# Final Recitation – Information Theory, Application.

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June 19, 2025

#### Introduction

- ▶ Brief overview of the topic
- ► Importance and relevance
- Objectives of the presentation

# **Key Points**

- ▶ Main point 1
- ► Main point 2
- ► Main point 3

#### Claim

Let Y be a bit given by moving X trough BSC(p). Then:

$$1 - H(Y) \le (1 - p^2) (1 - H(X))$$

Denote by  $\delta$  the parameter for which X distributed as  $\sim Bin(\frac{1+\delta}{2})$ . First notice that:

$$\Pr(Y = 1) = \frac{1+\delta}{2}(1-p) + \frac{1-\delta}{2}p = \frac{1-2\delta p}{2}$$

So 
$$Y \sim \text{Bin}(\frac{1-2\delta p}{2})$$
, Or  $\delta \mapsto -2p\delta$ .

Now expand 1 - H(X) to it's Taylor Servias at  $\delta$  gives:

$$1 - H(X) = 1 - \frac{1}{2} \left( (1 + \delta) \log \left( \frac{1 + \delta}{2} \right) + (1 - \delta) \log \left( \frac{1 - \delta}{2} \right) \right)$$
$$= -\frac{1}{2} \left( (1 + \delta) \log \left( \frac{1 + \delta}{2} \right) + (1 - \delta) \log \left( \frac{1 - \delta}{2} \right) \right)$$

Denote the above by  $K(\delta)$ 

Now, observes that:

$$1 - H(Y) = K(2p\delta) = \sum_{i=1}^{\infty} \frac{(2p\delta)^{2n}}{2n(2n-1)}$$
  
  $\leq 2p^2K(\delta) = 2p^2(1 - H(X))$ 

#### Claim

Let  $Y = (Y_1, Y_2, ..., Y_m)$  be a bit given by moving each of  $X_i \in X = (X_1, X_2, ..., X_m)$  trough BSC(p). Then:

$$m - H(Y) \le (1 - p^2) (m - H(X))$$

$$m - H(Y_1, Y_2, ..., Y_m) = m - \sum_{i} H(Y_i | Y_1, Y_2, ..., Y_{i-1})$$

$$\leq m - \sum_{i} H(Y_i | X_1, X_2, ..., X_{i-1})$$

$$\leq \sum_{i} 1 - H(Y_i | X_1, X_2, ..., X_{i-1})$$

$$\leq \sum_{i} (1 - p^2) (1 - H(X_i | X_1, X_2, ..., X_{i-1}))$$

$$\leq (1 - p^2) \sum_{i} (1 - H(X_i | X_1, X_2, ..., X_{i-1}))$$

$$= (1 - p^2) (m - H(X))$$