Magic States Distillation Using Quantum Expander Codes.

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1 Good Codes With Large Λ.

Definition 1.1. *Let $M \in \mathbb{F}_2^{k\times n}$ upper triangular matrix such that $k < n$. We say that M has the 1-stairs property if $M_{ij} = 1$ any $j < i$.*

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 1 & 1 & 1 & 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 0 & 1 & 1 & 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 0 & 0 & 1 & 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 0 & 0 & 0 & 1 & \cdot & \cdot & \cdot & \cdot & \cdot \end{bmatrix}$$

Claim 1.1. *Any $M \in \mathbb{F}_2^{k\times n}$ upper triangular matrix can be turn into upper triangular matrix that has the 1-stairs property by elementary operation.*

Proof. Consider the following algorithm: Let M be our initial matrix. We iterate over the rows from left to right. In the ith iteration, we check for any row $j < i$ if $M_{ji} = 1$. If not, we set M to be the matrix obtained by adding the ith row to the jth row. Since M is an upper triangular matrix, adding the ith row does not change any entry M_{js} for $s < i$. Therefore, the obtained matrix is still an upper triangular matrix and the entries at M_{js} for $j, s < i$ remain the same, namely 1 if and only if $j \leq s$.

Continuing with the process eventually yields, after k iterations, a matrix with the 1-stair property. □

Claim 1.2. *Let C be a $[n, k, d]$ binary linear code, and let Λ be subcode $\Lambda \subset C$ at dimension k'. Then there exists a code $C' = [\leq 2n, \geq k - k'/2, d]$ and a subcode of it Λ' in it at dimension $\geq k'/2$, such:*

1. *$x \in \Lambda'$ and $y \in C'$ $\ x \cdot y = 0$*
2. *$x \in \Lambda'$ and $y, z \in C'$ $\ x \cdot y \cdot z = 0$*

Proof. First, consider the upper triangular matrix obtained by applying Gaussian elimination on Λ that has the 1-stair property. Now, consider the following process: go uphill, from right to left, iterating over the matrix. Let $j = k$ be the first non-zero coordinate in the bottom row of the matrix. In the ith iteration, we ask how many rows u_m, such that $m < j$, satisfy $u_m u_j = 0$.

- If more than half of such u_m satisfy the equality, then we move on to the next iteration.
- Otherwise, we encode the jth coordinate by C_0, which maps $1 \to w$ such that $w \cdot w = 0$. This flips the value of $u_m u_j$ for any pair, so we get that the majority of pairs satisfy the equality.

Notice that because we iterate on the upper triangular matrix, we don't change the value of $u_m u_{j'}$ for any $j' > j$ (since its jth coordinate was 0 before the encoding, the encoded bit will also be 0, thus not affecting the multiplication).

Denote the set of the obtained vectors by Γ. Let $S \subset \Gamma$ be the group of vectors for which there exists at least one vector in Γ whose multiplication with them is not zero. Note that the total number of pairs with zero multiplication is greater than:

$$\frac{k'-1}{2} + \frac{k'-2}{2} + .. + \frac{2}{2} = \frac{1}{2}\frac{(k'-1)(k'-2)}{2}$$

So

$$|S| \cdot (k'-1) \leq \binom{k'}{2} - \frac{1}{2}\frac{(k'-1)(k'-2)}{2} < \frac{k'(k'-1)}{2} \Rightarrow |S| < \frac{k'}{2}$$

Set $\Lambda' \leftarrow \Gamma/S$. And we got what we wanted.

$\square$

Claim 1.3. *We can repeat Claim 1.2 by considering triple multiplications instead of pair multiplications. Let C_2 and C_3 be the codes obtained from this process. We can then guarantee the existence of $\Lambda_2 \in C_2$ and $\Lambda_3 \in C_3$ such that for any $x, y \in \Lambda_2$, $xy = 0$, and for any triple $x, y, z \in \Lambda_3$, $xyz = 0$. The code $C_2 \otimes C_3$ has a group of codewords Λ_{23} such that for any $x, y, z \in \Lambda_{23}$, $xy = 0$ and $xyz = 0$.*

Claim 1.4. *Suppose that a set of vectors $\Lambda \subset C$ satisfies the relation $xy = 0$ and $xyz = 0$ for any $x, y, z \in \Lambda$. Then, there exists a code C' with a code length roughly equal to C and a subset $\Lambda' \subset C'$ such that for any distinct $x, y, z \in \Lambda'$, $xy = 0$, $xyz = 0$, and $xx =_4 1$.*

Proof. We return to the process in Claim 1.2, but taking the standard upper triangular form of Λ instead the 1-stairs form. Notice that the rows are linear combinations of the original vectors in Λ and therefore also preserve the original relations. So now, for any $j < k$, we have that encoding the M_{jj} bit only affects the multiplication of $u_j u_j$. Thus, we will encode the jth coordinate such that the multiplication of a row by itself is 1 residue 4. $\square$

Claim 1.5. *We can repeat Claim 1.2 by flipping the bit, ensuring that the majority of pairs and triple multiplications are zero. In the end, we will have the following inequality:*

$$|S| \cdot \left(k + k^2\right) \leq \frac{1}{2}\left(k^2 + k^3\right)$$

And still we will get that $|S| \leq k/2$