华东师范大学期末试卷 (A)

2007—2008 学年第二学期 软件工程数学参考答案

(1) $(\exists x)A(x)$

前提引入

 $(2) \quad (\exists x) A(x) \to (\forall x) ((B(x) \lor A(x)) \to C(x))$

前提引入

 $(3) \quad (\forall x)((B(x) \lor A(x)) \to C(x))$

(1)(2)假言推理

(4) A(c)

(1)EI

 $(5) (B(c) \lor A(c)) \to C(c)$

(3)UI

(6) $A(c) \rightarrow (B(c) \lor A(c))$

附加

(7) $B(c) \vee A(c)$

(4)(6)假言推理

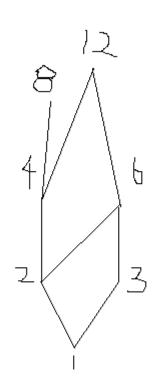
(8) C(c)

(7)(5) 假言推理

(9) $(\exists x)C(x)$

(8)EG

(1)



(2) 极大元: 8, 12 极小元: 1 最大元: 无 最小元: 无

 \equiv .

 $1.(a).5 \times 10 \times 10 = 500.$

- (b). $10 \times 5 \times 5 = 250$.
- (c). $500+250-5\times5\times5=625$.

2

$$P_5^5 = 120$$

四.

Write each of the n+1 integers $a_1, a_2, ..., a_{n+1}$ as a power of 2 times an odd integer. In other words, let $a_j = 2^{k_j} q_j$ for j = 1, 2, ..., n+1, where k_j is a nonnegative integer and q_j is odd. The integers $q_1, q_2, ..., q_{n+1}$ are all odd positive integers less than 2n. Since there are only n odd positive integers less than 2n, it follows from the pigeonhole priciple that two of the integers $q_1, q_2, ..., q_{n+1}$ must be equal. Therefore, there are integers i and j such that $q_i = q_j$. Let q be the common value of q_i and q_j . Then, $a_i = 2^{k_i} q$ and $a_i = 2^{k_i} q$ and $a_j = 2^{k_j} q$. It follows that if $k_i < k_j$, then a_i divides a_j ; while if $k_i > k_j$, then a_j divides a_i .

五.

$$\mathbb{R} + \mathbb{R} (x+y)^n = \sum_{k=0}^n C(n,k) x^{n-k} y^k$$

整理即可得到:
$$C(n,0) + C(n,2) + C(n,4) + \dots + C(n,n)$$

= $C(n,1) + C(n,3) + C(n,5) + \dots + C(n,n-1)$

六.

$$(a)C(8,3)=56$$

七.

$$(1) a_n = 8a_{n-1} + 10^{n-1}$$

(2)
$$a_1 = 9$$

八

1.
$$-\frac{1}{4} - \frac{1}{12} \times (-3)^n + \frac{1}{3} \times 3^n$$

2.a)
$$p_3 n^3 + p_2 n^2 + p_1 n + p_0$$

b)
$$n^2(p_2n^2 + p_1n + p_0)(-2)^n$$

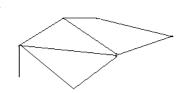
c)
$$n^2(p_4n^4 + p_3n^3 + p_2n^2 + p_1n + p_0)2^n$$

九.

$$(x+x^2+x^3+x^4)(1+x+x^2+x^3)(1+x+x^2+x^3+x^4+x^5+x^6)$$

+

1.



2.不同构,右图有一个4度顶点,而左图没有4度顶点

 $3.\{u1,u3,u5,u6\}, \{u2,u4\}$

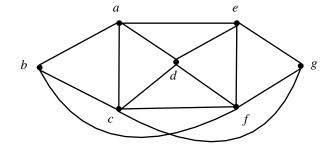
4. <u1,u2>,<u2,u4>,<u4,u7>,<u4,u5>,<u5,u8>,<u8,u9>

+--

- 1. n 是偶数, chromatic number 3 n 是奇数, chromatic number 4
- 2. 因为该连通图的每个顶点的度数都是偶数,所以一定存在欧拉回路可以从任意顶点出发,遍历每条边一次。 参考答案:

$$\begin{split} &\langle a,b \rangle \to \langle b,c \rangle \to \langle c,d \rangle \to \langle d,b \rangle \to \langle b,g \rangle \to \langle g,h \rangle \to \langle h,c \rangle \to \langle c,j \rangle \to \langle j,e \rangle \\ &\to \langle e,d \rangle \to \langle d,i \rangle \to \langle i,j \rangle \to \langle j,o \rangle \to \langle o,n \rangle \to \langle n,i \rangle \to \langle i,h \rangle \to \langle h,m \rangle \to \langle m,n \rangle \\ &\to \langle n,l \rangle \to \langle l,m \rangle \to \langle m,f \rangle \to \langle f,k \rangle \to \langle k,l \rangle \to \langle l,g \rangle \to \langle g,f \rangle \to \langle f,a \rangle \end{split}$$

+=



将边<a,b>和边<b,f>合并为<a,f>将边<e,g>和边<g,c>合并为<e,c>。 则 a,c,f,e,d 构成 K_5 。故该图是非平面图。