



```
graph TD
    submitted --> Bills
    submitted --> were
    submitted --> by
    Bills --> on
    on --> ports
    ports --> and
    ports --> immigration
    by --> Brownback
    Brownback --> Senator
    Brownback --> Republican
    Republican --> of
    of --> Kansas
```

submitted  
nsubjpass Bills  
prep on  
pobj ports  
cc and  
conj immigration  
auxpass were  
by  
pobj Brownback  
nn Senator  
appos Republican  
prep of  
pobj Kansas

submitted  
nsubjpass Bills  
auxpass were  
prep by  
pobj Brownback  
appos nn Senator  
prep of  
pobj Kansas  
cc ports  
conj and  
immigration  
Republican

- The idea of dependency structure goes back a long way
  - To Pāṇini's grammar (c. 5th century BCE)
  - Basic approach of 1st millennium Arabic grammarians
- Constituency is a new-fangled invention
  - 20th century invention (R.S. Wells, 1947)
- Modern dependency work often linked to work of L. Tesnière (1959)
  - Was dominant approach in "East" (Russia, China, ...)
    - Good for free-er word order languages
- Among the earliest kinds of parsers in NLP, even in the US:
  - David Hayes, one of the founders of U.S. computational linguistics, built early (first?) dependency parser (Hayes 1962)

# Relationship between phrase structure and dependency structure

- A dependency grammar has a notion of a head. Officially, CFGs don't.
- But modern linguistic theory and all modern statistical parsers (Charniak, Collins, Stanford, ...) do, via hand-written phrasal "head rules":
  - The head of a Noun Phrase is a noun/number/adj/...
  - The head of a Verb Phrase is a verb/modal/...
- The head rules can be used to extract a dependency parse from a CFG parse


The closure of dependencies give constituency from a dependency tree

But the dependents of a word must be at the same level (i.e., "flat") – there can be no VP!

```

graph TD
    swalked --> NP_Sue[NP Sue]
    swalked --> VP_walked[VP walked]
    NP_Sue --> NNP_Sue[NNP Sue]
    VP_walked --> VBD_walked[VBD walked]
    VP_walked --> PP_into[PP into]
    VBD_walked --> Sue1[Sue]
    VBD_walked --> walked[walked]
    PP_into --> P_into[P into]
    PP_into --> NP_store[NP store]
    P_into --> into[into]
    NP_store --> DT_the[DT the]
    NP_store --> NN_store[NN store]
    DT_the --> the[the]
    NN_store --> store[store]
  
```

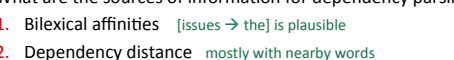
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## Dependency Conditioning Preferences

What are the sources of information for dependency parsing?

1. Bilexical affinities [issues → the] is plausible
2. Dependency distance mostly with nearby words
3. Intervening material  
Dependencies rarely span intervening verbs or punctuation
4. Valency of heads  
How many dependents on which side are usual for a head?



ROOT Discussion of the outstanding issues was completed .


# Dependency Parsing

- A sentence is parsed by choosing for each word what other word (including ROOT) that it is a dependent of.
- Usually some constraints:
  - Only one word is a dependent of ROOT
  - Don't want cycles  $A \rightarrow B, B \rightarrow A$
- This makes the dependencies a tree
- Final issue is whether arrows can cross (non-projective) or not

Diagram illustrating dependency parsing for the sentence "I'll give a talk tomorrow on bootstrapping". The words are: ROOT, I, 'll, give, a, talk, tomorrow, on, bootstrapping. Arrows show dependencies: ROOT to 'll, I to 'll, 'll to give, give to a, a to talk, talk to tomorrow, tomorrow to on, and on to bootstrapping. The arrows for 'll to give and talk to tomorrow cross, indicating a non-projective dependency.

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# Methods of Dependency Parsing

1. Dynamic programming (like in the CKY algorithm)

You can do it similarly to lexicalized PCFG parsing: an  $O(n^3)$  algorithm  
Eisner (1996) gives a clever algorithm that reduces the complexity to  $O(n^3)$ , by producing parse items with heads at the ends rather than in the middle

2. Graph algorithms

You create a Minimum Spanning Tree for a sentence  
McDonald et al.'s (2005) MSTParser scores dependencies independently using a ML classifier (he uses MIRA, for online learning, but it could be something else)


3. Constraint Satisfaction

Edges are eliminated that don't satisfy hard constraints. Karlsson (1990), etc.

4. "Deterministic parsing"


Greedy choice of attachments guided by good machine learning classifiers  
MaltParser (Nivre et al. 2008)

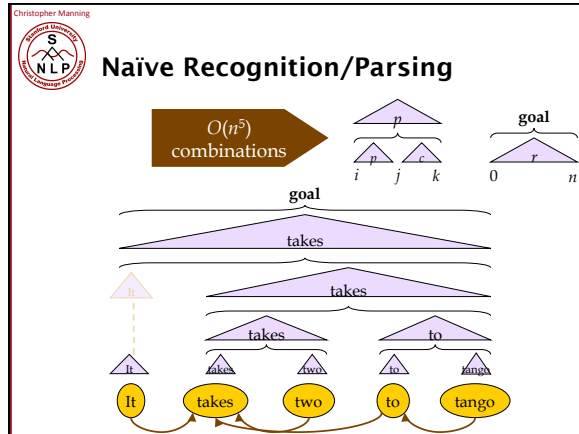
[illegible]




## Probabilistic dependency grammar: generative model


1. Start with left wall \$
2. Generate root  $w_0$
3. Generate left children  $w_{-1}, w_{-2}, \dots, w_{-\ell}$  from the FSA  $\lambda_{w_0}$
4. Generate right children  $w_1, w_2, \dots, w_r$  from the FSA  $\rho_{w_0}$
5. Recurse on each  $w_i$  for  $i$  in  $\{-\ell, \dots, -1, 1, \dots, r\}$ , sampling  $\alpha$  (steps 2-4)
6. Return  $\alpha_{-\ell} \dots \alpha_{-1} w_0 \alpha_1 \dots \alpha_r$








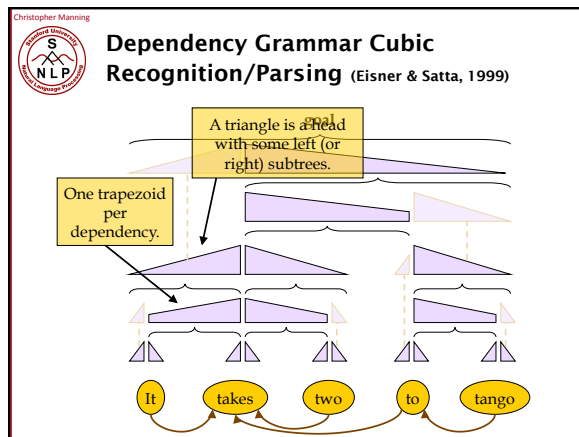
## Dependency Grammar Cubic Recognition/Parsing (Eisner & Satta, 1999)




• ***Triangles***: span over words, where tall side of triangle is the head, other side is dependent, and no non-head words expecting more dependents



• ***Trapezoids***: span over words, where larger side is head, smaller side is dependent, and smaller side is still looking for dependents on its side of the trapezoid



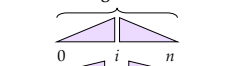


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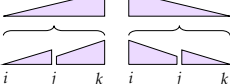
## Cubic Recognition/Parsing (Eisner & Satta, 1999)

$O(n)$   
combinations

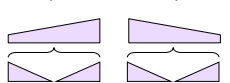
goal



$O(n^2)$   
combinations




$O(n^3)$   
combinations



Gives  $O(n^3)$  dependency grammar parsing

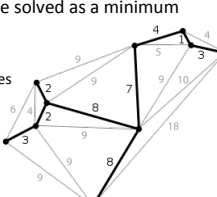
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# McDonald et al. (2005 ACL)

## Online Large-Margin Training of Dependency Parsers

- One of two best-known recent dependency parsers
- Score of a dependency tree = sum of scores of dependencies
- Scores are independent of other dependencies
- If scores are available, parsing can be solved as a minimum spanning tree problem
  - Chiu-Liu-Edmonds algorithm
  - One then needs a score for dependencies






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# McDonald et al. (2005 ACL):

## Online Large-Margin Training of Dependency Parsers

- Edge scoring is via a discriminative classifier
  - Can condition on rich features in that context
  - Each dependency is a linear function of features times weights
- Feature weights were learned by MIRA, an online large-margin algorithm
  - But you could use an SVM, maxent, or a perceptron
- Features cover:
  - Head and dependent word and POS separately
  - Head and dependent word and POS bigram features
  - Words between head and dependent
  - Length and direction of dependency

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# Basic transition-based dependency parser


**Start:**  $\sigma = [\text{ROOT}]$ ,  $\beta = w_1, \dots, w_n$ ,  $A = \varnothing$

1. Shift  $\sigma, w_i | \beta, A \rightarrow \sigma | w_i, \beta, A$
2. Left-Arc<sub>c</sub>  $\sigma | w_i, w_j | \beta, A \rightarrow \sigma, w_i | \beta, A \cup \{r(w_i, w_j)\}$
3. Right-Arc<sub>c</sub>  $\sigma | w_i, w_j | \beta, A \rightarrow \sigma, w_i | \beta, A \cup \{r(w_i, w_j)\}$

**Finish:**  $\beta = \varnothing$

Notes:

- Unlike the regular presentation of the CFG reduce step, dependencies combine one thing from each of stack and buffer



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# Chapter 10

## Dependency Parsing


### Actions (“arc-eager” dependency parser)

**Start:**  $\sigma = [\text{ROOT}]$ ,  $\beta = w_1, \dots, w_n$ ,  $A = \varnothing$

1. Left-Arc<sub>r</sub>     $\sigma | w_i, w_j | \beta$ ,  $A \rightarrow \sigma, w_j, A \cup \{r(w_i, w_j)\}$   
 Precondition:  $r' (w_i, w_j) \notin A$ ,  $w_i \neq \text{ROOT}$
2. Right-Arc<sub>r</sub>     $\sigma | w_i, w_j | \beta$ ,  $A \rightarrow \sigma | w_i, w_j, \beta$ ,  $A \cup \{r(w_i, w_j)\}$
3. Reduce     $\sigma | w_i, \beta$ ,  $A \rightarrow \sigma, \beta$ ,  $A$   
 Precondition:  $r' (w_i, w_i) \in A$
4. Shift     $\sigma, w_i | \beta$ ,  $A \rightarrow \sigma | w_i, \beta$ ,  $A$

**Finish:**  $\beta = \varnothing$

This is the common “arc-eager” variant: a head can immediately take a right dependent, before *its* dependents are found



## Example


1. Left-Arc:  $\alpha[w, w] \beta, A \Rightarrow \alpha, w] \beta, A U \{r(w, w)\}$   
 Precondition:  $\{w, r', w\} \in A, w, \# \text{ ROOT}$
2. Right-Arc:  $\alpha, w] \beta, A \Rightarrow \alpha[w] \beta, A U \{r(w, w)\}$   
 Precondition:  $\{w, r', w\} \in A$
3. Reduce  $\alpha[w, \beta, A \Rightarrow \alpha, \beta, A$   
 Precondition:  $\{w, r', w\} \in A$
4. Shift  $\alpha, w] \beta, A \Rightarrow \alpha, w] \beta, A$

*Happy children like to play with their friends .*

	[ROOT]	[Happy, children, ...]	$\varnothing$
Shift	[ROOT, Happy]	[children, like, ...]	$\varnothing$
$LA_{\text{amod}}$	[ROOT]	[children, like, ...]	$A_1$
Shift	[ROOT, children]	[like, to, ...]	$A_2$
$LA_{\text{nsubj}}$	[ROOT]	[like, to, ...]	$A_1 \cup \{\text{nsubj}(\text{like}, \text{children})\} = A_2$
$RA_{\text{root}}$	[ROOT, like]	[to, play, ...]	$A_2 \cup \{\text{root}(\text{ROOT}, \text{like})\} = A_3$
Shift	[ROOT, like, to]	[play, with, ...]	$A_3$
$LA_{\text{aux}}$	[ROOT, like]	[play, with, ...]	$A_3 \cup \{\text{aux}(\text{play}, \text{to})\} = A_4$
$RA_{\text{xcomp}}$	[ROOT, like, play]	[with their, ...]	$A_4 \cup \{\text{xcomp}(\text{like}, \text{play})\} = A_5$

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 **Example**

*Happy children like to play with their friends .*

RA <sub>xcomp</sub>	[ROOT, like, play]	[with their, ...]	A <sub>4</sub> ∪ {xcomp(like, play)} = A <sub>5</sub>
RA <sub>prep</sub>	[ROOT, like, play, with]	[their, friends, ...]	A <sub>5</sub> ∪ {prep(play, with)} = A <sub>6</sub>
Shift	[ROOT, like, play, with, their]	[friends, .]	A <sub>6</sub>
LA <sub>poss</sub>	[ROOT, like, play, with]	[friends, .]	A <sub>6</sub> ∪ {poss(friends, their)} = A <sub>7</sub>
RA <sub>pobj</sub>	[ROOT, like, play, with, friends]	[.]	A <sub>7</sub> ∪ {pobj(with, friends)} = A <sub>8</sub>
Reduce	[ROOT, like, play, with]	[.]	A <sub>8</sub>
Reduce	[ROOT, like, play]	[.]	A <sub>8</sub>
Reduce	[ROOT, like]	[.]	A <sub>8</sub>
RA <sub>punc</sub>	[ROOT, like, .]	[]	A <sub>8</sub> ∪ {punc(like, .)} = A <sub>9</sub>

You terminate as soon as the buffer is empty. Dependencies = A<sub>9</sub>

- Left-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma, w_l | \beta, A \cup \{r(w_p, w_l)\}$   
Precondition:  $(w_p, r', w_l) \in A, w_l \neq \text{ROOT}$
- Right-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma | w_l, w_p | \beta, A \cup \{r(w_p, w_l)\}$
- Reduce:  $\sigma | w_p, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_p, r', w_l) \in A$
- Shift:  $\sigma, w_l | \beta, A \rightarrow \sigma | w_p, \beta, A$

**Example**

- Left-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma, w_l | \beta, A \cup \{r(w_p, w_l)\}$   
Precondition:  $(w_p, r', w_l) \in A, w_l \neq \text{ROOT}$
- Right-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma | w_l, w_p | \beta, A \cup \{r(w_p, w_l)\}$
- Reduce:  $\sigma | w_p, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_p, r', w_l) \in A$
- Shift:  $\sigma, w_l | \beta, A \rightarrow \sigma | w_p, \beta, A$

$\left[ \text{\_ROOT\_} \right]_S \left\{ \text{Red figures on the screen indicated falling stocks} \right\}_Q$

28 Dependency Parsing (Prashanth Mannem) October 21, 2013

**Example**

- Left-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma, w_l | \beta, A \cup \{r(w_p, w_l)\}$   
Precondition:  $(w_p, r', w_l) \in A, w_l \neq \text{ROOT}$
- Right-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma | w_l, w_p | \beta, A \cup \{r(w_p, w_l)\}$
- Reduce:  $\sigma | w_p, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_p, r', w_l) \in A$
- Shift:  $\sigma, w_l | \beta, A \rightarrow \sigma | w_p, \beta, A$

$\left[ \text{\_ROOT\_} \text{Red} \right]_S \left\{ \text{figures on the screen indicated falling stocks} \right\}_Q$

**Shift**

29 Dependency Parsing (P. Mannem) October 21, 2013

**Example**

- Left-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma, w_l | \beta, A \cup \{r(w_p, w_l)\}$   
Precondition:  $(w_p, r', w_l) \in A, w_l \neq \text{ROOT}$
- Right-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma | w_l, w_p | \beta, A \cup \{r(w_p, w_l)\}$
- Reduce:  $\sigma | w_p, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_p, r', w_l) \in A$
- Shift:  $\sigma, w_l | \beta, A \rightarrow \sigma | w_p, \beta, A$

$\left[ \text{\_ROOT\_} \right]_S \text{Red} \left\{ \text{figures on the screen indicated falling stocks} \right\}_Q$

**Left-arc**

30 Dependency Parsing (P. Mannem) October 21, 2013

**Example**

- Left-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma, w_l | \beta, A \cup \{r(w_p, w_l)\}$   
Precondition:  $(w_p, r', w_l) \in A, w_l \neq \text{ROOT}$
- Right-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma | w_l, w_p | \beta, A \cup \{r(w_p, w_l)\}$
- Reduce:  $\sigma | w_p, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_p, r', w_l) \in A$
- Shift:  $\sigma, w_l | \beta, A \rightarrow \sigma | w_p, \beta, A$

$\left[ \text{\_ROOT\_} \text{Red figures} \right]_S \left\{ \text{on the screen indicated falling stocks} \right\}_Q$

**Shift**

31 Dependency Parsing (P. Mannem) October 21, 2013

**Example**

- Left-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma, w_l | \beta, A \cup \{r(w_p, w_l)\}$   
Precondition:  $(w_p, r', w_l) \in A, w_l \neq \text{ROOT}$
- Right-Arc:  $\sigma | w_p, w_l | \beta, A \rightarrow \sigma | w_l, w_p | \beta, A \cup \{r(w_p, w_l)\}$
- Reduce:  $\sigma | w_p, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_p, r', w_l) \in A$
- Shift:  $\sigma, w_l | \beta, A \rightarrow \sigma | w_p, \beta, A$

$\left[ \text{\_ROOT\_} \text{Red figures on} \right]_S \left\{ \text{the screen indicated falling stocks} \right\}_Q$

**Right-arc**

32 Dependency Parsing (P. Mannem) October 21, 2013

### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

$\left[ \text{\_ROOT\_} \text{Red figures on the} \right]_S \left[ \text{screen indicated falling stocks} \right]_Q$

**Shift**

► 33 Dependency Parsing (P. Mannem) October 21, 2013

### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

$\left[ \text{\_ROOT\_} \text{Red figures on } \right]_S \left[ \text{the screen indicated falling stocks} \right]_Q$

**Left-arc**

► 34 Dependency Parsing (P. Mannem) October 21, 2013

### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

$\left[ \text{\_ROOT\_} \text{Red figures on the screen} \right]_S \left[ \text{indicated falling stocks} \right]_Q$

**Right-arc**

► 35 Dependency Parsing (P. Mannem) October 21, 2013

### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

$\left[ \text{\_ROOT\_} \text{Red figures on the screen} \right]_S \left[ \text{indicated falling stocks} \right]_Q$

**Reduce**

► 36 Dependency Parsing (P. Mannem) October 21, 2013

### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

$\left[ \text{\_ROOT\_} \text{Red figures} \right]_S \left[ \text{on the screen indicated falling stocks} \right]_Q$

**Reduce**

► 37 Dependency Parsing (P. Mannem) October 21, 2013

### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i, \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

$\left[ \text{\_ROOT\_} \right]_S \left[ \text{Red figures on the screen indicated falling stocks} \right]_Q$

**Left-arc**

► 38 Dependency Parsing (P. Mannem) October 21, 2013

### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i] \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

**Right-arc**

► 39 Dependency Parsing (P. Mannem) October 21, 2013

### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i] \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

**Shift**

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### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i] \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

**Left-arc**

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### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i] \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

**Right-arc**

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### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i] \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

**Reduce**

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### Example

1. Left-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma, w_j \beta, A \cup \{r(w_i, w_j)\}$   
Precondition:  $(w_i, r', w_j) \in A, w_i \neq \text{ROOT}$
2. Right-Arc,  $\sigma[w_i, w_j] \beta, A \rightarrow \sigma[w_i] w_j \beta, A \cup \{r(w_i, w_j)\}$
3. Reduce,  $\sigma[w_i] \beta, A \rightarrow \sigma, \beta, A$   
Precondition:  $(w_i, r', w_j) \in A$
4. Shift,  $\sigma, w_j \beta, A \rightarrow \sigma[w_j] \beta, A$

**Reduce**


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 **MaltParser**  
[Nivre et al. 2008]

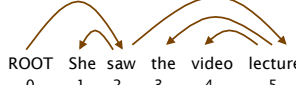
- We have left to explain how we choose the next action
- Each action is predicted by a discriminative classifier (often SVM, can be perceptron, maxent classifier) over each legal move
  - Max of 4 untyped choices, max of  $|R| \times 2 + 2$  when typed
  - Features: top of stack word, POS; first in buffer word, POS; etc.
- There is **NO** search (in the simplest and usual form)
  - But you could do some kind of beam search if you wish
- It provides **VERY** fast linear time parsing
- The model's accuracy is *slightly* below the best Lexicalized PCFGs (evaluated on dependencies), but
- It provides close to state of the art parsing performance

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 **Evaluation of Dependency Parsing:**  
(labeled) dependency accuracy


Acc =  $\frac{\# \text{ correct deps}}{\# \text{ of deps}}$

UAS =  $4 / 5 = 80\%$   
LAS =  $2 / 5 = 40\%$



Gold			Parsed		
1	2	She	1	2	She
2	0	saw	2	0	saw
3	5	the	3	4	the
4	5	video	4	5	video
5	2	lecture	5	2	lecture


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 **Representative performance numbers**


- The CoNLL-X (2006) shared task provides evaluation numbers for various dependency parsing approaches over 13 languages
  - MALT: LAS scores from 65–92%, depending greatly on language/treebank
- Here we give a few UAS numbers for English to allow some comparison to constituency parsing

Parser	UAS%
Sagae and Lavie (2006) ensemble of dependency parsers	92.7
Charniak (2000) generative, constituency, as dependencies	92.2
Collins (1999) generative, constituency, as dependencies	91.7
McDonald and Pereira (2005) – MST graph-based dependency	91.5
Yamada and Matsumoto (2003) – transition-based dependency	90.4


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 **Projectivity**

- Dependencies from a CFG tree using heads, must be **projective**
  - There must not be any crossing dependency arcs when the words are laid out in their linear order, with all arcs above the words.
- But dependency theory normally does allow non-projective structures to account for displaced constituents
  - You can't easily get the semantics of certain constructions right without these nonprojective dependencies



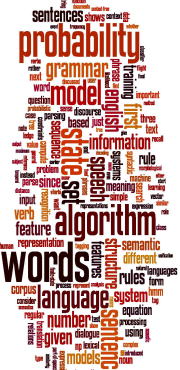
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 **Handling non-projectivity**

- The arc-eager algorithm we presented only builds projective dependency trees
- Possible directions to head:
  - Just declare defeat on nonprojective arcs
  - Use a dependency formalism which only admits projective representations (a CFG doesn't represent such structures...)
  - Use a postprocessor to a projective dependency parsing algorithm to identify and resolve nonprojective links
  - Add extra types of transitions that can model at least most non-projective structures
  - Move to a parsing mechanism that does not use or require any constraints on projectivity (e.g., the graph-based MSTParser)

**Dependencies encode relational structure**

**Relation Extraction with Stanford Dependencies**





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**Dependency paths identify relations like protein interaction**

[Erkan et al. EMNLP 07, Fundel et al. 2007]

KaiC ←<sub>nsbj</sub> interacts <sub>prep\_with</sub>→ SasA  
 KaiC ←<sub>nsbj</sub> interacts <sub>prep\_with</sub>→ SasA <sub>conj\_and</sub>→ KaiA  
 KaiC ←<sub>nsbj</sub> interacts <sub>prep\_with</sub>→ SasA <sub>conj\_and</sub>→ KaiB

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**Stanford Dependencies**

[de Marneffe et al. LREC 2006]

- The basic dependency representation is projective
- It can be generated by postprocessing headed phrase structure parses (Penn Treebank syntax)
- It can also be generated directly by dependency parsers, such as MaltParser, or the Easy-First Parser

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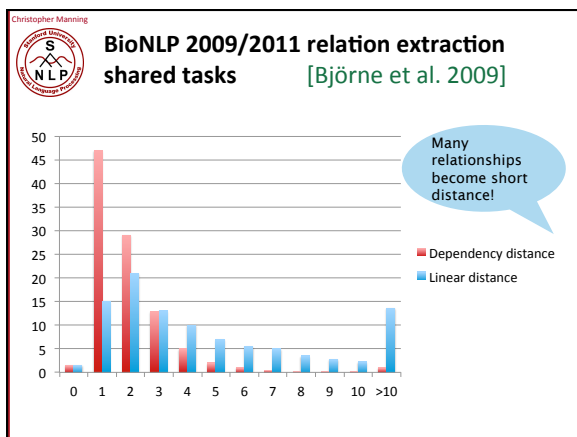
**Graph modification to facilitate semantic analysis**

Bell, based in LA, makes and distributes electronic and computer products.

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**Graph modification to facilitate semantic analysis**

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