## Magic States Distillation Using $\Delta$ -Toric (good qLDPC?).

## David Ponarovsky

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Let  $|f\rangle$  be a codeword in  $C_X$ , and let  $X_g$  be the indicator that equals 1 if f has support on  $X_g$ , and 0 otherwise. Observes that applying  $T^{\otimes}$  on  $|f\rangle$  yilds the state:

$$\begin{split} T^{\otimes n} \left| f \right\rangle &= T^{\otimes n} \left| \sum_{g} X_g g \right\rangle = \exp \left( i \pi / 4 \sum_{g} X_g |g| - 2 \cdot i \pi / 4 \sum_{g,h} X_g X_h |g \cdot h| \right. \\ &+ 4 \cdot i \pi / 4 \sum_{g,h} X_g X_h X_l |g \cdot h \cdot l| - 8 \cdot i \pi / 4 \cdot \text{ integers } \right) \left| f \right\rangle \\ &= \exp \left( i \pi / 4 \sum_{g} X_g |g| - 2 \cdot \pi / 4 \sum_{g,h} X_g X_h |g \cdot h| + 4 \cdot i \pi / 4 \sum_{g,h} X_g X_h X_l |g \cdot h \cdot l| \right) \left| f \right\rangle \end{split}$$

Now assume that the code  $C_X$  is the quantum Tanner code, denote by G,A,B the group and the two generator sets that are used for constructing the square complex. In addition, let us assume the existence of  $d \in G$  such that d is non-identity and commutes with any element in  $A \cup B$ . Then, observe that multiplying by d preserves adjacency on the complex. Namely, if  $\{u,v\} \in E$  then also  $\{du,dv\} \in E$ .

Consider  $|f\rangle$  such that if  $X_g$  is not zero, and g is associated with a local codeword  $c \in C_A \otimes C_B$  on vertex v, then the generator associated with the local codeword c on vertex  $d \cdot v$  also supports f, denoted by g'. Thus, the exponent above becomes:

$$\begin{split} &= \exp \left( i \pi / 4 \sum_{g} X_{g} |g| - 2 \cdot \pi / 4 \sum_{g,h \in G/a} X_{g} X_{h} |g \cdot h| + X_{g'} X_{h'} |g \cdot h| \right. \\ &+ 4 \cdot i \pi / 4 \sum_{g,h \in G/a} X_{g} X_{h} X_{l} |g \cdot h \cdot l| + X_{g'} X_{h'} X_{l'} |g \cdot h \cdot l| \right) |f\rangle \\ &= \exp \left( i \pi / 4 \sum_{g} X_{g} |g| - 2 \cdot 2 \cdot \pi / 4 \sum_{g,h \in G/a} X_{g} X_{h} |g \cdot h| + 2 \cdot 4 \cdot i \pi / 4 \sum_{g,h \in G/a} X_{g} X_{h} X_{l} |g \cdot h \cdot l| \right) |f\rangle \\ &= \exp \left( i \pi / 4 \sum_{g} X_{g} |g| - i \pi \sum_{g,h \in G/a} X_{g} X_{h} |g \cdot h| \right) |f\rangle \end{split}$$

Claim 0.1. The gate 
$$|f\rangle \mapsto \exp\left(-i\pi \sum_{g,h\in G/a} X_g X_h |g\cdot h|\right) |f\rangle$$
 is in the Clifford.

*Proof.* Just decode f and apply  $\mathbb{C}\mathbf{Z}$  between any pair of qubits corresponding to the generators g,h such that  $g \cap h = 1$ . Then encode the state again. Observes that  $\mathbb{C}\mathbf{Z}$  is a Clifford gate, and by the fact that the code is a CSS code then the decoder and the encoder are both in the Clifford.

Let's denote the circuit defined in Claim 0.1 by  $\Lambda$ . So we have that:

$$\Lambda^{\dagger} \exp\left(i\pi/4\sum_{g} X_{g}|g| - i\pi\sum_{g,h \in G/a} X_{g}X_{h}|g \cdot h|\right)|f\rangle$$
$$= \exp\left(i\pi/4\sum_{g} X_{g}|g|\right)|f\rangle$$

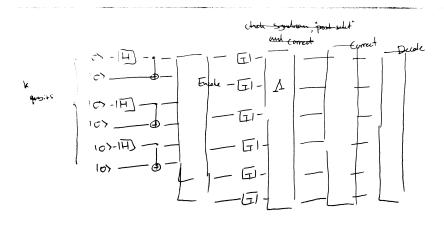


Figure 1: Quantum Circuit for distillation.