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CSC 225

1. $\frac{1}{n}$ 2^{100} $\log \log n$ $\sqrt{\log n}$ $\log^2 n$ $n^{0.01}$ $\sqrt{n}, 3n^{0.5}$ $2^{\log n}, 5n$ $n \log_4 n, 6n \log n$ $2n \log^2 n$ $4n^{\frac{3}{2}}$ $4^{\log n}$ $n^2 \log n$ n^3 2^n 4^n 2^{2^n} 2. $d(n)$ is $O(f(n))$ and $e(n)$ is $O(g(n))$ $\therefore d(n) = k \cdot f(n)$ for all $n \geq N$, $e(n) = L \cdot g(n)$ for all $n \geq M$. $\therefore d(n)e(n) = (k \cdot L) f(n)g(n)$ for all $n \geq n_0$ That is $d(n)e(n)$ is $O(f(n)g(n))$ 3. $\log_b f(n) = \log_2 f(n) / \log_2 b$ and $b > 1$ and b is a constant. $\therefore 1/\log_2 b$ is a constant. $\therefore \log_b f(n)$ is $O(\log_2 f(n))$ and $\log_2 f(n)$ is $O(\log_b f(n))$.That is if $b > 1$ is a constant, $\log_b f(n)$ is $\Theta(\log_2 f(n))$.

4. When $n=1$, $T(1)=1$

Assume when $n=k$, $T(k)=2^{k+1}-1$

When $n=k+1$,

$$T(k+1) = T(k) + 2^{k+1} = 2^{k+1} - 1 + 2^{k+1} = 2^{k+2} - 1 = 2^{(k+1)+1} - 1$$

Hence, $T(n) = 2^{n+1} - 1$

5. (a) T.b.f 4

$$T(n) = 1 + n \times (1+1+1) + 1 + 1 = 3n + 3$$

(b) Let S_k : arrayFind returns the index of element or -1 after k^{th} iteration.

Base: show true for S_0 .

Case 1: $X = A[0]$, returns 0 \rightarrow True.

Case 2: $X \neq A[0]$, returns -1 \rightarrow True.

I.H: Assume S_{i-1} is true.

Case 1: $X = A[0]$ or $A[1]$ or \dots or $A[i-1]$, returns index \rightarrow True.

Case 2: $X \neq A[0]$ or $A[1]$ or \dots or $A[i-1]$, returns -1 \rightarrow True.

I.S: Show S_i is true after i^{th} iteration.

Case 1: $X = A[0]$ or $A[1]$ or \dots or $A[i]$, returns index, \rightarrow True.

Case 2: $X \neq A[0]$ or $A[1]$ or \dots or $A[i]$, returns -1, \rightarrow True.

Therefore, arrayFind is correct.