

# CS 112 – Introduction to Computing II

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Digraphs and Graphs:

Basic notions, representations, examples

Graph Search: Depth-first, breadth-first, best-first



# Graphs: Basic Definitions

Graphs are the most basic model of collections of information, generalizing trees to allow arbitrary links between nodes. There are two flavors, Directed and Undirected.

A Directed Graph (or Digraph) is:

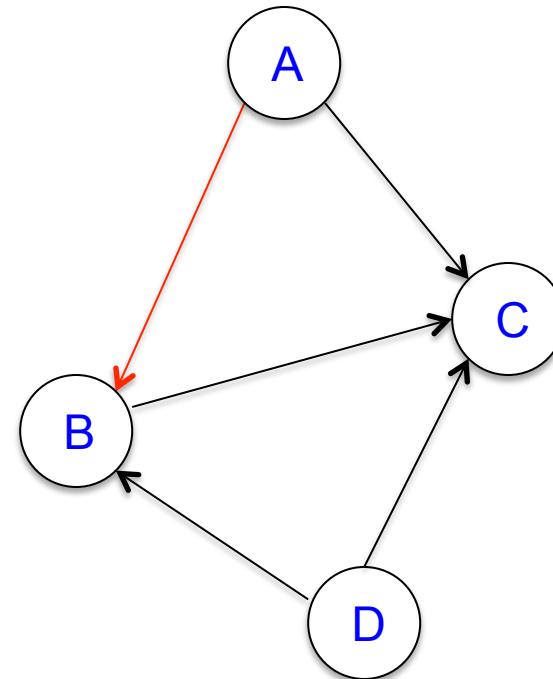
- A set  $V$  of Vertices (or: Nodes) containing (possibly):
  - A Label (1, 2, ... A, B, ... etc.)
  - Data fields (boolean flags, counters, etc. )
- A set  $E$  of Edges (links) connecting vertices; edges may have labels or data (e.g., weight or cost) associated with them.

$E$  is usually expressed as a relation on  $V$ , i.e.,  $E$  consists of pairs of vertices:

(source, target)

$E$  is a subset of  $V \times V$  (Cartesian Product of  $V$ )

Directed Graph  $G$



$$V = \{ A, B, C, D \}$$

$$E = \{ (A,B), (A,C), (B,C), (D,B), (D,C) \}$$

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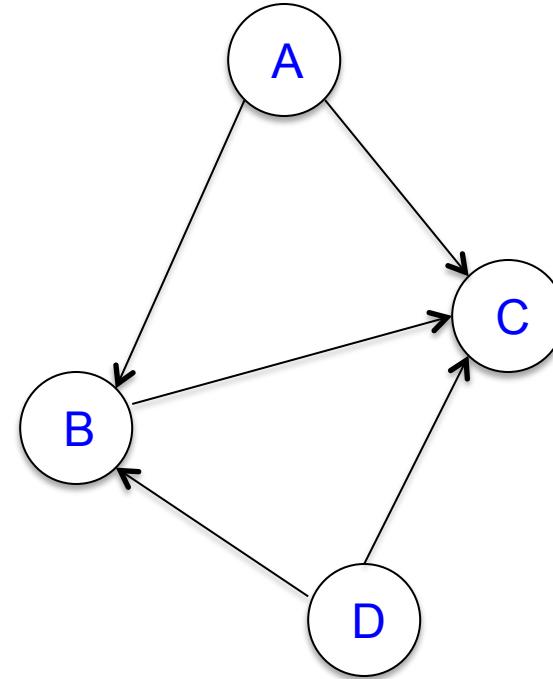
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$$V = \{ A, B, C, D \}$$

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Note: Like a tree, but

- Any vertex can be connected to any other vertex;
- There is no root.

# Graphs: Basic Definitions

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Edges have a direction (e.g., one way streets)

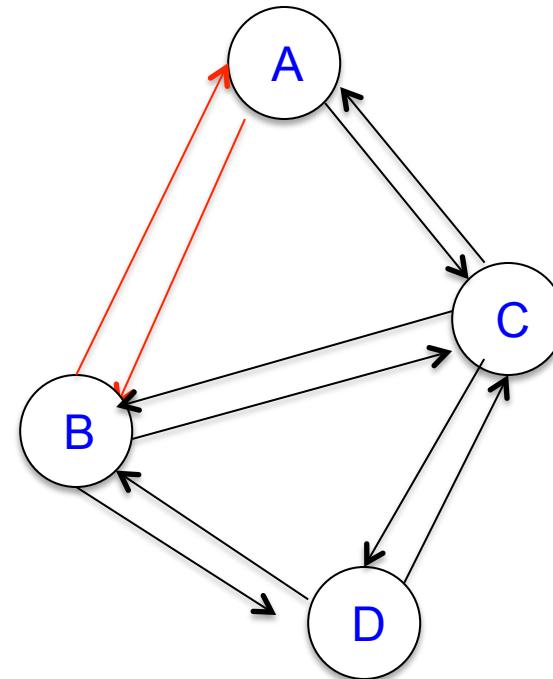
An Undirected Graph has the additional feature that all edges are two way, i.e., the relation  $E$  is symmetric:

$(A, B)$  is in  $E$  iff  $(B,A)$  is in  $E$

Edges do NOT have a direction (e.g., two-way streets).

“Graph” can mean either, but generally is assumed to be undirected.

Undirected Graph  $G$



$$V = \{ A, B, C, D \}$$

$$E = \{ (A,B), (A,C), (B,C), (D,B), (D,C), (B,A), (C,A), (C,B), (B,D), (C,D) \}$$

# Graphs: Basic Definitions

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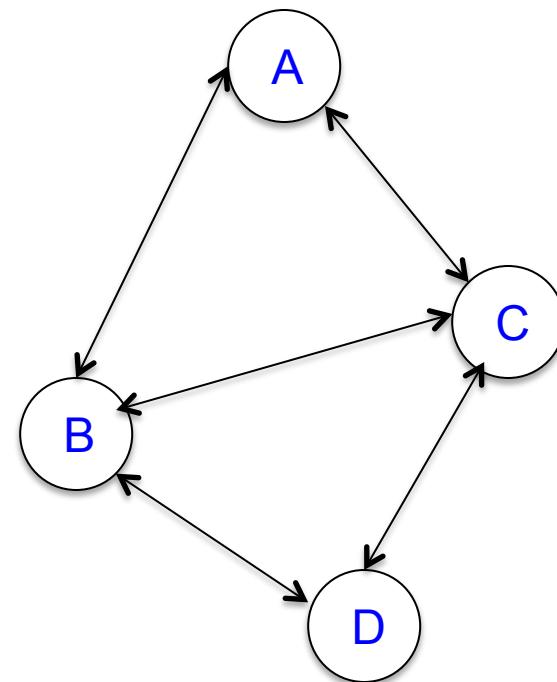
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## Undirected Graph $G$



Edges in an undirected graph may have two arrows

$$V = \{ A, B, C, D \}$$

$$E = \{ (A,B), (A,C), (B,C), (D,B), (D,C), (B,A), (C,A), (C,B), (B,D), (C,D) \}$$

# Graphs: Basic Definitions

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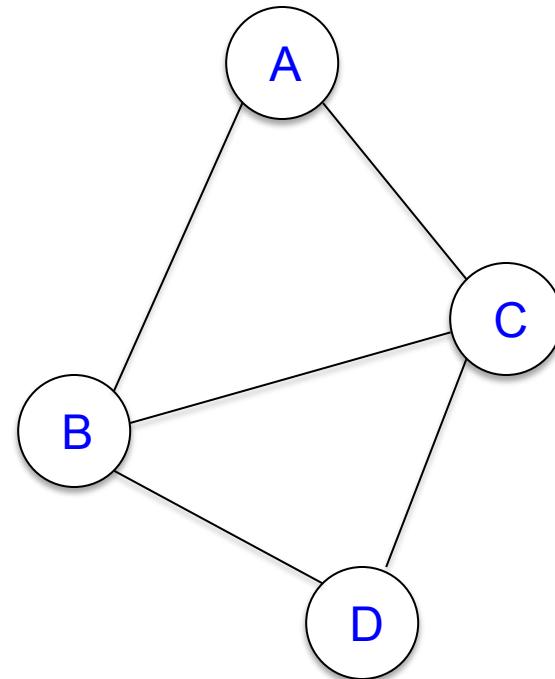
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Undirected Graph  $G$



But generally are drawn without arrow-heads.

$$V = \{ A, B, C, D \}$$

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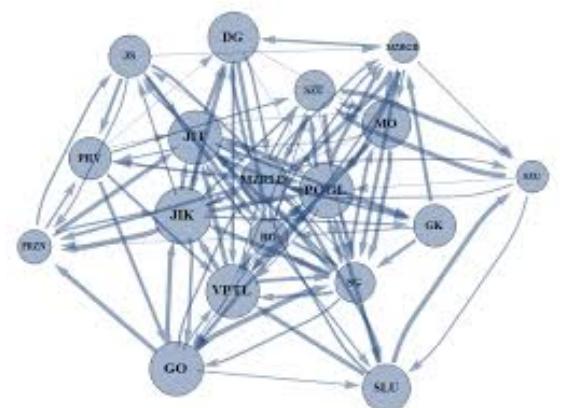
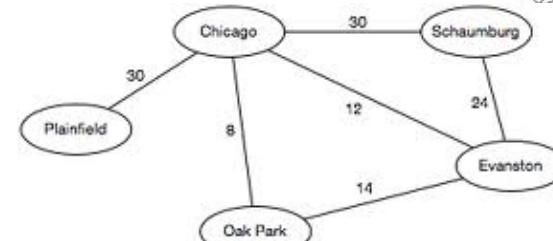
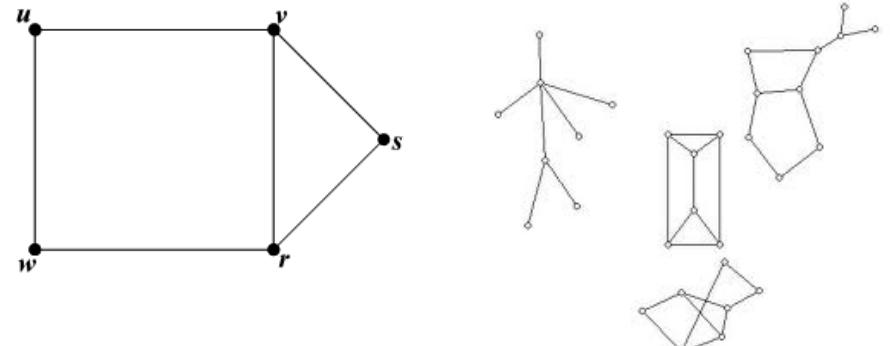
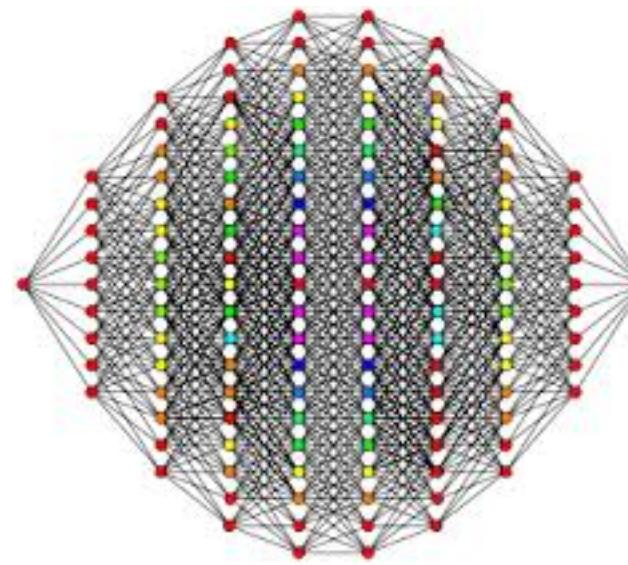
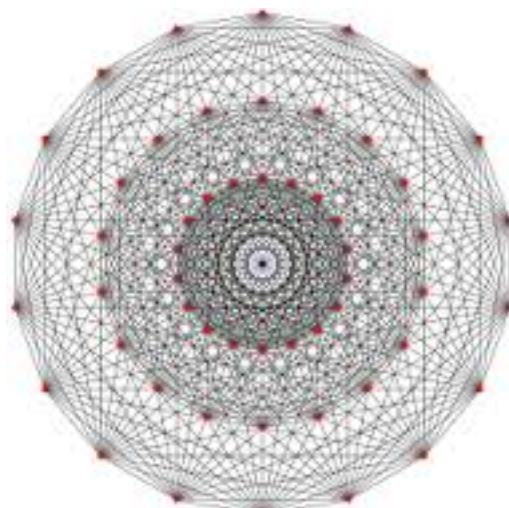
# Graphs: Basic Definitions



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We will focus on Undirected Graphs for this lecture.

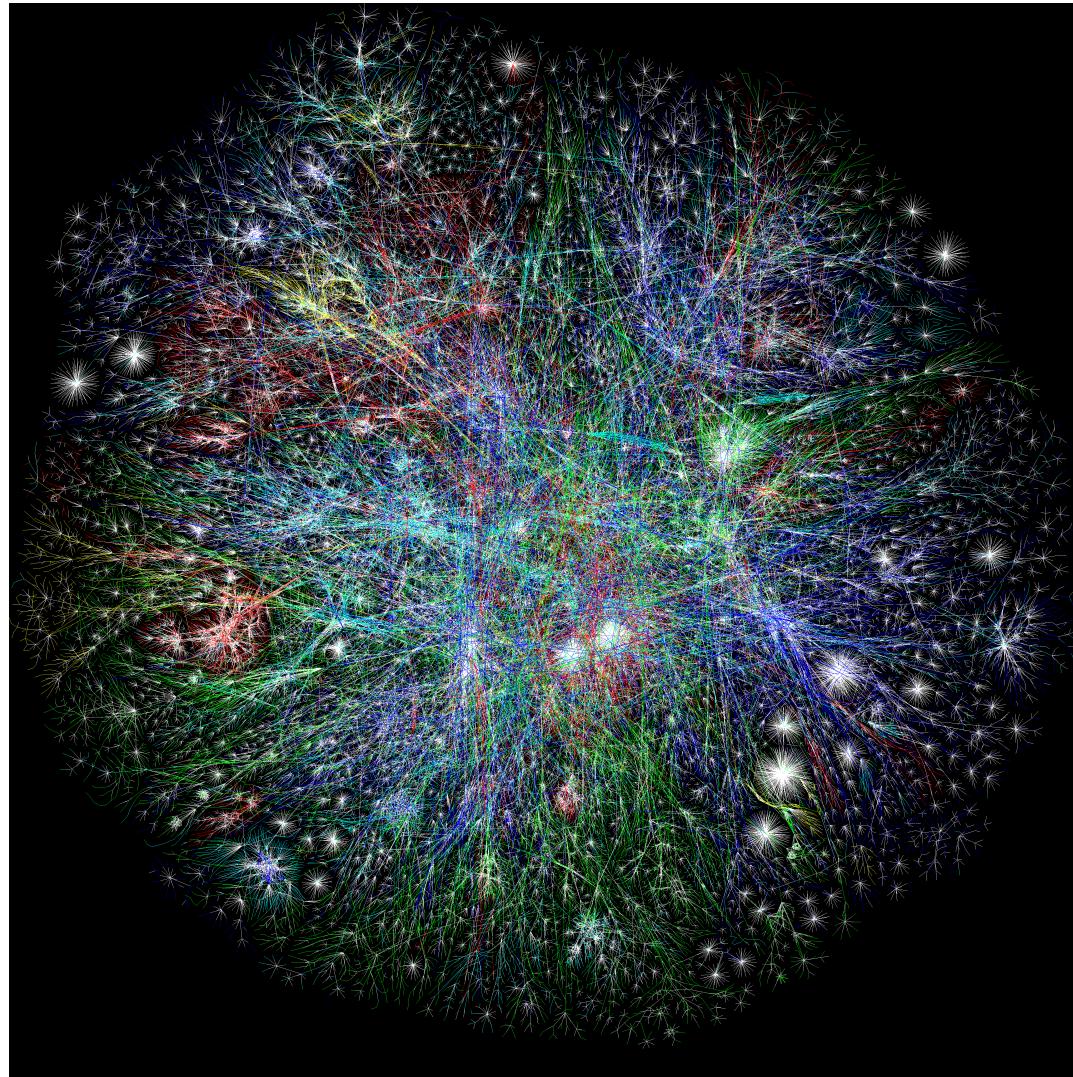
Graphs are EVERYWHERE in computer science and mathematics, and you have been looking at them for years.....



# Graphs: Basic Definitions



And you've been using graphs for years.....



# Graphs: Basic Definitions

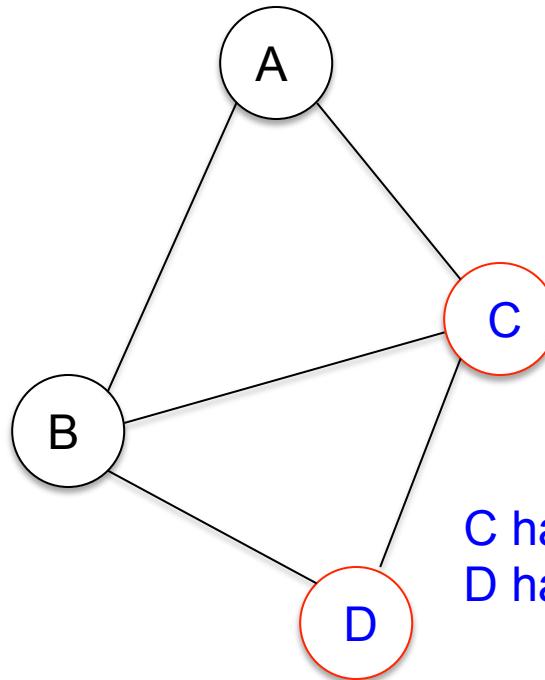
Basic Notions of Graphs:

Vertex (Vertex set V)

Edge (Edge set E)

The **degree** of a vertex is the number of edges it participates in .

Graph G



C has degree 3;  
D has degree 2

$$V = \{ A, B, C, D \}$$

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# Graphs: Basic Definitions

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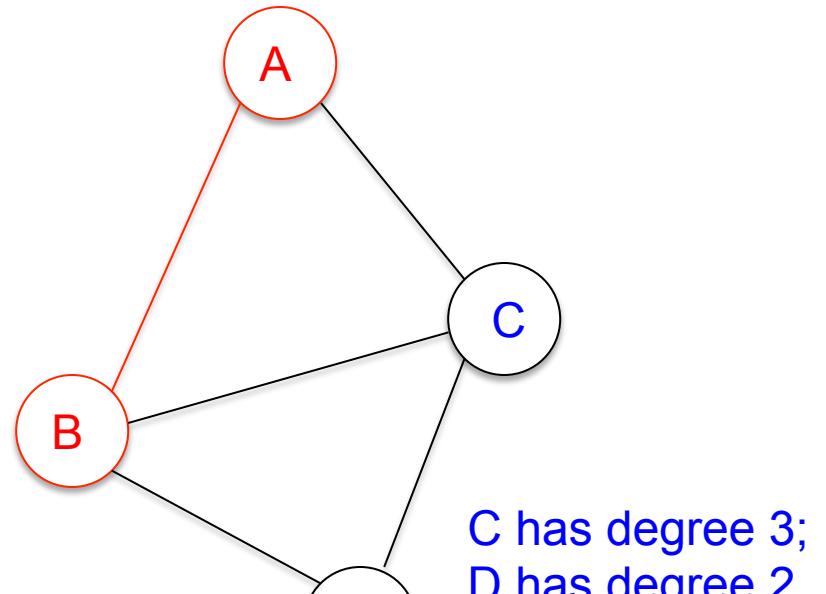
Vertex (Vertex set V)

Edge (Edge set E)

The degree of a vertex is the number of edges it participates in .

Two vertices are **adjacent** if there is an edge between them.

Graph G



$$V = \{ A, B, C, D \}$$

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# Graphs: Basic Definitions

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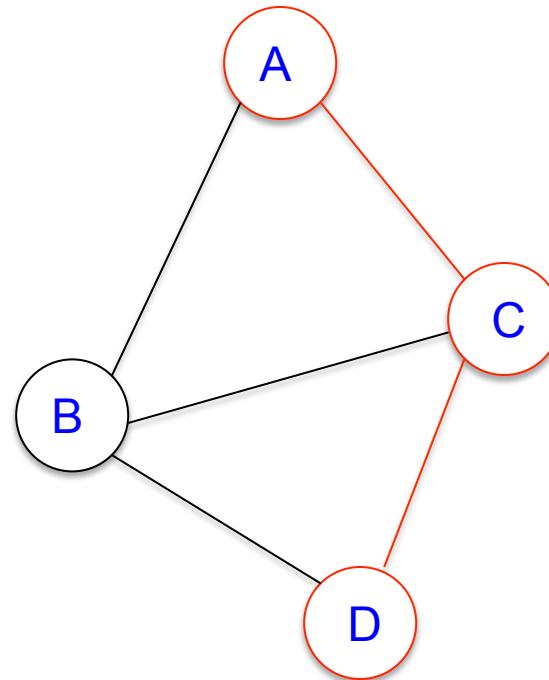
Edge (Edge set E)

The degree of a vertex is the number of edges it participates in .

Two vertices are adjacent if there is an edge between them.

A path is a sequence of adjacent edges.

Graph G



Path: A, C, D

$$V = \{ A, B, C, D \}$$

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# Graphs: Basic Definitions

Basic Notions of Graphs:

Vertex (Vertex set V)

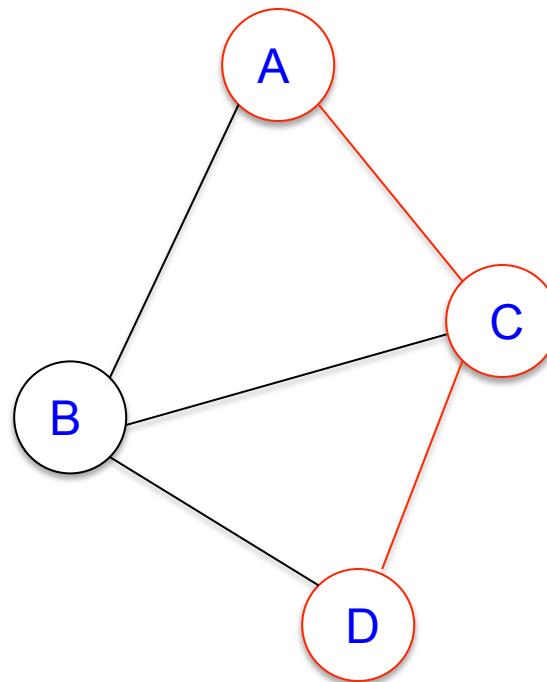
Edge (Edge set E)

The degree of a vertex is the number of edges it participates in .

Two vertices are adjacent if there is an edge between them.

A path is a sequence of adjacent edges. The length of a path is the number of edges.

Graph G



Path A, C, D has length 2.

$$V = \{ A, B, C, D \}$$

$$E = \{ (A,B), (A, C), (B,C), (D,B), (D,C), (B,A), (C,A), (C,B), (B,D), (C,D) \}$$

# Graphs: Basic Definitions

Basic Notions of Graphs:

Vertex (Vertex set V)

Edge (Edge set E)

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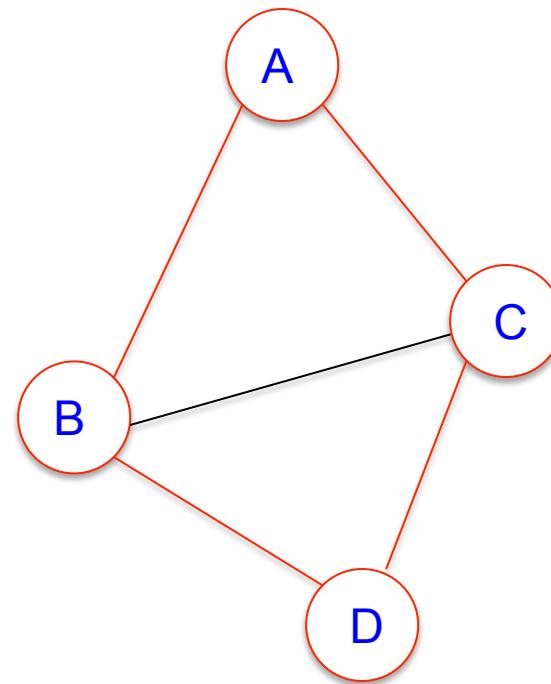
Two vertices are adjacent if there is an edge between them.

A path is a sequence of adjacent edges.

A **cycle** (or loop) is a path that begins and ends on the same vertex.

A, C, D, B, A

Graph G



$$V = \{ A, B, C, D \}$$

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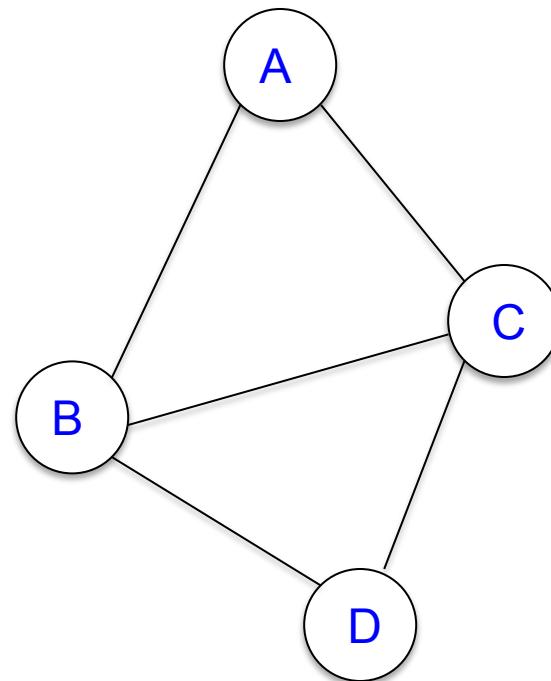
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A, C, D, B, A

A **simple cycle** has no other repeated vertices (i.e., does not cross itself).

A, C, B, A, C, D, B, A is not simple.

Graph G



$$V = \{ A, B, C, D \}$$

$$E = \{ (A,B), (A, C), (B,C), (D,B), (D,C), (B,A), (C,A), (C,B), (B,D), (C,D) \}$$

# Graphs: Basic Definitions



Computer Science

Basic Notions of Graphs:

Vertex (Vertex set V); Edge (Edge set E)

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A **simple cycle** has no other repeated vertices (i.e., does not cross itself).

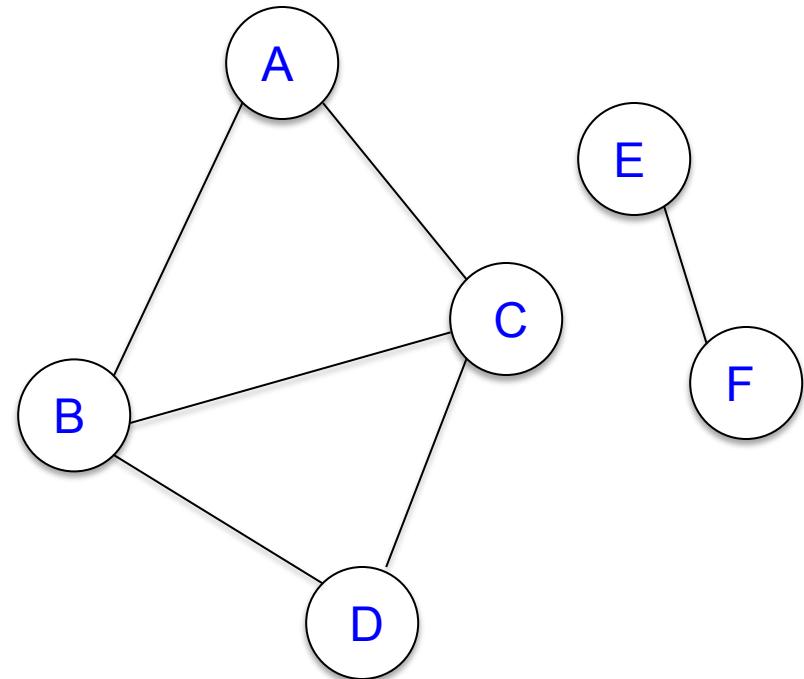
A, C, B, A, C, D, B, A is not simple.

Two vertices are connected if there is a path between them.

A **set of vertices is connected** if there is a path between each pair; a **graph** is connected if there is a path between any two nodes.

{A, B, C, D} is connected; so is {E, F}

Graph G



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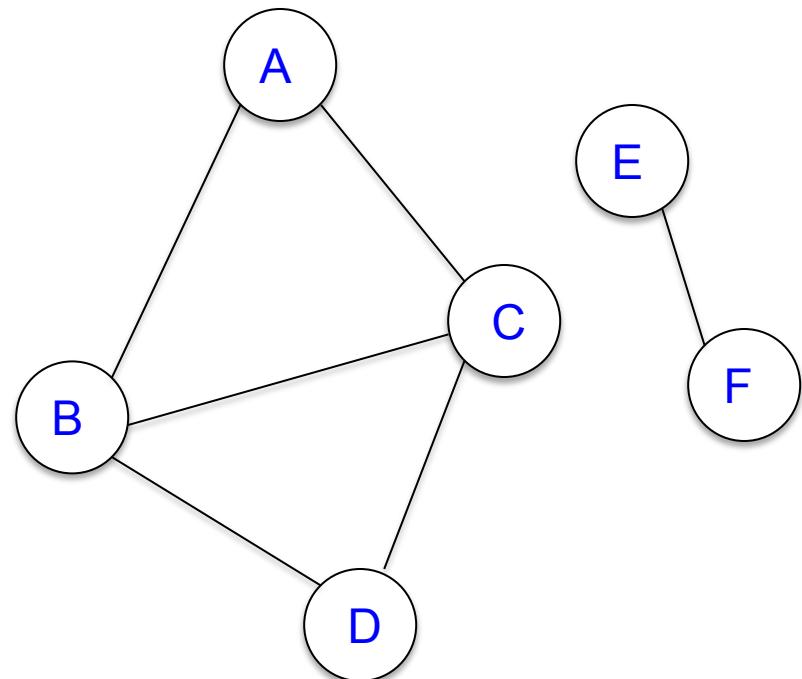
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The set of **connected components** of G is the collection of maximal connected subsets of vertices. The **connected components** of G are {A, B, C, D} and {E, F}

Graph G



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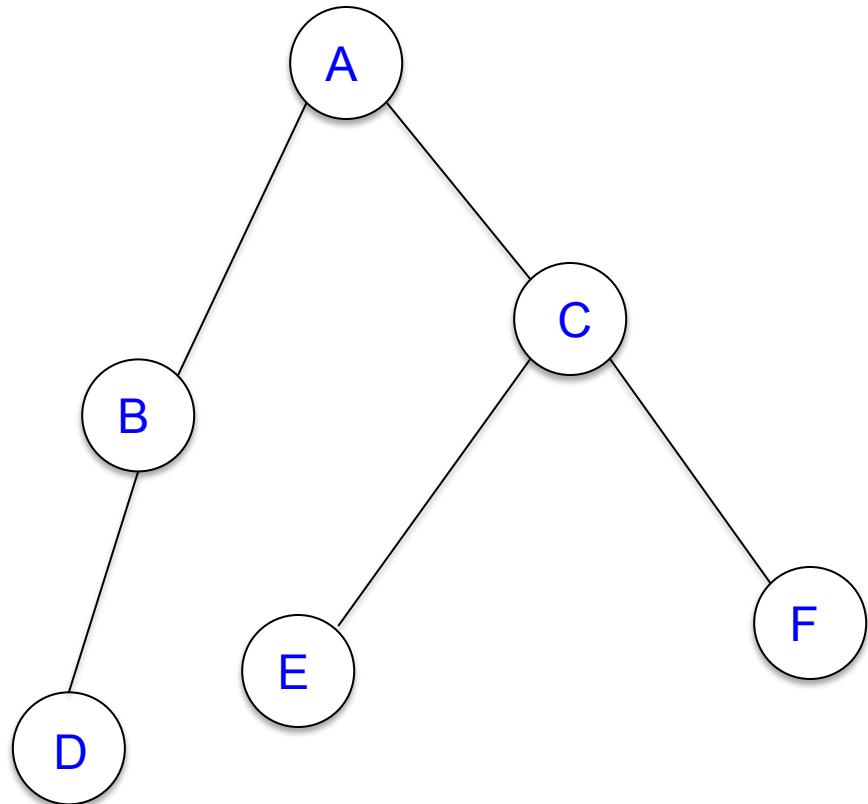
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A **tree** is a acyclic, connected graph.

Graph G



# Graphs: Basic Definitions



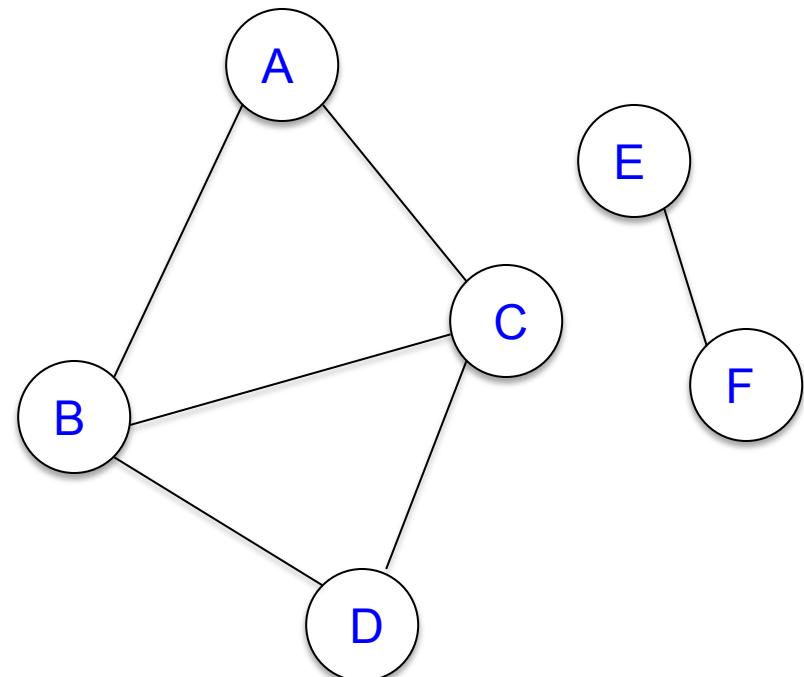
Implementing Graphs.

Vertex set V is a list of vertices:

$$V = (A, B, C, D, E, F)$$

Edge set is a boolean array:

	A	B	C	D	E	F
A		T	T			
B	T		T	T		
C	T	T		T		
D		T	T			
E						T
F				T		



# Graphs: Basic Definitions



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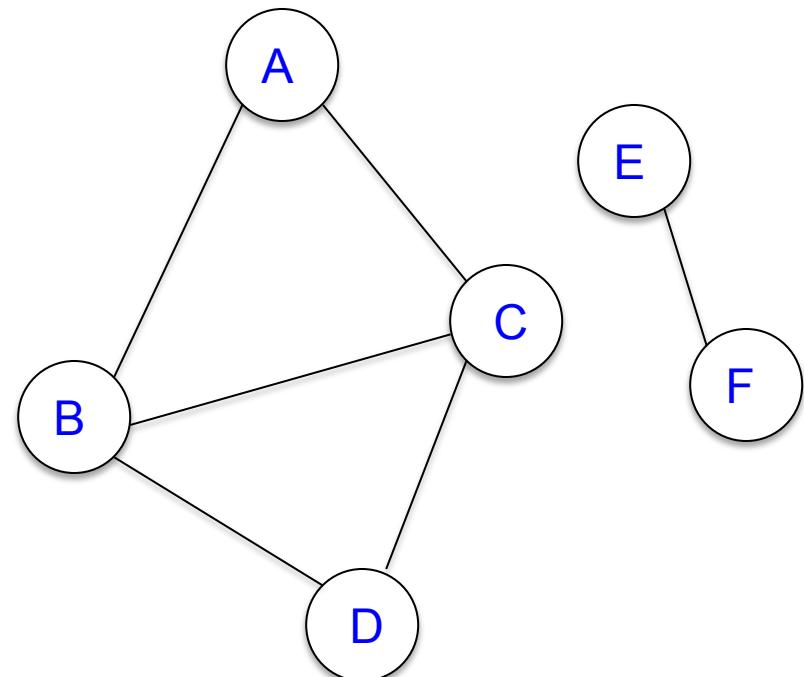
Implementing Graphs.

Vertex set V is a list of vertices:

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Edge set is a boolean array: Note the symmetry!

	A	B	C	D	E	F
A		T	T			
B	T		T	T		
C	T	T		T		
D		T	T			
E						T
F					T	



# Graphs: Basic Definitions

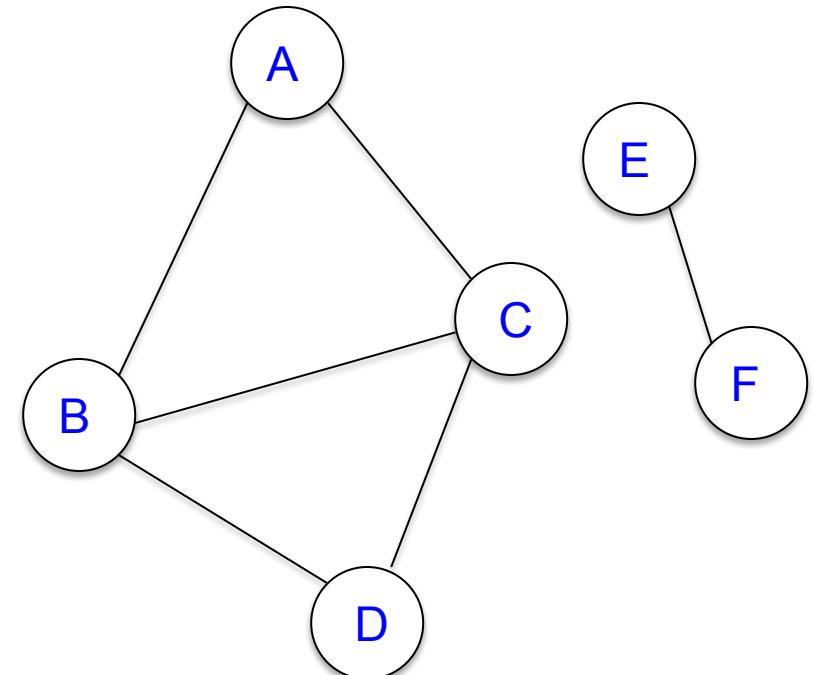
Implementing Graphs.

Vertex set  $V$  is a list of vertices:

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Edge set is a boolean array; you could optimize by only storing half!

	A	B	C	D	E	F
A		T	T			
B	T		T	T		
C	T	T		T		
D		T	T			
E					T	
F				T		



# Graphs: Basic Definitions



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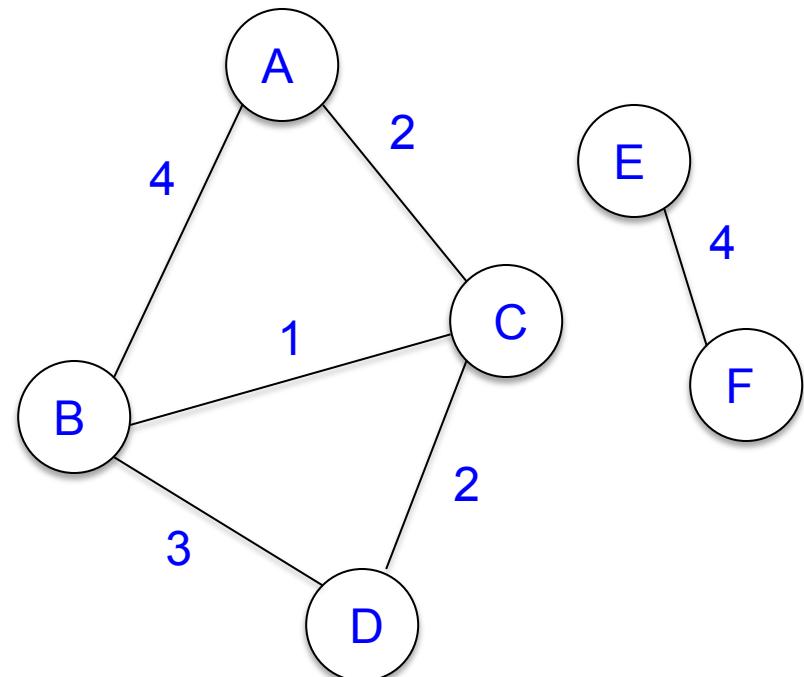
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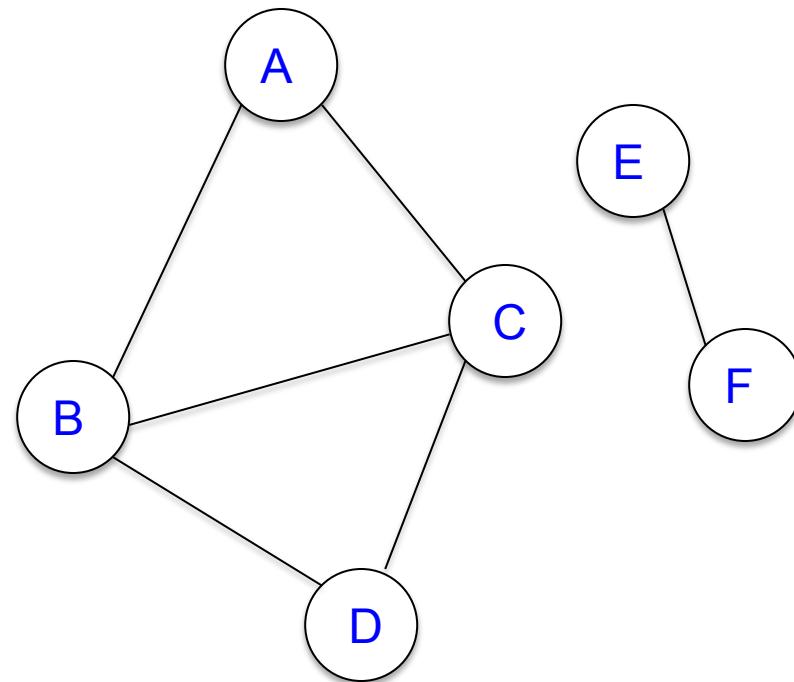
OR, edge set is an integer array giving the weights of the edges:

	A	B	C	D	E	F
A		4	2			
B	4		1	3		
C	2	1		2		
D		3	2			
E						4
F				4		



# Graphs: Basic Definitions

	A	B	C	D	E	F
A		T	T			
B	T		T	T		
C	T	T		T		
D		T	T			
E						T
F					T	



You need a number of methods:

**Adjacent(v)** = list of vertices adjacent to v, in some order (if scan across the row for v would be in order of set V, but could be otherwise)

**Degree(v)** = number of T values in row for v

Etc. (more as we develop the algorithms)

# Graphs: Basic Definitions

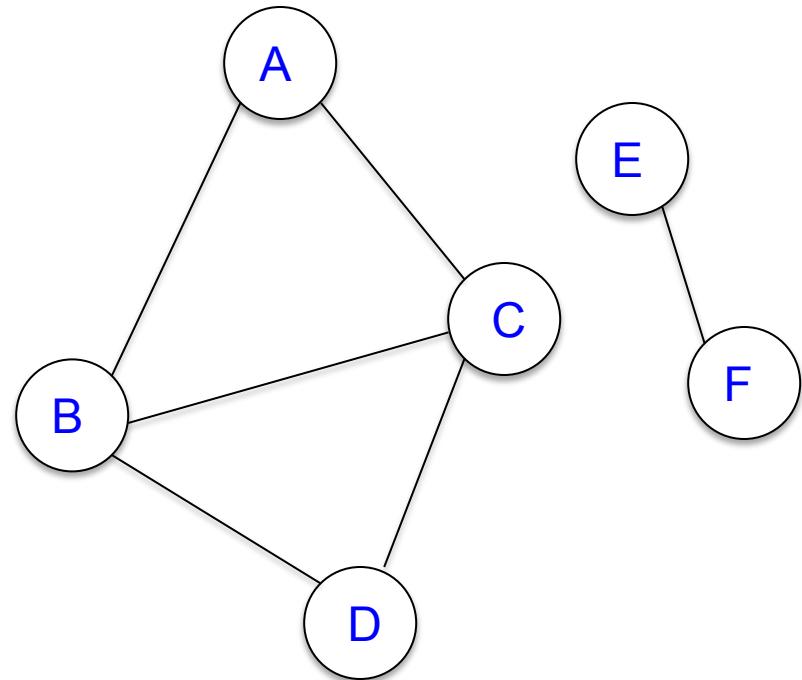
Graph Search Algorithms are all generalizations of the following algorithm for Depth-First Search

Suppose that vertices have a flag:

boolean visited;

```
searchGraph(V, E) {  
    foreach( v in V )           // initialize all vertices  
        v.visited = false;  
    foreach( v in V )  
        if(!v.visited)  
            DFS(v);  
}  
  
DFS( v ) {  
    if(!v.visited) {  
        visit(v);             // do something, e.g. print out  
        v.visited = true;  
        foreach( u in Adjacent(v) ) {  
            DFS( u );  
        }  
    }  
}
```

Graph G



# Graphs: Basic Definitions



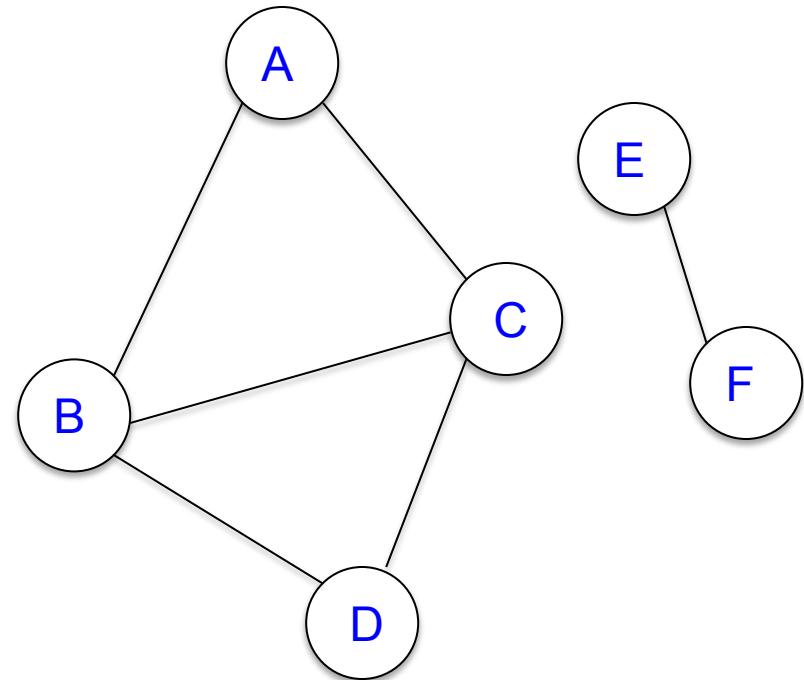
Computer Science

To consider the general search algorithms, we have to rephrase this as a non-recursive algorithm.

Here is depth-first search again:

```
searchGraph(V, E) {  
    foreach( v in V )           // initialize all vertices  
        v.visited = false;  
    foreach( v in V )  
        if(!v.visited)  
            DFS(v);  
}  
  
DFS( v ) {  
    Stack S = new Stack();  
    S.push(v);  
    while( ! S.isempty() ) {  
        Vertex u = S.pop();  
        if(!u.visited) {  
            visit(u);  
            u.visited = true;  
            foreach( w in Adjacent(u) )  
                if(!w.visited)  
                    S.push(w);  
        }  
    }  
}
```

Graph G



# Graphs: Basic Definitions

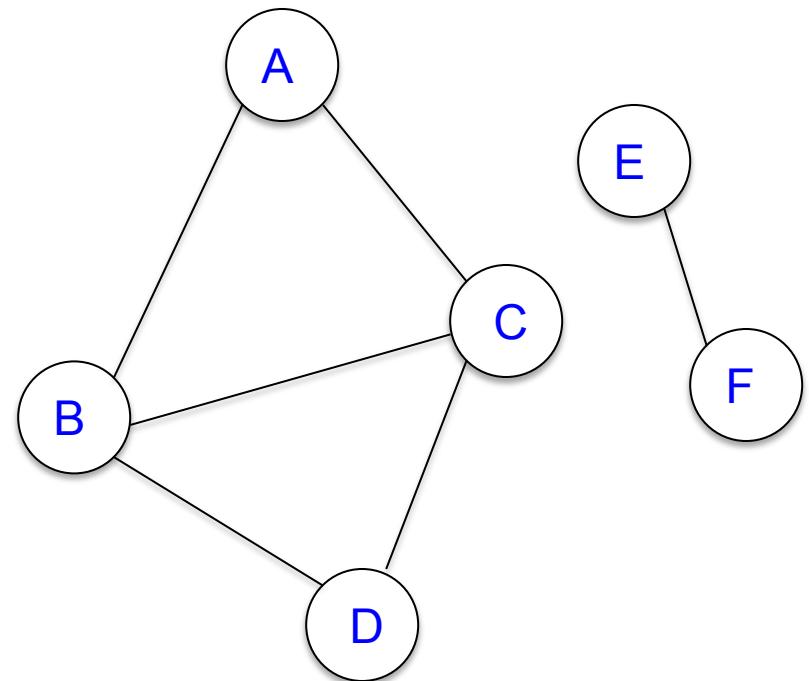
To consider the general search algorithms, we have to rephrase this as a non-recursive algorithm.

By replacing the Stack with a Queue, we have Breadth-First Search

```
boolean visited;
searchGraph(V, E) {
    foreach( v in V )           // initialize all vertices
        v.visited = false;
    foreach( v in V )
        if(!v.visited)
            BFS(v);
}
```

```
BFS( v ) {
    Queue S = new Queue();
    S.enqueue(v);
    while( ! S.isEmpty() ) {
        Vertex u = S.dequeue();
        if(!u.visited) {
            visit(u);
            u.visited = true;
            foreach( w in Adjacent(u) )
                if(!w.visited)
                    S.enqueue(w);
    }
}
```

Graph G



# Graphs: Basic Definitions



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To complete  
rephe

Optional Slide: You are not responsible for best-first search.

By re  
Best

bool  
sear  
fo  
fo  
fo  
}

BFS  
PC  
S.  
wh

```
vertex u = S.dequeue(),
if(!u.visited) {
    visit(u);
    u.visited = true;
    foreach( w in Adjacent(u) )
        if(!w.visited)
            S.enqueue(w);
}
```

F

A diagram showing a single circular node labeled 'F'. A single directed edge originates from the bottom of the node and points upwards towards another node that is partially visible at the top right of the slide area.