

Chapter 11

Minimum Spanning Tree Recitation.

11.1 The Spanning Tree Problem.

Definition 11.1.1. A spanning tree T of a graph $G = (V, E)$ is a subset of edges in E such that T is a tree (having no cycles), and the graph (V, T) is connected.

Problem 11.1.1 (MST). Let $G = (V, E)$ be a weighted graph with weight function $w : E \rightarrow \mathbb{R}$. We extend the weight function to subsets of E by defining the weight of $X \subset E$ to be $w(X) = \sum_{e \in X} w(e)$. The minimum spanning tree (MST) of G is the spanning tree of G that has the minimal weight according to w . Note that in general, there might be more than one MST for G .

Definition 11.1.2. Let $U \subset V$. We define the cut associated with U as the set of outer edges of U , namely all the edges $(u, v) \in E$ such that $u \in U$ and $v \notin U$. We use the notation $X = (U, \bar{U})$ to represent the cut. We say that $E' \subset E$ respects the cut if $E' \cap X = \emptyset$.

Lemma 11.1.1 (The Cut-Lemma). *Let T be an MST of G . Consider a forest $F \subset T$ and a cut X that respects F (i.e. $F \cap X = \emptyset$). Then $F \cup \arg \min_e w(e)$ is also contained in some MST. Note that it does not necessarily have to be the same tree T .*

Proof. If $e \in T$ then $F \cup \{e\} \subset T$ and we are done. So consider the second case $e \notin T$. The intersection $T \cap X$ must not be empty, otherwise there is no path connecting vertices on the opposite ends of the cut (In notations of Definition 11.1.2, no path connects vertices in U to vertices in \bar{U}), and that is a contradiction for T to be a spanning tree. Denote by $e' = T \cap X$. And consider $T' = T / \{e'\} \cup \{e\}$.

First assume that \square

In the lecture, you have seen the Kruskal algorithm for finding an MST. This algorithm constructs the MST iteratively, where in each step it holds a forest F contained in the MST and then looks for the minimal edge in a cut that it respects. Note that since F has no cycles, any edge $e \in E$ that does not close a cycle in F must belong to some cut X that is respected by F . By enforcing the order of edges being examined to be increasing in weight, it follows that the first edge that does not close a cycle is also the one with the minimum weight among them. Therefore,

by Lemma 11.1.1, we can conclude that the forest obtained by adding e into F is contained in the MST, and we can continue with it.

Result: Returns MST of given $G = (V, E, w)$

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1 sorts the  $E$  according to  $w$ 
2 define  $F_0 = \emptyset$  and  $i \leftarrow 0$ 
3 for  $e \in E$  in sorted order do
4   if  $F_i \cup \{e\}$  has no cycle then
5      $F_{i+1} \leftarrow F_i \cup \{e\}$ 
6      $i \leftarrow i + 1$ 
7   end
8 end
9 return  $F_i$ 
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Algorithm 1: Kruskal alg.