COMP 251 Assignment 3 - Written

3)

Given an array of integers A, there can be 3 cases regarding implementing active lists with A[i]:

- 1. A[i] is the smallest, so we create a new list of length 1
- 2. A[i] is the largest, so we clone existing list and add A[i] to end
- 3. A[i] is in the middle, so we find the list with the largest end element thats smaller than A[i], clone and add A[i] to end and we discard lists with the same length as our new list

Note: Here, we have the end element of smaller lists always smaller than the end element of larger lists

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//Pseudocode
longestIncrementalSubseunce(int[] s)
      // Keep track of predecessors and ends of each subsequence
      int[] preds = new int[s.length]
      int[] ends = new int[s.length + 1]
      // Start off with 0 length
      int length = 0;
      for i = 0 to s.length - 1
            // Binary search
            int low = 1
             int high = length
            while low <= high
                   int middle = Ceil((low + high) / 2)
                   if s[ends[middle] < s[i]]
                          low = middle + 1
                   else
                          high = middle - 1
            // Update preds, and replace ends
```

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preds[i] = ends[low]
  ends[low] = i

// Update length
  if (low > length)
        length = low

// Get longest increment subsequence
int[] longestIncSub = new int[length]
int j = ends[length]

for i = length - 1 to 0
        longestIncSub[i] = s[j]
        j = preds[j]
```

return longestIncSub

Binary search time is $O(\log n)$, which runs in a loop of n times. Therefore, run time is $O(n \log n)$

