

# Online Computation, Ex 3.

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**ex1.** Consider the experts setting with gains:  $g_{i,t} \in [0, 1]$  is the gain of expert  $i$  at step  $t$ . Hedge updates:

$$P_{i,t+1} = \frac{e^{\eta G_{i,t}}}{\sum_j e^{\eta G_{j,t}}}$$

where  $G_{i,t} = \sum_{s \leq t} g_{i,s}$ . Prove that the regret of Hedge at time  $T$  is  $O(\sqrt{T \log n})$ , for a good choice of the learning rate  $\eta$ , against the adaptive adversary.

**Solution.** Let us prove the following inequality first:

**Lemma.** Let  $w_0, w_1, \dots, w_n \in [0, 1]^n$  such that  $\sum_i w_i = 1$  then there exists  $\eta$  such that:

$$\sum_i w_i e^{\eta x_i} \leq e^{\eta \sum_i w_i x_i}$$

Define the potential  $\psi(t) = \sum_j e^{\eta G_{j,t}}$  and notice that:

1.  $\frac{\psi(t+1)}{\psi(t)} = \mathbf{E}[e^{\eta g_t}]$
2.  $\frac{\psi(t+1)}{\psi(t)} \leq e^{\eta}$

by the fact that  $e^x$  is positive function we have that  $\psi(t) \geq e^{\eta \max_j G_{j,t}}$ . In addition:

$$\begin{aligned} \frac{\psi(t+1)}{\psi(t)} &= \psi(t)^{-1} \sum_j e^{\eta G_{j,t} + \eta g_{j,t+1}} \\ &= \mathbf{E}[e^{\eta g_t}] \leq e^{\mathbf{E}[\eta g_t]} \Rightarrow \psi(T) \leq \psi(0) e^{\eta \mathbf{E}[G_t]} \\ \psi(t+1) &= \sum_j e^{\eta G_{j,t+1}} = \sum_j e^{\eta G_{j,t} + \eta g_{j,t+1}} \\ &\leq \sum_j e^{\eta G_{j,t}} (1 + c_1 \eta g_{j,t+1}) = \psi(t) + c_1 \eta \psi(t) \mathbf{E}[g_{t+1}] \\ &\leq e^{\eta} (1 + c_2) \psi(t) \mathbf{E}[g_{t+1}] \\ &\leq \prod_t (1 + c_1 \mathbf{E}[g_t]) (\psi(0)) \leq e^{c_2 \mathbf{E}[\sum g_t]} \psi(0) \\ &\leq e^{c_2 \mathbf{E}[G_t]} \cdot e^{\eta n} \end{aligned}$$

So after  $T$  steps, by taking the logarithm of both sides, we obtain that the regret is bounded by  $R_T \leq$ .

**ex2.** Show a lower bound of  $\Omega(\sqrt{T})$  in the experts setting on the regret of any online algorithm against the oblivious adversary.

**Solution.** solution.

**ex3.** Consider a system of linear inequalities  $Ax \geq b$ , where  $A \in [0, \infty]^{m \times n}$ ,  $b \in [0, \infty]^m$ , and unknown  $x \in [0, \infty]^n$ . (we are seeking a non-negative solution). An  $\varepsilon$ -approximate solution  $x \geq 0$  satisfies  $Ax \geq b - \varepsilon \mathbf{1}$ . Suppose we have an efficient procedure for following problem: Given  $p \in [0, 1]^m$ ,  $\sum_{i \in [m]} p_i = 1$ , decide if exists  $x \geq 0$ ,  $p^\top Ax \geq p^\top b$ . Show how to find an  $\varepsilon$ -approximate solution to  $Ax \geq b$ . Analyze the run-time.

**Solution.** solution.

**ex4.** Recall that we showed, for  $EXP$  updates, that w.p  $1 - \delta$

$$RT \leq \beta nT + \gamma T + (1 + \beta) \eta + \frac{\ln(\delta^{-1}n)}{\beta} + \frac{\ln n}{\eta}$$

Infer that for the right choice of  $\beta, \gamma, \eta$

$$\mathbf{E}[R_T] = O(\sqrt{Tn \ln n})$$