# Hardness of Computing Fault Tolerance.

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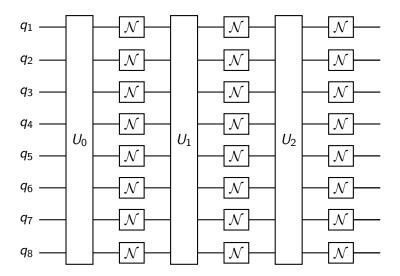
#### Introduction

- ► Brief overview of the topic
- ► Importance and relevance
- Objectives of the presentation

# **Key Points**

- ▶ Main point 1
- ► Main point 2
- ► Main point 3

# Nosiy Circuit.



Threshoold Theorem.

#### Pippenger's Construction.

Encode each bit with the repetition code  $0 \mapsto 0^m$ ,  $1 \mapsto 1^m$ . Now observe that any logical operation, without decoding, can be made in O(1) depth.

For example,  $OR(\bar{x}, \bar{y})$  can be computed by applying in parallel  $OR(x_i, y_i)$  for each i.

## The 'Decoding' trick.

Instead of completely decoding, we would apply only a single step of partial decoding. We assume that in each code block the bits are partitioned into random disjoint triples, and we will apply a local correction to each of the triples by majority.

#### Claim

There are constants  $\alpha, \eta \in (0,1)$  such that for any bit string x at distance  $\leq \alpha n$  from the code (Repetition Code), A one cycle of local correction on x yield x' such:

$$d(x',C) \leq d(x,C)$$

The Franch's Construction.