

# Quantum Information Theory - 67749

## Recitation 2, May 7, 2025

### 1 Overview - Quantum States as Computational Resources.

In the last lectures, we saw that quantum states can be considered as resources. In particular, we saw that shared **EPR** pair ( $\mathbf{Bell}_{00}$ ) enables one:

1. Transmit two classical bits by sending a single qubit, via the superdense-coding.
2. 'Teleoperate' a qubit by sending two classical bits. From an engineering point of view, it means that for having a complete quantum internet, it's enough to provide a mechanism to distribute **EPR** pairs.

### 2 Dense Encoding.

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### 3 Quantum Teleportation.

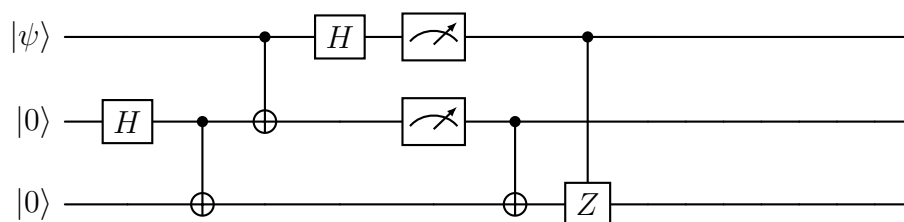


Figure 1: Measuring the single-qubit state  $|\psi\rangle$  at the  $\{|+\rangle, |-\rangle\}$  base.

## 4 Gate Teleportation.

Gate teleportation is a method to 'encode' operations by states. At the high level, given a precomputed state, it allows one to apply an operation (gate) by using (probably) simpler gates. The precomputed states are called **Magic States**.

### 4.1 Leading Example: $T$ -Teleportation.

Recall that the Clifford<sup>1</sup> +  $T$  is a universal quantum gate set. The Clifford group alone is considered from the computer science point of view a simple/weak computational class since it can be classically simulated<sup>2</sup>. Yet, we will see that given access to the magic  $|T\rangle = T|+\rangle$ , one can simulate the  $T$  gate using only Clifford gates and measurements.

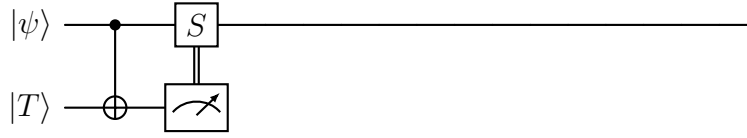


Figure 2: Measuring the single-qubit state  $|\psi\rangle$  at the  $\{|+\rangle, |-\rangle\}$  base.

$$\begin{aligned}
 \left( \sum_x \alpha_x |x\rangle \right) \otimes \frac{1}{\sqrt{2}} (|0\rangle + e^{i\frac{\pi}{4}} |1\rangle) &\xrightarrow{\text{CX}} \sum_{x,y} \frac{1}{\sqrt{2}} \alpha_x |x\rangle |x \oplus y\rangle e^{i\frac{\pi}{4}y} \\
 &\mapsto \begin{cases} \sum_x \alpha_x |x\rangle e^{i\frac{\pi}{4}x} = T|\psi\rangle & \text{measured 0} \\ \sum_x \alpha_x |x\rangle e^{i\frac{\pi}{4}\bar{x}} & \text{measured 1} \end{cases} \\
 &\xrightarrow{\text{CS}} \begin{cases} T|\psi\rangle \\ \sum_x \alpha_x |x\rangle e^{i(\frac{\pi}{4}\bar{x} + \frac{\pi}{2}x)} = \sum_x \alpha_x |x\rangle e^{i\frac{\pi}{4}} e^{i(\frac{\pi}{4}\bar{x} + \frac{\pi}{4}x)} \end{cases} \\
 &= \begin{cases} T|\psi\rangle \\ e^{i\frac{\pi}{4}} \sum_x \alpha_x |x\rangle e^{i\frac{\pi}{4}} = e^{i\frac{\pi}{4}} T|\psi\rangle \end{cases}
 \end{aligned}$$

### 4.2 Extends it.

Let's extend it to a general gate. First create  $|\mathbf{GHZ}_{2n}\rangle$  state, then

Let's split upon the measurement result.

1. If we measured 0, means the states 'agreed' in the computational base.

$$|\psi\rangle \otimes \left( \sum_x |x\rangle \otimes U|x\rangle \right)$$

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<sup>1</sup>Generated by  $H, S$  and  $CX$

<sup>2</sup>And conjectured to be strictly weaker than  $\mathbf{P}$

## 5 Uhlmann's theorem

**Claim 5.1.**

$$\langle \Omega | A \otimes B | \Omega \rangle = \text{Tr} AB^\dagger$$

*Proof.*

$$\begin{aligned} \langle \Omega | A \otimes B | \Omega \rangle &= \sum_{ij} \langle i; i | AB | j; j \rangle = \sum_{ij} \langle i | A | j \rangle \langle i | B | j \rangle = \sum_{ij} \langle i | A | j \rangle \langle j | B^\dagger | i \rangle \\ &= \sum_i \langle i | AB^\dagger | i \rangle = \text{Tr} AB^\dagger \end{aligned}$$

□

$$|\psi_\rho\rangle = \sum_i \left( \rho^{\frac{1}{2}} |\psi_i\rangle \right) |i\rangle = \sum_i \left( \rho^{\frac{1}{2}} U_\rho |i\rangle \right) |i\rangle = \left( \rho^{\frac{1}{2}} U_\rho \right) \otimes I |\Omega\rangle$$

$$|\psi_\sigma\rangle = \sum_i \left( \sigma^{\frac{1}{2}} |\psi'_i\rangle \right) |i'\rangle = \sum_i \left( \sigma^{\frac{1}{2}} U_\sigma |i\rangle \right) V |i\rangle = \left( \sigma^{\frac{1}{2}} U_\sigma \right) \otimes V |\Omega\rangle$$

**Claim 5.2.** For any square matrix  $A$ :

$$\max_{U \in \mathcal{U}} \text{Tr} AU = \text{Tr} \sqrt{A^\dagger A}$$

$$\begin{aligned} \max |\langle \psi_\rho | \psi_\sigma \rangle|^2 &= \max |\langle \Omega | \left( U_\rho^\dagger \rho^{\frac{1}{2}} \right) \otimes I \left( \sigma^{\frac{1}{2}} U_\sigma \right) \otimes V | \Omega \rangle|^2 \\ &= \max |\text{Tr} \left[ \left( U_\rho^\dagger \rho^{\frac{1}{2}} \sigma^{\frac{1}{2}} U_\sigma \right) V^\dagger \right]|^2 \\ &= \max |\text{Tr} \left[ \rho^{\frac{1}{2}} \sigma^{\frac{1}{2}} V^\dagger \right]|^2 \\ &\leq \left| \text{Tr} \sqrt{\rho^{\frac{1}{2}} \sigma^{\frac{1}{2}} \sigma^{\frac{1}{2}} \rho^{\frac{1}{2}}} \right|^2 = \left| \text{Tr} \sqrt{\rho^{\frac{1}{2}} \sigma \rho^{\frac{1}{2}}} \right|^2 \end{aligned}$$

## 6 Monotonicity of Fidelity.

Let  $\rho_{AB}, \sigma_{AB} \in \mathcal{D}(\mathcal{H}_A \otimes \mathcal{H}_B)$ . Then the fidelity is non-decreasing with respect to the partial trace:

$$F(\rho_{AB}, \sigma_{AB}) \leq F(\rho_A, \sigma_A),$$

where  $\rho_A = \text{Tr}_B \{\rho_{AB}\}$  and  $\sigma_A = \text{Tr}_B \{\sigma_{AB}\}$ .

*Proof.* Consider fixed purifications  $|\psi\rangle_{RAB}$  and  $|\phi\rangle_{RAB}$  of  $\rho_{AB}$  and  $\sigma_{AB}$ , respectively, which also purify  $\rho_A$  and  $\sigma_A$ . By Uhlmann's theorem,

$$F(\rho_{AB}, \sigma_{AB}) = \max_{U_R} |\langle \psi | U_R \otimes I_A \otimes I_B | \phi \rangle|^2.$$

On the other hand, since  $U_R \otimes I_A$  is a subset of the larger class of unitaries  $U_{RB}$  on  $RB$ ,

$$F(\rho_A, \sigma_A) = \max_{U_{RB}} |\langle \psi | U_{RB} \otimes I_A | \phi \rangle|^2 \geq F(\rho_{AB}, \sigma_{AB}).$$

Thus, we conclude that

$$F(\rho_{AB}, \sigma_{AB}) \leq F(\rho_A, \sigma_A).$$

□

Notice that  $|i\rangle\langle j|$  is unitray since.

## 7 $|\text{EPR}\rangle$ Distillation.

## 8 Magic State Distillation.

**Question.** Can we purify noisy magic states into high-fidelity ones, using only Clifford operations?

Magic state distillation is a procedure that uses many copies of noisy magic states, plus only Clifford gates and measurements, to produce fewer, higher-fidelity magic states.