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COMP5201

Assignment 2

Fall 2020

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Submit electronically. No extension will be granted.

1. [24 marks] Digital Logic.

- a) 'X' is the ternary connective such that 'Xpqr' is logically equivalent to ' $p \wedge (p + q + r)$ '. '+' is 'xor'. 'F' and 'T' denote the 0-ary connectives 'false' and true, respectively. Whenever possible, put letters in alphabetical order, and put letters before any 0-place connectives.

Using {'X', 'F', 'T'}, synthesize: $\sim p \mid = \mid X(T p F)$

p	$\sim p$	X(TpF)
0	1	X(100) = 1
1	0	X(110) = 0

Using {'X', 'T'}, synthesize: $p \wedge q \mid = \mid X(p q T)$

p	q	p and q	X(pqT)
0	0	0	X(001)=0
0	1	0	X(011)=0
1	0	0	X(101)=0
1	1	1	X(111)=0

Using {'X', ' \sim ', 'T'}, synthesize: $p \vee q \mid = \mid \sim X(\sim p \sim q T)$

p	q	q or p	$\sim X(\sim p \sim q T)$
0	0	0	X(111) = 0
0	1	1	X(101) = 1
1	0	1	X(011) = 1
1	1	1	X(011) = 1

b) 'Y' is the ternary connective such that 'Ypqr' is logically equivalent to ' $p \wedge (q \leftrightarrow r)$ '. ' $p \leftrightarrow q$ ' is true iff 'p' and 'q' have the same truth value. 'F' and 'T' denote the two 0-ary connectives 'false' and 'true', respectively. Whenever possible, put letters in alphabetical order, and put letters before any 0-place connectives.

Using {'Y', 'F', 'T'}, synthesize: $\sim p \mid = \mid Y(TpF)$

p	$\sim p$	Y(pqT)
0	1	Y(100) = 1
1	0	Y(110) = 0

Using {'Y', 'F', 'T'}, synthesize $p \wedge q \mid = \mid Y(pqT)$

p	q	p and q	Y(pqT)
0	0	0	Y(001)=0
0	1	0	Y(011)=0
1	0	0	Y(101)=0
1	1	1	Y(111)=0

Using {'Y', ' \sim ', 'F'}, synthesize: $p \vee q \mid = \mid \sim Y(\sim pqF)$

p	q	p and q	Y($\sim pqF$)
0	0	0	Y(100)=0
0	1	1	X(110)=1
1	0	1	X(000)=1
1	1	1	X(010)=1

2. [16 marks] Binary and Hexadecimal Numbers.

Convert each of the following six binary or hexadecimal natural numbers into decimal. Show work. For the last number, show either scientific notation with six decimal places, or the exact integer.

a) Binary numbers: 1111, 1111 1110, 1111 1110 1101 0111

signed:	-7	-126	-32 471
unsigned:	15	254	65 239

1111:

each digit represents an increasing power of 2:

4 bits: $2^4 \rightarrow$ We can represent up to 16 values

The values go from $(n/2)-1$ to $(n/2)-1$ with $n = 16$: $[-7:7]$

Most significant bit is 1 so if it's signed, it's a negative value.

The other bits are 111, so it's the highest negative value: -7

If it's an unsigned number (no negative values) it's $n-1$: 15

1111 1110:

8 bits: $2^8 \rightarrow$ 256 values:

if signed, values are in: $[-127:127]$

if unsigned, values are in $[0:255]$

First bit is 1, and all but the least significant bit are 1:

it's the $-[n-1]$ value where n is the lowest value possible: $-(127-1) = -126$

if it's unsigned, it's the $n-1$ highest value possible: $255-1$: 254

1111 1110 1101 0111:

16 bits: $2^{16} \rightarrow$ the values are in $[0:65\,535]$

$1111\ 1110\ 1101\ 0111 = 1111\ 1111\ 1111\ 1111 - 0000\ 0001\ 0010\ 1000$

$0000\ 0001\ 0010\ 1000$ // from previous question we know that $1\ 0000\ 0000 = 256$

$= 256 + 0010\ 1000$

$= 256 + 32 + 8$

$= 296$

----> $1111\ 1111\ 1111\ 1111 - 0000\ 0001\ 0010\ 1000$

is equal to $65\,535 - 296 = 65\,239$

if signed values:

$$1111\ 1111\ 1111\ 1111 = (65\ 536/2) - 1 = -32\ 767$$

$$- (32\ 767 - 296) = -32\ 471$$

b) Hexadecimal numbers: af, afba, afba 51de

af:

$$a = 10 * (16^1) = 160$$

$$f = 15 * (16^0) = 15$$

$$af = a + f = 160 + 15 = 175$$

afba:

$$a = 10 * (16^3) = 40\ 960$$

$$f = 15 * (16^2) = 3\ 840$$

$$b = 11 * (16^1) = 176$$

$$a = 10 * (16^0) = 10$$

$$afba = a + f + b + a = 44\ 986$$

afba 51de:

$$a = 10 * (16^7) = 2\ 684\ 354\ 560$$

$$f = 15 * (16^6) = 251\ 658\ 240$$

$$b = 11 * (16^5) = 11\ 534\ 336$$

$$a = 10 * (16^4) = 655\ 360$$

$$afba\ 0000 = 2\ 948\ 202\ 496$$

$$5' = 5 * (16^3) = 20\ 480$$

$$1' = 1 * (16^2) = 256$$

$$d = 13 * (16^1) = 208$$

$$e = 14 * (16^0) = 14$$

$$51de = 20\ 958$$

$$afba51de = 2\ 948\ 202\ 496 + 20\ 958 = 2\ 948\ 223\ 454$$

3. [25 marks] Fractional Numbers and Blackboard Notation.

Infinite binary expansions of rational numbers are either pure recurring or mixed recurring depending on whether the cycle starts immediately after the point.

This exercise is done assuming $6 \frac{4}{9}$ means - 6 is the integer part

and $\frac{4}{9}$ the fractional part

a) [math] Show the infinite binary expansion of $6 \frac{4}{9}$ without normalization.

6 in binary is 110

$6 \frac{4}{9}$ is equal to $\frac{58}{9}$ which is an endless series of 4 after the .

An approximation of $6 \frac{4}{9}$ equal like $6 + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-8} + 2^{-9} + 2^{-10} + 2^{-14} + 2^{-15} + 2^{-16}$

gives: 6.4444427490234375 in decimal and 110.0111000111000111 in binary

There is a regularity in the alternation between the 1s and 0s

and since the decimal expression of $\frac{58}{9}$ has no end, so does its binary expression

b) [math] Show this infinite binary expansion in hexadecimal without normalization.

With the longer decimal expression:

$$6 + 2^{-2} + 2^{-3} + 2^{-4} + 2^{-8} + 2^{-9} + 2^{-10} + 2^{-14} + 2^{-15} + 2^{-16} + 2^{-20} + 2^{-21} + 2^{-22} + 2^{-26} + 2^{-27} + 2^{-28} + 2^{-32} + 2^{-33} + 2^{-34} + 2^{-35}$$

we get a better approximation: 6.444444444438

Let's convert the integer part to hexadecimal:

6 -> 6

Let's convert the fractional part to hexadecimal:

int part -> hex

0.444444444438	x	16	=	7.11111111008	7
0.11111111008	x	16	=	1.777777776128	1
0.777777776128	x	16	=	12.444444418048	C
0.444444418048	x	16	=	7.111110688768	7
0.111110688768	x	16	=	1.777771020288	1
0.777771020288	x	16	=	12.444336324608	C
0.444336324608	x	16	=	7.109381193728	7
0.109381193728	x	16	=	1.750099099648	1
0.750099099648	x	16	=	12.001585594368	C

.....

We get 6.444444444438 is ~6.71C71C71C in hexadecimal.

Just like in the previous question we can see a repetition pattern for an infinite expansion

c) [math] Show the infinite binary expansion of $6 \frac{4}{9}$ with normalization. Do not forget the scale factor. (This is a scaled infinite binary expansion).

$(110.0111000111000111) * 2^0$ <identity> // for comparison

$(1.100111000111000111) * 2^2$ <normalized>

d) [math] Show this infinite binary expansion in hexadecimal. Again, do not forget the scale factor, which may be shown in decimal.

$1.8 <71C> 2^2$ <normalized>

e) Show the normalized (binary) blackboard notation that best approximates $6 \frac{4}{9}$. The fractional field is 16 bits. Show all 16 bits. Now, show just the 16-bit (4-hexit) fractional field in hexadecimal.

in 16bits:

$1.10<0111000111000111> [2] <"lazy">$ // for comparison

$1.10<0111000111000111> 2^2 <\text{normalized}>$

Fractional field in hexadecimal:

0111 0001 1100 0111 -> <71C7>

4. [15 marks] Integer Multiplication I.

a) Multiply the following 10-bit binary natural numbers. The multiplicand is 10011 11100 (27c hex) and the multiplier is 11010 (1a hex). Show, in hexadecimal, i) the initial value of the accumulator, and ii) each term added to the accumulator, and the partial sum after the addition. The last addition yields the final value.

carry

a----> +1 +4 +7

1----> +0 +0 +0

addition----> +1 +1

sum = 2 5 7

16³ 16² 16¹ 16⁰

2 7 c

1) $a * c = 10 * 12 = 120 = 7 * 16 + 8$

* 1 a

2) $a * 7 = 10 * 7 = 70 = 4 * 16 + 6$

3) $a * 2 = 10 * 2 = 20 = 1 * 16 + 4$

a * 27c = 4 6 8

4) $1 * c = 13 = c$

1 * 27c = 2 7 c

5) $1 * 7 = 7$

6) $1 * 2 = 2$

result = 2 b 2 8

7) $6 + c = 18 = 1 * 16 + 2$

8) $4 + 7 = 11 = b$

acc = 2 5 7

9) $2 + 7 = 9$

10) $b + 5 = 1 * 16 + 0$

2b28 + acc = 4 0 9 8

11) $2 + 2 = 4$

15) $2+2 = 4$

5. [20 marks] Integer Multiplication II.

a) Show that, regardless of the initial n -bit value of the accumulator, the fused multiply-add result of two n -bit natural-number operands is always representable in $2n$ bits. Now, suppose $n = 16$. Starting from the largest possible FMA result, what is the hexadecimal representation of the largest $n = 16$ -bit number that can _still_ be added without producing overflow?

If we're looking for the largest possible accumulator value that can be used without producing overflow:

For a 16 bits bit register:

$$\begin{array}{r} 1111\ 1111\ 1111\ 1111\ //\ 65\ 535 \\ *1111\ 1111\ 1111\ 1111\ //\ 65\ 535 \\ \hline =\ 1111\ 1111\ 1111\ 1110\ 0000\ 0000\ 0000\ 0001\ //\ 4\ 294\ 836\ 225 \end{array}$$

The value x must be such that $x < 2^n - 1$ to avoid overflow, the total value must be less than $2^{32} - 1$: 4 294 967 295

The difference between the max value and the value we can have by multiplying $6535 * 6535$ is:

$$4\ 294\ 967\ 295 - 4\ 294\ 836\ 225 = 131\ 070$$

$131070 > 2^{16} \Rightarrow$ The largest 16 bits number still fit without producing overflow

\Rightarrow 1111 1111 1111 1111 (binary) or FFFF (hexadecimal)

b) A modular-adder device 'M' operates with 16-bit registers. You give it two 16-bit natural numbers 'a' and 'b'. It adds them, divides by 2^{16} , keeps the quotient 'q' a secret, and publishes the remainder 'r'. Hint: Before answering, experiment with small addition tables.

i) If $a = 31,465$ and $r = 53,576$, what are 'b' and 'q'?

to binary:

$a \rightarrow 0111\ 1010\ 1110\ 1001$

$r \rightarrow 1101\ 0001\ 0100\ 1000$

solving: {

$$(31465 + b) / (2^{16}) = q$$

$$(31465 + b) \% (2^{16}) = 53576$$

}

solution: {

$$z = |53576 - 31465| \quad // \text{ absolute value}$$

$$b = z + n \cdot 65536$$

q is a function of b

}

assuming $n = 0$ we have:

$$b = r - a \iff 1101\ 0001\ 0100\ 1000 - 0111\ 1010\ 1110\ 1001$$

$$b \rightarrow 0101\ 0110\ 0101\ 1111$$

$$b = 22\ 111$$

$$a + b < 2^{16} \Rightarrow q = 0$$

assuming $n = 1$ we have:

$$b = 53576 - 31465 + 65536 = 87\ 647$$

$$b \rightarrow 1010\ 1011\ 0010\ 1111\ 1 \rightarrow \text{doesn't fit 16 bits address}$$

\implies We can't have any other solution than $b = 22\ 111$ so $q = 0$

ii) If $a = 35,492$ and $r = 11,087$, what are 'b' and 'q'?

to binary:

$$a \rightarrow 1000\ 1010\ 1010\ 0100$$

$$r \rightarrow 0010\ 1011\ 0100\ 1111$$

solving: {

$$(35492 + b)/(2^{16}) = q$$

$$(35492 + b) \% (2^{16}) = 11087$$

}

solution: {

$$b = 41\ 131 + n \cdot 65536$$

q is a function of b

}

assuming $n = 0$ we have:

$$b = 41\,131$$

$b \rightarrow 1010\,0000\,1010\,1011$

$$a+b > 2^{16}$$

$$(a+b)/(2^{16}) = 1 \quad // \text{ division}$$

$$(a+b)\%(2^{16}) = 11087$$

b is indeed 41131

3

assuming $n = 1$ we have:

$$b = 41\,131 + 65\,536 = 106\,667$$

$b \rightarrow 1101\,0000\,0101\,0101\,1 \rightarrow$ doesn't fit 16 bits address

\implies We can't have any other solution than $b = 41\,131$
