

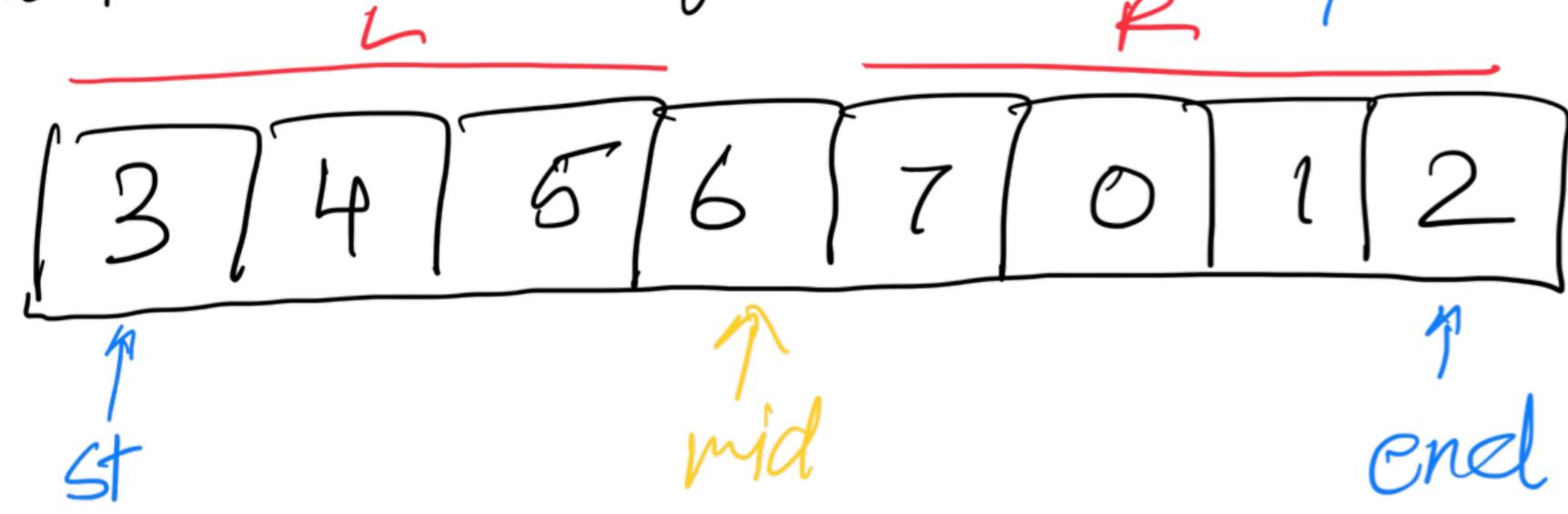
Search in rotated sorted array
ascending order; distinct values

[3 4 5 6 7 0 1 2] tar=0

Given array (rotated)

[0 1 2 3 4 5 6 7]

original array after being rotated
must be in $O(n \log n)$ TC \Rightarrow binary search



if you see one side of the array, one side WILL always be SORTED

abhi you know ki left side is sorted then use binary search apply bcz we can find

RS Array

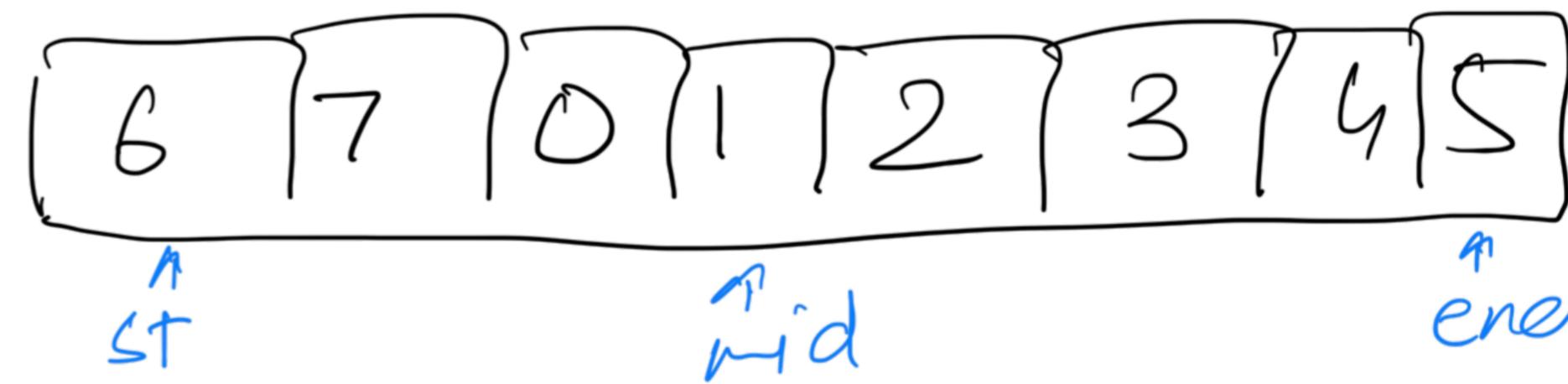
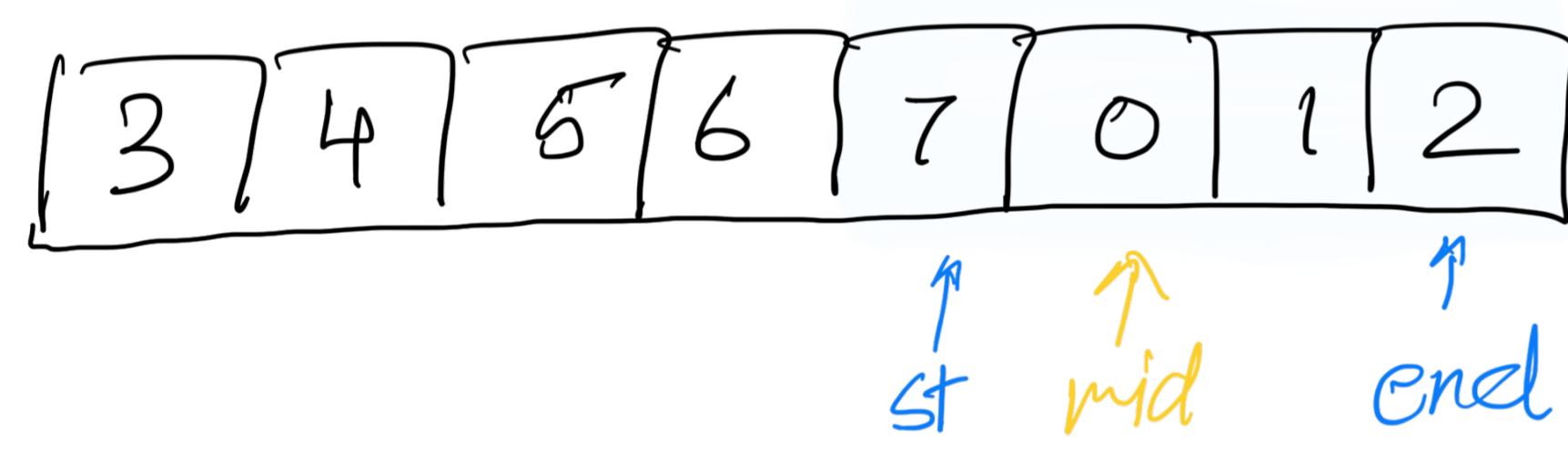
LS sorted RS sorted

to know if its LS then
 $A[st] \leq A[mid]$
if yes then LS

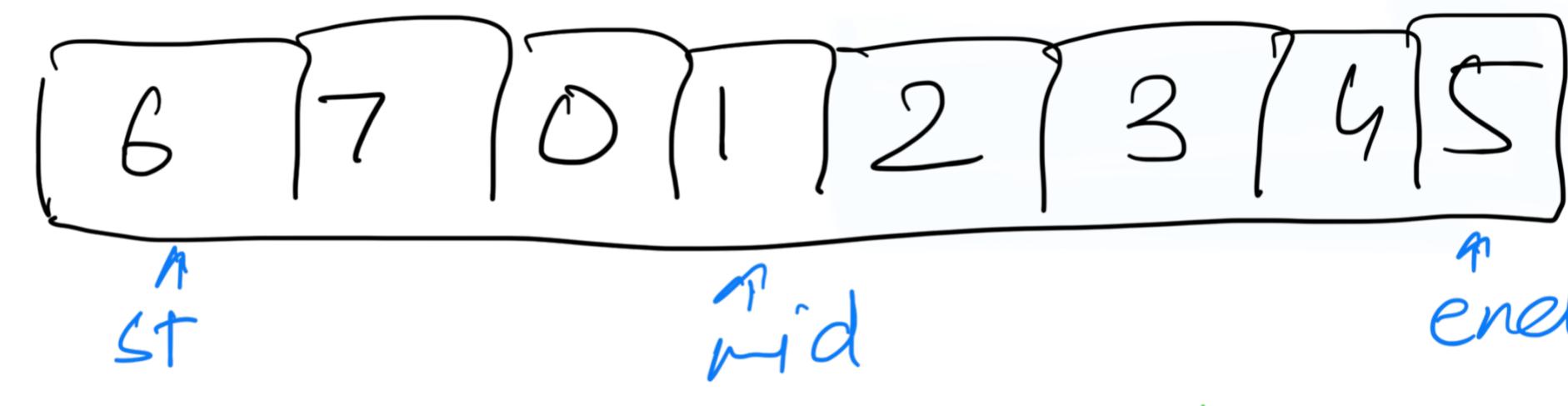
\hookrightarrow L-BS (left side search for tar)

$A[st] < tar < A[mid]$ left
else right

\therefore st, end
 \therefore (st, mid-1)
 \therefore (mid+1, end)



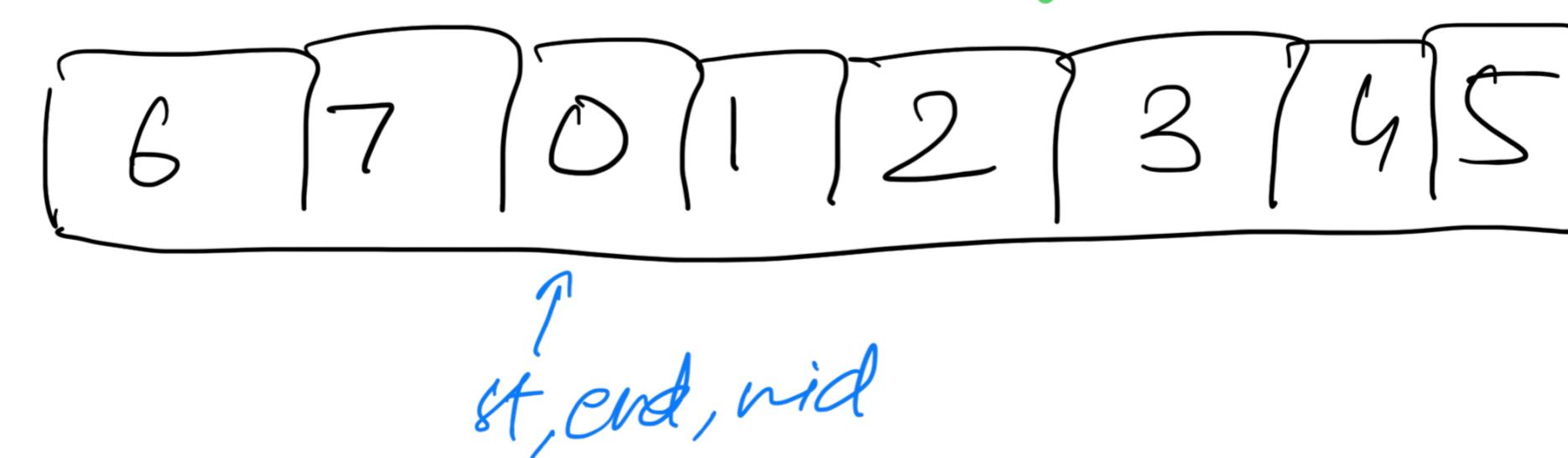
R-BS (right side search for tar)
 $A[mid] \leq tar \leq A[end]$ $\therefore (mid+1, end)$
else $\therefore (st, mid-1)$



Searching in right side but target doesn't exist so we go left side then we change the st & end pts



again we will find which side is sorted in this case left side sorted, we'll check the conditions again and move accordingly



Code

```
st = 0, end = n - 1;  
while (st <= end){  
    mid = st + (end - st) / 2;  
    if (A[mid] == tar)  
        return mid;  
    if (A[st] <= A[mid]) { // left sorted  
        if (A[st] <= tar <= A[mid])  $\Rightarrow$  left  $\Rightarrow$  end = mid - 1  
        else  $\Rightarrow$  right  $\Rightarrow$  st = mid + 1  
    } else { // right sorted  
        if (A[mid] <= tar <= A[end])  $\Rightarrow$  right  $\Rightarrow$  st = mid + 1  
        else  $\Rightarrow$  left  $\Rightarrow$  end = mid - 1  
    }  
}
```

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