内核笔记

段武杰

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VFS

1.1 注册一个文件系统

```
#include <linux/fs.h>
extern int register_filesystem(struct file_system_type *);
extern int unregister_filesystem(struct file_system_type *);
```

/proc/filesystems 下显示了所有已经注册的文件系统。

file_system_type 用于描述一个文件系统:

```
struct lock_class_key i_lock_key;
struct lock_class_key i_mutex_key;
struct lock_class_key i_mutex_dir_key;
};
```

The mount() method must return the root dentry of the tree requested by caller. An active reference to its superblock must be grabbed and the superblock must be locked. On failure it should return ERR_PTR(error).

The arguments match those of mount(2) and their interpretation depends on filesystem type. E.g. for block filesystems, dev_name is interpreted as block device name, that device is opened and if it contains a suitable filesystem image the method creates and initializes struct super_block accordingly, returning its root dentry to caller.

->mount() may choose to return a subtree of existing filesystem - it doesn't have to create a new one. The main result from the caller's point of view is a reference to dentry at the root of (sub)tree to be attached; creation of new superblock is a common side effect.

The most interesting member of the superblock structure that the mount() method fills in is the "s_op" field. This is a pointer to a "struct super_operations" which describes the next level of the filesystem implementation.

Usually, a filesystem uses one of the generic mount() implementations and provides a fill_super() callback instead. The generic variants are:

- mount_bdev: mount a filesystem residing on a block device
- mount_nodev: mount a filesystem that is not backed by a device
- mount_single: mount a filesystem which shares the instance between all mounts

A fill_super() callback implementation has the following arguments:

struct super_block *sb: the superblock structure. The callback must initialize this properly.

- void *data: arbitrary mount options, usually comes as an ASCII string (see
 "Mount Options" section)
- int silent: whether or not to be silent on error

VFS 使用面向对象的设计思路, VFS 中有 4 个主要的对象类型:

- 超级块对象 (super_block): 它表示一个具体的已安装的文件系统
- 索引节点对象 (inode): 它表示一个具体的文件
- 目录项对象 (dentry): 它表示一个目录项, 是路径的一个组成部分。
- 文件对象 (file): 它表示进程打开的文件。

VFS 将目录当作文件来处理,所以不存在目录对象,目录项代表的是路径中的一个组成部分。

1.2 super_block

A superblock object represents a mounted filesystem.

```
struct super_block {
struct list_head s_list; /* 指向的链表super_block */
dev_t s_dev; /* 设备标识符 */
unsigned char s_blocksize_bits;/* 以位为单位的块大小 */
unsigned long s_blocksize;/* 以字节为单位的块大小 */
loff_t s_maxbytes; /* Max file size */
struct file_system_type *s_type;/* Filesystem type */
const struct super_operations *s_op;/*超级块方法*/
const struct dquot_operations *dq_op;/*磁盘限额方法 */
```

```
*s_qcop;/* 限额控制方法 */
      const struct quotactl_ops
      const struct export_operations *s_export_op;/* 导出方法 */
      unsigned long
                       s_flags;/* 挂载标志 */
      unsigned long
                          s_iflags; /* internal SB_I_* flags */
      unsigned long
                         s_magic;/* 文件系统魔数 */
                         *s_root;/* 目录挂载点 */
      struct dentry
      struct rw_semaphore s_umount;/* 卸载信号量 */
                          s_count; /* 超级块引用计数 */
                          s_active;/* 活动引用计数 */
      atomic_t
#ifdef CONFIG_SECURITY
      void
                          *s_security;/* 安全模块 */
#endif
      const struct xattr_handler **s_xattr;/*扩展的属性操作*/
      struct hlist_bl_head s_anon;
                                       /* anonymous dentries for (nfs) exporting */
                         s_mounts; /* list of mounts; _not_ for fs use */
      struct list_head
      struct block_device *s_bdev;/*相关的块设备*/
      struct backing_dev_info *s_bdi;
      struct mtd info
                         *s_mtd;
      struct hlist_node
                       s_instances;
      unsigned int
                       s_quota_types; /* Bitmask of supported quota types */
      struct quota_info s_dquot;
                                      /* Diskquota specific options */
      struct sb_writers
                         s_writers;
      char s_id[32];
                                        /* Informational name */
      u8 s_uuid[16];
                                        /* UUID */
                                       /* 文件系统私有数据 */
                         *s_fs_info;
      unsigned int
                         s_max_links;
      fmode_t
                           s_mode;
      /* Granularity of c/m/atime in ns.
        Cannot be worse than a second */
      u32
                     s_time_gran;
       * The next field is for VFS *only*. No filesystems have any business
       * even looking at it. You had been warned.
      * Filesystem subtype. If non-empty the filesystem type field
       * in /proc/mounts will be "type.subtype"
       */
      char *s subtype:
       * Saved mount options for lazy filesystems using
       * generic_show_options()
```

```
char __rcu *s_options;
       const struct dentry_operations *s_d_op; /* default d_op for dentries */
        * Saved pool identifier for cleancache (-1 means none)
       int cleancache_poolid;
       struct shrinker s_shrink;
                                     /* per-sb shrinker handle */
       /* Number of inodes with nlink == 0 but still referenced */
       atomic_long_t s_remove_count;
       /* Being remounted read-only */
       int s_readonly_remount;
       /* AIO completions deferred from interrupt context */
       struct workqueue_struct *s_dio_done_wq;
       struct hlist_head s_pins;
        \boldsymbol{\ast} Keep the lru lists last in the structure so they always sit on their
        * own individual cachelines.
                            s_dentry_lru ____cacheline_aligned_in_smp;
       struct list_lru
       struct list_lru
                             s_inode_lru ____cacheline_aligned_in_smp;
       struct rcu_head
       struct work_struct destroy_work;
                        s_sync_lock; /* sync serialisation lock */
       struct mutex
        * Indicates how deep in a filesystem stack this SB is
       int s_stack_depth;
       /* \  \, s\_inode\_list\_lock \  \, protects \  \, s\_inodes \  \, */
                           s_inode_list_lock ____cacheline_aligned_in_smp;
       struct list_head
                             s_inodes; /* all inodes */
};
```

1.3 inode

内核处理文件的关键是 inode,每个文件(和目录)都有且只有一个对应的 inode,其中包含元数据(如访问权限,上次修改的日期,等等)和指向文件数据的指针。

```
* Keep mostly read-only and often accessed (especially for
* the RCU path lookup and 'stat' data) fields at the beginning
* of the 'struct inode'
*/
struct inode {/* fs.h */
      umode_t
                            i_mode;/* 文件访问权限和所有权 */
      unsigned short
                          i_opflags;
                          i_uid;/* uid about the file */
      kuid_t
                          i_gid;/* gid about the file */
      kgid_t
       unsigned int
                            i_flags;
#ifdef CONFIG_FS_POSIX_ACL
      struct posix_acl
                         *i_acl;
      struct posix_acl
                          *i_default_acl;
#endif
       /* 负责管理结构性操作(如删除一个文件)和文件相关的元数据例如属性()*/
       const struct inode_operations *i_op;
       struct super_block
       struct address_space *i_mapping;
#ifdef CONFIG_SECURITY
      void
                          *i_security;
#endif
       /* Stat data, not accessed from path walking */
       /* 对给定的文件系统, 唯一的编号标识 */
       unsigned long
                          i_ino;
       * Filesystems may only read i_nlink directly. They shall use the
       * following functions for modification:
           (set|clear|inc|drop)_nlink
           inode_(inc|dec)_link_count
       union {
              /* 记录使用该 inode 的硬链接总数 */
             const unsigned int i_nlink;
             unsigned int __i_nlink;
       dev_t
                            i_rdev;
                          i_size;/* 文件大小 */
      loff_t
                          i_atime;/* 最后访问时间 */
      struct timespec
                          i_mtime;/* 最后修改时间*/
       struct timespec
                            i_ctime;/* inode 最后修改时间 */
       struct timespec
       spinlock_t
                            i_lock; /* i_blocks, i_bytes, maybe i_size */
       unsigned short
                            i_bytes;
       unsigned int
                            i_blkbits;
```

```
i blocks:/*指定了按块存放的长度*/
       blkcnt t
#ifdef __NEED_I_SIZE_ORDERED
      seqcount_t
                            i_size_seqcount;
#endif
      /* Misc */
      unsigned long
                            i_state;
       struct mutex
                            i_mutex;
                          dirtied_when; /* jiffies of first dirtying */
      unsigned long
       unsigned long
                          dirtied_time_when;
       struct hlist_node
                          i_hash;
       struct list_head
                          i_io_list; /* backing dev IO list */
#ifdef CONFIG_CGROUP_WRITEBACK
       struct bdi_writeback *i_wb; /* the associated cgroup wb */
      /* foreign inode detection, see wbc_detach_inode() */
      int
                          i_wb_frn_winner;
                          i_wb_frn_avg_time;
      1116
                            i_wb_frn_history;
#endif
       struct list_head
                         i_lru;
                                        /* inode LRU list */
       struct list_head
                          i_sb_list;
       union {
            struct hlist_head
                                 i_dentry;
            struct rcu_head
                                 i_rcu;
      };
                          i_version;
      u64
       atomic_t
                           i_count;/* 访问该的进程数目inode */
       atomic_t
                            i_dio_count;
       atomic_t
                          i_writecount;
#ifdef CONFIG_IMA
      atomic_t
                          i_readcount; /* struct files open RO */
#endif
       const struct file_operations *i_fop; /* 用于操作文件中包含的数据 */
       struct file_lock_context
                                  *i_flctx;
       struct address_space i_data;
       struct list_head
                          i_devices;
       union {
              struct pipe_inode_info *i_pipe;
              struct block_device *i_bdev;
              struct cdev
                                   *i_cdev;
                                   *i_link;
      }:
      __u32
                          i_generation;
#ifdef CONFIG_FSNOTIFY
                            i_fsnotify_mask; /* all events this inode cares about */
       __u32
```

```
struct hlist_head i_fsnotify_marks;
#endif

void *i_private; /* fs or device private pointer */
};
```

1.4 inode_operations

大多数请况下,各个函数指针成员的意义可以根据其名称推断。它们与对应的系统调用和用户空间工具在名称方面非常相似。

```
struct inode_operations {
       /* lookup 根据文件系统对象的名称表示为字符串)查找其( inode 实例*/
       struct dentry * (*lookup) (struct inode *, struct dentry *, unsigned int);
       const char * (*follow_link) (struct dentry *, void **);
       int (*permission) (struct inode *, int);
       struct posix_acl * (*get_acl)(struct inode *, int);
       int (*readlink) (struct dentry *, char __user *,int);
       void (*put_link) (struct inode *, void *);
       int (*create) (struct inode *,struct dentry *, umode_t, bool);
       int (*link) (struct dentry *,struct inode *,struct dentry *);
       int (*unlink) (struct inode *,struct dentry *);
       int (*symlink) (struct inode *.struct dentry *.const char *);
       int (*mkdir) (struct inode *,struct dentry *,umode_t);
       int (*rmdir) (struct inode *,struct dentry *);
       int (*mknod) (struct inode *,struct dentry *,umode_t,dev_t);
       int (*rename) (struct inode *, struct dentry *,
                       struct inode *, struct dentry *);
       int (*rename2) (struct inode *, struct dentry *,
                       struct inode *, struct dentry *, unsigned int);
       int (*setattr) (struct dentry *, struct iattr *);
       int (*getattr) (struct vfsmount *mnt, struct dentry *, struct kstat *);
       int (*setxattr) (struct dentry *, const char *,const void *,size_t,int);
       ssize_t (*getxattr) (struct dentry *, const char *, void *, size_t);
       ssize_t (*listxattr) (struct dentry *, char *, size_t);
       int (*removexattr) (struct dentry *, const char *);
       int (*fiemap)(struct inode *, struct fiemap_extent_info *, u64 start,
                     u64 len):
       int (*update_time)(struct inode *, struct timespec *, int);
       int (*atomic_open)(struct inode *, struct dentry *,
                          struct file *, unsigned open_flag,
                          umode_t create_mode, int *opened);
       int (*tmpfile) (struct inode *, struct dentry *, umode_t);
       int (*set_acl)(struct inode *, struct posix_acl *, int);
```

```
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```

```
/* WARNING: probably going away soon, do not use! */
} ___cacheline_aligned;
```

1.5 dentry

```
struct dentry {
     /* RCU lookup touched fields */
     struct qstr d_name;
     struct inode *d_inode;
                          /* Where the name belongs to - NULL is
     /* Ref lookup also touches following */
     struct lockref d_lockref;
                          /* per-dentry lock and refcount */
     const struct dentry_operations *d_op;
     struct super_block *d_sb; /* The root of the dentry tree */
     unsigned long d_time;
                          /* used by d_revalidate */
     void *d_fsdata;
                          /* fs-specific data */
                          /* LRU list */
     struct list_head d_lru;
     struct list_head d_child;
                           /* child of parent list */
     struct list_head d_subdirs;  /* our children */
      * d_alias and d_rcu can share memory
     union {
           struct hlist_node d_alias;  /* inode alias list */
           struct rcu_head d_rcu;
     } d_u;
};
```

参考文献

[1] http://linux-mm.org/PageTableStructure.

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