

State Synthesis Using PRS.

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Abstract

We studies the complexity of synthesis quantum states using PRS, our reasch continues the work by [Ira+22], [Ros23], [RY21], [MY23], [Del+23].

1 Pseudorandomness.

Definition 1.1 (Pseudorandom Quantum states). *Let \mathcal{H}, \mathcal{K} be the Hilbert and the key spaces, their diminsions depeand on a security paramter n . A state famliiy $\{|\psi_k\rangle\}_{k \in \mathcal{K}}$ is a pseudiorandom, if the following hold:*

1. *Efficient generation. There is a polynomial-time quantum algorithm G that generates state $|\psi_k\rangle$ on input k .*
2. *Pseudorandomness. Any polynomially many copies of $|\phi_k\rangle$ with the same random $k \in K$ is computationally indistinguishable from the same number of copies of the Haar random state.*

Definition 1.2 (Pseudorandom Unitary Operators). *A famliiy of unitary operators $\{U_k \in U(\mathcal{H})\}_{k \in \mathcal{K}}$ is pseudorandom, if two conditions hold:*

1. *Efficient computation. There is an efficient quantum algorithm Q , such that for all k and any $|\psi\rangle \in \mathcal{H}$ $Q(k, |\psi\rangle) = U_k |\psi\rangle$.*
2. *Pseudorandomness. The uniform random distribution on U_k is computationally in distinguishable from a Haar random unitary operator.*

Definition 1.3 (The keeping setting). *Let $R^A \otimes R^B$ be a general two registers domain. We define the **keeping setting** to let one construct quntum/classical circuits¹ $G : R^A \otimes R^B \rightarrow R^A \otimes R^B$ such that it is gurnted that the register R^B can't be accsed after the computation.*

Claim 1.1. *Let G be a PRS generator, than under the the keeping setting one can assume that G takes as input two register, the first contains n ancille qubits initiliaized to $|0\rangle$ and the seconed contain a classic string initiliezied to be the seed k .*

Proof. Given a PRS $G : R^A \rightarrow R^A$ define $\tilde{G} : R^A \otimes R^B \rightarrow R^A \otimes R^B$ as follow, first \tilde{G} copy the calscial state in R^B (the k -length seed) to R^A and then appaly G on R^A , Hence on sampled seed $k \in R^B$ results the output $|\psi_k\rangle \otimes |k\rangle$. Under the keeping setting any polynomial distinguishier-canidate D has accses only for $|\psi_k\rangle$, So if D distinguish between the distrubition generated by \tilde{G} and the Haar measure then it also distinguish between G and Haar measure. \square

Claim 1.2. *Let $G : |0\rangle^n \otimes \mathbb{F}_2^k \rightarrow \{|\psi_k\rangle\}_{k \in \mathcal{K}}$ be a PRS generator uses n - ancilles and k classicl bits. Then for any unitery $V : \mathcal{H}_n \rightarrow \mathcal{H}_n$ it holds that $(V \otimes I^{\otimes k})G$ is also a PRS.*

Proof. \square

¹On which we think as a canidate for PRS/PRF/PRG generator.

Claim 1.3 (Levi's Lemma for PRS). *Let $f : \mathcal{H} \rightarrow \mathbb{R}$ be a **BQP**-computable function on the n -qubits Hilbert space, and let $g : (0, 1) \rightarrow \mathbb{R}$ be a function such that:*

$$\Pr_{|\psi\rangle \sim U} [f(|\psi\rangle) > \varepsilon] < g(\varepsilon)$$

Then, a similar inequality also holds for states sampled by the PRS, when the probability for the measure f -value greater than ε is bounded by $g(2\varepsilon)$. Namely,

$$\Pr_{|\psi\rangle \sim |\psi_k\rangle} [f(|\psi\rangle) > \varepsilon] < g(2\varepsilon)$$

In particular, Levi's lemma has a version that captures concentration of states sampled by PRS generator, states the following: Assume there exists K such that for any $|\psi\rangle, |\phi\rangle \in \mathcal{S}(\mathbb{C}^d)$ $|f(|\psi\rangle) - f(|\phi\rangle)| < K||\psi\rangle - |\phi\rangle|$. Then there exists a universal constant $C > 0$ such:

$$\Pr_{|\psi\rangle \sim |\psi_k\rangle} [|f(|\psi\rangle) - \mathbf{E}_{|\phi\rangle \sim U} [f(|\phi\rangle)]| > \varepsilon] < \exp\left(-\frac{Cd}{K^2}4\varepsilon^2\right)$$

Proof. □

Claim 1.4. *Probabilistic counting argument and ε -net over PRS.*

Claim 1.5. *existence of $\text{poly}(n)$ gates G_1, G_2, \dots such that, any G_i has a polynomial depth, $\langle p(G_i) | \tau \rangle > a$ and $\langle \tau^\perp | p(G_j) \rangle \langle p(G_i) | \tau^\perp \rangle < b$ for any $i \neq j$.*

Proof. □

Claim 1.6. *bla bla bla*

Definition 1.4. ε -biased test 2-degree for testing RPU/RPS. $f(\langle x_j | G_s | \theta \rangle) = 1$ For example ask if $\langle \psi_j, \tau^\perp \rangle \langle \tau^\perp | \psi_j \rangle$ what I can say about that quantity as polynomial?

2 What We Need for Synthesis.

Definition 2.1 (Pseudorandom Unitary for Synthesis). *A family of unitary operators $\{U_k \in U(\mathcal{H})\}_{k \in \mathcal{K}}$ is pseudorandom for synthesis, if two conditions hold:*

1. *Efficient computation.* There is an efficient quantum algorithm Q , such that for all k and any $|\psi\rangle \in \mathcal{H}$ $Q(k, |\psi\rangle) = U_k |\psi\rangle$.
2. *Pseudorandomness for synthesis.* Given a state $|\tau\rangle$ and polynomial number of samples U_1, U_2, \dots, U_m . Then:

$$(a) \quad |\langle \Phi(\tau, U_k) | U_k \tau \rangle|^2 > a$$

$$(b) \quad |\langle \Phi(\tau, U_k) | U_k \tau^\perp \rangle \langle \tau^\perp | U_j^\dagger | \Phi(\tau, U_j) \rangle|^2 < b$$

The uniform random distribution on U_k is computationally indistinguishable from a Haar random unitary operator.

What about, Assume that U is a quantum circuit such that $\log n$ qubits are initialized to some state and instead ancilla, we have noisy ancilla, can we show that circuit is equivalent to $\log n$ circuit? That will enable us to prove a quantum version for Nisan Wigderson PRG ($\text{BPP} = \text{P}$).

Problem. Let U be a quantum circuit which get $\log n$ stable qubits and $\text{poly}(n)$ more random qubits obtained from the random Haar measure, can we simulate the circuit in $\log n$ time?

approximate the absolute value function, For example, you can consider the binomial expansion of $\sqrt{1-y}$ on $[0, 1]$. Namely, setting $y = 1 - x^2$, we have $|x| = \sqrt{1-y} = \sum_{m=0}^{\infty} \binom{1/2}{m} (-y)^m$, $x \in [-1, 1]$. That will allow me to bound the k -design.

Denote by $q_d(x)$ the d -order approximation of $|x|$, Namely

$$q_d(x) = \sum_{m=0}^d \binom{1/2}{m} (-1)^m (1-x^2)^m$$

and as the series converges to any $x \in (-1, 1)$ we have that $|x| = q_d(x) + O(\binom{1/2}{d} (1-x^2)^d)$ which by the fact that $1-x^2 \in (-1, 1)$ can be simplified to $|x| = q_d(x) + O(\binom{1/2}{d}) = q_d(x) + O(1/d^{1+1/2})$.

$$\begin{aligned} \mathbf{E}_{U \sim D} [(\langle \Phi(\tau, U) | \text{Re } U\tau \rangle)^2] &= \mathbf{E}_{U \sim D} \left[\frac{1}{2^{n/2}} \sum_x (-1)^{\text{sign}(\text{Re} \langle x | U\tau \rangle)} \text{Re} \langle x | x \rangle \langle x | U\tau \rangle \right] \\ &= \mathbf{E}_{U \sim D} \left[\frac{1}{2^{n/2}} \sum_x |\text{Re} \langle x | U\tau \rangle| \right] \\ &= \mathbf{E}_{U \sim D} \left[\sum_x |\text{Re} \langle x | U\tau \rangle| / 2^{n/2} \right] \\ &\geq \mathbf{E}_{U \sim D} \left[\sum_x q_d \left(|\text{Im} \langle x | U\tau \rangle| / 2^{n/2} \right) - \binom{1/2}{d} \left(\frac{|\text{Im} \langle x | U\tau \rangle|}{2^{n/2}} \right)^d \right] \\ &\geq \mathbf{E}_{U \sim \text{Haar}} \left[\sum_x q_d \left(|\text{Im} \langle x | U\tau \rangle| / 2^{n/2} \right) - \binom{1/2}{d} \left(\frac{|\text{Im} \langle x | U\tau \rangle|}{2^{n/2}} \right)^d \right] - \delta \cdot 2^n \\ &\geq \mathbf{E}_{U \sim \text{Haar}} \left[\sum_x |\text{Re} \langle x | U\tau \rangle| / 2^{n/2} - 2 \cdot \binom{1/2}{d} \left(\frac{|\text{Im} \langle x | U\tau \rangle|}{2^{n/2}} \right)^d \right] - \delta \cdot 2^n \\ &\sim \mathbf{E}_{U \sim \text{Haar}} \left[\sum_x |\text{Re} \langle x | U\tau \rangle| / 2^{n/2} \right] - \delta \cdot 2^n \\ \mathbf{E}_{U, U_2 \sim D} [\langle \Phi(\tau, U) | U\tau^\perp \rangle \langle \tau^\perp U_2^\dagger | \Phi(\tau, U_2) \rangle] &= \end{aligned}$$

Claim 2.1. fix a state $|\tau\rangle$. Let U be a unitary sampled from k -design distribution D and denote by $|s\rangle$ the vector which U sends $|\tau\rangle$ to. Now, observe that U can be written as $U = |s\rangle \langle \tau| + V$ when V act on space orthogonal to $|\tau\rangle$ denote it by $|\tau^\perp\rangle$. Then the distribution over V is also a k -design relative to the Haar measure on $|\tau^\perp\rangle$.

Proof. □

Definition 2.2. Denote by

$$\begin{aligned} M(\tau, U)(x) &= \max \{ |\text{Re} \langle x | U\tau \rangle|, |\text{Im} \langle x | U\tau \rangle| \} \\ \bar{M}(\tau, U)(x) &= \min \{ |\text{Re} \langle x | U\tau \rangle|, |\text{Im} \langle x | U\tau \rangle| \} \end{aligned}$$

When it will be clear from the context we omit τ, U and use only $M(x), \bar{M}(x)$.

$$|\langle \Phi(\tau, U) | U\phi \rangle|^2 = |\langle \Phi(\tau, U) | \text{Re } U\phi \rangle|^2 + |\langle \Phi(\tau, U) | \text{Im } U\phi \rangle|^2$$

$$\begin{aligned}
\langle \Phi(\tau, U_k) | MU_k \phi \rangle &= \sum_x (-1)^{\text{sign } M(\langle x | U \tau \rangle)} \frac{1}{2^{n/2}} \langle x | U \phi \rangle \\
&= \sum_{\tau, \phi \text{ agree on } x} \left| \frac{1}{2^{n/2}} M(\langle x | U \phi \rangle) \right| - \sum_{\tau, \phi \text{ disagree on } x} \left| \frac{1}{2^{n/2}} M(\langle x | U \phi \rangle) \right| \\
&\approx \sum_{\tau, \phi \text{ agree on } x} q_d \left(\frac{1}{2^{n/2}} \bar{M}(\langle x | U \phi \rangle) \right) - \sum_{\tau, \phi \text{ disagree on } x} q_d \left(\frac{1}{2^{n/2}} \bar{M}(\langle x | U \phi \rangle) \right) \pm 2^n \zeta_d \left(\frac{1}{2^{n/2}} \right)
\end{aligned}$$

noitce that we obtained a d -degree polinomial, denote it by T_ϕ .

$$\begin{aligned}
|\langle \Phi(\tau, U) | MU \phi \rangle| &\approx q_{d'}(\langle \Phi(\tau, U) | U \phi \rangle) + \zeta_{d'}(\langle \Phi(\tau, U) | U \phi \rangle) \\
&\approx q_{d'}(\langle \Phi(\tau, U) | U \phi \rangle) + \zeta_{d'}(\langle \Phi(\tau, U) | U \phi \rangle) \\
&\approx q_{d'}(T_\phi) + \zeta_{d'}(T_\phi) \\
&\approx q_{d'}(T_\phi) + \zeta_{d'}(T_\phi)
\end{aligned}$$

Assume that our k -design collection is defined such that for any $|\varphi\rangle$ it holds that:

$$\mathbf{Pr}_{U_1, U_2 \sim D} [\text{sign}(\text{Re} \langle x | U_1 \varphi \rangle) = \text{sign}(\text{Re} \langle x' | U_2 \varphi \rangle)] = \frac{1}{2}$$

Claim 2.2. *left $f : N \rightarrow \{\pm\}$ then the set $(-1)^{f(x)} |x\rangle \langle x| U$ is a k -design.*

Proof.

$$\begin{aligned}
\text{tr}(U' V'^{\dagger}) &= \text{tr} \left((-1)^{f(x)} |x\rangle \langle x| U V^{\dagger} (-1)^{f(x)} |x\rangle \langle x| \right) \\
&= \text{tr} \left((-1)^{f(y)} |y\rangle \langle y| (-1)^{f(x)} |x\rangle \langle x| U V^{\dagger} \right) = \text{tr}(U V^{\dagger})
\end{aligned}$$

So, we get that:

$$\begin{aligned}
\frac{1}{|X|^{t,2}} \sum_{U, V \in X'} |\text{tr}(U V^{\dagger})|^{2t} &= \frac{1}{|X|^2} \sum_{U, V \in X} |\text{tr}(U V^{\dagger})|^{2t} \\
&= \int |\text{tr}(U)|^{2t} dU
\end{aligned}$$

□

Claim 2.3. *Assume f above is universal hash function. Then we have that :*

$$\mathbf{E}_{U, V \sim X} [|\langle \varphi V^{\dagger} | x \rangle \langle x | U \varphi \rangle|^2] \approx_{\delta} \mathbf{E}_{U, V \sim H_{aar}} [|\langle \varphi V^{\dagger} | x \rangle \langle x | U \varphi \rangle|^2]$$

Claim 2.4. $|\langle \Phi(\tau, U_k) | U_k \tau^{\perp} \rangle \langle \tau^{\perp} U_j^{\dagger} | \Phi(\tau, U_j) \rangle|^2 < b$

Proof.

$$\begin{aligned}
&\mathbf{E}_{U \sim D} \left[|\langle \Phi(\tau, U_k) | U_k \tau^{\perp} \rangle \langle \tau^{\perp} U_j^{\dagger} | \Phi(\tau, U_j) \rangle|^2 \right] \\
&\leq \mathbf{E}_{U \sim D} \left[|\langle \Phi(\tau, U_k) | U_k \tau^{\perp} \rangle|^2 \cdot |\langle \tau^{\perp} U_j^{\dagger} | \Phi(\tau, U_j) \rangle|^2 \right] \\
&= \mathbf{E}_{U \sim D} \left[|\langle \Phi(\tau, U_k) | U_k \tau^{\perp} \rangle|^2 \right]^2 \\
&= \mathbf{E}_{U \sim D} \left[\left| \sum_x \langle x | U_k \tau^{\perp} \rangle \right|^2 \right]^2 \\
&= \mathbf{E}_{U \sim D} \left[\sum_x |\langle x | U_k \tau^{\perp} \rangle|^2 \right]^2
\end{aligned}$$

□

References

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