## Online Computation, Ex 3.

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ex1. Consider the experts setting with gains:  $g_{i,t} \in [0,1]$  is ex4. Recall that we showed, for EXP updates, that w.p. the gain of expert i at step t. Hedge updates:

$$P_{i,t+1} = \frac{e^{\eta G_{i,t}}}{\sum_{j} e^{\eta G_{j,t}}}$$

where  $G_{i,t} = \sum_{s \leq t} g_{i,t}$ . Prove that the regret of Hedge at time T is  $O(\sqrt{T \log n})$ , for a good choice of the learning rate  $\eta$ , against the adaptive adversary.

**Solution.** Define the potential  $\psi(t) = \sum_{j} e^{\eta G_{i,t}}$  and notice that by the fact that  $e^x$  is positive function we have that  $\psi(t) > e^{\eta \max_{j} G_{i,j}}$ . In addition:

$$\frac{\psi(t+1)}{\psi(t)} = \psi(t)^{-1} \sum_{j} e^{\eta G_{j,t} + \eta g_{j,t+1}}$$

$$= \mathbf{E} \left[ e^{\eta g_{t}} \right] \leq e^{\mathbf{E} \left[ \eta g_{t} \right]} \Rightarrow \psi(T) \leq \psi(0) e^{\eta \mathbf{E} \left[ G_{t} \right]}$$

$$\psi(t+1) = \sum_{j} e^{\eta G_{j,t+1}} = \sum_{j} e^{\eta G_{j,t} + \eta g_{j,t+1}}$$

$$\leq \sum_{j} e^{\eta G_{j,t}} \left( 1 + c_{1} \eta g_{j,t+1} \right) = \psi(t) + c_{1} \eta \psi(t) \mathbf{E} \left[ g_{t+1} \right]$$

$$\leq e^{\eta} \left( 1 + c_{2} \right) \psi(t) \mathbf{E} \left[ g_{t+1} \right]$$

$$\leq \prod_{t} \left( 1 + c_{1} \mathbf{E} \left[ g_{t} \right] \right) \left( \psi(0) \right) \leq e^{c_{2} \mathbf{E} \left[ \sum g_{t} \right]} \psi(0)$$

$$< e^{c_{2} \mathbf{E} \left[ G_{t} \right]} \cdot e^{\eta} \eta$$

So after T steps, by taking the logaritm of both sides, we obtain that the regret is bounded by  $R_T \leq$ .

**ex2.** Show a lower bound of  $\Omega\left(\sqrt{T}\right)$  in the experts setting on the regret of any online algorithm against the oblivious adversary.

**Solution.** solution.

**ex3.** Consider a system of linear inequalities  $Ax \geq b$ , where  $A \in [0,\infty]^{m \times n}, b \in [0,\infty]^m$ , and unknown  $x \in [0,\infty]^n$ . (we are seeking a non-negative solution). An  $\varepsilon$ -approximate solution  $x \geq 0$  satisfies  $Ax \geq b - \varepsilon \mathbf{1}$ . Suppose we have an efficient procedure for following problem: Given  $p \in$  $[0,1]^m, \sum_{i \in [m]} p_i = 1$ , decide if exists  $x \geq 0, p^\top Ax \geq p^\top b$ . Show how to find an  $\varepsilon$ -approximate solution to  $Ax \geq b$ . Analyze the run-time.

**Solution.** solution.

$$RT \le \beta nT + \gamma T + (1+\beta)\eta + \frac{\ln(\delta^{-1}n)}{\beta} + \frac{\ln n}{\eta}$$

Infer that for the right choice of  $\beta$ ,  $\gamma$ ,  $\eta$ 

$$\mathbf{E}\left[R_T\right] = O\left(\sqrt{Tn\ln n}\right)$$