

# data\_structures(&algorithms, lecture07)

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# Plan

## **Queue ADT**

Introduction, queue operation, queue application

## **Array Implementation**

How to implement queue with Array, Circular Queue with array

## **Linked List Implementation**

How to implement queue with Linked List

## **Priority Queue**

What is Priority Queue, implement by Array and by Heap

0

1

0

2

0

3

0

4

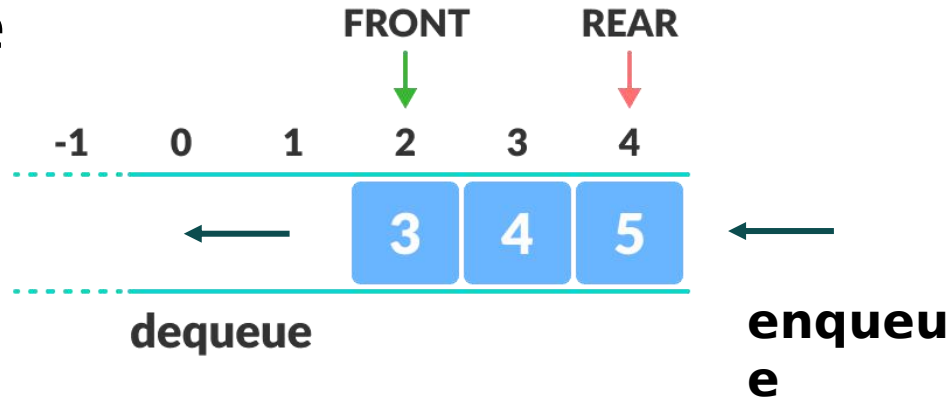


# Queue ADT

Introduce Queue data structure, some operations

# What Is A Queue

- ✓ A linear data structure that can be accessed in both begins and ends for storing and retrieving data
- ✓ A queue is a First In, First Out (FIFO) data structure



# Operations On Queue

- v clear: clear the queue
- v is empty: check if a queue is empty
- v enqueue: put element on the rear of the queue
- v dequeue: take the element on the top out of the queue
- v front: get first element without removing it
- v rear: get last element without removing it
- v size: get size of the queue

# Applications Of Queue

- v Any kind of waiting list
- v Operating System:
  - v CPU scheduling
  - v Disk scheduling
  - v Shared resources between processes
- v Traverse in tree / graph data structures
- v Turned-base games
- v Alt-Tab in Windows
- v ...



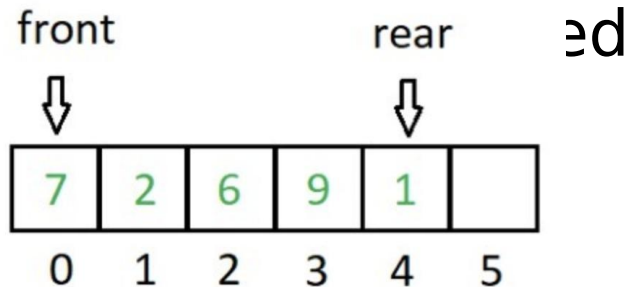
# Array Implementation

Implement Queue by Array, Circular Queue



# Queue Interface By Array

- Need global variables
- 2 indicators run forward
- Queue size is fixed
- Number of



```
#define QUEUE_SIZE 100

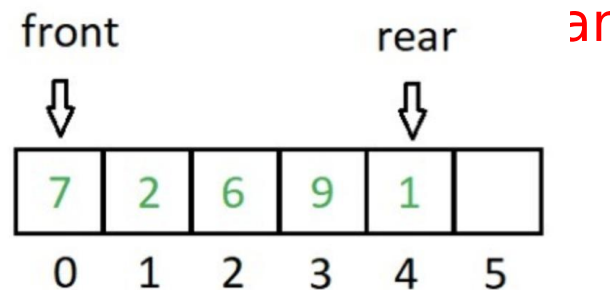
extern int front;
extern int rear;
extern int queue[];

void clear_queue(void);
int is_empty(void);
int is_full(void);
void enqueue(const int n);
int dequeue(void);
int size(void);
void overflow(char* msg);
```



# Queue Implementation By Array

- ✓ is\_empty: when **front** = -1
- ✓ is\_full: when **rear** runs to the end of array
- ✓ enqueue: need to check full queue before, then increase **rear** and assign new item to new **rear** position
- ✓ dequeue: need to check empty queue before, then increase **front** and return the item at old **front**
- ✓ size: distance between **front** and **rear**



# Queue Implementation By Array

- ✓ Number of enqueue is fixed

- ✓ => When dequeue, if queue has only one item, clear queue to reset front & rear to -1
- ✓ => Implement circular queue

- ✓ Size is fixed

- ✓ => Using dynamic array and grow size if queue is full

# Circular Queue By Array

✓ Main idea: front & rear run forward but go back (if possible) to 0 index.

✓ is\_empty: front = -1

✓ is\_full:

✓ front = 0, rear = max size - 1

✓ front = rear + 1

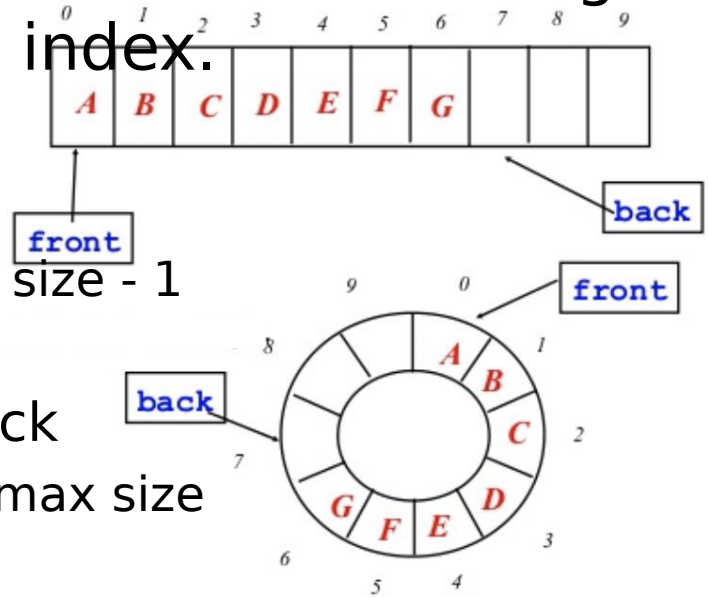
✓ Using modulo to go back

✓ front = (front + 1) % max size

✓ real size:

✓ if (front < rear): rear - front + 1

✓ else: max size - (front - rear - 1)



# Example: Queue of messages

- ✓ Simulation of sending message.

Commands:

- ✓ **+**message: enqueue message
  - ✓ When the queue is full, all messages will be sent
- ✓ **-**: dequeue a message to send
- ✓ **a**: Sending all messages by dequeue one by one
- ✓ **q**: Quit simulation, all messages will be sent
- ✓ Using Circular Queue will support unlimited enqueue operation, still having full condition

# Example: Queue of messages

v Queue: array of char pointers in heap memory

v Enqueue

Check if queue is full

Else

If queue is empty front = 0

Increase rear (go back if needed)

Add message to rear position

v Dequeue

Check if message is empty

Else

Get message at front

If there is one message, clear queue

Else increase front (go back if needed)

# Example: Queue of messages

## v Running simulation

### Message sending simulation

```
>> +hello
Saved hello
>> +world
Saved world
>> +stack and queue
Saved stack and queue
>> -
Sending hello ...
>> +good bye
Saved good bye
>> a
Sending world ...
Sending stack and queue ...
Sending good bye ...
>> q
No message to send!
Program ended with exit code: 0
```



# Linked List Implementation

Implement Queue by Linked List



# LinkedList vs Queue

- Queue is Linked List with limited operations and/or other interfaces

Linked List	Queue
Add to end	Enqueue
Remove from head	Dequeue
Clear list	Clear queue
Is empty	Is empty
Get size	Get size
	To array
	Front / Rear

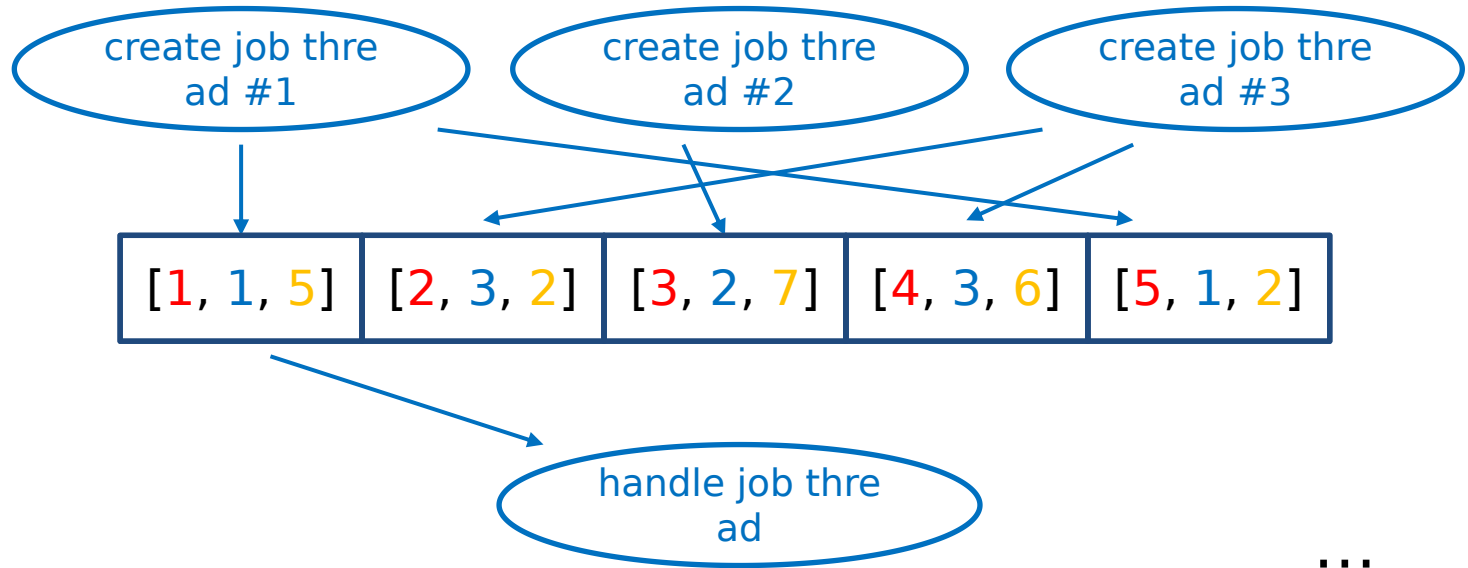
# Example: Queue of Jobs

- ✓ Simulation a queue of jobs
  - ✓ A job has an **id**, **agent id** (who creates this job) and **duration** (in seconds)
  - ✓ Several threads will create jobs in parallel and store them in a job queue
  - ✓ A thread will get job from that queue to execute it
  - ✓ Shared resources: queue of jobs and number of jobs

[1, 1, 5]	[2, 3, 2]	[3, 2, 7]	[4, 3, 6]	[5, 1, 2]	...
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# Example: Queue of Jobs

## Simulation a queue of jobs



...

# Example: Queue of Jobs

## ✓ Implement queue of jobs

✓ Job data

Job node in LL

```
typedef struct
{
    int job_id;
    int agent_id;
    int duration; // in second
} job_data;
```

```
typedef struct str_job job;

struct str_job
{
    job_data data;
    job* next;
};
```

# Example: Queue of Jobs

## v Implement queue of jobs

### v Enqueue operation

```
void enqueue(job** jobs_queue, job_data ajob)
{
    job* j = create_job(ajob);

    if (is_empty(*jobs_queue)) *jobs_queue = j;
    else
    {
        job* rear = *jobs_queue;
        while (rear->next != NULL) rear = rear->next;
        rear->next = j;
    }
}
```

# Example: Queue of Jobs

## v Implement queue of jobs

### v Dequeue operation

```
job_data dequeue(job** jobs_queue)
{
    if (is_empty(*jobs_queue))
    {
        printf("Jobs queue is empty!");
        exit(1);
    }
    else
    {
        job* j = *jobs_queue;
        *jobs_queue = j->next;
        job_data data = j->data;
        free(j);
        return data;
    }
}
```

# Example: Queue of Jobs

## v Implement parallel threads

- v pthread\_create: create a new thread to run in parallel

```
pthread_create(&tid, NULL, handle_jobs, NULL);
```

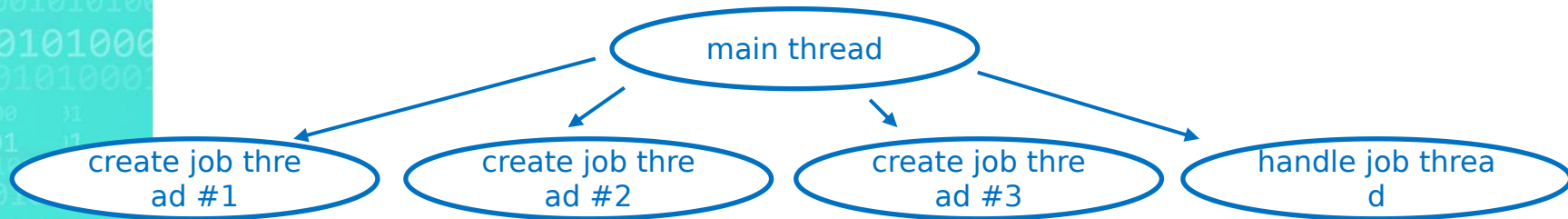
thread structure

thread attribute

callback function

arguments

- v pthread\_exit: terminate a thread
- v sleep: to make current thread idle for a while

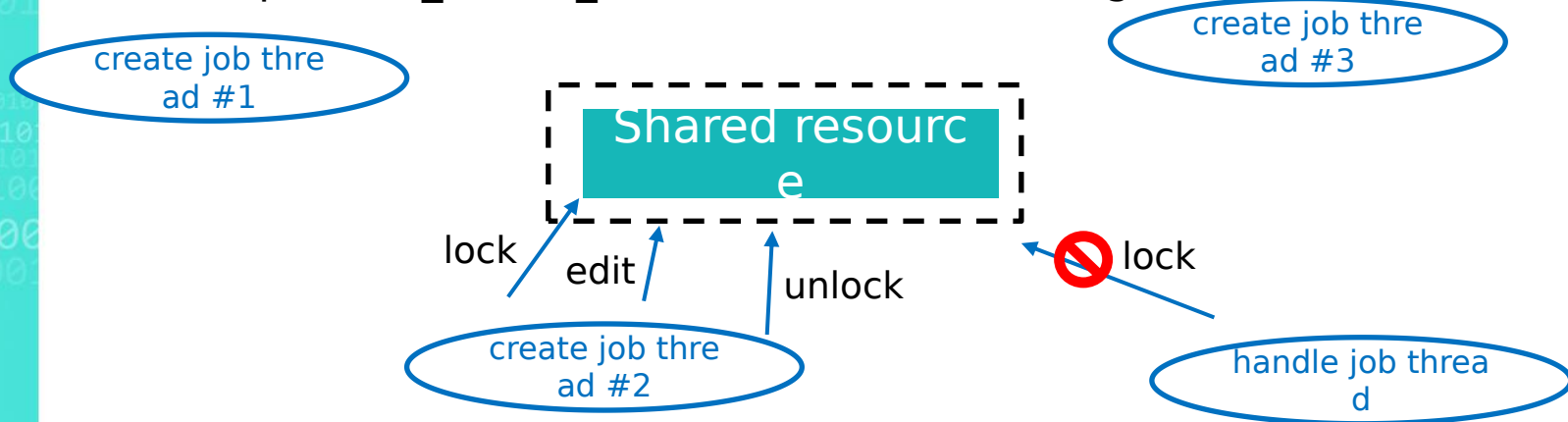




# Example: Queue of Jobs

## ✓ Shared resources

- ✓ `pthread_mutex_t locker`: declare a locker to lock shared resource
- ✓ `pthread_mutex_init(&locker, NULL)`: initialize the locker
- ✓ `pthread_mutex_lock(&locker)`: lock global resources
- ✓ `pthread_mutex_unlock(&locker)`: unlock global resources



# Example: Queue of Jobs

## v Putting all together

### v main thread

```
// init locker
pthread_mutex_init(&locker, NULL);

// create threads to generate jobs in parallel
for (int agent_id = 1; agent_id <= NAGENTS; agent_id++)
    pthread_create(&tid, NULL, generate_jobs, (void*)agent_id);
// create thread to handle generated jobs
pthread_create(&tid, NULL, handle_jobs, NULL);
```

# Example: Queue of Jobs

## v Putting all together

### v create job thread

```
int agent_id = (int) threadid;
for (int i = 1; i <= NJOBS; i++)
{
    pthread_mutex_lock(&locker); // lock the locker

    // increase number of jobs
    // create job data based on agent id, number of jobs
    // enqueue job data

    pthread_mutex_unlock(&locker); // unlock the locker

    sleep(rand() % WAIT_TIME);
}
pthread_exit(NULL);
```

# Example: Queue of Jobs

- ✓ Putting all together
  - ✓ handle job thread

```
Endless loop
  Lock the locker
  If jobs queue is empty and no more jobs in future
    Quit handle thread
  Else If queue is not empty
    Handle a job
  Else
    Wait for a job
```

# Example: Queue of Jobs

## v Putting all together

```
Agent id 1 create job 1
Number of waiting jobs: 1
Handle job 1 of agent 1 in 4s
Agent id 2 create job 2
Agent id 3 create job 3
Agent id 2 create job 4
Agent id 2 create job 5
Agent id 1 create job 6
Agent id 3 create job 7
Number of waiting jobs: 6
Handle job 2 of agent 2 in 0s
Number of waiting jobs: 5
Handle job 3 of agent 3 in 0s
Number of waiting jobs: 4
Handle job 4 of agent 2 in 4s
```

The background features two large, overlapping teal geometric shapes. The top shape is a parallelogram pointing towards the top right, and the bottom shape is a parallelogram pointing towards the bottom left. Both shapes are filled with a pattern of binary code (0s and 1s) in a lighter shade of teal. The overall background is a light teal color.

# Priority Queue

Introduce Priority Queue and how to implement it

# What is Priority Queue?

- ✓ Same as a normal queue but ...
  - ✓ Each element has a priority
  - ✓ Elements with higher priority will be dequeued before elements with lower priority
  - ✓ Elements with same priority will be dequeued in the order of enqueueing

**“All animals are equal, but some animals are more equal than others”**

*George Orwell – Animal Farm*



# Example of Priority Queue

- ✓ Enqueue following pairs (key, priority)
  - ✓ (1, 1), (2, 1), (3, 2), (4, 1), (5, 2), (6, 3), (7, 1)
- ✓ Actual queue in building:
  - ✓ (1, 1)
  - ✓ (1, 1), (2, 1)
  - ✓ (3, 2), (1, 1), (2, 1)
  - ✓ (3, 2), (1, 1), (2, 1), (4, 1)
  - ✓ (3, 2), (5, 2), (1, 1), (2, 1), (4, 1)
  - ✓ (6, 3), (3, 2), (5, 2), (1, 1), (2, 1), (4, 1)
  - ✓ (6, 3), (3, 2), (5, 2), (1, 1), (2, 1), (4, 1), (7, 1)

# Priority Queue by Array

## ✓ Enqueue algorithm:

- ✓ Insert new element at the rear
- ✓ If priority of new element is greater than the one ahead then move it up until it's less than or equal

✓ (3, 2), (1, 1), (2, 1), (4, 1)

✓ (3, 2), (1, 1), (2, 1), (4, 1), (5, 2)

✓ (3, 2), (1, 1), (2, 1), (5, 2), (4, 1)

✓ (3, 2), (1, 1), (5, 2), (2, 1), (4, 1)

✓ (3, 2), (5, 2), (1, 1), (2, 1), (4, 1)

## ✓ Complexity:

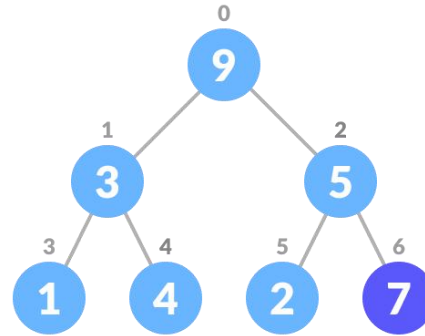
✓ Enqueue:  $O(n)$

✓ Dequeue:  $O(1)$

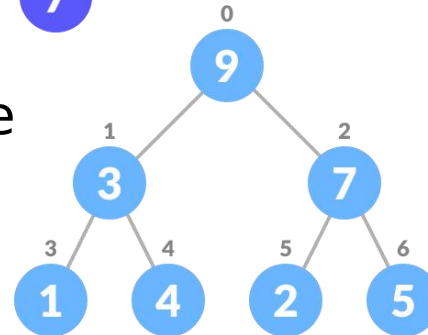
# Priority Queue by Heap

## ✓ Enqueue algorithm:

- ✓ Insert new element at the end of the heap



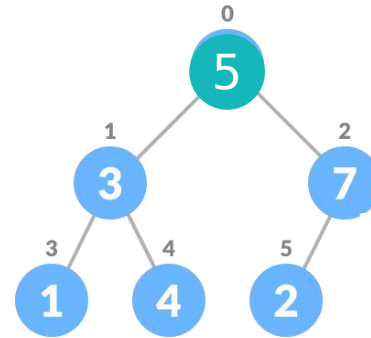
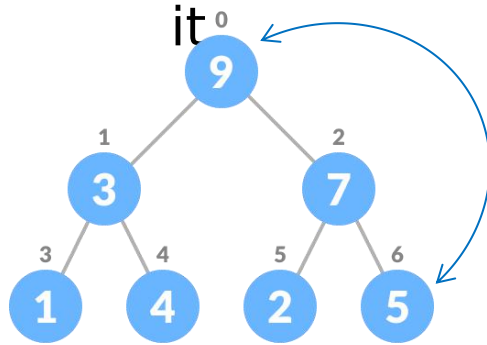
- ✓ Heapify the whole tree



# Priority Queue by Heap

## Dequeue algorithm:

- Swap the root with the last element then remove

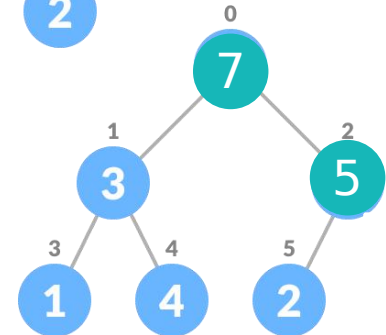


- Heapify the whole tree

## Complexity:

- Enqueue:  $O(\log n)$

- Dequeue:  $O(\log n)$



# Priority Queue by Heap

## ✓ Implement Priority Queue by Heap

```
typedef struct
```

```
{
```

```
    int key;
```

```
    int priority;
```

```
    int order;
```

```
} element;
```

```
extern int qsize;
```

```
void swap(element *a, element *b);
```

```
void print_queue(element* queue);
```

```
void heapify(element* queue, int i);
```

```
void enqueue(element* queue, element e);
```

```
element dequeue(element* queue);
```

```
int greater(element a, element b);
```

```
int is_empty(void);
```

Fix unstable characteristics of heap

# Priority Queue by Heap

## v Implement PriorityQueue by Heap

```
void heapify(element* queue, int i)
{
    int left = 2 * i + 1;
    int right = 2 * i + 2;
    int imax = i;

    if (left < qsize && greater(queue[left], queue[imax]))
        imax = left;
    if (right < qsize && greater(queue[right], queue[imax]))
        imax = right;

    if (imax != i)
    {
        swap(&queue[i], &queue[imax]);
        heapify(queue, imax);
    }
}
```

# Priority Queue by Heap

## v Implement Priority Queue by Heap

```
int greater(element a, element b)
{
```

```
    // First compare by priority
```

```
    // Then compare by order if same priority
```

```
}
```

```
void enqueue(element* queue, element e)
{
```

```
    // Add e to the end of the queue
```

```
    // Heapify the whole tree bottom-up
```

```
}
```

```
element dequeue(element* queue)
```

```
{
```

```
    // Get first element, save into e
```

```
    // Swap first element and last element
```

```
    // Heapify the new tree (exclude last element)
```

```
    // Return e
```

```
}
```