$\mathbf{QNC}_1 \subset$ noisy-**BQP**

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1 Notations.

We denote by C_g the good qLDPC code [Din+22] [PK21] [LZ22b], and by C_{ft} the concatenation code presented at [AB99] (ft stands for fault tolerance). For a code C_y, we use Φ_y, E_y, D_y to denote the channel maps circuits into the circuits computed in the code space, the encoder, and the decoder, respectively. We use Φ_U to denote the 'Bell'-state storing the gate U. We say that a state $|\psi\rangle$ is at a distance d from a quantum code C if there exists an operator U that sends $|\psi\rangle$ into C such that U is spanned on Paulis with a degree of at most d. Sometimes, when the code being used is clear from the context, we will say that a block B of qubits has absorbed at most d noise if the state encoded on B is at a distance of at most d from that code.

2 The Noise Model

3 Fault Tolerance (With Resets gates) at Linear Depth.

Claim 3.1. *There exists a value $p_{th} \in (0,1)$ such that if $p < p_{th}$, then any quantum circuit C with a depth of D and a width of W can be computed by a p-noisy circuit C', which allows for resets. The depth of C' is at most* $\max\{O(D), O(\log(WD))\}$.

3.1 Initializing Magic for Teleportation gates and encodes ancillaries.

The Protocol:

1. Initialization of zeros: The qubits are divided into blocks of size $|B|$. Each block is encoded in C_g using $D_{ft}\Phi_{ft}[E_g]\,|0^{|B|}\rangle$.
2. Initialization of Magic for Teleportation gates: The gates in the original circuit are encoded in C_g using $D_{ft}\Phi_{ft}[E_g]\,|\Phi_U\rangle$.
3. Gate teleportation: Each gate in the original circuit is replaced by a gate teleportation.
4. Error reduction: After the initialization step, at each time tick, each block runs a single round of error reduction.

Claim 3.2 (From [LZ22a]). *Assuming that an error $|e| = \gamma n$, i.e e is supported on less than γn bits, then a single correction round reduce e to an error e' such that $|e'| < \nu|e|$.*

Claim 3.3. *The gate $D_{ft}\Phi_{ft}[E_g]$ initializes states encoded in C_g subject to a $3p$-noise channel.*

Proof. Clearly, with high probability, $\Phi_{ft}[E_g]$ successfully encodes into $C_{ft} \circ C_g$, let's say with probability $1 - \frac{1}{poly(n)}$. Denote by E_i and D_i the encoder and decoder at the ith level of the concatenation construction. Recall that by definition, $D_iE_i = I$, or in other words, $D_i = E_i^{\dagger}$. Consider the decoder under $\mathcal{N}$ action: $P_2D_1P_2D_2, .., P_{i-1}D_iP_i$, by the fault-tolerance construction, a logical error at the ith stage occurs with probability p^{2^i}. Therefore, by the union bound, the probability that in one of the steps the circuit absorbs an error that is not corrected is less than $p + p^2 + p^4 + .. < 2p$. Hence, any decoded qubit absorbs noise with probability less than $2p$.

Thus, overall, we can bound the probability of a single qubit being faulty by:

$$\begin{aligned}\mathbf{Pr}\left[\text{fault}\right] &= \mathbf{Pr}\left[\text{fault}|\Phi_{ft}[E_g]\right] \cdot \mathbf{Pr}\left[\Phi_{ft}[E_g]\right] + \mathbf{Pr}\left[\text{fault}|\overline{\Phi_{ft}[E_g]}\right] \cdot \mathbf{Pr}\left[\overline{\Phi_{ft}[E_g]}\right] \\ &\leq \mathbf{Pr}\left[\text{fault}|\Phi_{ft}[E_g]\right] + \mathbf{Pr}\left[\overline{\Phi_{ft}[E_g]}\right] \leq 2p + \frac{1}{poly(n)} \leq 3p\end{aligned}$$

Remark 3.1. *In our construction, we use the concatenation code to encode blocks of length* $\log(n)$*. Therefore, any* $poly(n)$ *in the above should be replaced by* $\log(n)$*. However, this does not affect anything since the inequality does not depend on* n*.*

□

Claim 3.4. *With probability* $1 - \frac{WD}{|B|} \cdot D2e^{-2|B|(\beta-p)}$*, the total amount of noise been absorb in a block, in any time* t*, is less than* γn*.*

Proof. Consider the ith block, denoted by B_i. Using the Hoeffding's inequality we have that the probability that more than $\beta|B|$ bits are flipped at time t is less than $\leq 2e^{-2|B|(\beta-p)}$. Using the union bounds over all the blocks at all the different time locations we get that with probability $1 - \frac{WD}{|B|} \cdot D2e^{-2|B|(\beta-p)}$ the noise been absorbed in a block is less than $|\beta|B$ for the whole computation.

Denote by X_t the support's size of the error over B_i at time t. Now using Claim 3.2, given that $X_{t-1} \leq \gamma n$, it follows that the total amount of error absorbed by a block until time t can be bounded by:

$$X_t \leq \nu \cdot (X_{t-1} + \beta|B|) \leq \nu(\gamma + \beta)|B| \leq \gamma|B|$$

□

Claim 3.5. *The total depth of the circuit is* $O\left(\log n\right)$*.*

Proof. The gate for encoding $|B|$-length blocks in C_g, is a Clifford and therefore can be computed in $O(\log|B|)$ depth. The encoding of the magic/bell states, done by first compute them in the logical space (un-encoded qubits) and then by using the encoder. Hence it's foult-tolerence version of both initializing ancillaries and magic states /bell states cost $O((\log|B|) \cdot \log^c(|B|\log|B|))$ [1] depth [AB99]. Backing into C_g from C_{ft} by decoding the concatenation code takes exactly as the encoding namely.

Then using the bell measurements any of the logical gates takes $O(1)$ depth and since we use perform only a single round of error correction we get that the reaming computation till the last decoding stage is a at most constant time of the original depth. Finally we pay $O(\log|B|)$ for complete decoding. Summing all, we get:

$$\begin{aligned}&O(\log|B| \cdot \log^c(|B|\log|B|)) + O(\text{original depth}) + O(\log|B|) \\ =&O(\text{original depth}) + O(\log^c|B|)\end{aligned}$$

□

Taking the block length to be $|B| = \log((W \cdot D)^c)$ gives, by Claim 3.4, a linear[2] fault tolerance construction that success with probability $1 - \frac{1}{\log^{c_2}(W \cdot D)}$. Particularly, the fault tolerance version of circuits in $\mathbf{QNC}_1$ has logarithmic depth. Then using the construction in [Aha+96] yields a polynomial fault tolerance circuit, in the only reversible gates setting.

[1]The width of the original circuit is $|B|^2$ so the number of locations is $|B|^2 \cdot \log|B|$

[2]Assuming W is polynomial in D

References

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