

# Final Recitation – Information Theory, Application for Quantum Fault Tolerance.

David Ponarovsky

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- ▶ Noisy circuit and noisy computation.
- ▶ Classical no-go for reversible noisy computation at logarithmic depth (and polynomial space).
- ▶ Quantum case.
- ▶ Trading (local) entropy for space.

# Nosiy Circuit.



# Nosiy Circuit.

## Definition

$p$ - Depolarizing Channel. The qubit depolarizing channel with parameter  $p \in [0, 1]$  is the quantum channel  $\mathcal{D}_p$  defined by:

$$\mathcal{D}_p(\rho) = (1 - p)\rho + p \cdot \frac{I}{2}$$

where  $\rho$  is a single-qubit density matrix and  $I$  is the identity matrix.

## Definition

$p$ -Noisy Circuit. Given a circuit  $C$  (regardless of the model), its  $p$ -noisy version  $\tilde{C}$  is the circuit obtained by alternately taking layers from  $C$  and then passing each (qu)bit through a  $p$ -Depolarizing channel.

# Classical no-go.

## Claim

Let  $Y$  be a bit given by moving  $X$  through  $\text{BSC}(p)$ , Then there is  $\gamma_p < 1$  such :

$$1 - H(Y) \leq \gamma(1 - H(X))$$



## Classical no-go.

Denote by  $\delta$  the parameter for which  $X$  distributed as  $\sim \text{Bin}(\frac{1+\delta}{2})$ .  
First notice that:

$$\Pr(Y = 1) = \frac{1+\delta}{2}(1-p) + \frac{1-\delta}{2}p = \frac{1+\delta-2p\delta}{2}$$

So  $Y \sim \text{Bin}(\frac{1+\delta-2p\delta}{2})$ , Or  $\delta \mapsto 1 - 2p\delta$ .

## Classical no-go.

Now expand  $1 - H(X)$  to it's Taylor Seryias at  $\delta$  gives:

$$1 - H(X) = 1 - \frac{1}{2} \left( (1 + \delta) \log \left( \frac{1 + \delta}{2} \right) + (1 - \delta) \log \left( \frac{1 - \delta}{2} \right) \right)$$

Denote the above by  $K(\delta)$

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Now expand  $1 - H(X)$  to it's Taylor Series at  $\delta$  gives:

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## Classical no-go.

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Denote the above by  $K(\delta)$

## Classical no-go.

Now, observes that:

$$\begin{aligned} 1 - H(Y) &= K(2p\delta) = \sum_{i=1}^{\infty} \frac{(2p\delta)^{2n}}{2n(2n-1)} \\ &\leq (1-2p)^2 K(\delta) = (1-2p)^2 (1 - H(X)) \end{aligned}$$

And notice that since  $p < 1$  we have  $\gamma < 1$ , notice also that inequality is symmetric to  $p \mapsto 1-p$ , in particular the entropy is not increase if either  $p = 0$  or  $p = 1$ .

# Classical no-go.

## Claim

Let  $Y = (Y_1, Y_2, \dots, Y_m)$  be a bit given by moving each of  $X_i \in X = (X_1, X_2, \dots, X_m)$  through  $\text{BSC}(p)$ . Then:

$$m - H(Y) \leq \gamma(m - H(X))$$



## Classical no-go.

$$\begin{aligned} m - H(Y_1, Y_2, \dots, Y_m) &= m - \sum_i H(Y_i | Y_1, Y_2, \dots, Y_{i-1}) \\ &\leq m - \sum_i H(Y_i | X_1, X_2, \dots, X_{i-1}) \\ &\leq \sum_i 1 - H(Y_i | X_1, X_2, \dots, X_{i-1}) \\ &\leq \sum_i \gamma (1 - H(X_i | X_1, X_2, \dots, X_{i-1})) \\ &\leq \gamma \sum_i (1 - H(X_i | X_1, X_2, \dots, X_{i-1})) \\ &= \gamma (m - H(X)) \end{aligned}$$

# Classical no-go.

## Claim

Denote by  $X = (X_1, X_2, \dots, X_m)$  and  $Y = (Y_1, Y_2, \dots, Y_m)$  the input and the output distributions of reversible  $p$ -noisy computation at width  $m$  (bits) and depth  $d$ . Then, there:

$$m - H(Y) \leq \gamma^d (m - H(X))$$

In particular, for  $d = \Omega(\log m)$  we have  $H(Y) \rightarrow m$ .

# Your Title Here

# Quantum no-go.

## Claim

Let  $\rho_1$  be a reduce density matrix of  $\rho$  Then:

$$-\text{Tr } \rho \log (\rho_1 \otimes I) = S(\rho_1)$$

## Quantum no-go.

First consider the case in which  $\rho$  is a tensor of  $\rho_1$  namely  $\rho = \rho_1 \otimes \rho_2$ . Then clearly  $\rho$  and  $\log \rho_1 \otimes I$  commute. Denote by  $\lambda_1, \dots, \lambda_n$  and  $\mu_1, \dots, \mu_m$  the eigen values of  $\rho_1$  and  $\rho_2$ . So the trace equals:

$$\begin{aligned} \sum \lambda_i \mu_j \log(\lambda_i \cdot 1) &= \left( \sum \mu_j \right) \left( \sum_i \lambda_i \log \lambda_i \right) \\ &= (\text{Tr } \rho_2) \sum_i \lambda_i \log \lambda_i = -S(\rho_1) \end{aligned}$$

# Quantum no-go.

Let's use the notation  $\sum_{A_k} \rho|_{A_k}$  to denote the sum over all the reduced matrices over  $k$  qubits.

## Claim

Let  $\rho$  be a density matrix over  $n$  qubits then:

$$\binom{n}{k}^{-1} \sum_{A_k} I(\rho|_{A_k}) \leq \frac{k}{n} I(\rho)$$

## Quantum no-go.

Let  $\rho_1$  be  $\rho^{\otimes k \binom{n}{k}}$  and let  $\rho_2 = \left( \prod_{A_k} \rho|_{A_k} \right)^{\otimes n}$ .

$$\begin{aligned} 0 \geq S(\rho_2|\rho_1) &= \text{Tr} (\rho_1 (\log \rho_1 - \log \rho_2)) = -S(\rho_1) - \text{Tr} (\rho_1 \log \rho_2) \\ &= -k \binom{n}{k} S(\rho) - \sum_{A_k} \text{Tr} (\rho_1 \log (\rho|_{A_k})^n \otimes I^n) \end{aligned}$$

Now observes that  $\rho|_{A_K}^{\otimes n}$  is a reduced density matrix of  $\rho_1$ . So we get:

$$\begin{aligned} 0 &\leq -k \binom{n}{k} S(\rho) - \sum_{A_k} n S(\rho|_{A_k}) \\ \Rightarrow \sum_{A_k} I(\rho|_{A_k}) &\leq \frac{k}{n} \binom{n}{k} I(\rho) \end{aligned}$$

# Quantum no-go.

## Claim

Let  $\rho$  be a density matrix of  $n$  qubits. Let each qubit be replaced with independent probability  $p$  by a fully mixed qubit denoted by  $v$ , to give the density matrix  $\sigma$ . Then  $I(\sigma) \leq (1 - p) I(\rho)$ .



## Quantum no-go.

Let us write:

$$\sigma = \sum_{k=1}^n \sum_{A_k} p^{n-k} (1-p)^k \rho|_{A_k} \otimes v^{n-k}$$

By the concavity of the entropy (convexity of  $I$ ), We have:

$$\begin{aligned} I(\sigma) &\leq \sum_{k=1}^n \sum_{A_k} p^{n-k} (1-p)^k [I(\rho|_{A_k}) + (n-k)I(v)] \\ &= \sum_{k=1}^n \sum_{A_k} p^{n-k} (1-p)^k I(\rho|_{A_k}) \\ &\leq \sum_{k=1}^n \sum_{A_k} p^{n-k} (1-p)^k \frac{k}{n} \binom{n}{k} I(\rho) \\ &= (1-p) I(\rho) \end{aligned}$$

## Entropy-Space tradeoff.

Zeros Distillation. Consider the unitary majority gate over 3 qubits, which on the computational basis sets the last 3rd bit to be the majority, and acts as follows:

$$M |0, 0, 1\rangle \mapsto |1, 1, 0\rangle$$

$$M |1, 1, 0\rangle \mapsto |0, 0, 1\rangle$$

And acts trivially on every other configuration.

### Claim

Consider the noisy zero  $\rho = (1 - \frac{p}{2}) |0\rangle \langle 0| + \frac{p}{2} |1\rangle \langle 1|$ . Then:

$$\text{Tr}_{[1,2]} M \rho^3 = \left(1 - \frac{3}{4}p^2 + \frac{1}{4}p^3\right) |0\rangle \langle 0| + \dots$$

## Entropy-Space tradeoff.

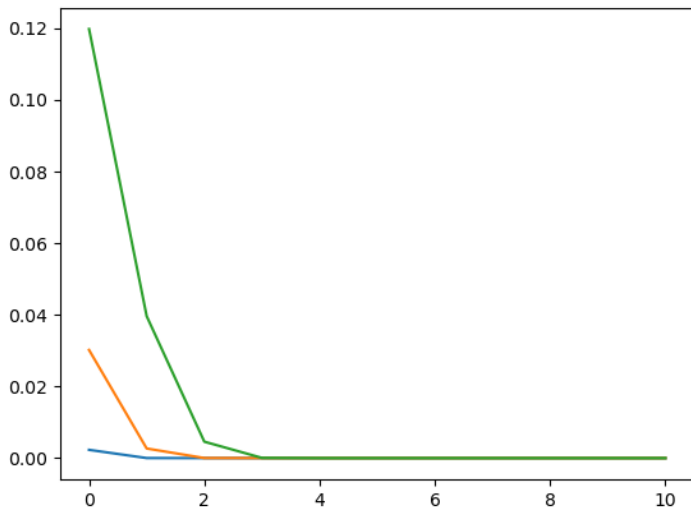
$$\begin{aligned}\rho^{\otimes 3} = & \left(1 - \frac{p}{2}\right)^3 |000\rangle \langle 000| \\ & + \left(1 - \frac{p}{2}\right)^2 \frac{p}{2} (|001\rangle \langle 001| + |010\rangle \langle 010| + |100\rangle \langle 100|) \\ & + \left(1 - \frac{p}{2}\right) \left(\frac{p}{2}\right)^2 (|110\rangle \langle 110| + |101\rangle \langle 101| + |011\rangle \langle 011|) \\ & + \left(\frac{p}{2}\right)^3 |111\rangle \langle 111|\end{aligned}$$

$$\begin{aligned}M_{\rho^{\otimes 3}} = & " \\ & + \left(1 - \frac{p}{2}\right)^2 \frac{p}{2} (|110\rangle \langle 110| + ") \\ & + \left(1 - \frac{p}{2}\right) \left(\frac{p}{2}\right)^2 (|001\rangle \langle 001| + ") \\ & + "$$

## Entropy-Space tradeoff.

$$\begin{aligned}\mathbf{Tr}_{[1,2]} M_{\rho^{\otimes 3}} &= \left( \left(1 - \frac{p}{2}\right)^3 + 3 \left(1 - \frac{p}{2}\right)^2 \frac{p}{2} \right) |0\rangle \langle 0| + \\ &\quad + \left( 3 \left(1 - \frac{p}{2}\right) \left(\frac{p}{2}\right)^2 + \left(\frac{p}{2}\right)^3 \right) |1\rangle \langle 1| \\ &= \left( 1 - \frac{3}{4}p^2 + \frac{1}{4}p^3 \right) |0\rangle \langle 0| + , ,\end{aligned}$$

# Entropy-Space tradeoff.



## Entropy-Space tradeoff.

$$\begin{aligned} & \left(1 - \frac{3}{4}p^2 + \frac{1}{4}p^3\right) |0\rangle \langle 0| + , , \\ \mapsto & (1-p) \left(1 - \frac{3}{4}p^2 + \frac{1}{4}p^3\right) |0\rangle \langle 0| + \frac{p}{2} |0\rangle \langle 0| \\ = & 1 - \frac{1}{2}p - \frac{3}{4}p^2 + p^3 - \frac{1}{4}p^4 \end{aligned}$$

# Entropy-Space tradeoff.

the probability of the third qubit to be 1 is still less than 1. There is  $p_0$  such that for any  $p < p_0$ , it follows that if each qubit has a probability less than  $p$  of being at  $|1\rangle$ , then after a cycle of Noise  $\rightarrow$  Majority, the probability of the third qubit being  $|1\rangle$  is still less than  $p$ .

## Entropy-Space tradeoff.

Proof. After the noise round, the density matrices of each qubit:

$$\begin{aligned} &\leq (1-p)((1-p)|0\rangle\langle 0| + p|1\rangle\langle 1|) + p\frac{1}{2}(|0\rangle\langle 0| + |1\rangle\langle 1|) \\ &= \left((1-p)^2 + \frac{1}{2}p\right)|0\rangle\langle 0| + \left((1-p)p + \frac{1}{2}p\right)|1\rangle\langle 1| \\ &= |0\rangle\langle 0| + (3/2p - p^2)|1\rangle\langle 1| \end{aligned}$$

And for convenience, let's bound by looking at the noisier state, where each qubit is in

$$(1-2p)|0\rangle\langle 0| + 2p|1\rangle\langle 1|$$

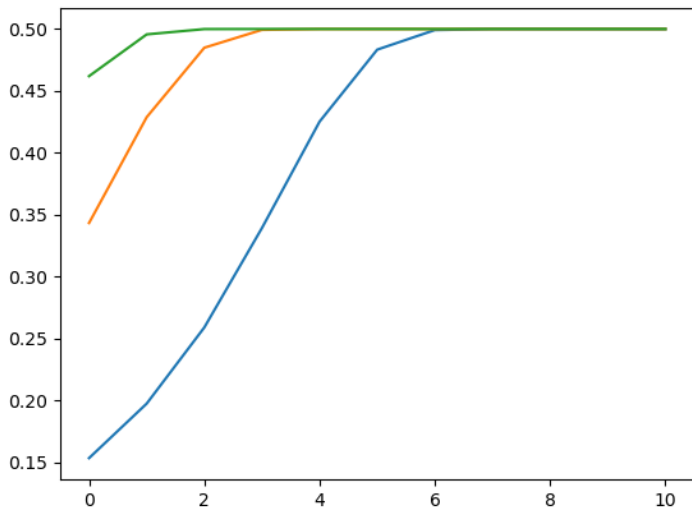
Now, after applying the majority gate, the kets whose third bit is 1 are kets in  $\rho^{\otimes 3}$  which absorb at least two flips. So, after tracing out the first two qubits, the coefficient of  $|1\rangle\langle 1|$  would be less than:

$$(2p)^3 + 3(1-2p)(2p)^2 \leq (2p)^3 + 3(2p)^2$$

And for  $p_0 = \frac{1}{6}$  we have that this probability is less than  $p$ .



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