

# Static Analysis

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Lecture #23 out of 24  
80 minutes

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STEVEN JOHNSON

“**Lint** is a command which examines C source programs, detecting a number of bugs and obscurities. It enforces the type rules of C more strictly than the C compilers. It may also be used to enforce a number of portability restrictions involved in moving programs between different machines and/or operating systems. Another option detects a number of wasteful, or error prone, constructions which nevertheless are, strictly speaking, legal.”

— Stephen C. Johnson. *Lint, a C Program Checker*. Bell Labs, 1977



“This is dryer lint, which is scraped out of a clothes dryer filter after it has dried a few loads. The idea of the Lint tool is to get this sort of stuff out of your code by being very pedantic about warnings and advice on possible bad code constructions.” — Quora

## Some Types of Bugs to Be Found by Static Analysis

### Unreachable Code:

```
1 | int a = 10;  
2 | if (a > 20) {  
3 |     a = a + 1;  
4 | }
```

### Uninitialized Variable:

```
1 | int x;  
2 | int y = x + 42;  
3 | print(y);
```

### Division by Zero:

```
1 | int f(int x) {  
2 |     return 42 / x;  
3 | }
```

### Integer Overflow:

```
1 | var x: u8 = 142;  
2 | x = x * 2;
```

### Endless Loop:

```
1 | int x = 5;  
2 | int y = 0;  
3 | while (x > 0) {  
4 |     y = y + x;  
5 | }
```

### Buffer Overflow:

```
1 | #include <stdio.h>  
2 | char buf[16];  
3 | fgets(buf, 1024, stdin);
```

## Inter-procedural Analysis

### Unused Global Var:

```
1 int x;  
2 int foo() {  
3     return 42;  
4 }  
5 int bar(int x) {  
6     return x + 1;  
7 }
```

### Endless Recursion:

```
1 int foo(int n) {  
2     return bar(n - 1);  
3 }  
4 int bar(int n) {  
5     return foo(n + 1);  
6 }
```

### Pointer Dereferencing:

```
1 int foo() {  
2     return *bar();  
3 }  
4 int* bar() {  
5     return 0;  
6 }
```

## Violations, Smells, Bugs

### Style Violation:

```
1 | int f
2 |   (int x)
3 | {
4 |     return 42/x;
5 | }
```

Line 2: Indentation  
Line 3: Curled bracket  
Line 4: Indentation

### Code Smell:

```
1 | int f(int x) {
2 |     return 42.0 / x;
3 | }
```

Line 2: Implicit type  
cast from float to int

### Bug:

```
1 | int f(int x) {
2 |     return 42 / x;
3 | }
```

Line 2: Division by zero



BRIAN CHESS

“Beware of any tool that says something like, ‘zero defects found, your program is, rather, now secure.’ The appropriate output is, ‘sorry, couldn’t find any more bugs.’”

— Brian Chess and Gary McGraw. Static Analysis for Security. *IEEE Security & Privacy*, 2(6):76–79, 2004

## False Negative vs. False Positive

```
1 | int f(int x) {  
2 |     return 42 / x;  
3 | }
```

**True Positive (TP):**

“Division by zero”

**False Positive (FP):**

“Integer overflow”

**True Negative (TN):**

“No buffer overflow”

**False Negative (FN):**

“No errors at all”





SUNGHUN KIM

“About 90% of warnings remain in the program or are removed during non-fix changes — likely false positive warnings.”

— Sunghun Kim and Michael D. Ernst. Which Warnings Should I Fix First? In *Proceedings of the the 6th Joint Meeting of the European Software Engineering Conference*, pages 45–54, 2007



STEVEN ARZT

“In our experiments on DroidBench examples, TASMAn reduces the number of false positives by about 80% without pruning any true positives.”

— Steven Arzt, Siegfried Rasthofer, Robert Hahn, and Eric Bodden. Using Targeted Symbolic Execution for Reducing False-Positives in Dataflow Analysis. In *Proceedings of the 4th International Workshop on State of the Art in Program Analysis*, pages 1–6, 2015



NACHIAPPAN NAGAPPAN

“Our results show that the static analysis defect density is correlated at statistically significant levels to the pre-release defect density determined by various testing activities. Further, the static analysis defect density can be used to predict the pre-release defect density with a high degree of sensitivity.”

— Nachiappan Nagappan and Thomas Ball. Static Analysis Tools as Early Indicators of Pre-release Defect Density. In *Proceedings of the 27th International Conference on Software Engineering*, pages 580–586, 2005

## My Favorite Static Analyzers

- Java: SpotBugs, Checkstyle, PMD, and Qulice for Java
- C++: Clang-Tidy
- Rust: clippy

## Some Static Analysis Mechanisms

- Data Flow Analysis
- Symbolic Execution
- Model Checking
- Taint Analysis

You may want to watch my “[Practical Program Analysis](#)” course.

## For some tools, you have to pay:

- Coverity by Synopsys (US)
- Klockwork by Perforce (US)
- Fortify by Micro Focus (UK)
- Checkmarx (US)
- Veracode (US)
- Snyk (US)
- PVS-Studio (Russia)

Usually, up to \$3,000 per developer per year.

## Why do JavaScript developers use linters?

- Prevent Errors
- Augment Test Suites
- Avoid Ambiguous and Complex Code
- Maintain Code Consistency
- Faster Code Review
- Spare Developers' Feelings
- Save Discussion Time
- Learn About JavaScript

Source: Kristín Fjóra Tómasdóttir, Mauricio Aniche, and Arie Van Deursen. Why and How JavaScript Developers Use Linters. In *Proceedings of the 32nd International Conference on Automated Software Engineering (ASE)*, pages 578–589. IEEE, 2017



KRISTÍN FJÓLA TÓMASDÓTTIR

“Every single interview participant mentioned that one of the reasons why they use a linter is to maintain code consistency.”

— Kristín Fjóra Tómasdóttir, Mauricio Aniche, and Arie Van Deursen. The Adoption of JavaScript Linters in Practice: A Case Study on ESLint. *IEEE Transactions on Software Engineering*, 46(8):863–891, 2018



| Category             | Description   | Available rules |
|----------------------|---|-----------------|
| Possible Errors      | Possible syntax or logic errors in JavaScript code        | 31              |
| Best Practices       | Better ways of doing things to avoid various problems     | 69              |
| Strict Mode          | Strict mode directives                                    | 1               |
| Variables            | Rules that relate to variable declarations                | 12              |
| Node.js and CommonJS | For code running in Node.js, or in browsers with CommonJS | 10              |
| Stylistic Issues     | Stylistic guidelines where rules can be subjective        | 81              |
| ECMAScript 6         | Rules for new features of ES6 (ES2015)                    | 32              |
| <b>Total</b>         |   | <b>236</b>      |

TABLE 1: ESLint rule categories with ordering and descriptions from the ESLint documentation [28]

Source: Kristín Fjóra Tómasdóttir, Mauricio Aniche, and Arie Van Deursen. The Adoption of JavaScript Linters in Practice: A Case Study on ESLint. *IEEE Transactions on Software Engineering*, 46(8):863–891, 2018

## SARIF

```
{
  "results": [
    {
      "ruleId": "CA2101",
      "message": {
        "text": "Variable '{0}' is uninitialized.",
        "arguments": [ "pBuffer" ]
      }
    }
  ]
}
```

“This document defines a standard format for the output of static analysis tools.”

Source: OASIS. Static Analysis Results Interchange Format (SARIF) Version 2.1.0 Plus Errata 01.  
<https://docs.oasis-open.org/sarif/sarif/v2.1.0/sarif-v2.1.0.html>, 2023. [Online; accessed 08-03-2024]



FLORIAN OBERMÜLLER

“We introduce the concept of code perfumes as the counterpart to code smells, indicating the correct application of programming practices considered to be good. Using a catalogue of 25 code perfumes for, we empirically demonstrate that these represent frequent practices in, and we find that better programs indeed contain more code perfumes.”

— Florian Obermüller, Lena Bloch, Luisa Greifenstein, Ute Heuer, and Gordon Fraser. Code Perfumes: Reporting Good Code to Encourage Learners. In *Proceedings of the 16th Workshop in Primary and Secondary Computing Education*, pages 1–10, 2021

# References

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- Kristín Fjóra Tómasdóttir, Mauricio Aniche, and Arie Van Deursen. The Adoption of JavaScript Linters in Practice: A Case Study on ESLint. *IEEE Transactions on Software Engineering*, 46(8): 863–891, 2018.