

# Rust Crypto For XChain

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Xuper-sdk-go vs rust-sgx for Crypto

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## Xuper-sdk-go vs rust-sgx for Crypto

	超级链	crate-rust-sgx
ecdsa	crypto/ecdsa: P256-SHA256-ANS1	ring::P256-SHA256-ASN1 ring::P256-SHA384-ASN1
hash	crypto/hmac crypto/sha512 crypto/sha256 "golang.org/x/crypto/ripemd160"	ring::{hmac,sha256,sha512} ripemd160
encode	self/base58 自己实现的	base58
bigint	math/bigint	num-bigint
rand	crypto/rand	rand
aes	crypto/aes (Rijndael 128, 192, 256)	ring
ecies	<a href="#">Kylom's implementation</a> curve: P256	需要实现
sign	multi_sign, schnorr_ring_sign, schnorr_sign	需要实现
hdwallet/keychain	hdwallet/keychain	需要实现

# ECC

## ECDSA

Parameter	
CURVE	the elliptic curve field and equation used
$G$	elliptic curve base point, a point on the curve that generates a <a href="#">subgroup of large prime order <math>n</math></a>
$n$	integer order of $G$ , means that $n \times G = O$ , where $O$ is the identity element.
$k$	the private key (randomly selected)
$P$	the public key (calculated by elliptic curve)
$M$	the message to send

### ECDSA签名：

$$\begin{aligned}P &= (x_1, y_1) = k \times G \\S &= k^{-1}(\text{Hash}(M) + k * x_1) \bmod p \\Signature &= (x_1, S)\end{aligned}$$

### ECDSA 验证签名

$$\begin{aligned}P' &= S^{-1} * \text{Hash}(M) \times G + S^{-1} * x_1 \times P \\&= P\end{aligned}$$

- 证明

$$\begin{aligned}P' &= S^{-1} * \text{Hash}(M) \times G + S^{-1} * k \times G \\&= (S^{-1} * \text{Hash}(M) + S^{-1} * k) \times G \\&= (\text{Hash}(M) + x_1) * S^{-1} \times G \\&= (\text{Hash}(M) + x_1) * (k^{-1}(\text{Hash}(M) + k))^{-1} \times G \\&= (\text{Hash}(M) + x_1) * k * (\text{Hash}(M) + k)^{-1} \times G \\&= k \times G \\&= (x_1, y_1)\end{aligned}$$

## ECIES算法

为了向Bob发送ECIES加密信息，Alice需要以下信息：

- 密码学套件（KDF，MAC，对称加密E）

- 椭圆曲线(p, a, b, G, n, h)
- Bob的公钥:

$$K_b, K_b = k_b G, k_b \in [1, n - 1]$$

- 共享信息

$$S_1, S_2$$

- 无穷远点O

## 加密

Alice使用Bob的公钥加密消息m:

*For random  $r \in [1, n - 1]$ , calculate  $R = rG$*   
*derive shared secret :  $S = P_x$ , where  $P = P(P_x, P_y) = rK_b, P \neq O$*   
*derive  $K_E || K_M = KDF(S || S_1)$*   
*encrypt message  $m : c = E(k_E; m)$*   
*calculate MAC :  $d = MAC(k_M; c || S_2)$*   
*output :  $R || c || d$*

## 解密

Bob解密密文  $R || c || d$  的步骤如下:

*derive shared secret :  $S = P_x, P = P(P_x, P_y) = k_B R$*   
*derive  $K_E || K_M = KDF(S || S_1)$*   
*verify MAC :  $d == MAC(k_M; c || S_2)$*   
*decrypt :  $m = E^{-1}(k_E; c)$*

## 证明过程

we need ensure S is really shared by Alice and Bob:

$$P = K_B r = k_B R$$

## Refer

1. [https://en.wikipedia.org/wiki/Integrated\\_Encryption\\_Scheme](https://en.wikipedia.org/wiki/Integrated_Encryption_Scheme)

## 有限域运算

## 表示

有限域:

$$GF(p^n)$$

这里专门针对p=2为特征的多项式进行计算。

## 多项式基表示法

## 正规基表示法

## 运算

### 加法

### 减法

### 乘法

### 除法

### 求逆