

Memory.

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1 Notations and Definitions.

Consider a code with a 2-colored (k -colored) Tanner graph, such that any two left bits of the same color share no stabilizer. For a subset of bits S , we denote by S_{c_1} its restriction to color c_1 . We use the integer Δ to denote half of the stabilizers connected to a single bit. (We assume fixed left and right degree in the graph). Our computation is subjected to p -depolarized noise. We denote by m the block length of the code. The decoder works as follows:

1. Pick a random color.
2. For any (q)bit at that color, check if flipping it decreases the syndrome. If so, then flip it.

We say that a density matrix ρ , induced on the m -length block, is a **good noisy distribution** if:

1. ρ is subjected to q - local stochastic noise.
2. Denote by S the support of an error occurring on ρ (S is a random variable). Then, with high probability¹, $|S_{c_1}| > \frac{1}{4}|S|$.

Claim 1.1. Given density ρ , which is a **good noisy distribution**, then with high probability, after correction and noise accumulation, it will remain a **good noisy distribution**.

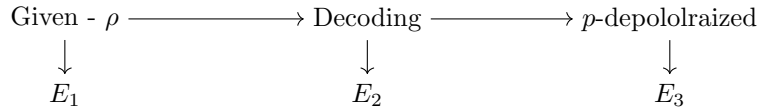


Figure 1: Illustration of the cycle.

1.1 Proof.

First, let's bound the probability that the error after the decoding round (E_2) is supported on S :

$$\Pr[\text{Sup}(E_2) = S] \leq \Pr[\text{any bit } v \in S_{c_1} \text{ sees majority of unsatisfied stabilizers}] \leq q^{\Delta|S|_{c_1}}$$

$$\begin{aligned}
 \Pr[\text{Sup}(E_3) = S] &= \sum_{S' \subset S} \Pr[\text{Sup}(E_2) = S' \cap \text{Sup}(E_3/E_2) = S/S'] \\
 &\leq \sum_{S' \subset S} q^{\Delta|S'_{c_1}|} p^{|S/S'_{c_1}|} \leq \sum_{S' \subset S} q^{\Delta|S'_{c_1}|} p^{|S_{c_1}| - |S'_{c_1}|} \\
 &\leq (q^\Delta + p)^{|S_{c_1}|} \leq \begin{cases} (q^\Delta + p)^{\frac{1}{4}|S|} & \text{if } |S_{c_1}| \geq \frac{1}{4}|S| \\ \star & \text{else} \end{cases}
 \end{aligned}$$

¹I'm leaving specifying what it is to later.

Let $S^t = \mathbf{Sup}(E)$ at time t and denote by \mathcal{P}_t the probability that $|S_{c_1}^t| > \frac{1}{4}|S_t|$. Then:

$$\begin{aligned}\mathcal{P}_{t+1} &\geq \mathbf{Pr} \left[|S_{c_1}^t| > \frac{1}{4}|S_t| \text{ and } |(S_{t+1}/S_t)_{c_1}| \geq \frac{1}{4}|S_{t+1}/S_t| \right] \\ &\geq \mathcal{P}_t \cdot (1 - e^{-\varepsilon} m) \geq \mathcal{P}_0 (1 - (t+1)e^{-\varepsilon m})\end{aligned}$$