Bucket Sort When You Know The Distribution.

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Abstract

We propose a new simple construction based on Tanner Codes, which yields a good LDPC code with testability query complexity of $\Theta\left(n^{1-\varepsilon}\right)$ for any $\varepsilon > 0$.

The problem. Let $f : [0,1] \to [0,1]$ a fixed distribution function. Write an algorithm that sort n draws $x_1...x_n$ at linear expectation time.

Solution. We will define a partition of the input into a seira of n buckets $\mathcal{B} = \{B_k = [t_k, t_{k+1}] : k \in [n]\}$ such that $\mathbf{Pr}\left[x \in B_i\right] = \frac{1}{n}$ for any bucket.

Claim. The probability that the size of the ith bucket excceds $t \in \mathbb{N}$ is bounded by: $\mathbf{Pr}\left[B_i \geq t\right] \leq kt^{-k}$ for every integer $k \leq n$.

Proof. Let the X_{ij} be the indecator of the event that x_j belongs to B_i. Then we have:

$$
\mathbf{E}\left[B_i^k\right] = \mathbf{E}\left[\left(\sum_j X_{ij}\right)^k\right] = \mathbf{E}\left[\sum_{J \in [n]^k} \prod_{l \in [k]} X_{iJ_l}\right]
$$

$$
= \mathbf{E}\left[\sum_{l \in [k]} \sum_{\substack{J \subset [n] \\ |J| = l}} \prod_{j \in J} X_{ij}\right]
$$

$$
= \sum_{l \in [k]} \binom{n}{l} \frac{l!}{n^l}
$$

And noitce that quantinue of sequance elements in summation is bounded by:

$$
\binom{n}{l+1} \frac{(l+1)!}{n^{l+1}} \Big/ \binom{n}{l} \frac{l!}{n^l} = \frac{n-l}{n} = 1 - \frac{l}{n} \leq 1
$$

Hence the sum over k elements is lower than k. Unsing the markov inequality we have that:

$$
\mathbf{Pr}\left[B_i \geq t\right] \leq \frac{\mathbf{E}\left[B_i^k\right]}{t^k} \leq kt^{-k}
$$

¡++¿

$$
\frac{1}{n} = \mathbf{Pr}\left[x \in B_k\right] = f\left(t_{k+1}\right) - f\left(t_k\right)
$$

$$
\Rightarrow t_{k+1} \leftarrow f^{-1}\left(\frac{1}{n} + f\left(t_k\right)\right)
$$