

$\sqrt{n} \mapsto \Theta(n)$ Magic States 'Distillation' Using Quantum LDPC Codes.

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1 The Construction.

Let x_0 be a codeword of C_X/C_Z^\perp . Denote by $w \in \mathbb{F}_2^n$ the binary string presents the Z -generator that anti commute with the X -generator corresponds to x_0 . Let $\mathcal{X} = \{x_0, x_1, \dots, x_{k'}\} \in \mathbb{F}_2^n$ be a subset of a base for the code C_X/C_Z^\perp . Such $(\text{span } \mathcal{X}/x_0)|_w$ is Triorthogonal code. Let us denote by \mathcal{X}' the base $\{y_1, y_2, \dots, y_{k'}\} \in \mathbb{F}_2^n$ defined such: $y_i = x_j + x_0$.

Denote by E the circuit that encodes the logical i th bit to y_i , by $T^{(w)}$ the application of T gates on the qubits for which w act non trivial, means $T^{(w)}$ is a tensor product of T 's and identity where on the i th qubit $T^{(w)}$ apply T if w_i is 1 and identity otherwise. And finally by D denote the gate that decode binary strings in \mathbb{F}_2^n back into the logical space.

2 Proof of Theorem 1.

Claim 2.1. *There exists family of non-trivial distance quantum LDPC codes Q such the codes span \mathcal{X}' chosen respect to them has a positive rate. Furthermore, the rate of span \mathcal{X}' is a asymptotically converges to Q rate:*

$$|\rho(Q) - \rho(\text{span } \mathcal{X}')| = o(1)$$

Proof. Let Δ be a constant integer, C_0, \tilde{C}_0 codes over Δ bits such \tilde{C}_0 is Triorthogonal and C_0 contains \tilde{C}_0 , C_0 has parameters $\Delta[1, \delta_0, \rho_0]$, and C_0^\top has relative distance greater than δ_0 . Let C_{Tanner} be a Tanner code, defined by taking an expander graph with good expansion and C_0 as the small code. Let C_{initial} be the dual-tensor code obtained by taking $(C_{\text{Tanner}}^\perp \otimes C_{\text{Tanner}}^\perp)^\perp$. Notes that first this code has positive rate and $\Theta(\sqrt{n})$ distance, second this code is an LDPC code as well. Notice also that C_{initial}^\top obtained by transporting the parity check matrix, and therefore equals to $(C_{\text{Tanner}}^{\top, \perp} \otimes C_{\text{Tanner}}^{\top, \perp})^\perp$. Hence C_{initial}^\top has a square root distance as well.

Let Q the CSS code, obtained by taking the Hyperproduct of C_{initial} with itself. So Q is an quantum qLDPC code with parameters $[n, \Theta(n^{\frac{1}{4}}), \Theta(n)]$. Pick x_0 and $w \in \mathbb{F}_2^n$, which correspond to the supports of anti commute X and Z generators, such that w can be obtains by setting a codeword of C_{Tanner} on the first $n^{\frac{1}{4}}$ bits and padding by zeros the rest. Clearly, $|w| = \Theta(n^{\frac{1}{4}})$.

Now for defying $\text{span } \mathcal{X}'$, we are going to consider the parity checks matrix obtained by adding restrictions to C_X restrictions as follows: Divide the first $n^{\frac{1}{4}}$ bits into Δ -size buckets $B_1 = \{x_1, x_2, \dots, x_\Delta\}$, $B_2 = \{x_{\Delta+1}, x_{\Delta+2}, \dots, x_{2\Delta}\}$, .., $B_i = \{x_{(i-1)\Delta+1}, x_{(i-1)\Delta+2}, \dots, x_{i\Delta}\}$, .. $B_{n^{\frac{1}{4}}/\Delta}$. Then let $\text{span } \mathcal{X}'$ be all the codewords of C_X/C_Z^\perp satisfying \tilde{C}_0 restrictions for each bucket. \square

Claim 2.2. *Let $|\mathcal{X}'\rangle \propto \sum_{x \in \text{span } \mathcal{X}'} |x\rangle$. Then $T^{(w)} |\mathcal{X}'\rangle \propto \sum_{x \in \text{span } \mathcal{X}'} x$*