

# Intersection-free spatial embeddings of graphs

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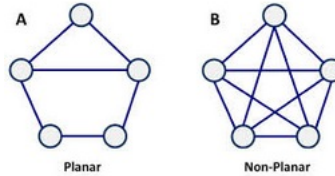
## 1 Introduction

In this paper, we will provide proof that any given graph containing edges crossing each others can be embedded in the space without said edges crossing any other.

## 2 Definitions

### 2.1 Planar/Non-planar graph:

A planar graph is a graph that can be embedded in the plane, i.e., it can be drawn on the plane in such a way that its edges intersect only at their endpoints. In other words, it can be drawn in such a way that no edges cross each other [1].



**Figure 1:** Planar & Non-planar Graphs

### 2.2 Crossing Number:

The crossing number of a graph  $G = (V, E)$  is the minimum number of crossings between edges of  $G$  among all the ways to draw  $G$  in the plane. It is denoted  $cr(G)$ .

The edges in a drawing of  $G$  need not be line segments, they are allowed to be arbitrary continuous curves. If one restricts to the straight-line drawings, then one obtains the rectilinear crossing number  $lin-cr(G)$ .

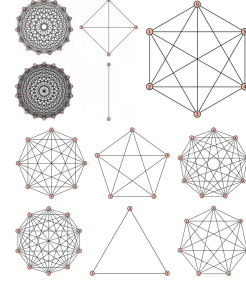
The space crossing number of  $G$ , denoted  $cr_4(G)$ , is the least number of crossings in any spatial drawing of  $G$ .

As in the planar case, the spatial rectilinear crossing number  $lin - cr_4(G)$  is obtained by restricting to straight-line spatial drawings. [2]

### 2.3 Complete Graph:

A complete graph is a simple undirected graph in which every pair of distinct vertices is connected by a unique edge. [3]

$$|V| = n \quad |E| = \frac{n(n+1)}{2}$$



**Figure 2:** A few complete graphs

## 3 Proposition

Any given graph with  $lin - cr(G) \geq 0$  crossings can be embedded in the euclidean space with no crossing edges,  
i.e.  $lin - cr_4(G) = 0$ .

## 4 Proof using Planes

A few lemmas needed for our proof.



**Lemma 4.1** *3 points in general position in space defines a plane.*

**Lemma 4.2** *The volume  $V$  of any given plane (2 dimensional space) is  $V = 0$ .*

**Lemma 4.3** *The number of 2 dimensional subspaces in 3 dimensional space is infinite.*

### 4.1 Proof by induction:

We will prove by induction on  $n$  that for any given graph  $G(V, E)$  embedded in space such as  $|V(G)| = n$  and  $lin - cr_4(G) = 0$ , we can add another vertex  $e'$ , thus creating  $G'$  such as  $lin - cr_4(G') = lin - cr_4(G) = 0$ .

### Base Case:

1. Let's assume the complete  $H(V, E)$  such as  $V(H) = 3$  embedded in space (it is trivial that  $lin - cr_4(H) = 0$ ).
2. Using **Lemma 1**, we can notice that  $H$  is also embedded into a single plane  $P$ .



3. Let's add a 4<sup>th</sup> vertex  $e$  so that  $e \notin P$ .

### Transition Step:

Let's assume the proposition is true for the complete spatial graph  $G(V, E)$  such as  $V(G) = n$  and  $lin - cr_4(G) = 0$ .

**Induction Step:** *Let's see that the proposition is also true for the complete spatial graph  $G'(V, E)$  such as  $V(G') = n + 1$*

1. Let's find a  $(n + 1)^{th}$  point  $e'$  in space. Using **Lemma 2** & **Lemma 3**,  $e'$  can be found on a plane not yet populated by 3 already existing vertices.
2. By connecting  $e'$  to all existing vertex of  $G$ , we obtain  $G'$ .

*Hence, we created a graph in which no 4 vertices are coplanar, thus avoiding any crossing edges in the resulting graph.*

## 5 Conclusion

By means of decomposition and recomposition, any given graph can be embedded in space without any crossing edges. ■

## References

- [1] Trudeau, Richard J. (1993) *Introduction to Graph Theory*. New York: Dover Pub. p.64. "Thus a planar graph, when drawn on a flat surface, either has no edge-crossings or can be redrawn without them."
- [2] Bukh, Boris & Hubard, Alfredo (2011) *Space Crossing Numbers*. arXiv:1102.1275v2. p.1.
- [3] Diestel, Reinhard (2010) *Graph Theory, Fourth Electronic Edition*. Springer. "If all the vertices of  $G$  are pairwise adjacent, then  $G$  is complete."