

# Memory.

August 10, 2025

## 1 Notations and Definitions.

Consider a code with a 2-colored ( $k$ -colored) Tanner graph, such that any two left bits of the same color share no stabilizer. For a subset of bits  $S$ , we denote by  $S_{c_1}$  its restriction to color  $c_1$ . We use the integer  $\Delta$  to denote half of the stabilizers connected to a single bit. (We assume fixed left and right degree in the graph). Our computation is subjected to  $p$ -depolarized noise. We denote by  $m$  the block length of the code. The decoder works as follows:

1. Pick a random color.
2. For any (q)bit at that color, check if flipping it decreases the syndrome. If so, then flip it.

We say that a density matrix  $\rho$ , induced on the  $m$ -length block, is a **good noisy distribution** if:

1.  $\rho$  is subjected to  $q$  - local stochastic noise.
2. Denote by  $S$  the support of an error occurring on  $\rho$  ( $S$  is a random variable). Then, with high probability<sup>1</sup>,  $|S_{c_1}| > \frac{1}{4}|S|$ .

**Claim 1.1.** Given density  $\rho$ , which is a **good noisy distribution**, then with high probability, after correction and noise accumulation, it will remain a **good noisy distribution**.

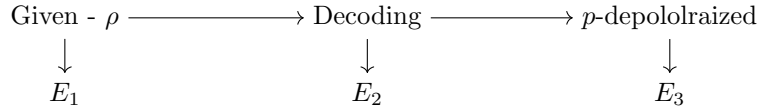


Figure 1: Illustration of the cycle.

### 1.1 Proof.

First, let's bound the probability that the error after the decoding round ( $E_2$ ) is supported on  $S$ . (We use here the fact that views of the bits through their stabilizer don't overlap since we took only bits of the same color for the decoding):

$$\Pr[\text{Sup}(E_2) = S] \leq \Pr[\text{any bit } v \in S_{c_1} \text{ sees majority of statisfied stabilizers}] \leq q^{\Delta|S|_{c_1}}$$

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<sup>1</sup>I'm leaving specifying what it is to later.

Now, for roughly analyzing the error after observing a round of  $p$ -depolarized noise, we consider a model in which new errors due to the depolarized channel don't correct previous errors. So we get:

$$\begin{aligned}
\Pr[\mathbf{Sup}(E_3) = S] &= \sum_{S' \subset S} \Pr[\mathbf{Sup}(E_2) = S' \cap \mathbf{Sup}(E_3/E_2) = S/S'] \\
&\leq \sum_{S' \subset S} q^{\Delta|S'_{c_1}|} p^{|S/S'|} \\
&= \sum_{S' \subset S \text{ and } |S'_{c_1}| \geq \frac{1}{4}|S'|} q^{\Delta|S'_{c_1}|} p^{|S/S'|} + \sum_{S' \subset S \text{ and } |S'_{c_1}| < \frac{1}{4}|S'|} q^{\Delta|S'_{c_1}|} p^{|S/S'_{c_1}|} \\
&\leq \sum_{S' \subset S \text{ and } |S'_{c_1}| \geq \frac{1}{4}|S'|} q^{\Delta|S'_{c_1}|} p^{|S/S'|} + \Pr\left[|S'_{c_1}| \geq \frac{1}{4}|S'|\right] \cdot 1 \\
&\leq \sum_{S' \subset S} q^{\Delta \frac{1}{4}|S'|} p^{|S/S'|} + \Pr\left[|S'_{c_1}| \geq \frac{1}{4}|S'|\right] \\
&\leq \left(q^{\frac{1}{4}\Delta} + p\right)^{|S|} + \Pr\left[|S'_{c_1}| \geq \frac{1}{4}|S'|\right]
\end{aligned}$$

So, it remains to show that property (2) still holds with high probability. The following is incorrect, yet almost correct. I want to say that a new error observed by the depolarized channel has to spread evenly on bits at color  $c_1$ , and by concentration get that they are far away from  $\frac{1}{4}$  with probability less than  $\exp(-\varepsilon m)$ .

Then, let  $S^t = \mathbf{Sup}(E)$  at time  $t$  and denote by  $\mathcal{P}_t$  the probability that  $|S^t_{c_1}| > \frac{1}{4}|S^t|$ . Then:

$$\begin{aligned}
\mathcal{P}_{t+1} &\geq \Pr\left[|S^t_{c_1}| > \frac{1}{4}|S^t| \text{ and } |(S_{t+1}/S^t)_{c_1}| \geq \frac{1}{4}|S_{t+1}/S^t|\right] \\
&\geq \mathcal{P}_t \cdot (1 - e^{-\varepsilon m}) \geq \mathcal{P}_0 (1 - e^{-\varepsilon m})^{t+1} \\
&\geq \mathcal{P}_0 (1 - (t+1)e^{-\varepsilon m})
\end{aligned}$$

There is a problem with the assumption that the new error spreads uniformly across the colors. In particular,  $m$  should be taken as the untapped qubits, so it changes over time and might not contain qubits of color  $c_1$  at all.