1 Ex 1.

Claim 1. *Let A be a random matrix in $M(\mathbb{F}_2^{k\times n})$ then for any non zero $x \in \mathbb{F}$ we have that Ax distributed uniformly.*

Proof. By the fact that $x \neq 0$ there exists at least one coordinate $i \in [k]$ such that $x_i \neq 0$. Thus we have

$$\begin{aligned}(Ax)_j &= \sum_k A_{jk}x_k = \sum_{i\ neqk} A_{jk}x_k + A_{ji}x_i \\ &= \sum_{i\ neqk} A_{jk}x_k + A_{ji}\end{aligned}$$

Notice that due to the fact that $\mathbb{F}_2$ is a field, there is exactly one assignment that satisfies the equation conditioned on all the values A_{jk} where $j \neq k$.

$$\begin{aligned}\mathbf{Pr}\left[(Ax)_j = 1\right] &= \sum_{A_{jk};k\neq i} \mathbf{Pr}\left[(Ax)_j = 1 | A_{jk}; k \neq i\right] \mathbf{Pr}\left[A_{jk}; k \neq i\right] \\ &= \frac{1}{2}\end{aligned}$$

therefore any coordinate of Ax distributed uniformly $\Rightarrow$ Ax distributed uniformly. $\square$

By the uniformity of Ax we obtain that the expected Hamming wight of Ax is :

$$\mathbf{E}\left[|Ax|\right] = \mathbf{E}\left[\sum_i^n (Ax)_i\right] = \frac{1}{2}n$$

As the coordinates of A_x are independent (each row of A is sampled separately) we can use the Hoff' bound to conclude that:

$$\mathbf{Pr}\left[\left||Ax| - \mathbf{E}\left[|Ax|\right]\right| \geq \left(\frac{1}{2} - \delta\right) n\right] \leq e^{-n\left(\frac{1}{2}-\delta\right)^2}$$

Now we will use the union bound to show that any $x \in \mathbb{F}_2^k$, Ax is at weight at least δ.

$$\mathbf{Pr}\left[|Ax| \geq \delta : \forall x \in \mathbb{F}_2^k\right] \geq 1 - |\mathbb{F}_2^k| \cdot e^{-n\left(\frac{1}{2}-\delta\right)^2}$$

Denote $k = \rho n$ and notice that the above holds when $\rho \geq \left(\frac{1}{2} - \delta\right)^2$

2 Ex 2.

Claim 2. *Let $v_1, v_2..v_m$ unit vectors in an inner-product space such that $\langle v_i, v_j\rangle \leq -2\varepsilon$ for all $i \neq j$, then $m \leq \frac{1}{2\varepsilon} + 1$.*

Proof. Let's us bound form both sides the norm of the summation $|\sum_i v_i|$. As the norm is by definition (construction) non-negative we are going to bound from the left by 0, on the other hand we have that:

$$0 \leq |\sum_{v_i} v_i| = m + 2\sum_{i,j} \langle v_i, v_j\rangle \leq m - 2 \cdot \frac{m(m-1)}{2} \cdot 2\varepsilon$$

Thus we obtain $m\left(2(m-1)\varepsilon - 1\right) \leq 0$ namely, $m \leq \frac{1}{2\varepsilon} + 1$ $\square$

Now, define the following product for $u, v \in \mathbb{F}_2^n$, $\langle v, u\rangle = \sum_i (-1)^{v_i}(-1)^{\bar{u}_i}$ observes that:

1. $\langle v, v\rangle = \sum_i 1 = n \geq 0$.
2. $\langle v, u\rangle = \langle u, v\rangle$.
3. $\langle ax + by, z\rangle = (-1)^a \langle x, z\rangle + (-1)^b \langle y, z\rangle$.

Now the v's corresponds to code with distance at least d then, i.e for any codewords v and u disagree on at least d coordinates, and therefore $\langle v, u\rangle \leq$ `agree`$-$ `disagree = n - 2 disagree` $= n - 2d$. Now consider the normal codewords $\tilde{v_1}..\tilde{v_n}$ and assume that

$$\langle \tilde{v_i}, \tilde{v_j}\rangle = (1 - 2\delta) = \frac{1}{n}\left(n - 2d(v_i, v_j)\right) \leq \varepsilon$$

So if $d \geq \frac{1}{2} + \varepsilon$ we obtain the condition of the above claim.

3 Ex 3.

Consider the following process for decoding a,

for $t \in [\tau]$ **do**
  for $i \in [n]$ **do**
    $x \sim_u \mathbb{F}_2^n$
    $a_i^{(t)} \leftarrow w\left(x\right) + w\left(\sigma_i(x)\right)$
  end
end
for $i \in [n]$ **do**
  $\hat{a}_i \leftarrow [\frac{1}{\tau}\sum_t^\tau a_i^{(t)}]$
end
return $\hat{a}_0, \hat{a}_1, \hat{a}_2..\hat{a}_n$

Claim 3. *For $\tau = \Omega\left(\frac{1}{\varepsilon^4}\log\left(n\right)\right)$ The above decoding success to decode $w\left(x\right)$ with probability $\geq 1 - \frac{1}{n}$.*

Proof. In this question we will say that w agree on $x, \sigma_i(x)$ if both $x, \sigma_i(x)$ were either filliped or unflipped. Clearly if $w(x)$ agree with $w(\sigma_i(x))$ than

$$\begin{aligned} w\left(x\right) + w\left(\sigma_i(x)\right) &= H_a(x) + H_a(\sigma_i(x)) \\ &= \sum_{i \neq j} a_j(x_j + x_j) + a_i(x_j + 1 + x_j) = a_j \quad (\text{ neither of them were flipped. }) \\ &= 1 + H_a(x) + 1 + H_a(\sigma_i(x)) = a_i \quad (\text{ both flipped.}) \end{aligned}$$

Thus we can bound the probability that $w\left(x\right) + w\left(\sigma_i(x)\right) \neq a_{_i}$ by the probability that w disagree on x and that append at probability:

$$\xi := \left(1 - f(x)\right) f(\sigma_i(x)) + \left(1 - f(\sigma_i(x))\right) f(x)$$

Now as we want bound ξ we could think about the maximization problem under the restrictions that $f(x), f(\sigma_i(x) \leq \frac{1}{2} - \varepsilon$. We know that the maximum lay on the boundary so we can assign $\frac{1}{2} - \varepsilon$ for each of the probabilities to obtain an upper bound. That will yield $\xi \leq 2 \cdot \left(\frac{1}{2} - \varepsilon\right)\left(\frac{1}{2} + \varepsilon\right)$, namely $\xi \leq \frac{1}{2} - 2\varepsilon^2$. Now the probability that a coordinate i will rounded to the opposite side, that it $\hat{a}_i \neq a_i$ mean that arithmetic mean over τ experiments were $2\varepsilon^2$ far from the expectation. Which by Hoff' bound is bounded by: $e^{\tau 4\varepsilon^4}$. So using the union bound we obtain:

$$\mathbf{Pr}\left[\text{ decoding success }\right] \geq 1 - n \cdot e^{\tau 4\varepsilon^4}$$

Therefore it's enough to take $\tau = O(\frac{1}{\varepsilon^4}\log(n))$ to obtain a decoder which run at time $O(\frac{1}{\varepsilon^4} n \log(n))$ and success with heigh probability. $\square$

4 Ex 4.

4.1 (a)

We will prove that if for any x, f interpolate well on $x, x+1, .., x+d+1$ than any it interpolate well on every coordinates set at size $d+1$. Denote by $J \subset \mathbb{F}_q$ at size $d+1$. Let's continue by induction on $\max J$. The base case $\max J = d+1 \Rightarrow J = \{1, 2.., d+1\}$ follow straightforwardly from the assumption. Assume the correctness for any J such that $\max J \leq x_0$ and consider J' such that $\max J' = x_0+$. Now it given that $S = \{x_0 - d, x_0 - d + 1, ..x_0 + 1\}$ is well interpolating set, so there exists coefficients a_1, a_{d+1} such that $a_{d+1} f(x_0+1) = \sum_{x_i \in S/(x_0+1)} a_i f(x_i)$ On the overhand, for ant any $x_i \in S/(x_0+1)$ the union $K = x_i \cup J/(x_0+1)$ is subset of $\mathbb{F}_q$ at size $d+1$ such that $\max K \leq x_0$. Hence by induction assumption K is well interpolating set and we can exchange any $f(x_i)$ for $x_i \in S$ by a linear combination of $f(x_i)$ for $x_i \in J/(x_0+1)$. So in overall we obtain that J is depended set, namely f is well interpolate on J.

4.2 (b)

We will prove the claim only for monoms $x^l \mapsto (tx+s)^l$ and that will sufficent by lineratiy arguments. Again we will show correction by induction on l, if $l = 1$ then it's trivial. Suppose the rightness of the claim for any degree lower then l and consider the monom x^{l+1}. Note that for any $t \neq 0, s \in \mathbb{F}_q$ we have that $(tx+s)^{l+1} = (tx+s)^l (tx+s) = s(tx+s)^l + tx(tx+s)^l$. Now applay the induction assumption to obtain that:

$$a_{d+1}(tx+s)^l(x_{d+1}) = \sum_i a_i (tx+s)^l (x_i) \Rightarrow \left(tx(tx+s)^l\right)(x_i) = tx_i \left(a_i^{-1}\left(a_{d+1} - a_j(\cdot)\right)\right)$$

¡++¿ Note that q is prime, thus $(\mathbb{F}_q/0, \cdot)$ and $(\mathbb{F}_q, +)$ are cyclic groups and therefore any $s \in (\mathbb{F}_q, +)$ and a subset of $J \subset \mathbb{F}_q$ of $d+1$ coordinates the set $s + J$ is also at size $d+1$ (no collision occurs). Therefore we could find

Now for any set $\{x_1 + s, x_2 + s, .., x_{d+1} + s\}$ we could find the interpolation coffetions relative to the function