

# Practical Work 3: MPI-Based RPC File Transfer

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## Abstract

This report details the implementation of a file transfer system using the Message Passing Interface (MPI). By utilizing MPI's rank-based architecture and custom data structures, we simulate Remote Procedure Call (RPC) semantics to transfer files between two processes. The system features a structured metadata handshake, synchronous acknowledgments, and chunk-based data streaming within an MPI communicator.

## 1 RPC Service Design

The system implements a single RPC-style method, `UploadFile`, which facilitates the transfer of binary or text data from a client process to a server process.

### `UploadFile(filename, filesize, data)`

To handle the communication of complex metadata, we defined a custom MPI structure that mirrors our C structure. This ensures that the metadata is correctly serialized and deserialized between processes.

```
1 typedef struct {  
2     char method[32];  
3     char filename[256];  
4     long long filesize;  
5 } Metadata;
```

Listing 1: MPI Metadata Structure

## Protocol Steps

1. **Handshake:** Client (Rank 1) sends the `Metadata` struct to the Server (Rank 0) using a custom MPI datatype.
2. **Verification:** Server receives the metadata and validates the requested "UploadFile" method.
3. **ACK:** Server sends an integer acknowledgment (code 200) to the Client.

4. **Streaming:** Client reads the file in 4096-byte chunks and sends them via `MPI_Send`.
5. **Reconstruction:** Server receives chunks, using `MPI_Get_count` to handle the exact number of bytes in the final chunk, and writes them to disk.
6. **Final Status:** Server sends a completion status (201 for success) back to the client.

## 2 System Organization

The system is organized into a single `MPI_COMM_WORLD` communicator where roles are assigned based on process rank.

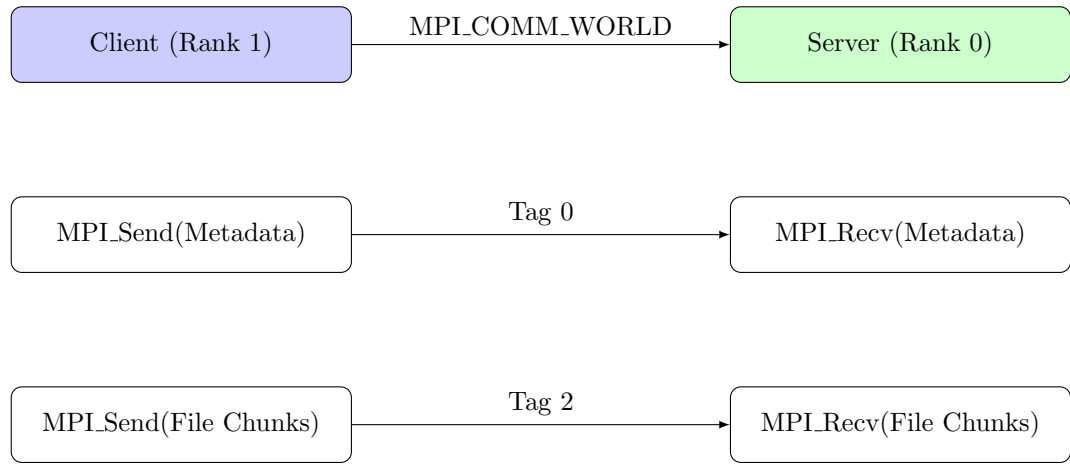


Figure 1: MPI-Based RPC File Transfer Architecture

## 3 Implementation

### 3.1 Server-Side (Rank 0)

The server utilizes `MPI_Recv` with `MPI_Status` to dynamically determine the size of the final incoming data chunk.

```

1 while (total_received < meta.filesize) {
2     int count;
3     MPI_Recv(buffer, CHUNK_SIZE, MPI_CHAR, CLIENT_RANK, 2,
4             MPI_COMM_WORLD, &status);
5     MPI_Get_count(&status, MPI_CHAR, &count); // Handle
6         last chunk size

```

```

7     fwrite(buffer, 1, count, f);
8     total_received += count;
9 }

```

Listing 2: Server Chunk Reception Loop

### 3.2 Client-Side (Rank 1)

The client is responsible for obtaining the file size via the `stat` system call and initiating the transfer.

```

1 FILE *f = fopen(filepath, "rb");
2 char buffer[CHUNK_SIZE];
3 size_t n;
4
5 while ((n = fread(buffer, 1, CHUNK_SIZE, f)) > 0) {
6     MPI_Send(buffer, n, MPI_CHAR, SERVER_RANK, 2,
7             MPI_COMM_WORLD);
8 }

```

Listing 3: Client Chunk Transmission Loop

## 4 Conclusion

By leveraging MPI’s point-to-point communication and structured datatypes, we successfully simulated an RPC file transfer system. This implementation demonstrates that while MPI is traditionally used for high-performance computing, its robust message-passing capabilities are well-suited for distributed system tasks such as reliable file transfer and remote service simulation.