

Homework #2 due October 25, Tuesday 2022, before recitation

(1) Try to simplify the following expressions and justify your simplification.

(i) $(0+1)^*.1.(0+1) + (0+1)^*.1.(0+1)$

(ii) $((0^*.1^*)+1)^*(0+1)^*$

(iii) $(L+M^*)^*$

(iv) $(L.M^*)^*$

(2) Convert the *regular expression* $((0.0^*(1.1))+0.1)^*$ into an ϵ -NFA

(3) Problems from the textbook

3.1.1 (b) and (c)

3.1.4 (b) and (c)

3.2.1 (c) and (d)

3.2.3

Homework 2 :

2- (i) $(0+1)^* \cdot 1 \cdot (0+1) + (0+1)^* \cdot 1 \cdot (0+1)$

using $A = \lambda$ and $\lambda = \lambda$

$$= (0+1)^* \cdot 1 \cdot (0+1) \cdot \lambda + (0+1)^* \cdot 1 \cdot (0+1) \cdot \lambda$$

using $A \cdot b + A \cdot c = A[b+c]$

$$= (0+1)^* \cdot 1 \cdot (0+1) [\lambda + \lambda]$$

$$\rightarrow \lambda + \lambda = \lambda$$

→ using $A = \lambda \cdot A$

$$\boxed{(0+1)^* \cdot 1 \cdot (0+1)}$$

$$(ii) (((0^*1^*)+1)^*)^* (0+1)^*$$

→ The first ^{portion} expression can generate all kind of strings from 0,1

→ The second portion of this expression can generate all kind of strings from 0,1

→ Hence, the given regex refers to all the strings of 0,1.

→ Most simplified version is $(0+1)^*$, which represents the set of all the strings of 0,1.

$$(iii) (L+M^*)^*$$

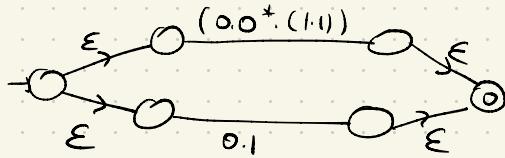
M^* → zero or more occurrences of M. Hence, we can write it as

$$(L+M^*)^* \Rightarrow (L+M)^*. \text{In other words, zero or more occurrences of } L \text{ or } M.$$

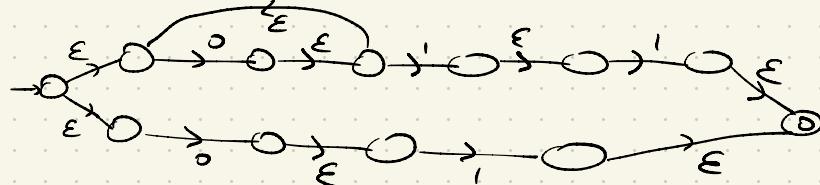
(iv) The answer will be same logic as (iii). Because, $(L \cdot M^*)^*$ and $(L \cdot M)^*$ are same. Hence, it can be simplified as $(L \cdot M)^*$

2- $((0.0^*.(1.1))+01)^*$ into an ϵ -NFA.

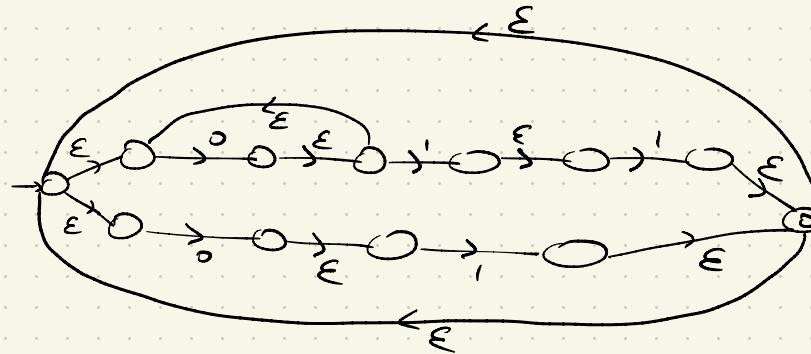
Step 1:



Step 2:



Step 3:



3.1.1

b) The set of strings of 0's and 1's whose tenth symbol from the right end is 1.

$$(0+1)^* 1 (0+1)^9$$

c) The set of strings of 0's and 1's with at most one pair of consecutive 1's.

$$(0+10)^* (11+\epsilon) (0+10)^*$$

3.1.4

b) $(0^* 1^*)^* 000 (0+1)^*$

c) $(0+10)^* 2^*$

Answer: This is the language of strings in which there are no two consecutive 1's, except for possibly a string of 1's at the end.

32.1 c) $R_{ij}^{(2)}$

$$R_{11}^{(2)} = R_{11}^{(1)} + R_{12}^{(1)}(R_{22}^{(1)})^* R_{21}^{(1)} \\ = 1^* + 1 * 0 (\epsilon + 1^+ 0)^* 1^+ = (1+0)^*$$

$$R_{12}^{(2)} = R_{12}^{(1)} + R_{11}^{(1)}(R_{22}^{(1)})^* R_{22}^{(1)} = R_{12}^{(1)}(R_{22}^{(1)})^* \\ = 1^* 0 (\epsilon + 1^+ 0)^* = (1+0)^* 0$$

$$R_{13}^{(2)} = R_{13}^{(1)} + R_{11}^{(1)}(R_{21}^{(1)})^* R_{23}^{(1)} = \emptyset + 1^* 0 (\epsilon + 1^+ 0)^* 0 = (1+0)^* 0 0$$

$$R_{21}^{(2)} = R_{21}^{(1)} + R_{22}^{(1)}(R_{21}^{(1)})^* R_{21}^{(1)} = (R_{21}^{(1)})^* R_{21}^{(1)} = \epsilon + 1^+ 0 | 1^+ = 1^+ (\epsilon + 0 | 1^+)$$

$$R_{22}^{(2)} = R_{22}^{(1)} + R_{21}^{(1)}(R_{21}^{(1)})^* R_{22}^{(1)} = (R_{21}^{(1)})^* R_{22}^{(1)} = (\epsilon + 1^+ 0)^* 0 = (1^+ 0)^*$$

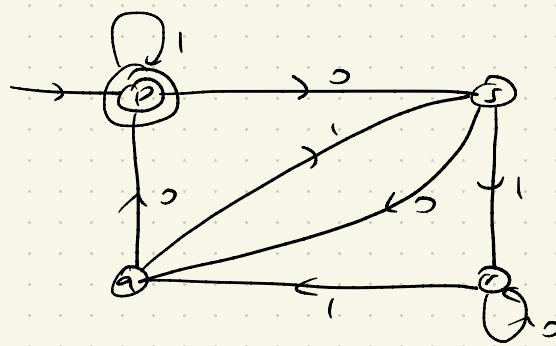
$$R_{23}^{(2)} = R_{23}^{(1)} + R_{21}^{(1)}(R_{22}^{(1)})^* R_{23}^{(1)} = (R_{21}^{(1)})^* R_{23}^{(1)} = (\epsilon + 1^+ 0)^* 0 = (1^+ 0)^* 0$$

$$R_{31}^{(2)} = R_{31}^{(1)} + R_{31}^{(1)}(R_{21}^{(1)})^* R_{21}^{(1)} = \emptyset + 1 (\epsilon + 1^+ 0)^* 1^+ = 1 (1+0)^* 1^+$$

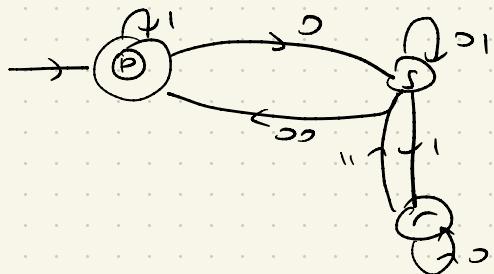
$$R_{32}^{(2)} = R_{32}^{(1)} + R_{31}^{(1)}(R_{21}^{(1)})^* R_{22}^{(1)} = 1 + 1 (\epsilon + 1^+ 0)^+ = 1 (1^+ 0)^+$$

$$R_{33}^{(2)} = R_{33}^{(1)} + R_{31}^{(1)}(R_{21}^{(1)})^* R_{23}^{(1)} = (0 + \epsilon) + 1 (\epsilon + 1^+ 0)^* 0 = 0 + 1 (1^+ 0)^* 0 \in$$

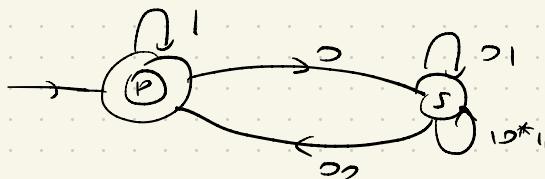
3.2.3



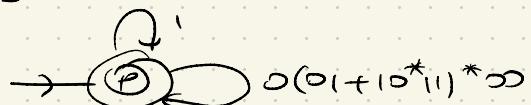
step 1: eliminate q :



step 2: eliminate r :



step 3: eliminate s :



Regular expression: $(1 + 0(01 + (0^*11)^*0)^*)^*$