

# Extended Operators of Relational Algebra

# Bags

- A **bag** is like a set, but an element may appear more than once.
  - *Multiset* is another name for “bag.”
- Example:
  - $\{1,2,1,3\}$  is a bag.
  - $\{1,2,3\}$  is also a bag that happens to be a set.
- Bags also resemble lists, but **order in a bag is unimportant**.
  - Example:
    - $\{1,2,1\} = \{1,1,2\}$  as bags, but
    - $[1,2,1] \neq [1,1,2]$  as lists.

# Why bags?

- SQL is actually a bag language.
- SQL will eliminate duplicates, but usually only if you ask it to do so explicitly
  - except for union, intersection, and difference where the default is “set mode”.

# Bag Union

- **Union, intersection, and difference** need new definitions for bags.
- An element appears in the **union** of two bags the **sum** of the number of times it appears in each bag.
- Example:  
$$\{1,2,1\} \cup \{1,1,2,3,1\}$$
$$= \{1,1,1,1,1,2,2,3\}$$

# Bag Intersection

- An element appears in the **intersection** of two bags the **minimum** of the number of times it appears in either.

- Example:

$$\{1,2,1\} \cap \{1,2,3\} \\ = \{1,2\}.$$

# Bag Difference

- An element appears in **difference**  $A - B$  of bags as many times as it appears in  $A$ , **minus** the number of times it appears in  $B$ .
  - But never less than 0 times.
- Example:  $\{1,2,1\} - \{1,2,3\}$   
 $= \{1\}.$

# Examples of Union, Intersection, Difference in SQL

- Remember, we need to have the same schema for the relations that we union, intersect, or take difference.

```
(SELECT * FROM R)
UNION
(SELECT * FROM S);
```

```
(SELECT * FROM R)
INTERSECT
(SELECT * FROM S);
```

```
(SELECT * FROM R)
EXCEPT                --MINUS in Oracle
(SELECT * FROM S);
```

- Add “**ALL**” for bag version of these operators. Only UNION ALL exists in Oracle.
- These are the only operators that work in ‘set mode’ by default. All the others work in ‘bag mode’ by default.

# The Extended Algebra

1.  $\delta$ : eliminates duplicates from bags.
2.  $\tau$ : sorts tuples.
3. **Extended projection**: arithmetic, duplication of columns.
4.  $\gamma$ : grouping and aggregation.
5. **OUTERJOIN**: includes “dangling tuples” = tuples that do not join with anything.



# Example: Duplicate Elimination

$$R_1 := \delta(R_2)$$

- $R_1$  consists of one copy of each tuple that appears in  $R_2$  one or more times.

R =

A	B
1	2
3	4
1	2

```
SELECT DISTINCT *  
FROM R;
```

$\delta(R)$  =

A	B
1	2
3	4

# Sorting

$R_1 := \tau_L(R_2).$

–  $L$  is a list of some of the attributes of  $R_2$ .

- $R_1$  is the list of tuples of  $R_2$  sorted first on the value of the first attribute on  $L$ , then on the second attribute of  $L$ , and so on.
- $\tau$  is the only operator whose result is neither a set nor a bag.

Example:

```
SELECT *  
FROM R  
ORDER BY A, B;
```

# Example: Extended Projection

Using the same  $\pi_L$  operator, we allow the list  $L$  to contain arbitrary expressions involving attributes, for example:

1. Arithmetic on attributes, e.g.,  $A+B$ .
2. Duplicate occurrences of the same attribute.

R =

A	B
1	2
3	4

```
SELECT A+B AS C, A AS A1, A AS A2  
FROM R;
```

$\pi_{A+B \rightarrow C, A \rightarrow A1, A \rightarrow A2}(R) =$

C	A1	A2
3	1	1
7	3	3

# Aggregation Operators

- They apply to entire columns of a table and produce a single result.
- The most important examples:
  - SUM
  - AVG
  - COUNT
  - MIN
  - MAX

# Example: Aggregation

R =

A	B
1	3
3	4
3	2

SUM(A) = 7

COUNT(A) = 3

MAX(B) = 4

MIN(B) = 2

AVG(B) = 3

```
SELECT SUM(A), COUNT(A), MAX(B), MIN(B), AVG(B)
FROM R;
```

# Grouping Operator

$$R_1 := \gamma_L (R_2)$$

$L$  is a list of elements that are either:

1. Individual (*grouping*) attributes.
2.  $AGG(A)$ , where  $AGG$  is one of the aggregation operators and  $A$  is an attribute.

# Example: Grouping/Aggregation

R =

A	B	C
1	2	3
4	5	6
1	2	5

Then, average *C* within groups:

A	B	AVG(C)
1	2	4
4	5	6

$\gamma_{A,B,AVG(C)}(R) = ??$

First, group *R* :

A	B	C
1	2	3
1	2	5
4	5	6

```
SELECT A,B,AVG(C)
FROM R
GROUP BY A,B;
```

## $\gamma_L(R)$ - Formally

- Group  $R$  according to all the grouping attributes on list  $L$ .
  - That is, form one group **for each distinct list** of values for those attributes in  $R$ .
- Within each group, compute  $AGG(A)$  for each aggregation on list  $L$ .
- Result has grouping attributes and aggregations as attributes.
- One tuple for each list of values for the grouping attributes **and** their group's aggregations.



# Example: Grouping/Aggregation

**StarsIn**(title, year, starName)

- For each star who has appeared in at least three movies give the earliest year in which he or she appeared.
  - First we group, using *starName* as a grouping attribute.
  - Then, we compute the  $\text{MIN}(\text{year})$  for each group.
  - Also, we need to compute the  $\text{COUNT}(\text{title})$  aggregate for each group, for filtering out those stars with less than three movies.
- $\pi_{\text{starName}, \text{minYear}}(\sigma_{\text{ctTitle} \geq 3}(\gamma_{\text{starName}, \text{MIN}(\text{year}) \rightarrow \text{minYear}, \text{COUNT}(\text{title}) \rightarrow \text{ctTitle}}(\text{StarsIn})))$

# Example: Grouping/Aggregation

```
SELECT starName, miny
FROM (SELECT starname, COUNT(title) AS cnt, MIN(year) AS miny
      FROM StarsIn
      GROUP BY starname)
WHERE cnt>=3;
```

Or (more concisely):

```
SELECT starname, MIN(year) AS miny
FROM StarsIn
GROUP BY starname
HAVING COUNT(title)>=3;
```

# Problems

**Product**(maker, model, type)

**PC**(model, speed, ram, hd, rd, price)

**Laptop**(model, speed, ram, hd, screen, price)

**Printer**(model, color, type, price)

Find the manufacturers who sell exactly three different models of PC.

Find those manufacturers of at least two different computers (PC or Laptops) with speed of at least 700.

# Outerjoin

## Motivation

- Suppose we join  $R \bowtie S$ .
- A tuple of  $R$  which doesn't join with any tuple of  $S$  is said to be *dangling*.
  - Similarly for a tuple of  $S$ .
  - **Problem:** We loose dangling tuples.

## Outerjoin

- Preserves dangling tuples by padding them with a special **NULL** symbol in the result.

## Example: Outerjoin

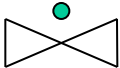
R =

A	B
1	2
4	5

S =

B	C
2	3
6	7

(1,2) joins with (2,3), but the other two tuples are dangling.

R  S =

A	B	C
1	2	3
4	5	NULL
NULL	6	7

```
SELECT *  
FROM R FULL OUTER JOIN S USING(B);
```

# Problems

- $R(A,B) = \{(0,1), (2,3), (0,1), (2,4), (3,4)\}$
- $S(B,C) = \{(0,1), (2,4), (2,5), (3,4), (0,2), (3,4)\}$
- $\gamma_{A, \text{SUM}(B)}(R)$
- $R \bowtie S$
- $R \bowtie_L S$  -- This left outerjoin: Only pad dangling tuples from the left table.
- $R \bowtie_R S$  -- -- This right outerjoin: Only pad dangling tuples from the right table.