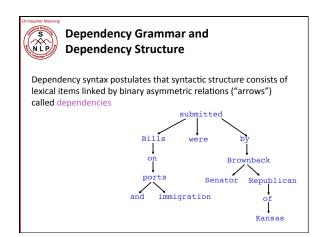


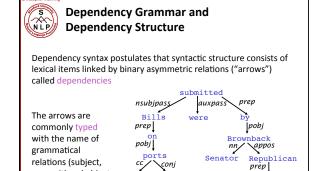
prepositional object,

apposition, etc.)

#### Dependency Grammar

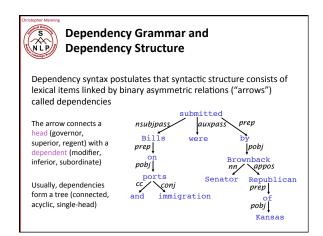
Introduction

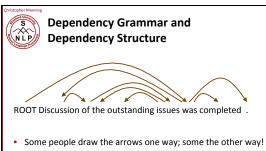




immigration

pobj |





- Tesnière had them point from head to dependent...
- Usually add a fake ROOT so every word is a dependent of precisely 1 other node



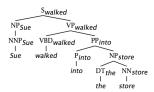
#### **Dependency Grammar/Parsing History**

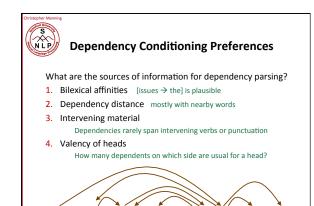
- The idea of dependency structure goes back a long way
  - To Pāṇini's grammar (c. 5th century BCE)
  - Basic approach of 1st millennium Arabic grammarians
- Constituency is a new-fangled invention
  - 20th century invention (R.S. Wells, 1947)
- Modern dependency work often linked to work of L. Tesnière (1959)
  - Was dominant approach in "East" (Russia, China, ...)
  - Good for free-er word order languages
- Among the earliest kinds of parsers in NLP, even in the US:
  - David Hays, one of the founders of U.S. computational linguistics, built early (first?) dependency parser (Hays 1962)



## Relation between phrase structure and dependency structure

- A dependency grammar has a notion of a head. Officially, CFGs don't.
- But modern linguistic theory and all modern statistical parsers (Charniak, Collins, Stanford, ...) do, via hand-written phrasal "head rules":
  - The head of a Noun Phrase is a noun/number/adj/...
  - The head of a Verb Phrase is a verb/modal/....
- The head rules can be used to extract a dependency parse from a CFG parse
- The closure of dependencies give constituency from a dependency tree
- But the dependents of a word must be at the same level (i.e., "flat") – there can be no



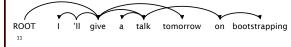


ROOT Discussion of the outstanding issues was completed  $% \left( x\right) =\left( x\right) +\left( x\right) +\left($ 



#### **Dependency Parsing**

- A sentence is parsed by choosing for each word what other word (including ROOT) that it is a dependent of.
- Usually some constraints:
  - Only one word is a dependent of ROOT
  - Don't want cycles A → B, B → A
- This makes the dependencies a tree
- Final issue is whether arrows can cross (non-projective) or not





#### **Methods of Dependency Parsing**

- Dynamic programming (like in the CKY algorithm)
   You can do it similarly to lexicalized PCF6 parsing: an O(n<sup>5</sup>) algorithm
   Eisner (1996) gives a clever algorithm that reduces the complexity to O(n<sup>3</sup>), by producing parse items with heads at the ends rather than in the middle
- 2. Graph algorithms

You create a Minimum Spanning Tree for a sentence McDonald et al.'s (2005) MSTParser scores dependencies independently using a ML classifier (he uses MIRA, for online learning, but it could be something else)

3. Constraint Satisfaction

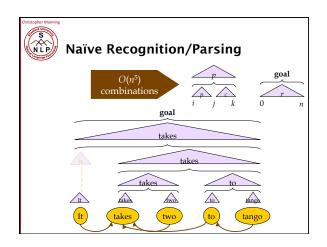
Edges are eliminated that don't satisfy hard constraints. Karlsson (1990), etc.

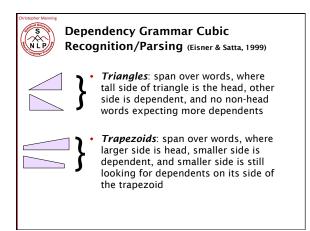
"Deterministic parsing"
 Greedy choice of attachments guided by good machine learning classifiers MaltParser (Nivre et al. 2008)

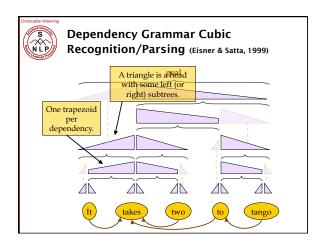


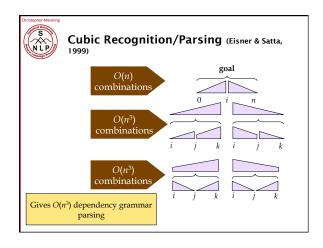
# DP for generative dependency grammars

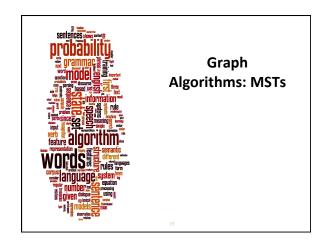
# Probabilistic dependency grammar: generative model 1. Start with left wall \$ 2. Generate root $w_0$ 3. Generate left children $w_1$ , $w_2$ , ..., $w_\ell$ from the FSA $\lambda w_0$ 4. Generate right children $w_1$ , $w_2$ , ..., $w_\ell$ from the FSA $\rho w_0$ 5. Recurse on each $w_\ell$ for $\ell$ in $\ell$ $\ell$ , ..., $\ell$ , $\ell$ , ..., $\ell$ , sampling of (steps 2-4) 6. Return $\alpha_\ell$ ... $\alpha_1 w_0 \alpha_1$ ... $\alpha_r$ These 5 slides are based on slides by Jason Eisner and Noah Smith

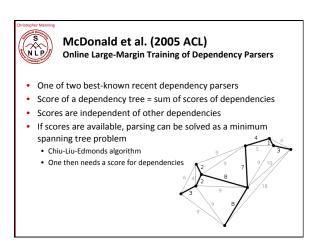














### McDonald et al. (2005 ACL): Online Large-Margin Training of Dependency Parsers

- Edge scoring is via a discriminative classifier
  - Can condition on rich features in that context
  - Each dependency is a linear function of features times weights
- Feature weights were learned by MIRA, an online large-margin algorithm
  - But you could use an SVM, maxent, or a perceptron
- Features cover:
  - · Head and dependent word and POS separately
  - · Head and dependent word and POS bigram features
  - · Words between head and dependent
  - · Length and direction of dependency



#### Greedy Transition-Based Parsing

MaltParser



#### MaltParser

[Nivre et al. 2008]

- A simple form of greedy discriminative dependency parser
- The parser does a sequence of bottom up actions
- Roughly like "shift" or "reduce" in a shift-reduce parser, but the "reduce" actions are specialized to create dependencies with head on left or right
- The parser has:
  - a stack σ, written with top to the right
  - which starts with the ROOT symbol
  - a buffer β, written with top to the left
     which starts with the input sentence
  - a set of dependency arcs A
  - · which starts off empty
  - a set of actions



#### Basic transition-based dependency parser

Start:  $\sigma = [ROOT], \beta = w_1, ..., w_n, A = \emptyset$ 

1. Shift  $\sigma, w_i | \beta, A \rightarrow \sigma | w_i, \beta, A$ 

2. Left-Arc<sub>r</sub>  $\sigma|w_i, w_j|\beta, A \rightarrow \sigma, w_j|\beta, A \cup \{r(w_i, w_i)\}$ 

3. Right-Arc<sub>r</sub>  $\sigma | w_i, w_j | \beta, A \rightarrow \sigma, w_i | \beta, A \cup \{r(w_i, w_j)\}$ 

Finish:  $\beta = \emptyset$ 

#### Notes:

 Unlike the regular presentation of the CFG reduce step, dependencies combine one thing from each of stack and buffer



#### Actions ("arc-eager" dependency parser)

Start:  $\sigma = [ROOT]$ ,  $\beta = w_1, ..., w_n$ ,  $A = \emptyset$ 

**1.** Left-Arc<sub>r</sub>  $\sigma(w_i, w_j | \beta, A \rightarrow \sigma, w_j | \beta, A \cup \{r(w_j, w_i)\}$ Precondition:  $r'(w_k, w_i) \notin A, w_i \neq ROOT$ 

2. Right-Arc<sub>r</sub>  $\sigma|w_i, w_i|\beta$ , A  $\rightarrow \sigma|w_i|w_i, \beta$ , A  $\cup \{r(w_i, w_i)\}$ 

3. Reduce  $\sigma | w_i, \beta, A \rightarrow \sigma, \beta, A$ Precondition:  $r'(w_k, w_i) \in A$ 

4. Shift  $\sigma, w_i | \beta, A \rightarrow \sigma | w_i, \beta, A$ 

Finish:  $\beta = \emptyset$ 

This is the common "arc-eager" variant: a head can immediately take a right dependent, before its dependents are found



#### Example

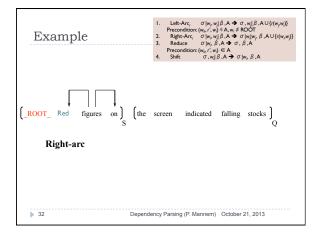
1. Left-Arc,  $\sigma(w_p, w_p|\beta, A \rightarrow \sigma, w_p|\beta, A \cup \{r(w_p, w_p)\}$ Precondition:  $(w_p, r', w_p) \in A, w_p \in ROOT$ 2. Right-Arc,  $\sigma(w_p, w_p|\beta, A \rightarrow \sigma, g_h)$ 3. Reduce  $\sigma(w_p, \beta, A \rightarrow \sigma, \beta, A)$ Precondition:  $(w_p, r', w_p) \in A$ 1. Shift  $\sigma(w_p|\beta, A \rightarrow \sigma) \in W_p, \beta, A$ 

Happy children like to play with their friends .

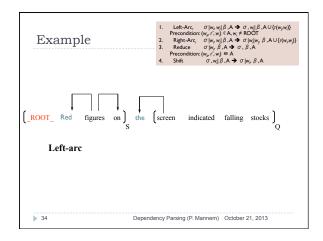
[ROOT] [Happy, children, ...] Shift [ROOT, Happy] [children, like, ...]  $\mathsf{LA}_{amod}$ [ROOT] [children, like, ...]  $\{amod(children, happy)\} = A_1$ Shift [ROOT, children] [like, to, ...]  $A_1 \cup \{nsubj(like, children)\} = A_2$  $\mathsf{LA}_{\mathit{nsubj}}$ [ROOT] [like. to. ...]  $\mathsf{RA}_{root}$  $A_2 \cup \{root(ROOT, like) = A_3\}$ [ROOT, like] [to, play, ...] [ROOT, like, to] [play, with, ...] Shift  $\mathsf{LA}_{\mathit{aux}}$ [ROOT, like] [play, with, ...]  $A_3 \cup \{aux(play, to) = A_4\}$ RA<sub>xcomp</sub> [ROOT, like, play] [with their, ...]  $A_4 \cup \{xcomp(like, play) = A_5\}$ 

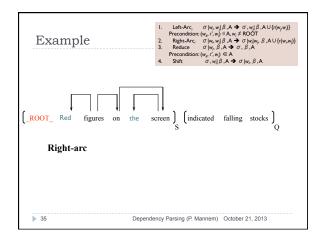
```
 \begin{split} 1. & \quad \text{Left-Arc}, \quad \sigma(w_s, w_s|\beta, A \to \sigma, w_s|\beta, A \cup \{r(w_s, w_s)\} \\ & \quad \text{Precondition:} & \langle w_s, r', w \rangle \in A, w_s \neq ROOT \\ 2. & \quad \text{Right-Arc}, & \sigma(w_s, w_s|\beta, A \to \sigma) \in w_s \mid w_s, \beta, A \cup \{r(w_s, w_s)\} \\ 3. & \quad \text{Reduce} \quad \sigma(w_s, f', w_s) \in A, \\ & \quad \text{Precondition:} & \langle w_s, r', w_s \rangle \in A, \\ & \quad \text{A}. & \quad \text{Shift} & \quad \sigma_s w_s|\beta, A \to \sigma|w_s, \beta, A \\ & \quad \text{A}. & \quad \text{Shift} & \quad \sigma_s w_s|\beta, A \to \sigma|w_s, \beta, A \\ \end{aligned} 
                            Example
 Happy children like to play with their friends .
                                                                      [with their, ...] A_4 \cup \{xcomp(like, play) = A_5\}
RA<sub>xcomp</sub> [ROOT, like, play]
RA<sub>prep</sub> [ROOT, like, play, with] [their, friends, ...] A_5 \cup \{\text{prep}(\text{play}, \text{with}) = A_6 \}
Shift [ROOT, like, play, with, their] [friends, .] A_6 \cup \{\text{prep}(\text{play}, \text{with}) = A_6 \}
LA_{poss} [ROOT, like, play, with] [friends, .] A_6 \cup \{poss(friends, their) = A_7\}
RA<sub>pobj</sub> [ROOT, like, play, with, friends] [.]
                                                                                                         A_7 \cup \{pobj(with, friends) = A_8
 Reduce [ROOT, like, play, with] [.]
                                                                                                               A<sub>8</sub>
Reduce [ROOT, like, play] [.]
                                                                                                              A_8
Reduce [ROOT, like]
                                                  [.]
[]
                                                                                                               A_8
 RA<sub>punc</sub> [ROOT, like, .]
                                                                                                             A_8 U \{punc(like, .) = A_9
You terminate as soon as the buffer is empty. Dependencies = A<sub>9</sub>
```

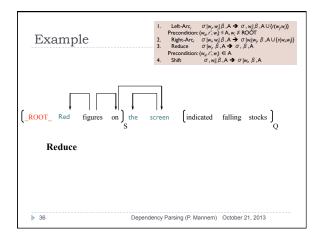
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Example \begin{tabular}{ll} $\mathbb{E}$ & & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ &
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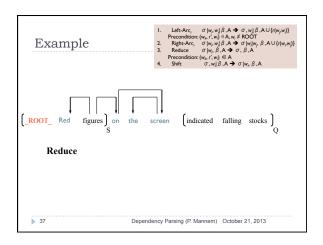


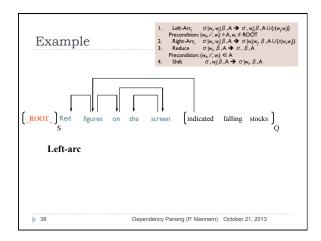
```
Example \begin{tabular}{lll} \hline \textbf{Example} \\ \hline & \textbf{Example} \\ \hline & \textbf{I. Left-Arc.} & \sigma | \textbf{w}, \textbf{w} | \beta, \textbf{A} \Rightarrow \sigma, \textbf{w} | \beta, \textbf{A} \cup \{r(\textbf{w}, \textbf{w})\} \\ & \textbf{Precondition:} (\textbf{w}, \textbf{r}', \textbf{w}) & \beta, \textbf{A} \neq \beta, \textbf{A} \cup \{r(\textbf{w}, \textbf{w})\} \\ & \textbf{2. Right-Arc.} & \sigma | \textbf{w}, \textbf{w} | \beta, \textbf{A} \Rightarrow \sigma | \textbf{w}, \beta, \textbf{A} \cup \{r(\textbf{w}, \textbf{w})\} \\ & \textbf{3. Reduce} & \sigma | \textbf{w}, \beta, \textbf{A} \Rightarrow \sigma, \beta, \textbf{A} \\ & \textbf{Precondition:} (\textbf{w}, \textbf{r}', \textbf{w}) \in \textbf{A} \\ & \textbf{4. Shift} & \sigma, \textbf{w} | \beta, \textbf{A} \Rightarrow \sigma | \textbf{w}, \beta, \textbf{A} \rightarrow \sigma | \textbf{w}, \beta, \textbf{A} \\ & \textbf{Shift} & \sigma, \textbf{w} | \beta, \textbf{A} \Rightarrow \sigma | \textbf{w}, \beta, \textbf{A} \rightarrow \sigma |
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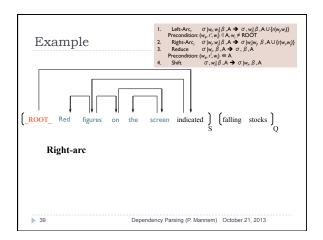


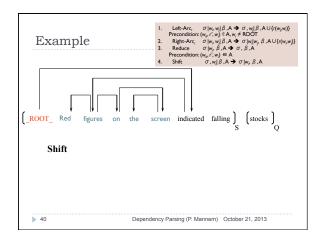


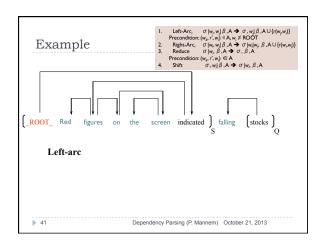


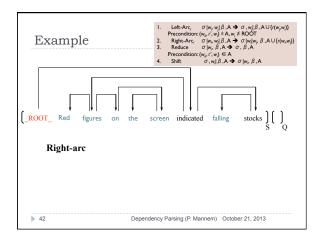


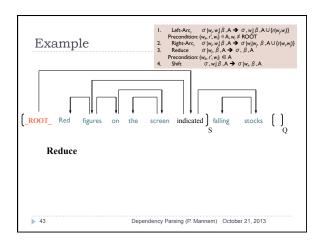


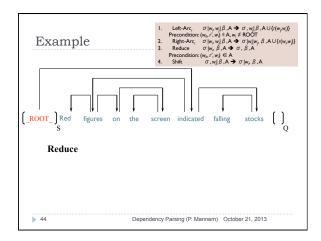










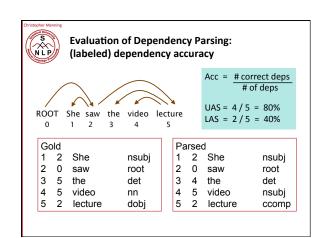




#### MaltParser

[Nivre et al. 2008]

- We have left to explain how we choose the next action
- Each action is predicted by a discriminative classifier (often SVM, can be perceptron, maxent classifier) over each legal move
  - Max of 4 untyped choices, max of |R| × 2 + 2 when typed
  - · Features: top of stack word, POS; first in buffer word, POS; etc.
- There is NO search (in the simplest and usual form)
  - · But you could do some kind of beam search if you wish
- · It provides VERY fast linear time parsing
- The model's accuracy is slightly below the best Lexicalized PCFGs (evaluated on dependencies), but
- It provides close to state of the art parsing performance





#### Representative performance numbers

- The CoNLL-X (2006) shared task provides evaluation numbers for various dependency parsing approaches over 13 languages
  - MALT: LAS scores from 65–92%, depending greatly on language/treebank
- Here we give a few UAS numbers for English to allow some comparison to constituency parsing

Parser	UAS%
Sagae and Lavie (2006) ensemble of dependency parsers	92.7
Charniak (2000) generative, constituency, as dependencies	92.2
Collins (1999) generative, constituency, as dependencies	91.7
McDonald and Pereira (2005) – MST graph-based dependency	91.5
Yamada and Matsumoto (2003) – transition-based dependency	90.4



#### **Projectivity**

- Dependencies from a CFG tree using heads, must be projective
  - There must not be any crossing dependency arcs when the words are laid out in their linear order, with all arcs above the words.
- But dependency theory normally does allow non-projective structures to account for displaced constituents
  - You can't easily get the semantics of certain constructions right without these nonprojective dependencies

Who did Bill buy the coffee from yesterday?



#### Handling non-projectivity

- The arc-eager algorithm we presented only builds projective dependency trees
- Possible directions to head:
  - 1. Just declare defeat on nonprojective arcs
  - 2. Use a dependency formalism which only admits projective representations (a CFG doesn't represent such structures...)
  - Use a postprocessor to a projective dependency parsing algorithm to identify and resolve nonprojective links
  - 4. Add extra types of transitions that can model at least most nonprojective structures
  - Move to a parsing mechanism that does not use or require any constraints on projectivity (e.g., the graph-based MSTParser)



## Dependencies encode relational structure

Relation Extraction with Stanford Dependencies

