Memory.

August 19, 2025

1 Notations and Definitions.

Consider a code with a 2-colorized (k-colorized) Tanner graph, such that any two left bits of the same color share no stabilizer (check). For a subset of bits S, we denote by S_{c_1} its restriction to color c_1. We use the integer Δ to denote half of the stabilizers connected to a single bit. (We assume fixed left and right degree in the graph). Our computation is subjected to p-depolarized noise. We denote by m the block length of the code. The decoder works as follows:

1. Pick a random color.
2. For any (q)bit at that color, check if flipping it decreases the syndrome. If so, then flip it.

We say that a density matrix ρ, induced on the m-length block, is a **good noisy distribution** if:

1. ρ is subjected to q - local stochastic noise.
2. Denote by S the support of an error occurring on ρ (S is a random variable). Then, with high probability[1], $|S_{c_1}| > \frac{1}{4}|S|$.

Claim 1.1. Given density ρ, which is a **good noisy distribution**, then with high probability, after correction and noise accumulation, it will remain a **good noisy distribution**.

$$\begin{tikzcd}
\text{Given - } \rho \arrow[r] \arrow[d] & \text{Decoding} \arrow[r] \arrow[d] & p\text{-depololraized} \arrow[d] \\
E_1 & E_2 & E_3
\end{tikzcd}$$

Figure 1: Illustration of the cycle.

1.1 Proof.

First, let's bound the probability that the error after the decoding round (E_2) is supported on S. (We use here the fact that views of the bits through their stabilizer don't overlap since we took only bits of the same color for the decoding):

$$\mathbf{Pr}\left[\mathbf{Sup}\left(E_2\right) = S\right] \leq \mathbf{Pr}\left[\text{any bit } v \in S_{c_1} \text{ sees majority of statisfied stabilizers }\right] \leq q^{\Delta |S|_{c_1}}$$

[1] I'm leaving specifying what it is to later.

Now, for roughly analyzing the error after observing a round of p-depolarized noise, we consider a model in which new errors due to the depolarized channel don't correct previous errors. So we get:

$$\begin{aligned}
\mathbf{Pr}\left[\mathbf{Sup}\left(E_3\right)=S\right] &\le \sum_{S'\subset S}\mathbf{Pr}\left[\mathbf{Sup}\left(E_2\right)=S'\cap\mathbf{Sup}\left(E_3/E_2\right)=S/S'\middle||S'_{c_1}|\ge\frac{1}{4}|S'|\right]\\
&+\mathbf{Pr}\left[|\mathbf{Sup}\left(E_2\right)_{c_1}|<\frac{1}{4}|\mathbf{Sup}\left(E_2\right)|\right]\\
&=\sum_{S'\subset S\text{ and }|S'_{c_1}|\ge\frac{1}{4}|S'|} q^{\Delta|S'_{c_1}|}p^{|S/S'|}+\mathbf{Pr}\left[|\mathbf{Sup}\left(E_2\right)_{c_1}|<\frac{1}{4}|\mathbf{Sup}\left(E_2\right)|\right]\\
&\le\sum_{S'\subset S} q^{\Delta\frac{1}{4}|S'|}p^{|S/S'|}+\mathbf{Pr}\left[|\mathbf{Sup}\left(E_2\right)_{c_1}|<\frac{1}{4}|\mathbf{Sup}\left(E_2\right)|\right]\\
&\le\left(q^{\frac{1}{4}\Delta}+p\right)^{|S|}+\mathbf{Pr}\left[|\mathbf{Sup}\left(E_2\right)_{c_1}|<\frac{1}{4}|\mathbf{Sup}\left(E_2\right)|\right]
\end{aligned}$$

So, it remains to show that property (2) still holds with high probability. The following is incorrect, yet almost correct. I want to say that a new error observed by the depolarized channel has to spread evenly on bits at color c_1, and by concentration get that they are far away from $\frac{1}{4}$ with probability less than $\exp(-\varepsilon m)$.

Then, let $S^t=\mathbf{Sup}\left(E\right)$ at time t and denote by $\mathcal{P}_t$ the probability that $|S^t_{c_1}|>\frac{1}{4}|S^t|$. Then:

$$\begin{aligned}
\mathcal{P}_{t+1} &\ge \mathbf{Pr}\left[|S^t_{c_1}|>\frac{1}{4}|S_t|\text{ and }|\left(S_{t+1}/S_t\right)_{c_1}|\ge\frac{1}{4}|S_{t+1}/S_t|\right]\\
&\ge\mathcal{P}_t\cdot\left(1-e^{-\varepsilon m}\right)\ge\mathcal{P}_0\left(1-e^{-\varepsilon m}\right)^{t+1}\\
&\ge\mathcal{P}_0\left(1-(t+1)e^{-\varepsilon m}\right)
\end{aligned}$$

There is a problem with the assumption that the new error spreads uniformly across the colors. In particular, m should be taken as the untapped qubits, so it changes over time and might not contain qubits of color c_1 at all.

(**[COMMENT]** See the comment in blue below, it gets complicated.)

Question. Consider the n-dimensional toric code, where qubits are placed on k-cells of the n-dimensional hypercubic lattice. For an i-cell, denote by Δ_i^+ the number of $(i+1)$-cells adjacent to it, and by Δ_i^- the number of $(i-1)$-cells adjacent to it. For which values of k do both of the following strict inequalities hold?

$$\Delta_k^+>\Delta_{k+1}^-,\qquad\Delta_k^->\Delta_{k-1}^+.$$

Answer. In an n-dimensional hypercubic lattice one has

$$\Delta_i^+=2\,(n-i),\qquad\Delta_i^-=2\,i.$$

Therefore, the two inequalities become

$$\begin{aligned}
2(n-k)>2(k+1)\qquad&\Longleftrightarrow\qquad k<\tfrac{n-1}{2},\\
2k>2\big(n-(k-1)\big)\qquad&\Longleftrightarrow\qquad k>\tfrac{n+1}{2}.
\end{aligned}$$

These conditions are mutually exclusive, since they require simultaneously

$$k<\frac{n-1}{2}\quad\text{and}\quad k>\frac{n+1}{2}.$$

Thus, there is no value of k (for any dimension n) for which both inequalities hold at once.

Yet, if one is willing to satisfy only the first inequality. Then:

$$1<\frac{\Delta_k^-}{\Delta_{k-1}^+}=\frac{2k}{2(n-(k-1))}\to k>\frac{2}{3}n$$

Should be verfied:

1. In addition the dimension of the code should be $\binom{n}{k}$. (Also known as the Betti numbers).
2. Numebr of k-cells shared by a j - cell and a i -cell. $\binom{j-i}{k-i}$.
3. The partiy of $\binom{2l}{l}$.

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