

Midterm Review

- 1. Branch instructions
 - BHI (unsigned), BGT (signed)
 - Take a look at the preceding comparison instruction.
 - Then, you can use this instead of using complex formula in the instruction reference.
 - BRANCH IF REGISTER IS HIGHER/GREATER/... THAN OPERAND
 - e.g. CMPA #\$D0
 - Branch if 'A' is Hlgher than 'operand.'
 - BMI, BVC, BNE
 - Check the result if it is Minus/signed oVerflow/Not Equal to zero
 - CMPA: A – (M)
 - TSTA: A – 0
 - V bit (signed overflow)
 - $N + N \rightarrow P$
 - $P + P \rightarrow N$
 - $N - P \rightarrow P$
 - $P - N \rightarrow N$

Midterm Review – cont'd

- 2. BRA and LBRA
 - Write a line of assembly code that begins in memory location \$2450.
 - BRA <8-bit signed offset value>
 - LBRA <16-bit signed offset value>
 - When the effective (destination) address is calculated in relative addressing mode, a value that you need to add to an offset is PC.
→ We have to use PC which is the next address of the line of assembly code of (L)BRA

Midterm Review

- 3. Mostly this is about calculating postbyte(s)
 - -20: You cannot use 5-bit offset. You have to use 9-bit one.

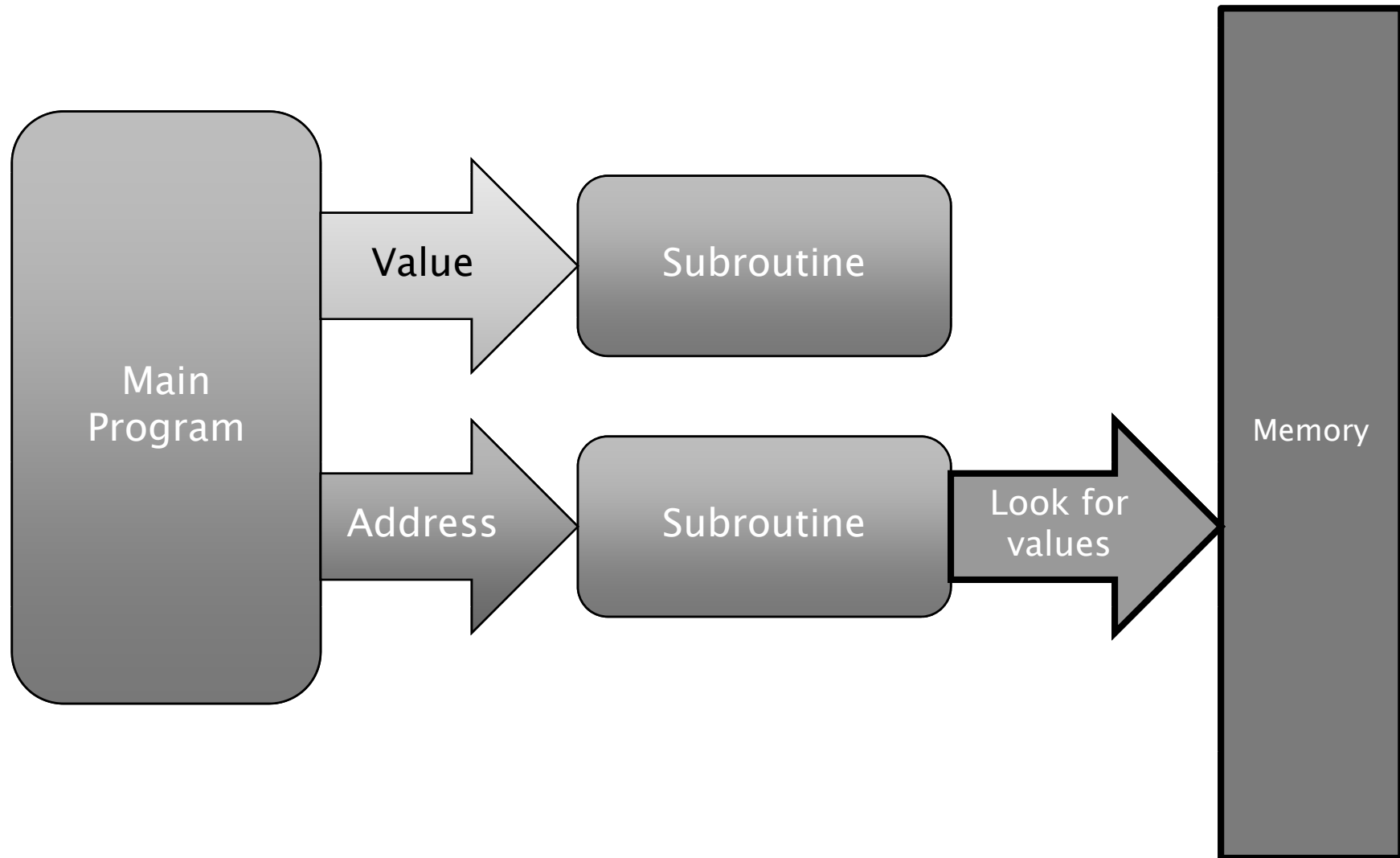
Lecture 16:

Parameter Passing

Today's Goals

- Parameter passing
- Understand how to pass parameters using the stack

Call by Value vs. Call by Reference



Passing Parameters in Registers

- Use registers to pass parameters
- The simplest type of parameter passing
- Pros vs. Cons
 - Data is immediately available to the subroutine
 - No extra preparations are needed.
 - Often fastest execution and smallest code size
 - Only a limited number of registers are available.
 - We can use 'pass-by-reference' to point a list of input/output values.

Example (Parameter Passing in Registers)

- Requirements
 - A subroutine that adds an array of two-byte numbers and a sample main program that calls it.
 - The array is passed by reference in X.
 - The length is passed by value in D.
 - The sum should be returned by value in D.
 - Do not worry about indicating signed or unsigned overflow.

Example (Parameter Passing in Registers)

```
Array    ORG      $3000
         dc.w     $1234,$5678,$ABCD
Length   dc.w     3

         ORG      $2000
         LDS      #$3600
         LDX      #Array    ; load X with address of list
         LDD      Length    ; Load D with actual length
         JSR      sumword
         SWI

sumword   TFR      D,Y
         LDDD     #0
         CPY      #0
Loop      BEQ      endsum
         ADDD     0,X
         INX
         INX
         DEY
         BRA      loop
endsum    RTS
```

Preserving Registers

- In the previous example, the subroutine uses Y. The value in Y is destroyed.
- If Y was in use by the main program, the main program would continue after the subroutine with an incorrect value in Y.
- To avoid this, registers used by a subroutine may be saved to the stack before a subroutine uses them.
- The saved registers are restored after the subroutine.
- We have two options
 - The caller (main program) does this.
 - The callee (subroutine) does this.

Preserving Registers

- Responsibility of the Caller:
 - The calling main program assumes all registers are destroyed by the subroutine
 - If the registers used by the subroutine are unknown (i.e. using a sub. provided by someone else), may save registers that the subroutine wouldn't affect
 - Code to save/restore registers duplicated with every subroutine call.
- Responsibility of the Callee:
 - Saves only the registers that will be used by the subroutine
 - Save/restore code only occurs once

Example – Caller Responsible

```

                                ORG      $3000
Array      dc.w      $1234,$5678,$ABCD
Length     dc.w      3

                                ORG      $2000
                                LDS      #$3600
                                LDX      #Array    ; load X with address of list
                                LDD      Length    ; Load D with actual length
                                PSHX
JSR        sumword
                                PULX
                                SWI

sumword     TFR        D,Y
                                LDDD      #0
                                CPY       #0
loop        BEQ        endsum
                                ADDD      0,X
                                INX
                                INX
                                DEY
                                BRA       loop
endsum     RTS
```

Stack Frame

RAH
RAL
XH
XL

Example – Callee Responsible

```

Array      ORG      $3000
           dc.w      $1234,$5678,$ABCD
Length     dc.w      3
           ORG      $2000
           LDS      #$3600
           LDX      #Array    ; load X with address of list
           LDD      Length    ; Load D with actual length
           JSR      sumword
           SWI

sumword     PSHX
           PSHY
TFR         D,Y
           LDDD      #0
           CPY       #0
loop        BEQ      endsum
           ADDD      0,X
           INX
           INX
           DEY
           BRA       loop
endsum      PULY
           PULX
RTS
```

Why not D?

It is used to return the answer!



Stack Frame

YH
YL
XH
XL
RAH
RAL

Passing Parameters in the Stack

- Pros and Cons
 - The stack pointer is already in use for the return address
 - Indexed addressing can easily access data stored on the stack
 - The amount of data passed is not limited by the register set
 - The data passed on the stack must be removed, and this is usually the responsibility of the caller (although this can be done by the callee)
- Note:
 - If the caller will save registers on the stack, it should be done before passing parameters.
 - Why? Those parameters are supposed to be used in the subroutine.