

## Abstract

We studies the complexity of synthesis quantum states using PRS, our reasch continues the work by [Ira+22], [Ros23], [RY21], [MY23], [Del+23].

**Claim 0.1.** *Let  $G$  be a PRS generator, than one can assume that  $G$  takes as input two register, the first contains  $n$  ancille qubits initiliazied to  $|0\rangle$  and the seconed contain a classic string initiliezied to be the seed  $k$ .*

**Claim 0.2.** *Let  $G : |0\rangle^n \otimes \mathbb{F}_2^k \rightarrow \{|\psi_k\rangle\}_{k \in \mathcal{K}}$  be a PRS generator uses  $n$ - ancilles and  $k$  classici bits. Then for any unitary  $V : \mathcal{H}_n \rightarrow \mathcal{H}_n$  it holds that  $(V \otimes I^{\otimes k})G$  is also a PRS.*

*Proof.* □

**Claim 0.3** (Levis Lemma for PRS). *Let  $f : \mathcal{H} \rightarrow \mathbb{R}$  be a **BQP**-computible fuction on the  $n$ -qubits hilbert space, and let  $g : (0, 1) \rightarrow \mathbb{R}$  a function such that:*

$$\Pr_{|\psi\rangle \sim U} [f(|\psi\rangle) > \varepsilon] < g(\varepsilon)$$

*Then, a similar inequality also holds for states sampled by the PRS, when the probability for the measure  $f$ -value grater than  $\varepsilon$  is bounded by  $g(2\varepsilon)$ . Namely,*

$$\Pr_{|\psi\rangle \sim |\psi_k\rangle} [f(|\psi\rangle) > \varepsilon] < g(2\varepsilon)$$

*In praticular, Levi's lemma has a version that capture consetration of states sampled by PRS generator, states the following: Assume there exists  $K$  such that for any  $|\psi\rangle, |\phi\rangle \in \mathcal{S}(\mathbb{C}^d)$   $|f(|\psi\rangle) - f(|\phi\rangle)| < K ||\psi\rangle - |\phi\rangle|$ . Then there exists a universal constant  $C > 0$  such:*

$$\Pr_{|\psi\rangle \sim |\psi_k\rangle} [|f(|\psi\rangle) - \mathbf{E}_{|\phi\rangle \sim U} [f(|\phi\rangle)]| > \varepsilon] < \exp\left(-\frac{Cd}{K^2}4\varepsilon^2\right)$$

*Proof.* □

**Claim 0.4.** *Probabilistic counting argument and  $\varepsilon$ -net over PRS.*

**Claim 0.5.** *exsistness of  $\text{poly}(n)$  gates  $G_1, G_2..$  such that, any  $G_i$  has a polynomial depth,  $\langle p(G_i) | \tau \rangle > a$  and  $\langle \tau^\perp | p(G_j) \rangle \langle p(G_i) | \tau^\perp \rangle < b$  for any  $i \neq j$ .*

*Proof.* □

**Claim 0.6.** *bla bla bla*

## References

- [RY21] Gregory Rosenthal and Henry Yuen. *Interactive Proofs for Synthesizing Quantum States and Unitaries*. 2021. arXiv: [2108.07192 \[quant-ph\]](#).
- [Ira+22] Sandy Irani et al. “Quantum Search-To-Decision Reductions and the State Synthesis Problem”. en. In: Schloss Dagstuhl - Leibniz-Zentrum für Informatik, 2022. DOI: [10.4230/LIPICS.CCC.2022.5](#). URL: <https://drops.dagstuhl.de/opus/volltexte/2022/16567/>.
- [Del+23] Hugo Delavenne et al. *Quantum Merlin-Arthur proof systems for synthesizing quantum states*. 2023. arXiv: [2303.01877 \[quant-ph\]](#).
- [MY23] Tony Metger and Henry Yuen. *stateQIP = statePSPACE*. 2023. arXiv: [2301.07730 \[quant-ph\]](#).
- [Ros23] Gregory Rosenthal. *Efficient Quantum State Synthesis with One Query*. 2023. arXiv: [2306.01723 \[quant-ph\]](#).