Final Recitation – Information Theory, Application for Quantum Fault Tolerance.

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Introduction

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- Nosiy circuit and noisy computation.

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- Classical no-go for reversible noisy computation at logarithmic depth (and polynomial space).

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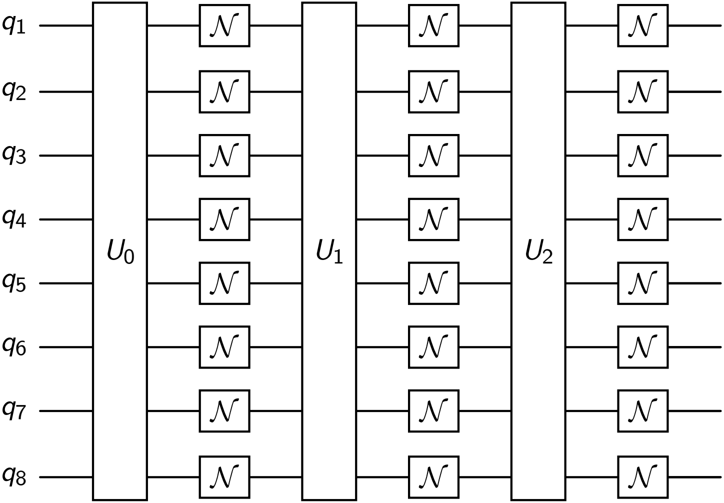
- Classical no-go for reversible noisy computation at logarithmic depth (and polynomial space).
- Quantum case.
- Trading (local) entropy for space.

Nosiy Circuit.



Nosiy Circuit.

Definition

p- Depolarizing Channel. The qubit depolarizing channel with parameter $p \in [0, 1]$ is the quantum channel $\mathcal{D}_p$ defined by:

$$\mathcal{D}_p(\rho) = (1 - p)\rho + p \cdot \frac{I}{2}$$

where ρ is a single-qubit density matrix and I is the identity matrix.

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Definition

p-Noisy Circuit. Given a circuit C (regardless of the model), its *p*-noisy version $\tilde{C}$ is the circuit obtained by alternately taking layers from C and then passing each (qu)bit through a *p*-Depolarizing channel.

Classical no-go.

Informal:

We can't compute when subjected to noise, after a while, the bits become garbage.

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Claim

Denote b $X = (X_1, X_2, .., X_m)$ and $Y = (Y_1, Y_2, .., Y_m)$ the input and the output distrubtions of reversible p-noisy compuation at widith m (bits) and depth d. Then:

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$$m - H(Y) \leq \gamma^d \left(m - H(X)\right)$$

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$$m - H(Y) \leq \gamma^d \left(m - H(X)\right)$$

In particular, for $d = \Omega(\log m)$ we have $H(Y) \to m$.

Classical no-go.

Claim

Let Y be a bit given by moving X trough BSC(p), Then there is $\gamma_p < 1$ such :

$$1 - H(Y) \leq \gamma\,(1 - H(X))$$

Classical no-go.

Denote by δ the parameter for which X distributed as $\sim \text{Bin}(\frac{1+\delta}{2})$. First notice that:

$$\mathbf{Pr}\left(Y = 1\right) = \frac{1+\delta}{2}(1-p) + \frac{1-\delta}{2}p = \frac{1+\delta-2\delta p}{2}$$

So $Y \sim \text{Bin}(\frac{1-\delta(1-2p)}{2})$, Or $\delta \mapsto 1 - 2p\delta$.

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Now expand $1 - H(X)$ to it's Taylor Seryias at δ gives:

$$1 - H(X) = 1 - \frac{1}{2}\left((1+\delta)\log\left(\frac{1+\delta}{2}\right) + (1-\delta)\log\left(\frac{1-\delta}{2}\right)\right)$$

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&= -\frac{1}{2} \cdot (1+\delta)\sum_{i=1}^{\infty}\frac{(-1)^{n+1}\delta^n}{n} + (1-\delta)\sum_{i=1}^{\infty}\frac{(-1)^{n+1}(-\delta)^n}{n}
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Denote the above by $K(\delta)$

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Now, observes that:

$$1 - H(Y) = K(2p\delta) = \sum_{i=1}^{\infty} \frac{(2p\delta)^{2n}}{2n(2n-1)}$$

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And notice that since $p < 1$ we have $\gamma < 1$, noitce also that inequlity is symmetric to $p \mapsto 1 - p$, in paritcular the entropy is not increase if eithrr $p = 0$ or $p = 1$.

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Claim

Let $Y = (Y_1, Y_2, .., Y_m)$ be a bit given by moving each of $X_i \in X = (X_1, X_2, .., X_m)$ trough BSC(p). Then:

$$m - H(Y) \leq \gamma\,(m - H(X))$$

Classical no-go.

Proof.

$$m - H\left(Y_1, Y_2, .., Y_m\right) = m - \sum_i H\left(Y_i \middle| Y_1, Y_2, .., Y_{i-1}\right)$$

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m - H\left(Y_1, Y_2, .., Y_m\right) &= m - \sum_i H\left(Y_i \middle| Y_1, Y_2, .., Y_{i-1}\right) \\
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\begin{aligned}
m - H\left(Y_1, Y_2, .., Y_m\right) &= m - \sum_i H\left(Y_i | Y_1, Y_2, .., Y_{i-1}\right) \\
&\leq m - \sum_i H\left(Y_i | X_1, X_2, .., X_{i-1}\right) \\
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&= \gamma (m - H(X))
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In particular, for $d = \Omega(\log m)$ we have $H(Y) \rightarrow m$.

Quantum no-go.

Let's continue to the quantum case. From now on we will denote by $I(\rho) = \dim \rho - S(\rho)$.

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Quantum no-go.

Claim

Let ρ_1 be a reduce density matrix of ρ Then:

$$-\mathbf{Tr}\ \rho \log\left(\rho_1 \otimes I\right) = S(\rho_1)$$

Quantum no-go.

First consider the case in which ρ is a tensor of ρ_1 namely $\rho = \rho_1 \otimes \rho_2$, Then clearly ρ and $\log \rho_1 \otimes I$ commute. Denote by $\lambda_1, ..\lambda_n$ and $\mu_1, ..\mu_m$ the eigen values of ρ_1 and ρ_2. So the trace equals:

$$\sum \lambda_i \mu_j \log(\lambda_i \cdot 1) = \left(\sum \mu_j\right) \left(\sum_i \lambda_i \log \lambda_i\right)$$

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$$\begin{aligned}\sum \lambda_i \mu_j \log(\lambda_i \cdot 1) &= \left(\sum \mu_j\right) \left(\sum_i \lambda_i \log \lambda_i\right) \\ &= (\mathbf{Tr}\ \rho_2) \sum_i \lambda_i \log \lambda_i = -S(\rho_1)\end{aligned}$$

Quantum no-go.

Let's use the notation $\sum_{A_k} \rho|_{A_k}$ to denote the sum over all the reduced matrices over k qubits.

Claim

Let ρ be a density matrix over n qubits then:

$$\binom{n}{k}^{-1} \sum_{A_k} I(\rho|_{A_k}) \leq \frac{k}{n} I(\rho)$$

Quantum no-go.

Proof

Let ρ_1 be $\rho^{\otimes k\binom{n}{k}}$ and let $\rho_2 = \left(\prod_{A_k} \rho|_{A_k}\right)^{\otimes n}$.

$$0 \geq S(\rho_2|\rho_1) = \mathbf{Tr}\ \left(\rho_1 \left(\log \rho_1 - \log \rho_2\right)\right) = -S(\rho_1) - \mathbf{Tr}\ \left(\rho_1 \log \rho_2\right)$$

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Now observes that $\rho|_{A_K}^{\otimes n}$ is a reduced density matrix of ρ_1. So we get:

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$$\Rightarrow$$

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Now observes that $\rho|_{A_K}^{\otimes n}$ is a reduced density matrix of ρ_1. So we get:

$$0 \leq -k\binom{n}{k} S(\rho) - \sum_{A_k} nS\left(\rho|_{A_k}\right)$$
$$\Rightarrow \sum_{A_k} I\left(\rho|_{A_k}\right) \leq \frac{k}{n}\binom{n}{k} I\left(\rho\right)$$

Quantum no-go.

Claim

Let ρ be a density matrix of n qubits. Let each qubit be replaced with independent probability p by a fully mixed qubit denoted by υ, to give the density matrix σ. Then $I(\sigma) \leq (1-p)\, I(\rho)$.

Quantum no-go.

Let us write:

$$\sigma = \sum_{k=1}^{n} \sum_{A_k} p^{n-k} (1-p)^k \rho|_{A_k} \otimes \upsilon^{n-k}$$

By the concavity of the entropy (convexity of I), We have:

$$I(\sigma)$$

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$$I\left(\sigma\right) \leq \sum_{k=1}^{n} \sum_{A_k} p^{n-k} \left(1-p\right)^k \left[I\left(\rho|_{A_k}\right) + (n-k) I(\upsilon)\right]$$

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I(\sigma) &\le \sum_{k=1}^{n} \sum_{A_k} p^{n-k} (1-p)^k \left[I\left(\rho|_{A_k}\right) + (n-k) I(\upsilon) \right] \\
&= \sum_{k=1}^{n} \sum_{A_k} p^{n-k} (1-p)^k \, I\left(\rho|_{A_k}\right) \\
&\le \sum_{k=1}^{n} \sum_{A_k} p^{n-k} (1-p)^k \, \frac{k}{n} \binom{n}{k} I(\rho) \\
&= (1-p)\, I(\rho)
\end{aligned}$$

Entropy-Space tradeoff.

Zeros Distillation. Consider the unitary majority gate over 3 qubits, which on the computational basis sets the last 3rd bit to be the majority, and acts as follows:

$$M|0,0,1\rangle \mapsto |1,1,0\rangle$$
$$M|1,1,0\rangle \mapsto |0,0,1\rangle$$

And acts trivially on every other configuration.

Entropy-Space tradeoff.

Claim

There is a p_0 such that for any $p < p_0$, it follows that if each qubit has a probability less than p of being at $|1\rangle$, then after a cycle of Noise $\rightarrow$ Majority, the probability of the third qubit being $|1\rangle$ is still less than p.

Entropy-Space tradeoff.

Proof.

After the noise round, the density matrices of each qubit[1]:

$$\leq (1-p)\left((1-p)\left|0\right\rangle\left\langle 0\right| + p\left|1\right\rangle\left\langle 1\right|\right) + p\frac{1}{2}\left(\left|0\right\rangle\left\langle 0\right| + \left|1\right\rangle\left\langle 1\right|\right)$$

[1] The inequality refers to the $\left|1\right\rangle\left\langle 1\right|$ coefficient.

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&= \left((1-p)^2 + \frac{1}{2}p\right)\left|0\right\rangle\left\langle 0\right| + \left((1-p)p + \frac{1}{2}p\right)\left|1\right\rangle\left\langle 1\right|
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$$= \left((1-p)^2 + \frac{1}{2}p\right)\left|0\right\rangle\left\langle 0\right| + \left((1-p)p + \frac{1}{2}p\right)\left|1\right\rangle\left\langle 1\right|$$
$$= '' \left|0\right\rangle\left\langle 0\right| + \left(3/2p - p^2\right)\left|1\right\rangle\left\langle 1\right|$$

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&= \left((1-p)^2 + \frac{1}{2}p\right)\left|0\right\rangle\left\langle 0\right| + \left((1-p)\,p + \frac{1}{2}p\right)\left|1\right\rangle\left\langle 1\right| \\
&= \text{''}\left|0\right\rangle\left\langle 0\right| + \left(3/2p - p^2\right)\left|1\right\rangle\left\langle 1\right|
\end{aligned}
$$

And for convenience, let's bound it by looking at the following noisier state, where each qubit is in:

$$
(1-2p)\left|0\right\rangle\left\langle 0\right| + 2p\left|1\right\rangle\left\langle 1\right|
$$

[1] The inequality refers to the $\left|1\right\rangle\left\langle 1\right|$ coefficient.

Entropy-Space tradeoff.

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And for $p_0 = \frac{1}{6}$ we have that this probability is less than p.

The End.

June 25, 2025