

Quicksort And Liner Time Sorts - Recitation 6

Quicksort, Countingsort, Radixsort And Bucketsort.

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Till now we have quantified the algorithm performance against the worst case scenerio. And we saw that according to that measure, In the comparisons model, one can not sorting in less then $\Theta(n \log n)$ time. In this recitation we present a two new main concepts that, in certain cases, achieve better analyses. The first one is the Expectation Complexity, By Letting the algorithm to behave undeterministically, we might obtain an algorithm that most of the time run compute the task fast. Yet we will not succumbed get down the $\Theta(n \log n)$ lower bound, but we will back to use that concept in the pending of the course. The second concept is to restrict ourselves to deal only in particular type of inputs. For example We will see that if we suppose that the a given array contains only integer in bounded domain then we can sort it in linear time.

0.1 Quicksort.

The quicksort algorithm is a good example for a **non-deterministic** algorithm that has a worst-case running time of $\Theta(n^2)$. Yet its expected running time is $\Theta(n \log n)$. Namely fix an array of n numbers, the runnings of Quicksort over that array might be different, each of them is an different event in probability space, and the running time of the algorithm is a random variable defined over that space. Saying that the algorithm has worst space complexity of $\Theta(n^2)$ means that there exist event in which it runs $\Theta(n^2)$ time with non-zero probability. But practically the interesting question is not the existence of such event but how likely that it happens. It turns out that expectation of the running time is actually $\Theta(n \log n)$.

What is the exactly reason that happens? By giving up on the algorithm behavior certainty we are going to turn the task of engineering bad input impossible.

Our study of quicksort is broken into four sections. Section 7.1 describes the algorithm and an important subroutine used by quicksort for partitioning. Because the behavior of quicksort is complex, we'll start with an intuitive discussion of its performance in Section 7.2 and analyze it precisely at the end of the chapter. Section 7.3 presents a randomized version of quicksort. When all elements are distinct,¹ this randomized algorithm has a good expected running time and no particular input elicits its worst-case behavior. (See Problem 7-2 for the case in which elements may be equal.) Section 7.4 analyzes the randomized algorithm, showing that it runs in $\Theta(n^2)$ time in the worst case and, assuming distinct elements, in expected $O(n \log n)$ time.

randomized-partition(A, p, r)

randomized-quicksort (A, p, r)

$i++$

Partitioning the array The key to the algorithm is the PARTITION procedure on the next page, which rearranges the subarray $A[p : r]$ in place, returning the index of the dividing point between the two sides of the partition. Figure 7.1 shows how PARTITION works on an 8-element array. PARTITION always selects the element $x = A[r]$ as the pivot. As the procedure runs, each element falls into exactly one of four regions, some of which may be empty. At the start of each iteration of the for loop in lines 3–6, the regions satisfy certain properties, shown in Figure 7.2. We state these properties as a loop invariant:

Partition(A, p, r)

At the beginning of each iteration of the loop of lines 3–6, for any array index k , the following conditions hold: • if $p \leq k < i$, then $A[k] \leq x$ (the tan region of Figure 7.2); • if $i \leq k < j$, then $A[k] > x$ (the blue region); • if $k = r$, then $A[k] = x$ (the yellow region). We need to show that this loop invariant is true prior to the first iteration, that each iteration of the loop maintains the invariant, that the loop terminates, and that correctness follows from the invariant when the loop terminates. Initialization: Prior to the first iteration of the loop, we have $i = p - 1$ and $j = p$. Because no values lie between p and i and no values lie between $i + 1$ and $j - 1$, the first two conditions of the loop invariant are trivially satisfied. The assignment in line 1 satisfies the third condition. Maintenance: As Figure 7.3 shows, we consider two cases, depending on the outcome of the test in line 4. Figure 7.3(a) shows what happens when $A[j] \leq x$: the only action in the loop is to increment j . After j has been incremented, the second condition holds for $A[j - 1]$ and all other entries remain unchanged. Figure 7.3(b) shows what happens when $A[j] > x$: the loop increments i , swaps $A[i]$ and $A[j]$, and then increments j . Because of the swap, we now have that $A[i] \leq x$, and condition 1 is satisfied. Similarly, we also have that $A[j - 1] > x$, since the item that was swapped into $A[j - 1]$ is, by the loop invariant, greater than x . Termination: Since the loop makes exactly $r - p$ iterations, it terminates, whereupon $j = r$. At that point, the unexamined subarray $A[j : r - 1]$ is empty, and every entry in the array belongs to one of the other three sets described by the invariant. Thus, the values in the array have been partitioned into three sets: those less than or equal to x (the low side), those greater than x (the high side), and a singleton set containing x (the pivot).

1 Appendix. Exercise from last year

Question. Consider the sets $X = \{x_1, x_2, \dots, x_n\}$, $Y = \{y_1, y_2, \dots, y_n\}$. Assume that each of the values x_i, y_i is unique. Write an Algorithm which compute the k most small items in $X \oplus Y = \{x_i + y_j : x_i \in X, y_j \in Y\}$ at $O(n + k \log k)$ time.

Solution. Notice that If $a \in X$ is greater than i elements of X and $b \in Y$ greater than j elements of Y . Then, $a + b$ greater than $i \cdot j$ elements of $X \oplus Y$. Denote by $X' = \{x'_1, \dots, x'_n\}$ ($Y' = \{y'_1, \dots, y'_n\}$) The elements of X in sorted order. So it's clear that if $x_i + y_j = x'_{i'} + y'_{j'}$ is one of the k smallest elements of $X \oplus Y$ then $i'j' \geq k$. So we will create a heap of elements that respect that inequality and then query that heap.

Heappush.

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1  $H_X \leftarrow \text{build}(X)$ 
2  $H_Y \leftarrow \text{build}(Y)$ 
3  $S_X \leftarrow \text{extract-}k(H_X)$ 
4  $S_Y \leftarrow \text{extract-}k(H_Y)$ 
5  $H_{XY} \leftarrow \text{Heap}(\{\})$ 
6 for  $i \in [k]$  do
7   for  $j \in [k/i]$  do
8      $\text{Heappush}(H_{XY}, S_{X,i} + S_{Y,j})$ 
9 return  $\text{extract-}k(H_{XY})$ 
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