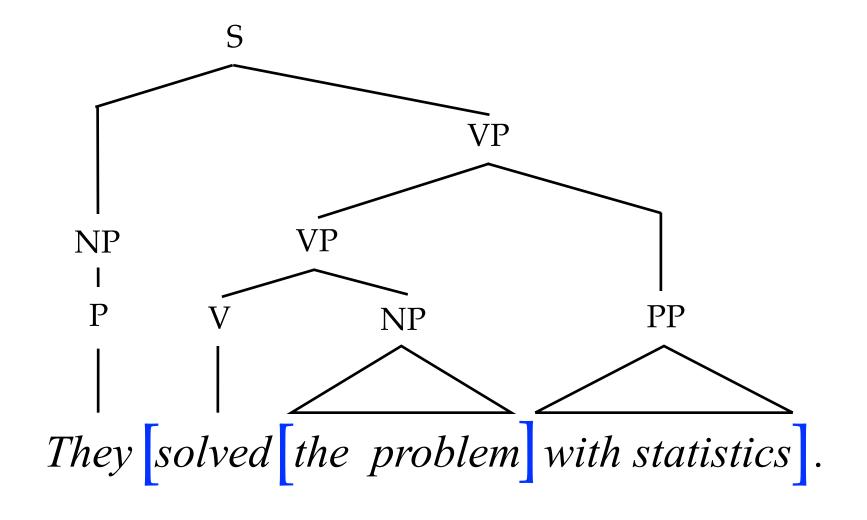
# Syntax and Parsing

#### Slav Petrov – Google Research

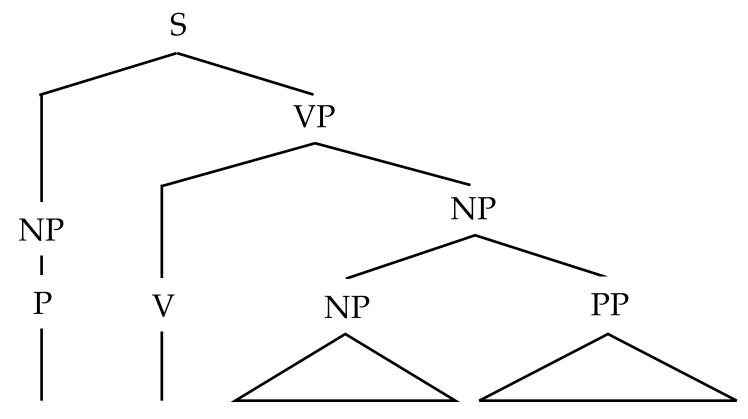
Thanks to: Dan Klein, Ryan McDonald

**Lisbon Machine Learning School** 

### Analyzing Natural Language

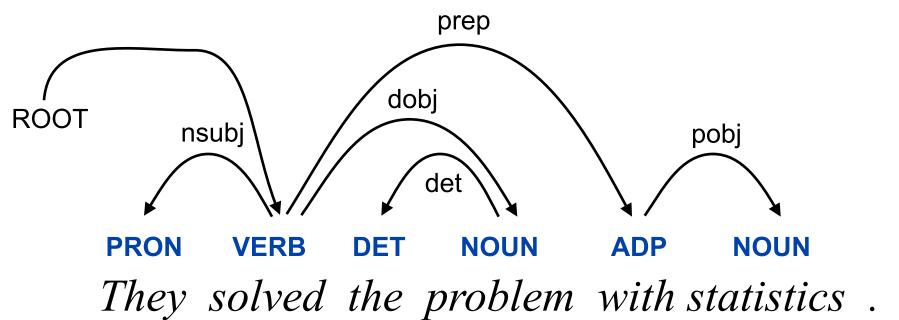


### Syntax and Semantics

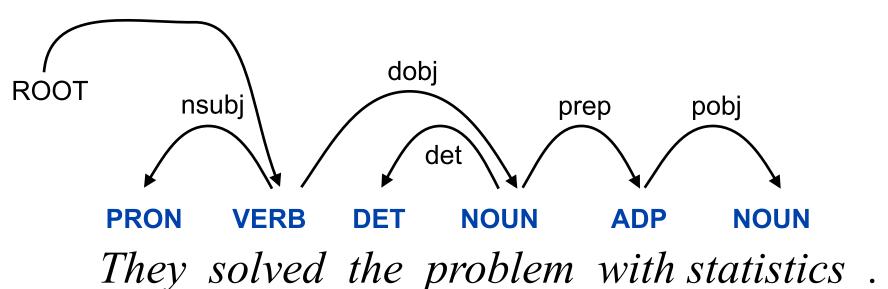


They solved the problem with statistics.

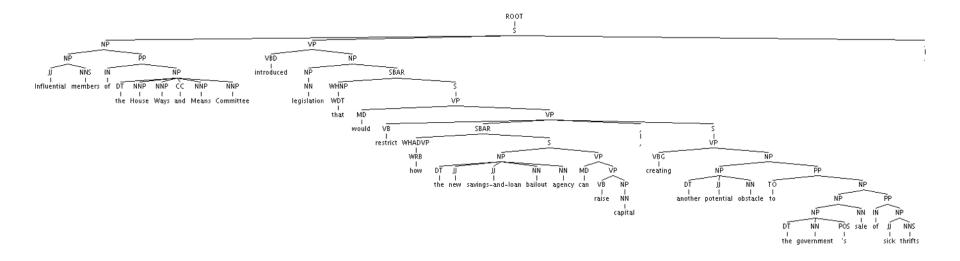
### Constituency and Dependency



### Constituency and Dependency



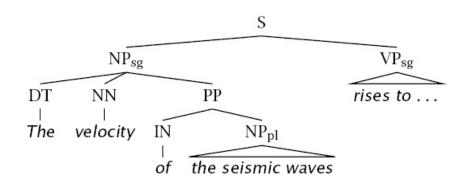
### A "real" Sentence



Influential members of the House Ways and Means Committee introduced legislation that would restrict how the new savings-and-loan bailout agency can raise capital, creating another potential obstacle to the government's sale of sick thrifts.

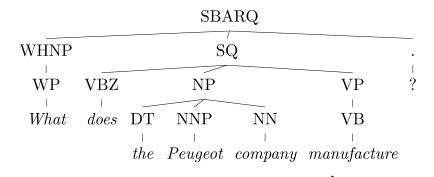
### Phrase Structure Parsing

- Phrase structure parsing organizes syntax into constituents or brackets
- In general, this involves nested trees
- Linguists can, and do, argue about details
- Lots of ambiguity
- Not the only kind of syntax...
- First part of today's lecture

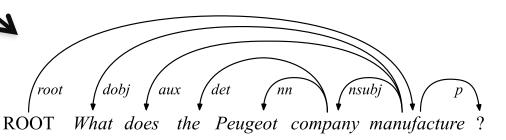


new art critics write reviews with computers

### Dependency Parsing



- Directed edges between pairs of word (head, dependent)
- Can handle free word-order languages
- Very efficient decoding algorithms exist
- Second part of today's lecture



### Classical NLP: Parsing

Write symbolic or logical rules:

```
VBD VB
VBN VBZ VBP VBZ
NNP NNS NN NNS CD NN
Fed raises interest rates 0.5 percent
```

- Use deduction systems to prove parses from words
  - Minimal grammar on "Fed raises" sentence: 36 parses
  - Real-size grammar: many millions of parses
- This scaled very badly, didn't yield broad-coverage tools

#### **Attachments**

I cleaned the dishes from dinner

I cleaned the dishes with detergent

I cleaned the dishes in my pajamas

I cleaned the dishes in the sink

#### Probabilistic Context-Free Grammars

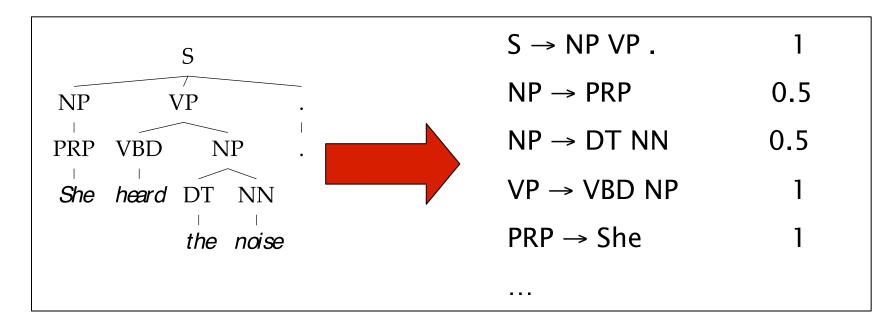
- A context-free grammar is a tuple <N, T, S, R>
  - N: the set of non-terminals
    - Phrasal categories: S, NP, VP, ADJP, etc.
    - Parts-of-speech (pre-terminals): NN, JJ, DT, VB
  - T: the set of terminals (the words)
  - S: the start symbol
    - Often written as ROOT or TOP
    - Not usually the sentence non-terminal S
  - R: the set of rules
    - Of the form  $X \rightarrow Y_1 Y_2 \dots Y_k$ , with  $X, Y_i \in N$
    - Examples: S → NP VP, VP → VP CC VP
    - Also called rewrites, productions, or local trees

#### A PCFG adds:

A top-down production probability per rule P(Y<sub>1</sub> Y<sub>2</sub> ... Y<sub>k</sub> | X)

#### Treebank Grammars

- Need a PCFG for broad coverage parsing.
- Can take a grammar right off the trees (doesn't work well):



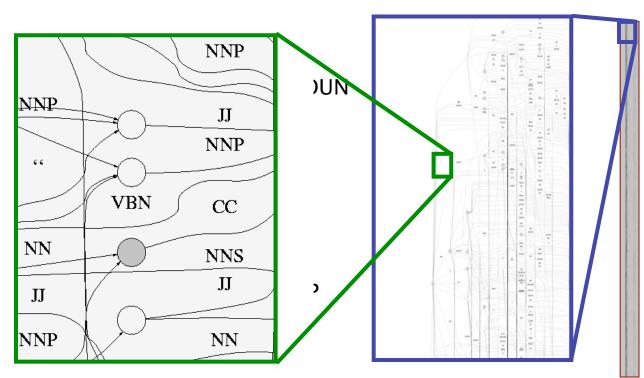
- Better results by enriching the grammar (e.g., lexicalization).
- Can also get reasonable parsers without lexicalization.

### Treebank Grammar Scale

Treebank grammars can be enormous

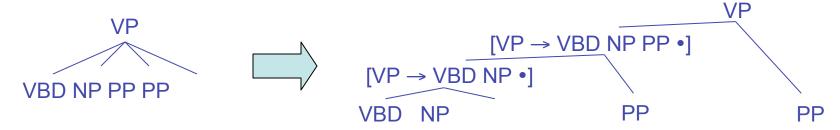
NP

- As FSAs, the raw grammar has ~10K states, excluding the lexicon
- Better parsers usually make the grammars larger, not smaller



### **Chomsky Normal Form**

- Chomsky normal form:
  - All rules of the form  $X \rightarrow Y Z$  or  $X \rightarrow w$
  - In principle, this is no limitation on the space of (P)CFGs
    - N-ary rules introduce new non-terminals



- Unaries / empties are "promoted"
- In practice it's kind of a pain:
  - Reconstructing n-aries is easy
  - Reconstructing unaries is trickier
  - The straightforward transformations don't preserve tree scores
- Makes parsing algorithms simpler!

#### A Recursive Parser

- Will this parser work?
- Why or why not?
- Memory requirements?

#### A Memoized Parser

One small change:

```
bestScore(X,i,j,s)
  if (scores[X][i][j] == null)
   if (j = i+1)
      score = tagScore(X,s[i])
  else
      score = max     score(X->YZ) *
            bestScore(Y,i,k) *
            bestScore(Z,k,j)
      scores[X][i][j] = score
  return scores[X][i][j]
```

# A Bottom-Up Parser (CKY)

Can also organize things bottom-up

```
bestScore(s)
 for (i : [0,n-1])
      for (X : tags[s[i]])
      score[X][i][i+1] =
           tagScore(X,s[i])
 for (diff : [2,n])
                                      i k
    for (i : [0,n-diff])
      j = i + diff
      for (X->YZ : rule)
         for (k : [i+1, j-1])
              score[X][i][j] = max score[X][i][j],
                             score(X->YZ) *
                             score[Y][i][k] *
                             score[Z][k][j]
```

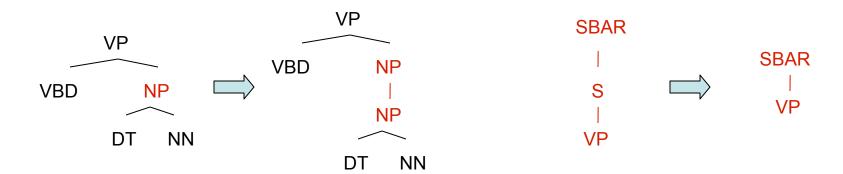
### **Unary Rules**

#### • Unary rules?

```
bestScore(X,i,j,s)
    if (j = i+1)
        return tagScore(X,s[i])
    else
        return max max score(X->YZ) *
        bestScore(Y,i,k) *
        bestScore(Z,k,j)
        max score(X->Y) *
        bestScore(Y,i,j)
```

# CNF + Unary Closure

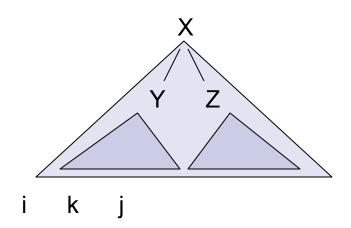
- We need unaries to be non-cyclic
  - Can address by pre-calculating the unary closure
  - Rather than having zero or more unaries, always have exactly one



- Alternate unary and binary layers
- Reconstruct unary chains afterwards

### Time: Theory

- How much time will it take to parse?
  - For each diff (<= n)</p>
    - For each i (<= n)</p>
      - For each rule  $X \rightarrow Y Z$ 
        - For each split point k
           Do constant work



- Total time: |rules|\*n³
- Something like 5 sec for an unoptimized parse of a 20-word sentences, or 0.2sec for an optimized parser

# Agenda-Based Parsing

- Agenda-based parsing is like graph search (but over a hypergraph)
- Concepts:
  - Numbering: we number fenceposts between words
  - "Edges" or items: spans with labels, e.g. PP[3,5], represent the sets of trees over those words rooted at that label (cf. search states)
  - A chart: records edges we've expanded (cf. closed set)
  - An agenda: a queue which holds edges (cf. a fringe or open set)

PP



#### Word Items

- Building an item for the first time is called discovery.
   Items go into the agenda on discovery.
- To initialize, we discover all word items (with score 1.0).



critics[0,1], write[1,2], reviews[2,3], with[3,4], computers[4,5]

#### **CHART [EMPTY]**

0 1 2 3 4 5

critics write reviews with computers

#### Item Successors

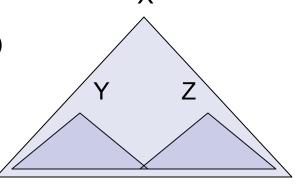
- When we pop items off of the agenda:
  - Graph successors: unary projections (NNS → critics, NP → NNS)

$$Y[i,j]$$
 with  $X \rightarrow Y$  forms  $X[i,j]$ 

Hypergraph successors: combine with items already in our chart

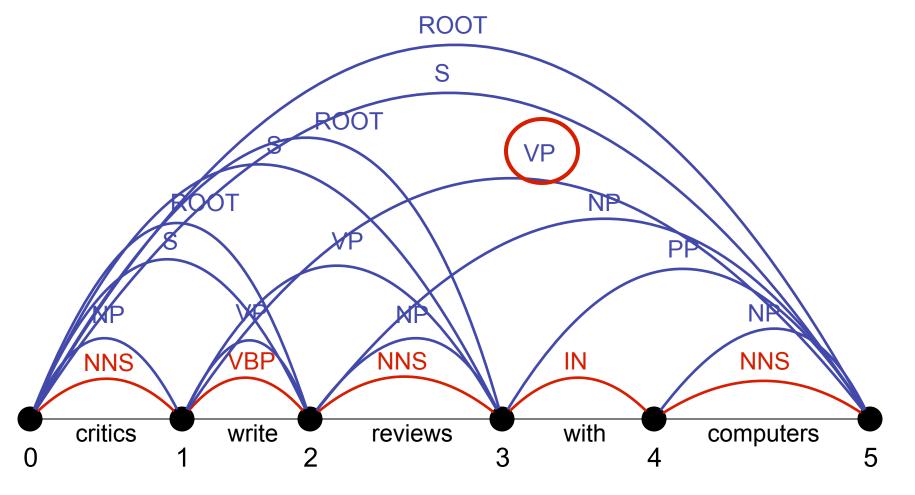
$$Y[i,j]$$
 and  $Z[j,k]$  with  $X \rightarrow Y Z$  form  $X[i,k]$ 

- Enqueue / promote resulting items (if not in chart already)
- Record backtraces as appropriate
- Stick the popped edge in the chart (closed set)
- Queries a chart must support:
  - Is edge X:[i,j] in the chart? (What score?)
  - What edges with label Y end at position j?
  - What edges with label Z start at position i?



### An Example

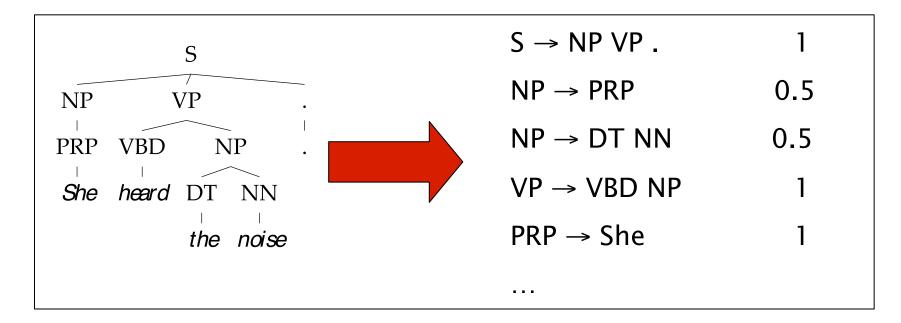
NNS[0,1] VBP[1,2] NNS[2,3] IN[3,4] NNS[3,4] NP[0,1] VP[1,2] NP[2,3] NP[4,5] S[0,2] VP[1,3] PP[3,5] ROOT[0,2] S[0,3] VP[1,5] NP[2,5] ROOT[0,3] S[0,5] ROOT[0,5]



#### Treebank Grammars

[Charniak '96]

- Need a PCFG for broad coverage parsing.
- Can take a grammar right off the trees (doesn't work well):

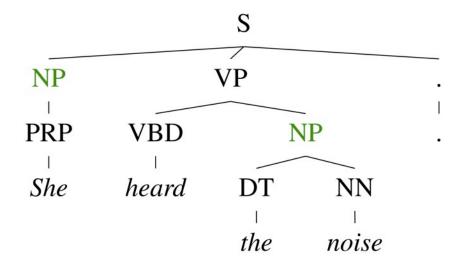


- Better results by enriching the gramma
- Can also get reasonable parsers witho

Model	F1
Charniak '96	72.0

### Conditional Independence?

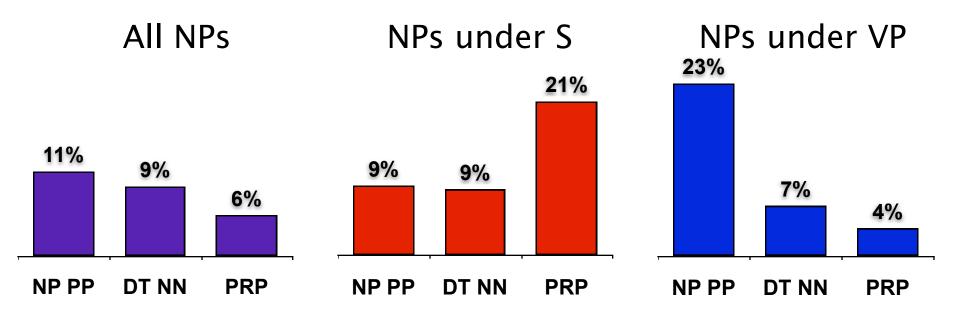
Not every NP expansion can fill every NP slot



- A grammar with symbols like "NP" won't be context-free
- Statistically, conditional independence too strong

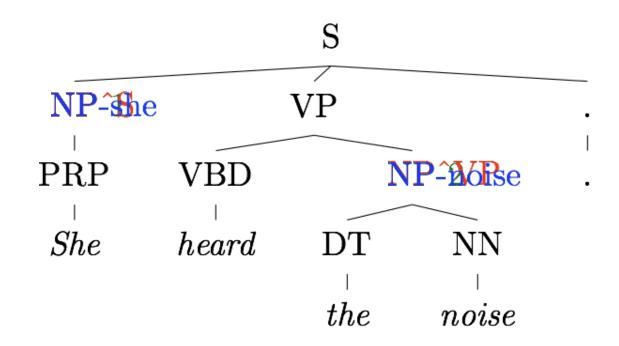
### Non-Independence

Independence assumptions are often too strong.



- Example: the expansion of an NP is highly dependent on the parent of the NP (i.e., subjects vs. objects).
- Also: the subject and object expansions are correlated!

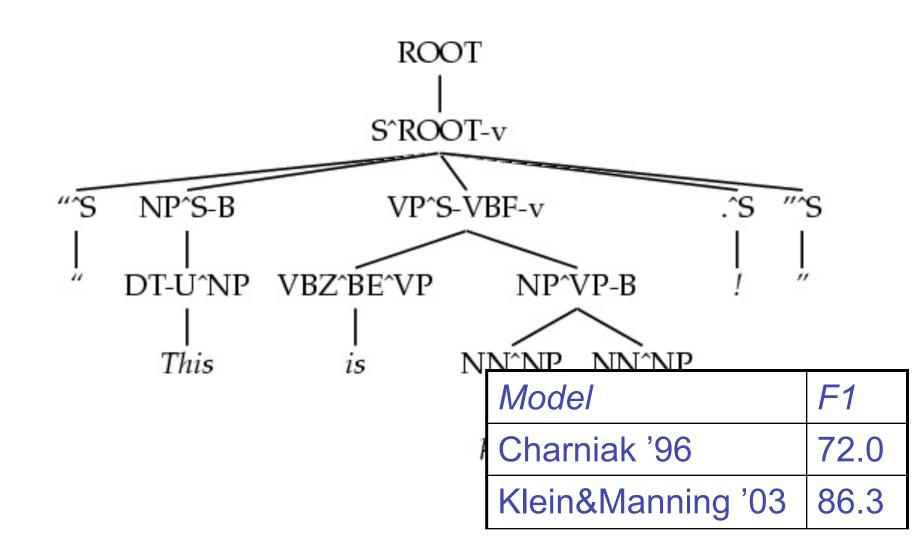
#### The Game of Designing a Grammar



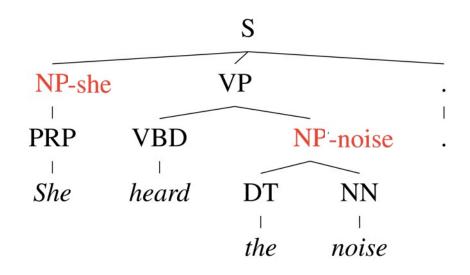
- Structure Annotation [Johnson '98, Klein&Manning '03]
- Lexicalization [Collins '99, Charniak '00]
- Latent Variables [Matsuzaki et al. 05, Petrov et al. '06]

#### A Fully Annotated (Unlexicalized) Tree

[Klein & Manning '03]

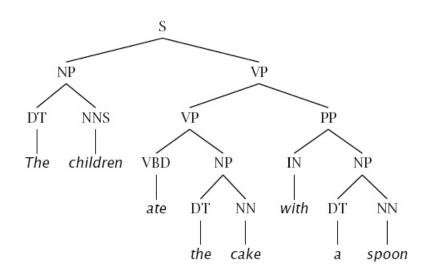


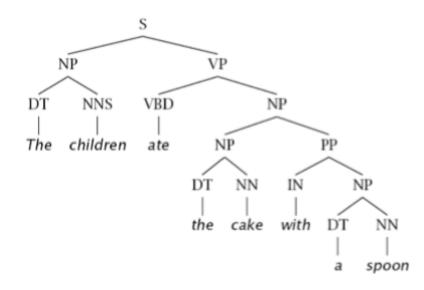
### The Game of Designing a Grammar



- Annotation refines base treebank symbols to improve statistical fit of the grammar
  - Head lexicalization [Collins '99, Charniak '00]

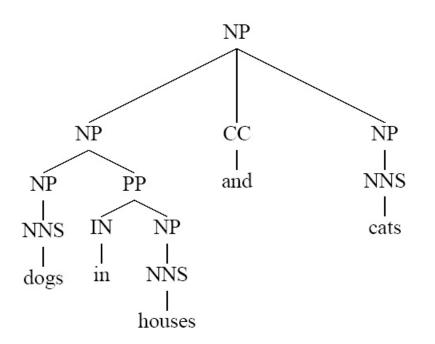
#### Problems with PCFGs

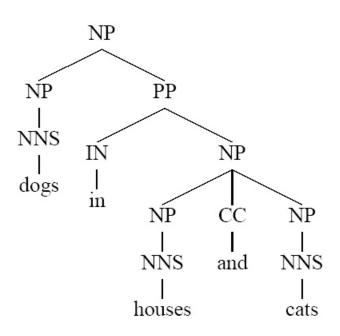




- If we do no annotation, these trees differ only in one rule:
  - VP → VP PP
  - NP → NP PP
- Parse will go one way or the other, regardless of words
- We addressed this in one way with unlexicalized grammars (how?)
- Lexicalization allows us to be sensitive to specific words

#### Problems with PCFGs



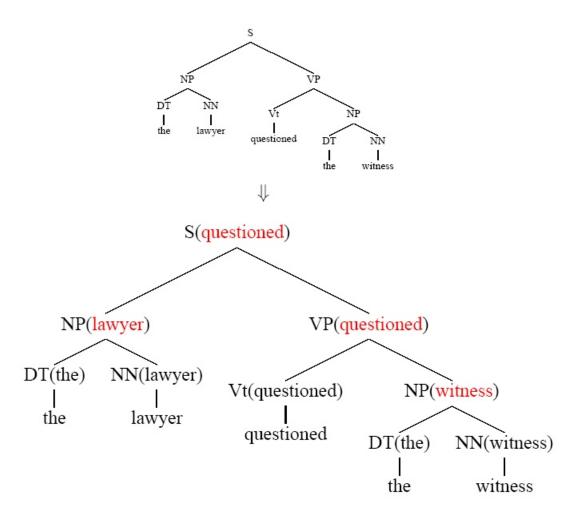


- What's different between basic PCFG scores here?
- What (lexical) correlations need to be scored?

#### Lexicalized Trees

# [Charniak '97, Collins '97]

- Add "headwords" to each phrasal node
  - Syntactic vs. semantic heads
  - Headship not in (most) treebanks
  - Usually use head rules, e.g.:
    - NP:
      - Take leftmost NP
      - Take rightmost N\*
      - Take rightmost JJ
      - Take right child
    - VP:
      - Take leftmost VB\*
      - Take leftmost VP
      - Take left child



### Lexicalized PCFGs?

Problem: we now have to estimate probabilities like

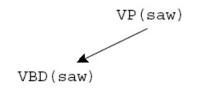
```
VP(saw) -> VBD(saw) NP-C(her) NP(today)
```

- Never going to get these atomically off of a treebank
- Solution: break up derivation into smaller steps

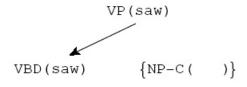


### **Lexical Derivation Steps**

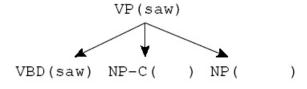
A derivation of a local tree [Collins '99]



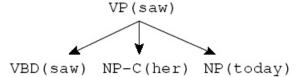
Choose a head tag and word



Choose a complement bag



Generate children (incl. adjuncts)

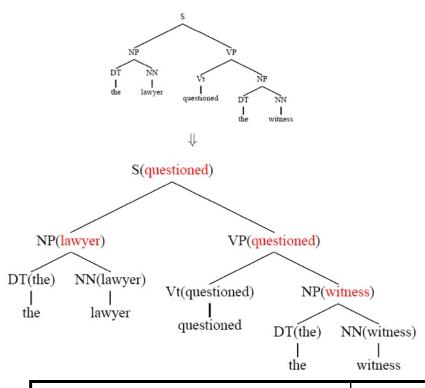


Recursively derive children

#### Lexicalized Grammars

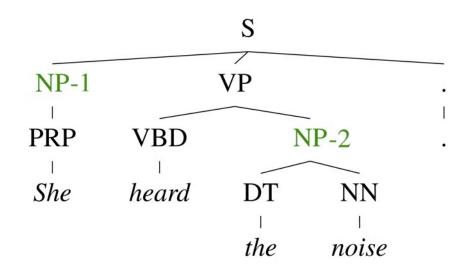
#### Challenges:

- Many parameters to estimate: requires sophisticated smoothing techniques
- Exact inference is too slow: requires pruning heuristics
- Difficult to adapt to new languages: At least head rules need to be specified, typically more changes needed



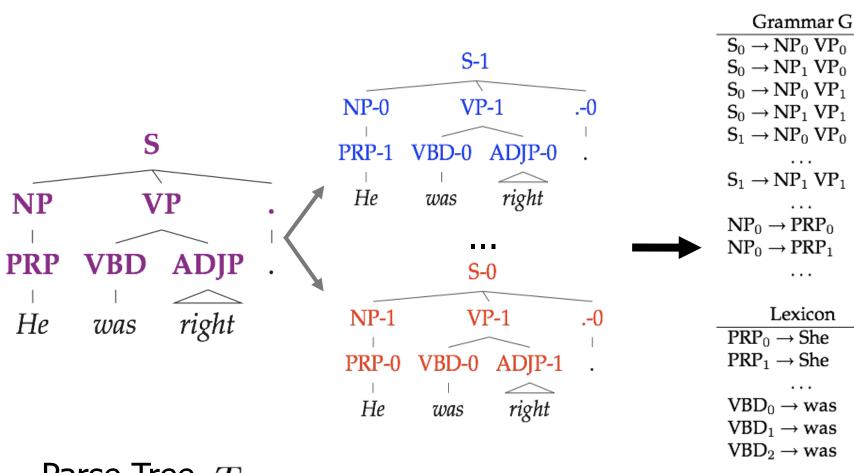
Model	F1
Klein&Manning '03	86.3
Charniak '00	90.1

#### The Game of Designing a Grammar



- Annotation refines base treebank symbols to improve statistical fit of the grammar
  - Automatic clustering

#### Latent Variable Grammars



Parse Tree T Sentence  $\eta \eta$ 

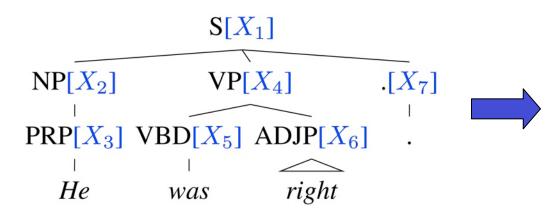
Derivations t:T

Parameters  $\theta$ 

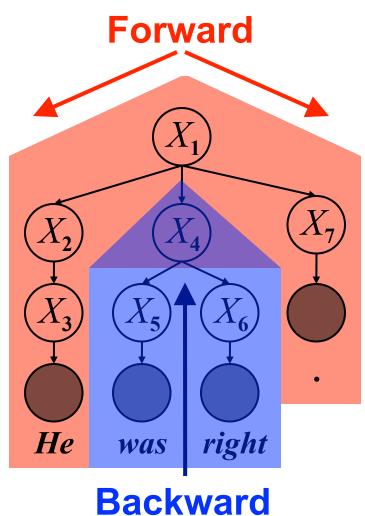
#### **Learning Latent Annotations**

#### EM algorithm:

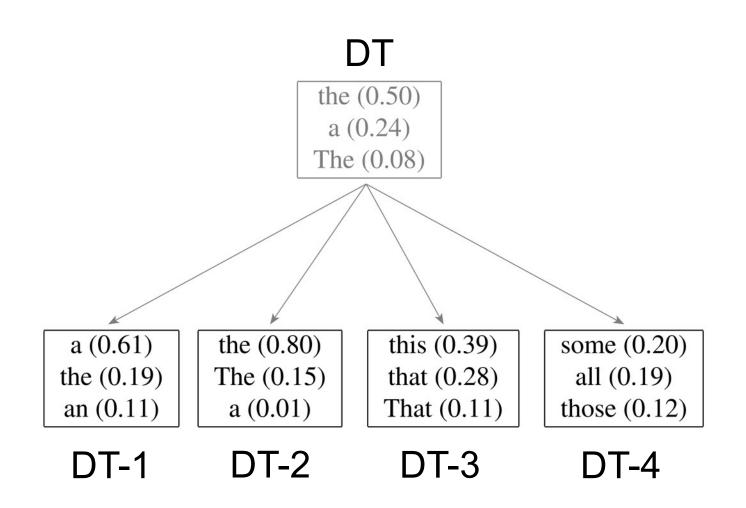
- Brackets are known
- Base categories are known
- Only induce subcategories



Just like Forward-Backward for HMMs.

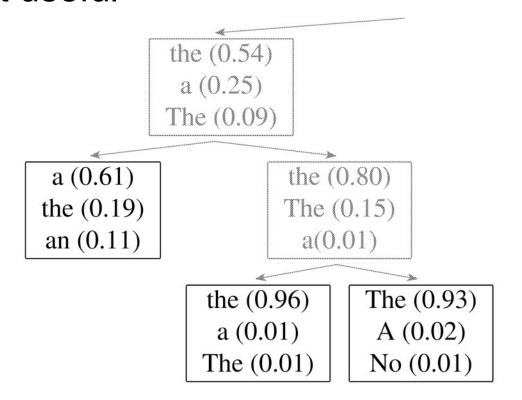


#### Refinement of the DT tag



### Adaptive Splitting

- Want to split complex categories more
- Idea: split everything, roll back splits which were least useful



#### Learned Splits

Proper Nouns (NNP):

NNP-14	Oct.	Nov.	Sept.
NNP-12	John	Robert	James
NNP-2	J.	E.	L.
NNP-1	Bush	Noriega	Peters
NNP-15	New	San	Wall
NNP-3	York	Francisco	Street

Personal pronouns (PRP):

PRP-0	It	He	
PRP-1	it	he	they
PRP-2	it	them	him

#### Learned Splits

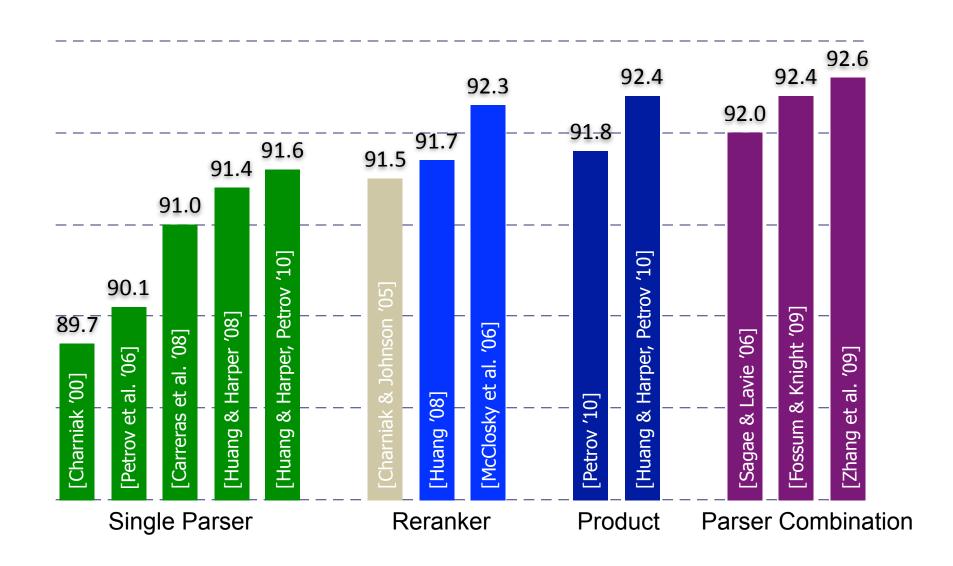
Relative adverbs (RBR):

RBR-0	further	lower	higher
RBR-1	more	less	More
RBR-2	earlier	Earlier	later

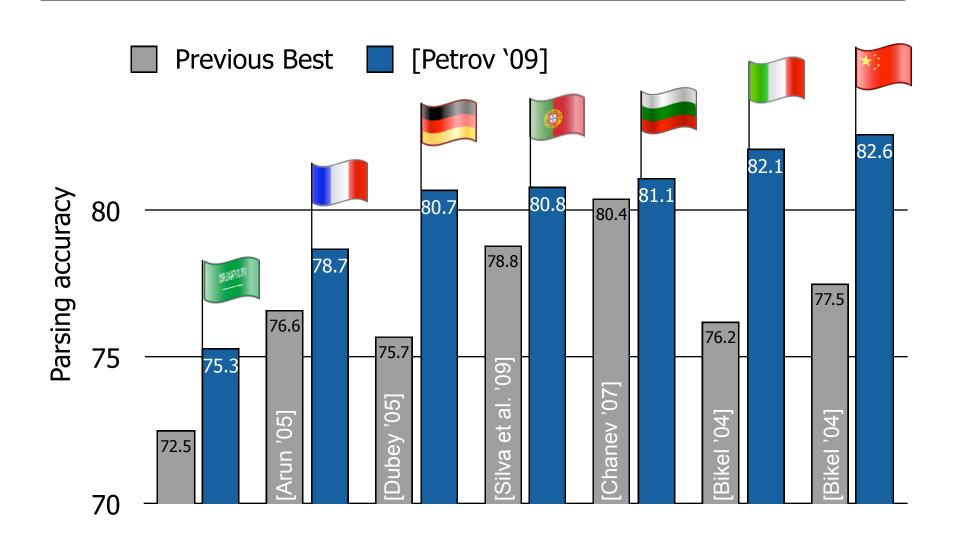
Cardinal Numbers (CD):

CD-7	one	two	Three
CD-4	1989	1990	1988
CD-11	million	billion	trillion
CD-0	1	50	100
CD-3	1	30	31
CD-9	78	58	34

#### Detailed English Results



### Multi-Lingual Results

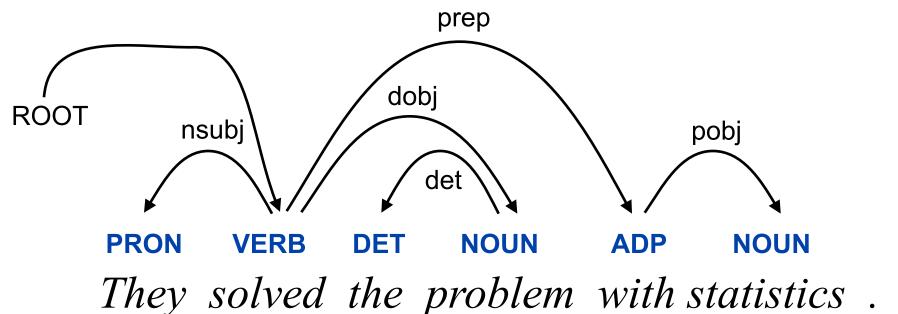


## Syntax and Parsing

Slav Petrov – Google Research

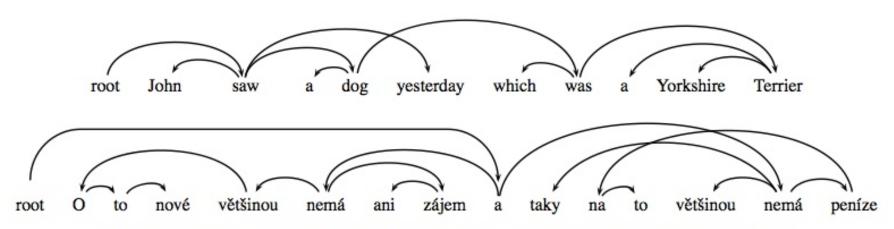
**Lisbon Machine Learning School** 

#### Dependency Parsing



### (Non-)Projectivity

- Crossing Arcs needed to account for nonprojective constructions
- Fairly rare in English but can be common in other languages (e.g. Czech):



He is mostly not even interested in the new things and in most cases, he has no money for it either.

#### **Formal Conditions**

- For a dependency graph G = (V, A)
- ▶ With label set  $L = \{l_1, \ldots, l_{|L|}\}$
- ► *G* is (weakly) connected:
  - ▶ If  $i, j \in V$ ,  $i \leftrightarrow^* j$ .
- ▶ G is acyclic:
  - ▶ If  $i \rightarrow j$ , then not  $j \rightarrow^* i$ .
- ► *G* obeys the single-head constraint:
  - ▶ If  $i \rightarrow j$ , then not  $i' \rightarrow j$ , for any  $i' \neq i$ .
- G is projective:
  - ▶ If  $i \rightarrow j$ , then  $i \rightarrow^* i'$ , for any i' such that i < i' < j or j < i' < i.

#### **Arc-Factored Models**

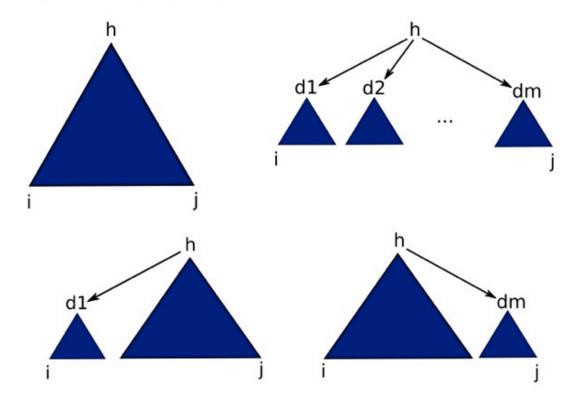
Assumes that the score / probability / weight of a dependency graph factors by its arcs

$$w(G) = \prod_{(i,j,k)\in G} w_{ij}^k$$
 look familiar?

- $w_{ij}^k$  is the weight of creating a dependency from word  $w_i$  to  $w_j$  with label  $I_k$
- Thus there is an assumption that each dependency decision is independent
  - Strong assumption! Will address this later.

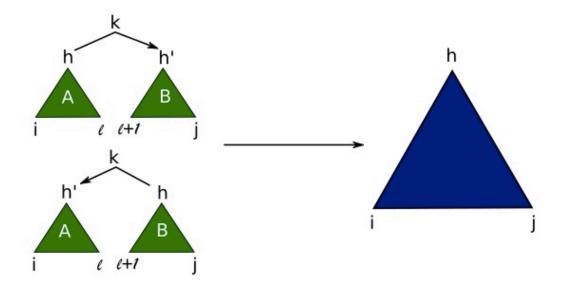
# Arc-factored Projective Parsing

All projective graphs can be written as the combination of two smaller adjacent graphs



### Arc-factored Projective Parsing

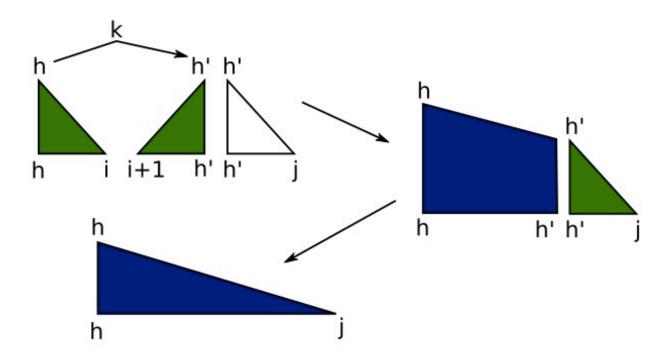
- Chart item filled in a bottom-up manner
  - ▶ First do all strings of length 1, then 2, etc. just like CKY



- ▶ Weight of new item:  $\max_{l,j,k} w(A) \times w(B) \times w_{hh'}^{k}$
- ▶ Algorithm runs in  $O(|L|n^5)$
- Use back-pointers to extract best parse (like CKY)

#### Eisner Algorithm

- $ightharpoonup O(|L|n^5)$  is not that good
- ▶ [Eisner 1996] showed how this can be reduced to  $O(|L|n^3)$ 
  - Key: split items so that sub-roots are always on periphery

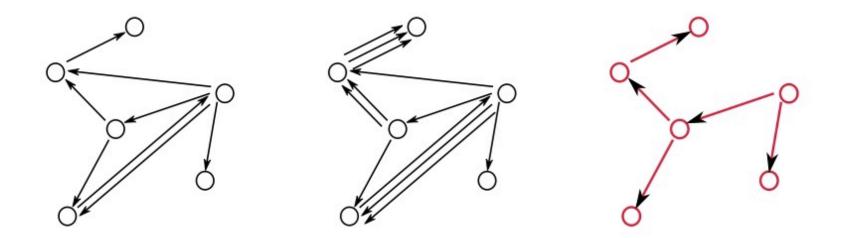


#### Eisner Algorithm PseudoCode

```
Initialization: C[s][s][d][c] = 0.0 \quad \forall s, d, c
for k:1..n
  for s:1..n
    t = s + k
    if t > n then break
     % First: create incomplete items
    C[s][t][\leftarrow][0] = \max_{s \le r < t} (C[s][r][\rightarrow][1] + C[r+1][t][\leftarrow][1] + s(t,s))
    C[s][t][\to][0] = \max_{s \le r \le t} (C[s][r][\to][1] + C[r+1][t][\leftarrow][1] + s(s,t))
     % Second: create complete items
    C[s][t][\leftarrow][1] = \max_{s \le r \le t} (C[s][r][\leftarrow][1] + C[r][t][\leftarrow][0])
    C[s][t][\to][1] = \max_{s < r < t} (C[s][r][\to][0] + C[r][t][\to][1])
  end for
end for
```

### Maximum Spanning Trees (MSTs)

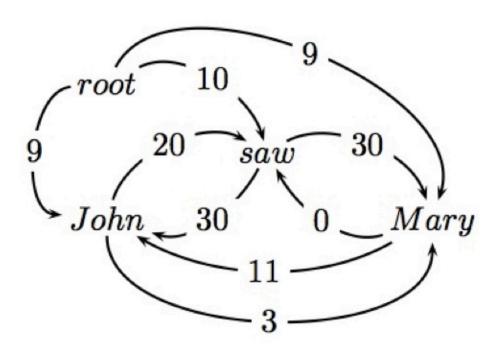
- A directed spanning tree of a (multi-)digraph G = (V, A), is a subgraph G' = (V', A') such that:
  - V' = V
  - $ightharpoonup A' \subseteq A$ , and |A'| = |V'| 1
  - ► G' is a tree (acyclic)
- A spanning tree of the following (multi-)digraphs



Can use MST algorithms for nonprojective parsing!

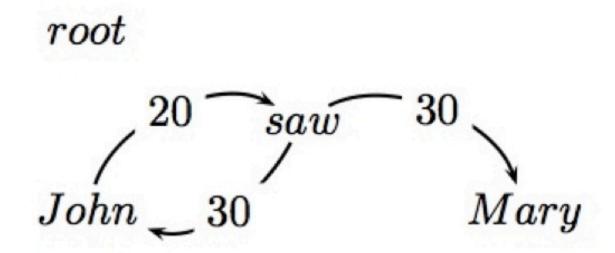
#### Chu-Liu-Edmonds

 $\triangleright x = \text{root John saw Mary}$ 



#### Chu-Liu-Edmonds

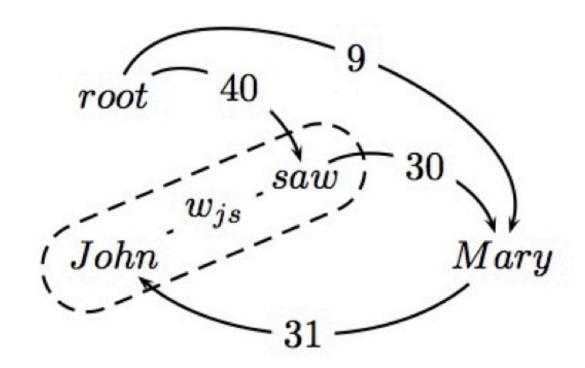
Find highest scoring incoming arc for each vertex



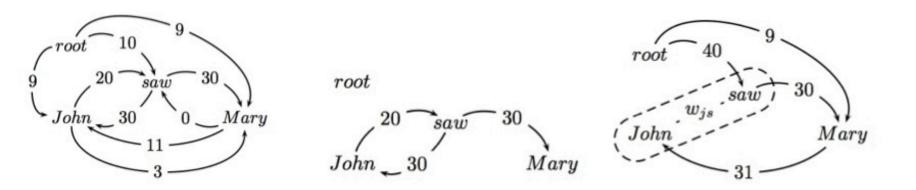
If this is a tree, then we have found MST!!

### Find Cycle and Contract

- ▶ If not a tree, identify cycle and contract
- Recalculate arc weights into and out-of cycle



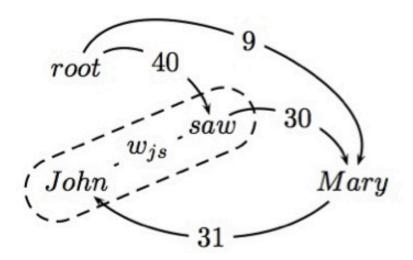
### Recalculate Edge Weights



- Incoming arc weights
  - Equal to the weight of best spanning tree that includes head of incoming arc, and all nodes in cycle
  - root → saw → John is 40 (\*\*)
  - root → John → saw is 29

#### Theorem

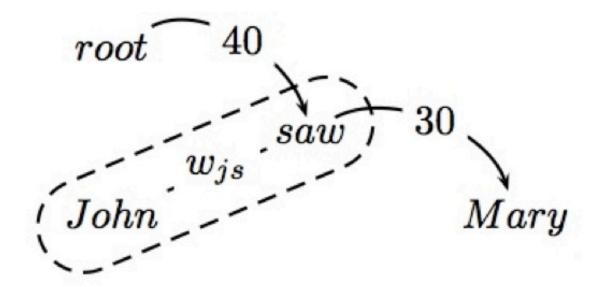
The weight of the MST of this contracted graph is equal to the weight of the MST for the original graph



Therefore, recursively call algorithm on new graph

#### Final MST

This is a tree and the MST for the contracted graph!!



▶ Go back up recursive call and reconstruct final graph

#### Chu-Liu-Edmonds PseudoCode

#### Chu-Liu-Edmonds( $G_x, w$ )

- 1. Let  $M = \{(i^*, j) : j \in V_x, i^* = \arg\max_{i'} w_{ij}\}$
- 2. Let  $G_M = (V_x, M)$
- 3. If  $G_M$  has no cycles, then it is an MST: return  $G_M$
- Otherwise, find a cycle C in G<sub>M</sub>
- 5. Let  $\langle G_C, c, ma \rangle = \text{contract}(G, C, w)$
- 6. Let  $G = \text{Chu-Liu-Edmonds}(G_C, w)$
- 7. Find vertex  $i \in C$  such that  $(i', c) \in G$  and ma(i', c) = i
- 8. Find arc  $(i'', i) \in C$
- 9. Find all arc  $(c, i''') \in G$
- 10.  $G = G \cup \{(ma(c, i'''), i''')\}_{\forall (c, i''') \in G} \cup C \cup \{(i', i)\} \{(i'', i)\}$
- 11. Remove all vertices and arcs in G containing c
- return G
  - ▶ Reminder:  $w_{ij} = \arg \max_k w_{ij}^k$

#### Chu-Liu-Edmonds PseudoCode

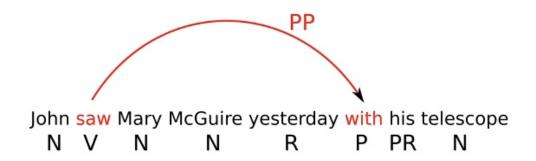
```
contract(G = (V, A), C, w)
     Let G_C be the subgraph of G excluding nodes in C
   Add a node c to G_C representing cycle C
3.
     For i \in V - C: \exists_{i' \in C}(i', i) \in A
        Add arc (c, i) to G_C with
           ma(c, i) = \arg \max_{i' \in C} score(i', i)
           i' = ma(c, i)
           score(c, i) = score(i', i)
    For i \in V - C: \exists_{i' \in C}(i, i') \in A
        Add edge (i, c) to G_C with
           ma(i, c) = \arg \max_{i' \in C} [score(i, i') - score(a(i'), i')]
           i' = ma(i, c)
           score(i, c) = [score(i, i') - score(a(i'), i') + score(C)]
              where a(v) is the predecessor of v in C
              and score(C) = \sum_{v \in C} score(a(v), v)
5.
      return \langle G_C, c, ma \rangle
```

#### **Arc Weights**

$$w_{ij}^k = e^{\mathbf{W} \cdot \mathbf{f}(i,j,k)}$$

- Arc weights are a linear combination of features of the arc, f, and a corresponding weight vector w
- Raised to an exponent (simplifies some math ...)
- What arc features?
- ► [McDonald et al. 2005] discuss a number of binary features

### Arc Feature Ideas for f(i,j,k)



- Identities of the words wi and wj and the label lk
- Part-of-speech tags of the words wi and wj and the label lk
- Part-of-speech of words surrounding and between wi and wj
- Number of words between wi and wj, and their orientation
- Combinations of the above

### (Structured) Perceptron

```
Training data: T = \{(x_t, G_t)\}_{t=1}^{|T|}
1. \mathbf{w}^{(0)} = 0; i = 0
2. for n : 1..N
3. for t:1...T
           Let G' = \arg \max_{G'} \mathbf{w}^{(i)} \cdot \mathbf{f}(G')
5.
           if G' \neq G_t
              \mathbf{w}^{(i+1)} = \mathbf{w}^{(i)} + \mathbf{f}(G_t) - \mathbf{f}(G')
6.
       i = i + 1
7.
     return wi
8.
```

#### Partition Function

Partition Function: 
$$Z_x = \sum_{G \in T(G_x)} w(G)$$

▶ Lapacian Matrix Q for graph  $G_x = (V_x, A_x)$ 

$$Q_{jj} = \sum_{i \neq j, (i,j,k) \in A_x} w_{ij}^k$$
 and  $Q_{ij} = \sum_{i \neq j, (i,j,k) \in A_x} -w_{ij}^k$ 

Cofactor Q<sup>i</sup> is the matrix Q with the i<sup>th</sup> row and column removed

> The Matrix Tree Theorem [Tutte 1984] The determinant of the cofactor  $Q^0$  is equal to  $Z_x$

- ▶ Thus  $Z_x = |Q^0|$  determinants can be calculated in  $O(n^3)$
- ▶ Constructing Q takes  $O(|L|n^2)$
- ▶ Therefore the whole process takes  $O(n^3 + |L|n^2)$

#### **Arc Expectations**

$$\langle i,j,k\rangle_{\times} = \sum_{G\in\mathcal{T}(G_{\times})} w(G) \times \mathbb{1}[(i,j,k)\in A]$$

Can easily be calculated, first reset some weights

$$w_{i'j}^{k'} = 0 \ \forall i' \neq i \ \text{and} \ k' \neq k$$

- Now,  $\langle i, j, k \rangle_x = Z_x$
- Why? All competing arc weights to zero, therefore every non-zero weighted graph must contain (i, j, k)
- ▶ Naively takes  $O(n^5 + |L|n^2)$  to compute all expectations
- ▶ But can be calculated in  $O(n^3 + |L|n^2)$  (see [McDonald and Satta 2007, Smith and Smith 2007, Koo et al. 2007])

#### Summary

- Constituency Parsing
  - CKY Algorithm
  - Lexicalized Grammars
  - Latent Variable Grammars

- Dependency Parsing
  - Eisner Algorithm
  - Maximum Spanning Tree Algorithm

# There is lots more and these models are being actively used in practice!