

Deterministic Edge Connectivity and Max Flow using Subquadratic Cut Queries

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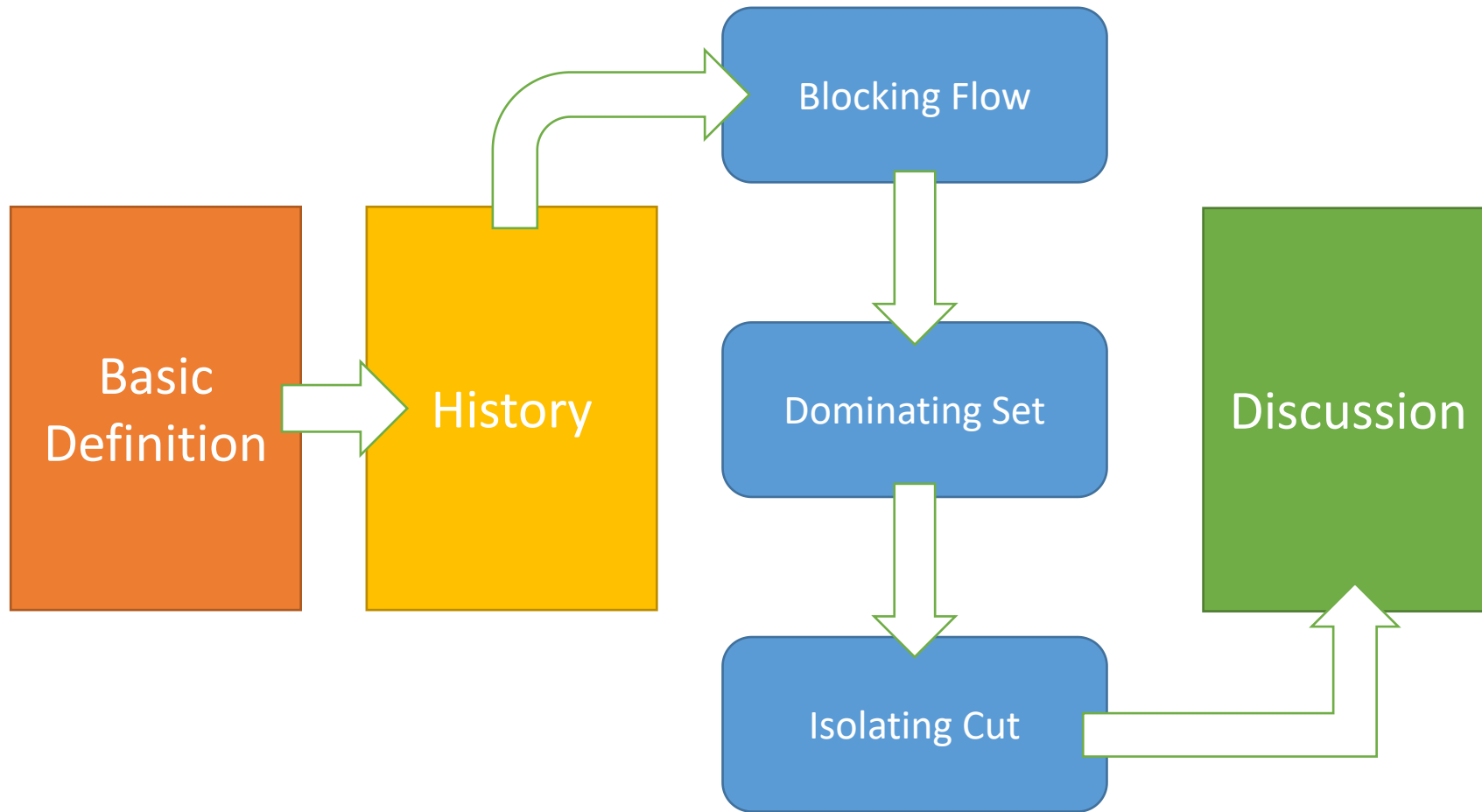
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Outline



Basic Definition

- **Cut Query Model**



Q: How many **query** do we need to compute the **global minimum cut**?

Sense of the problem

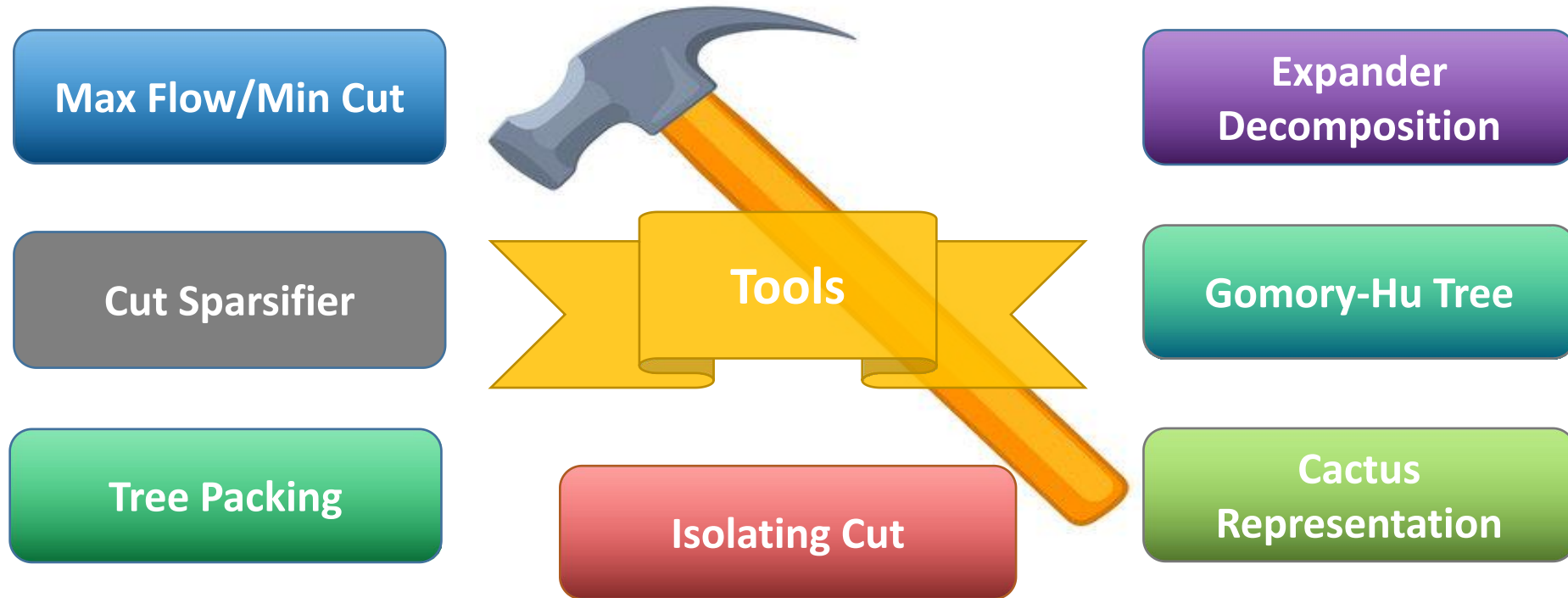


“Learn the whole graph”

- **A Naive Algorithm**
- Learn each edge (u, v) by computing $\text{Cut}(\{u\}) + \text{Cut}(\{v\}) - \text{Cut}(\{u, v\})$, which requires $O(n^2)$ queries [**Nearly Optimal**]
- For dense graph, one needs at least $\sim \Omega(n^2)$ **query** to learn the whole graph [**Rubinstein, Schramm & Weinberg, ITCS 2018**]
- One can use $O(n^2/\log n)$ **query** to learn the whole graph [**Grebinski & Kucherov, Algorithmica 2000**]

Motivation

- **Q:** Can we use **less information** to **represent** the **global minimum cut**?



History

- **Edge Connectivity in Cut Query Model**

	Det./ Rand.	Simple/ Weighted	Query
Rubinstein, Schramm& Weinberg[RSW18]	Rand.	Simple	$O(n \log^3 n)$
Mukhopadhyay & Nanongkai[MN20]	Rand.	Weighted	$O(n \log^{O(1)} n)$
Apers et al.[AEG ⁺ 22]	Rand.	Simple	$O(n)$
This paper	Det.	Simple	$\tilde{O}(n^{\frac{5}{3}})$

State-of-the-Art

- Different Settings in **Cut Query Model**

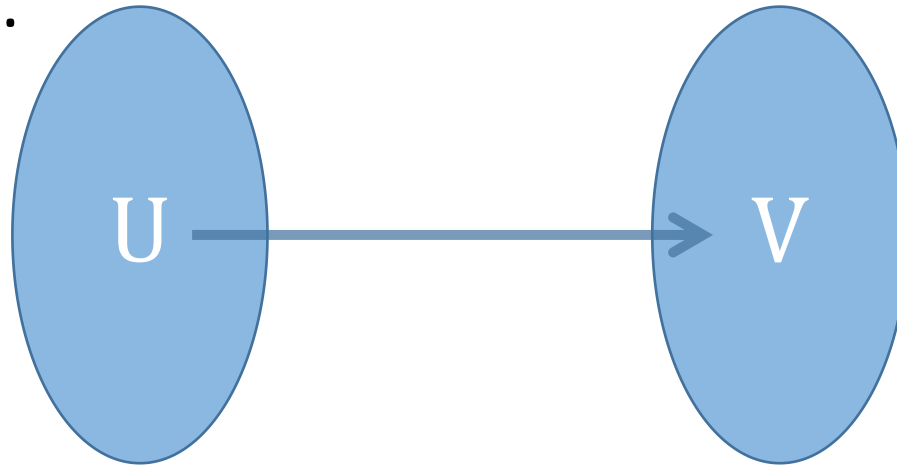
	Connectivity		Edge Connectivity	
	Lower	Upper	Lower	Upper
Deterministic	$\Omega(n)$ [HMT88]	$O(\frac{n \log n}{\log \log n})$ [LC24]	$\Omega(n)$ [HMT88]	$\tilde{O}(n^{5/3})$ (This paper)
Zero-error, Randomized	$\Omega(\frac{n \log \log(n)}{\log n})$ [RS95]	$O(n)$ [AEG ⁺ 22]	$\Omega(n)$	$\tilde{O}(n^{5/3})$ (This paper)
Bounded Error, Randomized	$\Omega(\frac{n}{\log n})$ [BFS86]	$O(n)$ [AEG ⁺ 22]	$\Omega(\frac{n \log \log(n)}{\log n})$ [AD21]	$O(n)$ [AEG ⁺ 22]
Quantum	$\Omega(1)$	$O(\log^5(n))$ [AL21]	$\Omega(1)$	$\tilde{O}(\sqrt{n})$ [AEG ⁺ 22]

Max Flow/Min Cut

- **Max Flow Min Cut Theorem**
- No Duality Gap for $s - t$ **max flow** and $s - t$ **min cut**
max flow value = min cut capacity
- **Q:** Why should we consider **Max Flow** as a start point?
- **Inspiration:** [RSW18] shows that, for graph G with integral weights from $[0, W]$, every $s - t$ flow of value f can be covered by edges of at most $O(n\sqrt{fW})$ total weight
- **We can use $O(n\sqrt{n})$ edges to cover any $s-t$ flow in simple graph!**

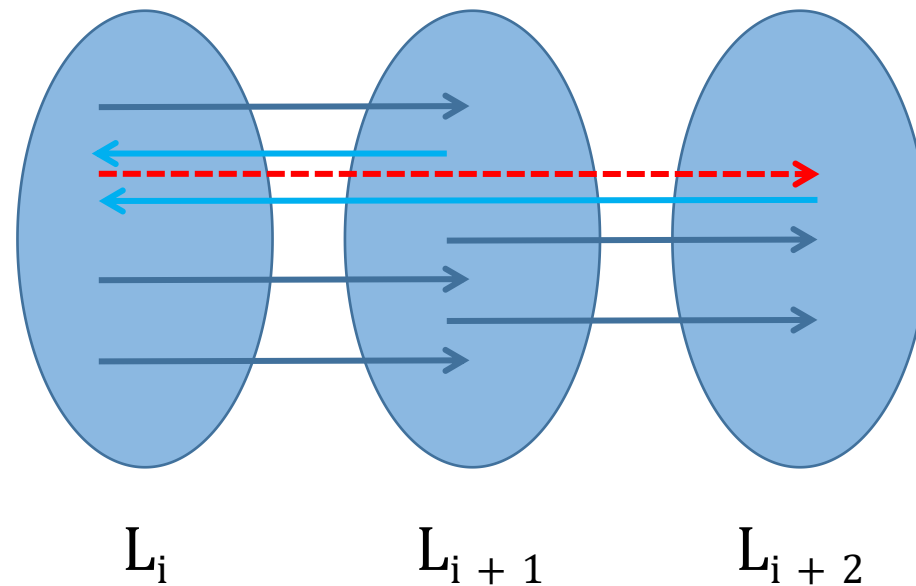
BIS/Cross Query

- A **BIS (Bipartite Independent Set)** or **Cross Query** asks whether there exists an edge between two sets U and V . In other words, it checks if there is an edge (u, v) such that $u \in U, v \in V$.
- **Fact:** A **BIS/Cross query** can be replaced by $O(1)$ Cut Query in undirected graph.

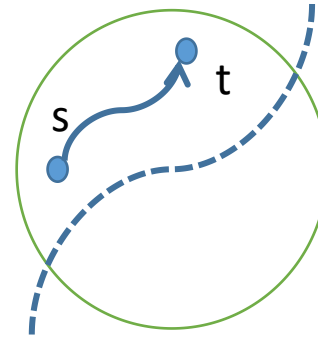
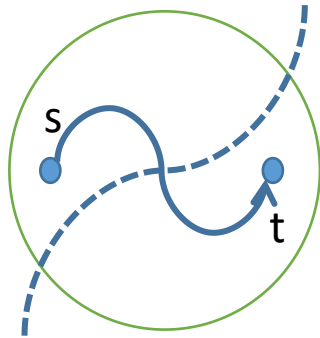


Blocking Flow

- **Theorem 1:** We can use $\sim O(n^{5/3})$ **BIS/Cross query** to obtain an **explicit** s-t max flow in simple graph [**Idea: Dinitz's algorithm**]

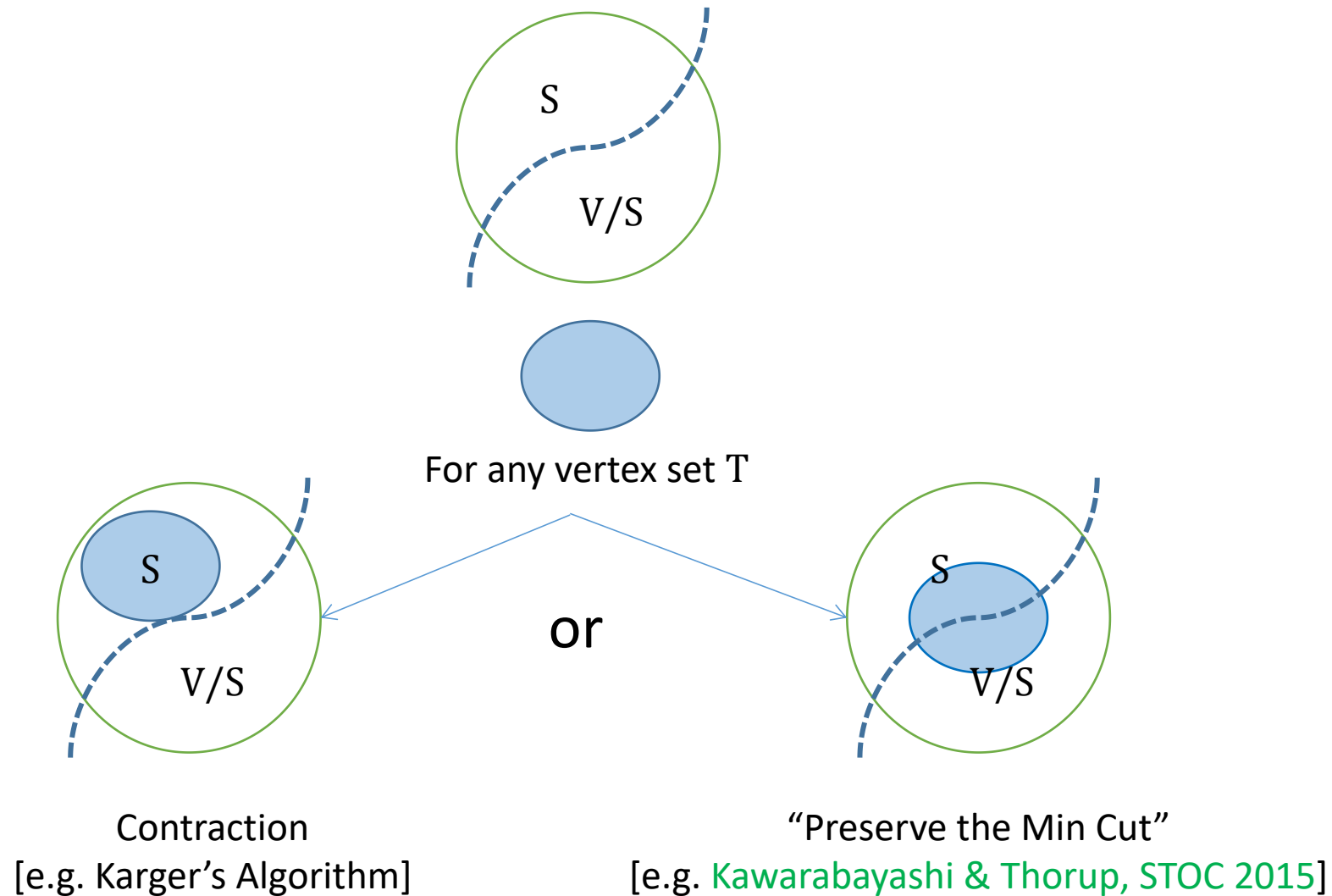


From Flow to Cut



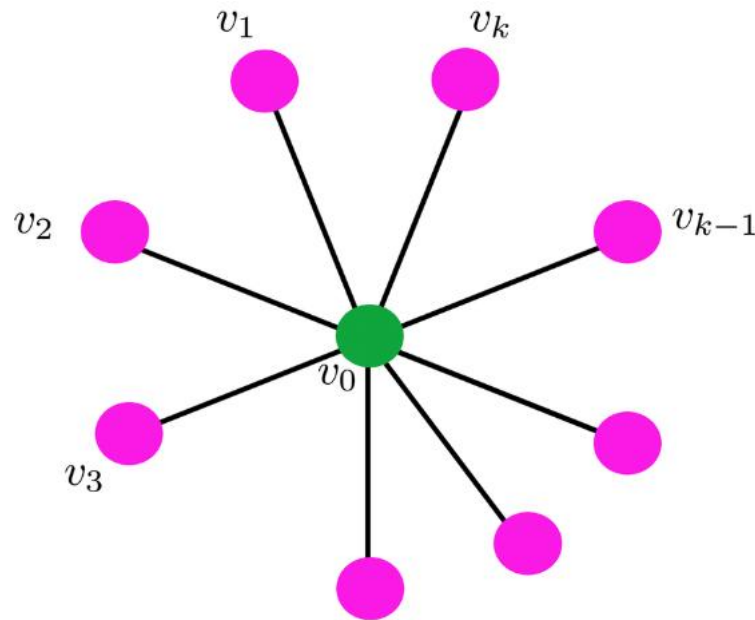
How can we guarantee that s and t are on **different** sides of the minimum cut?

“Preserve the Minimum Cut”



Dominating Set

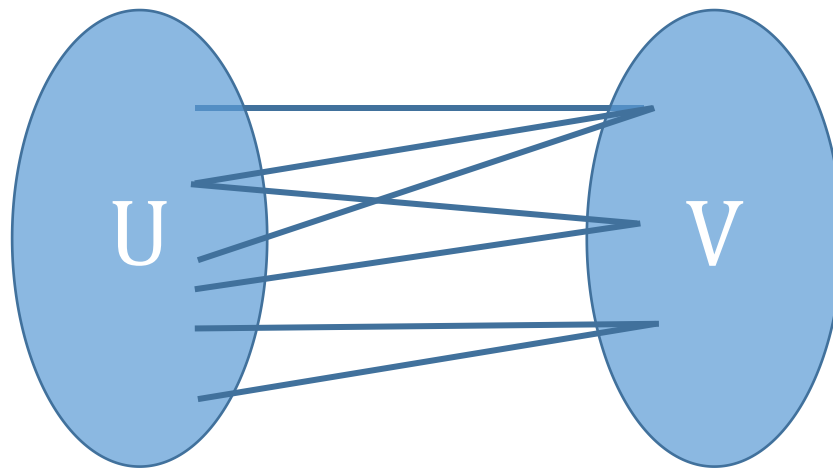
- **A Key Observation:** In simple graph, a dominating set can “preserve all non-trivial minimum cut”



Special Case: Star Graph[No non-trivial minimum cut]

Dominating Set

- **Theorem 2:** If the minimum degree is δ , then we can find a dominating set D with size at most $O(\frac{n}{\delta} \log \frac{n}{\delta})$ with $\sim O(n)$ cut **query**.
- Existence: Sample each vertex with probability $\sim \frac{1}{\delta}$.
- **De-randomize Idea:** Finding an element **above the average**

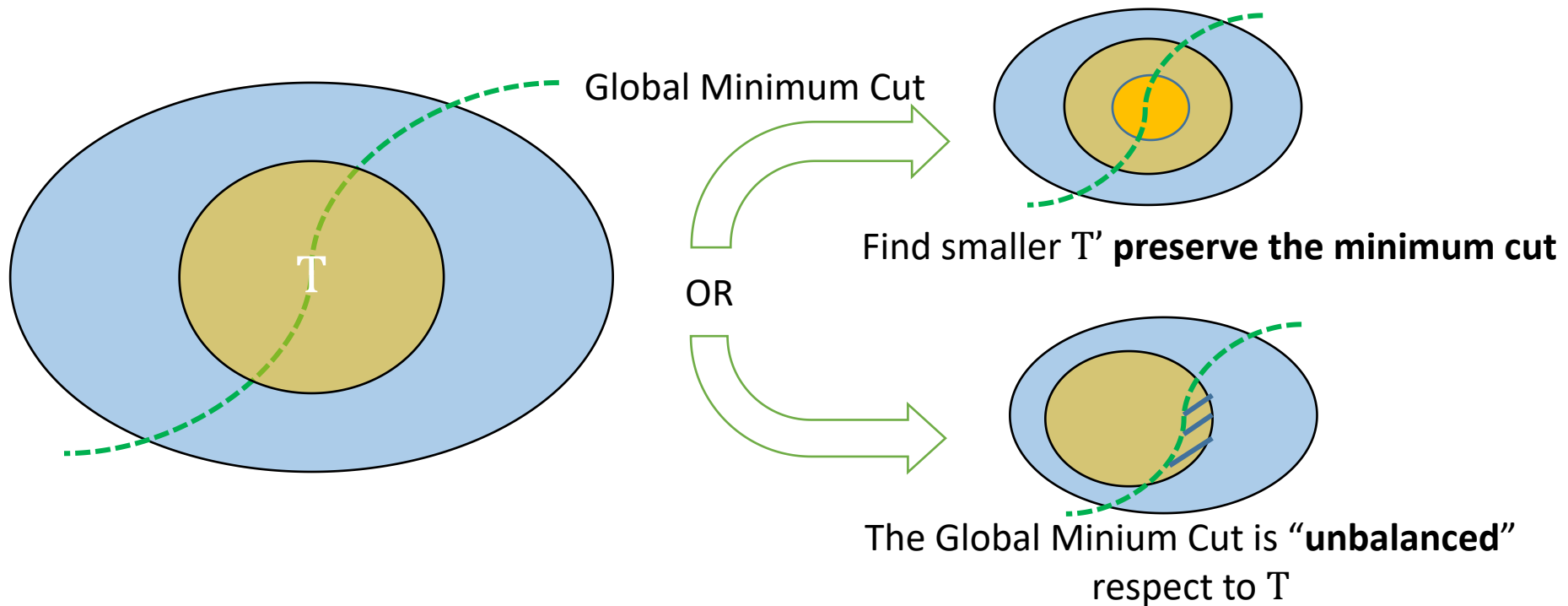


In $O(\log n)$ queries

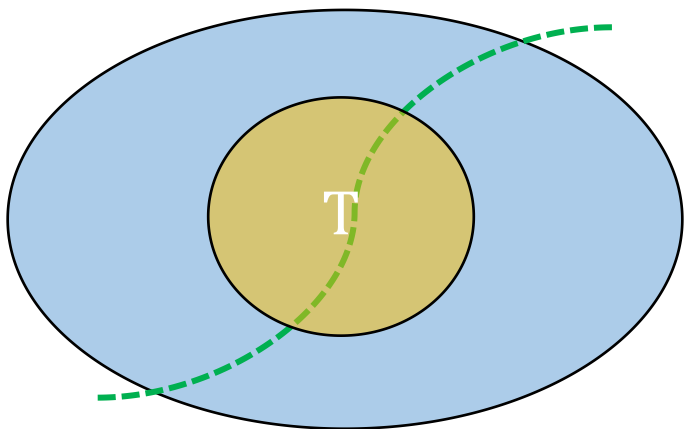
Find $u \in U$ such that $\deg_V(u) \geq |E(U, V)|/|U|$

Framework of Isolating Cut

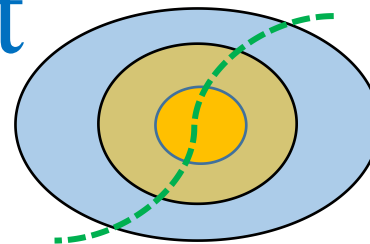
- Isolating Cut[Li & Panigrahi, FOCS 2020]



Framework of Isolating Cut

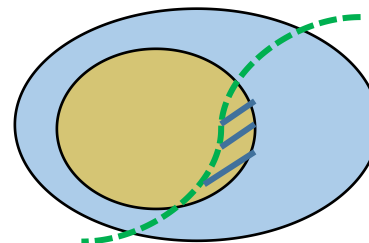


Global Minimum Cut



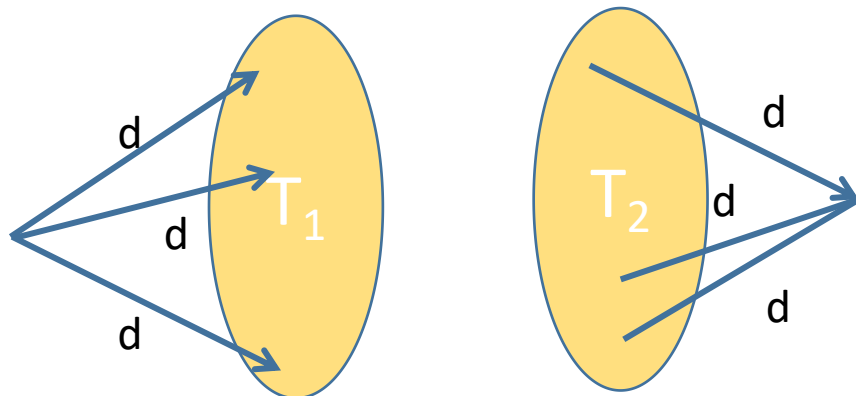
Find T' preserve the minimum cut

OR



The Global Minimum Cut is “unbalanced”

Cut Matching Game



Almost every edge saturated

OR

Find balanced sparse cut



Polylog times

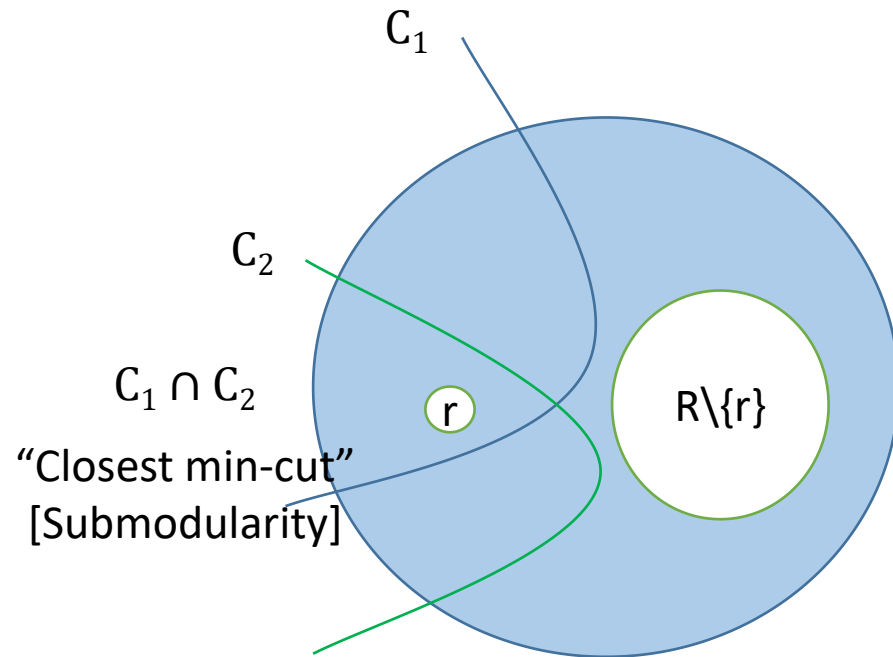
Guarantee “Almost Expander”

Demand $d \leq \delta$ [Total Flow Size $\leq \sim O(n)$]

NOT exactly the same as [LP20], we don't compute $T_1 - T_2$ min-cut

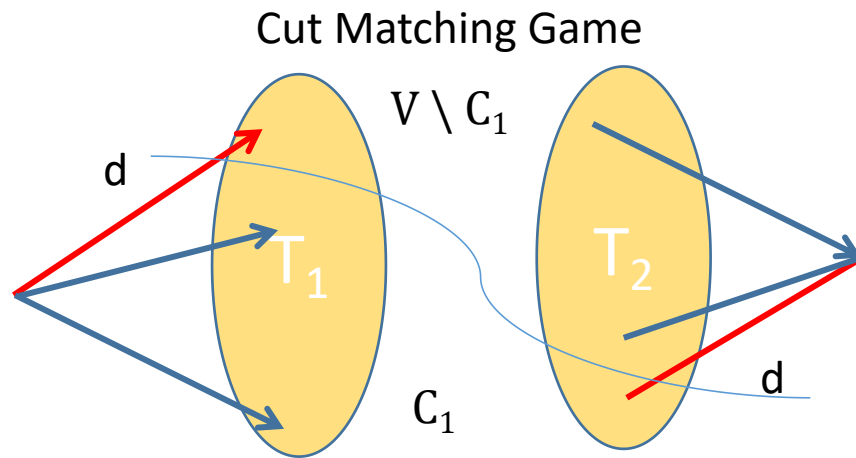
Minimum Isolating Cut

- For any set of vertices $R \subseteq V$, $r \in R$, the minimum isolating cut of r is an $\{r\} - \{R \setminus \{r\}\}$ min-cut



Subroutine

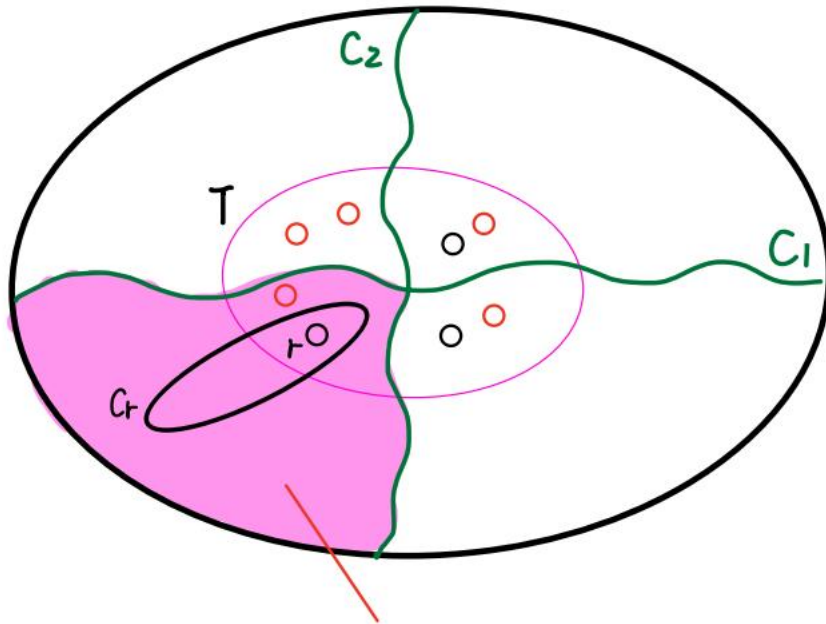
- Let $d = \tau + 1$, we will either output an **isolating cut** of R of size at most τ ,
- or certifies that the **minimum isolating cut** of R has a size larger than τ .



If an edge is **saturated**, then the corresponding minimum isolating cut has size at least d .

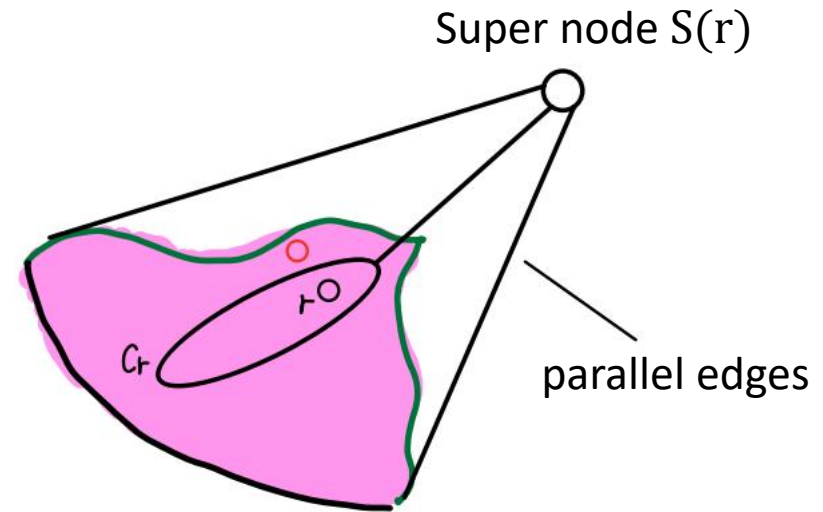
If the minimum isolating cut of R is C_r less than d , then we must have $C_r \subseteq C_1$ or $C_r \subseteq V/C_1$

“From Global to Local”



Guarantee: each part contains at most 1 black vertex

- not saturated in all max-flow call
- saturated at least once



Call $r - S(r)$ max flow

Discussion

