good qLTC

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August 11, 2022

preamble. preamble.

Claim *for any ? $[[n, k, d]]$ **CSS** code property 1 holds*. **Proof.** let $y \in \{0,1\}^n$ be a vector such $y \in G_z^{\delta}$, let assume that $|y|_{c^{x\perp}} \leq C_2 d$ then for any $c \in C_x^{\perp}$:

$$\delta r_z \geq |H_z y| = |H_z (y + c)|$$

Robusstness *Let $\omega \leq \Delta^2$. Let C_A and C_B be codes of length Δ with minimum distance d_A and d_B. We shall say that the dual tensor code $C = C_A \otimes \mathbb{F}_2^B + \mathbb{F}_2^A \otimes C_B$ is ω-robust, if for any codeword $c \in C$ of Hamming weight $|c| \leq \omega$, there exist $A' \subset A, B' \subset B, |A'| \leq |c|/d_B, |B'| \leq |c|/d_A$, such that $c_{ab} = 0$ whenever $a \notin A', b \notin B'$.*

Claim. Subcode Robusstness. *Consider the subspaces $C^{A'} \subset C^A, C^{B'} \subset C^B$, such that the dual tensor of C_A, C_B is ω-robust then the dual tensor of $C_{A'}, C_{B'}$ is also ω-robust.*

Proof. Let c be a codeword in the dual tensor of $C_{A'}, C_{B'}$ then it's clear that c is also in the dual tensor of C_A, C_B and therfore there exists V, Usubsets of A, B respectivly such that c supported only on them, and their size is less then $|c|/d_B, |c|/d_A$.