Security in Computing & Information Technology

COSC2536/2537

Assignment 1

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Question 1

a) [Double Transposition Cipher]

Step 1: Find a row with "T", "H", "I", and "S"

	1	2	3	4	5	6	7	
1	S	0	G	I	U	Ν	В	<- has "I"
2	С	Ν	U	0	D	L	М	<- has none
3	I	Χ	0	Т	S	1	Ν	<- has "T", "I" and "S"
4	С	R	Н	1	Ν	Р	Ε	<- has "H" and "I"
5	Н	Ν	W	R	Т	0	Α	<- has "T" and "H"
6	R	0	S	Α	Т	Ν	Р	<- has "T" and "S"
7	Н	Α	I	I	Т	S	S	<- has "T", "H", "I" and "S" (2x)
8	0	В	0	F	Ε	D	U	<- has none
9	0	0	Т	S	Р	1	1	<- has "T", "I" and "S"
10	Ε	L	М	Χ	Ν	Α	Р	<- has none
11	Ε	S	Α	Т	L	R	Ν	<- has "T", and "S"

Row 7 is the only row with the right combination of letters. However, because there are duplicates of some letters, we will have to order the columns to satisfy all possible combinations of T H I S.

Step 2: Move row 7 to the top.

	1	2	3	4	5	6	7
7	Н	Α	1	1	Т	S	S
1	S	0	G	I	U	Ν	В
2	С	Ν	U	0	D	L	М
3	ı	Χ	0	Т	S	1	Ν
4	С	R	Н	1	Ν	Р	Ε
5	Н	Ν	W	R	Т	0	Α
6	R	0	S	Α	Т	Ν	Р
8	0	В	0	F	Ε	D	U
9	0	0	Т	S	Р	1	1
10	Ε	L	М	Χ	Ν	Α	Р
11	Ε	S	Α	Т	L	R	Ν

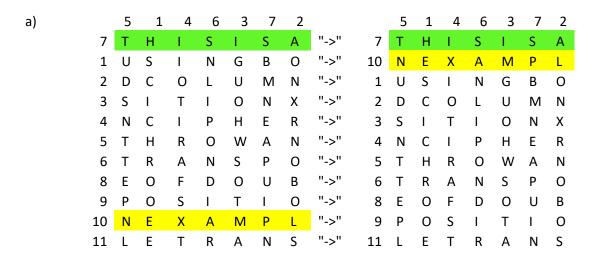
Step 3: Creating all column combinations to have "T H I S" as the first word.

	5	1	4	6	3	7	2		5	1	4	7	3	6	2
7	Т	Н	1	S	I	S	Α	7	Т	Н	1	S	1	S	Α
1	U	S	I	Ν	G	В	О	1	U	S	I	В	G	N	0
2	D	С	0	L	U	М	Ν	2	D	С	0	М	U	L	N
3	S	I	Т	I	0	N	Χ	3	S	I	Т	N	0	I	Χ
4	N	С	I	Р	Н	Ε	R	4	N	С	I	Ε	Н	Р	R
5	T	Н	R	0	W	Α	Ν	5	T	Н	R	Α	W	0	N
6	Т	R	Α	N	S	Р	О	6	Т	R	Α	Р	S	N	0
8	Ε	0	F	D	0	U	В	8	Ε	0	F	U	0	D	В
9	Р	0	S	I	Т	I	О	9	Р	0	S	I	T	I	0
10	N	E	Χ	Α	M	Р	L	10	N	Ε	Χ	Р	M	Α	L
11	L	Ε	Т	R	Α	N	S	11	L	Ε	Т	N	Α	R	S
	5	1	3	7	4	6	2		5	1	3	6	4	7	2
7	Т	Н	1	S	1	S	Α	7	Т	Н	1	S	1	S	Α
1	T U	H S	l G	S B	1 1	S N	A O	1	T U	H S	l G	S N	I I	S B	A O
1 2	T U D	H S C	G U	S B M	I I O	S N L	A O N	1 2	T U D	H S C	I G U	S N L	I I O	S B M	A O N
1 2 3	T U D S	H S C	G U O	S B M N	I I O T	S N L I	A O N X	1 2 3	T U D S	H S C	I G U O	S N L I	I I O T	S B M N	A O N X
1 2 3 4	T U D S N	H S C I	G U O H	S B M N	I O T I	S N L I	A O N X R	1 2 3 4	T U D S N	H S C I	G U O H	S N L I	I O T I	S B M N	A O N X R
1 2 3 4 5	T U D S N T	H S C I C	G U O H W	S B M N E	I O T I R	S N L I P	A O N X R	1 2 3 4 5	T U D S N T	H S C I C	G U O H W	S N L I P	I O T I R	S B M N E	A O N X R
1 2 3 4 5 6	T U D S N T	H S C I C H R	G U O H W S	S B M N E A	I O T I R	S N L I P O N	A O N X R O	1 2 3 4 5 6	T U D S N T	H S C I C H R	I G U O H W S	S N L I P O N	I O T I R A	S B M N E A	A O N X R O
1 2 3 4 5 6 8	T U D S N T T	H S C I C H R	G U O H W S	S B M N E A P	I O T I R A F	S N L I P O N D	A O N X R N O B	1 2 3 4 5 6 8	T U D S N T T	H S C I C H R	G U O H W S	S N L I P O N D	I O T I R A F	S B M N E A P	A O N X R O B
1 2 3 4 5 6 8 9	T U D S N T T E	H S C I C H R O	G U O H W S O T	S B M N E A P U I	I O T I R A F S	S N L I P O N D	A O N X R N O B O	1 2 3 4 5 6 8 9	T U D S N T T E	H S C I C H R O	G U O H W S O T	S N L I P O N D	I O T I R A F S	S B M N E A P U I	A O N X R O B O
1 2 3 4 5 6 8	T U D S N T T	H S C I C H R	G U O H W S	S B M N E A P	I O T I R A F	S N L I P O N D	A O N X R N O B	1 2 3 4 5 6 8	T U D S N T T	H S C I C H R	G U O H W S	S N L I P O N D	I O T I R A F	S B M N E A P	A O N X R O B

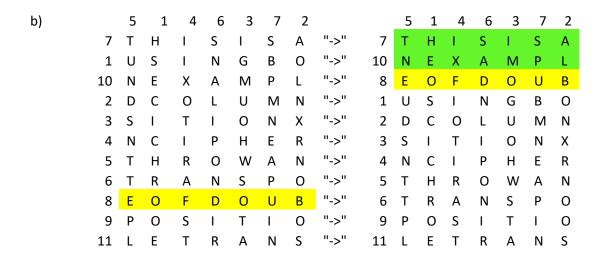
From the first combination of "T H I S", the first line is clearly "This is a". Thus, I only need to do is swap the Is and Ss.

After looking at all the combinations, the first is the only one with legible words. (E.g. Column, row)

Step 4: Changing the row order to make words and sentences



Adding row 10 to the end of 7 will create "This is an exampl"



Row 8 is the only row that can finish the "E" in example. It also adds "of", which makes logical sense in a sentence.

c)		5	1	4	6	3	7	2			5	1	4	6	3	7	2
	7	Т	Н	1	S	1	S	Α	"->"	7	Т	Н	1	S	1	S	Α
	10	Ν	Ε	Χ	Α	M	Р	L	"->"	10	Ν	Ε	Χ	Α	M	Р	L
	8	Ε	0	F	D	0	U	В	"->"	8	Ε	0	F	D	0	U	В
	1	U	S	1	Ν	G	В	0	"->"	11	L	Ε	Т	R	Α	N	S
	2	D	С	0	L	U	М	N	"->"	1	U	S	1	N	G	В	0
	3	S	1	Т	I	0	Ν	Χ	"->"	2	D	С	0	L	U	М	Ν
	4	Ν	С	1	Р	Н	Ε	R	"->"	3	S	I	Т	I	0	N	Χ
	5	Τ	Н	R	0	W	Α	Ν	"->"	4	Ν	С	I	Р	Н	Ε	R
	6	Τ	R	Α	Ν	S	Р	0	"->"	5	Т	Н	R	0	W	Α	Ν
	9	Р	0	S	I	Т	I	0	"->"	6	Т	R	Α	N	S	Р	0
	11	1	E	т	D	Λ	NI	C	"_<"	۵	D	\circ	c	1	т	1	\circ

Row 11 is the only row that can finish the "LE" in double.

d)	5	1	4	6	3	7	2			5	1	4	6	3	7	2
7	Т	Н	1	S	1	S	Α	"->"	7	Т	Н	1	S	1	S	Α
10	N	Е	Χ	Α	M	Р	L	"->"	10	N	Ε	Χ	Α	М	Р	L
8	Ε	0	F	D	0	U	В	"->"	8	Ε	0	F	D	0	U	В
11	L	Е	Т	R	Α	N	S	"->"	11	L	Ε	Т	R	Α	N	S
1	U	S	1	Ν	G	В	0	"->"	9	Р	0	S	1	Т	1	O
2	D	С	Ο	L	U	M	Ν	"->"	1	U	S	1	Ν	G	В	0
3	S	1	Т	I	О	Ν	Χ	"->"	2	D	С	0	L	U	М	Ν
4	Ν	С	1	Р	Н	Ε	R	"->"	3	S	I	Т	I	0	Ν	Χ
5	Т	Н	R	0	W	Α	Ν	"->"	4	Ν	С	I	Р	Н	Ε	R
6	Т	R	Α	Ν	S	Р	0	"->"	5	Т	Н	R	0	W	Α	Ν
9	Р	0	S	ı	Т	T	0	"->"	6	Т	R	Α	Ν	S	Р	0

Row 11 finishes with "TRANS", the only remaining row that completes an actual word is row 9. It makes "TRANSPOSITION".

Row 4 is the only one with the necessary "N" to finish "POSITION". It also adds a full word afterwards (CIPHER) that makes logical sense in the sentence.

f)	5	1	4	6	3	7	2			5	1	4	6	3	7	2
7	Т	Н	1	S	1	S	Α	"->"	7	Т	Н	1	S	1	S	Α
10	Ν	Ε	Χ	Α	М	Р	L	"->"	10	N	Ε	Χ	Α	М	Р	L
8	Ε	0	F	D	0	U	В	"->"	8	Е	0	F	D	0	U	В
11	L	Е	Т	R	Α	Ν	S	"->"	11	L	Ε	Т	R	Α	Ν	S
9	Р	0	S	1	Т	1	0	"->"	9	Р	0	S	1	Т	1	О
4	N	С	1	Р	Н	Е	R	"->"	4	N	С	1	Р	Н	Е	R
1	U	S	1	N	G	В	0	"->"	1	U	S	1	N	G	В	O
2	D	С	0	L	U	M	Ν	"->"	2	D	С	0	L	U	М	Ν
3	S	1	Т	I	0	Ν	Χ	"->"	3	S	I	Т	I	0	Ν	Χ
5	Т	Н	R	0	W	Α	Ν	"->"	5	Т	Н	R	0	W	Α	Ν
6	Т	R	Α	N	S	Р	0	"->"	6	Т	R	Α	Ν	S	Р	0

Row 4 ends with "CIPHER", so now we must start a new word. Row 1, 5 and 6 make words at the start of the row. Row 1 makes the most sense as it progresses the sentence logically (Adds "USING").

g) 4 3 7 4 6 3 7 2 5 1 6 2 5 1 S S "->" S 7 Н 7 Н S Т Α Α "->" 10 Р Р Ε Χ Α Μ L 10 Ν Ε Χ Α Μ L Ν Ε F D 0 U В "->" 8 E F D 0 U В 8 0 0 "->" Т S E Т S 11 L E R Α Ν 11 L R Α Ν Р "->" Р Т 0 9 0 S Т Т 0 9 0 S 1 Т "->" 4 Р 4 Р Ν C T Н Ε R Ν C Н Ε R "->" S 1 U S Ν G В 0 1 U N G В 0 "->" 2 D C 0 L U Μ 5 Т Н R 0 W Α Ν Ν "->" 3 S 2 D C Τ ı 0 Ν Χ 0 L U Ν 1 M Т 0 "->" 5 Н R W Α Ν 3 S T 0 Ν Χ 1 1 "->" 6 Т R Α Ν S Ρ 0 6 Т R Α Ν S Ρ 0

Row 1 ends with an unfished word. Row 5 is the only row that makes a logical word.

h) 5 4 3 7 5 4 6 3 7 2 1 6 2 1 "->" S 7 Т Н S Α 7 Т Н S S Α 10 N Ε Χ Α Μ Р L "->" 10 Ν Ē Χ Α Μ P L "->" F 0 U F U 8 Ε 0 D В 8 Ε 0 D 0 В "->" S S 11 L Ε Т R Α N 11 L Ε Т R Α N Р S Т "->" 9 Р S Т 0 9 0 L 0 0 Τ Ī "->" Р C R 4 N C Τ Н Ε R Ν P Н E 4 "->" 1 U S Т N G В 0 1 U S Ν G В 0 5 "->" 5 T Т Н R 0 W Α Ν Н R 0 W Α Ν "->" 2 L D C 0 U 2 D C 0 L U M N M Ν "->" 3 S Τ ı 0 Ν Χ 3 S ı Т ı 0 Χ ı Ν Т Α Р "->" Т R S Р 0 6 R Ν S 0 6 Α Ν

Row 2 would complete a phrase; "row and column".

i) 5 1 4 6 3 7 2 5 1 4 6 3 7 2 "->" 7 Т Н Ī S S Α 7 Т Н S S Α "->" Ē Ē P 10 N Χ Α Μ P L 10 Ν Χ Α Μ L "->" 8 Ε 0 F D 0 U В 8 Ε 0 F D 0 U В "->" E Т R N S L Ε Т R S 11 L Α 11 Α N 9 Р S "->" 9 Р S T 0 0 Т 0 0 Ī "->" 4 C L Р Н Ε R 4 Ν C Р н E R Ν "->" U S Τ В U S G В 0 1 Ν G 0 1 N "->" 5 Т 5 Т R 0 W Н R W Α Ν Н Α Ν 0 "->" 2 D C 0 U Ν 2 D C 0 L U Ν M M "->" Τ 0 Т Ν S Р 3 S ı ı Ν Χ 6 R Α 0 "->" Ν S S 6 Т R Р 0 3 Т 0 Χ Α Ν

Row 2 ends with "COLUMN", so we will have to start a new word. Row 6 is the only one to start an actual word. This also makes sense because row 3 would finish off row 6's word.

j) Fully decrypted and translated cipher text:

	5	1	4	6	3	7	2
7	Т	Н	1	S	1	S	Α
10	N	Ε	Χ	Α	М	Р	L
8	Ε	0	F	D	0	U	В
11	L	Ε	Т	R	Α	Ν	S
9	Р	0	S	1	Т	1	0
4	N	С	1	Р	Н	Ε	R
1	U	S	1	N	G	В	0
5	Т	Н	R	0	W	Α	N
2	D	С	0	L	U	М	N
6	Т	R	Α	N	S	Р	0
3	S	1	Т	1	0	N	Χ

[&]quot;This is an example of a double transposition cipher using both row and column transposition." $\label{eq:contraction}$

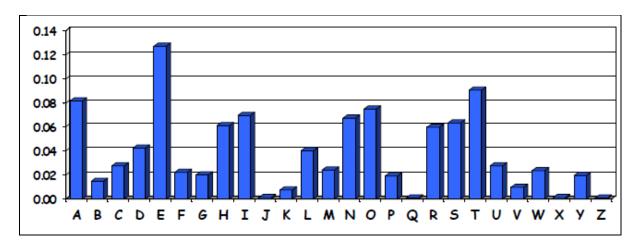
b) [Cryptanalysis on Simple Substitution Cipher]

Cipher Text:	U	М	Υ	Т	М	Н	U	Z	S	R	G	Z
n = 25	٧	N	Z	U	N	1	٧	Α	Т	S	Н	Α
n = 24	W	0	Α	٧	0	J	W	В	U	Т	1	В
n = 23	Χ	Р	В	W	Р	K	Χ	С	٧	U	J	С
n = 22	Υ	Q	С	Χ	Q	L	Υ	D	W	V	K	D
n = 21	Z	R	D	Υ	R	М	Z	Ε	Χ	W	L	Ε
n = 20	Α	S	Ε	Z	S	N	Α	F	Υ	Χ	М	F
n = 19	В	Т	F	Α	Т	0	В	G	Z	Υ	Ν	G
n = 18	С	U	G	В	U	Р	С	Н	Α	Z	0	Н
n = 17	D	٧	Н	С	٧	Q	D	I	В	Α	Р	I
n = 16	Ε	W	I	D	W	R	Ε	J	С	В	Q	J
n = 15	F	Χ	J	Ε	Χ	S	F	K	D	С	R	K
n = 14	G	Υ	K	F	Υ	T	G	L	Ε	D	S	L
n = 13	Н	Z	L	G	Z	U	Н	M	F	Ε	Т	M
n = 12	1	Α	M	Н	Α	٧	1	N	G	F	U	N
n = 11	J	В	Ν	1	В	W	J	0	Н	G	V	0
n = 10	K	С	0	J	С	Χ	K	Р	1	Н	W	Р
n = 9	L	D	Р	K	D	Υ	L	Q	J	I	Χ	Q
n = 8	M	Ε	Q	L	Ε	Z	М	R	K	J	Υ	R
n = 7	Ν	F	R	M	F	Α	N	S	L	K	Z	S
n = 6	0	G	S	Ν	G	В	0	Т	М	L	Α	Т
n = 5	Р	Н	Т	0	Н	С	Р	U	N	M	В	U
n = 4	Q	I	U	Р	1	D	Q	V	0	N	С	٧
n = 3	R	J	٧	Q	J	Ε	R	W	Р	0	D	W
n = 2	S	K	W	R	K	F	S	Χ	Q	Р	Ε	Χ
n = 1	Т	L	Χ	S	L	G	Т	Υ	R	Q	F	Υ
Plain Text:	l am	n hav	ing f	un								

c) [Cryptanalysis on Substitution Cipher]

XDJ MJXDGSGQGVK FJDHNS OBJECJNTK UNUQKWHW BJQHJW GN XDJ OUTX XDUX HN UNK QUNVCUVJ, JUTD QJXXJB DUW HXW GRN ZJBWGNUQHXK. XDJ MGWX GFIHGCW XBUHX XDUX QJXXJBW DUIJ HW XDJ OBJECJNTK RHXD RDHTD XDJK UZZJUB HN U QUNVCUVJ. TQJUBQK HN JNVQHWD XDJ QJXXJB "L" UZZJUBW OUB QJWW OBJECJNXQK XDUN, WUK, "U". HN XHMJW VGNJ FK, HO KGC RUNXJS XG OHNS GCX XDJ OBJECJNTHJW GO QJXXJBW RHXDHN U QUNVCUVJ, KGC DUS XG OHNS U QUBVJ ZHJTJ GO XJAX UNS TGCNX JUTD OBJECJNTK. NGR, DGRJIJB, RJ DUIJ TGMZCXJBW XDUX TUN SG XDJ DUBS RGBP OGB CW. FCX HN OUTX, RJ SGN'X JIJN NJJS XG SG XDHW WXJZ, UW OGB MGWX QUNVCUVJW XDJBJ UBJ SUXUFUWJW GO XDJ QJXXJB OBJECJNTHJW, RDHTD DUIJ FJJN TUQTCQUXJS FK QGGPHNV UX MHQQHGNW GO XJAXW. UNS UBJ XDCW IJBK DHVDQK UTTCBUXJ.

Letter Frequency										
A = 2	H = 31	O = 18	V = 14							
B = 30	I = 7	P = 2	W = 31							
C = 21	J = 80	Q = 25	X = 60							
D = 36	K = 17	R = 11	Y = 0							
E = 6	L = 1	S = 15	Z = 8							
F = 7	M = 6	T = 20								
G = 34	N = 39	U = 54								



- Based on the frequency of letters in the English language, "J" (Which has a frequency of 80 in the passage) is most likely "E"

XDE MEXDGSGQGVK FEDHNS OBEECENTK UNUQKWHW BEQHEW GN XDE OUTX XDUX HN UNK QUNVCUVE, EUTD QEXXEB DUW HXW GRN ZEBWGNUQHXK. XDE MGWX GFIHGCW XBUHX XDUX QEXXEBW DUIE HW XDE OBEECENTK RHXD RDHTD XDEK UZZEUB HN U QUNVCUVE. TQEUBQK HN ENVQHWD XDE QEXXEB "L" UZZEUBW OUB QEWW OBEECENXQK XDUN, WUK, "U". HN XHMEW VGNE FK, HO KGC RUNXES XG OHNS GCX XDE OBEECENTHEW GO QEXXEBW RHXDHN U QUNVCUVE, KGC DUS XG OHNS U QUBVE ZHETE GO XEAX UNS TGCNX EUTD OBEECENTK. NGR, DGREIEB, RE DUIE TGMZCXEBW XDUX TUN SG XDE DUBS RGBP OGB CW. FCX HN OUTX, RE SGN'X EIEN NEES XG SG XDHW WXEZ, UW OGB MGWX QUNVCUVEW XDEBE UBE SUXUFUWEW GO XDE QEXXEB OBEECENTHEW, RDHTD DUIE FEEN TUQTCQUXES FK QGGPHNV UX MHQQHGNW GO XEAXW, UNS UBE XDCW IEBK DHVDQK UTTCBUXE.

- From some observation, it appears "XDE" is a common word in the passage, likely being "The", thus "X" and "D" could be "T" and "H" respectively
- Furthermore, "X" in the passage has a frequency of 60 and "T" is the second most common letter in the English language

THE METHGSGQGVK FEHHNS OBECCENTK UNUQKWHW BEQHEW GN THE OUTT THUT HN UNK QUNVCUVE, EUTH QETTEB HUW HTW GRN ZEBWGNUQHTK. THE MGWT GFIHGCW TBUHT THUT QETTEBW HUIE HW THE OBECCENTK RHTH RHHTH THEK UZZEUB HN U QUNVCUVE. TQEUBQK HN ENVQHWH THE QETTEB "L" UZZEUBW OUB QEWW OBECCENTQK THUN, WUK, "U". HN THMEW VGNE FK, HO KGC RUNTES TG OHNS GCT THE OBECCENTHEW GO QETTEBW RHTHHN U QUNVCUVE, KGC HUS TG OHNS U QUBVE ZHETE GO TEAT UNS TGCNT EUTH OBECCENTK. NGR, HGREIEB, RE HUIE TGMZCTEBW THUT TUN SG THE HUBS RGBP OGB CW. FCT HN OUTT, RE SGN'T EIEN NEES TG SG THHW WTEZ, UW OGB MGWT QUNVCUVEW THEBE UBE SUTUFUWEW GO THE QETTEB OBECCENTHEW, RHHTH HUIE FEEN TUQTCQUTES FK QGGPHNV UT MHQQHGNW GO TEATW, UNS UBE THCW IEBK HHVHQK UTTCBUTE.

- "U" has a frequency of 54 in the passage and "A" is the next most common letter in the English language, thus "U" could be "A"

THE METHGSGQGVK FEHHNS OBECENTK ANAQKWHW BEQHEW GN THE OATT THAT HN ANK QANVCAVE, EATH QETTEB HAW HTW GRN ZEBWGNAQHTK. THE MGWT GFIHGCW TBAHT THAT QETTEBW HAIE HW THE OBECENTK RHTH RHHTH THEK AZZEAB HN A QANVCAVE. TQEABQK HN ENVQHWH THE QETTEB "L" AZZEABW OAB QEWW OBECENTQK THAN, WAK, "A". HN THMEW VGNE FK, HO KGC RANTES TG OHNS GCT THE OBECCENTHEW GO QETTEBW RHTHHN A QANVCAVE, KGC HAS TG OHNS A QABVE ZHETE GO TEAT ANS TGCNT EATH OBECCENTK. NGR, HGREIEB, RE HAIE TGMZCTEBW THAT TAN SG THE HABS RGBP OGB CW. FCT HN OATT, RE SGN'T EIEN NEES TG SG THHW WTEZ, AW OGB MGWT QANVCAVEW THEBE ABE SATAFAWEW GO THE QETTEB OBECCENTHEW, RHHTH HAIE FEEN TAQTCQATES FK QGGPHNV AT MHQQHGNW GO TEATW, ANS ABE THCW IEBK HHVHQK ATTCBATE.

- "GO" appears a lot, this could be a common two letter word like "Of"

THE METHOSOQOVK FEHHNS FBECCENTK ANAQKWHW BEQHEW ON THE FATT THAT HN ANK QANVCAVE, EATH QETTEB HAW HTW ORN ZEBWONAQHTK. THE MOWT OF HOCW TBAHT THAT QETTEBW HAIE HW THE FBECCENTK RHTH RHHTH THEK AZZEAB HN A QANVCAVE. TQEABQK HN ENVQHWH THE QETTEB "L" AZZEABW FAB QEWW FBECCENTQK THAN, WAK, "A". HN THMEW VONE FK, HF KOC RANTES TO FHNS OCT THE FBECCENTHEW OF QETTEBW RHTHHN A QANVCAVE, KOC HAS TO FHNS A QABVE ZHETE OF TEAT ANS TOCNT EATH FBECCENTK. NOR, HOREIEB, RE HAIE TOMZCTEBW THAT TAN SO THE HABS ROBP FOB CW. FCT HN FATT, RE SON'T EIEN NEES TO SO THHW WTEZ, AW FOB MOWT QANVCAVEW THEBE ABE SATAFAWEW OF THE QETTEB FBECCENTHEW, RHHTH HAIE FEEN TAQTCQATES FK QOOPHNV AT MHQQHONW OF TEATW, ANS ABE THCW IEBK HHVHQK ATTCBATE.

- "FATT" could become "Fact". Thus "T" could be "C"

THE METHOSOQOVK FEHHNS FBEECENCK ANAQKWHW BEQHEW ON THE FACT THAT HN ANK QANVCAVE, EACH QETTEB HAW HTW ORN ZEBWONAQHTK. THE MOWT OF HOCW TBAHT THAT QETTEBW HAIE HW THE FBEECENCK RHTH RHHCH THEK AZZEAB HN A QANVCAVE. CQEABQK HN ENVQHWH THE QETTEB "L" AZZEABW FAB QEWW FBEECENTQK THAN, WAK, "A". HN THMEW VONE FK, HF KOC RANTES TO FHNS OCT THE FBEECENCHEW OF QETTEBW RHTHHN A QANVCAVE, KOC HAS TO FHNS A QABVE ZHECE OF TEAT ANS COCNT EACH FBEECENCK. NOR, HOREIEB, RE HAIE COMZCTEBW THAT CAN SO THE HABS ROBP FOB CW. FCT HN FACT, RE SON'T EIEN NEES TO SO THHW WTEZ, AW FOB MOWT QANVCAVEW THEBE ABE SATAFAWEW OF THE QETTEB FBEECENCHEW, RHHCH HAIE FEEN CAQCCQATES FK QOOPHNV AT MHQQHONW OF TEATW, ANS ABE THCW IEBK HHVHQK ACCCBATE.

- The last word in the passage looks to be "Accurate". Thus "C" could be "U" and "B" could be "R"

THE METHOSOQOVK FEHHNS FREEUENCK ANAQKWHW REQHEW ON THE FACT THAT HN ANK QANVUAVE, EACH QETTER HAW HTW ORN ZERWONAQHTK. THE MOWT OF HOUW TRAHT THAT QETTERW HAIE HW THE FREEUENCK RHTH RHHCH THEK AZZEAR HN A QANVUAVE. CQEARQK HN ENVQHWH THE QETTER "L" AZZEARW FAR QEWW FREEUENTQK THAN, WAK, "A". HN THMEW VONE FK, HF KOU RANTES TO FHNS OUT THE FREEUENCHEW OF QETTERW RHTHHN A QANVUAVE, KOU HAS TO FHNS A QARVE ZHECE OF TEAT ANS COUNT EACH FREEUENCK. NOR, HOREIER, RE HAIE COMZUTERW THAT CAN SO THE HARS RORP FOR UW. FUT HN FACT, RE SON'T EIEN NEES TO SO THHW WTEZ, AW FOR MOWT QANVUAVEW THERE ARE SATAFAWEW OF THE QETTER FREEUENCHEW, RHHCH HAIE FEEN CAQCUQATES FK QOOPHNV AT MHQQHONW OF TEATW, ANS ARE THUW IERK HHVHQK ACCURATE.

- **"FREEUENCJEW"** seems to match "frequencies". Thus "E" would be "Q", "N" would be "N", "H" would be "I", and "W" would be "S"

THE METHOSOQOVK FEHINS FREQUENCK ANAQKSIS REQIES ON THE FACT THAT IN ANK
QANVUAVE, EACH QETTER HAS ITS ORN ZERSONAQITK. THE MOST OFIIOUS TRAIT THAT QETTERS
HAIE IS THE FREQUENCK RITH RHICH THEK AZZEAR IN A QANVUAVE. CQEARQK IN ENVQISH THE
QETTER "L" AZZEARS FAR QESS FREQUENTQK THAN, SAK, "A". IN TIMES VONE FK, IF KOU RANTES
TO FINS OUT THE FREQUENCIES OF QETTERS RITHIN A QANVUAVE, KOU HAS TO FINS A QARVE
ZIECE OF TEAT ANS COUNT EACH FREQUENCK. NOR, HOREIER, RE HAIE COMZUTERS THAT CAN SO
THE HARS RORP FOR US. FUT IN FACT, RE SON'T EIEN NEES TO SO THIS STEZ, AS FOR MOST
QANVUAVES THERE ARE SATAFASES OF THE QETTER FREQUENCIES, RHICH HAIE FEEN
CAQCUQATES FK QOOPINV AT MIQQIONS OF TEATS, ANS ARE THUS IERK HIVHQK ACCURATE.

- "CAQCUQATES" and "QETTERS" indicates that "Q" is "L"

THE METHOSOLOVK FEHINS FREQUENCK ANALKSIS RELIES ON THE FACT THAT IN ANK LANVUAVE, EACH LETTER HAS ITS ORN ZERSONALITK. THE MOST OFIIOUS TRAIT THAT LETTERS HAIE IS THE FREQUENCK RITH RHICH THEK AZZEAR IN A LANVUAVE. CLEARLK IN ENVLISH THE LETTER "L" AZZEARS FAR LESS FREQUENTLK THAN, SAK, "A". IN TIMES VONE FK, IF KOU RANTES TO FINS OUT THE FREQUENCIES OF LETTERS RITHIN A LANVUAVE, KOU HAS TO FINS A LARVE ZIECE OF TEAT ANS COUNT EACH FREQUENCK. NOR, HOREIER, RE HAIE COMZUTERS THAT CAN SO THE HARS RORP FOR US. FUT IN FACT, RE SON'T EIEN NEES TO SO THIS STEZ, AS FOR MOST LANVUAVES THERE ARE SATAFASES OF THE LETTER FREQUENCIES, RHICH HAIE FEEN CALCULATES FK LOOPINV AT MILLIONS OF TEATS, ANS ARE THUS IERK HIVHLK ACCURATE.

- "K" is at the end of words that would spell "Frequently, "Any" etc. if it was "Y"

THE METHOSOLOVY FEHINS FREQUENCY ANALYSIS RELIES ON THE FACT THAT IN ANY LANVUAVE, EACH LETTER HAS ITS ORN ZERSONALITY. THE MOST OFIIOUS TRAIT THAT LETTERS HAIE IS THE FREQUENCY RITH RHICH THEY AZZEAR IN A LANVUAVE. CLEARLY IN ENVLISH THE LETTER "L" AZZEARS FAR LESS FREQUENTLY THAN, SAY, "A". IN TIMES VONE FY, IF YOU RANTES TO FINS OUT THE FREQUENCIES OF LETTERS RITHIN A LANVUAVE, YOU HAS TO FINS A LARVE ZIECE OF TEAT ANS COUNT EACH FREQUENCY. NOR, HOREIER, RE HAIE COMZUTERS THAT CAN SO THE HARS RORP FOR US. FUT IN FACT, RE SON'T EIEN NEES TO SO THIS STEZ, AS FOR MOST LANVUAVES THERE ARE SATAFASES OF THE LETTER FREQUENCIES, RHICH HAIE FEEN CALCULATES FY LOOPINV AT MILLIONS OF TEATS, ANS ARE THUS IERY HIVHLY ACCURATE.

- "R**HICH HAIE FEEN"** seems to be "Which have been". Thus, "R" would be "W", "I" would be "V" and "F" would be "B"

THE METHOSOLOVY BEHINS FREQUENCY ANALYSIS RELIES ON THE FACT THAT IN ANY LANVUAVE, EACH LETTER HAS ITS OWN ZERSONALITY. THE MOST OBVIOUS TRAIT THAT LETTERS HAVE IS THE FREQUENCY WITH WHICH THEY AZZEAR IN A LANVUAVE. CLEARLY IN ENVLISH THE LETTER "L" AZZEARS FAR LESS FREQUENTLY THAN, SAY, "A". IN TIMES VONE BY, IF YOU WANTES TO FINS OUT THE FREQUENCIES OF LETTERS WITHIN A LANVUAVE, YOU HAS TO FINS A LARVE ZIECE OF TEAT ANS COUNT EACH FREQUENCY. NOW, HOWEVER, WE HAVE COMZUTERS THAT CAN SO THE HARS WORP FOR US. BUT IN FACT, WE SON'T EVEN NEES TO SO THIS STEZ, AS FOR MOST LANVUAVES THERE ARE SATABASES OF THE LETTER FREQUENCIES, WHICH HAVE BEEN CALCULATES BY LOOPINV AT MILLIONS OF TEATS, ANS ARE THUS VERY HIVHLY ACCURATE.

- "SO" appears a lot, which would be fine if "S'ONT" didn't appear. Thus, "S" would be "D"

THE METHODOLOVY BEHIND FREQUENCY ANALYSIS RELIES ON THE FACT THAT IN ANY LANVUAVE, EACH LETTER HAS ITS OWN ZERSONALITY. THE MOST OBVIOUS TRAIT THAT LETTERS HAVE IS THE FREQUENCY WITH WHICH THEY AZZEAR IN A LANVUAVE. CLEARLY IN ENVLISH THE LETTER "L" AZZEARS FAR LESS FREQUENTLY THAN, SAY, "A". IN TIMES VONE BY, IF YOU WANTED TO FIND OUT THE FREQUENCIES OF LETTERS WITHIN A LANVUAVE, YOU HAD TO FIND A LARVE ZIECE OF TEAT AND COUNT EACH FREQUENCY. NOW, HOWEVER, WE HAVE COMZUTERS THAT CAN DO THE HARD WORP FOR US. BUT IN FACT, WE DON'T EVEN NEED TO DO THIS STEZ, AS FOR MOST LANVUAVES THERE ARE DATABASES OF THE LETTER FREQUENCIES, WHICH HAVE BEEN CALCULATED BY LOOPINV AT MILLIONS OF TEATS, AND ARE THUS VERY HIVHLY ACCURATE.

- "M" and "P" make words such as "Methodology", "Looking", "Most" and "Computers" work; thus, they stay the same
- In saying that, if that is true; we can see that "Z" would be "P" for "Computers" and "Personality", "V" would be "G" for "Looking", "English" and "Methodology", and "P" would be "K" for "Looking and "Work"

THE METHODOLOGY BEHIND FREQUENCY ANALYSIS RELIES ON THE FACT THAT IN ANY LANGUAGE, EACH LETTER HAS ITS OWN PERSONALITY. THE MOST OBVIOUS TRAIT THAT LETTERS HAVE IS THE FREQUENCY WITH WHICH THEY APPEAR IN A LANGUAGE. CLEARLY IN ENGLISH THE LETTER "L" APPEARS FAR LESS FREQUENTLY THAN, SAY, "A". IN TIMES GONE BY, IF YOU WANTED TO FIND OUT THE FREQUENCIES OF LETTERS WITHIN A LANGUAGE, YOU HAD TO FIND A LARGE PIECE OF TEAT AND COUNT EACH FREQUENCY. NOW, HOWEVER, WE HAVE COMPUTERS THAT CAN DO THE HARD WORK FOR US. BUT IN FACT, WE DON'T EVEN NEED TO DO THIS STEP, AS FOR MOST LANGUAGES THERE ARE DATABASES OF THE LETTER FREQUENCIES, WHICH HAVE BEEN CALCULATED BY LOOKING AT MILLIONS OF TEATS, AND ARE THUS VERY HIGHLY ACCURATE.

- In logical sense "L" would be "Z" for the A to Z comparison in the sentence, that leaves "A" in the passage, which fits into "TEAT(S)", making "A" "X"

THE METHODOLOGY BEHIND FREQUENCY ANALYSIS RELIES ON THE FACT THAT IN ANY LANGUAGE, EACH LETTER HAS ITS OWN PERSONALITY. THE MOST OBVIOUS TRAIT THAT LETTERS HAVE IS THE FREQUENCY WITH WHICH THEY APPEAR IN A LANGUAGE. CLEARLY IN ENGLISH THE LETTER "Z" APPEARS FAR LESS FREQUENTLY THAN, SAY, "A". IN TIMES GONE BY, IF YOU WANTED TO FIND OUT THE FREQUENCIES OF LETTERS WITHIN A LANGUAGE, YOU HAD TO FIND A LARGE PIECE OF TEXT AND COUNT EACH FREQUENCY. NOW, HOWEVER, WE HAVE COMPUTERS THAT CAN DO THE HARD WORK FOR US. BUT IN FACT, WE DON'T EVEN NEED TO DO THIS STEP, AS FOR MOST LANGUAGES THERE ARE DATABASES OF THE LETTER FREQUENCIES, WHICH HAVE BEEN CALCULATED BY LOOKING AT MILLIONS OF TEXTS, AND ARE THUS VERY HIGHLY ACCURATE.

Question 2 [Application of Hash Algorithm]

- Since Trudy knows the range of the bids, she can calculate the hash values of \$95 to \$103 using the SHA-256 Algorithm as well.
 - 2) She can then compare her calculated hash values of \$95, \$103 and the intervals between with Alice and Bob hash values.
 - 3) Once she determines the bids of Alice and Bob, all Trudy needs to do is bit slightly higher than both
 - 4) Trudy wins the auction
- To prevent people like Trudy, the auction process can be modified in how they encode the bids into has values. Have bidders add a number to the bid amount (E.g. \$100 + 51 = 151) and encode that using SHA-256. Once bidding is over, have the bidders send their additional number to the auctioneer; who will then subtract the number from the hashed value to find the bid amount. This way, if people like Trudy figure out bidder's hashed values; they still won't know the actual bid without knowledge of the additional umber

Question 3 [RSA Encryption algorithm]

a) Generating the Public Key

$$p = 443$$

$$q = 503$$

$$n = 443 \times 503 = 222,829$$

$$\varphi(n) = (443 - 1) \times (503 - 1) = 442 \times 502 = 221,884$$

$$e = 3$$
 (221,884 ÷ 3 = 73,961.33) : $gcd(221884, 3) = 1$

Public Key: (222829, 3)

b) Generating the Private Key

$$\varphi(n) = 221,884$$

$$e = 3$$

 $d \times 3 = 1 \mod 221,884$

Modular Multiplicative Inverse

 Integer
 Modulo

 3
 221884

CALCULATE

Modular Multiplicative Inverse

147923

$$d = 147,923$$

c) Encrypting the Message with the Public Key

$$n = 222,829$$

$$e = 3$$

$$M = 7415$$

 $C = 7415^3 \mod 222,829$

MATH 139 SPRING 2003

PowerMod Calculator

Computes $(base)^{(exponent)} \mod (modulus)$ in log(exponent) time.

Base: 7415	Exponent: 3	Modulus: 222829
Compute	$b^e \text{ MOD } m =$	134908

The program is written in JavaScript, and runs on the client computer. Most implementations seem to handle numbers of up to 16 digits correctly

C = 134908

d) Decrypting the Message with the Private Key

$$n = 222,829$$

$$C = 8166$$

$$d = 147,923$$

 $M = 134908^{147,923} \mod 222,829$

MATH 139					Ģ	SPRING 2003
PowerMod Calculator						
Computes (base) ^(exponent) mod (modulus) ir	n log(exponent) time.					
	Base: 134908	Exponent: 147923		Modulus: 222829		
	Compute		$b^e \text{ MOD } m =$	7415		
The program is written in JavaScript, and runs on the client computer.	Most implementations seem to handle numbers of up t	to 16 digits correctly:				

M = 7415

Question 4 [Breaking RSA Encryption algorithm]

a) Determining p and q

$$C = 1602$$

$$n = 3901$$

$$e = 11$$

$$p \times q = 3901 : p = 47, q = 83 \text{ (Both prime)}$$

b) Determining $\phi(n)$

$$p = 47$$

$$q = 83$$

$$\varphi(n) = (47 - 1) \times (83 - 1) = 46 \times 82 = 3,772$$

c) Determining the Private Key

$$\varphi(n) = 3,772$$

$$e = 11$$

$$d \times 11 = 1 \mod 3,772$$

Modular Multiplicative Inverse

 Integer
 Modulo

 11
 3772

CALCULATE

 $\begin{array}{c} \text{Modular Multiplicative Inverse} \\ 343 \end{array}$

d) Decrypt the Message using the Private Key n = 3901C = 1602d = 343 $M = 1602^{343} \mod 3901$ MATH 139 SPRING 2003 PowerMod Calculator mod (modulus) in log(exponent) time. Computes (base)(Base: 1602 Modulus: 3901 Exponent: 343 Compute $b^e \text{ MOD } m = 3$ ster. Most implementations seem to handle numbers of up to 16 digits correctly M = 3**Question 5 [ElGamal Encryption algorithm]** a) Calculating y M = 59p = 8467g = 7919x = 91r = 51 $y = 7919^{91} \mod 8467$ SPRING 2003 PowerMod Calculator ent) mod (modulus) in log(exponent) time Base: 7919 Modulus: 8467 Exponent: 91 $b^{\varepsilon} \text{ MOD } m = 3249$ Compute y = 3249b) Calculating k r = 51y = 3249p = 8467 $k = 3249^{51} \mod 8467$ MATH 139 SPRING 2003 PowerMod Calculator Computes (base)(exponent) mod (modulus) in log(exponent) time. Exponent: 51 Compute $b^e \text{ MOD } m = 2145$

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M = 59

p = 8467

g = 7919

r = 51

k = 2145

 $C_1 = 7919^{51} \mod 8467$

MATH 139 SPRING 2003 PowerMod Calculator Computes (base)(exponent) mod (modulus) in log(exponent) time.

Base: 7919	Exponent: 51	Modulus: 8467
Compute	$b^e \text{ MOD } m =$	6276

 $C_1 = 6276$

 $C_2 = 59 \times 2145 \mod 8467$

MATH 139 SPRING 2003

PowerMod Calculator

 $Computes\ (base)^{(exponent)}\ mod\ (modulus)\ in\ log(exponent)\ time.$

Base: 126555	Exponent: 1	Modulus: 8467	
Compute	$b^e \text{ MOD } m =$	8017	

The program is written in JavaScript, and runs on the client computer. Most implementations seem to handle numbers of up to 16 digits correctly.

 $C_2 = 8017\,$

x = 91 $C_1 = 6276$ p = 8467 $k = 6276^{91} \mod 8467$ MATH 139 SPRING 2003 Computes (base)^(exponent) mod (modulus) in log(exponent) time Modulus: 8467 Base: 6276 Exponent: 91 $b^e \text{ MOD } m = 2145$ Compute rs of up to 16 digits correctly k = 2145 $k^{-1} = 2145^{-1} \mod 8467$ 🔋 Modular Multiplicative Inverse Integer 2145 8467 Modular Multiplicative Inverse 4271 $k^{-1} = 4271$ e) Decrypting the encrypted message $C_2 = 8017$ $k^{-1} = 4271$ p = 8467 $m = 4271 \times 8017 \mod 8467$ MATH 139 SPRING 2003 PowerMod Calculator Computes (base)(exponent) mod (modulus) in log(exponent) time. Base: 34240607 Modulus: 8467 Exponent: 1 $b^e \text{ MOD } m = 59$ Compute m = 59

d) Decrypting the message – calculating k and its multiplicative inverse

Question 6 [Paillier Encryption algorithm]

a) Calculating n

M = 1234

p = 61

q = 97

g = 119

r = 67

 $n = 61 \times 97 = 5917$

b) Calculating 3

p = 61

q = 97

a = lcm(61 - 1, 97 - 1) = lcm(60, 96) = 480

c) Calculating k

g = 119

n = 5917

a = 480

 $k = L(119480 \mod 5917^2) = L(119480 \mod 35010889)$



$$k = L(24508215) = (24508215 - 1) \div 5917 = 4142$$

d) Calculating µ

k = 4142

n = 5917

 $\mu = 4142^{-1} \mod 5917$

📵 Modular Multiplicative Inverse

Integer Modulo 4142 5917

Modular Multiplicative Inverse

10

 $\mu = 10$

e) Encrypting the message

g = 119

M = 1234

r = 67

n = 5917

 $c = 119^{1234} \times 67^{5917} \mod 5917^2 = (119^{1234} \mod 35010889 \times 67^{5917} \mod 35010889) \mod 35010889$ 35010889

MATH 139 SPRING 2003 PowerMod Calculator Computes (base)^(exponent) mod (modulus) in log(exponent) time Base: 119 Exponent: 1234 Modulus: 35010889 $b^e \text{ MOD } m = 29522434$ SPRING 2003 PowerMod Calculator Computes (base)(exponent) mod (modulus) in log(exponent) time. Exponent: 5917 Modulus: 35010889 $b^{\epsilon} \text{ MOD } m = 7510496$ Compute $c = (29522434 \times 7510496) \mod 35010889$

ATH 139				SPRING 2003
PowerMod Calculator Computes (base)(exponent) mod (modulus) in	n log(exponent) time.			
	Base: 221728122467264	Exponent: 1	Modulus: 35010889	
	Compute	$b^e \text{ MOD } m =$	31145362	
The program is written in JavaScript, and runs on the client computer.	Most implementations seem to handle numbers of up	to 16 digits correctly.		

c = 31145362

Decrypting the encrypted message $\lambda = 480$ $\mu = 10$ c = 31145362n = 5917k = 4142 $M = L(31145362^{480} \mod 5917^2) \times 10 \mod 5917$ MATH 139 SPRING 2003 PowerMod Calculator Computes (base)(exponent) mod (modulus) in log(exponent) time. Modulus: 35010889 Base: 31145362 Exponent: 480 $b^e \text{ MOD } m = 28738870$ The program is written in JavaScript, and runs on the client computer. Most impl stations seem to handle numbers of up to 16 digits correctly. $M = L(28738870) \times 10 \mod 5917$ $M = ((28738870 - 1) \div 5917) \times 10 \mod 5917$ $M = 4857 \times 10 \mod 5917$ $M = 48570 \mod 5917$ MATH 139 SPRING 2003 PowerMod Calculator Computes (base)(exponent) mod (modulus) in log(exponent) time.

 $b^e \text{ MOD } m = 1234$

M = 1234

Compute