Online Computation, Ex 3.

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ex1. Consider the experts setting with gains: $g_{i,t} \in [0,1]$ is **ex2.** Show a lower bound of $\Omega\left(\sqrt{T}\right)$ in the experts setting the gain of expert i at step t. Hedge updates:

$$P_{i,t+1} = \frac{e^{\eta G_{i,t}}}{\sum_{j} e^{\eta G_{j,t}}}$$

where $G_{i,t} = \sum_{s < t} g_{i,t}$. Prove that the regret of Hedge at time T is $O(\sqrt{T \log n})$, for a good choice of the learning rate η , against the adaptive adversary.

Solution. Let g_t be the random variable which count the gain at time step t and by $G_t = \sum_t^T g_t$. Recall that for any pair of random variable X, Y such that $X \geq Y$ it holds that $\mathbf{E}[X] \geq \mathbf{E}[Y]$. Also notice that for x restricted to some range [-r, r] there are constants c_+, c_- depeand on r such that $c_{-}x^{2} \leq e^{x} - 1 - x \leq c_{+}x^{2}$. Namely, the exponent is bounded by quaderic approximation (second tylor series order). By the montenus property of the expection, for any random variable X that maps to bounded range [-r, r] it holds that:

$$c_{-}\mathbf{E}\left[x^{2}\right] \leq \mathbf{E}\left[e^{x} - x - 1\right] \leq c_{+}\mathbf{E}\left[x^{2}\right]$$

Define the potential $\psi(t) = \sum_{i} e^{\eta G_{i,t}}$ and notice that:

- 1. $\frac{\psi(t+1)}{\psi(t)} = \mathbf{E}\left[e^{\eta g_t}\right]$ relatives to the distribution $P_{i,t+1}$.
- 2. $\psi(t) \geq e^{\eta G_{t,j}}$ for any t and j in particular the j which maximizes the gain.

Therefore we obtain that:

$$\psi\left(T\right) = \frac{\psi\left(T\right)}{\psi\left(0\right)}\psi\left(0\right) = \prod_{t=0}^{T} \frac{\psi\left(t+1\right)}{\psi\left(t\right)}\psi\left(0\right)$$

$$n\prod_{t=0}^{T} \mathbf{E}\left[e^{\eta g_{t}}\right] \le n\prod_{t=0}^{T} \mathbf{E}\left[1+\eta g_{t}+c_{\pm}\left(\eta g_{t}\right)^{2}\right]$$

$$n\prod_{t=0}^{T} 1+\mathbf{E}\left[\eta g_{t}+c_{\pm}\left(\eta g_{t}\right)^{2}\right] \le n\prod_{t=0}^{T} e^{\mathbf{E}\left[\eta g_{t}+c_{\pm}\left(\eta g_{t}\right)^{2}\right]} \le ne^{\mathbf{E}\left[\sum \eta g_{t}+c_{\pm}\left(\eta g_{t}\right)^{2}\right]} \le ne^{\mathbf{E}\left[\sum \eta g_{t}+c_{\pm}\left(\eta g_{t}\right)^{2}\right]}$$

On the other hand by the second property it follows that for any j:

$$e^{\eta G_{j,T}} < ne^{\mathbf{E}\left[\sum \eta g_t\right] + \mathbf{E}\left[c_{\pm}(\eta g_t)^2\right]}$$

By dividing at $e^{\mathbf{E}[\sum \eta g_t]}$, extracting the logarithm and combine the fact that $g_t^2 = g_t$ (indicator) we have that:

$$R_T \le \frac{1}{n} \log\left(n\right) + c_+ \eta T$$

And by choosing $\eta = \sqrt{\log(n)/T}$ we complete the proof.

on the regret of any online algorithm against the oblivious adversary.

Solution. Consider an adversarial which draw the values of $g_{i,t}$ uniformly random, in particular $g_{i,t}$'s are independent. Fix an online algorithm for the problem and denote by g_t the gain that earn by it at time step t. As $g_{i,t}$ are independent, the sum $G_T = \sum g_t$ is a sumation of indeapndent variables with the same expection and variance. Therefore we know that $(G_T - T\mu)/\sqrt{T} \sim G(0, \sigma)$ where μ and σ do not deapnds on T. In the otherhand run in which the optimal gain $T(\mu + \frac{1}{2})$ might occurred with positive probability (if $\mu > \frac{1}{2}$). Hence, the regret is at least:

$$R_T \ge T\left(\frac{1}{2} + \mu\right) - \mathbf{E}\left[G_t\right]$$

ex3. Consider a system of linear inequalities $Ax \geq b$, where $A \in [0,\infty]^{m\times n}, b \in [0,\infty]^m$, and unknown $x \in [0,\infty]^n$. (we are seeking a non-negative solution). An ε -approximate solution $x \geq 0$ satisfies $Ax \geq b - \varepsilon \mathbf{1}$. Suppose we have an efficient procedure for following problem: Given $p \in$ $[0,1]^m, \sum_{i\in[m]} p_i = 1$, decide if exists $x \geq 0, p^\top Ax \geq p^\top b$. Show how to find an ε -approximate solution to $Ax \geq b$. Analyze the run-time.

Solution. solution.

ex4. Recall that we showed, for EXP updates, that w.p.

$$RT \le \beta nT + \gamma T + (1+\beta)\eta + \frac{\ln(\delta^{-1}n)}{\beta} + \frac{\ln n}{\eta}$$

Infer that for the right choice of β , γ , η

$$\mathbf{E}\left[R_T\right] = O\left(\sqrt{Tn\ln n}\right)$$