF CHAR.

Comments	The string will be right-justified, and the left-most characters padded with blanks.
	If the number is too large to fit into the string, the exception <code>TooLarge</code> will be raised.
See Also	<code>CardToStr</code> <code>Convert</code> <code>Doubles</code> <code>LongToStr</code> <code>RealToStr</code> <code>StrToInt</code>
Example	Put the string '5000' into <i>s</i> :

```
i := 5000;  
IntToStr (i,s);
```

IOTRANSFER procedure (SYSTEM) L

Description	<i>IOTRANSFER</i> sets up an interrupt vector to point to the next line of code and transfers control to another coroutine.
Declaration	PROCEDURE IOTRANSFER(VAR source, dest: PROCESS; n: CARDINAL);
Usage	IOTRANSFER(inthandler, background, intvector) <i>IOTRANSFER</i> must be imported from the library module SYSTEM . <i>inthandler</i> and <i>background</i> must be of type PROCESS imported from SYSTEM . <i>intvector</i> must be of type CARDINAL .

Comments	<i>inhandler</i> and <i>background</i> must have been previously set up with calls to <i>NEWPROCESS</i> .
	<i>intvector</i> must be a legal and available interrupt vector number.
See Also	SYSTEM TRANSFER
Example	IOTRANSFER(KbdIntHandler,MainRoutine,KbdIntVector);

legal variable (Doubles) L

Description	<i>legal</i> monitors the input of a LONGREAL during a call <i>ReadDouble</i> .
Declaration	VAR <i>legal</i> : BOOLEAN;
Usage	finished := legal; <i>finished</i> must be of type BOOLEAN .
Comments	The variable is usually checked after input of a LONGREAL . <i>ReadDouble</i> is the only procedure that can affect it.
See Also	Texts done
Example	MODULE InvalidInput; FROM Doubles IMPORT <i>legal</i> , <i>ReadDouble</i> ; FROM Texts IMPORT <i>input</i> ; VAR <i>d</i> : LONGREAL ; BEGIN REPEAT <i>ReadDouble</i> (<i>input</i> , <i>d</i>); UNTIL <i>legal</i> ; END <i>InvalidInput</i> .

Length procedure (Strings) L

Description	<i>Length</i> finds the current length of a string.
Declaration	PROCEDURE Length(VAR str: ARRAY OF CHAR): CARDINAL;
Usage	<code>size := Length(str)</code>
	<i>Length</i> must be imported from the library module <i>Strings</i> .
	<i>size</i> must be of type CARDINAL.
	<i>str</i> must be a character array variable.
Comments	<i>Length</i> returns the position of the first byte containing a zero, which is defined to be the end of the string.
	<i>size</i> will not be greater than the originally declared size of <i>str</i> .
See Also	Append Copy Delete Insert Pos
Example	Find the length of <i>str</i> : <code>size := Length(str)</code>

LoadError exception (Loader) L/E

Description	LoadError is raised by the module <i>Loader</i> when the loading of an overlay is unsuccessful.
Declaration	EXCEPTION LoadError;

Usage	EXCEPTION LoadError: WRITELN('Overlay Error');
LoadError must be imported from the module <i>Loader</i> .	
Comments	This exception is raised if the overlay file is not found, if there is not enough memory to hold the overlay, or if a version conflict occurs.
See Also	Loader
Example	<pre>MODULE LoadOverlay; FROM Loader IMPORT firstDrive, Call, LoadError; BEGIN firstDrive := 13; (* Search ram disk (M:) first *) Call('overlay.fil') EXCEPTION LoadError: WRITELN(' Load error occurred '); RAISE END LoadOverlay.</pre>

LongMath module L

Description	<i>LongMath</i> provides commonly used double-precision mathematical functions.
Declaration	<pre>DEFINITION MODULE LongMath; PROCEDURE Sqrt (x:LONGREAL):LONGREAL; PROCEDURE Exp (x:LONGREAL):LONGREAL; PROCEDURE Ln (x:LONGREAL):LONGREAL; PROCEDURE Cos (x:LONGREAL):LONGREAL; PROCEDURE Arctan (x:LONGREAL):LONGREAL; PROCEDURE Entier (x:LONGREAL):LONGINT; EXCEPTION ArgumentError; END LongMath.</pre>

Comments These routines are identical to the routines for single-precision real numbers.

The ArgumentError exception is raised when an argument is outside the permitted range. The following ranges apply:

Sqrt Argument must be positive or zero

Exp Argument must be less than 710.47

Ln Argument must be positive (greater than zero)

If you want to use the same routine from *MathLib* and *LongMath*, the statements

```
FROM MathLib IMPORT Sqrt;  
FROM LongMath IMPORT Sqrt;
```

will lead to a conflict. Thus, qualified import must be used.

```
IMPORT MathLib,LongMath;
```

```
BEGIN
```

```
  dSqr := LongMath.Sqrt( arg );  
  rSqr := MathLib.Sqrt( arg );
```

LONGREAL standard type S

Description LONGREAL is a standard type with variables that can assume any value between -3.5953862697246D+308 and +3.5953862697246D+308.

Usage **VAR** x,y: LONGREAL;

Comments

LONGREALs are particularly useful in arithmetic expressions that concern a small difference between large numbers or a sum that results from many (a 1000 or so) additions and subtractions. In both cases, the round-off error tends to accumulate; however, it can be kept smaller by using a LONGREAL. Don't use them indiscriminately, though, since your program will pay the price in slower execution times. You may use any arithmetic operator in LONGREAL expressions, and LONGREAL variables may take part in relational expressions.

The *DOUBLE* standard function will convert other numeric data types into LONGREAL.

String and I/O support for LONGREALs are provided by the library module *Doubles*.

Mathematical support similar to *MathLib* for REALs is provided by the library module *LongMath*.

Turbo Modula-2 uses the IEEE 8-byte, double-precision real-number format; that is, in the order of their significance, 1 sign bit, 11 exponent bits, and 52 mantissa bits:

63	62		52	51		0
	s		exponent		mantissa	

The mantissa is in binary format but is offset by 1.0, representing only those digits to the right of the decimal point.

The bytes are stored with the least-significant bytes in the lower-numbered addresses.

If the exponent is zero, the floating-point value is considered zero.

Sample hex real numbers:

3FF0000000000000H = 1.0
4000000000000000H = 2.0

See the example in *WORD* for printing out other values.

Example

```
VAR
  SquareArea, SquareSide: LONGREAL;
BEGIN
  SquareSide:= 4.21D0;
  SquareArea:= SquareSide * SquareSide;
```

Ln function (MathLib, LongMath) L

Description Ln returns the natural logarithm of X.

Declaration PROCEDURE Ln(x: REAL): REAL;
PROCEDURE Ln(x: LONGREAL): LONGREAL;

Usage Y:= Ln(X)

Argument X and result Y are either both REAL or both LONGREAL.

The argument X must be positive.

Comments You must always make sure that the argument X is a positive value; otherwise, the ArgumentError exception is raised. The inverse of the logarithmic function is the natural exponential function *Exp*.

See Also Exp
LONGMATH
MATHLIB

Example

```
MODULE LnTest;
FROM MathLib IMPORT Ln;
VAR
  Argument, Result: REAL;
BEGIN
  IF Argument > 0.0 THEN
    Result:= Ln(Argument)
  END
END LnTest;
```

Loader module L

Description *Loader* is a library module that handles overlays.

Declaration **DEFINITION MODULE** Loader;

```
VAR
  firstDrive: [0..15];

PROCEDURE Call(modName: ARRAY OF CHAR);

EXCEPTION LoadError;

END Loader.
```

Comments *Loader* is used by a program to explicitly overlay certain modules. The overlay code may not export anything to the caller; however, the overlay code may communicate with the caller by importing from the caller a buffer where data and messages can be left.

See Also Refer to »The Loader Module« in Chapter 11.

LONG standard function S

Description *LONG* converts the argument to type LONGINT.

Usage *Y := LONG(X)*

Argument *X* can be of type INTEGER, CARDINAL, LONGINT, REAL, or LONGREAL.

The argument must be in the range from -(2^{31} -1) to 2^{31} -1.

The result *Y* is type LONGINT.

Comments When you use this function with a REAL argument, the value is truncated (the fractional part removed); thus, the value 12.35 becomes 12.

Do not confuse this function with the LONGINT-type transfer function, which transfers data without converting it.

See Also FLOAT
INT
TRUNC

Example

```
VAR
  Convert: LONGINT;
BEGIN
  Convert:= 2000000L * LONG(-12);
(* Convert is -24000000L *)
END
```

LONGINT standard type S

Description	LONGINT is a standard type with variables that can assume whole values between -2147483648 to 2147483647.
Comments	You can use LONGINT types when you need a larger range of whole numbers than that offered by the CARDINAL and INTEGER types. You may use any arithmetic operator in LONGINT expressions, and LONGINT variables may take part in relational expressions.
	The LONG standard function will convert other numeric data to type LONGINT. Note that even INTEGERtypes must be converted.
	A LONGINT value requires 4 bytes for storage. The bytes are stored as two's complement, with the least-significant byte stored first.
	LONGINT is compatible with ARRAY OF WORD, but not with WORD or BYTE.
See Also	CARDINAL INTEGER WORD
Example	MODULE Longs; CONST N = 1000000L; VAR BigNumber:LONGINT; BEGIN BigNumber:= N; BigNumber:= 2L * BigNumber; END Longs.

LongMath module

Description Provides commonly used mathematical functions with double precision accuracy.

Declaration **DEFINITION MODULE** LongMath;

```
PROCEDURE Sqrt    (x:LONGREAL): LONGREAL;
PROCEDURE Exp     (x:LONGREAL): LONGREAL;
PROCEDURE Ln      (x:LONGREAL): LONGREAL;
PROCEDURE Sin     (x:LONGREAL): LONGREAL;
PROCEDURE Cos     (x:LONGREAL): LONGREAL;

PROCEDURE Arctan  (x:LONGREAL): LONGREAL;
PROCEDURE Entier  (x:LONGREAL): LONGINT;

EXCEPTION ArgumentError;

END LongMath.
```

For a description of each of these procedures and exceptions, see Chapter 11 or the appropriate reference in this chapter.

Comments The *LongMath* module exports a number of mathematical procedures and functions. You can import any of the procedure and exception identifiers into your own modules, which allows you to use the predefined facilities.

The *ArgumentError* exception is raised when an argument is outside the permitted range. The following ranges apply:

Sqrt Argument must be positive or zero.

Exp Argument must be less than 710.475D00

Ln Argument must be positive (greater than zero).

Example If you want to call the *Cos* and *Sin* functions from a module called *Problem*, you would start the module with:

```
MODULE Problem;
FROM LongMath IMPORT Cos,Sin;
```

If *Y,Z*, *Angle*, and *Radians* are defined as LONGREAL variables then the library functions can then be used normally

```
Y:= Cos(Angle);
Z:= Sin(Radians);
```

LongToStr procedure (Convert) L

Description *LongToStr* converts a long integer variable to a string.

Declaration PROCEDURE LongToStr(*l*: LONGINT; VAR *s*: ARRAY OF CHAR);

Usage LongToStr(*l*, *s*)

LongToStr must be imported from the library module *Conv*.

l must be of type LONGINT.

s must be a character array variable.

Comments The string will be right-justified, and the left-most character will be padded out with blanks.

If the number is too large to fit into the string, the exception *TooLarge* will be raised.

See Also Doubles
InToStr
CardToStr
RealToStr
Strings
StrToInt

Example Put the string '300000' into s:

```
VAR
  s : ARRAY [0..6] OF CHAR;
  l : LONGINT;
BEGIN
  l := 300000;
  LongToStr(l,s);
END
```

LOOP statement R

Description **LOOP** repeatedly executes a statement sequence until terminated by an **EXIT** statement.

Usage **LOOP** StatementSequence **END**

StatementSequence = Statement {; Statement}

Comments There are two main reasons for using a **LOOP** statement: (1) when programming a continually cycling process, such as many real-time systems and concurrent processes, or (2) when the **FOR**, **REPEAT** and **WHILE** statements are unsuitable because loop termination can only be determined in the middle of the loop.

Termination of the loop requires an **EXIT** statement.

See Also **EXIT**

Example **BEGIN**
LOOP
 IF Found THEN **EXIT** **END**
END ;

MARK procedure (STORAGE) L

Description	LOOP marks the beginning of a block of dynamic variable.
Declaration	PROCEDURE MARK(VAR a: ADDRESS);
Usage	MARK(a)
	<i>MARK</i> must be imported from the library module STORAGE .
	<i>a</i> must be compatible with type ADDRESS (any pointer).
Comments	<i>MARK</i> and <i>RELEASE</i> provide an alternate method of handling heap data. The use of these procedures allows the heap to be treated in a stack-like manner.
	A call to <i>MARK</i> creates a new heap that consists of the space between the previous top of Heap and the stack pointer. All variables allocated on the heap prior to the call to <i>MARK</i> are no longer accessible until a <i>RELEASE</i> with the same pointer is called.
See Also	DISPOSE NEW RELEASE
Example	Allocate a mess of variables and then throw them away: MARK(MessStart); NEW(Mess1); NEW(Mess2); NEW(Mess3); RELEASE(MessStart);

MathLib module L

Description *MathLib* provides commonly used mathematical functions.

Declaration

```
DEFINITION MODULE MathLib;
```

```
PROCEDURE Sqrt(x:REAL):REAL;
PROCEDURE Exp(x:REAL):REAL;
PROCEDURE Ln(x:REAL):REAL;
PROCEDURE Sin(x:REAL):REAL;
PROCEDURE Cos(x:REAL):REAL;
PROCEDURE Arctan(x:REAL):REAL;
PROCEDURE Entier(x:REAL):INTEGER;
```

```
PROCEDURE Randomize(n:CARDINAL);
PROCEDURE Random( ):REAL;
```

```
EXCEPTION ArgumentError;
```

```
END MathLib.
```

For a description of each of these procedures and exceptions, see Chapter 11 or the appropriate reference in this chapter.

Comments The *MathLib* module exports a number of mathematical procedures and functions. You can import any of the procedure and exception identifiers into your own modules, allowing you to use the predefined facilities.

The exception *ArgumentError* is raised when an argument is outside the permitted range. The following ranges apply:

Sqrt Argument must be positive or zero.

Exp Argument must be less than 87.4.

Ln Argument must be positive (greater than zero).

Example

If you want to call the *Cos* and *Sin* functions from a called *Problem*, you would start the module with the following:

```
MODULE Problem;
FROM MathLib IMPORT Cos,Sin;
```

The predefined functions can then be used normally:

```
Y:= Cos(Angle);
Z:= Sin(Radians);
```

MAX standard function S**Description**

MAX returns the largest element of type *T*.

Usage

```
Y:= MAX(T);
```

Argument *T* is of type CARDINAL, INTEGER, BOOLEAN, CHAR, enumeration, LONGINT, REAL, or LONGREAL, or any unstructured scalar type.

Result *Y* is of the same type.

Comments

This function allows you to find the last element of an unstructured type. Note that the argument is the type of a variable of the type.

The maximum values of the predefined types are:

CARDINAL	65535
INTEGER	32767
BOOLEAN	TRUE
CHAR	377C
REAL	2147483647
LONGINT	3.5953862697246D+308
LONGREAL	6.80565E+38

The smallest element of a type is returned by the **MIN** function.

module
l owing:
like so:
()

Example

```

MODULE Max;
TYPE
  Weather = (Clear,Rain,Wind,Snow);
VAR
  MaxWeather: Weather;
  MaxCardinal: CARDINAL
BEGIN
  MaxWeather:= MAX(Weather);
  (* MaxWeather is value Snow *)
  MaxCardinal:= MAX(CARDINAL);
  (* MaxCardinal is value 65535 *)
END Max.
```

MIN standard function S

Description *MIN* returns the smallest element of type *T*.

Usage *Y:= MIN(T);*

Argument *T* is of type CARDINAL, INTEGER, BOOLEAN, CHAR, enumeration, LONGINT, REAL, or LONGREAL; in short, any unstructured type.

Result *Y* is of the same type.

Comments This function allows you to find the first element in any unstructured type. Note that the argument is the type and not a variable of the type.

The minimum values of the predefined types are

CARDINAL	0
INTEGER	-32768
BOOLEAN	FALSE
CHAR	0C
REAL	-6.80565E+38
LONGINT	-2147483648
LONGREAL	-3.5953862697246D+308

The largest element of a type is returned by the **MAX** function.

Example

```
MODULE Min;
TYPE
  Weather = (Clear,Rain,Wind,Snow);
VAR
  MinWeather: Weather;
  MinCardinal:CARDINAL
BEGIN
  MinWeather:= MIN(Weather);
  (* MinWeather is value Clear *)
  MinCardinal:= MIN(CARDINAL);
  (* MinCardinal is value 0 *)
END Min.
```

MOD standard operator R

Description MOD is an integer modulus operator; it is a binary operator that works on operands of the same type. The MOD operator returns the remainder of integer division. The operands may be both CARDINAL, both INTEGER, or both LONGINT.

Usage

i := i MOD j;

where *i* and *j* are of the same type.

The result of the expression is the same type as the operands.

Comments

The modulus operator can be used to insure that a variable remains within a certain defined range.

See Also

DIV

Example **MODULE** Modulus;
 VAR
 i: INTEGER;
 BEGIN
 LOOP
 i := (i + 1) **MOD** 100; (* i will range from 0 to 99 *)
 END
 END Modulus.

(**MODULE** declaration R

Description **MODULE** serves to declare the beginning of a module.

Usage **MODULE** moduleName

or

DEFINITION MODULE libraryName

or

IMPLEMENTATION MODULE libraryName

Comments The first form declares a main module, which is a stand-alone program. It can also start the beginning of a local module.

A local module limits visibility of global objects (variables, types, constants, procedures). Thus, an object buried inside a local module can retain its value throughout the program's lifetime, while simultaneously remain safe from outside tampering.

The second and third forms are for **libraries**, separately compiled sets of frequently used routines.

See Also **DEFINITION**
IMPLEMENTATION

Example

The following file is an example of a main module in 'MAIN.MOD':

```
MODULE main;
FROM SimpleMath IMPORT add;
BEGIN
  WRITELN(' 2 + 2 = ',add(2,2));
END main.
```

The next example is a local module in a main module:

```
MODULE Main;

  MODULE local;
  EXPORT p1;

  PROCEDURE p1;
  BEGIN
  END p1;

  PROCEDURE p2;
  BEGIN
  END p2;

  END local;

BEGIN (* p1 is visible here, but not p2 *)
END Main.
```

NEW standard procedure (STORAGE) S**Description**

NEW allocates a dynamic variable from the heap.

Usage **NEW(p);**

or

NEW(p, tag1, tag2, ..., tagn);

p must be of type **POINTER**.

ALLOCATE must be imported from the pseudomodule **STORAGE**.

tag1, ..., tagn are allowed in the case where *p* points to a variant record, and the tags are defined in that record.

Comments **NEW** is translated into an appropriate call to **ALLOCATE**.

Turbo Modula-2 expects **ALLOCATE** to be defined the way it is in the pseudomodule **STORAGE**; however, you may substitute your own allocation scheme by importing **ALLOCATE** from your own module instead of **STORAGE**.

See Also **ALLOCATE**
DISPOSE
STORAGE

Example Allocate a new Name:

```
MODULE Allocate;
FROM STORAGE IMPORT ALLOCATE;
VAR
  n: POINTER TO Name;
BEGIN
  NEW(n);
END Allocate.
```

The **NEW(n)** statement would be translated into a call of the form **ALLOCATE(n, TSIZE(Name))**.

NEWPROCESS procedure (SYSTEM) L

Description	NEWPROCESS sets up an independent process (or coroutine).
Declaration	PROCEDURE NEWPROCESS(p: PROC; a: ADDRESS; n: CARDINAL; VAR q: PROCESS);
Usage	NEWPROCESS (newProc, wrkAdr, wrkSize, P);
	NEWPROCESS must be imported from the pseudomodule SYSTEM .
	<i>newProc</i> must be of the standard type PROC (a parameterless procedure).
	<i>wrkAdr</i> must be compatible with type ADDRESS (any pointer).
	<i>wrkSize</i> must be of type CARDINAL . It specifies the number of bytes of storage.
	<i>P</i> must be of type PROCESS imported from SYSTEM .
Comments	The procedure <i>newProc</i> contains the code for the new process. The workspace must be big enough to accommodate any local stack and heap activity; 100 words is a good size to start with. Local stack activity consists of local variables, procedure parameters, and some space for working out expressions. Local heap activity consists of the user's calls to NEW and ALLOCATE and any file control blocks that are active. If the allocated workspace is not enough, the exception OutOfMemory will be raised. NEWPROCESS allocates space and sets the corresponding PROCESS variable, but it does not actually start up the routines-- TRANSFER accomplishes that.

The procedure *StartProcess* in *Processes* is implemented using this routine.

See Also Processes
 PROCESS
 PROC
 TRANSFER
 SYSTEM

Example Start up a driver called *Driver*:

```
MODULE NewProcess;
FROM SYSTEM IMPORT ADR, PROCESS, NEWPROCESS,
    TRANSFER;
VAR
    work: ARRAY [ 0..100 ] OF WORD;
    M,P: PROCESS;

    PROCEDURE Driver;
    BEGIN
        (* Driver's code *)
    END Driver;

    BEGIN
        NEWPROCESS(Driver, ADR(work), SIZE(work), P);
        TRANSFER(M,P); (* Driver now has control *)
    END NewProcess.
```

NextPos procedure (Files) L

Description *NextPos* returns the current byte position in a file.

Declaration PROCEDURE NextPos(f: FILE): LONGINT;

Usage

```
n := NextPos(f);
```

NextPos must be imported from the library module *Files*.

n must be of type LONGINT.

f must be of type FILE imported from *Files*.

Comments

The result of *NextPos* is a LONGINT because file can be larger than MAX(CARDINAL).

See Also

FileSize

SetPos

Example

Keep track of the percentage of file you have completed:

```
MODULE NextPosPercent;
FROM Files IMPORT FILE, Open, NextPos, ReadRec;
VAR
  f:File;
  totsize: LONGINT;
  wreck: SomeRecord;
BEGIN
  IF NOT Open(f, "reckfile") THEN HALT END ;
  totsize := FileSize(f);
  WHILE NOT EOF(f) DO
    ReadRec(f, wreck);
    n := (100L*NextPos(f)) / totsize;
    WRITE('Finished ',n,' % of the file ')
    (* Process wreck *)
  END
END NextPosPercent.
```

NIL standard value S

Description	<i>NIL</i> designates an unused pointer variable.
Usage	<code>BufferPtr := NIL;</code> <i>BufferPtr</i> is any pointer type.
	<i>NIL</i> is compatible with all pointer variables.
Comments	The actual value of <i>NIL</i> is zero; thus, <code>WRITE(CARDINAL(NIL))</code> will print a 0.
See Also	POINTER
Example	<pre>MODULE NilList; FROM STORAGE IMPORT ALLOCATE; TYPE NodePtr = POINTER TO Node; Node = RECORD element: CARDINAL; Next: NodePtr; END ; VAR l: NodePtr; BEGIN NEW(l); l^.Next := NIL; END NilList.</pre>

Normal procedure (Terminal) L

Description	<i>Normal</i> turns off highlighting.
Declaration	PROCEDURE Normal;
Usage	<code>Normal;</code> <i>Normal</i> must be imported from the library module <i>Terminal</i> .

Comments The highlighting will only work if Turbo Modula-2 has been properly installed. If highlighting has not been turned on, then the effect of *Normal* is undefined (it depends on your terminal, but it is probably harmless).

Example Highlight the word 'high':

```
WRITE( 'This is a example of ');
Highlight;
WRITE( 'High' );
Normal;
WRITE( 'lighting' );
```

NoTrailer procedure (Files) L

Description *NoTrailer* turns off the special Turbo Modula-2 end-of-file handling.

Declaration **PROCEDURE** NoTrailer(*f*: FILE);

Usage NoTrailer(*f*);

NoTrailer must be imported from the library module *Files*.

f must be of type *FILE* imported from *Files*.

Comments Turbo Modula-2 handles the last byte of the last record of a CP/M file so that a program can determine exactly how many bytes are in a file.

The last byte of the last record contains the length of the last record plus 128. Thus, if the last byte is 127 or less, then the record is completely filled; otherwise, the number of bytes in the last record is found by subtracting 128 from the last byte. This is how the procedure *FileSize* works.

If *NoTrailer* is used on a file, then *FileSize* will always return a multiple of 128.

NoTrailer should be used immediately after *Open* when reading; it can be used immediately after *Create* as well.

In general, this function should only be used if you are having trouble with file communication involving a non- Turbo Modula-2 program.

See Also

Files
FileSize

Example

Declare a file to have no trailer:

```
IF Open(f,filename) THEN
  NoTrailer(f);
  (* process the file *)
END ;
```

numCols variable (Terminal) L**Description**

numCols holds the number of columns available on the currently installed terminal.

Declaration

```
VAR
  numCols: CARDINAL;
```

Usage

```
IF CurCol > numCols THEN CurCol := 0 END ;
```

numCols must be imported from the module *Terminal*.

Comments

This variable is useful for writing terminal-independent programs. The value of *numCols* is set by the installation program INSTM2.

See Also

numRows
Terminal

Example

```
MODULE ColumnNumber;
FROM Terminal IMPORT numCols;
VAR
  I: CARDINAL;
BEGIN
  I := 0;
  WHILE I < numCols DO WRITE(' - ') END ;
END ColumnNumber.
```

numRows variable (Terminal) L

Description	<i>numRows</i> holds the number of rows available on the current installed terminal.
Declaration	VAR <i>numRows</i> : CARDINAL;
Usage	IF CurRow1 > numRows THEN CurRow := 0 END ; <i>numRows</i> must be imported from the module <i>Terminal</i> .
Comments	This variable is useful for writing terminal-independent programs. The value of <i>numRows</i> is set by the installation program INSTM2.
See Also	<i>numCols</i> <i>Terminal</i>
Example	<pre>MODULE RowNumber; FROM Terminal IMPORT numRows; VAR I: CARDINAL; BEGIN I := 0; WHILE I < numRows DO WRITELN(' !') END ; END RowNumber.</pre>

ODD standard function S

Description *ODD* returns TRUE if the ordinal value of *X* is odd.

Usage *Y := ODD(X);*

Argument *X* is of type CARDINAL or INTEGER.

Result *Y* is of type BOOLEAN.

Comments This function allows you to find out if a value is odd.

Note that for BOOLEAN types, the first value is FALSE with an ordinal value of 0.

Example **MODULE** Odd;

TYPE

 Shift = (Early,Late,Night);

VAR

 TimeSheet : ARRAY Shift OF CARDINAL;

 ShiftToday : Shift;

 OddNumber : BOOLEAN;

BEGIN

 TimeSheet[Early]:= 11;

 OddNumber:= ODD(TimeSheet[Early]);

 (* OddNumber is TRUE*)

END Odd.

Open procedure (Files) L

Description *Open* opens a binary disk file for input or output.

Declaration **PROCEDURE** Open(**VAR** f: FILE; name: ARRAY OF CHAR):
 BOOLEAN;

Usage	<code>okay := Open(f, name);</code>
	<i>Open</i> must be imported from the library module <i>Files</i> .
	<i>okay</i> must be declared as type BOOLEAN.
	<i>f</i> must be declared as type FILE imported from <i>Files</i> .
	<i>name</i> must be declared as ARRAY OF CHAR.
Comments	<p><i>f</i> is the internal file identifier. Further operations on this file must specify <i>f</i>.</p> <p><i>name</i> must represent a legal CP/M file name that is on the disk you are using or a standard CP/M device like RDR: or CON:.</p> <p><i>okay</i> is returned TRUE when the file can be opened; it is FALSE if the file could not be found.</p> <p>Use <i>Create</i> to start a new file.</p> <p>Readable files are better handled using routines from the module <i>Texts</i>.</p>
See Also	Close Create Files
Example	<p>Open and close a disk file 'SCRATCH.SCR':</p> <pre> MODULE openandclose; FROM Files IMPORT FILE, Open, Close; VAR f : FILE; BEGIN IF Open(f, 'SCRATCH.SCR') THEN Close(f) END END openandclose. </pre>

OpenInput Procedure (InOut) L

Description *OpenInput* asks the user for a disk file name, then opens it as standard input.

Declaration **PROCEDURE** OpenInput(*defext*: **ARRAY OF CHAR**);

Usage OpenInput(*ext*);

OpenInput must be imported from the library module *InOut*.

ext must be a character array.

Comments Notice that calling this procedure causes the program to halt and ask the user for a file name. The *ext* parameter will be used only if the user does not specify an extension.

Example **MODULE** Input;
 FROM InOut **IMPORT**
 OpenInput, CloseInput, ReadString,
 WriteString, WriteLn;
 VAR
 s: **ARRAY [0..255] OF CHAR;
 BEGIN
 (* Prompt user for filename *)
 (* Use default extention MOD *)
 OpenInput("MOD");
 (* Read a string from the file *)
 ReadString(*s*);
 WriteString(*s*);
 (* Display it *)
 WriteLn;
 CloseInput;
 (* Sever the link to the file *)
 (* Now input comes from the console *)
 ReadString(*s*);
 WriteString(*s*);
 (* Display it *)
 WriteLn;
 END Input.**

OpenOutput procedure (InOut) L

Description	<i>OpenOutput</i> asks the user for a disk file name, then opens it as the standard output.
Declaration	PROCEDURE OpenOutput(defext: ARRAY OF CHAR);
Usage	OpenOutput(ext); <i>OpenOutput</i> must be imported from the library module <i>InOut</i> . <i>ext</i> must be a character array.
Comments	Notice that calling this procedure causes the program to halt and ask the user for a file name. The <i>ext</i> parameter will be used only if the user does not specify an extension.
Example	MODULE Output; FROM InOut IMPORT OpenOutput, CloseOutput, ReadString, WriteString, WriteLn; VAR s: ARRAY [0..255] OF CHAR ; BEGIN WriteString('Enter output file and then the data:'); WriteLn; OpenOutput("MOD"); (* Read data from the console *) ReadString(s); (* Place it in the file *) WriteString(s); WriteLn; CloseOutput; (* Sever the link to the file *) WriteString('The file has been closed.');// WriteLn; END Output.

OpenText procedure (Texts) L

Description *OpenText* associates a *TEXT* with an external medium and opens it for reading.

Declaration **PROCEDURE** OpenText(**VAR** *t*: *TEXT*;
 name: **ARRAY OF CHAR**): **BOOLEAN**;

Usage *okay* := OpenText(*t*, *fname*);

OpenText must be imported from the library module *Texts*.

t must be declared as type *TEXT* imported from *Texts*.

fname must be a character array.

okay must be declared as type **BOOLEAN**.

Comments *fname* must either be a valid CP/M file name and be present on the appropriate disk or represent an existing CP/M device like 'CON:'.

okay is returned FALSE if the file is not found; otherwise, the file is successfully opened and *okay* is returned TRUE.

t can only be used for input.

See Also CreateText
 TEXT
 Texts

Example Open a text file 'MYTEXT.TXT' and then close it:

```
MODULE openandclose;
FROM Texts IMPORT TEXT, OpenText, CloseText;
VAR
  t : TEXT;
BEGIN
  IF OpenText(t, 'MYTEXT.TXT') THEN
    CloseText(t)
  END
END openAndClose.
```

OpSet type (Terminal) L

Description *OpSet* holds the set of screen operations available on the currently installed terminal.

Declaration **TYPE**

```
SpecialOps = (clearEol, insertDelete,
highlightnor mal);
OpSet : SET OF SpecialOps;
```

Usage **VAR**

```
available : OpSet;
```

available is a predefined variable declared as type *OpSet*. It must be imported from the module *Terminal*.

Comments

This variable is useful for writing terminal-independent programs. The value of *available* is set by the installation program INSTM2.

See Also

SpecialOps
available
Terminal

Example

```
MODULE TermStuff;
FROM Terminal IMPORT
  available, SpecialOps, Highlight, Normal;
BEGIN
  IF highlightNormal IN available THEN
    Highlight
  ELSE
    WRITE(' ');
  END ;
  WRITELN('This is highlighted text!');
  IF highlightNormal IN available THEN
    Normal
  ELSE
    WRITE(' ');
  END ;
END TermStuff.
```

OR R

Description	OR is a binary logical operator resulting in a Boolean value.
Usage	<pre>IF InputError OR TimeIsUp THEN Beep END ;</pre> OR is a reserved word. <i>InputError</i> and <i>TimeIsUp</i> are Boolean expressions.
Comments	OR is used to concatenate two logical expressions. An OR expression is TRUE if at least one of its subexpressions is TRUE. Note that the second subexpression in an OR expression is only evaluated if the first subexpression is FALSE. This is called short-circuit evaluation.
See Also	AND NOT

Example REPEAT statement using OR to test either of two terminating conditions.

```
MODULE Delay;
FROM Terminal IMPORT BusyRead;
VAR
  ch: CHAR;
  x : CARDINAL;
BEGIN
  x := 0;
REPEAT
  INC(x);
  BusyRead(ch);
UNTIL ch#0C OR x>1000;
END Delay.
```

ORD standard function S

Description *ORD* returns the ordinal number of *X*.

Usage Y:= ORD(X);

Argument *X* is of type enumeration, INTEGER, CARDINAL, BOOLEAN, or CHAR.

Result is of type CARDINAL.

Comments Ordinal types have an ordered sequence of values, with position in the sequence being the ordinal number; thus, first value has an ordinal of 0, the next has ordinal of 1, so forth.

Note that INTEGER values are written in two's complement format so that the lower ordinals correspond to positive values and higher ordinals to negative values.

Ordinals	0 to 32767	Values	0 to 32
Ordinals	32768 to 65535	Values	-32768 to -0

For BOOLEAN types, the first value is FALSE with ordinal 0.

The inverse operation from ordinal to value uses the *VAL* function, or you can use the *CHR* function for character variables.

Example

```
TYPE Direction: {Forward,Backward,Up,Down};
VAR Move: Direction;
      Value: INTEGER
      Ordinal: CARDINAL

BEGIN
  Move:= Up;
  Ordinal:= ORD(Move);
(* Ordinal is value 2 *)
  Value:= 10; Ordinal:= ORD(Value);
(* Ordinal is value 10 *)

  Ordinal:= ORD("?");
(* Ordinal is value 63 *)
```

OUT procedure (SYSTEM) L

Description *OUT* outputs a word or a byte to a given I/O port.

Declaration PROCEDURE OUT(*port*,*data*: WORD);

Usage OUT(*port*, *outWord*);

OUT must be imported from the library module *SYSTEM*.

outWord must be compatible with type *WORD* imported from *SYSTEM*.

port can be any scalar type compatible with *WORD*.

Comments The value given by *outWord* is output to the port whose number is given by *port*.

See Also	INP SYSTEM
Example	Output a byte to port 23 hex: OUT(023H, outWord);
outName variable (ComLine) L	
Description	<i>outName</i> is a string variable that contains a string found at the output redirection symbol (>) on the command line.
Declaration	VAR outName: ARRAY [0..19] OF CHAR ;
Usage	WRITE('Output redirected to ',outName);
Comments	The redirection symbol and <i>outName</i> will not appear in text <i>commandLine</i> . <i>outName</i> contains the standard output string »CON:<« if redirection argument has been found on the command line.
See Also	ComLine inName RedirectOutput
Example	MODULE RedirectOutput; FROM ComLine IMPORT outName, RedirectOutput; FROM Texts IMPORT console; BEGIN RedirectOutput; IF outName < > "CON:" THEN WRITE(console,'Output redirected to ',outName) END END RedirectOutput

OUTOFGMEMORY exception (SYSTEM) L/E

Description	OUTOFGMEMORY is raised by the module SYSTEM when an allocation operation makes the stack and heap collide.
Declaration	EXCEPTION OUTOFGMEMORY;
Usage	EXCEPTION OUTOFGMEMORY: WRITELN('Out of memory.');
	END
	OUTOFGMEMORY must be imported from the module SYSTEM .
Comments	This exception may occur during explicit allocation or procedure calls.
See Also	SYSTEM BADOVERLAY OVERFLOW REALOVERFLOW
Example	MODULE OutOfMemory; FROM SYSTEM IMPORT OUTOFGMEMORY; FROM STORAGE IMPORT ALLOCATE, FREEMEM; VAR p: POINTER TO CHAR ; BEGIN ALLOCATE(p, FREEMEM()+100); EXCEPTION OUTOFGMEMORY: WRITELN('Out of memory.');
	END OutOfMemory.

output (Texts) L

Description	<i>output</i> is a standard <i>TEXT</i> variable that can be redirected.
Declaration	VAR output: TEXT;
Usage	<i>output</i>
	<i>output</i> must be imported from the library module <i>Texts</i> and can be used anywhere a <i>TEXT</i> can be used.
Comments	<i>output</i> is declared in <i>Texts</i> as a variable of type <i>TEXT</i> and write-only.
	<i>output</i> is always open and by default connected to the terminal. It is as if the call <i>CreateText(output, »CON:«)</i> has taken place before the execution of every main module. The call <i>CloseText(output)</i> reestablishes this connection.
	<i>WRITE</i> and <i>WRITELN</i> always use <i>output</i> as their default text stream.
See Also	console input <i>Texts</i>
Example	Write the number 5A5A (=23130 decimal) to the output:
	MODULE StandardOut; FROM Texts IMPORT output; BEGIN WRITELN(output,5A5AH); END StandardOut.

OVERFLOW exception (SYSTEM) L/E

Description **OVERFLOW** is raised by the module **SYSTEM** when an integer or cardinal expression results in a value that is too large or too small to be represented by the result type.

Declaration **EXCEPTION**
 OVERFLOW;

Usage **EXCEPTION**
 OVERFLOW:
 WRITELN('Overflow.');
END

OVERFLOW must be imported from the module **SYSTEM**.

Comments This exception occurs only during evaluation of integer expressions. Real expressions generate the exception **REALOVERFLOW**.

See Also **SYSTEM**
 BADOVERLAY
 OUTOFGMEMORY
 REALOVERFLOW

Example **MODULE** Overflow;
 FROM SYSTEM IMPORT OVERFLOW;
 VAR
 r: REAL;
 BEGIN
 r := MAX(REAL) * 2.2;
 EXCEPTION
 OVERFLOW:
 WRITELN('Overflow');
 END Overflow.

POINTER type R

Description A pointer references a dynamic variable that can be created, operated on, and discarded during program execution.

Usage **Pointer Type = POINTER OF Referenced Type;**

Referenced Type can be any type.

Comments If the precise data structure can only be defined while running a program, you could use pointers to enable the program to create the data structure. For example, suppose you want to record the name and job of each employee in a company. Assuming there will never be more than 100 employees, you could declare a fixed data structure of 100 records, or you could use pointer variables. The program creates new employee records as you enter data at the keyboard.

A pointer is bound to the referenced type and cannot point to any other type. For example, if a pointer type is declared as

ArrayPointer = POINTER TO ARRAY [1..10] OF CHAR

then a pointer of type ArrayPointer can only create, reference, and discard arrays of 10 CHAR components.

The standard procedures NEW and DISPOSE create and discard referenced variables.

It should be noted that NEW and DISPOSE are translated into calls to the standard procedures ALLOCATE and DEALLOCATE, which must be imported from the module STORAGE.

The referenced variable is designated by the ^ symbol, the dereferencing operator.

A pointer's reference can be assigned to another identical pointer; in fact, several pointers can reference the same variable.

A pointer can also be given the value NIL, which means the pointer does not reference anything; this is not the same as discarding variables.

A pointer requires 2 bytes of memory. The value NIL corresponds to 2 zero bytes.

Dynamic variables are stored in the heap area that is handled by procedures in module STORAGE.

Example

```
MODULE Company;
FROM STORAGE IMPORT ALLOCATE,DEALLOCATE;

TYPE
PersonPointer = POINTER TO Person;
Person = RECORD
    Name: ARRAY [1..20] OF CHAR;
    Job: ARRAY [1..20] OF CHAR;
    Next: PersonPointer;
END ;

VAR
    FirstPerson,LastPerson,NewPerson:
    PersonPointer;
    Name: ARRAY [1..20] OF CHAR;

BEGIN
    FirstPerson:= NIL;
REPEAT
    WRITE('Enter name: ');
    READLN(Name);

    IF Name < > '' THEN
        (* Create a person record if name entered *)
        NEW(NewPerson);
        NewPerson^.Name:= Name;
        WRITE('Enter profession: ');
        READLN(NewPerson^.Job);
        WRITELN;
```

```

IF FirstPerson = NIL THEN
  FirstPerson:= NewPerson
  (* FirstPerson references first in list *)
  ELSE
    LastPerson ^ .Next:= NewPerson
  (* Reference to next in list *)
    END ;
    LastPerson:= NewPerson;
  (* LastPerson references last in list *)
    LastPerson ^ .Next:= NIL
  END
  UNTIL Name = '';
  WRITELN;
  (* Write all records on screen *)
  WHILE FirstPerson < > NIL DO
    WITH FirstPerson ^ DO
      WRITELN(Name,' is a ',Job);
      FirstPerson:= Next;
    END
  END
END Company.

```

Pos procedure (Strings) L

Description *Pos* finds the position of a substring in a string.

Declaration **PROCEDURE** Pos(substr,str: **ARRAY OF CHAR**): CARDINAL

Usage	<code>p := Pos(substr, str);</code>
	<i>Pos</i> must be imported from the library module <i>Strings</i> .
	<i>substr</i> and <i>str</i> must be character arrays.
	<i>p</i> must be of type CARDINAL.
Comments	The first character has the position 0, the second 1, and so on. If the substring is not found, the value <i>HIGH(str)+1</i> is returned.
See Also	Copy Delete Insert Length Strings
Example	<pre>VAR str: ARRAY [0..30] OF CHAR; index:CARDINAL BEGIN str := 'This is a little test string'; index := Pos('This', str); (* index now equals 0 *) index := Pos('little', str); (* index now equals 10 *) index := Pos('big', str); (* index now equals 31 *)</pre>

PROC standard type S

Description	<i>PROC</i> is a predefined parameterless procedure type
Usage	<code>PROC = PROCEDURE ;</code>

Comments

PROCESS is an opaque type used to represent an independent process. In reality, it is nothing more than a storage area for a program counter and local heap and stack pointers. The user can manipulate variables of type *PROCESS* only through routines exported from *SYSTEM* and *Processes*. Notice that the module *Processes* is implemented using the primitives defined in *SYSTEM*.

The procedure *NEWPROCESS* sets up variables of type *PROCESS*, and the routines *TRANSFER* and *IOTRANSFER* provide cooperative and interrupt driven control functions.

See Also

TRANSFER
NEWPROCESS
Processes

Example

```
MODULE Process;  
VAR  
  M,P: PROCESS;  
  WorkSpace: ARRAY [0..99] OF CARDINAL;  
  PROCEDURE PP;  
  VAR  
    i: CARDINAL  
  BEGIN  
    FOR i:= 1 TO 10000 DO END ;  
    TRANSFER (P,M);  
  END PP;  
  BEGIN  
    NEWPROCESS(PP,ADR(WorkSpace),SIZE(WorkSpace),P)  
    TRANSFER (M,P);  
  END Process.
```

Comments PROCEDURE may be nested to any depth and may enclose modules. If a procedure has a local module, then that module is only in existence when the procedure is active. Thus, the module's initialization part is executed every time the procedure is called.

Example A procedure:

```
PROCEDURE AmI;
BEGIN
  WRITELN ( 'I am' );
END AmI;
```

A function:

```
PROCEDURE NotZero ( i :INTEGER ) :BOOLEAN
BEGIN
  RETURN ( i # 0 );
END NotZero;
```

PROCEDURE type R

Description PROCEDURE allows procedures to be assigned to variables.

Usage TYPE = PROCEDURE [Parameter Types] [:Result Type];

Parameter types and the result type can be any type.

Comments Procedure types allow you to assign arguments to a variable. You can use procedure-type variables if you want to call a procedure that has another procedure as a parameter. For example, suppose you define a procedure type as

```
TYPE StringProc = PROCEDURE(ARRAY OF CHAR);
```

This declares a procedure type *StringProc*. Any variable of type *StringProc* can take as argument an ARRAY OF CHAR. Note that nothing is mentioned about the contents of the procedure.

Suppose you have a procedure with *StringProc* as an argument:

```
PROCEDURE Example(X:CARDINAL P: StringProc);  
  :  
END Example;
```

This procedure expects two arguments: one is a CARDINAL, the other is a procedure of type *StringProc* with ARRAY OF CHAR parameter.

The procedure *Example* can be called with any parameters of these types. For example:

```
Example(12,OpenInput( ".INP" ));  
Example(21,OpenOutput( ".DTA" ));  
Example(3,WriteString( "Input data: " ))
```

These three examples all call procedure *Example* with the two expected arguments: a CARDINAL and a procedure with an ARRAY OF CHAR argument. The procedures *OpenInput*, *OpenOutput*, and *WriteString* are from the module *InOut*.

Note that standard procedures such as *EXCL* or *HALT* cannot be assigned to procedure variables, nor can procedures that are local to other procedures.

The predefined procedure type *PROC* denotes a procedure with no parameters.

Example

```

MODULE Procedures;
FROM InOut IMPORT WriteCard,WriteHex;
TYPE
  ListPointer = POINTER TO List;
  List = RECORD
    Key:CARDINAL
    Data:CARDINAL
    Next: ListPointer
  END ;
  ProcType = PROCEDURE (CARDINAL,CARDINAL);
VAR First: ListPointer;
PROCEDURE Search(X:CARDINAL Q: ProcType);
(* Search for record with Key X *)
  VAR P: ListPointer;
BEGIN
  P:= First;
  WHILE (P # NIL) AND (P^.Key # X) DO
    P:= P^.Next
  END ;
  IF P # NIL THEN
    Q(P^.Data,6)
  END
END Search;
BEGIN
  Search(3,WriteCard);
(* Write data from record with Key 3 asCARDINAL*)

  Search(3,WriteHex);
(* Write same data as a Hexadecimal value *)
END Procedures.
```

PROCESS type (SYSTEM) L

Description *PROCESS* declares processes or coroutines.

Usage p : PROCESS;

PROCESS must be imported from the library module **SYSTEM**.

Comments

PROCESS is an opaque type used to represent an independent process. In reality, it is nothing more than a storage area for a program counter and local heap and stack pointers. The user can manipulate variables of type *PROCESS* only through routines exported from *SYSTEM* and *Processes*. Notice that the module *Processes* is implemented using the primitives defined in *SYSTEM*.

The procedure *NEWPROCESS* sets up variables of type *PROCESS*, and the routines *TRANSFER* and *IOTRANSFER* provide cooperative and interrupt driven control functions.

See Also

TRANSFER
NEWPROCESS
Processes

Example

```
MODULE Process;
VAR
  M,P: PROCESS;
  WorkSpace: ARRAY [0..99] OF CARDINAL;
  PROCEDURE PP;
  VAR
    I: CARDINAL;
  BEGIN
    FOR I:= 1 TO 10000 DO END ;
    TRANSFER (P,M);
  END PP;
  BEGIN
    NEWPROCESS(PP,ADR(WorkSpace),SIZE(WorkSpace),P);
    TRANSFER (M,P);
  END Process.
```

Processes module L

Description *Processes* provides standard multiprogramming capabilities.

Declaration **DEFINITION MODULE Processes;**

TYPE SIGNAL;

PROCEDURE StartProcess(P: PROC; n:CARDINAL;

PROCEDURE SEND(VAR s: SIGNAL);

PROCEDURE WAIT(VAR s: SIGNAL);

PROCEDURE Awaited(s: SIGNAL):BOOLEAN;

PROCEDURE Init(VAR s: SIGNAL);

EXCEPTION DeadLock;

END Processes.

Comments Notice that there are no scheduling facilities associated with this module; the burden of providing these falls on the programmer.

Processes is implemented using the type **PROCESS** and the routines **NEWPROCESS** and **TRANSFER** in pseudomodule **SYSTEM**.

See Also **TRANSFER**

PROCESS

NEWPROCESS

Also see the individual identifiers exported from *Processes*.

progName variable (ComLine) L

Description *progName* is a string variable that contains the name of the program currently running.

Declaration **VAR**
progName: ARRAY [0..7] OF CHAR;

Usage **WRITE('The currently running program is ',progNa**

Comments *progName* allows for an eight-character name (assuming .COM extension).

progName is only valid if the program is started from within the Turbo Modula-2 shell. If the program is started from another operating system, the string is empty.

See Also ComLine

Example **MODULE RunFromTheShell;**
FROM ComLine IMPORT progName;
VAR
x,y: REAL;
BEGIN
IF progName < > '' THEN
WRITELN('This program is called ',progName);
END
END RunFromTheShell.

PromptFor procedure (ComLine) L

Description *PromptFor* looks for an input string from the command line. If there is not one present, it prompts the user to provide one.

Declaration **PROCEDURE PromptFor(prompt: ARRAY OF CHAR; VAR s: ARRAY OF CHAR);**

Usage	<pre>PromptFor(prompt, answer);</pre>
	<p><i>PromptFor</i> must be imported from the library module <i>ComLine</i>.</p>
	<p><i>prompt</i> and <i>answer</i> must be of type ARRAY OF CHAR.</p>
Comments	<p>This procedure essentially performs a <i>ReadString</i> from the TEXT <i>commandLine</i> exported by <i>ComLine</i>. If an <i>EOT</i> (32C) is encountered in <i>commandLine</i>, then the user is prompted with <i>prompt</i> to provide additional input.</p>
	<p>Note that in the Turbo Modula-2 system a command line can be specified by the user with the <i>Run</i> command from inside the Turbo Modula-2 shell.</p>
Example	<pre>MODULE ptest; FROM ComLine IMPORT PromptFor; BEGIN FOR i := 1 TO 3 DO BEGIN WRITE(i:3, ' '); PromptFor('Need more>', str); WRITELN(str); END ; END ptest.</pre>
	<p>If the user starts the program with</p>
	<p>Run MCD file: Ptest arg1 arg2</p>
	<p>then the output will be as follows:</p>
	<pre>1 ARG1 2 ARG2 3 Need more> arg3</pre>

QUALIFIED R

Description	QUALIFIED forces an identifier exported from a local module to be qualified in the surrounding environment.
Usage	EXPORT QUALIFIED this, that, theNextThing;
Comments	When an identifier is qualified, it means that every reference to it must be prefixed by the name of the module it has been defined in, followed by a period (.). For example, these identifiers are qualified: <i>Texts.ReadChar</i> , <i>MathLib.Cos</i> , <i>Files.EOF</i> . These are not: <i>ReadChar</i> , <i>Cos</i> , <i>EOF</i> .
	This command only makes sense for local modules, since main modules and implementation modules can't export anything, and in definition modules everything is exported by default.
Example	<i>p1</i> in <i>local</i> is EXPORT QUALIFIED :
	<pre>MODULE main; MODULE local; EXPORT QUALIFIED p1; PROCEDURE p1; BEGIN END ; END local; BEGIN local.p1; (* p1 is visible, but must be qualified *) END main.</pre>

RAISE statement R/E

Description *RAISE* causes an exception to be raised, passing control up the calling chain to an exception handler or the runtime system.

Usage **RAISE** [SomeError] [, message];

message must be a character array.

Comments When an exception is raised, the current procedure is searched for the appropriate handler. If one is found, the specified code is executed. If a handler for the exception is not found, the caller of the current procedure is searched. This search process continues until a handler is found or the runtime supervisor is reached. The supervisor prints the exception name on the terminal along with the optional message and the calling chain, control is then returned to the operating system.

The short form of the exception statement is **RAISE**, which may only be used within an exception handler. If this statement occurs in an exception handler, then the same exception that the exception handler is currently handling will be raised again in the next highest procedure of the calling sequence.

See Also **EXCEPTION**

Example Raise an exception to be caught somewhere:

RAISE MyError

Raise an exception with a message:

RAISE TooLateAtNight, 'Turn me off and go home !'

Do a little bit of exception-handling and then pass it on:

```
EXCEPTION
| DiskFull:
  WRITELN('DiskFull in thisProc');
  Close(f);
  RAISE ;
ELSE
  RAISE ;
END thisProc.
```

Random function (MathLib) L

Description *Random* returns a random number between 0 and 1.

Declaration PROCEDURE Random(): REAL;

Usage Y:= Random();

Result *Y* is REAL.

There is no argument, but the parentheses are required.

Comments The function returns the next element from a pseudo-random number sequence. The same sequence is repeated each time your program is executed unless you select a new sequence with the procedure *Randomize*.

Example Two ways to randomize the random number generator:

```
FROM MathLib IMPORT Randomize;  
VAR Sequence:CARDINAL  
BEGIN  
  WRITELN('Input sequence number, 0 to 65535');  
  READ(Sequence);  
  Randomize(Sequence);
```

or

```
i := 0;  
WRITE('Enter any character');  
REPEAT  
  i := i + 1;  
  BusyRead( ch );  
UNTIL ch # OC;  
Randomize( i );
```

Randomize procedure (MathLib) L

Description	<i>Randomize</i> initializes a pseudo-random-number sequence. The same random sequence can be reproduced by passing the same seed value to <i>Randomize</i> .
Declaration	PROCEDURE Randomize(n:CARDINAL);
Usage	<i>Randomize</i> (1986);
	<i>Randomize</i> must be imported from the module <i>MathLib</i> .
Comments	After randomizing the sequence, call <i>Random</i> to get the random numbers.
See Also	<i>MathLib</i> <i>Random</i>

Example

```
MODULE RandomMath;
FROM MathLib IMPORT Randomize, Random;
VAR x: REAL;
BEGIN
  Randomize(1986);
  x := Random( )*10.0/Random( )*20.0;
END RandomMath.
```

READ standard procedure S/E

Description *READ* is a generalized read statement.

Usage *READ(t, vlist)*

or

READ(vlist);

t is of type *TEXT* imported from *TEXTS*. *vlist* is simply a list of variables and constants.

Comments

READ will be translated by the compiler into the appropriate calls to the module *Texts*.

Notice that you do not have to import procedures from *Texts* in order to use these commands, but *Texts* must be somewhere on disk.

Although *READ*, *READLN*, *WRITE*, and *WRITELN* are not part of the standard Modula-2 definition, programming without them is much more tedious. However, if a program is to be portable, do not use these commands. A good strategy would be to develop the program using these statements for debugging and explicitly typed calls to *Texts* for the user's output. When the program is moved to another environment, only the module *Texts* must be duplicated in the new environment.

See Also	READLN Texts WRITE WRITELN
Example	Read an integer <i>i</i> :
	READ(i)
	Read a string <i>str</i> followed by 4 integers <i>i,j,k,l</i> :
	READ(str, i,j,k,l)
	Read a real number <i>r</i> from TEXT <i>t</i> :
	READ(t, r)

ReadAgain procedure (Terminal) L

Description	<i>ReadAgain</i> causes the last character read from the terminal to be repeated.
Declaration	PROCEDURE ReadAgain;
Usage	ReadAgain;
	<i>ReadAgain</i> must be imported from the library module <i>Terminal</i> .
Comments	The last character read before this procedure is called will be picked up by the next read operation from the terminal. This procedure works with all the input procedures in <i>Terminal</i> , but only works if at least one character has already been processed since the program began. With this routine you can achieve a certain amount of lookahead without having to set up a global buffer and complicated logic to use it.

See also	BusyRead ReadChar Terminal
Example	Read a character and then read it again: ReadChar(ch); ReadAgain; ReadChar(ch);

ReadAgain procedure (Texts) L

Description	<i>ReadAgain</i> causes the last character read from a <i>TEXT</i> to be repeated.
Declaration	PROCEDURE ReadAgain(<i>t</i> : TEXT);
Usage	<i>ReadAgain(t);</i> <i>ReadAgain</i> must be imported from the library module <i>T</i> . <i>t</i> must be of type <i>TEXT</i> imported from <i>Texts</i> .
Comments	The last character read from <i>t</i> before this procedure is called will be picked up by the next read operation from <i>t</i> . This procedure works with all the input procedures in <i>T</i> , but only works if at least one character has already been processed since the program began. With this routine, you can achieve a certain amount of lookahead without having to set up a global buffer and complicated logic to use it.
See Also	ReadChar Texts

Example Read a character from the *TEXT t* and then read it again:

```
ReadChar (t, ch);
ReadAgain(t);
ReadChar (t, ch);
```

Or, using the *ReadAgain*'s potential lookahead to good advantage, you can read a value as an integer or a string depending on the first nonblank character:

```
REPEAT
  ReadChar(t,ch);
UNTIL ch # " ";
ReadAgain(t);
IF (ch >= "0") and (ch <= "9") THEN
  ReadInt(t,x)
ELSE
  ReadString(t,s)
END ;
```

ReadByte procedure (Files) L

Description *ReadByte* reads a variable of any type compatible with *BYTE* from *FILE*.

Declaration **PROCEDURE** ReadByte(*f*: FILE; **VAR** *ch*: BYTE);

Usage *ReadByte(f, ch);*

ReadByte must be imported from the library module *Files*.

f must be of type *FILE* imported from *Files*.

ch must be of a type compatible with *BYTE*.

Comments	Types compatible with <i>BYTE</i> are those whose storage uses exactly 1 byte. These include CHAR, BOOLEAN, and enumeration types with at most 256 elements; subrange types with bounds within the range 0 to 255, and word-sized variables as well (but only 1 byte is read).
See Also	Files ReadWord ReadRec WriteByte
Example	Open a file 'ONEBYTE.DAT' and write out a byte; close it, and read it back:

```

MODULE writebyte;
FROM Files IMPORT FILE, Create, Close, ReadByte,
          WriteByte;
VAR   f : FILE;
        ch : CHAR;
BEGIN
  Create(f, 'ONEBYTE.DAT');
  WriteByte(f,69);
  Close(f);
  IF Open(f, 'ONEBYTE.DAT') THEN
    ReadByte(f, ch);
    (* ch now contains the character 'E'
     ( ascii code 69 ) *)
    Close(f);
  END writebyte.
```

ReadBytes procedure (Files) L

Description	<i>ReadBytes</i> reads a number of bytes into a block of memory beginning at an address.
Declaration	PROCEDURE ReadBytes(f: FILE; buf: ADDRESS; nbytes CARDINAL):CARDINAL;

Usage	<pre>nread := ReadBytes(f, bufadr, nbytes);</pre>
	<p><i>ReadBytes</i> must be imported from the library module <i>Files</i>.</p>
	<p><i>f</i> must be of type <i>FILE</i> imported from <i>Files</i>.</p>
	<p><i>buf</i> must be compatible with type <i>ADDRESS</i> (any pointer).</p>
	<p><i>nread</i> and <i>nbytes</i> must be of type <i>CARDINAL</i>.</p>
Comments	<p><i>ReadBytes</i> will attempt to read <i>nbytes</i> from <i>f</i>. It returns the number of bytes successfully read. It will not raise the exception <i>EndError</i>, even if the end of the file is encountered before <i>nbytes</i> bytes are processed.</p> <p>The variables are read into memory starting at the location specified in <i>bufadr</i>.</p> <p><i>ReadBytes</i> and <i>WriteBytes</i> are low-level routines with which you can take snapshots of memory to and from the disk, copy files, and so on. They must be used with care, since you can easily destroy your program or its variables with them. When transferring variables, we recommend using an explicitly typed procedure like <i>ReadRec</i>, <i>ReadWord</i>, and the like.</p>
See Also	<p>FILE Files WriteBytes</p>
Example	<p>Read 1024 bytes from <i>f</i> into memory starting at 8000 Hex:</p> <pre>uread := ReadBytes(f, 8000H, 1024)</pre> <p>The previous example is not very safe since the programmer doesn't always know what is at a given address in memory. The following is more useful:</p> <pre>nbytes := ReadBytes(f, ADR(area), SIZE(area));</pre>

ReadCard procedure (Texts) L

Description	<i>ReadCard</i> reads a cardinal number from a <i>TEXT</i> file.
Declaration	PROCEDURE ReadCard(<i>t</i> : TEXT ; VAR <i>c</i> : CARDINAL);
Usage	ReadCard(<i>t</i> , <i>card</i>);
	<i>ReadCard</i> must be imported from the library module <i>Texts</i> .
	<i>t</i> must be of type <i>TEXT</i> .
	<i>card</i> must be of type CARDINAL .
Comments	In general, <i>READ</i> is easier to use. This function works by reading a string into an integer variable and then converting it via <i>StrToCard</i> in <i>Conversion</i> . Leading blanks are skipped, and the range must be within the legal range of <i>MAX(CARDINAL)</i> .
	The BOOLEAN variable <i>Done</i> (exported by <i>Texts</i>) can be used to determine whether input was syntactically correct within the legal range. If <i>Done</i> is FALSE, the number returned by <i>ReadCard</i> is undefined.
See Also	READ Texts WriteCard
Example	Read a cardinal number called <i>c</i> from the file <i>t</i> : ReadCard(<i>t</i> , <i>c</i>)

ReadChar procedure (Terminal) L

Description	<i>ReadChar</i> reads a character from the terminal without echo.
Declaration	PROCEDURE ReadChar(VAR ch: CHAR);
Usage	ReadChar(c);
	<i>ReadChar</i> must be imported from the library module <i>Terminal</i> .
	c must be of type CHAR .
Comments	The character is not echoed to the screen.
	Line ends are returned as carriage returns, CHR(13).
See Also	BusyRead Terminal WriteChar
Example	Read a character ch from the terminal without echo: ReadChar(ch)

ReadChar procedure (Texts) L

Description	<i>ReadChar</i> reads a character from a <i>TEXT</i> file.
Declaration	PROCEDURE ReadChar(t: TEXT; VAR ch: CHAR);
Usage	ReadChar(t, c);
	<i>ReadChar</i> must be imported from the library module <i>Texts</i> .
t	must be of type TEXT .
c	must be of type CHAR .

Comments	In general, <i>READ</i> is easier to use than this procedure.
	A single character is read in; if the stream is <i>input</i> , then the character is echoed to the screen. For input without echo, use <i>ReadChar</i> from the module <i>Terminal</i> .
	Input from the console must always be terminated with a return.
See Also	<i>READ</i> <i>WriteChar</i>
Example	Read a character <i>c</i> from the terminal with echo: <code>ReadChar(input, c)</code> Read a character from the <i>TEXT</i> file <i>t</i> : <code>ReadChar(t, c)</code>

ReadDouble procedure (Doubles) L

Description	<i>ReadDouble</i> reads a double-precision real variable from a <i>TEXT</i> file.
Declaration	PROCEDURE <i>ReadDouble</i> (<i>t</i> : <i>TEXT</i> ; VAR <i>d</i> : <i>LONGREAL</i>);
Usage	<code>ReadDouble(t, d);</code>
	<i>ReadDouble</i> must be imported from the library module <i>Doubles</i> . <i>t</i> must be of type <i>TEXT</i> imported from <i>Texts</i> . <i>d</i> must be of type <i>LONGREAL</i> .

Comments	This procedure works by reading a string from <i>t</i> and then converting it to a LONGREAL using <i>StrToDouble</i> . The Boolean variable <i>legal</i> (exported from <i>Doubles</i>) is set to TRUE if a legal double-precision number is read from <i>t</i> .
	The <i>READ</i> and <i>READLN</i> procedures use <i>ReadDouble</i> to read LONGREALs.
See Also	<i>WriteDouble</i>

Example Read a double-precision *d* out of a text *t*:

```
ReadDouble(t, d);
```

ReadInt procedure (Texts) L

Description	<i>ReadInt</i> reads an integer from a <i>TEXT</i> file.
Declaration	PROCEDURE <i>ReadInt</i> (<i>t</i> : <i>TEXT</i> ; VAR <i>i</i> : <i>INTEGER</i>);
Usage	<i>ReadInt</i> (<i>t</i> , <i>i</i>); <i>ReadInt</i> must be imported from the library module <i>Texts</i> . <i>t</i> must be of type <i>TEXT</i> . <i>i</i> must be of type <i>INTEGER</i> .

Comments	In general, <i>READ</i> is easier to use than this procedure.
	This routine works by reading a string from text <i>t</i> and converting it to an INTEGER using <i>IntToStr</i> in <i>Convert</i> .
	Leading blanks are skipped, and the range must be within <i>MIN(INTEGER)</i> to <i>MAX(INTEGER)</i> .
	The BOOLEAN variable <i>Done</i> (exported by <i>Texts</i>) can be used to determine whether input was syntactically correct and within the legal range. If <i>Done</i> is FALSE, the number returned by <i>ReadInt</i> is undefined.
See Also	<i>READ</i> <i>Texts</i> <i>WriteInt</i>
Example	Read an integer from the text <i>t</i> : <i>ReadInt(t, i)</i>

ReadLine procedure (Texts) L

Description	<i>ReadLine</i> reads a whole line of text into a string variable.
Declaration	PROCEDURE <i>ReadLine(t: TEXT; VAR s: ARRAY OF CHAR);</i>
Usage	<i>ReadLine(t, line);</i>
	<i>ReadLine</i> must be imported from the library module <i>Texts</i> .
	<i>t</i> must be of type <i>TEXT</i> imported from <i>Texts</i> .
	<i>line</i> must be an ARRAY OF CHAR .
Comments	<i>ReadLine</i> and <i>ReadString</i> are similar except that <i>ReadString</i> skips over leading blanks and <i>ReadLine</i> doesn't, and <i>ReadLine</i> is only terminated by an EOL or an EOT.

See Also **ReadString**
 Texts

Example Read from *TEXT t* a string *str* with leading blanks:

ReadLine(t, str)

ReadLine procedure (Terminal) L

Description *ReadLine* reads a line of text into a string variable.

Declaration **PROCEDURE ReadLine(VAR s: ARRAY OF CHAR);**

Usage **ReadLine(line);**

ReadLine must be imported from the library module *Terminal*.

line must be an **ARRAY OF CHAR**.

Comments *ReadLine* is similar to *Texts.ReadLine*, but *ReadLine*'s input is terminated by a carriage return instead of an EOT character.

See Also **Terminal**
 Texts

Example Read a string *str* from the console:

```
MODULE ReadLineEx;
FROM Terminal IMPORT ReadLine;
VAR
  s: ARRAY [0..20] OF CHAR;
BEGIN
  WRITELN('Enter a string and press return ');
  ReadLine(str);
  WRITELN('The string you entered is ',str);
END ReadLineEx;
```

READLN standard procedure S/E

Description *READLN* is a generalized read statement that terminates with a new line.

Usage *READLN*(*{t}*, *vlist*)

or

READLN(vlist);

t is an optional *TEXT* file; the default is *input*.

vlist is simply a list of variables.

Comments If *t* is omitted, the standard input will be used.

READLN will be translated by the compiler into the appropriate calls to the module *Texts*. Thus do not change the definition of the module *Texts* if you want to use this sort of input and output.

READLN works just like *READ* except it skips over the text file being read until an *EOL* is encountered.

See Also *READLN*
Texts
WRITE
WRITELN

Example Get a number from the user at the terminal:

```
WRITE('Gimme a number, doesn't matter which one >');
READLN(number);
```

Read a bunch of numbers from a *TEXT* file:

```
READLN(t, num1, num2, num3, numlast);
```

ReadLong procedure (Texts) L

Description *ReadLong* reads a long integer from a *TEXT* file.

Declaration **PROCEDURE** ReadLong(*t*: TEXT; **VAR** *l*: LONGINT);

Usage *ReadLong(t, l);*

ReadLong must be imported from the library module *Texts*.

t must be of type *TEXT*.

l must be of type LONGINT.

Comments In general, *READ* is easier to use than this procedure.

ReadLong works by reading a string from a *TEXT* file and converting it to a LONGINT using the procedure *StrToLong* from *Convert*.

Leading blanks are skipped, and the range must be within *MIN(LONGINT)* to *MAX(LONGINT)*.

The BOOLEAN variable *Done* (exported by *Texts*) can be used to determine whether input was syntactically correct and within the legal range. If *Done* is FALSE, the number returned by *ReadLong* is undefined.

See Also **READ**

Texts

WriteLong

Example Read from *TEXT* file *t* a long integer *l*:

ReadLong(t, l)

ReadReal procedure (Texts) L

Description	<i>ReadReal</i> reads a real number from a <i>TEXT</i> file.
Declaration	PROCEDURE ReadReal(<i>t</i> : TEXT; VAR <i>r</i> : REAL);
Usage	<i>ReadReal(t, r);</i>
	<i>ReadReal</i> must be imported from the library module <i>Texts</i> .
	<i>t</i> must be of type <i>TEXT</i> .
	<i>r</i> must be of type <i>REAL</i> .
Comments	In general, <i>READ</i> is easier to use than this procedure. Leading blanks are skipped, and the range must be within the legal real-number range <i>MIN(REAL)</i> to <i>MAX(REAL)</i> . The BOOLEAN variable <i>Done</i> (exported by <i>Texts</i>) can be used to determine whether input was syntactically correct and within the legal range. If <i>Done</i> is FALSE, the number returned by <i>ReadLong</i> is undefined.
See Also	<i>READ</i> <i>Texts</i> <i>WriteReal</i>
Example	Read a real number <i>r</i> from text file <i>t</i> . <i>ReadReal(t, r)</i>

ReadRec procedure (Files) L

Description	<i>ReadRec</i> reads a record, array, or real from a <i>FILE</i> (a binary disk file).
Declaration	PROCEDURE ReadRec(<i>f</i> : FILE; VAR <i>rec</i> : ARRAY OF WORD);

Usage `ReadRec(f, r);`

ReadRec must be imported from the library module *Files*.

f must be of type *FILE* imported from *Files*.

r can be of any type.

Comments This routine can read any variable not covered by *ReadByte* and *ReadWord*. It reads variables compatible with **ARRAY OF WORD**; for example, **REAL**, **LONGREAL**, **LONGINT**, **ARRAY**, or **RECORD**.

See Also *Files*
ReadByte
ReadWord
WriteRec

Example Open a file 'SOMEWREK.DAT' and write out this record, then read it back:

```
MODULE writefile;
FROM Files IMPORT FILE, Create, Close, WriteRec,
ReadRec;
FROM wreck IMPORT wreckRec;
VAR
  f : FILE;
  wr : wreckRec;
BEGIN
  Create(f, 'SOMEWREK.DAT');
  WriteRec(f, wr);
  Close(f);
  IF Open(f, 'SOMEWREK.DAT') THEN
    ReadRec(f, wr);
    (* wr has whatever it had before in it *)
    Close(f);
  END ;
END writefile.
```

ReadString procedure (Texts) L

Description *ReadString* reads a string, which in this case is a sequence of characters that does not contain blanks or control characters.

Usage `ReadString(t, s);`

ReadString must be imported from the library module *Texts*.

t must be of type *TEXT* imported from *Texts*.

s must be of type **ARRAY OF CHAR**.

Comments In general, *READ* is easier to use than this procedure.

The routine works by skipping over blanks until it gets to a nonblank character, at which point it starts to read. Input is then terminated by a blank or a control character (like *EOL* or *EOT*). Thus, strings cannot be separated by commas.

If you want to read in a line as it is, with leading blanks and all, use *ReadLine*.

See Also **ReadLine**
TEXTS
WriteString

Example Read a string, ignoring leading blanks, from a *TEXT* file *t*:

`ReadString(t, str)`

ReadWord procedure (Files) L

Description *ReadWord* reads a variable of any type compatible with *WORD* from a *FILE*.

Declaration **PROCEDURE** `ReadWord(f: FILE; VAR w: WORD);`

Usage	ReadWord(f, w);
	<i>ReadWord</i> must be imported from the library module <i>Files</i> .
	<i>f</i> must be of type <i>FILE</i> imported from <i>Files</i> .
	<i>w</i> must be of a type compatible with <i>WORD</i> .
Comments	Types compatible with <i>WORD</i> are those whose storage uses exactly one word (2 bytes); these include <i>INTEGER</i> , <i>CARDINAL</i> , <i>BITSET</i> and all pointers (<i>ADDRESS</i> is a pointer).
(
See Also	Files ReadByte ReadRec WriteWord
Example	Open a file 'SOMEYEAR.DAT' and write out this year for later:
	<pre>MODULE writefile; FROM Files IMPORT FILE, Create, Open, Close, WriteWord, ReadWord; VAR f : FILE; okay : BOOLEAN; y : CARDINAL; BEGIN Create(f, 'SOMEYEAR.DAT'); WriteWord(f, 1985); Close(f); IF Open(f, 'SOMEYEAR.DAT') THEN ReadWord(f, y); (* y now has the value 1985 *) END ; END writefile.</pre>

REAL standard type S

Description REAL is a standard type with variables that can assume any value between -6.80565E+38 and 6.80565E+38.

Comments You will use REAL variables whenever your computation is likely to produce a fraction. For example, to record the mass of liquid emitting from a pipe, use *MassLiquid*:= *RateOfFlow***Density*.

In this case, none of the variables are likely to be whole numbers so REAL types are appropriate.

It should be remembered that REAL computations are not always exact; a certain amount of rounding occurs. For example, an expression such as *Y*:= 3.0 * (1.0/3.0) will not give the expected result of 1.0, but a value slightly less. In practice, the error is so small that it can usually be ignored. But don't make the following mistake:

```
Y:= 3.0 * (1.0/3.0);
IF Y = 1.0 THEN
  :
```

The correct solution is

```
IF (Abs(Y) - 1.0) < 0.01 THEN
  :
```

If *Y* is between 0.99 to 1.01, the statement sequence will be executed.

You may use any arithmetic operator in REAL expressions, and REAL variables may take part in relational expressions.

The *FLOAT* standard function will convert other numeric data types into REAL.

It should be noted that the REAL type is not an ordinal type and cannot be used to index arrays, define subranges or sets,

control FOR or CASE statements, or take arguments in any procedure expecting an ordinal type.

Turbo Modula-2 uses the IEEE 4-byte real-number format; that is, in the order of their significance, 1 sign bit, 8 exponent bits, and 23 mantissa bits.

31	30	23	22	0
s	exponent		mantissa	

The mantissa is in binary format but offset by 1.0, representing only those digits to the right of the decimal point. Thus,

000000H = 1.0
400000H = 1.5
7FFFFFFH = 1.9999999

The exponent is in binary format but with a 80H offset. For example, an exponent 84H indicates the mantissa is to be multiplied by

$$2^{(84H - 80H)} = 2^4 = 16$$

If the exponent is zero, the floating-point value is considered zero.

The least-significant bytes are stored first, at the lower addresses.

Sample Hex real numbers:

3F800000H = 1.0
40000000H = 2.0
40490FDBH = 3.141592

See the example in WORD for printing out other values.

Example **VAR**
 SquareArea, SquareSide: REAL;
BEGIN
 SquareSide:= 4.21;
 SquareArea:= SquareSide * SquareSide;

REALOVERFLOW exception (SYSTEM) L/E

Description REALOVERFLOW is raised by the module SYSTEM when a real expression results in a value that is too large or too small to be represented by the result type.

Declaration **EXCEPTION**
 REALOVERFLOW;

Usage **EXCEPTION**
 REALOVERFLOW:
 WRITELN('Overflow.');//
END

REALOVERFLOW must be imported from the module SYSTEM.

Comments This exception occurs only during evaluation of real expressions. Integer and cardinal expressions generate the exception OVERFLOW.

See Also SYSTEM
 BADOVERLAY
 OUTOFGMEMORY
 OVERFLOW

Example

```
MODULE RealOverflow;
FROM SYSTEM IMPORT REALOVERFLOW;
VAR
  r: REAL;
BEGIN
  r := MAX(REAL) * 2.2;
EXCEPTION
  REALOVERFLOW:
    WRITELN('Real Overflow');
END RealOverflow.
```

RealToStr procedure (Convert) L

Description *RealToStr* writes a real number into a string.

Declaration PROCEDURE RealToStr(*r*: REAL; VAR *s*: ARRAY OF CHAR; *digits*: INTEGER);

Usage RealToStr(*r*, *s*, *digits*);

RealToStr must be imported from the library module *Convert*.

r must be of type *Real*.

s must be of type **ARRAY OF CHAR**.

digits must be of type **CARDINAL**.

Concerning its The number *r* will be converted to the string *s*. This procedure is similar in form and action to *WriteReal* in *Texts*. Here, the declared length of *s* is substituted for the *width* parameter in *WriteReal*.

The *digits* parameter indicates the number of digits to the right of the decimal point. The sign of *digits* determines whether or not scientific notation will be used, with plus for normal notation and minus for scientific.

If *digits* is zero, the decimal point will be dropped.

If *s* is too small to contain the string representing *r*, the exception *TooLarge* will be raised.

See Also

CardToStr
Convert
Doubles
IntToStr
LongToStr
StrToReal

Example

Convert a real to a string:

```
VAR sstr : ARRAY [0..4] OF CHAR;
  lstr : ARRAY [0..9] OF CHAR;
  r :REAL
BEGIN
  r := 3.14159;
  RealToStr( r, sstr, 3 );
  (* sstr now has the value '3.141' *)

  RealToStr( r, sstr, 1 );
  (* sstr now has the value '3.1' *)

  RealToStr( r, sstr, 0 );
  (* sstr now has the value '3' *)

  RealToStr( r, sstr, -3 );
  (* causes the exception TooLarge to be raised *)

  RealToStr( r, lstr, -3 );
  (* lstr now has the value '3.141E+00' *)
```

RECORD type R

Description A record consists of a fixed number of bytes that contain *fields* of possibly different types. Each field has a name, the *field identifier*, which is used to select the field. A record may have variant parts that allow the record to have more than one field occupying the same memory.

Usage RecordType = RECORD
 Field sequence
 END

Field sequence = Field { ; Field }

A *field* can be either a *fixed field* or a *variant field*.

Fixed field = Field identifier: Field type

Variant field = CASE [Tag identifier]: Tag type OF
 [Case label: Field sequence] [|
 [Case label: Field sequence]]

 [ELSE Field sequence]
 END

The *field type* can be any type.

The *tag type* can be of type INTEGER, CARDINAL, BOOLEAN, CHAR, or enumeration.

Each *case label* can be either a single value or a range of values of the tag variable separated by commas.

Comments You can use a record to handle several related items of different types; for example, a record of the date could be

```
Date : RECORD
    Day: [1..31];
    Month: (Jan,Feb,Mar,Apr,May,June,
              July,Aug,Sept,Oct,Nov,Dec);
    Year: [1900..1999]
END
```

Each field in a record has an identifier that is used to access the field. Thus, in the example, a field is specified as *Date.Day*, *Date.Month*, and so on.

A field can take values defined by the field type. In the previous example, the field *Month* can take any value from the enumeration *Jan* to *Dec*.

For variant records, the tag variable is used to select the different cases depending on its value. For example:

```
TYPE
    Lights = (Red,Orange,Yellow,Green,Blue);
VAR
    Example: RECORD
        Name: ARRAY [1..20] OF CHAR;
        CASE Color: Lights OF
            Red..Yellow: Test : CARDINAL |
            Green, Blue: Today: Date |
        END
    END
```

In this example, the value of *Color* selects the form of the record. Thus, if *Select* is a record of type *Example*, altering the value of the tag *Select.Color* changes the record's structure. If the tag has a value of *Red*, *Orange*, or *Yellow* and the record is accessed, the field *Select.Test* exists and can be given a value.

No operators are directly applicable to records (except assignment), but any operator applicable to a field type can operate on that field. Thus, if a record *Value* has a field *RealNumbers* of type REAL, then the following is valid:

```
Value.RealNumbers := 2.31 * Value.RealNumbers
```

The assignment of values to record fields can either be explicitly stated, specifying both the record identifier and the field:

```
Value.CardNumbers := 34
```

or can use a **WITH** statement:

```
WITH Value DO
  CardNumbers := 34;
  RealNumbers := 2.31 * RealNumbers
END
```

The memory requirements for a record with fixed fields is the sum of the individual field lengths, but the fields that have 1 byte types are allocated in 2 bytes. For example, a record with two character fields actually occupies 4 bytes. A variant field requires enough memory for the largest of the cases.

Example

```
MODULE Records;
TYPE
  Origin = (Citizen, Alien);
  Name   = ARRAY [1..20] OF CHAR;
  Date   =RECORD
    Day: [1..31];
    Month: (Jan, Feb, Mar, Apr, May, June,
              July, Aug, Sept, Oct, Nov, Dec);
    Year: [1900..1999]
  END ;
  Person = RECORD
    PersonName: Name;
    BirthDate: Date;
    CASE CitizenShip: Origin OF
      Citizen: BirthPlaceName: Name
      !
      Alien: CountryOfOrigin: Name;
      DateOfEntry: Date
    END
  END ;
```

```

VAR
  Passenger: Person;
BEGIN
  WITH Passenger DO
    WITH BirthDate DO
      Day:= 15;
      Month:= Sept;
      Year:= 1958
    END ;
    CASE CitizenShip OF
      | Citizen: BirthPlaceName:= 'New Orleans'
      | Alien: CountryOfOrigin:= 'UK';
        WITH DateOfEntry DO
          Day:= 22;
          Month:= Aug;
          Year:= 1984
        END
      END
    END ;
    WRITELN(Passenger.BirthDate.Year);
  END Records.

```

RedirectInput procedure (ComLine) L

Description *RedirectInput* allows a filename defined in the command line to be designated as standard output.

Declaration **PROCEDURE** RedirectInput;

Usage **RedirectInput;**

RedirectInput must be imported from the library module *ComLine*.

Comments	The file name preceded by a >< < in the command line will be reassigned to the standard stream <i>input</i> . No space is allowed between the >< < and the file name.
	If no standard input file is specified in the command line, then the standard input is defined to be at the terminal console.
	The name of the file specified to be the standard input stream is retrievable in the variable <i>inName</i> exported by <i>ComLine</i> .
See Also	Comline RedirectOutput

Example Redirect the input:

```
MODULE RedirectTheInput;
FROM ComLine IMPORT RedirectInput,RedirectOutput;
BEGIN
  RedirectInput;
  (* Now input may be redirected from the
     command line *)
END RedirectTheInput.
```

RedirectOutput procedure (ComLine) L

Description	<i>RedirectOutput</i> allows a filename defined in the command line to be designated as standard output.
Declaration	PROCEDURE RedirectOutput;
Usage	RedirectOutput;

RedirectOutput must be imported from the library module *ComLine*.

Comments	The file name preceded by a >> in the command line will be reassigned to the standard stream <i>output</i> .
	If no standard output file is specified in the command line, then output is defined to be the terminal screen.
	The name of the file specified to be the standard output stream is retrievable in the variable <i>outName</i> exported by <i>ComLine</i> .
See Also	Comline RedirectInput

Example Redirect the output:

```
MODULE RedirectTheOutput;
FROM ComLine IMPORT RedirectInput,RedirectOutput;
BEGIN
  RedirectOutput;
  (* Now output can be redirected from the
     command line *)
```

RELEASE procedure (STORAGE) L

Description	<i>RELEASE</i> releases a block of allocated memory from the heap.
Declaration	PROCEDURE RELEASE(VAR a: ADDRESS);
Usage	<i>RELEASE(a);</i>
	<i>RELEASE</i> must be imported from the pseudomodule <i>STORAGE</i> . <i>a</i> must be compatible with type <i>ADDRESS</i> (any pointer).

Comments	All of the memory from address <i>a</i> to the top of the heap will be released; <i>a</i> will become the new top of the heap.
	You can only release a block of memory that has been previously specified with <i>MARK</i> .
	After the call to <i>RELEASE</i> , the value of <i>a</i> will be set to <i>NIL</i> , and any variable allocated on the heap prior to the <i>MARK</i> call will again be accessible.
	<i>RELEASE</i> should be used with care in conjunction with <i>DISPOSE</i> , <i>DEALLOCATE</i> , and file handling.
See Also	MARK NEW STORAGE
Example	Allocate a block of variables and then throw them all away: MARK(blockStart); NEW(blockVar1); NEW(blockVar2); NEW(blockVar3); (* Now use these variables, and when done: *) RELEASE(blockStart);

Rename procedure (Files) L

Description	<i>Rename</i> changes the name of an open file and then closes it.
Declaration	PROCEDURE Rename(VAR f: FILE; name: ARRAY OF CHAR);

Usage `Rename(f, newname);`

Rename must be imported from the library module *Files*.

f must be a variable of type *FILE*, imported from the module *Files*.

newname must be a character array.

Comments The file *f* must have already been opened under its old name.

newname must be a valid CP/M file name.

This function has the side effect of closing the renamed file. If you want to use it again, you must reopen it with the new name.

See Also `Close`
`Delete`
`Files`
`Open`

Example `MODULE renameExample;`
`FROM Files IMPORT FILE, Open, Rename, Close`
`VAR`
 `f : FILE;`
 `newname : ARRAY [0..12] OF CHAR;`
 `oldname : ARRAY [0..12] OF CHAR;`
`BEGIN`
 `oldname := 'oldtest.txt';`
 `newname := 'newtest.txt';`
 `IF Open(f, oldname) THEN`
 `Rename(f, newname)`
 `END ;`
`END renameExample.`

REPEAT statement R

Description	REPEAT repeatedly executes a sequence of statements until a Boolean expression returns a TRUE result.
Usage	REPEAT Statement sequence UNTIL BooleanExpression Statement sequence = Statement {; Statement}
Comments	The Boolean expression is evaluated only after the statement sequence has been executed at least once. The repetition is terminated when the expression gives a TRUE result.
Example	MODULE Repeats; VAR Answer: CHAR; BEGIN REPEAT WRITE('Enter "Q" and <RET> to continue . . . ');\br/> READ(Answer); UNTIL CAP(Answer) = 'Q'; END Repeats.

ResetSys procedure (Files) L

Description	<i>ResetSys</i> resets the disk drive system.
Declaration	PROCEDURE ResetSys();
Usage	<i>ResetSys</i> ;
	<i>ResetSys</i> must be imported from the library module <i>Files</i> .

Comments	<i>ResetSys</i> asks CP/M to update the file directory. Thus, after a call to <i>ResetSys</i> , it is possible to write to drives where disks have been swapped, but all open output files will then be lost. Consequently, make sure all your disk files are closed before using this procedure.
See Also	Files
Example	Get the user to change disks:

```
WRITELN('Change disk and type return when ready');
READLN(ch);
ResetSys;
```

RETURN statement R

Description	RETURN terminates execution and specifies the return value of a function procedure.
Usage	RETURN ; or RETURN answer;
	<i>answer</i> must be compatible with the type with which the function procedure has been declared.
Comments	Normally, there will be only one RETURN in a function procedure, just before the final END . However, it is allowable to put as many RETURNS as you like throughout the procedure. When a RETURN is encountered, the procedure will terminate, and the value specified at that point will be returned to the calling procedure.

Example

```

PROCEDURE maxi3(i1, i2, i3 :INTEGER) :INTEGER
BEGIN
  IF ((i1>i2) & (i1>i3)) THEN
    RETURN i1
  ELSIF (i2>i3) THEN
    RETURN i2
  ELSE
    RETURN i3
  END
END maxi3;

```

SEND procedure (Processes) L

Description	<i>SEND</i> transfers control to any process waiting for a given signal.
Declaration	PROCEDURE SEND(VAR s: SIGNAL);
Usage	SEND(s);
	<i>SEND</i> must be imported from the library module <i>Processes</i> .
	<i>s</i> must be of type <i>SIGNAL</i> imported from <i>Processes</i> .
Comments	<p><i>SEND</i> is used in conjunction with the other routines in the module <i>Processes</i>: <i>Init</i>, <i>StartProcesses</i>, <i>Wait</i>, and <i>Awaited</i>.</p> <p>If no process is waiting for signal <i>s</i>, then <i>SEND</i> has no effect.</p> <p>If a process is waiting for <i>s</i> (that is, has terminated execution with the routine <i>WAIT</i>), then control will be transferred to it.</p> <p>If several processes are waiting, then the first one in the queue will get control.</p> <p>Before a signal is used it must be initialized using the procedure <i>Init</i> in <i>Processes</i>.</p> <p>To find out if processes are waiting for a signal, use the routine <i>Awaited</i> in <i>Processes</i>.</p>

See Also Awaited
 Init
 Processes
 StartProcess
 WAIT

Example Send the signal done:

```
SEND(done);
```

SET type R

Description A set-type variable can assume as a value a set of up to 16 different elements with ordinal values in the range of 0 to 15.

Usage Set Type = **SET OF** Base Type;

Set Assignment := Base Type **o{** Base Type Elements **o}** ;

Base type can be enumeration or subrange.

Base-type elements are elements of the base type.

Comments A *set* is a collection of *elements* that you handle as a whole.

Besides assigning a value to a set, you can also include another element from the base type into the set, or exclude (or remove) an element from the set using the standard procedures *INCL* and *EXCL*.

The operators applicable to set expressions are the set operators and the relational operator *IN*. These operators allow you to check if a particular element is present in a set, to find the difference between two sets, and so forth.

There is a predefined set type called BITSET that is defined as *BITSET* = *SET OF {0..15}*. With this set type, you can define sets with values made up of combinations of CARDINAL numbers in the range 0 to 15. A set type value requires 2 bytes for storage in memory.

See Also**INCL****EXCL****Example**

Suppose you are an automobile manufacturer and wish to record the state of completion of each vehicle. You could define a set of the main parts required, like so:

TYPE**Components** = (Body, Engine, Chassis);**Compset** = **SET OF** Components**VAR**

State: Compset

The enumeration *Components* is called the base type of the set *State*. The set *State* can take as a value any combination of elements from the base type. For example, all of the following are possible assignments:

```

State:= Compset(Body,Chassis);
(* State has two elements Body and Chassis *)

State:= Compset(Body..Chassis);
(* State has all three elements *)

State:= Compset{ };
(* State has no elements *)

MODULE Sets;
TYPE
  Components = (Body,Engine,Chassis);
  Compset= SET OF Components;
VAR
  State: Compset;
BEGIN
  State := Compset(Body,Engine);
END Sets.

```

SetCol procedure (*Texts*) L

Description	<i>SetCol</i> advances the write position of a <i>TEXT</i> file to a given column.
Declaration	PROCEDURE SetCol(<i>t</i> : TEXT; <i>column</i> : CARDINAL);
Usage	<i>SetCol(t, column);</i>
	<i>SetCol</i> must be imported from the library module <i>Texts</i> . <i>t</i> must be of type <i>TEXT</i> imported from <i>Texts</i> . <i>column</i> must be of type <i>CARDINAL</i> .

Comments *SetCol* has no effect on input. The current column can be queried by use of the function *Col* exported from *Texts*.

See Also	Col Texts
Example	Cause the string »Hello world !« to be output in the twentieth column:
<pre>SetCol(output, 20); WRITE("Hello world !");</pre>	
<p>Example of tabbing to next multiple of 10:</p> <pre>SetCol(output, (Col(output) DIV 10 + 1) * 10)</pre>	

SetPos procedure (Files) L

Description	Sets the current byte position of a file.
Declaration	PROCEDURE SetPos(<i>f</i> : FILE; <i>pos</i> : LONGINT);
Usage	<i>SetPos(f, bytupos);</i>
<p><i>SetPos</i> must be imported from the library module <i>Files</i>.</p>	
<p><i>f</i> must be of type <i>FILE</i> imported from <i>Files</i>.</p>	
<p><i>bytupos</i> must be of type <i>LONGINT</i>.</p>	
Comments	The procedure is used for random access of data. If all the file elements of a certain file are of type <i>t</i> , then the file can be positioned at the <i>n</i> th element by the call
<pre>SetPos(<i>f</i>, LONG(<i>n</i>)*LONG(SIZE(<i>t</i>)))</pre>	
See Also	Files NextPos

Example Read the fifteenth widget from a file that has a header:

```
MODULE RandomAccess;
TYPE
  widget = RECORD .... END ;
  header = RECORD .... END ;
VAR
  currentWidget : widget;
  pos : LONGINT;
BEGIN
  pos := 15L*LONG(SIZE(widget)) + LONG(SIZE(header));
  SetPos(widgetFile, pos);
  ReadRec(widgetFile, currentWidget);
END RandomAccess.
```

SIGNAL type (Processes) L

Description *SIGNAL* references control loosely coupled processes or coroutines. A *SIGNAL* may be *Sent* or *Awaited*.

Declaration **TYPE**
 SIGNAL;

Usage **VAR**
 s: SIGNAL;

SIGNAL must be imported from the module *Processes*.

Comments It is mandatory that a *SIGNAL* variable be initialized by a call to the procedure *Init* in *Processes*.

See Also Processes
 SYSTEM

Example

```
MODULE Signals;
FROM Processes IMPORT Init, SIGNAL;
VAR s: SIGNAL;
BEGIN
  Init(s);
END Signals.
```

Sin function (MathLib, LongMath) L

Description	<i>Sin</i> returns the sine of <i>X</i> , where <i>X</i> is expressed in radians.
Declaration	PROCEDURE Sin(<i>x</i> : REAL): REAL; PROCEDURE Sin(<i>x</i> : LONGREAL): LONGREAL;
Usage	<i>Y</i> := Sin(<i>X</i>);
	Argument <i>X</i> and result <i>Y</i> are either both REAL or both LONGREAL.
Comments	If you want the argument expressed in degrees instead of radians, use the expression Result:= Sin(<i>X</i> * 3.141592/180.0) If you need to calculate the inverse of sine, you can use the <i>Arctan</i> function in the following expression: Arcsin:= Arctan(<i>X</i> /Sqrt(1.0 - <i>X</i> * <i>X</i>)).
See Also	Cos MathLib

Example

```
MODULE Sine;
FROM MathLib IMPORT Sin;
CONST
  Degrads = 3.141592/180.0;
VAR
  Degrads, Radians, sinAngle: REAL;
BEGIN
  Radians:= Degrads * Degrads;
  sinAngle:= Sin(Radians);
END Sine.
```

SIZE standard procedure S

Description *SIZE* returns the minimum required storage in bytes of a type or variable.

Usage *SIZE(x);*

x can be a variable of any type or a type identifier.

Comments The *SIZE* of a CHAR, BOOLEAN, or BYTE variable is 1; WORD, INTEGER, and CARDINAL variables are 2; REAL and LONGINT variables are 4; LONGREAL variables are 8.

Note that byte-size variables are only packed in arrays and thus always take up one word of storage, unless they are in an array.

See Also TSIZE

Example Multiply 2 x 4:

```
MODULE Sizes;
TYPE
  rec = RECORD
    a,b: CHAR;
  END ;
  array= ARRAY [0..1] OF CHAR;
VAR
  r: rec;
  a: array;
BEGIN
  WRITELN(SIZE(r), SIZE(a));
END Sizes.
```

The output is

4 2

which shows that arrays are packed and records are not.

SpecialOps type (Terminal) L

Description *SpecialOps* is an enumeration of terminal operations that are not available on all terminals. It is used to describe a set of operations available on the currently installed terminal.

Declaration **TYPE**
SpecialOps = (clearEol, insertDelete, highlightnormal);
OpSet : **SET OF** *SpecialOps*;

Usage **VAR**
available : *OpSet*;

available is a predefined variable declared as type *OpSet*; it must be imported from the module *Terminal*.

Comments	This variable is useful for writing terminal-independent programs. The value of <i>available</i> is set by the installation program INSTM2.
See Also	<i>available</i> <i>OpSet</i> <i>Terminal</i>
Example	<pre>MODULE TermStuff; FROM Terminal IMPORT available, SpecialOps, ClearToEOL, GotoXY, numCols; VAR I, CurCol, CurRow: CARDINAL; BEGIN IF clearEol IN available THEN ClearToEOL ELSE FOR I := CurCol TO numCols DO WRITE(' ') END ; GotoXY(CurCol,CurRow) END ; WRITE('A line of text.'); END TermStuff.</pre>

Sqrt function (MathLib, LongMath) L

Description	<i>Sqrt</i> returns the square root of positive <i>X</i> .
Declaration	PROCEDURE Sqrt(x: REAL): REAL; PROCEDURE Sqrt(x: LONGREAL): LONGREAL;
Usage	<i>Y:= Sqrt(X);</i> Argument <i>X</i> and result <i>Y</i> are either both REAL or both LONGREAL. The argument must be positive or zero.
Comments	The argument <i>X</i> must be positive or zero; otherwise, the exception ArgumentError is raised.

See Also **MathLib**
 LongMath

Example **MODULE** SquareRoot;
 FROM MathLib **IMPORT** Sqrt;
 CONST
 Pi = 3.141592;
 VAR
 CircleArea, CircleRadius: REAL;
 BEGIN
 READ (CircleArea);
 IF CircleArea >= 0.0 **THEN**
 CircleRadius:= Sqrt(CircleArea/Pi)
 END
 END SquareRoot.

StartProcess procedure (Processes) L

Description *StartProcess* allocates space to a process and transfers control to it.

Declaration **PROCEDURE** StartProcess(P: PROC; n: CARDINAL);

Usage StartProcess(proc, workSpaceSize);

StartProcess must be imported from the library module *Processes*.

proc must be of type *PROC*.

workSpaceSize must be of type *CARDINAL*.

Comments	<p><i>proc</i> is a procedure without RETURNS. Normally, its statements will be enclosed in a LOOP statement without EXITS. Control passes to and from the process by calls to <i>SEND</i> and <i>WAIT</i>.</p>
	<p>This routine is implemented using the procedures <i>NEWPROCESS</i> and <i>TRANSFER</i> of the pseudomodule <i>SYSTEM</i>.</p>
See Also	Awaited Init Processes SEND WAIT

Example Start the process *service1*:

```
StartProcess(service1,200);
```

StatusError exception (Files) L/E

Description	StatusError is raised by the <i>Files</i> module when a <i>FILE</i> variable is used incorrectly; for example, when reading a unopened file or opening a file twice.
Declaration	EXCEPTION StatusError;
Usage	EXCEPTION StatusError: WRITELN('File error'); END
	StatusError must be imported from the module <i>Files</i> .
Comments	Use the StatusError exception to prevent your program from crashing due to user input.

See Also	EndError UseError DeviceError DiskFull
Example	<pre>MODULE ReadFile; FROM Files IMPORT FILE, Open, ReadWord, Close, EndError, StatusError; VAR f: FILE; c: CARDINAL; BEGIN IF Open(f,'testfile.dat') THEN ReadWord(f,c); END ; ReadWord(f,c); Close(f) EXCEPTION EndError: WRITELN('End of file reached.');?> StatusError: WRITELN('Error file not opened'); END ReadFile.</pre>

STORAGE pseudomodule L

Description	<i>STORAGE</i> is a pseudomodule for managing heap memory in general and dynamic variables in particular.
Declaration	<pre>DEFINITION MODULE STORAGE; FROM SYSTEM IMPORT ADDRESS; PROCEDURE ALLOCATE(VAR A:ADDRESS; Size: CARDINAL); PROCEDURE DEALLOCATE(VAR A:ADDRESS; Size: CARDINAL); PROCEDURE MARK(VAR A: ADDRESS); PROCEDURE RELEASE(VAR A: ADDRESS); PROCEDURE FREEMEM():CARDINAL END STORAGE.</pre>

Comments	The procedures <i>NEW</i> and <i>DISPOSE</i> can be considered abbreviated forms of <i>ALLOCATE</i> and <i>DEALLOCATE</i> . To use the standard procedures <i>NEW</i> and <i>DISPOSE</i> , <i>ALLOCATE</i> and <i>DEALLOCATE</i> , respectively, must be imported.
See Also	MARK NEW RELEASE

(String type (Strings) L

Description	<i>String</i> is a predefined array of characters.
Declaration	TYPE <i>String</i> = ARRAY [0..80] OF CHAR;
Usage	VAR <i>s1, s2</i> : <i>String</i> ;
	<i>String</i> must be imported from the library module <i>Strings</i> .
Comments	Since <i>String</i> is a character array, a variable of this type may be assigned and compared with any other character array variable.
	Even though the module <i>Strings</i> exports the type <i>String</i> none of the procedures require this type. Instead, they use open array parameters that will accept any array of characters.
See Also	<i>Strings</i>
Example	<i>v1</i> and <i>v2</i> are defined as strings: FROM <i>Strings IMPORT</i> <i>String</i> ; VAR <i>v1,v2</i> : <i>String</i> ;

StringError exception (Strings) L/E

Description **StringError** is raised by the module *Strings* when the destination of a string operation is not large enough to hold the result.

Declaration

```
EXCEPTION
  StringError;
```

Usage

```
EXCEPTION
  StringError WRITELN('String error');
END
```

StringError must be imported from the module *Strings*.

Comments This exception occurs only during calls to the procedures exported by *Strings*.

See Also *Strings*

Example

```
MODULE StringTooShort;
FROM Strings IMPORT StringError, Append;
VAR
  s1, s2 : ARRAY [0..12] OF CHAR;
BEGIN
  s1 := 'Hello there';
  s2 := 'everyone';
  Append(s2,s1);
EXCEPTION
  StringError:
    WRITELN('Destination string too short.');
END StringTooShort.
```

Strings module L

Description *Strings* is a module for the manipulation of strings.

Declaration

```
DEFINITION MODULE Strings;
TYPE String = ARRAY [0..80] OF CHAR;
PROCEDURE Length  (VAR str: ARRAY OF CHAR):
                  CARDINAL;
PROCEDURE Pos      (substr,str: ARRAY OF CHAR):
                  CARDINAL;
PROCEDURE Insert   (substr: ARRAY OF
                     CHAR; VAR str: ARRAY OF CHAR;
                     inx:CARDINAL);
PROCEDURE Delete   (VAR str: ARRAY OF CHAR; inx,len:
                     CARDINAL);
PROCEDURE Append   (substr: ARRAY OF
                     CHAR; VAR str: ARRAY OF CHAR);
PROCEDURE Copy     (VAR str: ARRAY OF CHAR; inx,len:
                     CARDINAL;
                     VAR result: ARRAY OF CHAR);
PROCEDURE CAPS     (VAR str: ARRAY OF CHAR);
EXCEPTION StringError;
END Strings.
```

Comments These procedures exported from *Strings* work on any array of characters, not just on variables declared as type *String*.

These procedures use open arrays; that is, arrays of any size. You should, however, restrict yourself to arrays with indexing that start at zero, since the various *inx* parameters listed earlier work relative to zero.

See Also *String*

StrToCard procedure (Convert) L

Description *StrToCard* performs conversion from a string variable to a cardinal variable.

Declaration **PROCEDURE** StrToCard(**VAR** s: **ARRAY OF CHAR**; **VAR** CARDINAL): BOOLEAN;

Usage *okay* := StrToCard(s, c);

StrToCard must be imported from the library module *Convert*.

okay must be of type BOOLEAN.

s must be a character array.

c must be of type CARDINAL.

Comments *okay* is assigned TRUE if the string represents a valid cardinal number; otherwise, it is assigned FALSE.

See Also CardToStr
 Strings

Example Read a cardinal *c* out of a string *str*:

```
str := '12323';
IF NOT StrToCard(str, c) THEN
  WRITELN(str,'is not convertible.');
END ;
```

StrToDouble procedure (Doubles) L

Description *StrToDouble* performs conversion from a string variable to a double-precision real variable.

Declaration **PROCEDURE** StrToDouble(**VAR** s: **ARRAY OF CHAR**;
 VAR d: **LONGREAL**): **BOOLEAN**;

Usage **okay** := StrToDouble(s, d);

StrToDouble must be imported from the library module *Doubles*.

okay must be of type **BOOLEAN**.

s must be a character array.

d must be of type **LONGREAL**.

Comments *okay* is assigned TRUE if the string represents a valid real number; otherwise, it is assigned FALSE.

See Also Doubles
 DoubleToStr

Example Read a double-precision *d* out of a string *str*:

```
str := '2.99D10';
IF NOT StrToDouble(str, d) THEN
  Writeln (str,'is not convertible');
END ;
```

StrToInt procedure (Convert) L

Description *StrToInt* performs conversion from a string variable to an integer variable.

Declaration **PROCEDURE StrToInt (VAR s: ARRAY OF CHAR;**
 VAR i: INTEGER): BOOLEAN;

Usage *okay* := **StrToInt(s, i);**

StrToInt must be imported from the library module *Conver*

okay must be of type BOOLEAN.

s must be a character array.

i must be of type INTEGER.

Comments *okay* is assigned TRUE if the string represents a valid integer otherwise, it is assigned FALSE.

See Also **IntToStr**
Strings

Example Read an integer called *i* out of a string:

```
str := " 30000 - This number is almost too big ";
okay := StrToInt(str, i);
```

(* *i* has got 30000 now *)

StrToInt procedure (Convert) L

Description *StrToInt* performs conversion from a string variable to a long integer variable.

Declaration **PROCEDURE StrToInt(s: ARRAY OF CHAR;**
 VAR l: LONGINT): BOOLEAN;

Usage *okay* := *StrToInt*(*s*, *l*);

StrToInt must be imported from the library module *Convert*.

okay must be of type BOOLEAN.

s must be a character array.

l must be of type LONGINT.

Comments *okay* is assigned TRUE if the string represents a valid long integer; otherwise, it is assigned FALSE.

See Also [LongToStr](#)
[Strings](#)

Example Read a longer integer (see the example in *StrToInt*) out of a string:

```
VAR
  str: ARRAY [0..20] OF CHAR;
  l: LONGINT;
BEGIN
  str := "40000";
  okay := StrToInt(str, l);
END
```

StrToReal procedure (Convert) L

Description *StrToReal* performs conversion from a string variable to a real-number variable.

Declaration **PROCEDURE** StrToReal(**VAR** s: **ARRAY OF CHAR**;
VAR r: **REAL**): **BOOLEAN**;

Usage *okay* := StrToReal(*s*, *r*);

StrToReal must be imported from the library module *Convert*.

okay must be of type **BOOLEAN**.

s must be a character array.

r must be of type **REAL**.

Comments *okay* is assigned TRUE if the string represents a valid real; otherwise, it is assigned FALSE.

See Also **RealToStr**
Strings

Example Make Pi out of a string:

```
stringyPi := "3.14159265";
okay := StrToReal(stringyPi, Pi);
(* Pi is now Pi *)
```

SYSTEM pseudomodule L

Description **SYSTEM** is a pseudomodule that provides low-level support.

Declaration **DEFINITION MODULE SYSTEM**

```
TYPE
  WORD; BYTE; ADDRESS; PROCESS;
  VAR
    IORESULT, HLRESULT: CARDINAL;

  PROCEDURE ADR(VAR v: AnyType): ADDRESS;
  PROCEDURE TSIZE(AnyType): CARDINAL;

  PROCEDURE TRANSFER(VAR source, dest: PROCESS);
  PROCEDURE IOTRANSFER(VAR source, dest: PROCESS;
    n: CARDINAL);
  PROCEDURE NEWPROCESS(p: PROC; a: ADDRESS;
    n: CARDINAL; VAR q: PROCESS);

  PROCEDURE BIOS(n:CARDINAL w: WORD): CARDINAL;
  PROCEDURE BDOS(n:CARDINAL w: WORD): CARDINAL;
  PROCEDURE CODE(AnyStringLiteral);

  PROCEDURE MOVE(source, destin: ADDRESS;
    nbytes: CARDINAL);
  PROCEDURE FILL(start: ADDRESS; nbytes: CARDINAL;
    ch: BYTE);
  PROCEDURE INP(port: WORD): CARDINAL;
  PROCEDURE OUT(port: WORD; outByte: WORD);

  EXCEPTION OVERFLOW, READOVERLOW, OUTOFCMEMORY,
    BADOVERLAY;

  END SYSTEM.
```

Comments

For explanations of the various procedures, see the specific entries. In general, use of these procedures make a program much less portable.

termCH variable (InOut) L

Description *termCH* always contains the last character read by the module *InOut*.

Declaration **VAR**
 termCH: CHAR;

Usage *LastChar* := *termCH*;

termCH must be imported from the module *InOut*.

Comments *termCH* contains the last character read by any of the five read procedures in *InOut*.

See Also *InOut*
 Done
 Texts

Example **MODULE** *LastChar*;
 FROM *InOut* **IMPORT** *ReadCard*, *termCH*;
 VAR *ch*: CHAR; *c*: CARDINAL;
 BEGIN
 ReadCard(*c*);
 WRITELN('The cardinal read is ',*c*);
 WRITELN('The last character of the cardinal is
 ',*termCH*);
 END *LastChar*.

Terminal module L

Description *Terminal* provides input/output to the terminal.

```
Declaration      DEFINITION MODULE Terminal;

(* Terminal Input *)

PROCEDURE ReadChar(VAR ch: CHAR);
PROCEDURE BusyRead(VAR ch: CHAR);
PROCEDURE ReadAgain;
PROCEDURE ReadLine(VAR s: ARRAY OF CHAR);

(* Screen output *)

VAR numRows, numCols:CARDINAL

PROCEDURE WriteChar(ch: CHAR);
PROCEDURE WriteLn;
PROCEDURE WriteString(s: ARRAY OF CHAR);
PROCEDURE ClearScreen;
PROCEDURE GotoXY(col,row:CARDINAL);

PROCEDURE InitScreen;
PROCEDURE ExitScreen;

TYPE
  SpecialOps = (clearEol, insertDelete,
  highlightNormal);
  OpSet    = SET OF SpecialOps;

VAR
  available: OpSet;

PROCEDURE ClearToEOL;
PROCEDURE InsertLine;
PROCEDURE DeleteLine;
PROCEDURE Highlight;
PROCEDURE Normal;

END Terminal.
```

Comments This module, or at least the screen manipulation procedures, will work correctly only if Turbo Modula-2 has been properly installed on your terminal.

TEXT data structure (Texts) L

Description *TEXT* is a type that represents legible disk files.

Declaration **TYPE** TEXT = [1..16];

Usage **VAR** t: TEXT;

TEXT must be imported from the library module **Texts**.

Comments A *TEXT* is a character stream. For example, a WordStar document file is better handled as a *FILE*, while a nondocument file can be manipulated as a *TEXT*.

There are a variety of procedures to handle *TEXT* objects.

See Also **Texts**

Example Simulate the CP/M TYPE utility:

```

MODULE type;
FROM ComLine IMPORT PromptFor;
FROM Texts IMPORT TEXT, ReadString, OpenText,
              EOT, ReadChar, WriteChar,
              output;
VAR
  textName: ARRAY [0..20] OF CHAR;
  t      : TEXT;
  c      : CHAR;
  okay   : BOOLEAN;
BEGIN
  PromptFor( "File to type: ",textName);
  IF OpenText(t, textName) THEN
    WHILE NOT EOT(t) DO
      ReadChar(t, c);
      WriteChar(output, c);
    END
  ELSE
    WRITELN(textName,' not found!');
  END
END type.
```

TextDriver type (Texts) L

Description *TextDriver* must be used when installing your own character device drivers. A procedure of this type may be passed to *ConnectDriver*.

Declaration **TYPE** TextDriver = **PROCEDURE** (TEXT,**VAR** CHAR);

Usage **VAR**
 MyDriver: TextDriver;

TextDriver must be imported from the module *Texts*.

Comments You do not need to import the type *TextDriver*, but the procedure you pass to *ConnectDriver* must match it.

See Also	Texts ConnectDriver
Example	<pre>MODULE CharDriver; FROM Texts IMPORT TEXT, ConnectDriver, TextDriver; VAR t:TEXT; PROCEDURE MyDriver(t: TEXT; VAR ch: CHAR); BEGIN END MyDriver; BEGIN ConnectDriver(t,MyDriver); END CharDriver.</pre>

TextFile procedure (Texts) L

Description	<i>TextFile</i> returns a value of type <i>FILE</i> , as described in the module <i>Files</i> . Since a text file is implemented as a <i>Files.FILE</i> , you may obtain this handle to perform »file« operations on a <i>TEXT</i> .
Declaration	PROCEDURE TextFile(<i>T</i> :TEXT): FILE;
Usage	<i>f</i> := TextFile(<i>t</i>);
	<i>TextFile</i> must be imported from the module <i>Texts</i> .
	<i>t</i> must be of type <i>TEXT</i> .
	<i>f</i> must be of type <i>FILE</i> .
Comments	If the file has not been opened, <i>TextFile</i> returns <i>NIL</i> .
See Also	Texts Files

Example

```
MODULE SizeOfAText;
FROM Texts IMPORT TEXT, TextFile;
FROM Files IMPORT FileSize;
VAR t: TEXT;
BEGIN
  WRITE('The size of the text file is ');
  WRITELN(FileSize(TextFile(t)), ' bytes.');
END SizeOfAText.
```

(TextNotOpen exception (Texts) L/E)

Description	TextNotOpen is raised by the module <i>Texts</i> when an attempt is made to use an unopened text.
Declaration	EXCEPTION TextNotOpen;
Usage	EXCEPTION TextNotOpen: WRITELN('Text not open'); END
	TextNotOpen must be imported from the module <i>Texts</i> .
Comments	This exception can be raised by all input/output procedures and the procedure <i>SetCol</i> .
See Also	Texts TooManyTexts

Example

```

MODULE TextNotOpen;
FROM Texts IMPORT TEXT, OpenText, TextNotOpen;
VAR
    t      : TEXT;
    filename : ARRAY [0..15] OF CHAR;
    x      : REAL;
BEGIN
    IF OpenText(t,filename) THEN END ;
    READ(t,x);
EXCEPTION
    TextNotOpen: WRITELN( "Can't read an unopen text." );
END TextNotOpen.
```

Texts module L

Description *Texts* manipulates legible disk files.

Declaration

```

DEFINITION MODULE Texts;
FROM Files IMPORT FILE;
```

```

CONST EOL=36C;
TYPE TEXT = [1..16];
VAR input,output,console: TEXT;

PROCEDURE ReadChar   (t: TEXT; VAR ch: CHAR);
PROCEDURE ReadString (t: TEXT; VAR s: ARRAY OF
                      CHAR);
PROCEDURE ReadInt    (t: TEXT; VAR i: INTEGER);
PROCEDURE ReadCard   (t: TEXT; VAR c: CARDINAL);
PROCEDURE ReadLong   (t: TEXT; VAR l: LONGINT);
PROCEDURE ReadReal   (t: TEXT; VAR r: REAL);
PROCEDURE ReadLine   (t: TEXT; VAR s: ARRAY OF
                      CHAR);
PROCEDURE ReadAgain  (t: TEXT);
PROCEDURE ReadLn    (t: TEXT);

PROCEDURE WriteChar  (t: TEXT; ch: CHAR);
PROCEDURE WriteString(t: TEXT; s: ARRAY OF CHAR);
```

```
PROCEDURE WriteInt      (t: TEXT; i:INTEGER;
                        n:CARDINAL);
PROCEDURE WriteCard     (t: TEXT; c, n:CARDINAL);
PROCEDURE WriteLong     (t: TEXT; l:LONGINT;
                        n:CARDINAL);
PROCEDURE WriteReal     (t: TEXT; r:REAL; n:CARDINAL;
                        digits:INTEGER);
PROCEDURE WriteLn       (t: TEXT);

PROCEDURE Done           (t: TEXT):BOOLEAN;
PROCEDURE EOLN           (t: TEXT):BOOLEAN;
PROCEDURE EOT            (t: TEXT):BOOLEAN;

PROCEDURE OpenText        (VAR t: TEXT; name: ARRAY OF
                           CHAR):BOOLEAN;
PROCEDURE CreateText      (VAR t: TEXT; name: ARRAY OF
                           CHAR);
PROCEDURE CloseText       (VAR t: TEXT);

PROCEDURE Col             (t: TEXT): CARDINAL;
PROCEDURE SetCol          (t: TEXT; column:CARDINAL);

PROCEDURE TextFile        (t: TEXT): FILE;

TYPE TextDriver = PROCEDURE(TEXT, VAR CHAR);

PROCEDURE ConnectDriver(VAR t: TEXT; p: Text
                        Driver);

PROCEDURE Init; (* used only by system *)

VAR haltOnControlC:BOOLEAN; (*TRUE by default*)

EXCEPTION TextNotOpen, TooManyTexts;

END Texts.
```

Comments Rather than use the various read and write procedures in this module directly, the user is encouraged to use the standard procedures *READ*, *READLN*, *WRITE*, and *WRITELN*. These are translated by the compiler into calls to the appropriate procedures in *Texts*.

TooLarge exception (Convert) L/E

Comments Rather than use the various read and write procedures in this module directly, the user is encouraged to use the standard procedures *READ*, *READLN*, *WRITE*, and *WRITELN*. These are translated by the compiler into calls to the appropriate pro-

Description *TooLarge* indicates that a conversion from a number to a string cannot be done because the string representation of the number is too long for the destination string.

Declaration **EXCEPTION** TooLarge;

Usage **EXCEPTION**
 TooLarge:
 WRITELN('could not convert the number to string');
END

Comments This exception eliminates the necessity to check user input after every user entry. A common exception handler can be written that will trap all such conversion errors.

See Also Convert

Example **MODULE** InputTooLarge;
FROM Convert **IMPORT** IntToStr, TooLarge;
VAR
 i: CARDINAL; *s*: **ARRAY** [0..2] **OF** CHAR;

```

BEGIN
  WRITE('Give me an cardinal less than 100');
  READLN(i);
  CardToStr(i,s);
  WRITELN('The string is ',s);
EXCEPTION
  Toolarge: WRITELN('That was too large!');
END InputTooLarge.

```

oManyTexts exception (Texts) L/E

Description TooManyTexts is raised by the module *Texts* when an attempt is made to open a text and there are already 16 open.

Declaration **EXCEPTION**
 TooManyTexts;

Usage **EXCEPTION**
 TooManyTexts: WRITELN('Too many text files');
END

TooManyTexts must be imported from the module *Texts*.

Comments There are three texts that are always open: *console*, *input*, and *output*. These three texts are part of the 16 allowable texts.

See Also Texts
 TextNotOpen

Example MODULE TooManyTextFiles;
 FROM Texts IMPORT TEXT, CreateText, TooManyTexts;
 VAR
 t: ARRAY [0..15] OF RECORD
 t: TEXT;
 filename: ARRAY [0..15] OF
 CHAR;
 END ;
 I: CARDINAL;

```

BEGIN
  FOR I := 0 TO 15 DO
    CreateText(t[I].t,t[I].filename);
  END
EXCEPTION
  TooManyTexts: WRITELN("Can't open another text.");
END TooManyTextFiles.

```

TRANSFER procedure (SYSTEM) L

Description	<i>TRANSFER</i> allows cooperative control between processes.
Declaration	PROCEDURE TRANSFER(VAR source, dest: PROCESS);
Usage	TRANSFER(source, destination);
	<i>TRANSFER</i> must be imported from the pseudomodule SYSTEM .
	<i>source</i> must be of type PROCESS imported from SYSTEM .
	<i>destination</i> must be of type PROCESS imported from SYSTEM .
Comments	The <i>source</i> variable will be assigned to the process where the <i>TRANSFER</i> call is made. This permits the called process to return control to the caller, if desired. Before you transfer to a process it must have first been set up by a call to <i>NEWPROCESS</i> .
	We recommend using the module <i>Processes</i> where possible, since it provides rudimentary signaling support upon which to base scheduling.
See Also	IOTRANSFER NEWPROCESS PROCESS

Example Mark time:

```
MODULE TickTock;
FROM SYSTEM IMPORT
  NEWPROCESS, TRANSFER, ADDRESS, PROCESS, WORD;
VAR
  tickProcess, tockProcess, main: PROCESS;
  tickWork, tockWork : POINTER TO ARRAY [0..99] OF
  WORD;
PROCEDURE tick;
BEGIN
  WRITELN('tick');
  TRANSFER(tickProcess, tockProcess);
END ;

PROCEDURE tock;
BEGIN
  WRITELN('tock');
  TRANSFER(tockProcess, tickProcess);
END tock;

BEGIN
  NEWPROCESS(tick, tickWork, SIZE(tickWork ^ ), 
    tickProcess);
  NEWPROCESS(tock, tockWork, SIZE(tockWork ^ ), 
    tockProcess);
  TRANSFER(main, tickProcess); (* Start of
  tick-tock *)
  (* This routine can be hard to stop *)
END TickTock.
```

TRUE standard value S

Description TRUE denotes the Boolean state of truth.

Usage Finished := TRUE;

Finished is a variable of type BOOLEAN.

Comments	The ordinal value of TRUE is 1; thus, <i>WRITE(CARDINAL(TRUE))</i> will print a 1. In contrast, the ordinal value of FALSE is 0; thus, truth is greater than falsity.
See Also	BOOLEAN FALSE
Example	MODULE Truth; VAR b: BOOLEAN; BEGIN REPEAT b := TRUE UNTIL b; (* does not repeat *) END Truth.

TRUNC standard function S

Description	<i>TRUNC</i> converts from REAL type to CARDINAL type.
Usage	$Y := \text{TRUNC}(X);$ Argument <i>X</i> is of type REAL. The argument must be in the CARDINAL (0 to 65535) range. Result <i>Y</i> is of type CARDINAL.
Comments	You must make sure that the argument is in the CARDINAL range, otherwise an overflow error will occur. The argument is truncated (any fractional part is removed); for example, the number +16.74 will become a CARDINAL +16.
	Note that truncating a number is not the same as rounding the number to the nearest whole value; for example, 9.99999 will be truncated to 9, not rounded to 10. If you wish to round a REAL value to the nearest whole value, use the expression <i>Rounded:= TRUNC(X + 0.5)</i> .
	The inverse conversion from CARDINAL or INTEGER to REAL uses the <i>FLOAT</i> function.

TRUNC returns a CARDINAL result; therefore, its argument must be nonnegative. The result can be assigned to INTEGERS as well, since INTEGER and CARDINAL are assignment compatible.

Example

```
VAR
  Values: ARRAY 1..10] OF REAL;
  Rounded: ARRAY [1..10] OF CARDINAL;
  I: CARDINAL;
BEGIN
  FOR I:= 1 TO 10 DO
    Rounded[I]:= TRUNC(Values[I] + 0.5)
  END ;
```

TSIZE procedure (SYSTEM) L

Description TSIZE returns the storage requirements of a data type in bytes.

Declaration PROCEDURE TSIZE(AnyType): CARDINAL;

Usage bytes := TSIZE(t);

TSIZE must be imported from the pseudomodule *SYSTEM*.

bytes must be of type CARDINAL.

t may be any type identifier.

Comments This function returns the size of a type; if you want the size of a specific variable, use *SIZE*.

See Also SIZE
 SYSTEM

Example realSize := TSIZE(REAL); (* realSize = 4 *)

TYPE declaration R

Description **TYPE** signals the start of a data-type definition section in the declaration part of a procedure or module.

Usage **TYPE newtype = oldTypeList;**

newtype can be any valid identifier.

oldTypeList is a combination of previously defined data types.

Comments This declaration allows you to define objects whether they comprise sets like Red, Green, Blue; subranges of basic types like all integers between 1950 and 2000; or combinations of any length.

Example Define an age type:

```
TYPE
  ageRange = [0..114];
```

Define a customer record:

```
TYPE
  personRec = RECORD
    firstName,
    middleName: ARRAY [0..12] OF CHAR;
    lastName : ARRAY [0..20] OF CHAR;
    age      : ageRange;
    married   : BOOLEAN;
    company   : ARRAY [0..30] OF CHAR;
  END ;
```

Define colors for an RGB monitor:

```
TYPE
  pixelColors = (redPix, greenPix, bluePix);
  color       = SET OF pixelColors;
```

```

CONST
  black  = color{ };
  red    = color{ redPix   };
  green  = color{ greenPix };
  blue   = color{ bluePix  };
  yellow = color{ redPix,  greenPix };
  purple = color{ redPix,  bluePix };
  cyan   = color{ greenPix, bluePix };
  white  = color{ redPix,  greenPix, bluePix };

VAR
  currentColor: color;
BEGIN
  currentColor := red; (* etc *)

```

UseError exception (Files) L/E

Description	UseError is raised by the <i>Files</i> module when an operation on a file is impossible because the disk is write-protected.
Declaration	EXCEPTION UseError;
Usage	EXCEPTION UseError: WRITELN('File error'); END
	UseError must be imported from the module <i>Files</i> .
Comments	Use the UseError exception to prevent your program from crashing when the user swaps disks unexpectedly.
See Also	ErrorHandler StatusError DeviceError DiskFull

Example

```

MODULE WriteFile;
FROM Files IMPORT
  FILE, Create, WriteWord, Close, EndError,
  StatusError, UseError;
VAR f: FILE; c: CARDINAL;
BEGIN
  Create(f,'testfile.dat');
  WriteWord(f,c);
  Close(f)
EXCEPTION
  | EndError: WRITELN('End of file reached.');?>
  | StatusError: WRITELN('Error file not opened');
  | UseError : WRITELN('Disk not logged in.');
END WriteFile.
```

VAL standard function S

Description *VAL* returns the value of type *T* with ordinal number *X*.

Usage *Y:= VAL(T,X);*

Argument *X* is of CARDINAL type.

Argument *T* is a CARDINAL, INTEGER, CHAR, BOOLEAN, or enumeration-type identifier.

Result *Y* is of type *T*.

Comments Note that *VAL(T,ORD(X)) = X* for any type *T*.

Example

```
TYPE
  Direction: (Forward,Backward,Up,Down);
VAR
  Move: Direction;
  Value: CARDINAL;
  Character: CHAR;
BEGIN
  Move:= VAL(Direction,2);
(* Move is value Up *)

  Value:= 43;
  Character:= VAL(CHAR,Value);
(* Character is value '+'*)
```

Notice that this example is equivalent to type conversion. The last example could have been accomplished just as well with `Character:= CHR(Value);`). `VAL` provides a generic method for all ordinal types.

VAR declaration R

Description `VAR` signals the start of a declaration section for variables inside the declaration part of a module or a procedure.

In the parameter list of a procedure, `VAR` specifies that a variable is a variable parameter.

Usage `VAR vlist : anytype;`

vlist is simply a list of identifiers.

anytype is any standard, imported, or user-defined type.

Comments Variable declarations may appear anywhere in the declaration section of procedures and modules.

Example Declare an integer *i*:

```
VAR i : INTEGER;
```

Declare a couple of characters:

```
VAR : carney,gleason : CHAR;
```

WAIT procedure (Processes) L

Description *WAIT* causes the currently executing process to stop and wait for a signal.

Usage `WAIT(siggy);`

WAIT must be imported from the library module *Processes*.

siggy must be of type *SIGNAL* imported from *Processes*.

Comments *WAIT* is used in conjunction with the other routines in *Processes*.

When *WAIT* is called from an executing *Process*, it causes the execution of that *Process* to pause. It is then inserted into a queue associated with the signal *siggy*. It will not be revived again until the signal *siggy* is sent with the procedure *SEND*.

See Also

Awaited

Init

SEND

StartProcess

WAIT

Example Wait for a signal:

```
WAIT(signal)
```

WHILE statement R

Description **WHILE** repeatedly executes a sequence of statements as long as a Boolean expression returns a TRUE result.

Usage **WHILE** BooleanExpression **DO**
 Statement sequence
END ;

Statement sequence = Statement {; Statement};

Comment The Boolean expression is evaluated before the statement sequence is executed; thus, the sequence is not necessarily executed.

The repetition is terminated by an expression that returns a FALSE result.

Example **MODULE** While;
CONST
 N = 10;
VAR
 I, Find: CARDINAL;
 DataArray: **ARRAY** [0..N] **OF** CARDINAL;

BEGIN
 WRITE("Input value to find: ");
 READ(Find);
 I := 0;
 WHILE (I < N) & (DataArray[I] # Find) **DO**
 I := I + 1
 END
 END While.

WITH statement R

Description **WITH** aids access to record fields.

Usage **WITH** Designator **DO**
 Statement sequence
END ;

Designator is of record type.

Statement sequence = Statement {; Statement} ;

Comments There are two ways to access the values of record fields.

The first method states explicitly both the record identifier and the field:

Person.Name := "Alfred Smith"

The other method uses the **WITH** statement:

```
WITH Person DO
  Name := "Alfred Smith";
  Age := 34;
END ;
```

This second method provides a clearer notation, especially for large records or records with nested record fields.

Example

```

MODULE With;
TYPE
  PersonRecord = RECORD
    Name: ARRAY [1..20] OF CHAR;
    Age: [0..200]
  END ;

VAR
  Patient: PersonRecord;
BEGIN *
  WITH Patient DO
    WRITE("Enter patient's name: ");
    READ(Name);
    WRITE("Enter patient's age: ");
    READ(Age)
  END ;
  WITH Patient DO
    WRITE(Name, " is ", Age, " years old.")
  END
END With.

```

WORD type (SYSTEM) L

Description *WORD* is a low-level type with variables that occupy one word (2 bytes) of memory. All variables occupying one word of storage are assignment-compatible with type *WORD*.

Usage **PROCEDURE** p1 (w : WORD);

WORD must be imported from the pseudomodule *SYSTEM*.

Comments *WORD* is compatible with **INTEGER**, **CARDINAL**, **BITSET**, **SET**, enumeration types, **POINTER** types (including **ADDRESS**), **BOOLEAN**, and **CHAR**.

Because *BYTE* is a subrange of *WORD*, those variables compatible with *BYTE* are also compatible with *WORD*.

ARRAY OF WORD is compatible with every type. It may be used as a special open array parameter type that allows any variable to be passed as the actual parameter.

Only assignment operations are allowed on variables of type **WORD**; however, type-transfer functions may be applied to type **WORD**.

Example Dump a variable:

```

MODULE LowLevelDump;
FROM InOut IMPORT WriteHex;
FROM SYSTEM IMPORT WORD,ADR;

PROCEDURE DumpVar(VAR v:ARRAY OF WORD);
VAR
  i :INTEGER
BEGIN
  WRITELN('There are ',HIGH(v)+1,' words allocated
for v');
  FOR i := 0 TO HIGH(v) DO
    WRITE(' Address: $');
    WriteHex(ADR(v[i]),4);
    WRITE(i:3,' Value: $');
    WriteHex(v[i],4);
    WRITELN;
  END ;
END DumpVar;

VAR
  r: REAL;
  rec: RECORD
    a,b,c,d: REAL;
  END ;
BEGIN
  DumpVar(r);
  DumpVar(rec);
END LowLevelDump.

```

WRITE standard procedure S/E

Description *WRITE* is a generalized write statement.

Usage *WRITE(vlist)*

or

WRITE(t, vlist);

t must be of type *TEXT* imported from *Texts*.

vlist is simply a list of variables or constants and optional formatting parameters.

Comments If *t* is not specified, the standard output text *output* will be used.

WRITE will be translated by the compiler into the appropriate calls to the module *Texts*.

WRITELN is the same as *WRITE*, except *WRITELN* starts a new line when the output is complete.

Although *WRITE*, *WRITELN*, *READ*, and *READLN* are not included in the Modula-2 definition, programming without them can be very tedious; however, they are not portable.

One strategy might be to develop the code using these statements for debugging and development output, and use the explicit calls to *Texts* for the actual programmed output.

See Also **READ**
READLN
Texts
WRITELN

Example Write out the variable *nthingys* with an explanation:

```
WRITE('Number of things: ',nthingys)
```

In scientific notation, write out *pi* in a field 12 spaces wide with 7 digits in the mantissa.

```
pi := 4.0*ArcTan(1.0);
WRITE(' pi =', pi:12:-7);
```

WriteByte procedure (Files) L

Description *WriteByte* writes one byte to a file, which is compatible with the type *BYTE*.

Declaration **PROCEDURE** WriteByte(*f*: FILE; *ch*: BYTE);

Usage `WriteByte(f, b);`

WriteByte must be imported from the library module *Files*.

f must be of type *FILE* imported from *Files*.

b must be of a type compatible with *BYTE* imported from *SYSTEM*.

Comments With this routine you can write out a file in small pieces.

WriteByte is appropriate with types that require only 1 byte of storage, namely BOOLEAN, CHAR, subrange types in the range 0 to 255, and enumeration types with at most 256 elements.

Example Write out a boolean value *okay* to a file *f*:

```
WriteByte(f, okay)
```

WriteBytes procedure (Files) L

Description	<i>WriteBytes</i> writes a specified number of bytes to a file.
Declaration	PROCEDURE WriteBytes(<i>f</i> : FILE; <i>buf</i> : ADDRESS; <i>nbytes</i> : CARDINAL);
Usage	<code>WriteBytes(<i>f</i>, <i>bufStart</i>, <i>nbytes</i>);</code>
	<i>WriteBytes</i> must be imported from the library module <i>Files</i> .
	<i>f</i> must be of type <i>FILE</i> imported from <i>Files</i> .
	<i>bufStart</i> must be of type <i>ADDRESS</i> imported from <i>SYSTEM</i> .
	<i>nbytes</i> must be of type <i>CARDINAL</i> .
Comments	<i>nbytes</i> from the file <i>f</i> will be written from the block of memory starting from the address <i>bufStart</i> and extending to <i>bufStart</i> + <i>nbytes</i> - 1. This is a low-level routine that should only be used for special purposes. When writing out variables, we recommend using the routines <i>WriteByte</i> , <i>WriteWord</i> , and <i>WriteRec</i> .

Example	Take a snapshot of the entire memory (except for the last byte):
	<code>WriteBytes(snap, 0, 65355);</code>

(**WriteCard** procedure (Texts) L

Description	<i>WriteCard</i> writes out a cardinal number to a <i>TEXT</i> file.
Declaration	PROCEDURE WriteCard(<i>t</i> : TEXT; <i>c</i> , <i>n</i> : CARDINAL);

Usage	<code>WriteCard(t, c, width);</code>
	<i>WriteCard</i> must be imported from the library module <i>Texts</i> .
	<i>t</i> must be of type <i>TEXT</i> imported from <i>Texts</i> .
	<i>c</i> must be of type <i>CARDINAL</i> .
	<i>width</i> must be of type <i>CARDINAL</i> .
Comments	The number is output right-justified and blank-padded in a field <i>width</i> characters wide.
See Also	<i>CARDINAL</i> <i>ReadCard</i> <i>Texts</i> <i>WRITE</i>
Example	Write the number <i>n</i> to the standard output: <code>WriteCard(output, n, 6)</code>

WriteChar procedure (Terminal) L

Description	<i>WriteChar</i> writes a character to the console.
Declaration	PROCEDURE <code>WriteChar(ch: CHAR);</code>
Usage	<code>WriteChar(ch);</code>
	<i>WriteChar</i> must be imported from the library module <i>Terminal</i> .
	<i>ch</i> must be of type <i>CHAR</i> .
Comments	<i>WriteString</i> does the same thing for a string.
Example	Write out a character X to the screen: <code>WriteChar("X")</code>

WriteChar procedure (Texts) L

Description *WriteChar* writes a character to a *TEXT* file.

Declaration **PROCEDURE** WriteChar(*t*: TEXT; *ch*: CHAR);

Usage *WriteChar(t, ch);*

WriteChar must be imported from the library module *Texts*.

t must be of type *TEXT* imported from *Texts*.

ch must be of type *CHAR*.

Comments In general, the *WRITE* statement is easier to use than *WriteChar*.

WriteString will do the same thing for a whole string.

See Also *ReadChar*

Texts

WRITE

WriteString

Example Write out character *MrMouse* to text *comix*:

```
MODULE WriteCharacter;
FROM Texts IMPORT TEXT, WriteChar;
VAR
  MrMouse: CHAR;
  comix: TEXT;
BEGIN
  (* other code *)
  WriteChar(comix, MrMouse)
END WriteCharacter.
```

WriteDouble procedure (Doubles) L

Description *WriteDouble* writes a double-precision real variable to a *TEXT* file.

Declaration **PROCEDURE** WriteDouble(*t*: TEXT; *d*: LONGREAL;
n: CARDINAL; *m*: INTEGER);

Usage WriteDouble(*t*, *d*, *width*, *digits*);

WriteDouble must be imported from the library module *Doubles*.

t must be of type *TEXT* imported from *Texts*.

d must be of type *LONGREAL*.

width must be of type *CARDINAL*.

digits must be of type *INTEGER*.

Comments The number *d* is output right-justified in a field *width* characters wide with *digits* digits after the decimal point.

If *digits* is negative, then the number is output in scientific notation with *digits* digits after the decimal point.

If *digits* is zero, the decimal point is omitted.

The *WRITE* and *WRITELN* procedures use this procedure to write *LONGREALS* with a default for *width* of 21 and a default for *digits* of -14.

See Also Doubles
ReadDouble

Example Write a double-precision *d* to a text *t*:

```
WriteDouble(t, 2.99, 5, 2); yields " 2.99"  
WriteDouble(t, 2.99, 5, 0); yields "    3"  
WriteDouble(t, 2.99, 9, -2); yields " 2.99D+01"
```

WriteHex procedure (InOut) L

Description WriteHex writes a hexadecimal number out to output.

Declaration PROCEDURE WriteHex(x,n: CARDINAL);

Usage WriteHex(*num*, *digits*);

WriteHex must be imported from *InOut*.

num must be of type CARDINAL.

digits must be of type CARDINAL.

Comments The number will be converted to hexadecimal and printed out in a field that is at least *digits* wide. If the number doesn't need that many digits, it will be blank-padded.

This and *WriteOct* are useful procedures to print out addresses and storage requirements in their natural form.

See Also InOut

 WriteOct

(
Example Write in hex the number 32 in a field four digits wide:

```
WriteHex(32, 4) yields " 20"
```

WriteInt procedure (Texts) L

Description	WriteInt writes out an integer number to a TEXT file.
Declaration	PROCEDURE WriteInt(<i>t</i> : TEXT; <i>i</i> : INTEGER; <i>n</i> : CARDINAL)
Usage	WriteInt(<i>t</i> , <i>i</i> , <i>width</i>);
	<i>WriteInt</i> must be imported from the library module <i>Texts</i> .
	<i>t</i> must be of type <i>TEXT</i> imported from <i>Texts</i> .
	<i>i</i> must be of type INTEGER.
	<i>width</i> must be of type CARDINAL.
Comments	The number is output right-justified and blank-padded in field <i>width</i> characters wide.
See Also	INTEGER ReadInt Texts WRITE
Example	Write the number <i>n</i> to the standard output: WriteInt(output,n,6)

WRITELN standard procedure S/E

Description	<i>WRITELN</i> is a generalized write statement that starts a new line when it is done.
-------------	---

Usage `WRITELN(elist);`

or

`WRITELN(t, elist);`

t must be of type *TEXT* imported from *Texts*.

elist is simply a list of expressions and optional formatting parameters.

Comments If *t* is not specified, the standard output stream *output* will be used.

WRITELN will be translated by the compiler into the appropriate calls to the module *Texts*. Thus, a change in the definition module of *Texts* could lead to *WRITELN* not working.

WRITELN is the same as *WRITE*, except *WRITELN* starts a new line when the output is complete.

Although *WRITE*, *WRITELN*, *READ*, and *READLN* are not included in the Modula-2 definition, programming without them can be very tedious; however, they are not portable.

A way around this is to develop code using these statements for debugging and development output, and use the explicit calls to *Texts* for the actual programmed output.

See Also `READ`
 `READLN`
 `Texts`
 `WRITE`

Example Write out the variable *nthingys* with an explanation, then start a new line:

```
WRITELN('Number of things: ',nthingys)
```

In scientific notation, write out *e* in a field 12 spaces wide with 7 digits in the mantissa, and start a new line:

```
e := Exp(1.0);  
WRITELN(' e =', e:12:-7);
```

WriteLn procedure (Terminal) L

Description *WriteLn* starts a new line on the screen.

Declaration PROCEDURE WriteLn;

Usage WriteLn;

WriteLn must be imported from the library module *Terminal*.

Comments The *Terminal* routines provide input and output to the screen without the overhead of the stream abstraction in *Texts*.

See Also Terminal
 WriteChar
 WriteString

Example Write out an error message:

```
WriteString('Help !'); WriteLn;
```

WriteLn procedure (Texts) L

Description *WriteLn* writes an *EOL* to a *TEXT* file.

Declaration **PROCEDURE** WriteLn(*t*: TEXT);

Usage *WriteLn(t);*

WriteLn must be imported from the library module *Texts*.

(*t* must be of type *TEXT* imported from *Texts*.

Comments The *WRITELN* statement is easier to use than this procedure.

WriteLn(t) is equivalent to *WriteChar(t,EOL)*.

See Also *Texts*
WRITELN

Example Write an *EOL* to the text *console*:

WriteLn(console);

WriteLn procedure (InOut) L

Description Starts a new line on the current output device.

Declaration **PROCEDURE** WriteLn;

Usage *WriteLn;*

WriteLn must be imported from the library module *InOut*.

Comments *WriteLn* provides a standard new line routine.

This procedure is defined in terms of the *Texts* module.

See Also Terminal
Texts
WRITELN

Example Write out an error message:

```
MODULE WriteMessage;  
FROM InOut IMPORT WriteString, WriteLn;  
BEGIN  
  WriteString('Help !'); WriteLn;  
END WriteMessage.
```

WriteLong procedure (Texts) L

Description *WriteLong* writes out a long integer number to a *TEXT* file.

Declaration PROCEDURE WriteLong(t: TEXT; l: LONGINT;
n: CARDINAL);

Usage WriteLong(t, l, width);

WriteLong must be imported from the library module *Texts*.

t must be of type *TEXT* imported from *Texts*.

l must be of type *LONGINT*.

width must be of type *CARDINAL*.

Comments In general, the *WRITE* statement is easier to use than *WriteLong*.

The number is output right-justified and blank-padded in field *width* characters wide.

See Also **LONGINT**
 ReadLong
 Texts
 WRITE

Example Write the number *n* to the standard output:

```
WriteLong(output, n, 6)
```

WriteOct procedure (InOut) L

Description *WriteOut* writes an octal number out to *output*.

Declaration **PROCEDURE** WriteOct(*x,n*: CARDINAL);

Usage WriteOct(*num, digits*);

WriteOct must be imported from *InOut*.

num must be of type **CARDINAL**,

digits must be of type **CARDINAL**.

Comments The number will be converted to octal and printed out in a field at least *digits* wide. If the number doesn't need that many digits, it will be blank-padded.

This and *WriteHex* are useful procedures to print out addresses and storage requirements in their natural form.

See Also **InOut**
 WriteHex

Example Write out the number 8 in octal in a field four digits wide:

```
WriteOct(8, 4) yields '10'
```

WriteReal procedure (*Texts*) L

Description	WriteReal writes out a real number to a TEXT file.
Declaration	PROCEDURE WriteReal(<i>t</i> : TEXT; <i>r</i> : REAL; <i>n</i> : CARDINAL; <i>digits</i> : INTEGER);
Usage	WriteReal (<i>t</i> , <i>r</i> , <i>width</i> , <i>digits</i>); <i>WriteReal</i> must be imported from the library module <i>Texts</i> . <i>t</i> must be of type <i>TEXT</i> imported from <i>Texts</i> . <i>r</i> must be of type <i>REAL</i> . <i>width</i> must be of type <i>CARDINAL</i> . <i>digits</i> must be of type <i>INTEGER</i> .
Comments	In general, the <i>WRITE</i> statement is easier to use than <i>WriteReal</i> . The number is output in a field <i>width</i> characters wide, and the mantissa is <i>digits</i> long.
See Also	REAL ReadReal Texts WRITE
Example	Write the number <i>pi</i> to the <i>TEXT</i> file <i>doc</i> : WriteReal(doc, pi, 5, 2) yields " 3.14" WriteReal(doc, pi, 12, 2) yields "3.14" WriteReal(doc, pi, 5, 0) yields " 3" WriteReal(doc, pi, 10, -3) yields " 3.141E+00"

WriteRec procedure (Files) L

Description WriteRec writes a record or an array to a disk file declared as FILE.

Declaration **PROCEDURE** WriteRec(*f*: FILE; **VAR** *rec*: **ARRAY OF WORD**);

Usage WriteRec(*f*, *r*);

WriteRec must be imported from the library module *Files*.

f must be of type *FILE* imported from *Files*.

r is any type.

Comments This routine should be used for writing any structure larger than a byte or a word (see *WriteByte* and *WriteWord*).

See Also Files,
ReadRec,
WriteByte
WriteWord

Example Write the record *widget[i]* to *newFile*:

WriteRec(*newFile*, *widget[i]*)

(WriteString procedure (Terminal) L

Description *WriteString* writes a string to the terminal at the current cursor position.

Declaration **PROCEDURE** WriteString(*s*: **ARRAY OF CHAR**);

Usage	<code>WriteString(s);</code>
	<i>WriteString</i> must be imported from the library module <i>Terminal</i> .
	<i>s</i> must be of type ARRAY OF CHAR .
Comments	There is no stream (like console, input, or output) associated with <i>Terminal</i> procedures; thus, output goes directly to the console and redirection is not possible.
See Also	<code>ReadString</code> <code>Terminal</code>
Example	Write a string <i>prompt</i> to the terminal: <code>WriteString(prompt)</code>

WriteString procedure (Texts) L

Description	<i>WriteString</i> writes out a string to a file declared as <i>TEXT</i> .
Declaration	PROCEDURE <code>WriteString (t: TEXT; s: ARRAY OF CHAR);</code>
Usage	<code>WriteString(t, s);</code>
	<i>WriteString</i> must be imported from the library module <i>Texts</i> .
	<i>t</i> must be of type <i>TEXT</i> imported from <i>Texts</i> .
	<i>s</i> must be a character array.
Comments	In general, the <i>WRITE</i> statement offers easier output than <i>WriteString</i> . All of the characters in the string are output up until the first zero byte or the end of the string is encountered.

See Also	ReadString Texts WRITE
Example	Write the string <i>name</i> to the <i>TEXT</i> file <i>doc</i> : <code>WriteString(doc, name)</code>
<hr/>	
WriteWord procedure (Files) L	
Description	<i>WriteWord</i> writes a variable compatible with <i>WORD</i> to a disk file declared as <i>FILE</i> .
Declaration	PROCEDURE WriteWord(<i>f</i> : FILE; <i>w</i> : WORD);
Usage	<code>WriteWord(f, w);</code> <i>WriteWord</i> must be imported from the library module <i>Files</i> . <i>f</i> must be of type <i>FILE</i> imported from <i>Files</i> . <i>w</i> must be of a type compatible with <i>WORD</i> , which is imported from the pseudomodule <i>SYSTEM</i> .
Comments	Types compatible with <i>WORD</i> are those with storage of no more than one word (2 bytes); these include INTEGER, CARDINAL, BITSET, all pointers (<i>ADDRESS</i> is a pointer).
See Also	Files ReadWord
Example	Write the cardinal <i>nwidgets</i> to <i>newFile</i> : <code>WriteWord(newFile, nwidgets)</code>

APPENDIX A

Turbo Modula-2 and Turbo Pascal

This appendix examines the differences between Turbo Pascal and Turbo Modula-2 in a detailed manner. First, we will point out some of the features found only in Turbo Modula-2 and then we will compare the differences in the features of both languages.

In general, Turbo Modula-2 and Turbo Pascal are quite similar. Some features found in Turbo Modula-2 will be more familiar to Turbo Pascal programmers than they will be to regular Pascal programmers. This is because Turbo Pascal provides a rich set of standard procedures and useful extensions that are defined directly in the Turbo Modula-2 language.

Modula-2's separate compilation and library facilities provide a full and extensible set of primitives. The differences between the language implementations range from simple items like case sensitivity to complex issues involving type- and version-checking across compilation units.

Modula-2 can be used for large program development and for expressing operating system concepts like concurrency and interrupt handling. Rather than using Include files, Modula-2 separates programs textually into modules. Like Pascal procedures, modules can be used to control the scope of identifiers. Unlike procedures, the walls around a module are opaque in both directions, with only explicitly defined changes in scope. This control of scope occurs in both local (nested) and library (separately compiled) modules. In the case of a library identifier, any identifier needed from another module must be requested. And both local and library modules must explicitly make the identifier visible.

Library modules can be separately compiled. Though some other languages allow separate compilation, Modula-2 does full type- and version-checking. This helps prevent errors that occur because the program is textually broken into pieces. Thus you have the same safety as Turbo Pascal's Include files without having to recompile support procedures all the time.

Two new low-level features found in Turbo Modula-2 are coroutines and inter-

rupt handlers. As described in Chapter 8, coroutines simulate concurrent processes by using their own data space and sharing the processor.

While some people have expressed interrupt handlers in Turbo Pascal using inline code, Modula-2 allows interrupt handlers to be defined in a high-level manner. This is done by initializing the machine and then having the interrupt handler install itself at some predetermined vector. The rest is handled by Turbo Modula-2.

What's the Difference?

The following overview compares the elements of Turbo Modula-2 to Turbo Pascal and should quickly acquaint the Turbo Pascal programmer with Turbo Modula-2. (For a complete explanation of any feature, refer to the body of this manual.)

To start, let's consider the following Pascal and Modula-2 programs:

```
PROGRAM prime( OUTPUT );
CONST
  size = 8190;
VAR
  i,k,prime,count : INTEGER;
  flags : ARRAY [0..size] OF
BOOLEAN;
BEGIN
  count := 0;
  FOR i := 0 TO size DO
    flags[i] := TRUE;
  FOR i := 0 TO size DO BEGIN
    IF flags[i] THEN BEGIN
      prime := i + i + 3;
      k := i + prime;
      WHILE k <= size DO BEGIN
        flags[k] := FALSE;
        k := k + prime;
      END ;
      count := count + 1;
    END ;
  END;
```

```
MODULE prime;
CONST
  size = 8190;
VAR
  i,k,prime,count : INTEGER;
  flags : ARRAY [0..size] OF
BOOLEAN;
BEGIN
  count := 0;
  FOR i := 0 TO size DO
    flags[i] := TRUE;
  FOR i := 0 TO size DO
    IF flags[i] THEN
      prime := i + i + 3;
      k := i + prime;
      WHILE k <= size DO
        flags[k] := FALSE;
        k := k + prime;
      END ;
      count := count + 1;
```

```
END;  
END;  
WRITELN(count,' Primes');  
END (* prime *).  
END;  
END;  
WRITELN(count,' Primes');  
END prime.
```

Vocabulary

There are minor differences in the vocabularies of the two language implementations. For the most part, it is simple to convert a Turbo Pascal program into Turbo Modula-2.

Identifier Names

Turbo Modula-2's identifiers are written in the same manner as Pascal's. There are two major differences: (1) Turbo Modula-2 does not allow underscores in identifier names, and (2) it is case-sensitive. The latter has been made optional with a compiler switch. The following are some legal and illegal modula identifiers:

Legal	Illegal
AnotherIndent	An__Indent
n	
N	

Characters

Character constants may either be denoted by the character enclosed in single or double quotes or by the character's ordinal number written in octal notation and followed by the letter *C*. For example:

32C	A character constant representing a Control-Z
101C	The letter 'A'
'B'	The letter 'B'
'C'	The letter 'C'

Numbers

The numbers of Turbo Modula-2 are a superset of those allowed in Pascal. In ad-

dition to the integers and reals in Turbo Pascal, Turbo Modula-2 has an unsigned integer, called a CARDINAL, and two double-precision types, an integer (LONGINT) and a real (LONGREAL).

Like Turbo Pascal, Turbo Modula-2 allows you to specify the base and type of a constant number. In Turbo Pascal, a hexadecimal number is written starting with a dollar sign, such as \$0F; but in Turbo Modula-2 hex numbers are written starting with a decimal digit and ending with the letter *H*, such as 0FH. The various classes are Octal, Hexadecimal, Character, Real, and Long. The following are examples of numeric constants in Turbo Modula-2:

23B	Octal cardinal	(19 decimal)
023H	Hexadecimal cardinal	(35 decimal)

Real constants are written the same as they are in Turbo Pascal--a decimal point must be present. The scale factor is preceded by *E*. For example:

1.1414	A single-precision real approximating the square root of 2
3.02E9	A single-precision real in scientific notation

Double-precision constants are just like real constants with a scale factor, but with a *D* in place of the *E*. For example:

0.0D0	A double-precision real that is a very precise zero
-------	---

In the same way a *C* following a octal number makes a CHAR constant, an *L* following a integer makes a LONGINT constant. Of course, the integer can be larger than normal integers if followed by an *L*, as shown in the following:

1483236283L	A long integer over a trillion (1,483,236,283).
-1L	Internally, this long integer is 4 bytes of ones.

Strings

The two differences between strings in Turbo Pascal and Turbo Modula-2 are (1) strings can be enclosed in either single or double quotes, and (2) the quote used to enclose the string may not appear in the string.

In Pascal, to have a quote within a string you would write the quote twice. In

Turbo Modula-2, you may include whichever quote you are not using to enclose the string. Thus, the following strings yield the same result:

Modula-2 Pascal

“ ”	‘ ’
“ ”	‘ ’

Here are other examples of legal strings in Modula-2

(" ` bba " ' The language " Modula-2 " ' " Peter ' s programs ")

Set Constants

Sets in Turbo Modula-2 are substantially different than those in Turbo Pascal. In Turbo Modula-2, sets are represented by one machine word and thus are very efficient. Tests for set membership are much faster. However, there is one drawback: A machine word has only enough bits to represent a set of 16 elements. Pascal programmers who are used to using SET OF CHAR may miss it; however, there are other ways to obtain the same results.

The syntactic differences are that sets are delimited by curly brackets instead of the square brackets used in Pascal. A set must be preceded by its type identifier when used in expressions, otherwise, it is assumed to be the predefined set type BITSET. For example:

```
TYPE
  UserSet = SET OF (red, green, blue);
CONST
  hibit = [7];
  CurrentColor = UserSet{green};
```

Comments

In Turbo Pascal, comments can only be nested one level deep by using one of the two forms within the other. In Turbo Modula-2, the only form of comment uses the delimiters (* and *). (As previously shown, the curly braces, { and }, are used to denote sets.) In contrast to Turbo Pascal, Turbo Modula-2 comments may be nested to any depth. This is especially useful for debugging programs. For example:

```
c = 2.997925E8; (* the speed of light in meters per second *)
(*D WriteReal(c); (* printout for debugging *) D*)
```

Declarations

In contrast to standard Pascal, Turbo Modula-2's order of declarations (like Turbo Pascal's) is not fixed. Constant, type, variable, and procedure declarations can be written in any order. This gives the programmer more freedom to group related items together. Of course, every declaration section must be preceded by the appropriate CONST, TYPE, VAR, or PROCEDURE symbol.

Constant Declarations

There are two major differences between Turbo Pascal's constants and Turbo Modula-2's: (1) Turbo Modula-2 allows constant expressions in declarations. (2) There is no equivalent for Turbo Pascal's typed constants (which are actually pre-initialized variables).

Turbo Modula-2's constant expressions may be used anywhere Pascal allows only constants. Constant expressions consist of constants connected by the usual operators. Some examples of declarations follow (constant expressions are seen in lines 2, 5 and 8):

CONST

```
PiByTwo = 3.141592 / 2.0;
version = "1.6 last changes: Oct 84";
mask = [0..3,8];
Truth = NOT FALSE;
Size = 1000;
```

TYPE

```
a = ARRAY [0..Size-1] OF CHAR;
```

Turbo Pascal's typed constants can be simulated with normal Turbo Modula-2 variables and the initialization part of modules. The only difference is that the values are initialized at runtime instead of at load-time.

Type Declarations

Turbo Modula-2 has all of the types offered by Turbo Pascal, whether defined directly as standard identifiers or as library types. In addition, Turbo Modula-2 has procedure types. In defining types in Turbo Modula-2, a constant expression may be used anywhere a constant is used to define a type in Pascal.

Standard types in Turbo Modula-2 are CHAR, BOOLEAN, BITSET, INTEGER, CARDINAL, REAL, LONGINT, LONGREAL, and PROC. Of these, INTEGER, BOOLEAN, CHAR, and REAL are used exactly as they are in Turbo Pascal. We have already mentioned the addition numeric type, CARDINAL, LONGINT, and LONGREAL; the remaining types, BITSET and PROC, will be explained shortly. The user-defined types of both languages are similar, but have slight differences.

Subrange and enumeration types are the same for both languages except Turbo Modula-2 encloses in square brackets ([]) the values defining a subrange. For example:

```
TYPE
  bitnumber = [0..wordlength-1];
  smallint = [0..255];
```

The presence of the types INTEGER and CARDINAL causes a slight ambiguity when defining a subrange. In the previous example, it's not known if the base type of the subrange is INTEGER or CARDINAL. This problem is resolved by the convention that the base type is assumed to be CARDINAL if the lower bound is not negative; otherwise, it is INTEGER. You may override this convention by explicitly specifying the base type as follows:

```
TYPE
  ismall = INTEGER[0..255];
```

Arrays

As with subranges, the syntax of an array is slightly different in each language. When the bounds of an array are specified by a subrange, brackets are not needed; for example:

```
TYPE
  color = (brown,purple,orange);
  c = ARRAY color OF CHAR;
  a = ARRAY [0..9],[0..9] OF REAL;
  charkind = ARRAY OF CHAR (letter,digit,special,illegal);
```

Records

In Turbo Modula-2, records without variant parts are identical to those of Turbo Pascal. Turbo Modula-2 variant parts have a somewhat different syntax. If the name of the tag field is omitted, the colon and the type must still be written. The vertical bar serves to separate cases. Like Turbo Pascal, label ranges (for example, 0..5) and an optional ELSE part may be used. More than one variant part is allowed and it need not be written at the end of the record, as in Turbo Pascal. Variant parts in Turbo Modula-2 must have an END statement; for example:

```
complex = RECORD x, y : REAL END ;

sneaky = RECORD
  CASE : BOOLEAN OF
    FALSE : c : CARDINAL           |
    TRUE  : p : POINTER TO CARDINAL   |
  END
END ;

demo = RECORD
  a, b : CARDINAL;
  x : REAL;
  CASE t1 : CARDINAL OF
    0..3,7  : f1,f2 : File          |
    4,6      : name   : ARRAY [0..7] OF CHAR   |
  ELSE
    link    : POINTER TO sneaky
  END ;
  d, e : BOOLEAN ;
  CASE c : color OF
```

```
    red, blue :      |
    green : g : date |
END ;
last : BOOLEAN ;
END
```

Procedure Types

Procedure types are new to Turbo Pascal programmers. They can be used to define the interface of a procedure that is passed to some other procedure as a parameter. This is useful for allowing one data access routine to perform many different functions on the data.

For example, you may pass a tree-traversal procedure a procedure parameter that prints the node, uses the node for a calculation, or performs some other function on the nodes of the tree. The point is that each of these operations have the same interface, which can be defined globally and exported.

The type PROC is a predefined procedure type with no parameters. Variables of this type may receive assignments from procedures declared as

```
PROCEDURE Foo;
BEGIN
  (* Statements *)
END Foo;
```

Procedure variables are defined with the reserved word PROCEDURE, followed by a formal type list. In contrast to normal parameter lists, the names of the parameters are not given. Types are separated by commas and may be preceded by a VAR to indicate variable parameters. Function procedures declare a result type (see the section, "Function Procedures").

Variables of a procedure type may assume as their values procedures whose formal parameter list is compatible with the formal type list of the procedure type. However, procedures local to another procedure and standard procedures may not be assigned to procedure variables. Note that arithmetic and file-handling procedures (for example, *Sin* and *Open*) are not standard procedures in Turbo Modula-2. They are library procedures and can be assigned to procedure variables of the correct type. For example:

```
MODULE ProcedureVars;
FROM Texts IMPORT TEXT;
TYPE
  RealFunc = PROCEDURE (REAL): REAL;
  TextDriver = PROCEDURE(VAR Text, CHAR);
VAR
  MyExp: RealFunc;
  MyWriteChar: TextDriver;
  MyClearScreen: PROC;    (* Predefined parameterless procedure *)
PROCEDURE MyExponentiation(r:REAL): REAL;
BEGIN
  (* Statements *)
END MyExponentiation.

PROCEDURE MyWriteCharacter(VAR t: TEXT; ch: CHAR);
BEGIN
  (* Statements *)
END MyWriteCharacter;

PROCEDURE MyClearTheScreen;
BEGIN
  (* Statements *)
END MyClearTheScreen;

PROCEDURE ExecuteP(p:PROC);
BEGIN p;
END ExecuteP;

BEGIN (* Assign procedures to the procedure variables *)
  MyExp := MyExponentiation;
  MyWriteChar := MyWriteCharacter;
  MyClearScreen := MyClearTheScreen;
  (* Execute procedure variables *)
  MyClearScreen;
  WRITE(MyExp(29.0));
  MyWriteChar(output "A");
  ExecuteP(MyClearScreen);
END ProcedureVars.
```

Variable Declarations

Turbo Modula-2's variable declarations are identical to those of Pascal: A list of variable identifiers are given (separated by commas), along with a colon and the variable's type. For example:

```
i, j, k    : INTEGER;  
printer    : textwriter;
```

Like Turbo Pascal, Turbo Modula-2 offers a facility to specify the address of a variable. This must be considered a low-level facility and must be used with care. The address is specified in brackets after the variable identifier. Unlike Turbo Pascal, Turbo Modula-2 restricts absolute variables from assuming dynamic values. Thus when a Turbo Pascal program uses a local variable "absolutized" with a parameter to the procedure, the Turbo Modula-2 equivalent would simply be pointer assignments that may require type coercion. Variables in Turbo Modula-2 are made absolute as shown here.

```
maskregister [OFFCDH]: BITSET;
```

Procedure Declarations

Procedures in both languages are declared in much the same way. Unlike Turbo Pascal but available in other Pascals, Turbo Modula-2's procedures and functions may be passed as parameters. The parameter must be declared as an already defined procedure type (shown in the previous section, "Variable Declarations").

Open Array Parameters

Open array parameters allow arrays declared of different length to be passed to the same procedure. In Turbo Pascal, this can be done only with string parameters, and only when the *V* compiler option is turned on. In Turbo Modula-2, there is a provision for open arrays of any type and there is a mechanism to dynamically obtain the upper bound of open array parameters.

Within the procedure, the lowest array element always has an index of zero. The index of the highest array element can be obtained as *HIGH(a)*, *a* being specified as **ARRAY OF <some type>**. For example:

```
PROCEDURE writevector(v : ARRAY OF REAL);
VAR
  i : CARDINAL;
BEGIN
  FOR i := 0 TO HIGH(v) DO WRITE(v[i]) END
END writevector;
```

Untyped Parameters

Turbo Pascal allows you to skip type-checking with untyped parameters by simply leaving off the type specification. This is useful for writing generic procedures. Of course, Turbo Modula-2 also allows this with an explicit declaration that uses the *SYSTEM* type WORD. Turbo Modula-2 goes one step further by allowing you to dynamically determine the size of the object passed with the standard procedures *SIZE* and *HIGH*. This is shown in the following example:

```
MODULE Untyped;

PROCEDURE foo(object: ARRAY OF WORD);
BEGIN
  WRITELN(SIZE(object), HIGH(object));
END foo;

VAR
  c: CARDINAL;
  r: REAL;
BEGIN
  foo(c);
  foo(r);
END Untyped.
```

The first call to the procedure *foo* results in the output of 2 and 0. Thus, the procedure knows the object passed is 2 bytes and the highest index into the array of words is 0. During the second call, this procedure outputs a 4 and a 1, indicating the object is 4 bytes and occupies positions 0 and 1 of the word array.

Function Procedures

In Turbo Modula-2, **FUNCTION** is no longer a reserved word; function declarations differ from procedure declarations only by the indication of a result type.

The following is an example of a function declaration that finds the length of a string:

```
PROCEDURE len(s: ARRAY OF CHAR): CARDINAL;
(* return length of string s *)
VAR
  i : CARDINAL;
BEGIN  i := 0;
  WHILE (i <= HIGH(s)) & (s[i] # 0C) DO i := i+1 END ;
  RETURN i
(  ID len;
```

Expressions

Turbo Modula-2 expressions are very similar to those of Turbo Pascal. The usual operators, +, -, *, /, DIV, and MOD, are available for operands of type INTEGER, CARDINAL, and REAL, (DIV and MOD apply to INTEGER and CARDINAL, / applies to REAL). Of course, those operators applicable to INTEGER also work on LONGINT, and the same is true for REAL and LONGREAL.

There are, however, no implicit conversions. The following example would therefore be legal in Turbo Pascal but illegal in Turbo Modula-2, because REAL and CARDINAL are not compatible:

```
x := 1.0 + 1;
```

The logical operators AND (also written `&`), OR, and NOT (also written `~`), are available. If the first operand of an AND evaluates to FALSE, the second is not evaluated. Similarly, the second operand of an OR is not evaluated if the first one is TRUE. This rule sometimes shortens programs by eliminating a Boolean flag and a GOTO statement. For example:

```
WHILE (I#0) AND (s[I]>0) DO
  (* Something *)
END ;
```

would have to be translated to the following Pascal statement if *s* is not defined for *I* equal to 0:

```
while I < > 0 do begin
  if s[I]<=0 then goto endwhile;
  (* Something *)
end ;
endwhile:
```

Set Operators

The major benefit of Turbo Modula-2 sets is the ability to treat bits in a word as elements in a set. This makes bit manipulation very easy and defines it directly in the language instead of as an extension as in Turbo Pascal.

The type BITSET is a special predefined set type. It is declared as

```
TYPE
  BITSET = SET OF [0..wordlength-1];
```

where wordlength is the word length of the computer, which is 16 for Turbo Modula-2.

Since Turbo Modula-2 sets only take one word, the operations defined for sets can be viewed as equivalent to Turbo Pascal bitwise operations. Thus Turbo Modula-2 set operators can be viewed as abstract set operations or low-level bitwise operations as seen in Turbo Pascal equivalents.

Operation	Modula Symbol	Pascal Symbol
Union	+	OR
Difference	-	XOR and then AND
Intersection	*	AND
Symmetric Difference	/	XOR

In terms of set operations available in Turbo Pascal, only / is new. It is called symmetric set difference and is an exclusive OR. The resulting set contains all elements that are in either the first set operand or the second, but not in both. The following is an example using sets:

```
MODULE ExampleSet;
TYPE
  colors = (Red,White,Blue,Orange,Purple,Black,Yellow,Green,Cyan);
  flagColors = SET OF colors;
```

```
CONST
  frenchFlag = flagColors[Red,White];
VAR
  currentColors: flagColors;
BEGIN
  currentColors := flagColors[Red];
  currentColors := currentColors + flagColors[White];
  IF (currentColors=frenchFlag) THEN
    WRITELN('Viva la France !');
  END ;
END ExampleSet;
```

Familiar to Turbo Pascal programmers is Turbo Modula-2's alternate use for the preceding logical operators (AND, OR, XOR); in Turbo Pascal, these are also bitwise operators. As previously shown, Turbo Modula-2 has set operators that double as bitwise operators. The only bitwise operators defined in Turbo Pascal and not in Turbo Modula-2 are the shl and shr operators. However this is not a problem because Turbo Modula-2 translates multiplication and division by 2 into machine-language shifts in the appropriate direction. Thus, we have the following equivalents for the integer variable *I*:

Modula	Pascal
<i>I</i> * 2	<i>I</i> shl 1
<i>I</i> * 2 * 2	<i>I</i> shl 2
<i>I</i> * 256	<i>I</i> shl 8
<i>I</i> DIV 16	<i>I</i> shr 4

Turbo Modula-2 and Turbo Pascal use the same relational operators. The only difference is that Turbo Modula-2 provides an additional inequality operator, the pound sign (#), which has the same effect as < >. Both symbols are allowed in Turbo Modula-2.

Operands in expressions in Turbo Modula-2 are very similar to those found in Pascal. The familiar operations of indexing, field selection, dereferencing, and function invocation are available. One difference is seen in expressions that include parameterless function calls. Turbo Modula-2 requires the empty parameter list to be specified in function calls. This helps distinguish function identifiers from variable and constant identifiers.

The following are examples of expressions in Turbo Modula-2:

<code>c + b*3</code>	Integer or cardinal expression
<code>list IN options</code>	Boolean expression
<code>sum + a[i,k]*a[k,j]</code>	Numeric expression using arrays
<code>(ch >= "A") & (ch <= "Z")</code>	Boolean expression with relations
<code>s * {0..3}</code>	Set expression
<code>Exp(Random())</code>	Real expression with nested function calls

Statements

Like Pascal, the most elementary statement in Turbo Modula-2 is the assignment statement. It is stricter in Turbo Modula-2 than in Turbo Pascal in that no implicit type conversions are made. This means you may not assign an integer expression to a real variable, as is allowed in Turbo Pascal. Of course, Turbo Modula-2 provides a mechanism to do this with explicit type-transfer functions; thus we have the following equalities (where *x* is a real and *i* is an integer):

Modula

```
x := FLOAT(i);
i := TRUNC(x * FLOAT(i));
```

Pascal

```
x := i;
i := trunc(x * i);
```

Turbo Modula-2's structured statements have a more modern syntax that does away with Turbo Pascal's compound statement (`begin ... end`). Where Pascal requires the compound statement, Turbo Modula-2 allows a statement sequence to be terminated by an `END` statement.

Procedure Calls

The procedure call statement remains essentially unaltered in Turbo Modula-2. If a procedure has no parameters, empty parentheses are allowed but not required, as in function procedures.

Looping Statements

Turbo Modula-2's `WHILE` and `REPEAT` statements are essentially the same as Turbo Pascal's.

Turbo Modula-2 has an `EXIT` statement, but it is used to terminate the `LOOP` statement. The `LOOP` statement is the same as a Turbo Pascal's `repeat until`

FALSE or while TRUE do begin end, except that instead of using a GOTO to exit the endless loop, Turbo Modula-2 uses the explicit EXIT statement. EXIT causes control to pass to the statement directly after the END that matches the LOOP statement.

The FOR statement has been slightly changed. A step value may now be given using a BY clause (default is +1 if the BY part is left out). The step value must be a constant expression. The Turbo Pascal downto symbol is no longer used. You now simply specify a step value of -1. For example:

```
( FOR i := 0 TO size DO
  flags[i] := TRUE           (* Notice the explicit END *)
  END ;
  FOR i := n-2 DOWNTO 2 BY -1 DO      (* Pascal's DOWNTO *)
    s[i] := (b[i] -
    a2[i]*s[i+1])/a1[i];          (* Any number of statements *)
    j := j + i                  (* between DO and END *)
  END ;
```

CASE Statements

Turbo Modula-2's CASE statement has a different syntax than Pascal's. The vertical bar | is used to separate cases, thus eliminating the need for begin end in the statement part of the case statement. Like Turbo Pascal, Turbo Modula-2 allows ranges of values for CASE labels (like "A".."Z") and an optional ELSE part. For example:

```
CASE ch OF
  | "A".."Z", "a".."z"           : chartype := letter
  | "0".."9"                     : chartype := digit
  | " ", ",", "#", "&".."~/",
  ":"..">","[".."^","[".."~" : chartype := special
  |
ELSE      WRITELN('illegal
                      character'); chartype := illegal
END ;
```

Note that the first and last vertical bars in this example are optional.

WITH Statements

The WITH statement has remained essentially unchanged. While in Turbo Pascal a list of record variables is allowed after WITH, Turbo Modula-2 makes you nest WITH statements to achieve the same effect. For example:

```
WITH p ^ DO      (* in Pascal, only one WITH would have been used *)
  WITH valu ^ DO
    typ := reel; rval := nxrval;
  END ;
  link := head
END ;
```

RETURN Statements

There are two reserved words that take the place of Turbo Pascal's goto statement: the EXIT statement (already discussed) and the RETURN statement. The RETURN statement serves to terminate procedures the same way as the exit statement in Turbo Pascal. In addition, it is used to return function results instead of an assignment to the Pascal function identifier.

Standard Procedures in Turbo Modula-2

Several standard functions and procedures found in Turbo Pascal are no longer directly available in Turbo Modula-2: Some standard procedures have been added and some have been left to be implemented in library modules. Only the most important and commonly used standard procedures are still within the language.

Mathematical functions like *Sqrt*, *Sin*, *Exp*, and so on, are no longer standard functions. They must now be imported from the module *MathLib* (see Chapter 11). File-handling procedures must also be imported from the appropriate library modules (*Files*, *Texts*, or *InOut*).

The Turbo Pascal functions *pred* and *succ* can be replaced by the procedures *INC* and *DEC*. These procedures accept one argument of any scalar type and one count argument that determines the size of the increment or decrement.

Thus, Turbo Pascal statements that look like the following:

```
i := i + 1;
color := pred(Green);
```

will look like this in Turbo Modula-2:

```
INC(i);  
DEC(color);
```

NEW and *DISPOSE* still exist and are translated by the compiler into calls to procedures *ALLOCATE* and *DEALLOCATE*, which must be imported from the module *STORAGE* but may also be redefined within the user program. In contrast to Pascal, a *FLOAT* function exists in Turbo Modula-2. You may be surprised to find that there is no *ROUND* or *EVEN* function; instead, use *Entier(Value + 0.5)* for negative numbers and *TRUNC(Value + 0.5)* for positive numbers.

Tables A-1 through A-3 show which identifiers remain defined within the language. If you don't find what you're looking for here, try Chapter 12, "Turbo Modula-2 Reference Directory." Table A-1. Standard Modula-2 Functions

ABS(x)

Returns absolute value of *x*; result type is the same as argument type.

CAP(c)

Returns argument and result type are of type **CHAR**, the capital letter corresponding to *c*.

CHR(x)

Returns the character with ordinal number *x*.

FLOAT(x)

Converts value *x* of type **CARDINAL** to type **REAL**.

HIGH(a)

Returns high index bound of array *a*.

MAX(T)

Returns the largest element of the argument type.
T is **CARDINAL**, **INTEGER**, **BOOLEAN**, **CHAR**, **REAL**, **LONGREAL**, **LONGINT**, any enumeration or scalar type.

MIN(T)

Returns the smallest element of type *T*.

ODD(x)

Returns **TRUE** if *x* is odd; otherwise it returns **FALSE**.

ORD(x)

Returns the ordinal value of *x*, where *x* is of type **BOOLEAN**, **CHAR**, **INTEGER**, **CARDINAL**, or every enumeration type.

<i>SIZE(T)</i>	Returns the storage requirements of type <i>T</i> in bytes.
<i>SIZE(x)</i>	Returns the storage requirements of variable <i>x</i> in bytes.
<i>TRUNC(x)</i>	Returns <i>CARDINAL</i> result; <i>x</i> of type <i>REAL</i> truncated to integral part.
<i>VAL(T,x)</i>	Returns the value of type <i>T</i> , which has ordinal number <i>x</i> . <i>T</i> is <i>BOOLEAN</i> , <i>CHAR</i> , <i>INTEGER</i> , <i>CARDINAL</i> or every enumeration type.

Table A-2. Additional Functions Offered by Turbo Modula-2

<i>LONG(x)</i>	Converts its argument to <i>LONGINT</i>
<i>INT(x)</i>	Converts its argument to <i>INTEGER</i>
<i>CARD(x)</i>	Converts its argument to <i>CARDINAL</i>
<i>FLOAT(x)</i>	Converts its argument to <i>REAL</i>
<i>DOUBLE(x)</i>	Converts its argument to <i>LONGREAL</i>

Note: These functions only work on numeric types.

Table A-3. Standard Procedures in Turbo Module-2

<i>DEC(x)</i>	$x := x - 1$
<i>DEC(x,n)</i>	$x := x - n$
<i>EXCL(s,i)</i>	$s := s - [i]$
<i>HALT</i>	Terminates program execution
<i>INC(x)</i>	$x := x + 1$
<i>INC(x,n)</i>	$x := x + n$
<i>INCL(s,i)</i>	$s := s + [i]$

Note that type identifiers may be used like function identifiers denoting so-called type-transfer functions. (This is considered a low-level facility and is discussed in Chapter 8.)

APPENDIX B

Installation Procedures

The installation program INSTM2 has two functions. Its primary function is to allow you to configure the Turbo Modula-2 system to your hardware and to configure the Turbo Modula-2 editor to your own taste. Thus, you may tell the system what type of terminal you have and which disk drives you wish it to search for needed files. You may also change the editing commands to more familiar ones.

Note: The installation must be performed if you want Turbo Modula-2 to work the way it is described in this manual.

The second function of INSTM2 is to allow programs compiled on your setup to be re-installed on different terminals. This allows you to distribute .COM files (executable files) to your customers without knowing what type of terminal they have. All you need to do is include INSTM2 files along with your executable program.

In this chapter we will describe how to use INSTM2 to install Turbo Modula-2, and then discuss programs in general.

Installing M2

By this point we assume you have made a backup copy of the Turbo Modula-2 distribution disk, as well as a working disk. (If not, these procedures are described in Chapter 1, »Getting Started.«)

Place your working disk in the logged drive. It should contain all the installation files and the M2.COM file, listed as follows:

INSTM2.COM
INSTM2.OVR
INSTM2.DTA
M2.COM

Type INSTM2 to start the installation program. After the prompt, press **[S]** and the following menu will appear:

Modula-M2 system installation menu.

Choose installation item from the following:

[S]creen installation	[C]ommand installation
[M]iscellaneous	[Q]uit

Enter S, C, M, or Q:

There are four possible choices at this menu: Screen installation, Command installation, Miscellaneous, and Quit. To select a choice, enter one of the highlighted letters shown in square brackets. Any legal selection except *Q* will bring up a new menu: *Quit* will return you to the operating system (unless it is being run from the M2 shell).

Screen installation allows proper functioning of Turbo Modula-2's WordStar-like editor and the various screen manipulation functions in the module *Terminal*.

Command installation is performed if you wish to alter the WordStar-style editing commands. We recommend that you become familiar with the capabilities of the editor before you change any of the editing keys. Generally, you only need to go through the Command installation if your keyboard is missing the necessary keys or if you want your Turbo Modula-2 editor to be compatible with a non-WordStar-like editor. Miscellaneous installation is used in two instances. First, if your keyboard does not have certain display characters used in standard Modula-2, this command will allow you to substitute an alternate character or sequence of characters for those missing from your keyboard. Second, this command allows you to tell Turbo Modula-2 where it should look for files when it is compiling and linking.

Screen Installation

When you press **[S]** to perform Screen installation, the following menu will appear:

- 1) ADDS 20/25/30
- 2) ADDS 40/60
- 3) ADDS Viewpoint-1A
- 4) ADM-3A
- 5) DEC Rainbow, 8 bit
- 6) Ampex D80
- 7) ANSI
- 8) Morrow MDT-20
- 9) Hazeltine 1500
- 10) Hazeltine Esprit
- 11) Kaypro with hilite
- 12) Kaypro, no hilite
- 13) Lear-Siegler ADM-20
- 14) Lear-Siegler ADM-31
- 15) Liberty
- 16) Osborne 1
- 17) Otronix Attache
- 18) Qume
- 19) RC-855 (ITT)
- 20) Soroc 120/Apple CP/M
- 21) Soroc new models
- 22) SSM-UB3
- 23) Tandberg TDV 2215
- 24) Teleray series 10
- 25) Teletex 3000
- 26) Televideo 912/920/92
- 27) Texas Instruments
- 28) Visual 200
- 29) Wyse WY-100/200/300
- 30) Zenith
- 31) None of the above
- 32) Delete a definition

Which terminal? (Enter no. or ^Q to exit):

Select the terminal type you will be using while running Turbo Modula-2 by entering the appropriate number for your terminal at the prompt. Several things are dependent on the terminal you are using. Of course, the editor needs to know about the terminal, but your programs must also know the capabilities of the terminal they are running on. This information is available from the standard module *Terminal*.

Before installation is actually performed, you are asked the question:

Do you want to modify this definition before installation? (Y/N)?

This allows you to modify one or more of the values being installed (described in the next section). If you do not want to modify the terminal definition, type *N* to complete the installation and be returned to the Main Installation Menu.

Manual Installation

If your terminal is *not* on the menu, however, you must define the required values yourself (refer to your terminal manual for these values). An alternative method exists if you also own a copy of Turbo Pascal and have installed it for your configuration.

Using the Turbo Pascal TINST.DTA File

Many people with unusual terminals have Turbo Pascal because (like Turbo Modula-2) it runs on almost anything with a Z80 in it. These people have already gone through a painstaking installation procedure that required searching obscure manuals for cryptic codes. If you are one of these people, relax; you can use your existing TINST.DTA file simply by renaming it to INSTM2.DTA. (Note that the .DTA file only contains data for screen commands not keyboard commands.)

Use the renamed file as if it were the one on the distribution disk. When the terminal menu comes up you will see your terminal listed with any others in the TINST.DTA file.

Entering Terminal Codes

Enter the number corresponding to »None of the above« and answer the questions one by one as they appear on the screen.

The following section describes each installable command in detail. If your terminal does not support a command, press **[Delete]** at the prompt. If *Delete line*, *Insert line*, or *Erase to end of line* are not installed, they can be emulated in the editor. However, these emulation functions are not available in the module *Terminal*; the module *Terminal* only provides a method to check if the functions are available so that your programs can emulate them if necessary. Commands must be entered by pressing the appropriate keys or by entering the decimal or hexadecimal ASCII value of the command. If a command requires the two characters [‘ESCAPE’] and [‘=’], you may

- First press the [Esc] key, then the [=], and the entry will be echoed with appropriate labels.
- Or you can enter the decimal or hexadecimal values separated by spaces. (Hexadecimal values must be preceded by a dollar sign (\$); for example, 61, \$1B 61, or \$1B \$3B, which are all equivalent.) These two methods can be mixed since once you have entered a specific character (except a dollar sign), the rest of the command must be defined in that mode, and vice versa. If you need to delete an entry, you may enter a hyphen as the first character on the line; the text *Nothing* is echoed. This is only effective on commands that expect more than a single number of a Yes/No answer.

Terminal Properties

The following is an explanation of each question asked by the terminal installation program.

Terminal type:

Enter the name of the terminal you are about to install. When you complete INSTM2, the values will be stored and the terminal name will appear on the initial list of terminals. If you later need to re-install Turbo Modula-2 to this terminal, you can do so by choosing the terminal from the list.

Send an initialization string to the terminal?

If you want to initialize your terminal when Turbo Modula-2 starts (for example, to download commands to programmable function keys), press to answer this question; otherwise, press .

Send a reset string to the terminal?

Define a string to be sent to the terminal when Turbo Modula-2 terminates. The description of the preceding initialization command also applies here.

CURSOR LEAD-IN command:

Cursor Lead-in is a special sequence of characters that tells your terminal that the following characters comprise an address on the screen on which the cursor should be placed. When you define this command, you are asked the following supplemental questions:

CURSOR POSITIONING COMMAND to send between line and column:

Some terminals need a command *between* the two numbers defining the row and column cursor address.

CURSOR POSITIONING COMMAND to send after line and column:

Some terminals need a command *after* the two numbers defining the row and column cursor address.

Column first?

Most terminals require the address in the format: first ROW, then COLUMN. If this is the case on your terminal, press . If your terminal requires COLUMN first, then ROW, press .

OFFSET to add to LINE

Enter the number to add to the LINE (ROW) address.

OFFSET to add to COLUMN

Enter the number to add to the COLUMN address.

Binary address?

Most terminals need the cursor address sent in binary form. If this is true for your terminal, press **[Y]**. If your terminal expects the cursor address as ASCII digits, press **[N]**. You are then asked the supplemental question

2 or 3 ASCII digits?

Enter the number of digits in the cursor address for your terminal.

CLEAR SCREEN command:

Enter the command that will clear the entire contents of your screen, both foreground and background, if applicable.

Does CLEAR SCREEN also HOME cursor?

This is normally the case; if it is not so on your terminal, press **[N]** and define the cursor *HOME* command.

DELETE LINE command:

Enter the command that deletes the entire line at the cursor position.

INSERT LINE command:

Enter the command that inserts a line at the cursor position.

ERASE TO END OF LINE command:

Enter the command that erases the line at the cursor position, starting from cursor position through to the end of the line.

START OF 'LOW VIDEO' command:

If your terminal supports different video intensities, then define the command that initiates *dim video*.

START OF 'NORMAL VIDEO' command:

Define the command that sets the screen to show characters in 'normal video'. Some terminals have one command that toggles the video mode; thus this command may be the same as the last one.

Number of rows (lines) on your screen:

Enter the number of horizontal lines on your screen.

Number of columns on your screen:

Enter the number of vertical column positions on your screen.

Delay after CURSOR ADDRESS (0-255 ms):

Delay after CLEAR, DELETE, and INSERT (0-255 ms):

Delay after ERASE TO END OF LINE and HIGHLIGHT On/Off (0-255 ms):

Enter the delay in milliseconds required after the functions specified. RETURN means 0 (no delay).

Is this definition correct?

If you have made any errors in the definitions, press [N] to be returned to the terminal selection menu. The installation data you have just entered will be included in the installation data file and appear on the terminal selection menu, but installation will not be performed.

When you press Y in response to this question, installation is completed, the installation data is written to M2.COM, and you are returned to the main menu. Installation data is also saved in the installation data file and the new terminal will appear on the terminal selection list when you run INSTM2 in the future.

Installation of Editing Commands

The built-in editor performs a number of commands. Each function may be activated by either of two commands: a primary command and a secondary command. Primary commands are undefined, but can be easily defined to fit your taste or your keyboard. The secondary commands are installed by Borland and comply with the standard set by WordStar. From the installation main menu, press C for Command installation. The first command will appear like the following:

CURSOR MOVEMENTS:

1: Character left Nothing ->

This means that no primary commands have been installed to move the cursor one character left. If you want to install a primary command (in addition to the secondary WordStar-like Control-S, which is not shown here), you may enter the desired command after the -> prompt in either of two ways:

1. Simply press the key you want to use. It can be a function key; for example, a left-arrow key or any other key or sequence of keys that you choose (the maximum is 2).

The installation program responds with a mnemonic of each character it receives. If you have a left-arrow key that transmits an Escape character followed by a lower case *a*, and you press this key in the situation given previously, your screen will look like the following:

CURSOR MOVEMENTS:

1: Character left Nothing -> <ESC> a

2. Instead of pressing the key you want to use, you may enter the ASCII value(s) of the character(s) in the command. The values of multiple characters are entered and separated by spaces. Hexadecimal values are prefixed by a dollar sign (\$1B).

This may be useful to install commands not presently available on your keyboard; for example, if you want to install the values of a new terminal while still using the old one.

In both cases, terminate your input by pressing **[Esc]**. Note that the two methods cannot be mixed within one command. If you have started defining a command sequence by pressing keys, you must define all characters in that command by pressing keys and vice versa.

You may enter a minus (-) sign to remove a command from the list or a *B* to back through the list one item at a time.

The editor accepts a total of 45 commands, all of which may be installed to your specifications. If you make an error in the installation, such as defining the same command for two different purposes, a self-explanatory error message is issued. You must correct the error before terminating the installation. If a primary com-

mand conflicts with one of the WordStar® compatible secondary commands, it will render the secondary command inaccessible.

Table B-1 lists the secondary commands, and allows you to enter any primary commands you have installed.

Note: Items 2 and 28 let you define alternative commands to *Character Left* and *Delete Left Character* commands. Normally, the Back Space key is the alternative to **[Left] [Space]**, while there is no defined alternative to **[Del]**. You may install primary commands to suit your keyboard; for example, to use the Back Space key as an alternative to **[Del]** if the Back Space key is more conveniently located. Of course, the two alternative commands must be unambiguous like all other commands.

Table B-1 Turbo Modula-2 Secondary Commands

Cursor Movements

- 1: Character left
- 2: Alternative
- 3: Character right
- 4: Word left
- 5: Word right
- 6: Line up
- 7: Line down
- 8: Scroll up
- 9: Scroll down
- 10: Page up
- 11: Page down
- 12: To left on line
- 13: To right on line
- 14: To top of page
- 15: To bottom of page
- 16: To top of file
- 17: To end of file
- 18: To beginning of block
- 19: To end of block
- 20: To last cursor position

Ctrl	S
Ctrl	H
Ctrl	D
Ctrl	A
Ctrl	F
Ctrl	E
Ctrl	X
Ctrl	W
Ctrl	Z
Ctrl	R
Ctrl	C
Ctrl	Q
Ctrl	S
Ctrl	Q
Ctrl	D
Ctrl	Q
Ctrl	E
Ctrl	Q
Ctrl	X
Ctrl	Q
Ctrl	R
Ctrl	Q
Ctrl	C
Ctrl	Q
Ctrl	B
Ctrl	Q
Ctrl	K
Ctrl	Q
	P

APPENDIX C

Summary of Compiler Directives

This appendix contains a list of available compiler directives and a brief description of each. For a detailed explanation of compiler directives, refer to Chapter 10, »System Operations.«

There are two ways to set the current compiler options. The first is with the global switches in the Options submenu in the Turbo Modula-2 shell. These global switches have an effect over an entire program text. However, the second method, embedded switches, overrides the first method. Embedded switches appear directly in the program text within comments. In the following list of compiler options, the switch settings on the left are embedded and the ones on the right are global (as they appear in the options menu).

List Source Option

Embedded Global

\$L+	List(ON)	Generates a listing
\$L-	List(OFF)	Generates none

Overflow Check Option (INTEGER and CARDINAL)

Embedded Global

\$O+	Overflow(ON)	Checks for overflows
\$O-	Overflow(OFF)	Does not check for overflow

Test Range Option

Embedded Global

\$T+	Test(ON)	Checks array indices and subrange variables
\$T-	Test(OFF)	Does not check them

The Upper=Lower Case Sensitivity Option**Embedded Global**

\$U+	Upper=lower(ON)	Does not distinguish between uppercase and lowercase
\$U-	Upper=lower(OFF)	Distinguishes between type case

Extension Option**Embedded Global**

\$X+	eXtensions(ON)	Allows Turbo Modula-2 extensions
\$X-	eXtensions(OFF)	Flags extension as nonstandard

Native Code Option**Embedded Global**

N/A	Native(ON)	Generates native code for the Z80
N/A	Native(OFF)	Generates M-code

Note: Embedded native code switches have no effect.

APPENDIX D

Error Diagnosis

This section lists most messages that Turbo Modula-2 issues for execution errors. Usually, an error message is caused by a faulty program; however, some messages can also result from illegal input to an otherwise correctly executing program.

Format of a Runtime Error Message

While a program is executing, the Turbo Modula-2 system checks for a number of runtime error conditions. The various error sources are listed in the next two sections.

An error message can be issued by the runtime system, as well as by some support module or application program. The process is the same in every case: The executing program is stopped and a message of the following form is given:

(Name of Error) in module (Name of Module)
(possibly a further explaining message)
Press "C" for calling chain >

For example, in module *BadStuff* the value of an expression computes to 11 where the maximum allowed is 10. The following message is then given:

BoundsError in module BADSTUFF
0 to 10 is legal range, but 11 was evaluated
Press "C" for calling chain >

In another example, an attempt was made to read data from file IN.DAT when the end of that file had already been reached. This error is not detected by the interpreter, but rather by the module *Files*.

EndError in module FILES
While processing file B:IN.DAT
Press "C" for calling chain >

In Turbo Modula-2, an error constitutes an *exception*. Exceptions can be issued by the Turbo Modula-2 interpreter and by any Modula-2 program unit. They can be trapped in an application program with an exception handler. The raising and handling of an exception is explained further in Chapter 9, »Turbo Modula-2 Extensions.«

Errors Detected by the Interpreter

Errors detected by the interpreter are often caused by faulty program logic, and with some of these errors there is no reasonable way to continue a program. Therefore, only part of the following errors can be caught explicitly (by name) in an exception handler. These exceptions are the ones exported by the pseudomodule **SYSTEM**: **OVERFLOW**, **REALOVERFLOW**, and **OUTOFMEMORY**. The remainder can only be caught with the **ELSE** clause of the **EXCEPTION** statement: **BoundsError**, **DivisionByZero**, **StringTooLong**, **FunctionReturnsNoResult**, **EndOfCoroutine**, **CaseSelectError**, **PointerError**, and **IllegalInstruction**.

BoundsError. The value of an arithmetic expression lies outside of its admissible bounds. This can happen when

- A value is assigned to a variable of some subrange type, some user-defined scalar type, or one of the types **CHAR**, **BOOLEAN**, **CARDINAL**, or **INTEGER**. The value lies outside of the bounds defined by the type used.
- A function procedure returns a value exceeding the bounds defined by the type of the function.
- An array index lies outside of the admissible bounds.

A check for this error can be suppressed by the compiler option (* \$T- *). This leads to somewhat shorter object modules and slightly faster execution times. As a rule, however, the efficiency gains do not make up for the resulting safety loss; therefore, bounds-checking is turned on by default.

Additional messages: (*number*) to (*number*) is legal range, but (*number*) evaluated. Indicates the maximum or minimum value allowed, as well as the computed value. This can occur if a negative **INTEGER** value is assigned to a **CARDINAL**, or if a **CARDINAL** above 32767 (the maximum **INTEGER** value) is

assigned to an INTEGER. The system does not know whether the left-hand side of the assignment is of type INTEGER or CARDINAL; therefore, both possibilities are displayed. For example:

65535 is assigned to an INTEGER

or

-1 is assigned to a CARDINAL

(Note that the same internal bit value denotes 65535 when interpreted as a CARDINAL and -1 when regarded as an INTEGER.

DivisionByZero. An attempt has been made to divide by 0.

Additional message: None

Overflow. The result of some computation involving CARDINALS, INTEGERS, or LONGINTs that became too large or too small to be represented in the computer's memory. If the truncated value does not fit into the range of the destination type, an OVERFLOW can also result from a truncation function such as *INT*, *CARD*, *LONG*, or *TRUNC*. The ranges of the three types in question are CARDINAL, 0 to 65535; INTEGER, -32768 to 32767; and LONGINT, -2147483648 to 2147483647.

For operations involving INTEGER and CARDINAL arithmetic and CARDINAL multiplication, overflow checking can be suppressed by turning the compiler option \$O off.

(*Additional message:* None

RealOverflow. A real value has been computed that exceeds the admissible range for REALS (-10³⁸ to 10³⁸).

Additional message: None

StringTooLong. A string expression has been assigned to an array with elements of type CHAR. The string is too long to be held in its full length in the array. The string is not truncated to fit; instead, StringTooLarge is raised.

Additional message: None

FunctionReturnsNoResult. A function procedure reaches its end before executing a RETURN statement.

Additional message: None

EndOfCoroutine. A coroutine reaches its end before being left by a TRANSFER or IOTRANSFER statement.

Additional message: None

OUTOFMEMORY. The program has run out of main memory.

Additional message: Stack = (number), Heap = (number). Storage is occupied from two sides: Stack and Heap mark the boundaries of the upper and lower occupied part of memory, respectively. The out-of-memory condition arises when Stack and Heap meet.

CaseSelectError. A CASE statement without an ELSE clause has executed where no case alternatives apply.

Additional message: None.

PointerError. Dereferencing of some pointer with the value NIL has been attempted. For example, given the declaration VAR cp: POINTER TO RECORD x,y: REAL END; and the program fragment cp:=NIL; IF cp ^ .x > 0.0 THEN ..., the exception PointerError is raised upon execution of the last statement.

Additional message: None

IllegalInstruction. The interpreter finds a code byte that does not represent a legal command. This is a rare but severe error. It can occur if some parts of code have been overwritten by the executing program; for example, if the user erroneously addresses computations.

Additional message: (1) Absolute PC = (number), and (2) prints the value of the instruction counter.

Errors Detected by Support Modules

This section concentrates on exceptions that result from presenting illegal input or exceeding the computer's limits. All exceptions signaling a programming error are discussed in their respective modules description in Chapter 11, »The Standard Library.«

The modules in the Turbo Modula-2 Library recognize some error conditions. An error issued by a library module can be caught in an exception handler, provided the exception is imported from the library module.

Exceptions Issued by Module Files

UseError. This usually occurs if writing to a write-protected disk has been attempted. This exception is most often caused by a disk swap. You can solve the situation by resetting the disk before trying again.

Additional message: Drive (Drive-Code) is read-only.

DiskFull. Writing to disk is impossible, presumably because the disk or disk directory is full.

Additional message: While processing file, (Name): indicates the name of the file involved in the write operation.

EndError. An attempt has been made to read data after the end of a file has been reached. This usually constitutes a programming error.

Additional message: While processing file, (Name): indicates the name of the file involved in the read operation.

DeviceError. The disk cannot be read correctly. This is equivalent to the dreaded BDOS ERROR/BAD-SECTOR message of CP/M.

Exceptions Issued by Module Loader

LoadError. An error has occurred during the loading of a program. The cause of the error is given in an additional message.

Additional messages: (1) File not found : (*File Name*); (2) Read Error : (*File Name*); (3) Out of memory; (4) Version Conflict : (*File Name*).

The first three messages are self-explanatory. The fourth, version conflict, results when the version of an object module does not coincide with the version required by an importing module. (For a more detailed discussion, refer to the section about *Loader* in Chapter 11.)

The Calling Chain

If a program exhibits a runtime error, the Turbo Modula-2 system offers a powerful diagnosis. Turbo Modula-2 helps to localize the error and gives some clues about the state of the program at the time of the error's occurrence. This information is contained in the calling chain, which displays all active procedures at the point where they were called. The chain is displayed on the screen if you press [C] after the following prompt:

Press "C" for calling chain >

This appears next to the actual error message. If you do not want a calling chain, simply press or any letter except C. This will return you to the main menu.

The calling chain consists of one or several lines of the following form:

(Name of Module) (Name of Procedure) (Offset Number) (Program Counter)

The first line states the module, the procedure, the offset number, and the program counter where the error occurred. Note that only the first six characters of the module and procedure are displayed. The left column specifies the module enclosing the procedure in which the error occurred. The second column shows the procedure where the error occurred. The offset number corresponds to the numbers produced in the compiler listing. It specifies the offset (in bytes) of the error point from the beginning of the enclosing procedure or module, whose name is listed in the second column. The program counter is used in the Options menu to find where the error occurred in the source file.

The second line of the calling chain specifies the point where the procedure in the first line was called. The same specification holds for the lines that follow.

The chain finishes when the main program is reached. At that point, the name of the main program is in both the module and the procedure column. The calling chain also terminates if a procedure representing a coroutine is reached. Since coroutines are not called by any other part of the program, they are considered the same as main programs.

If you trace the calling chain from bottom to top, you will experience the same sequence of situations that your program went through before the point of the runtime error.

A calling chain, caused by some hypothetical program, is shown in the following example:

```
BoundsError in module BADSTUFF
0 to 10 is legal range, but 11 was evaluated
Press "C" for calling chain > C
```

Module	Procedure	offset	PC
BADSTUFF	SNEAKY	17	101
BADSTUFF	SUB	113	123
TEST	Q	55	154
TEST	TEST	20	164

>

We can follow the events by starting at the bottom with the main program *Test*. *Test* executes until it reaches the offset of 20 in the main program, then it calls the procedure *Q*. The call to procedure *Q* occurs when the program counter is equal to 164.

Looking at the next line we see that procedure *Q* exists in the module *Test*. This procedure then executes until it reaches an offset of 55 from its beginning. At this point the program counter is 154, which is pointing at a call to the procedure *Sub*.

The next line shows that we have called the procedure *Sub* in module *BadStuff*. The procedure *Sub* executes until it calls *Sneaky*. When *Sneaky* is called, the program counter is 123.

The top line of the calling chain shows that the error occurred in the procedure *Sneaky* in module *BadStuff*. The offset of the error in *Sneaky* is 17. The program counter where the error occurred is 101. We know from the error message that there was some type of bounds error; so we can look at that line for assignments or indexing problems.

Finding RunTime Errors

To find where this error occurred in the source code without having to guess with offsets into the procedures or look at a listing, you can use the program counter and the Find RunTime Error option found in the Options menu.

At the main menu, press **O** to enter the Options menu. Then press **F** for Find RunTime Error. You should know the program counter and the name of the module in which the error occurred. First you are prompted with the name of the main module that was running. If this is not the module where the error occurred, then backspace over it and enter the correct name. If it is the correct module name, then just press **J**.

Next you are prompted to enter the PC. At this point you should enter the number you obtained from the calling chain in the program counter column. When the compiler has compiled to this point, you will be thrown into the editor with the cursor at the corresponding runtime error position. This utility is helpful for finding errors quickly and reducing program development time.

Compiler Error Messages

Error messages sent by the compiler can be synthesized or fixed. The »fixed« messages are read from a file named ERRMSG.SXT, which may or may not be online during compilation. If it is not, only the error number, not the corresponding message, can be displayed. Table D-1 provides all of the messages contained in ERRMSG.SXT.

Some messages will not be found in this list. For example:

" :=" expected, but "=" found

is not in the list because it is synthesized from parts of the faulty program together with predefined pieces of text. Other messages are contained only partial

ly by ERRMSG.S.TXT; the rest are supplied from the faulty program. Items supplied by the compiler appear in parentheses.

Table D-1 Messages in ERRMSG.S.TXT

Error in Identifier

- 0 (*Identifier*) is undeclared.
- 1 (*Identifier*) is declared twice.
- 2 (*Identifier*) is not field of this record.
- 3 (*Identifier*) is not exported by this module.
- 4 (*Identifier*) already exists outside of module.
- 5 Unresolved FORWARD reference: (*Identifier*).
- 6 Unresolved export: (*Identifier*).
- 7 (*Identifier*) is not readable.
- 8 (*Identifier*) is not printable.
- 9 (*Identifier*) expected.
- 10 (*File name*) not found.
- 11 CODE must be imported from SYSTEM.
- 12 Two different versions of (*Module name*).SYM imported.
- 13 No standard procedure allowed here.
- 14 Procedure must be declared at outer-most level.
- 15 (*Identifier*) is exported twice.

Error in Syntax

- 20 Illegal key word at start of statement.
- 21 Illegal start of statement.
- 22 Identifier, literal, or »(*) expected.
- 23 (Not used).
- 24 Loop counter may not be external or parameter.
- 25 Expression must have constant value.
- 26 No enclosing LOOP for EXIT.
- 27 No RETURN from module allowed.
- 28 String literal spans over several lines.
- 29 Past end of file.
- 30 Badly formed number.
- 31 Illegal symbol.
- 32 End of file expected.
- 33 String literal expected.

34 CODE must follow procedure heading; it is not allowed here.

General Errors

- 41 Set elements may only range from 0 to 15.
- 42 (Not used).
- 43 Modules with unqualified export list may not be **exported**.
- 44 Module does not export in **QUALIFIED** mode.
- 45 Length of actual string does not match.
- 46 Second value must be greater.
- 47 Case label occurs twice.
- 50 String assignments not allowed in standard Modula-2.
- 51 String comparisons not allowed in standard Modula-2.
- 52 Not allowed in standard Modula-2.
- 53 Procedure **ALLOCATE** not found.
- 54 Procedure **DEALLOCATE** not found.
- 55 Illegal definition of procedure **ALLOCATE**.
- 56 Illegal definition of procedure **DEALLOCATE**.

Error in Type

- 60 Function procedure required.
- 61 No function procedure allowed here.
- 62 Actual parameter is byte-packed, but **Formal VAR parameter assigns a word**.
- 63 No such variant exists.
- 64 Element types of actual and formal arrays differ.
- 65 Types of actual and formal parameters differ.
- 66 Sizes of type and argument differ.
- 67 Declarations of procedure variable and procedure differ.
- 68 Two different declarations of same procedure.

Error in Constant

- 70 Integer required, but large cardinal value computed.
- 71 Cardinal required, but negative value computed.
- 72 Value lies outside of subrange bounds.
- 73 Number too large.
- 74 Real number too large.
- 75 Overflow in constant expression.

Error/Compiler Limit Exceeded

- 80 Case label must not be greater than 32767.
- 81 Too many local variables.
- 82 Too many procedures, strings, and exceptions.
- 84 Too many imported modules.
- 85 Insufficient space for import/export: name-table overflow.
- 86 Insufficient space for import/export: type-table overflow.
- 87 Insufficient space for import/export: identifier-table overflow.
- 88 Insufficient space for import/export.
- 90 Expression too complex or too many parameters.
- 91 Boolean expression too long.
- 92 Too many nested function calls.

Error Implementation Restriction

- 95 Byte-sized array elements cannot be substituted for VAR parameters.
- 96 Byte-sized array elements cannot be used as arguments of INC or DEC.

Warning

- 50 String assignments not allowed in standard Modula-2.
- 51 String comparisons not allowed in standard Modula-2.
- 52 & not allowed in standard Modula-2.

APPENDIX E

BNF Syntax for Turbo Modula-2

The syntax of the Turbo Modula-2 language is presented here using the formalism known as Backus-Naur Form (BNF). The following symbols are meta symbols belonging to the BNF formalism; they are not symbols of the language.

<term>	Names of language constructs are surrounded by "<" and ">".
{ X }*	Represents zero or more repetitions of X.
[X]	Means X is optional.
X	Means X is mandatory.
X Y	Indicates that X and Y are alternatives and that either X or Y must be used.
"X". or 'X'	Means X is written exactly as shown.
(X)	Means X must be chosen.

All other symbols are part of the language (reserved words of Turbo Modula-2 are in boldface type). For easy reference, the syntactic constructs are listed alphabetically.

```

<ActualParameters> ::= "(" [ <ExpList> ] ")"
<AddOperator> ::= "+" | "-" | "OR"
<ArrayType> ::= "ARRAY" <SimpleType> {",", <SimpleType>}*
"OF"
<type>
<assignment> ::= <designator> ":=" <expression>
<block> ::= { <declaration> }* [ "BEGIN" <StatementSequence> ]
[ <ExceptionHandler> ] "END"
<case> ::= [ <CaseLabelList> ":" <StatementSequence> ]

```

```

<CaseLabelList> ::= <CaseLabels> [ , " <CaseLabels>)*

<CaseLabels> ::= <ConstExpression> [ .." <ConstExpression> ]

<CaseStatement> ::= "CASE" <expression> "OF" <case>
    { "|" [ <case> ] }
    [ "ELSE" <StatementSequence> ]
    "END"

<character> ::= <letter> | <digit> | " " | ":" | ";" | "#" |
    "$" | "%" | "&" | ":" | "(" | ")" | "*" | "+" |
    "," | "-" | "." | "/" | ":" | ";" | "<" | "=" |
    ">" | "?" | "@" | "[" | "\\" | "]" | "^" | "_" |
    "{" | "[" | "]" | "}" | "~"

<CompilationUnit> ::= <DefinitionModule> | ["IMPLEMENTATION"]
                     <ProgramModule>

<ConstantDeclaration> ::= <ident> "=" <ConstExpression>

<ConstExpression> ::= <expression>

<declaration> ::= "CONST" { <ConstantDeclaration> ";" }* |
    "TYPE" { <TypeDeclaration> ";" }* |
    "VAR" { <VariableDeclaration> ";" }* |
    <ExceptionDeclaration> ";" |
    <ProcedureDeclaration> ";" |
    <ModuleDeclaration> ";"

<definition> ::= "CONST" { <ConstantDeclaration> ";" }* |
    "TYPE" { <ident> [=] <type> ";" }* |
    "VAR" { <VariableDeclaration> ";" }* |
    <ExceptionDeclaration> ";" |
    <ProcedureHeading> ";" |

<DefinitionModule> ::= "DEFINITION MODULE" <ident> ";" |
                      { <import> }* { <definition> }* "END"
                      <ident> "."

```

```
<designator> ::= qualident { "." <ident> |
  "[" "ExpList" "]"] | "^" ]*
```

```
<digit> ::= <octalDigit> | "8" | "9"
```

```
<element> ::= <expression> [ ".." <expression> ]
```

```
<enumeration> ::= "(" <IdentList> ")"
```

```
<exception> ::= [ <IdentList> ":" <StatementSequence> ]
```

```
( <export> ::= "EXPORT" [ "QUALIFIED" ] <IdentList> ";"
```

```
<expression> ::= <SimpleExpression> { <relation>
<SimpleExpression> }*
```

```
<ExceptionDeclaration> ::= "EXCEPTION" <ident> { "," <ident> }*
```

```
<ExceptionHandler> ::= "EXCEPTION" <exception> { "|" <exception> }*
[ "ELSE" <StatementSequence> ]
```

```
<ExpList> ::= <expression> { "," <expression> }*
```

```
<factor> ::= <number> | <string> | <set> | <designator>
[ <ActualParamters> ] |
  "(" <expression> ")" | "NOT" <factor>
```

```
( <FieldList> ::= [ <IdentList> ":" <type> |
  "CASE" [ <ident> ] ":" <qualident>
  "OF" <variant> [ "!" <variant> ]* |
  "ELSE" <FieldListSequence> ] "END" ]
```

```
<FieldListSequence> ::= <FieldList> [ ";" <FieldList> ]*
```

```
<ForStatement> ::= "FOR" <ident> ":" <expression> "TO"
  <expression>
  [ "BY" <ConstExpression> ]
  "DO" <StatementSequence> "END"
```

```

<FormalParamenters> ::= "(" [ <FPSection>
                                { ";" <FPSection> }* ] ")"
                                [ ":" <qualident> ]

<FormalType> ::= [ "ARRAY OF" <qualident> ]

<FormalTypeList> ::= "(" [ ["VAR"] <FormalType>
                            { "," ["VAR"] <FormalType>}* ] ")" [ ":" <qualident> ]

<FPSection> ::= [ "VAR" ] <IdentList> ":" <FormalType>

<hexDigit> ::= <digit> | "A" | "B" | "C" | "D" | "E" | "F"

<ident> ::= letter {letter | <digit>}*

<IdentList> ::= <ident> {"," <ident>}*

<IfStatement> ::= "IF" <expression> "THEN"
                  <StatementSequence>
                  ["ELSIF" <expression> "THEN"
                   <StatementSequence>}*
                  ["ELSE <StatementSequence>"]
                  "END"

<import> ::= [ "FROM" <ident> ] "IMPORT" <IdentList> ";"

<InlineCode> ::= "CODE" "(" <string> ")" "END"

<integer> ::= <digit> {<digit>}* |
                  {<octalDigit>}* ( "B" | "C" ) |
                  <digit> {<hexDigit>}*

<letter> ::= "A" | "B" | "C" | "D" | "E" | "F" | "G" | "H" |
                  "I" | "J" | "K" | "L" | "M" | "N" | "O" | "P" |
                  "Q" | "R" | "S" | "T" | "U" | "V" | "W" | "X" |
                  "Y" | "Z" | "a" | "b" | "c" | "d" | "e" | "f" |
                  "g" | "h" | "i" | "j" | "k" | "l" | "m" | "n" |
                  "o" | "p" | "q" | "r" | "s" | "t" | "u" | "v" |
                  "w" | "x" | "y" | "z"

```

```
<LoopStatement> ::= "LOOP" <StatementSequence> "END"

<ModuleDeclaration> ::= "MODULE" <ident> [ <priority> ] ";"  
                      [ <import> ] [ <export> ] <block>  
                      <ident>

<MulOperator> ::= "*" | "/" | "DIV" | "MOD" | "AND"

<number> ::= <integer> | <real>

<octalDigit> ::= "0" | "1" | "2" | "3" | "4" | "5" | "6" | "7"

<PointerType> ::= "POINTER TO" <type>

<priority> ::= "[" <ConstExpression> "]"

<ProcedureCall> ::= <designator> [ <ActualParameters> ]

<ProcedureDeclaration> ::= <ProcedureHeading> ";"  
                           ( <block> | <InlineCode> )  
                           <ident>

<ProcedureHeading> ::= "PROCEDURE" <ident>  
[ <FormalParameters> ]

<ProcedureType> ::= "PROCEDURE" [FormalTypeList]

<ProgramModule> ::= "MODULE" <ident> [ <priority> ] ";"  
                     { <import> }* <block> <ident> "."

<qualident> ::= <ident> {"." <ident>}*

<RaiseStatement> ::= "RAISE" [ <ident> [ "," <expression> ] ]

<real> ::= <digit> {<digit>}* "." {<digit>}* [ ScaleFactor ]

<RecordType> ::= "RECORD" <FieldListSequence> "END"

<relation> ::= "=" | ">" | "<" | ">=" | "<=" | "#" | "<>"  
| "IN"
```

```
<RepeatStatement> ::= "REPEAT" <StatementSequence> "UNTIL"
<expression>

<ScaleFactor> ::= ( "E" | "D" ) [ "+" | "-" ] <digit>
{<digit>}*
```

```
<set> ::= [ <qualident> ] "(" [ <element> { "," <element> }* ] ")"
```

```
<SetType> ::= "SET OF" <SimpleType>
```

```
<SimpleExpression> ::= [ "+" | "-" ] <term> { <AddOperator> <term> }*
```

```
<SimpleType> ::= <qualident> | <enumeration> | <SubrangeType>
```

```
<statement> ::= [ assignment | ProcedureCall | IfStatement |
CaseStatement | WhileStatement | RepeatStatement |
LoopStatement | ForStatement | WithStatement |
RaiseStatement | "EXIT" | "RETURN" { expression } ]
```

```
<StatementSequence> ::= <statement> { ";" <statement> }*
```

```
<string> ::= "" {<character>}* "" | '' {<character>}* ''
```

```
<SubrangeType> ::= [ <qualident> ] "[" <ConstExpression> ".." <ConstExpression> "]"
```

```
<term> ::= <factor> { <MulOperator> <factor> }*
```

```
<type> ::= <SimpleType> | <ArrayType> | <RecordType> |
<SetType> | <PointerType> | <ProcedureType>
```

```
<TypeDeclaration> ::= <ident> "=" <type>
```

```
<VariableDeclaration> ::= <ident> [ "[" <ConstExpression> "]"
    { "," <ident> [ "[" <ConstExpression>
        "]"] }*
    ":" <type>

<variant> ::= [ <CaseLabelList> ":" <FieldListSequence> ]

<WhileStatement> ::= "WHILE" <expression> "DO"
    <StatementSequence>
    "END"

<WithStatement> ::= "WITH" <designator> "DO"
    <StatementSequence>
    "END"
```

Index

- Abort command, 150
ABS, 273
ABS standard function, 273
Absolute Addresses, 117
Acknowledgements, 19
ADDRESS, 252
(ADDRESS type, 250
ADDRESS type (SYSTEM), 273
ADR, 250
ADR procedure (SYSTEM), 275
ALLOCATE, 260
ALLOCATE procedure (STORAGE), 276
AND operator, 277
Append, 239
APPEND procedure (Strings), 278
Arctan, 233
Arctan function (MathLib
 LongMath), 278
ArgumentError, 280
Array of Word, 248
ARRAY standard type, 281
Array type, 62
Arrays, 555
Assembler interface, 254
Assignment Statements, 72
available variable, 284
(voiding the Menu, 137
Awaited, 245, 285
Awaited procedure (Processes), 285

BADOVERLAY, 286
BDOS, 253
BDOS procedure (SYSTEM), 287
BEGIN, 288
Begin block, 141
BIOS, 254

BIOS procedure (SYSTEM), 288
BITSET, 55
BITSET standard type, 290
Block Commands, 141, 145
 Begin block, 141, 145
 Copy block, 141, 146
 Delete block, 141, 146
 End block, 141, 146
 Hide/display block, 141, 146
 Move block, 141, 146
 Read block from disk, 141, 146
 Write block to disk, 141, 146
BNF Syntax, 597
BOOLEAN standard type, 291
Boolean type, 55
BusyRead, 224
BusyRead procedure (Terminal), 292
BYTE type (SYSTEM), 294

Call procedure (Loader), 295
Calling Chain, 590
CAP standard function, 296
CAPS, 239
CAPS procedure (Strings), 296
CARD standard function, 297
CARDINAL standard type, 298
Cardinal type, 55
CardToStr procedure (Convert), 299
CASE statement, 76, 300, 565
CaseSelectError, 588
CHAR standard type, 303
CHAR type, 54, 122
Character left, 140, 141
Character right, 140, 141
Character Set, 37
CHR standard function, 304
ClearEol, 305
ClearScreen, 226, 576
ClearScreen procedure (Terminal), 306
ClearToEOL procedure (Terminal), 307
Close, 215
Close procedure (Files), 308
CloseInput procedure (InOut), 309

- CloseOutput procedure (InOut), 310
CloseText, 193
CloseText procedure (Texts), 311
CODE procedure (SYSTEM), 312
Col, 200
Col function (Texts), 313
ComLine library module, 182
ComLine Module, 229, 314
commandLine variable, 317
Comments, 42, 553
Compiler, 152
Compiler Directives, 583
Compiler Error Messages, 592
Compiler Installation, 581
Compiler Options and Switches, 163
Compress, 34
Conditional Statements, 75
ConnectDriver procedure (Texts), 318
console (Texts), 319
CONST declaration, 320
Constant Declaration, 57
Control-C, 140
Convert library module, 183
Convert Module, 241, 321
Copy, 34, 239
Copy block, 140, 146
Copy procedure (Strings), 322
Coroutines, 118,
Coroutines and Interrupts, 118, 251
Cos function (MathLib LongMath), 323
Create procedure (Files), 324
CreateText, 193
CreateText procedure (Texts), 325
Ctrl-A, 140
Ctrl-D, 140
Ctrl-E, 140
Ctrl-F, 140
Ctrl-G, 140
Ctrl-K-B, 141
Ctrl-K-C, 141
Ctrl-K-D, 141
Ctrl-K-H, 141
Ctrl-K-K, 141
Ctrl-K-R, 141
Ctrl-K-V, 141
Ctrl-K-W, 141
Ctrl-K-Y, 141
Ctrl-L, 141
Ctrl-N, 140
Ctrl-Q-A, 141
Ctrl-Q-B, 140
Ctrl-Q-C, 140
Ctrl-Q-D, 140
Ctrl-Q-E, 140
Ctrl-Q-F, 141
Ctrl-Q-I, 140
Ctrl-Q-K, 140
Ctrl-Q-P, 140
Ctrl-Q-R, 140
Ctrl-Q-S, 140
Ctrl-Q-X, 140
Ctrl-Q-Y, 140
Ctrl-R, 140
Ctrl-S, 140
Ctrl-T, 140
Ctrl-U, 141
Ctrl-V, 140
Ctrl-W, 140
Ctrl-X, 140
Ctrl-Y, 140
Ctrl-Z, 140
Cursor, 579
Cursor Movement Commands, 140
Character left, 140
Character right, 140
Line down, 140
Line up, 140
Page down, 140
Page up, 140
Scroll down, 140
Scroll up, 140
To top of window, 140
Word left, 140
Word right, 140

Data Types, 54
Data

- identifier, 53,
type, 53
value, 53,
DeadLock, 245, 326
DEALLOCATE, 260
DEALLOCATE procedure (STORAGE),
 327
DEC standard procedure, 328
Declarations, 56, 554
Declarations of Exceptions, 124
DEFINITION declaration, 330
Delete, 239
Delete file, 150
Delete block, 141, 146
Delete character under cursor, 140
Delete commands, 140
 Delete left character, 140
 Delete line, 140
 Delete right word, 140
 Delete to end of line, 140
Delete left character, 140
Delete line, 140
Delete procedure (Files), 332
Delete procedure (Strings), 333
Delete right word, 140
Delete to end of line, 140
DeleteLine, 334
DeleteLine procedure (Terminal), 334
Delimiters, 41
DeviceError, 335, 589
Dir, 132
directory, 34
DiskFull, 336, 589
DISPOSE standard procedure
 (**STORAGE**), 337
DIV operator, 339
Done, 197
Done procedure (InOut), 339
Done procedure (Texts), 340
Double standard function, 341
Double-precision, 39
Doubles, 342
Doubles Module, 183, 242
DoubleToStr procedure (Doubles), 343
Dynamic Variable Errors, 263
Editing Commands, 139
 Character left, 140
 Character right, 140
 Line down, 140
 Line up, 140
 Page down, 140
 Page up, 140
 Scroll down, 140
 Scroll up, 140
 To beginning of block, 140
 To bottom of window, 140
 To end of block, 140
 To end of file, 140
 To last position, 140
 To left on line, 140
 To right on line, 140
 To top of file, 140
 To top of window, 140
 Word left, 140
 Word right, 140
END, 343
End block, 140
EndError, 344, 589
EndOfCoroutine, 588
Entier, 235, 345
Enumeration type, 59,
EOF, 218
EOF procedure (Files), 346
EOL, 347
EOLN, 198
EOLN procedure (Texts), 348
EOT, 199
EOT procedure (Texts), 349
Error/Compiler Limit Exceeded, 598
Error Correction, 157
Error diagnosis, 585
Errors detected by Support Modules, 589
Errors Detected by the Interpreter, 586
Errors during File Handling, 210
Error, General, 594
Error Implementation Restriction, 595

- Error in constants, 594
in Identifier, 593
in Type, 594
in Syntax, 593
Error Implementation, 595
EXCEPTION, 350
EXCEPTION Handling, 123
Exception Issued by Module Loader, 589
EXCEPTION
handlers, 123, 128
propagation, 127
raised from another exception handler, 124
Exceptions Issued by Module, 589
EXCL standard procedure, 351
EXIT statement, 80, 352
ExitScreen procedure (Terminal), 353
Exp, 233
Exp function (MathLib LongMath), 353
EXPORT, 354
Export
opaque, 105
Expressions, 49, 561
Extended Movement Commands, 140, 143
To beginning of block, 140, 143
To bottom of window, 140, 143
To end of block, 140, 144
To end of file, 140, 143
To last position, 140, 144
To left on line, 140, 143
To right on line, 140, 143
To top of file, 140, 143
Extensions, 121
- FALSE**, 356
Field type, (Files), 356
Filecopy, 35, 133
File Management Utilities, 132
File Processing, 216
Files library module, 182
Files Module, 210, 357
Files On Your Diskette, 21
- Files with Elements of Mixed Types, 222
FileSize, 359
FileSize procedure (Files), 359
FILL, 253
FILL procedure (SYSTEM), 360
Find, 141, 147
Find and replace, 141, 148
Find And Replace Commands, 141, 147
Find run-time error, 136, 592
firstDrive, 360
FLOAT standard function, 361
Flush procedure (Files), 362
FOR statement, 77, 363
Format of a Run-Time Error message, 585
Forward statement, 93, 364
FREEMEM, 263
FREEMEM procedure (STORAGE), 365
Function Procedure, 89, 560
FunctionReturnsNoResult, 588
- GetName** procedure (Files), 366
Getting Started, 21
GotoXY, 226
GoToXY procedure (Terminal), 367
- HALT** standard procedure, 368
haltOnControlC, 368
heap-pointer, 260
Hide/display block, 141, 146
HIGH standard function, 370
highlightNormal, 372
HighLight procedure (Terminal), 371
HRESULT, 373
- Identifiers, 44, 551
Identifiers,
list of library, 46
IF statement, 75, 374
IllegalInstruction, 588

- IMPLEMENTATION MODULE, 252
Implementation declaration, 375
IMPORT declaration, 378
INC standard procedure, 379
INCL standard procedure, 382
Indent On/Off, 150
Init, 245
Init procedure (Processes), 382
Initialization string, 575
InitScreen procedure (Terminal), 383
inName variable, 384
InOut module, 182, 207, 385
INP, 253
INP procedure (SYSTEM), 387
Input Output, 184
Input and Output Extensions, 121
Input and Output Modules, 182
input (Texts), 388
Insert, 138
insertDelete, 389
Insert and Delete Commands, 140
Insert commands, 144
 Insert line, 140, 145
Insert line, 140, 145
Insert mode, 140, 145
Insert mode on/off switch, 140
Insert procedure (Strings), 389
InsertLine, 145
InsertLine procedure (Terminal), 391
Installation of Editing Commands, 577
Installation of screen, 572
Installation Procedures, 571
INT standard function, 392
Integer Numbers, 39
INTEGER standard type, 393
Integer type, 55
Interface to CP/M, 253
Interrupts, 119
IntToStr procedure (Convert), 394
IORESULT, 386
IOTRANSFER, 252
IOTRANSFER procedure (SYSTEM), 395
Kill, 35, 133
Language Elements, 37
legal variable, 396
Length, 237
Length procedure (Strings), 397
Librarian, 151
Library, 181
Library modules, 100
Library Identifiers, 45
 list of, 45
Line down, 140, 142
Line up, 140, 142
Linker, 167, 173
Linking Microsoft .REL-Files, 175
Linking with overlays, 169
Linking
 Microsoft, 175
Listings, 154
Ln, 233
Ln function (MathLib)
 LongMath, 401
Loader module, 184, 264, 402
LoadError, 397, 589
Local Modules, 106
Logical devices, 187
LONG standard function, 403
LONGINT standard type, 56, 404
LongMath, 233
LongMath module, 183, 233, 405
LONGREAL standard type, 56, 399
LongToStr procedure (Convert), 406
LOOP statement, 80, 407
Low Level Access to Data, 248
Low Level Facilities, 111
Low Level Types, 115

Main Module, 97
Manual Installation, 573
MARK, 262
MARK procedure (STORAGE), 408
MathLib, 233

- MathLib module, 183, 233, 409
MAX standard function, 410
Memory Management, 258
Menu
 avoiding it, 137
 System, 29
MIN standard function, 411
Miscellaneous commands, 141, 149
 Abort command, 141
 Auto tab on/off, 141
 Save file, 141
 Tab, 141
Miscellaneous Editing Commands, 141, 149
MOD standard operator, 412
MODULE declaration, 413
Module Library, 100, 184
Modules, 97, 104
MOVE, 253
Move block, 141, 146
Multidimensional Array, 122
- Nested Procedures, 91
NEW, 260
NEW standard procedure, 414
NEWPROCESS, 251
NEWPROCESS procedure (SYSTEM), 416
NextPos, 219
NextPos procedure (Files), 417
NIL, 419
Normal procedure (Terminal), 419
NoTrailer, 216
NoTrailer procedure (Files), 420
Numbers, 38
numCols, 421
numRows, 422
- ODD, standard function, 423
Opaque export, 105
Open array parameters, 86, 559
open arrays, 122
- Open procedure (Files), 423
Opening, Creating, and Closing a Text, 193
OpenInput procedure (InOut), 425
OpenOutput procedure (InOut), 426
OpenText, 193
OpenText procedure (Texts), 427
Operands, 50
Operating the Compiler, 153
Operating the Editor, 139
Operations on Entire Files, 214
Operator, 42, 50
Operator Precedence, 51
Operator
 arithmetic, 51
 logical, 51
 precedence, 51
 relational, 50
 set, 51
OpSet, 428
Options, 135, 583
OR, 429
ORD standard function, 430
OUT, 253
OUT procedure (SYSTEM), 431
outName variable, 432
OUTOFMEMORY, 433, 588
output (Texts), 434
OVERFLOW, 435, 587
- Page down, 140, 142
Page up, 140, 142
Parameter, 84
Pascal, 398
Path to search, 136
Pointer type, 61, 436
PointerError, 588
Pos procedure (Strings), 438
PROC standard type, 439
Procedural Statements, 81
Procedure call, 82
Procedure declarations, 82, 440, 559
PROCEDURE type, 68, 441, 557

- PROCESS type (SYSTEM), 443
Processes module, 183, 243, 445
Profile, 178
progName variable, 446
PromptFor, 229
PromptFor procedure (ComLine), 446
Pseudo-module, 115, 248
- QUALIFIED**, 448
Quit, 137
 no Save, 150
- RAISE**, 124
RAISE Statement, 449
Random, 236
Random function (MathLib), 450
Randomize, 236, 451
READ, 452
READ and WRITE Statements, 201,
Read block from disk, 141, 146
ReadAgain procedure (Terminal), 453
ReadAgain procedure (Texts), 454
ReadAgain, 198
 Terminal, 224
ReadByte, 455
ReadBytes, 218
ReadByte procedure (Files), 456
ReadCard, 196
ReadCard procedure (Texts), 458
ReadChar procedure (Terminal), 459
ReadChar procedure (Texts), 459
ReadChar, 196
 Terminal, 224
ReadDouble procedure
 (Doubles), 460
Reading and Writing, 195
ReadInt, 196
ReadInt (Texts) procedure, 461
ReadLine procedure (Texts), 462
ReadLine, 198, 463
 Terminal, 224
READLN, 198, 164
- ReadLong, 197
ReadLong procedure (Texts), 465
ReadReal, 196
ReadReal procedure (Texts), 466
ReadRec procedure (Files), 466
ReadString procedure (Texts), 468
ReadString, 196
ReadWord procedure (Files), 468
Real Numbers, 40
REAL standard type, 470
Real type, 56
REALOVERFLOW, 472, 587
RealToStr procedure (Convert), 473
Record type, 64, 475
Records, 556
RedirectInput, 229
RedirectInput procedure
 (ComLine), 479
RedirectOutput, 229
RedirectOutput procedure
 (ComLine), 480
RELEASE, 262
RELEASE procedure
 (STORAGE), 481
Rename, 36, 134
Rename procedure (Files), 482
Renaming, Deleteing and other File
operations, 195
REPEAT last find, 141, 149
REPEAT statement, 79, 484
Repetitive Statements, 77
Reserved Words, 43
 list of, 43
Reset Options, 575
ResetSys, 214
ResetSys procedure (Files), 484
RETURN statement, 79, 90, 485, 566
RETURN
 statement, 90
Routine Statement, 79
Running Out of Memory, 159

- Save and Edit, 150
and Quit, 150
Save current selection, 136
Save file, 141
Scope and local modules, 107
Scope of visibility, 92
Screen Installation, 572
Scroll down, 140, 142
Scroll up, 140, 142
Search, 141
Searching libraries, 151
SEND, 245
SEND procedure (Processes), 486
Set constants, 553
Set Operators, 562
SET TYPE, 67, 487
Set Types, 67
SetCol, 200
SetCol procedure (Texts), 489
SetPos, 219
SetPos procedure (Files), 490
SIGNAL, 491
Sin function (MathLib
 LongMath), 492
SIZE procedure, 493
Special Ops, 494
Specification of MathLib and LongMath,
 234
Specification of the Module Files, 212
Specification of the module Loader, 269
Specification of the Module
 Processes, 244
Specification of the Module
 STORAGE, 261
Specification of the Module
 Strings, 238
Specification of the module
 Terminal, 225
Specification of the Module Texts,
 190
Specification of the Pseudo-
 Module SYSTEM, 249
Sqrt, 233
- Sqrt function (MathLib
 LongMath), 495
Standard Identifiers, 44
Standard Library, 181
Standard Procedures, 94, 566
Standard text stream, 188
StartProcess, 243
StartProcess procedure
 (Processes), 496
Statements, 71, 564
 assignment, 72
 CASE, 76
 EXIT, 80
 FOR, 77
 IF, 75
 LOOP, 80
 REPEAT, 79
 repetitive, 77
 WHILE, 79
 WITH, 73
StatusError, 497
STORAGE Module, 184, 260
STORAGE pseudo-module, 498
String, 499, 552
StringError, 500
Strings, 40
Strings module, 183, 237, 501
StringTooLong, 587
String Extensions, 122
StrToCard procedure (Convert), 502
StrToDouble procedure
 (Doubles), 503
StrToInt procedure (Convert), 504
StrToLong procedure (Convert), 505
StrToReal procedure (Convert), 506
Structure of This Manual, 17
Structured Type, 62
Subrange type, 60
Summary of Compiler Directives, 583
Symbol files, 162
Syntax and Semantics of Exception
handling, 122
System and Low Level Modules, 183
SYSTEM Module, 184

- SYSTEM pseudomodule, 248, 506
Tab, 150
termCH, 508
Terminal Codes, 574
Terminal module, 182, 224, 508
Terminal properties, 575
Terminal Type, 575
TEXT data structure (Texts), 510
TextDriver, 511
TextFile, 512
(.tNotOpen, 513
Text module, 183, 188, 514
The Library Module Convert, 183
The Library Module Doubles, 183
The Library Module Texts, 183, 188
The Library Modules MathLib and
 LongMath, 183, 233
The Linker, 167
The Module Loader, 264
The Module STORAGE, 261
The Module Strings, 237
The Pseudo-Module SYSTEM, 248
TINST.DTA, 574
To beginning of block, 144
To bottom of window, 143
To end of block, 144
To end of file, 143
To last position, 144
To left on line, 143
To right on line, 143
To top of file, 143
(, top of window, 143
TooManyTexts, 517
TRANSFER, 252
TRANSFER procedure (System), 518
TRUE standard value, 519
TRUNC standard function, 520
TSIZE procedure (SYSTEM), 521
Turbo Editor, 138
TURBO Modula Extensions, 121
TURBO Modula
 files, 21, 29
load, 26
Type, 36, 39, 134, 237
Type declaration, 58, 522, 555
Type transfer, 113
Type transfer and type conversion, 114
Type
 array, 62
 Boolean, 55
 cardinal, 55
 char, 54
 enumeration, 59
 integer, 55
 pointer, 61
 procedure, 68
 real, 56
 record, 64
 set, 67
 structured, 62
 subrange, 60
 unstructured, 59
Typography, 18

Unstructured Types, 59
UseError, 523, 589
User-Defined Identifiers, 45
Untyped parameters, 116, 560
Utilities, 175
Utility Modules, 183, 233

VAL standard function, 524
VAR declaration, 525
Variable Declaration, 58, 559
Variant record, 66
Version Control, 174
Vocabulary, 38

WAIT, 245
WAIT procedure (Processes), 526
WHILE statement, 79, 527
WITH statement, 73, 528, 566
WORD, 92,

WORD type (SYSTEM), 529
WRITE, 531
Write block to disk, 141, 146
WriteByte procedures (Files), 532
WriteBytes, 218
WriteBytes procedures (Files), 533
WriteCard, 199
WriteCard procedure (Texts), 533
WriteChar, 199
WriteChar procedure (Terminal), 534
WriteChar procedure (Texts), 535
WriteChar
 Terminal, 224
WriteDouble procedure
 (Doubles), 536
WriteHex procedure (InOut), 537
WriteInt, 199
WriteInt procedure (Texts), 538
WRITELN, 538
WriteLn procedure (Terminal), 540
WriteLn procedure (Texts), 541
WriteLn, 200, 224
WriteLn procedure (InOut), 541
WriteLong, 200
WriteLong procedure (Texts), 542
WriteOct procedure (InOut), 543
WriteReal, 199
WriteReal procedure (Texts), 544
WriteRec procedure (Files), 545
WriteString, 199
WriteString procedure (Texts), 546
WriteString procedure (Terminal), 545
WriteString
 Terminal, 224, 545
 Texts, 546
WriteWord, 547
WriteWord procedure (Files), 547

Z80 Specific Procedures, 253

Reflex

Reflex ist mehr als nur die außergewöhnliche Kombination aus Datenbank und Tabellenkalkulation. Reflex ist Datenanalyse. Der direkte Weg zu den versteckten Wahrheiten hinter reiner Information.

In der Verbindung aus grafischen Eigenschaften, übersichtlicher Fenster-technik, hoher Programmgeschwindigkeit und einem vollständigen Satz von Funktionen schafft Reflex eine Transparenz, die Ihnen bisher kein anderes Datenbankprogramm geboten hat.

Reflex macht schnell

Da Reflex's Daten im Hauptspeicher gehalten werden, sind die meisten Operationen wie sortieren, filtern etc. sehr schnell. Durch die Above-Board-Unterstützung kann Reflex aber auch mit größeren Datenmengen umgehen, ohne Sie auf die Folter zu spannen.

Reflex kennt fünf verschiedene Ansichten: Die Listenansicht, die Formularansicht, die Kreuztabelle, die Grafik- und die Berichtsansicht. Während Sie in der Listenansicht alle Daten auf einen Blick sehen und erfassen, lässt sich der vollständige Datensatz im Formular darstellen. Was Sie gerade in der einen

Ansicht bearbeiten, sehen Sie sofort auch als Grafik in der Grafikansicht - als Balken-, Kreis- oder Liniendiagramm.

Reflex kann mehr

Mit Hilfe der Kreuztabelle bekommen Sie ein Instrument, von dem viele Spreadsheet-Eigner träumen: Denn hier können Sie verschiedene Datensätze und Datenfelder miteinander in Beziehung setzen,

und zwar so lange, bis Sie nur das sehen, worauf es Ihnen ankommt.

Reflex ist durch seine grafische Benutzeroberfläche leicht zu beherrschen. Pull-down-Menüs und Dialogboxen führen Sie sicher durch die Vielfalt von Funktionen. Ihren ersten Bericht oder Ihre erste Grafik erzeugen Sie schon nach den ersten 15 Minuten.

Mit unzähligen Funktionen aus den Bereichen Betriebswirtschaft, Finanzen, Statistik und Mathematik lösen Sie die meisten kalkulatorischen Probleme. Diese Kombination finden Sie bei kaum einer anderen Datenbank.

Technische Daten

- 250 Felder pro Datensatz
- 254 Zeichen pro Datenfeld
- 64.000 Datensätze pro Datenbank
- Feldformate: numerisch, Ganzzahl, Datum oder Text
- suchen und filtern in den verschiedensten booleschen Verknüpfungen.
- automatisch berechnete Felder
- Logik- und Datumsfunktionen,
- VARY-Funktion generiert Datensätze nach bestimmten Verlaufsvorgaben
- Fremdatenübernahme aus Lotus 1-2-3, dBase II, DIF (Multiplan) und ASCII

Systemvoraussetzungen

IBM PC bzw. kompatibel,
oder Siemens PCD
mindestens 384 KByte Hauptspeicher,
Farbgrafikadapter, Herculeskarte oder
EGA-Karte

HEIMSOETH & BORLAND