1. encode k qubit in $[n, 10k, d]$ good qLDPC code. With a reducing Lemma for threshold l.

2. implement X, Z, H, T in the straightforward way.

3. The CX, need more attention. Denote by g_i a generator of C and notice that we took only 1/10-fraction of the generator in the encoding process. Now, any CX will be followed by correction step. The idea we stretch a wire according predetermined match between the qubits in the support of g_i and the qubits in the support of g_j.

4. As we took only a fraction of the code space, we can require that any codeword spanned by the g_j's has an overlap with g_i which less than $l/3$. Or in other words, the decoder can correct a non desire CX in single step $\sim$.

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