## Quantum LTC With Positive Rate

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preamble. preamble.

Claim for any ? [[n,k,d]] CSS code property 1 holds . **Proof.** let  $y \in \{0,1\}^n$  be a vector such  $y \in G_z^{\delta}$ , let assume that  $|y|_{c^{x^{\perp}}} \leq C_2 d$  then for any  $c \in C_x^{\perp}$ :

$$\delta r_z \ge |H_z y| = |H_z \left( y + c \right)|$$

**Robusstness** Let  $\omega \leq \Delta^2$ . Let  $C_A$  and  $C_B$  be codes of length  $\Delta$  with minimum distance  $d_A$  and  $d_B$ . We shall say that the dual tensor code  $C = C_A \otimes \mathbb{F}_2^B + \mathbb{F}_2^A \otimes C_B$  is  $\omega$ -robust, if for any codeword  $c \in C$  of Hamming weight  $|c| \leq \omega$ , there exist  $A' \subset A, B' \subset B, |A'| \leq |c|/d_B, |B'| \leq |c|/d_A$ , such that  $c_{ab} = 0$  whenever  $a \notin A', b \notin B'$ .

**Definition.** Sub-Tensor Pair We will say that  $C'_A, C'_B$  are sub-tensor pair of  $C_A, C_B$  if each of the code is subspace of  $C_A, C_B$  respeactively and in addition one of the minimal codeword in  $C_A$  is also contained in  $C'_A$  (and similar to  $C'_B$ .

Note that the distance of each subcode is equal to the one from which its drived. And also such code can be generated efficintly by choosing  $\Delta$  non trival coordinate of one of the minimal codewords and sets a check nodes over them. (Assumming that  $\Delta$  is even and that there is at least one diffrent codeword in the code wich has an overlap with that minimal codewoed.

Claim. Subcode Robusstness. Consider the subtensor pair  $C'_A \subset C_A$ ,  $C'_B \subset C_B$ , such that the dual tensor of  $C_A$ ,  $C_B$  is  $\omega$ -robust then the dual tensor of  $C'_A$ ,  $C'_B$ is also  $\omega$ -robust.

**Proof.** Let c be a codeword in the dual tensor of  $C'_A, C'_B$  then it's clear that c is also in the dual tensor of  $C_A, C_B$  and therfore there exists V, U subsets of A, B respectively such that c supported only on them, and their size is less then  $|c|/d_B, |c|/d_A$ . As the length's space of the each of the subcode is indentical to his container, and by the fact that the distance of each of the subcode is equal to one which contain it, It's follow that  $(1) \ U \subset A' = A$  and  $(2) \ |c|/d_A = |c|/d_{A'}$ .

Existance Of Sub-Tensor Pair [COMMENT] Try to prove existance by the probablistic method.

**Theorem 1.** Let  $C_0 = C_A \otimes C_B$ , and  $C_1 = C_A' \otimes C_B'$  scuh that  $C_A', C_B'$  are sub-tensor pair of  $C_A, C_B$ , and each of the code has length  $\Delta$  and relative distance  $\delta$ . Consider the G-blance product of graph with good algebraic expansion  $\Gamma_0^\square, \Gamma_1^\square$ . Then the pair of the tanner codes  $\mathcal{T}\left(\Gamma_0^\square, C_0\right)$  and  $\mathcal{T}\left(\Gamma_1^\square, C_1\right)$  define a CSS code with linear distance, positive rate, and local testbility for some constant  $\kappa$ .