, Theorem Definition Claim Theorem Lemma

## Bucket Sort When You Know The Distribution.

## David Ponarovsky

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## Abstract

None

**The problem.** Let  $f:[0,1] \to [0,1]$  a fixed distribution function. Write an algorithm that sorts n draws  $x_1...x_n$  at linear expectation time.

**Solution.** We will define a partition of the input into a series of n buckets  $\mathcal{B} = \{B_k = [t_k, t_{k+1}] : k \in [n]\}$  such that  $x \in B_i = \frac{1}{n}$  for any bucket. Assume that we succeed in computing the buckets efficiently. Let the be the indicator of the event that  $x_j$  falls to  $B_i$ . Then we have:

$$\sum_{i} |B_{i}|^{2} \ge t = \sum_{i} \left( \sum_{j} X_{ij} \right)^{2} \ge t$$

$$= \sum_{i,j,j'} X_{i,j} X_{i,j'} \ge t = \sum_{i,j \ne j'} X_{i,j} X_{i,j'} \ge t - n$$

$$\le \frac{\sum_{i,j \ne j'} X_{ij} X_{ij'}}{t - n} = \frac{n}{(t - n) n^{2}} 2 \binom{n}{2} \le \frac{n}{t - n}$$

It follows that for any function  $t: \mathbb{N} \to \mathbb{R}$ , such that n = o(t), sorting quadric each bucket at turn would last almost surely less than t(n). It shows that knowing the distribution enables one to compute the buckets efficiently. Ensuring the uniform partitioned property leads to the following recursive relation:

$$\frac{1}{n} = x \in B_k = f(t_{k+1}) - f(t_k)$$
$$\Rightarrow t_{k+1} \leftarrow f^{-1}\left(\frac{1}{n} + f(t_k)\right)$$

Hence, if f can be computed in sublinear time, then we obtained an expected linear time algorithm for sorting  $\Box$  The result above demonstrates a case when knowing how the input is distributed turns the problem equivalent to facing a uniform distributed input.