1. encode k qubit in $[n, 10k, d]$ good qLDPC code. With a reducing Lemma for threshold l.

2. implement X, Z, H, T in the straightforward way.

3. The CX, need more attention. Denote by g_i a generator of C and notice that we took only $1/10$-fraction of the generator in the encoding process. Now, any CX will be followed by correction step. The idea we stretch a wire according predetermined match between the qubits in the support of g_i and the qubits in the support of g_j.

4. As we took only a fraction of the code space, we can require that any codeword spanned by the g_j's has an overlap with g_i which less than $l/3$. Or in other words, the decoder can correct a non desire CX in single step $\sim$.

5. key point of the bullet above is that by overlap we mean to the bitwise AND, otherwise , saying that we have such C is equivalence to say that we found a code with positive rate and distance greater than $\frac{1}{2}$. What we are actually want is that any code word will has a small weight, and there fore can't inferred as other codeword when we apply the CX gate.