

MIS9

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1.

1. From equation 15.10 that:

$$A^T (AA^T + \lambda I)^{-1} = (A^T A + \lambda I)^{-1} A^T$$

$$\hat{x} = A^T (A^T A + \lambda I)^{-1} b$$

First, form $AA^T + \lambda I$ ($m^2 n$ flops)

Then, $LL^T = AA^T + \lambda I$ ($\frac{1}{3} m^3$ flops)

For $L^{-1} b$ and $L^{-T} L^{-1} b$, each costs n^2 flops

Last for $A^T L^{-T} L^{-1} b$ ($2mn$ flops)

So, the result based on first two steps, which is $O(m^2 n + m^3)$

2.

2. (a) The characteristic polynomial of C is $p(x) = \det(xI - C)$

$$\det(xI - C) = \begin{vmatrix} x & 0 & \cdots & 0 & C_0 \\ -1 & x & & 0 & C_1 \\ 0 & -1 & & 0 & C_2 \\ \vdots & & \ddots & & \\ 0 & 0 & & -1 & x + C_{n-1} \end{vmatrix}$$

$$= x \begin{vmatrix} x & 0 & \cdots & 0 & C_1 \\ -1 & x & & 0 & C_2 \\ \vdots & & \ddots & & \\ 0 & 0 & \cdots & -1 & x + C_{n-1} \end{vmatrix} + (-1)^{n+1} C_0 \begin{vmatrix} -1 & x & \cdots & 0 \\ 0 & -1 & \cdots & 0 \\ \vdots & & \ddots & \\ 0 & 0 & \cdots & -1 \end{vmatrix}$$

$$= x(x^{n-1} + C_{n-1}x^{n-2} + \cdots + C_2x + C_1) + (-1)^{n+1} C_0 \cdot (-1)^{n-1}$$

$$= x^n + C_{n-1}x^{n-1} + \cdots + C_1x + C_0$$

$$= x^n + \sum_{k=0}^{n-1} C_k x^k$$

(b) For C has n distinct eigenvalues, we have $\lambda_i^n = \sum_{k=0}^{n-1} C_k \lambda_i^k = 0 \Rightarrow \lambda_i^n = -\sum_{k=0}^{n-1} C_k \lambda_i^k$

Then we choose the i th row of V

$$V_i^T = [1 \ \lambda_i \ \lambda_i^2 \ \cdots \ \lambda_i^{n-1}]$$

$$V_i^T C = [\lambda_i \ \cdots \ \lambda_i^{n-1} \ -\sum_{k=0}^{n-1} C_k \lambda_i^k]$$

$$= [\lambda_i \ \lambda_i^2 \ \cdots \ \lambda_i^n]$$

$$= \lambda_i V_i^T$$

So, we can have $VC = \Lambda V$

3.

(a) For $C = F^* \Lambda F$, we have $FC = \Lambda F$,

and in the j th row, $f_j^T C = \lambda_j \cdot f_j^T$

$$\frac{1}{\sqrt{n}} \sum_{d=1}^{k-1} w^{(j-1)(d-1)} \cdot C_{n-k+d} + \sum_{d=k}^n w^{(j-1)(d-1)} C_{d-k+1}$$

As we know that $w = e^{\frac{-2\pi i}{n}}$, so $w^d = w^{d+n}$

$$\text{and } \frac{1}{\sqrt{n}} \sum_{d=k}^{n+k-1} w^{(j-1)(d-1)} C_{d-k+1}$$

Set d with $d-k+1$, we have

$$\frac{1}{\sqrt{n}} w^{k-1} \sum_{d=1}^n w^{(j-1)(d-1)} C_d, \text{ which is}$$

the k th entry of $f_j^T C$

Based on $\Lambda = \text{diag}(F_c)$,

$$\lambda_j = \sum_{d=1}^n w^{(j-1)(d-1)} C_d$$

Then we can have the k th entry

of $\lambda_j f_j^T$ is

$$\frac{1}{\sqrt{n}} \lambda_j w^{k-1} = \frac{1}{\sqrt{n}} w^{k-1} \sum_{d=1}^n w^{(j-1)(d-1)} C_d$$

In the end of all $j, k = 1, \dots, n$,

$$[FC]_{jk} = [\Lambda F]_{jk}, \text{ so } C = F^* \Lambda F.$$

(b) To solve $Cx = b$, we have $x = C^{-1}b$

, so $x = F^* \Lambda^{-1} F b$, which $\Lambda = \text{Diag}(F_c)$

① Compute Fb via FFT ($n \log n$)

② Compute F_c via FFT ($n \log n$)

③ Compute $\Lambda^{-1} Fb$ (n)

④ Compute $F^* \Lambda^{-1} Fb$ ($n \log n$)

Therefore, the complexity is $O(n \log n)$

4.

4.

(a) Replace A with $U\Sigma V^T$ in $A^T A \hat{x} = A^T b$

$$A^T A \hat{x} = V \Sigma U^T U \Sigma V^T V \Sigma^T U^T b$$

$$= V \Sigma \Sigma \Sigma^T U^T b$$

$$= V \Sigma U^T b$$

$$= A^T b$$

(b) For \hat{x} that satisfies $A^T A \hat{x} = A^T b$

$$\|Ax - b\|^2 = \|Ax - A\hat{x} + A\hat{x} - b\|^2$$

$$= \|Ax - A\hat{x}\|^2 - 2(Ax - A\hat{x})^T (A\hat{x} - b) + \|A\hat{x} - b\|^2$$

$$= \|Ax - A\hat{x}\|^2 + \|A\hat{x} - b\|^2 - 2(x - \hat{x})^T (A^T A \hat{x} - A^T b)$$

As we know $A^T A x - A^T b = 0$ and $\|Ax - A\hat{x}\|^2$

≥ 0

So, $\|Ax - b\|^2 \geq \|A\hat{x} - b\|^2$, any \hat{x} satisfies

$A^T A \hat{x} = A^T b$ is an optimal solution that minimizes $\|Ax - b\|^2$.