

Heaps - Recitation 4

Correctness, Loop Invariant And Heaps.

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Apart from quantifying the resource requirement of our algorithms, we are also interested in proving that they indeed work. In this Recitation, we will demonstrate how to prove correctness via the notation of loop invariant. In addition, we will present the first (non-trivial) data structure in the course and prove that it allows us to compute the maximum efficiently.

Correctness And Loop Invariant.

In this course, any algorithm is defined relative to a task/problem/function, And it will be said correctly if, for any input, it computes desirable output. For example, suppose that our task is to extract the maximum element from a given array. So the input space is all the arrays of numbers, and proving that a given algorithm is correct requires proving that the algorithm's output is the maximal number for an arbitrary array. Formally:

Correctness.

We will say that an algorithm \mathcal{A} (an ordered set of operations) computes $f : D_1 \rightarrow D_2$ if for every $x \in D_1 \Rightarrow f(x) = \mathcal{A}(x)$. Sometimes when it's obvious what is the goal function f , we will abbreviate and say that \mathcal{A} is correct.

Other functions f might be including any computation task: file saving, summing numbers, posting a message in the forum, etc. Let's dive back into the maximum extraction problem and see how correctness should be proven in practice.

Task: Maximum Finding. Given $n \in \mathbb{N}$ numbers $a_1, a_2, \dots, a_n \in \mathbb{R}$ write an Algorithm that returns their maximum.

Consider the following suggestion. How would you prove it correct?

Maximum finding.

Result: returns the maximum of $a_1 \dots a_n \in \mathbb{R}^n$

```
1 let  $b \leftarrow a_1$ 
2 for  $i \in [2, n]$  do
3    $b \leftarrow \max(b, a_i)$ 
4 return  $b$ 
```

Usually, it will be convenient to divide the algorithms into subsections and then characterize and prove their correctness separately. One primary technique is using the notation of Loop Invariant. Loop Invariant is a property that is characteristic of a loop segment code and satisfies the following conditions:

Loop Invariant.

1. Initialization. The property holds (even) before the first iteration of the loop.
2. Conservation. As long as one performs the loop iterations, the property still holds.
3. Termination. Exiting from the loop carrying information.

Let's denote by $b^{(i)}$ the value of b at line number 2 at the i th iteration for $i \geq 2$ and define $b^{(1)}$ to be its value in the its initialization. What is the Loop Invariant here? **Claim.** "at the i -th iteration, $b^{(i)} = \max \{a_1 \dots a_i\}$ ".

Proof. Initialization, clearly, $b^{(1)} = a_1 = \max \{a_1\}$. Conservation, by induction, we have the base case from the initialization part, assume the correctness of the claim for any $i' < i$ and consider the i th iteration (ofcourse, assume that $i < n$). Then:

$$b^{(i)} = \max \{b^{(i-1)}, a_i\} = \max \{\max \{a_1, \dots, a_{i-1}\}, a_i\} = \max \{a_1, \dots, a_i\}$$

And that complete the Conservation part. Termination, follows by the conservation, at the n iteration, $b^{(i)}$ is seted to $\max \{a_1, a_2 \dots a_n\}$.

Task: Element finding. Given $n \in \mathbb{N}$ numbers $a_1, a_2, \dots, a_n \in \mathbb{R}$ and additional number $x \in \mathbb{R}$ write an Algorithm that returns i s.t $a_i = x$ if there exists such i and False otherwise.

Element finding.

Result: returns the maximum of $a_1 \dots a_n \in \mathbb{R}^n$

```
1 for  $i \in [n]$  do
2   if  $a_i = x$  then
3     return  $i, a_i$ 
4 return  $\Delta$ 
```

Correctness Proof. First, let's prove the following loop invariant.

Claim Suppose that the running of the algorithm reached the i 'th iteration, then $x \notin \{a_1 \dots a_{i-1}\}$.

Proof. Initialization, for $i = 1$ the claim is trivial, let's use that as the induction base case for proving Conservation. Assume the correctness of the claim for any $i' < i$. And consider the i th iteration. By the induction assumption we have that $x \notin \{a_1 \dots a_{i-1}\}$, and by the fact that we reached the i th iteration we have that in the $i - 1$ iteration, at line (2) the conditional weren't satisfied (otherwise, the function would returned at line (3) namely $x = a_{i-1}$. Hence, it follows that $x \notin \{a_1, a_2 \dots a_{i-1}\}$.

Seperate to cases. First consider, the case that given the input $a_1 \dots a_n$ the algorithm return Δ . In this case we have by the termination property that $x \notin \{a_1 \dots a_n\}$. Now, Suppose that the algorithm return the pair (i, x) then it's mean that the conditional at line (2) were satisfied at the i th iteration. So, indeed $a_i = x$ and the algorithm returns the expected output.

Task: The Superpharm Problem. You are requested to maintain a pharmacy line. In each turn, you get one of the following queries, either a new customer enters the shop, or the pharmacist requests the next person to stand in front. In addition, different customers have different priorities, So you are asked to guarantee that in each turn, the person with the height priority will be at the front.

Before we consider a sophisticated solution, What is the running time for the naive solution? (maintaining the line as a linear array) ($\sim O(n^2)$).

Heaps.

Heaps are structures that enable computing the maximum efficiently in addition to supporting adding and removing elements.

We have seen in the Lecture that no Algorithm can compute the max function with less than $n - 1$ comparisons. So our solution above is indeed the best we could expect for. The same is true for the search problem. Yet, we saw that if we are interested in storing the numbers, then, by keeping them according to sorted order, we could compute each query in logarithmic time via binary search. That raises the question, is it possible to have a similar result regarding the max problem?

Heap

Let $n \in \mathbb{N}$ and consider the sequence $H = H_1, H_2 \dots H_n \in \mathbb{R} (*)$. we will say that H is a Heap if for every $i \in [n]$ we have that: $H_i \leq H_{2i}, H_{2i+1}$ when we think of the value at indices greater than n as $H_{i>n} = -\infty$.

\Leftrightarrow

That definition is equivalent to the following recursive definition: Consider a binary tree, that we associate a number for each node. Then, we will say that this binary tree is a heap if the root's value is lower than the values of its sons and also each of the subtrees defined by its childrens is also a heap.



Checking vital signs. Are the following sequences are heaps?

1. 1,2,3,4,5,6,7,8,9,10 (Y)
2. 1,1,1,1,1,1,1,1,1,1 (Y)
3. 1,4,3,2,7,8,9,10 (N)
4. 1,4,2,5,6,3 (Y)



How much is cost (running time) to compute the min of H ? (without change the heap). ($O(1)$). Assume that option 4 is our Superpharm Line, let's try to imagine how should we maintain the line. After serving the customer at top, what can be said on $\{H_2, H_3\}$? or $\{H_{i>3}\}$? (the second highest value is in $\{H_2, H_3\}$.)

Subtask: Extracting Heap's Minimum. Let H be an Heap at size n , Write algorithm which return H_1 , erase it and returns H' , an Heap which contain all the remain elements. **Solution:**

Heappop.

Result: returns the maximum of $a_1 \dots a_n \in \mathbb{R}^n$

```

1 ret  $\leftarrow H_1$ 
2  $H_1 \leftarrow \infty$ 
3 Heapify-down(1)
4 return ret
```

Heapify-down.

Result: returns the maximum of $a_1 \dots a_n \in \mathbb{R}^n$

```

1 next  $\leftarrow i$ 
2 left  $\leftarrow 2i$ 
3 right  $\leftarrow 2i + 1$ 
4 if left <  $n$  and  $H_{\text{next}} \leq H_{\text{left}}$  then
5   | next  $\leftarrow$  left
6 if right <  $n$  and  $H_{\text{next}} \leq H_{\text{right}}$  then
7   | next  $\leftarrow$  right
8 if  $i \neq \text{next}$  then
9   |  $H_i \leftrightarrow H_{\text{next}}$ 
10  | Heapify-down(next)
```

Claim. Assume that H satisfies the Heap inequality for all the elements except the root. Namely for any $i \neq 1$ we have that $H_i \leq H_{2i}, H_{2i+1}$. Then applying Heapify-down on H at index 1 returns an heap.

Proof. By Induction on the heap size.

- Base, Consider an heap at size at most 3, and prove for each by consider each case seaprtly. (lefts as exersice).
- Assumption, assume the correctness of the claime for any tree, that satisfies the heap inequalititys except the root, at size $n' < n$.
- Induction step. Cnsider a tree at size n which and assume w.l.g (why could we?) that the right chiled of the root is the minmum between the triple. Then by the defination of the algortihm, at line (9) the values of the right child and the root are sawped. Given that before the swapping all the elements of the heap, except the root, had satisfied the heap inequality, we have that after the exchange, all the elements in right subtree, except the root of that subtree (the orignal root's right chiled) still satisfy the inequality. As the size of the right subtree is at most $n - 1$ we could use the assumption and have that after line (10) the right subtree is an heap.

Now, as the left subtree remains the same, (the values of the nodes of the left side weren't change) we have that this subtree is also an heap. So it's left to show that the root of the new ttree is smaller than both it's children. Suppose that is not the case, then it's clear that the root of the right subtree (heap) is smaller than the new root. Combine the fact that it origin must be the right subtree we have contrudiction to the fact that the orignal right subtree was an heap (as it's root wasn't the minimum elmenent in that subtree).

Question. How to construct an heap? and how much time it will take?

Build.

Result: returns the maximum of $a_1 \dots a_n \in \mathbb{R}^n$

- 1 leftHeap \leftarrow Build ($x_2 \dots x_{n/2}$)
- 2 rightHeap \leftarrow Build ($x_{n/2+1} \dots x_n$)
- 3 Heapify-down(1, $x_1 \dots x_n$)
- 4 return $x_1 \dots x_n$

Let's compute the running time of the construction algorithm. The algorithm makes two recursive calls on input at half of the original size, and then pays $\Theta(\log n)$ time to drop the first node. Hence by using the Master Theorem we have obtained that:

$$T(n) = \Theta(\log n) + 2T\left(\frac{n}{2}\right)$$

$$T(n) = \Theta(n)$$

Heapify-up.

Result: returns the maximum of $a_1 \dots a_n \in \mathbb{R}^n$

- 1 parent $\leftarrow \lfloor i/2 \rfloor$
- 2 **if** parent > 0 and $H_{\text{parent}} \leq H_i$ **then**
- 3 $H_i \leftrightarrow H_{\text{parent}}$
- 4 Heapify-up(parent)

Heappush.

Result: returns the maximum of $a_1 \dots a_n \in \mathbb{R}^n$

- 1 $H_n \leftarrow v$
- 2 Heapify-up(n)
- 3 $n \leftarrow n + 1$

1 Appendix. Exercise from last year

Question. Consider the sets $X = \{x_1, x_2 \dots x_n\}$, $Y = \{y_1, y_2 \dots y_n\}$. Assume that each of the values x_i, y_i is unieq. Write an Algorithm which compute the k most small items in $X \oplus Y = \{x_i + y_j : x_i \in X, y_j \in Y\}$ at $O(n + k \log k)$ time.

Solution. Notice that If $a \in X$ is greater than i elements of X and $b \in Y$ greater than j elemnts of Y . Then, $a + b$ greater than $i \cdot j$ elements of $X \oplus Y$. Denote by $X' = \{x'_1 \dots x'_n\}$ (Y') The elements of X in sorted order. So it's clear that if $x_i + y_j = x'_{i'} + y'_{j'}$ is one of the k smallest elements of $X \oplus Y$ then $i'j' \geq k$. So we are going to create an heap of elements which respects the that inequality, and then query that heap.

Heappush.

Result: returns the maximum of $a_1 \dots a_n \in \mathbb{R}^n$

```

1  $H_X \leftarrow \text{build}(X)$ 
2  $H_Y \leftarrow \text{build}(Y)$ 
3  $S_X \leftarrow \text{extract-}k(H_X)$ 
4  $S_Y \leftarrow \text{extract-}k(H_Y)$ 
5  $H_{XY} \leftarrow \text{Heap}(\{\})$ 
6 for  $i \in [k]$  do
7   for  $j \in [k/i]$  do
8      $\text{Heappush}(H_{XY}, S_{X,i} + S_{Y,j})$ 
9 return  $\text{extract-}k(H_{XY})$ 
```