

# Bucket Sort When You Know The Distribution.

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## Abstract

We propose a new simple construction based on Tanner Codes, which yields a good LDPC code with testability query complexity of  $\Theta(n^{1-\varepsilon})$  for any  $\varepsilon > 0$ .

**The problem.** Let  $f : [0, 1] \rightarrow [0, 1]$  a fixed distribution function. Write an algorithm that sort  $n$  draws  $x_1 \dots x_n$  at linear expectation time.

**Solution.** We will define a partition of the input into a serie of  $n$  buckets  $\mathcal{B} = \{B_k = [t_k, t_{k+1}] : k \in [n]\}$  such that  $\Pr[x \in B_i] = \frac{1}{n}$  for any bucket. Assume that we seccused to compute the buckets effciently. Let the  $X_{ij}$  be the indecator of the event that  $x_j$  fall to  $B_i$ . Then we have:

$$\begin{aligned} \Pr \left[ \sum_i |B_i|^2 \geq t \right] &= \Pr \left[ \sum_i \left( \sum_j X_{ij} \right)^2 \geq t \right] \\ &= \Pr \left[ \sum_{i,j,j'} X_{i,j} X_{i,j'} \geq t \right] = \Pr \left[ \sum_{i,j \neq j'} X_{i,j} X_{i,j'} \geq t - n \right] \\ &\leq \frac{\sum_{i,j \neq j'} \mathbf{E}[X_{i,j} X_{i,j'}]}{t - n} = \frac{n}{(t - n) n^2} 2 \binom{n}{2} \leq \frac{n}{t - n} \end{aligned}$$

It follows that for any function  $t : \mathbb{N} \rightarrow \mathbb{R}$ , such that  $n = o(t)$ , sorting quaderic each bucket at turn would last almost surly less than  $t(n)$ . It lefts to show that knowing the distribution enables to compute effciently the buckets. Ensuring the uniform partitonized prtroperty leads for the follow requisive relation:

$$\begin{aligned} \frac{1}{n} &= \Pr[x \in B_k] = f(t_{k+1}) - f(t_k) \\ &\Rightarrow t_{k+1} \leftarrow f^{-1} \left( \frac{1}{n} + f(t_k) \right) \end{aligned}$$

Hence, if  $f$  can be computed in sublinear time then we obtained am expected linear time algortihm for sorting.