State Synthesis Using PRS.

David Ponarovsky

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Abstract

We studies the complexity of synthesis quantum states using PRS, our reasch continues the work by [Ira+22], [Ros23], [RY21], [MY23], [Del+23].

Definition 0.1 (The keeping setting). Let $R^A \otimes R^B$ be a general two registers domain. We define the **keeping setting** to let one construct quntum/classical circuits¹ $G: R^A \otimes R^B \to R^A \otimes R^B$ such that it is gurnted that the register R^B cann't be acssed after the computation.

Claim 0.1. Let G be a PRS generator, than under the the keeping setting one can assume that G takes as input two register, the first contains n ancille qubits initilizated to $|0\rangle$ and the seconed contain a classic string initilized to be the seed k.

Proof. Given a PRS $G: R^A \to R^A$ define $\tilde{G}: R^A \otimes R^B \to R^A \otimes R^B$ as follow, first \tilde{G} copy the calsical state in R^B (the k-length seed) to R^A and then appaly G on R^A , Hence on sampled seed $k \in R^B$ results the output $|\psi_k\rangle \otimes |k\rangle$. Under the keeping setting any polynomial distingushier-canidate D has acsses only for $|\psi_k\rangle$, So if D distinguish between the distrubition generated by \tilde{G} and the Haar measure then it also distingush between G and Haar measure.

Claim 0.2. Let $G: |0\rangle^n \otimes \mathbb{F}_2^k \to \{|\psi_k\rangle\}_{k \in \mathcal{K}}$ be a PRS generator uses n- ancilles and k classic bits. Then for any unitary $V: \mathcal{H}_n \to \mathcal{H}_n$ it holds that $(V \otimes I^{\otimes k})G$ is also a PRS.

Proof.

Claim 0.3 (Levis Lemma for PRS). Let $f: \mathcal{H} \to R$ be a **BQP**-computible function on the n-qubits hilbert space, and let $g: (0,1) \to \mathbb{R}$ a function such that:

$$\mathbf{Pr}_{|\psi\rangle\sim U}\left[f\left(|\psi\rangle\right)>\varepsilon\right]< g(\varepsilon)$$

Then, a similar inequality also holds for states sampled by the PRS, when the probability for the measure f-value grater than ε is bounded by $g(2\varepsilon)$. Namely,

$$\mathbf{Pr}_{|\psi\rangle\sim|\psi_{k}\rangle}\left[f\left(|\psi\rangle\right)>\varepsilon\right]< g(2\varepsilon)$$

In praticular, Levi's lemma has a version that capture consetration of states sampled by PRS generator, states the following: Assume there exsists K such that for any $|\psi\rangle$, $|\phi\rangle \in \mathcal{S}(\mathbb{C}^d)$ $|f(|\psi\rangle) - |f(|\phi\rangle)| < K||\psi\rangle - |\phi\rangle|$. Then there exsists a universal constant C > 0 such:

$$\mathbf{Pr}_{|\psi\rangle\sim|\psi_{k}\rangle}\left[\left|f\left(\left|\psi\right\rangle\right) - \mathbf{E}_{\left|\phi\right\rangle\sim U}\left[f\left(\left|\phi\right\rangle\right)\right]\right| > \varepsilon\right] < exp\left(-\frac{Cd}{K^{2}}4\varepsilon^{2}\right)$$

Proof.

Claim 0.4. Probablisite counting argument and ε -net over PRS.

Claim 0.5. exsistness of poly(n) gates G_1, G_2 .. such that, any G_i has a polynomial depth, $\langle p(G_i)|\tau\rangle > a$ and $\langle \tau^{\perp}|p(G_j)\rangle \langle p(G_i)|\tau^{\perp}\rangle < b$ for any $i \neq j$.

$$\square$$

Claim 0.6. bla bla bla

Definition 0.2. ε -bised test 2-degree for testing RPU/RPS. $f(\langle \theta | G_s | \theta \rangle) = 1$.

¹On which we think as a canidate for PRS/PRF/PRG generator.

References

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