Quantum LTC With Positive Rate

David Ponarovsky

August 11, 2022

preamble. preamble.

Claim *for any ? $[[n, k, d]]$ **CSS** code property 1 holds*. **Proof.** let $y \in \{0,1\}^n$ be a vector such $y \in G_z^\delta$, let assume that $|y|_{c^{x\perp}} \leq C_2 d$ then for any $c \in C_x^\perp$:

$$\delta r_z \geq |H_z y| = |H_z (y + c)|$$

Robusstness *Let $\omega \leq \Delta^2$. Let C_A and C_B be codes of length Δ with minimum distance d_A and d_B. We shall say that the dual tensor code $C = C_A \otimes \mathbb{F}_2^B + \mathbb{F}_2^A \otimes C_B$ is ω-robust, if for any codeword $c \in C$ of Hamming weight $|c| \leq \omega$, there exist $A' \subset A, B' \subset B, |A'| \leq |c|/d_B, |B'| \leq |c|/d_A$, such that $c_{ab} = 0$ whenever $a \notin A', b \notin B'$.*

Definition. Sub-Tensor Pair *We will say that C_A', C_B' are sub-tensor pair of C_A, C_B if each of the code is subspace of C_A, C_B respeactivly and in addition one of the minimal codeword in C_A is also contained in C_A' (and similar to C_B'.*

Note that the distance of each subcode is eqaul to the one from which its drived. And also such code can be generated efficintly by choosing Δ non trival coordinate of one of the minimal codewords and sets a check nodes over them. (Assunming that Δ is even and that there is at least one diffrent codeword in the code wich has an overlap with that minimal codewoed.

Claim. Subcode Robusstness. *Consider the sub-tensor pair $C_A' \subset C_A, C_B' \subset C_B$, such that the dual tensor of C_A, C_B is ω-robust then the dual tensor of C_A', C_B' is also ω-robust.*

Proof. Let c be a codeword in the dual tensor of C_A', C_B' then it's clear that c is also in the dual tensor of C_A, C_B and therfore there exists V, U subsets of A, B respectivly such that c supported only on them, and their size is less then $|c|/d_B, |c|/d_A$. As the length's space of the each of the subcode is indentical to his container, and by the fact that the distance of each of the subcode is equal to one which contain it, It's follow that (1) $U \subset A' = A$ and $|c|/d_A = |c|/d_{A'}$.