# Math 356: Number Theory (Rutgers)

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#### Abstract

These notes arise from lecture videos of Math 356: Number Theory (Quadratic Forms), originally taught by Professor Alex Kontorovich, at Rutgers University. I am grateful to Professor Kontorovich for releasing the lecture videos online. I am responsible for all faults in this document, mathematical or otherwise; any merits of the material here should be credited to Professor Kontorovich and not to me. Feel free to message me with any suggestions or corrections at dyang5@swarthmore.edu.

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# 1 Pythagorean Triples

**Theorem 1.1** (Fermat's Last Theorem, 1637). There are no positive integers x, y, z that satisfy

$$x^n + y^n = z^n$$

for integer n > 2.

An equivalent formulation is the one discussed in lecture:

$$x^n + y^n = z^n$$

where n > 2 is an integer, has no non-trivial solutions x, y, and z in  $\mathbb{Q}, \mathbb{Z}$ .

**Definition 1.1** (Diophantine Problem). A **Diophantine problem** is a polynomial equation solved in  $\mathbb{Z}$  and  $\mathbb{Q}$ .

Fermat's Last Theorem is an example of a Diophantine problem.

Historically, in 1993, Andrew Wiles published a proof of Fermat's Last Theorem. Later, Wiles and his student Richard Taylor rectified the 1993 proof and published the first successful proof. For these efforts, Wiles won the 2016 Abel Prize and the 2017 Copley Medal.

With some extra background, we can prove Fermat's Last Theorem for n = 4.

#### 1.1 Parameterization of Pythagorean Triples

**Theorem 1.2** (Primitive Pythagorean Triple Generation). If (x, y, z) is a **primitive** pythagorean triple<sup>a</sup> ( $\nexists d$  such that  $d \mid x, d \mid y$  and  $d \mid z$ ), then z is odd. Furthermore, assuming that x is odd and y is even, then there exists coprime  $r, s \in \mathbb{Z}$  with such that

$$x = r^2 - s^2$$
,  $y = 2rs$ ,  $z = r^2 + s^2$ .

**Proof Sketch (Theorem 3).** Since y is even, we can rewrite  $y = 2y_1$  for some integer  $y_1$ . Since (x, y, z) are a Pythagorean triple, it follows that  $y^2 = z^2 - x^2 = (z - x)(z + x)$ . Substituting and expanding, we get that

$$4y_1^2 = (z - x)(z + x).$$

Since x, z are both odd, we can rewrite z - x = 2a and z + x = 2b, so

$$y_1^2 = ab$$

where (a,b)=1. Consequently, a,b must themselves be perfect squares. We conclude that  $y_1=rs$ , where  $r^2=a$ ,  $s^2=b$ , and so it follows that y=2rs for coprime r, s.

<sup>&</sup>lt;sup>a</sup>if (x, y, z) is a primitive Pythagorean triple, then it follows that x, y, and z are pairwise primitive.

## 1.2 Pythagorean Varieties

**Definition 1.2** (Variety). A variety is a set of solutions to a system of polynomial equations.

Consider the variety  $\mathcal{V}: x^2 + y^2 - z^2$ , which is a Pythagorean variety. Note that  $\mathcal{V}(\mathbb{R})$  is a structure consisting of two cones which are reflections about the z-axis, and  $\mathcal{V}(\mathbb{Z})$ ,  $\mathcal{V}(\mathbb{Q})$  are subsets of  $\mathcal{V}(\mathbb{R})$ .

**Lemma 1.1.** For every  $\vec{v} \in \mathcal{V}(\mathbb{Q})$ , there exists a unique primitive  $\vec{w} \in \mathcal{V}(\mathbb{Z})^0_+$  with  $\vec{w} \sim \vec{v}$ , where  $\vec{u} \sim \vec{v} \iff \vec{u} = \lambda \vec{v}, \ \lambda \in \mathbb{R} \setminus \{0\}.$ 

Note that  $\mathcal{V}(\mathbb{R})_+/\sim$  can be thought of as  $S^1$ ; for a point  $(x,y,z)\in\mathcal{V}(\mathbb{R})_+\mapsto\left(\frac{x}{z},\frac{y}{z},1\right)$ . Furthermore,  $\mathcal{V}(\mathbb{R})/\sim=S^1\cup\{0\}$ .

Furthermore, note that another set of representatives of  $\mathcal{V}(\mathbb{Q})_+/\sim$  is  $S^1(\mathbb{Q})$ . This gives a notion for why Pythagorean triples are fundamental: they parameterize rational points on the unit circle. Formally,

$$\left(\frac{r}{u}, \frac{t}{u}\right) \in S^1(\mathbb{Q}) \iff (r, t, u) \in \mathcal{V}(\mathbb{Z})^0_+$$

#### 1.3 Proof of Fermat's Last Theorem for n = 4

Statement.  $x^4 + y^4 = z^4$  in  $\mathbb{Z}$  implies xyz = 0. Stronger Statement.  $x^4 + y^4 = z^2$  in  $\mathbb{Z}$  implies xyz = 0.

**Proof.** Suppose there is a solution to  $x^4 + y^4 = z^2$  in  $\mathbb{Z}$  with  $xyz \neq 0$ . Then  $(x^2, y^2, z)$  is a Pythagorean triple. Furthermore, since we can scale Pythagorean triples by their GCD to get a primitive triple, we can assume that  $(x^2, y^2, z) \in \mathcal{V}(\mathbb{Z})^0_+$  is rather a primitive Pythagorean triple. By the Parameterization of primitive Pythagorean triples, there exists coprime  $r, s \in \mathbb{Z}$  such that

$$x^2 = r^2 - s^2$$
,  $y^2 = 2rs$ ,  $z = r^2 + s^2$ .

Furthermore, since  $x^2$  is odd and  $y^2$  is even by assumption, r and s also have opposite parity. From the first equation, we know that

$$x^2 + s^2 = r^2,$$

where x is odd, forcing s to be even, and r to be odd. Once again, (x, s, r) is a primitive Pythagorean triple, so there exists coprime  $u, v \in \mathbb{Z}$  such that

$$x = u^2 - v^2$$
,  $s = 2uv$ ,  $r = u^2 + v^2$ 

where (r, s) = 1. We also know that  $y^2 = 2rs$  from our first primitive Pythagorean triple  $(x^2, y^2, z)$ . Since y is even, we can express it as  $y = 2y_1$  for some  $y_1 \in \mathbb{Z}$ . Since  $r = u^2 + v^2$ , s = 2uv, and  $y = 2y_1$ , we have that

$$4y_1^2 = 2(u^2 + v^2)(2uv)$$

and so

$$y_1^2 = uv(u^2 + v^2).$$

Note however that (u, v) = 1. Consequently, u, v, and  $u^2 + v^2$  are three pairwise coprime numbers whose product is a perfect square. Thus, u, v, and  $u^2 + v^2$  must themselves be a perfect square. We now write

$$u = a^2$$
,  $v = b^2$ ,  $u^2 + v^2 = c^2$ 

From the third equation, we have that  $a^4 + b^4 = c^2$ .

Note that if  $xyz \neq 0$ , then

$$c < c^2 = u^2 + v^2 < (u^2 + v^2)^2 = r^2 < r^2 + s^2 = z.$$

Thus, we have found a smaller, in the sense that c < z, primitive solution to  $x^4 + y^4 = z^2$  with  $abc \neq 0$ .

We can continue this process of finding primitive solutions until we arive at a trivial solution. Thus, by the principle of infinite descent<sup>a</sup>, we conclude that  $x^4 + y^4 = z^2$  in  $\mathbb{Z}$  implies xyz = 0.

<sup>&</sup>lt;sup>a</sup>or Fermat's descent

# 2 Number Theory Basics

## 2.1 Modular Arithmetic

Note: Basic modular arithmetic notes skipped due to prior knowledge. Important theorems/results added for completeness.

**Theorem 2.1** (Fermat's Little Theorem). If p is a prime with  $p \nmid a$ ,

$$a^{p-1} \equiv 1 \pmod{p}$$
.

**Proof.** Note that  $\{1 \cdot a, 2 \cdot a, \dots, (p-1) \cdot a, p \cdot a\}$  is a complete residue system modulo p. It follows that

$$(1 \cdot a)(2 \cdot a) \dots ((p-1) \cdot a) \equiv 1 \cdot 2 \cdot 3 \dots (p-1) \pmod{p}.$$

Simplifying, we have that

$$(p-1)!a^{p-1} \equiv (p-1) \pmod{p}.$$

Equivalently, we have that  $p \mid (p-1)! (a^{p-1}-1)$ . Since  $p \nmid (p-1)!$ , it follows that  $p \mid a^{p-1}-1$ , so

$$a^{p-1} \equiv 1 \pmod{p}$$
.

2.2 Linear Diophantine Equations (and Euclidean Algorithm)

Let  $S = \{an + bm \mid n, m \in \mathbb{Z}\}...$ 

**Theorem 2.2.**  $S = d\mathbb{Z}$ , where  $d = \gcd(n, m)$ .

**Proof.** First, note that  $S \subseteq d\mathbb{Z}$ . If  $d \mid n$  and  $d \mid m$ , then  $d \mid an$  and  $d \mid bm$ , so  $d \mid an + bm$ . It remains to show that an + bm = d has a solution. It is possible to find such a solution by reversing the steps of the Euclidean Algorithm using n and m.

**Theorem 2.3** (Euclidean Algorithm). The **Euclidean Algorithm** can be used to find the GCD of two integers n > m > 0 as follows:

- Apply the 'division' algorithm to rewrite  $n = m \cdot q_1 + r_1$ , where  $0 \le r_1 < m$ .
- Continue this process, finding  $q_i, r_i$  such that  $r_{i-2} = r_{i-1}q_i + r_i$ , with  $0 \le r_i < r_{i-1}$ .
- Stop at some finite  $J \leq m$  with  $r_J = 0$ .

It follows that  $r_{J-1} = \gcd(n, m)$ .

**Proof.** First, let's show that  $gcd(n, m) \mid r_{J-1}$ . Note that if  $l \mid n$  and  $l \mid n$ ,  $l \mid n - mq_1$ , so  $l \mid r_1$ . Similarly, if  $l \mid r_{j-2}$  and  $l \mid r_{j-1}$ , then  $l \mid r_{j-2} - r_{j-1}q_j$ , so  $l \mid r_j$ . Thus, it follows, that  $l \mid r_{J-1}$ ,

<sup>&</sup>lt;sup>a</sup>for reference, we can express n as  $r_{-1}$  and m as  $r_0$ .

and so  $gcd(n, m) \mid r_{J-1}$ .

It remains to show that  $r_{J-1} \mid \gcd(n, m)$ . Note that  $r_{J-2} = r_{J-1}q_J + r_J$ , and  $r_J = 0$ . It follows that  $r_{J-1} \mid r_{J-2}$ . Following our steps backwards, we will find that  $r_{J-1} \mid r_j$  for all  $j \in [-1, J-2]$ . Recall that  $r_0 = m$  and  $r_{-1} = n$ . Consequently, since  $r_{J-1} \mid r_j$  for all j, we have that  $r_{J-1} \mid n$  and  $r_{J-1} \mid m$ , so  $r_{J-1} \mid \gcd(n, m)$ .

Since  $\gcd(n,m) \mid r_{J-1}$  and  $r_{J-1} \mid \gcd(n,m)$ , it follows that  $r_{J-1} = \gcd(n,m)$ , as desired.  $\square$ 

## 2.3 Connections to Algebra

When we studied Linear Diophantine Equations, we looked at sets  $S = \{xn + ym \mid x, y \in \mathbb{Z}\}$ . Such a set S is an example of an ideal.

**Definition 2.1** (Ideals). An **ideal** of the ring  $\mathbb{Z}$  satisfies

- For  $z \in S$ ,  $rz \in S$  for all  $r \in \mathbb{Z}$ .
- For  $z_1, z_2 \in S$ ,  $z_1 \pm z_2 \in S$ .

**Definition 2.2.** An ideal S of the ring  $\mathbb{Z}$  is **principal** if there exists  $d \in \mathbb{Z}$  such that

$$S = (d) = \{d \cdot r \mid r \in \mathbb{Z}\} = d\mathbb{Z}.$$

**Theorem 2.4.**  $\mathbb{Z}$  is a **Principal Ideal Domain** (every ideal is principal).

Equivalently, suppose that  $S = (n_1, \ldots, n_k)$  is an ideal. Then there exists  $d \in \mathbb{Z}$  such that S = (d).

<sup>a</sup>in fact,  $d = \gcd(n_1, \ldots, n_k)$ .

**Proof.** Suppose that  $S = (n_1, \ldots, n_k)$ . If  $l \mid n_1, \ldots, l \mid n_k$ , then for all  $z \in S, l \mid z$ . Thus,  $d = \gcd(n_1, \ldots, n_k) \mid z$  for all  $z \in S$ , so  $S \subseteq d\mathbb{Z}$ .

It remains to show that  $d = \gcd(n_1, \ldots, n_k) \in S$ . Let r be the smallest positive element in S. Clearly,  $r\mathbb{Z} \subseteq S$ . It follows, by the Euclidean Algorithm that  $d \mid r$ , so d = r. Thus,  $d\mathbb{Z} \subseteq S$ .  $\square$ 

**Definition 2.3** ((Loose Definitions of) Groups, Rings, and Fields). A **group** is a set S with an operation +, 0, and inverses.

A ring is a set S with operations +,  $\times$  that supports addition, subtraction, and multiplication.

A field is a set S with operations  $+, \times$  that supports addition, subtraction, multiplication,

and divison<sup>a</sup>

<sup>a</sup>for all  $x \in S \setminus \{0\}$ ,  $\exists y \in S$  with xy = 1.

**Theorem 2.5.**  $(\mathbb{Z}/n\mathbb{Z}, +, \times)$  is a field if and only if n is prime.

**Proof.** For prime p and  $a \not\equiv 0 \pmod{p}$ ,  $a^{p-1} \equiv 1 \pmod{p}$  by Fermat's Little Theorem, so  $a^{p-2} = \overline{a}$ , the modular inverse of  $a \pmod{p}$ . Consequently, we have found an explicit inverse for each  $a \not\equiv 0 \pmod{p}$ .

On the other hand, suppose that n is not prime and let  $a \in \mathbb{Z}/n\mathbb{Z}$  with  $a \not\equiv 0 \pmod{n}$  and  $\gcd(a,n) > 1$ . Then ab + nm = 1 has no solutions, so  $ab \equiv 1 \pmod{n}$  has no solutions b.  $\square$ 

**Definition 2.4** (Units). Let u be an element of a ring R. u is a **unit** of R if and only if there exists  $v \in R$  such that uv = 1.

The units of  $\mathbb{Z}/n\mathbb{Z}$ , expressed as  $(\mathbb{Z}/n\mathbb{Z})^{\times}$ , is

$$(\mathbb{Z}/n\mathbb{Z})^{\times} = \{ a \pmod{n} \mid \exists b : ab = 1 \}.$$

**Exercise.** The Gaussian Integers  $\mathbb{Z}[i] = \{x + iy \mid x, y \in \mathbb{Z}\}$  form a ring.

Let  $R = \text{linear polynomials}/\mathbb{Z}$  is not a ring (product of linear polynomials is quadratic).

Related Observation: If R is a PID, then it is a UFD.

Going back to the Gaussian integers, it turns out that we can apply the Euclidean Algorithm using the norms of the Gaussian Integers as part of the division algorithm. On the other hand, note that to avoid the issue of gcd's being unit multiples of each other, we say that two Gaussian integers are coprime if the ideal generated by them is the entire ring.

**Lemma 2.1** (Division Algorithm for Gaussian Integers). For all  $n, m \neq 0$  in the Gaussian integers, there exists q, r in Gaussian integers such that

$$n = qm + r$$
 and  $0 < \mathcal{N}(r) < \mathcal{N}(m)$ 

where  $\mathcal{N}(x + iy) = |x + iy|^2 = x^2 + y^2$ .

**Proof.** Consider

$$\frac{n}{m} = z = \frac{n\overline{m}}{\mathcal{N}(m)}.$$

There exists a nearby close point q in  $\mathbb{Z}[i]$  such that

$$\left| q - \frac{n}{m} \right| \le \frac{\sqrt{2}}{2}$$

Multiplying both sides by |m|, it follows that

$$|mq - n| \le |m| \frac{\sqrt{2}}{2}$$

We define -r = mq - n, so that r = n - mq. It follows from the above equation that

$$\mathcal{N}(r) = |r|^2 \leq \frac{1}{2}|m|^2 = \frac{1}{2}\mathcal{N}(m) < \mathcal{N}(m).$$

Thus, the division algorithm holds for Gaussian integers, as desired.

We can also prove a number of simple division rules, using our knowledge of ideals and Linear Diophantine Equations...

#### 2.4 Division Rules

**Theorem 2.6.** If  $a \mid bc$  and gcd(a, b) = 1, then  $a \mid c$ .

**Proof.** The ideal  $(a, b) = \mathbb{Z}$ . Equivalently, the Euclidean Algorithm gives us a solution to ax + by = 1. Multiplying both sides of this equation by c, we get that

$$acx + bcy = c$$
.

Since  $a \mid bc$ ,  $a \mid bcy$ ; furthermore,  $a \mid acx$ . Thus, a divides the left hand side and so a must also divide the right hand side, giving  $a \mid c$ .

**Corollary 2.1.** If p is a prime and  $p \mid bc$ , then  $p \mid b$  or  $p \mid c$ .

**Proof.** This follows directly from Theorem 2.6, with a = p and a = c for the two cases.

**Definition 2.5** (Order). The **order of** p **in** a, ord<sub>p</sub>(a), is

$$\operatorname{ord}_{p}(a) = \max\{k \mid (p^{k} \mid a)\}.$$

**Corollary 2.2.** If p is a prime and  $a, b \in \mathbb{Z}$ , then

$$\operatorname{ord}_n(ab) = \operatorname{ord}_n(a) + \operatorname{ord}_n(b).$$

**Proof.** Let  $\operatorname{ord}_p(a) = \alpha$ , so that  $p = p^{\alpha}a_1$ , with  $p \nmid a_1$ . Similarly, let  $\operatorname{ord}_p(b) = \beta$ , so that  $p = p^{\beta}b_1$ , with  $p \nmid b_1$ .

Consider

$$ab = (p^{\alpha}a_1)\left(p^{\beta}b_1\right) = p^{\alpha+\beta}a_1b_1.$$

By the contrapositive of Corollary 2.1, since  $p \nmid a_1$  and  $p \nmid b_1$ , so  $p \nmid a_1b_1$ . Thus,

$$\operatorname{ord}_p(ab) = \alpha + \beta = \operatorname{ord}_p(a) + \operatorname{ord}_p(b).$$

**Theorem 2.7** (Fundamental Theorem of Arithmetic). If n is a nonzero integer, then there exists  $\varepsilon(n) = 0$  or -1 such that

$$n = (-1)^{\varepsilon(n)} \prod_{p} p^{\operatorname{ord}_{p}(n)}$$

and this decomposition is unique.

**Proof.** Assume for the sake of contradiction that n has a second factorization

$$n = (-1)^{\varepsilon(n)} \prod_{p} p^{e_p}.$$

Fix p, and apply ord<sub>p</sub> to both sides. We have that

$$\operatorname{ord}_{p}(n) = \operatorname{ord}_{p}\left((-1)^{\varepsilon(n)} \prod_{q} q^{e_{q}}\right)$$
$$= \operatorname{ord}_{p}\left((-1)^{\varepsilon(n)}\right) + \sum_{q} \operatorname{ord}_{p}(q^{e_{q}}).$$

Since  $\operatorname{ord}_p\left((-1)^{\varepsilon(n)}\right)=0$  and  $\operatorname{ord}_p(q^{e_q})=0$  for  $q\neq p$ , it follows that  $\operatorname{ord}_p(n)=p^{e_p}=e_p$ .

Thus, the second factorization must be the same as the first one, and so the decomposition is unique.  $\Box$ 

Note that the Fundamental Theorem of Arithmetic does not hold in the ring  $\mathbb{Z}[\sqrt{5}i]$ . For example, in this ring,

$$6 = 2 \cdot 3 = (1 + \sqrt{5}i)(1 - \sqrt{5}i).$$

## 2.5 Rings, Units, and Connections to Pythagorean Triples

So where does the above proof fail for this ring? Note that  $\operatorname{ord}_2(6) = 1$ ,  $\operatorname{ord}_{1+\sqrt{5}(i)} = \operatorname{ord}_{1-\sqrt{5}(i)} = 0$ , and so the additive property of orders for "relatively prime" numbers in the integers does not hold for the ring  $\mathbb{Z}[\sqrt{5}i]$ .

**Exercise.** The set of units of  $\mathbb{Z}[i]$  are  $\pm 1, \pm i$ . The set of units of  $\mathbb{Z}[\sqrt{5}i]$  are  $\pm 1$ .

Thus far, we've studied the similarities between viewing Pythagorean triples as rational points on a circle. Similarly, we studied the question of which numbers appear as the hypotenuses of triples.<sup>1</sup> Note that this is also equivalent to finding possible norms of Gaussian integers; for example, circles centered at (0,0) with radius  $\sqrt{n}$  for  $n \equiv 3 \pmod{4}$  pass through no lattice points. In general, we may also want to study which n arise as the values of a general binary quadratic form  $Ax^2 + Bxy + Cy^2$ .

<sup>&</sup>lt;sup>1</sup>numbers that are 3 (mod 4) cannot appear, due to the fact that squares are 0,1 (mod 4)

**Definition 2.6** (Irreducible). An element r in the ring R is irreducible if, for an element b in R,  $b \mid r$  implies that b is a unit or an associate to r.

**Definition 2.7** (Associates). Two elements r and s in the ring R are **associates** if there exists an element u in the set of non-units of R,  $R^{\times}$ , such that

r = us.

**Definition 2.8** (Prime). An element r in the ring R is **prime** if

 $r \mid s \cdot t \text{ implies } r \mid s \text{ or } r \mid t.$ 

**Remark.** In  $R = \mathbb{Z}[\sqrt{5}]$ , 2 is irreducible but is not prime.

**Definition 2.9** (Integral Domain). A ring R is an integral domain if

$$xy = 0$$
 implies  $x = 0$  or  $y = 0$ 

or equivalently, if R has no zero divisors.

**Lemma 2.2** (Prime implies irreducible in integral domains). A prime element is irreducible over an integral domain R.

**Proof.** Let p be a prime element in R. Suppose that  $s \mid p$ . We want to show that s is a unit or s is an associate of p.

Since  $s \mid p$ , there exists some  $k \in R$  such that sk = p. It follows that  $p \mid sk$ . Since p is a prime, we know by definition that  $p \mid s$  or  $p \mid k$ .

In the former case, it follows that s = lp and p = sk. Thus, p = sk = (lp)k, and so p(lk-1) = 0. Since  $p \neq 0$ , it follows that lk = 1, meaning that k is a unit and s is by definition an associate of p. The latter case follows similarly, but will yield the result that s is a unit of R.

Thus, prime elements are irreducible over integral domains, as desired.

**Lemma 2.3.** Let R be a PID. Then an element of R that is irreducible is also prime.

**Proof.** Let r be an element of R that is irreducible. Assume that  $r \mid ab$ , and  $r \nmid a$ . We want to show that  $r \mid b$ .

Consider the ideal (r, a). Since R is a PID, (r, a) = (d), where gcd(r, a) = d. Since  $d \mid r$  and r

is irreducible, we know that d is an associate of r or d is a unit of R.

If d is an associate of r, it follows that  $d = r \cdot u$  for a unit u. Since  $a \in (d)$ , we must have that  $d = r \cdot u \mid a$ , meaning that  $r \mid a$ . However, by assumption,  $r \nmid a$ , so d cannot be an associate of r.

It follows that d is a unit of R, so (r, a) = (d) = R. Note that  $(r \cdot b, a \cdot b) = (b)$ . Furthermore,  $(r \cdot b) \subset (r)$ , and  $(a \cdot b) \subset (r)$ . It follows that  $(r \cdot b, a \cdot b) \subset (r)$ . Thus,  $(b) \subset (r)$  and so  $r \mid b$ .  $\square$ 

**Exercise.**  $R^{\times}$  is a group.

**Remark.** If R is a PID, then an element is irreducible if and only if it is prime.

## 2.6 Quadratic Residues and Nonresidues

Recall that we studied the question of which numbers are represented as sums of squares. From our modular arithmetic section, we know that no prime  $p \equiv 3 \pmod 4$  can be expressed as a sum of squares.

**Lemma 2.4.** Every prime  $p \equiv 1 \pmod{4}$  is a sum of two squares.

**Lemma 2.5** (Quadratic Residues modulo p). Modulo any prime p, half of the numbers (excluding 0) are quadratic residues.

Equivalently, the number of numbers  $n \equiv \pmod{p}$  where  $n \neq 0$  and there exists some x such that  $x^2 \equiv n \pmod{p}$  is  $\frac{p-1}{2}$ .

**Proof.** Note that  $(x)^2 \equiv (-x)^2 \pmod{p}$ . It follows that there must be at most half the numbers among the squares.

Furthermore, note that if  $x^2 \equiv y^2 \pmod{p}$ , then  $p \mid x^2 - y^2$ , so  $p \mid (x - y)(x + y)$ . By Corollary 2.1, it follows that  $p \mid x - y$  or  $p \mid x + y$ , meaning  $x \equiv y \pmod{p}$  or  $x \equiv -y \pmod{p}$ . Thus, all squares in the first half are distinct and there are exactly  $\frac{p-1}{2}$  quadratic residues modulo p.

**Remark.**  $a^{\frac{p-1}{2}} = 1 \pmod{p} \iff a \text{ is a quadratic residue.}$ 

We can now find a procedure to write a prime  $p \equiv 1 \pmod{4}$  as a sum of two squares.

- Find quadratic non-residue a modulo p.
- Let  $z = a^{\frac{p-1}{4}} \pmod{p}$ . It follows that

$$z^2 = a^{\frac{p-1}{2}} \equiv -1 \pmod{p},$$

so  $z^2 + 1 \equiv 0 \pmod{p}$ , and  $z^2 + 1 = pk$  in the integers.

- In  $\mathbb{Z}[i]$ ,  $z^2 + 1 = (z+i)(z-i)$ , so (z+i)(z-i) = pk.
- To find the common factor, apply the Euclidean algorithm to z+i and p in  $\mathbb{Z}[i]$ .

Put more succinctly, this procedure has two steps:

- Find quadratic non-residue a modulo p.
- Let  $z = a^{\frac{p-1}{4}} \pmod{p}$ . Find  $\gcd(z+i,p)$  in  $\mathbb{Z}[i]$ .

**Example.** Let p=17. Then a=3 is a quadratic non-residue modulo p, and  $z=3^{\frac{17-1}{4}}\pmod{17}=13$ . Note that  $13^2+1^2=170=17\cdot 10$ .

We apply the Euclidean Algorithm to 13 + i and 17 in  $\mathbb{Z}[i]$ :

$$17 = (13+i) \cdot 1 + (4-i)13 + i = (4-i) \cdot (3+i) + 0$$

so gcd(13 + i, 17) = 4 - i.

Thus,  $17 = (4+i)(4-i) = 4^2 + 1^2$ .

**Example** (Super Large Example, all calculations computed using WolframAlpha). Take p=55,497,159,953. Let a=21,345 (randomly picked quadratic non-residue), so that  $a^{\frac{p-1}{2}}\equiv -1\pmod{p}$ .

Let  $z = a^{\frac{p-1}{4}} = 4,334,787,849$ . Since  $\gcd(z+i,p) = 235,532 + 4673i$ , we conclude that  $p = 235,532^2 + 4,673^2$ .

# 3 Norm Forms

# 3.1 Introduction to Eisenstein Integers

**Definition 3.1.** For  $\omega = e^{\frac{2\pi i}{3}}$ , the ring  $\mathbb{Z}[\omega]$  consists of the **Eisenstein Integers**:

$$\mathbb{Z}[\omega] = \{a_0 + a_1\omega \mid a_0, a_1 \in \mathbb{Z}\}.$$

Since  $\omega$  satisfies  $\omega^2 + \omega + 1 = 0$ ,  $\omega^2 = -\omega - 1$ , so  $\mathbb{Z}[\omega] = \mathbb{Z}[x]/(x^2 + x + 1)$ .

Remark. The Eisenstein integers form a triangular/"hexagonal" lattice in the complex plane.

**Exercise.** What is the "norm form" for  $\mathbb{Z}[\omega]$ ?

Note that

$$\mathcal{N}(a+b\omega) = (a+b\omega)(a+b\overline{\omega}) = a^2 - ab + b^2.$$

**Definition 3.2.** A ring R is a **Euclidean Domain** if there is a norm function  $\mathcal{N}: R \setminus \{0\} \mapsto \mathbb{N}$  such that for all n, m in R with  $m \neq 0$ , there exists q, r in R such that

$$n = mq + r$$

with  $\mathcal{N}(r) < \mathcal{N}(m)$ .

**Remark.** The ring  $R = \mathbb{R}[x]$  is a Euclidean domain, where the "norm function" is the degree of a polynomial.<sup>a</sup>

On the other hand,  $R = \mathbb{Z}[\sqrt{-5}]$  is not a Euclidean domain, since the division algorithm does not work:

if 
$$q = \lfloor \frac{n}{m} \rfloor$$
, then

$$|r|^2 = |n - mq|^2 = \left| \left( \frac{n}{m} - q \right) \right|^2 |m|^2.$$

However, note that  $\left|\left(\frac{n}{m}-q\right)\right|^2<\frac{3}{2}$ , and so we can't conclude that  $\mathcal{N}(r)<\mathcal{N}(m)$ .

**Theorem 3.1.** If R is a Euclidean domain, it is a principal ideal domain.

**Proof.** Let  $\mathcal{I}$  be an ideal of R. Either  $\mathcal{I} = (0)$ , which is principal, or  $\mathcal{I}$  contains some nonzero element a. Consider the latter case. Let

$$S = {\mathcal{N}(i) : i \in \mathcal{I} \setminus {0}} \subset \mathbb{N}.$$

Let s be the least element of S and let  $d \in \mathcal{I}$  satisfy  $\mathcal{N}(d) = s$ . By definition,  $(d) \subseteq \mathcal{I}$ . We want to show that, in fact,  $(d) = \mathcal{I}$ . Suppose that there is some nonzero element  $m \in \mathcal{I} \setminus (d)$ .

By the Division Algorithm, there must exist some q, r in R such that m = dq + r, and since  $m \notin (d)$ ,  $r \neq 0$  and  $\mathcal{N}(r) < \mathcal{N}(d)$ , which contradicts the fact that d has the smallest norm in S. Thus, R must be a principal ideal domain.

# 3.2 Reformulating with Ideals

Let's revisit some of our definitions for units, associates, and irreducible elements of a ring R in terms of the ideals.

#### Definition 3.3.

•  $a \mid b$  if and only if  $(b) \subseteq (a)$ .

<sup>&</sup>lt;sup>a</sup>we can define degree(0) =  $-\infty$ .

- a is an **associate** of b if and only if (a) = (b).
- u is a **unit of** R if and only if (u) = R.
- r is irreducible in R if  $r = s \cdot t$  for s, t in R implies s is a unit of (s) = (r).
- p is a **prime in** R if and only if for  $a \cdot b \in (p)$ , either  $a \in (p)$  or  $b \in (p)$ .
- a, b in R are **coprime** for R a PID if and only if (a, b) = R.
- d is a **gcd of** a, b **in** R if and only if  $(a) \subseteq (d)$  and  $(b) \subseteq (d)$  and if  $(a) \subseteq (d')$  and  $b \subseteq (d')$ , then  $(d') \subseteq (d)$ .

**Lemma 3.1.** If R is a PID, and a and b are two elements in R, then there exists a gcd d of a and b such that (a,b)=(d).

**Proof.** Note that there is some d such that (d) = (a, b). It remains to show that d is a gcd of a and b. Note that  $b \in (a, b) \subseteq (d)$ . Thus,  $b = d \cdot k$ , and  $d \mid b$ . By similar logic, we know  $d \mid a$ .

Suppose that  $d' \mid a$  and  $d' \mid b$ . Then  $a = d' \cdot l$  and  $b = d' \cdot k$ , so  $a \in (d')$  and  $b \in (d')$ . It follows that  $(d) = (a, b) \subseteq (d')$ , so  $d \in (d')$ . Thus,  $d = d' \cdot r$  and  $d' \mid d$ . Thus, d is the gcd of a and b as desired.

**Exercise.** Is  $x^2 + 1$  irreducible in  $R = (\mathbb{Z}/5)[x]$ ? How about in  $R = (\mathbb{Z}/7)[x]$ ?

**Answer.**  $x^2 + 1$  is not irreducible in  $(\mathbb{Z}/5)[x]$ , as

$$(x+2)(x+3) = x^2 + 5x + 6 \equiv x^2 + 1 \pmod{5}.$$

Note that this is equivalent to the question of whether there exists a solution x to  $x^2 \equiv -1 \pmod 5$ . In general, recall that  $x^2 + y^2 \equiv 0 \pmod p$  has a solution if and only if  $p \equiv 1 \pmod 4$  or p = 2.

Thus, since  $7 \equiv 3 \pmod{4}$ ,  $x^2 + 1$  is irreducible in  $R = (\mathbb{Z}/7)[x]$ .

**Definition 3.4** (Noetherian Ring). A ring R is **Noetherian**<sup>a</sup> if

$$\mathcal{I}_1 \subseteq \mathcal{I}_2 \subseteq \dots$$

is a chain of ideals, then the chain stops in finite time (there exists N such that for all n > N,  $I_N = I_n$ ).

**Remark.**  $\mathbb{Z}$  is Noetherian (which also follows from the theorem below, as  $\mathbb{Z}$  is a PID.)

<sup>&</sup>lt;sup>a</sup>named after Emily Noether

**Theorem 3.2.** If R is a PID, then it is Noetherian.

**Proof.** Let  $\mathcal{I}_1 \subseteq \mathcal{I}_2 \subseteq \dots$  be an increasing chain of ideals in R. Let

$$\mathcal{I} = \bigcup_{k=1}^{\infty} \mathcal{I}_k.$$

 $\mathcal{I}$  is an ideal, and since R is a PID, it follows that  $\mathcal{I} = (a)$  for some a. There must exist a N such that  $a \in \mathcal{I}_N$ , so  $(a) \subseteq \mathcal{I}_N$ . Since for any n > N,  $\mathcal{I}_n \subseteq \mathcal{I}_N$ , and they can't add any additional elements to  $\mathcal{I}_N$ , all  $\mathcal{I}_n = \mathcal{I}_N$ . It follows that R is Noetherian.

**Lemma 3.2.** If R is Noetherian, then any element r in R is a finite product of irreducibles.

**Proof.** If r is not irreducible, then there exists an element  $r_1$  in R such that  $r_1 \mid r$ ,  $r_1$  is not a unit,  $(r_1) \neq (r)$ , and  $(r) \subset (r_1)$ .

If  $r_1$  is itself not irreducible, then there exists some  $r_2$  such that  $r_2 \mid r_1, r_2$  is not a unit,  $(r_2) \neq (r_1)$ , and  $(r_1) \subset (r_2)$ .

Continuing this process, we arrive at an increasing chain of ideals:

$$(r) \subsetneq (r_1) \subsetneq (r_2) \dots$$

Since R is Noetherian, this process cannot go on forever. Suppose that it stops at  $r_j$ ; this means that  $l_1 = r_j$  is an irreducible factor of r.

We can continue this process on  $s_1$ , where  $r = l_1 \cdot s_1$ ; we will find that there exist some  $l_2$ , an irreducible factor of  $s_1$  and consequently, of r. Consider  $s_2$ , where  $s_1 = l_2 \cdot s_2$ , and so on and so forth. We once again find an increasing chain of ideals:

$$(r) \subsetneq (s_1) \subsetneq (s_2) \subsetneq \dots$$

Since R is Noetherian, this chain must stabilize, giving us an irreducible element  $s_J$ . It follows that

$$r = l_1 \cdots l_J \cdot s_J$$

giving us a finite product of irreducible elements.

**Lemma 3.3.** Let R be a Noetherian ring. Given a prime element p (equivalently, irreducible) in R and an element a in R, there exists some natural number n such that  $\operatorname{ord}_p(a) = n$ .

**Proof.** If  $p \mid a$ , then  $a = pa_1$ . If  $p \mid a_1$ , then  $a_1 = pa_2$ , and so on and so forth. We have an increasing chain of ideals

$$(a) \subseteq (a_1) \subseteq (a_2) \subseteq \dots$$

This process must stop, so there must be some n such that  $p^n \mid a$  and  $p^{n+1} \nmid a$ .

**Lemma 3.4.** If R is Noetherian and p is prime in R, then

$$\operatorname{ord}_p(a+b) = \operatorname{ord}_p(a) + \operatorname{ord}_p(b).$$

**Proof.** Same as for  $\mathbb{Z}$ .

**Theorem 3.3** (Noetherian implies UFD). Any element r in a Noetherian ring R is uniquely expressed as

$$r = u \prod_{p} p^{e_p}$$

where u is a unit, p is prime, and  $e_p = \operatorname{ord}_p(r)$ .

Remark (Important Takeaway). Note that Euclidean Domain implies PID which implies Noetherian ring which implies UFD.

Recall from earlier that  $\mathbb{Z}[\sqrt{-5}]$  is not Euclidean; this also follows from the fact that it is not a UFD:

$$6 = 2 \cdot 3 = (1 + \sqrt{-5})(1 - \sqrt{-5}).$$

## 3.3 A Return to Norm Forms, for Eisenstein Integers

Recall that earlier in the course, we wanted to know which n could be expressed as a sum of squares:  $n = x^2 + y^2$ . We answered this question by studying  $\mathbb{Z}[i]$ , since the quadratic form  $Q(x, y) = x^2 + y^2$  is the norm form for  $\mathbb{Z}[i]$ .

In the Eisenstein integers  $\mathbb{Z}[\omega]$ , the norm form is  $\mathcal{N}(x+\omega y)=x^2-xy+y^2$ .

**Exercise.** Which integers n can be expressed as  $n = x^2 - xy + y^2$ , for integers x and y?

**Answer.** We see  $n = Q(x, y) \in \{0, 1, 3, 4, 7, 9, 12, 13, ...\}$ , and no values that are  $\equiv 2 \pmod{3}$ ! Additionally, for every prime  $p \equiv 1 \pmod{3}$ , there exists x, y such that  $x^2 - xy + y^2 = p$  (i.e. all primes equivalent to  $1 \pmod{3}$  are represented – and the x, y can be determined efficiently.)

**Proof.** Let p be a prime equivalent to 1 (mod 3). Set  $z = a^{\frac{p-1}{3}} \pmod{p}$ , for  $z \not\equiv 1 \pmod{p}$  (there will be  $\frac{2(p-1)}{3}$  possible z).

It follows by Fermat's Little Theorem that  $z^3 - 1 \equiv 0 \pmod{p}$ , so  $z^3 - 1 = pk$  for some  $k \in \mathbb{Z}$ . Since  $z^3 - 1$  factors completely into  $(z - 1)(z - \omega)(z - \overline{\omega})$  in  $\mathbb{Z}[\omega]$ , we know that

$$(z-1)(z-\omega)(z-\overline{\omega}) = pk$$

in  $\mathbb{Z}[\omega]$ . Since  $z \not\equiv 1 \pmod{p}$ , we know that  $z-1 \mid k$ . Say that k = (z-1)l. Then

$$(z - \omega)(z - \overline{\omega}) = pl.$$

Note that  $p \nmid z - \omega$  and  $p \nmid z - \overline{\omega}$ . <sup>a</sup> Since  $p \nmid z - \omega$  and  $p \nmid z - \overline{\omega}$  but  $p \mid (z - \omega)(z + \overline{\omega})$ , p is not prime. Suppose that  $p = (x + y\omega)(x + y\overline{\omega}) = x^2 - xy + y^2$ . It follows that the component  $x + y\omega$  must divide at least one of  $z - \omega$  or  $z - \overline{\omega}$ .

Take the greatest common divisor, of p and  $z - \omega$  in  $\mathbb{Z}[\omega]$ ; if this is  $a + b\omega$ , it follows that

$$p = a^2 - ab + b^2$$

for integers a, b.

<sup>a</sup>if  $p \mid z - \omega$ , then  $z - \omega = (c + d\omega)p$ , which will only be satisfied if pd = -1. A similar line of reasoning can be used to show  $p \nmid z - \overline{\omega}$ .

In the above proof, we actually find an algorithm to solve  $p = \mathcal{N}(x + y\omega) = x^2 - xy + y^2$  for  $p \equiv 1 \pmod{3}$ .

- In  $\mathbb{Z}/p$ , find an a such that  $a^{\frac{p-1}{3}} \not\equiv 1$  and set  $z = a^{\frac{p-1}{3}} \pmod{p}$ .
- Find  $gcd_{\mathbb{Z}[\omega]}(p, z \omega) = x + y\omega$ . Then  $x^2 xy + y^2 = p$ .

**Example.** Find integers x, y for p = 3571 such that

$$3571 = x^2 - xy + y^2$$

using the Eisenstein integers.

**Answer.** We can check that p = 3571 is a prime equivalent to 1 (mod 3). For ease of reference, let's define  $l = \frac{p-1}{3} = 1190$ .

• We can find an a such that  $a^l \not\equiv 1 \pmod{p}$ : there are 2l such a, and one such a is a = 2, as

$$2^{1190}\pmod{3571}=3467\neq 1.$$

Now, let z = 3467. We know that  $z^3 - 1 \equiv 0 \pmod{p}$ .

• We now find  $gcd_{\mathbb{Z}[\omega]}(3571, 3467 - \omega)$ :

$$3571 = (3467 - \omega)(1) + (104 + \omega)$$
$$3467 - \omega = (104 + \omega)(33) + (35 - 34\omega)$$
$$104 + \omega = (35 - 34\omega)(2 + \omega) + 0.$$

Thus,  $gcd_{\mathbb{Z}[\omega]}(3571, 3467 - \omega) = 35 - 34\omega$  and so  $\mathcal{N}(35 - 34\omega) = 3571$  and

$$3571 = 35^2 + 34 \cdot 35 + 34^2$$

**Remark.** Furthermore, for general  $n = p_1 \cdots p_k$ , where each  $p_j$  for j from 1 to k is  $\equiv 1 \pmod 3$ , we can simply find a representation of each  $p_j$  and multiply the resulting Eisenstein integers to find the representation for n.

# 4 Binary Quadratic Forms

Thus far, we've studied  $x^2 + y^2 = n$  and  $x^2 - xy + y^2 = n$ , the norm forms of the Gaussian and Eisenstein integers. We know for what values of n there are integer solutions, as well as how to efficiently find prime solutions to these equations.

However, how might we generalize this to a general binary quadratic form

$$Q(x,y) = Ax^2 + Bxy + Cy^2.$$

For ease of reference, we will use the coefficients [A, B, C] to reference the above quadratic form.

**Remark.** Q = [1, 2, 1] is degenerate since its discriminant is 0.

Similarly, Q = [1, 3, 2] can be factored as (x-2y)(x-y). Using the change of variables z = x-y and w = -y, the quadratic form becomes Q(z, w) = z(z + y). Note that any value n can be achieved by setting z = 1 and y = n - 1.

In general, if the quadratic form factorizes, or has discriminant zero, the values of the quadratic form can be rather uninteresting (either every integer or only perfect squares, respectively, can be achieved).

# 4.1 Types of Quadratic Forms

**Definition 4.1** (Degenerate Quadratic Forms). If a quadratic form Q has discriminant zero, it is a **degenerate** quadratic form.

If a quadratic form Q has a discriminant equal to a perfect square, it splits and is reducible (as Q is a product of two linear polynomials).

**Definition 4.2** (Definite and Indefinite Quadratic Forms). A quadratic form Q is **definite** if its sign never changes.

A quadratic form Q is **indefinite** if its sign does change.

**Remark.** In general, the **discriminant**  $B^2 - 4AC$  of a quadratic form  $Ax^2 + Bxy + Cy^2$  tells us whether a quadratic form is definite or indefinite:

$$\begin{cases} \text{indefinite} & \text{if } B^2 - 4AC > 0; \\ \text{definite} & \text{if } B^2 - 4AC < 0; \\ \text{definite/degenerate quadratic form (only squares)} & \text{if } B^2 - 4AC = 0 \end{cases}$$

After our brief interlude on degeneracy, discriminants, definite and indefinite quadratic forms, let's revisit the question of what n are represented in the quadratic form  $Q(x,y) = Ax^2 + Bxy + Cy^2$ .

**Definition 4.3** (Primitive Quadratic Form). A quadratic form Q = [A, B, C] is **primitive** if gcd(A, B, C) = 1.

A similar yet more specific question may be the following: which n are primitively represented (meaning Q(x,y) = n where gcd(x,y) = 1) by Q?

**Answer.** For a specific quadratic Q, we can figure out all values Q takes on mod p by simply calculating  $p^2$  values and taking them mod p. In this strategy, we keep Q constant a; nd change values of x and y, calculating the respective modulos.

A dual approach is to fix x and y and change Q. To illustrate this idea, consider

$$Q_1(x,y) = x^2 + y^2$$

and replace x with x + y to create the new quadratic form

$$Q_2(x,y) = (x+y)^2 + y^2.$$

Note that  $Q_2$  represents the same solutions as  $Q_1$ , any values x = a, y = b in  $Q_1$  have identical values x = a - b, y = b in  $Q_2$ . The main takeaway here is that  $Q_1$  and  $Q_2$  differ by  $\mathbb{Z}$ -invertible linear change of variables.

We can work more generally using the group  $GL_2(\mathbb{Z})$ , the group of two-by-two invertible matrices (which must have determinant  $\pm 1$ ) in the integers.

Let Q = [A, B, C], and let  $\gamma = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$  be an element in  $GL_2(\mathbb{Z})$ . Then  $Q_1 = Q \circ \gamma$  represents the same n as Q. Note that

$$Q_1(x,y) = Q\left(\begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}\right)$$

$$= Q(ax + by, cx + dy)$$

$$= A(ax + by)^2 + B(ax + by)(cx + dy) + C(cx + dy)^2$$

$$= A_1x^2 + B_1xy + C_1y^2$$

where  $A_1 = Aa^2 + Bac + Cc^2 = Q(a,c), B_1 = A(2ab) + B(ad + bc) + C(2cd), and C_1 = Ab^2 + C(2cd)$ 

$$Bbd + Cd^2 = Q(b, d).$$

**Example.** Suppose we start with the quadratic form  $Q_0 = [7, -11, 5]$ . Let us change  $Q_0$  to a new quadratic form  $Q = Q_0 \circ \gamma$  and evaluate at (1, 0). Suppose we want

$$Q(1,0) = Q_0(1,1).$$

Then we need some matrix in  $GL_2(\mathbb{Z})$  such that

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \end{pmatrix}.$$

Solving gives us one such solution:  $\begin{pmatrix} a & b \\ c & d \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$ .

Consequently, we have that

$$Q = Q_0 \circ \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} = [7, -11, 5] \circ \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}.$$

From our above formula, we know that  $Q = [A_1, B_1, C_1]$ , where  $A_1 = Q(a, c) = Q(1, 1) = 1$ ,  $B_1 = 0 - 11 + 10 = -1$ , and  $C_1 = Q(b, d) = Q(0, 1) = 5$ .

Thus,  $Q(1,0) = Q_0(1,1)$ , for Q = [1, -1, 5].

We can formalize this duality process of fixing x and y and varying Q as follows.

**Definition 4.4** (Equivalent Quadratic Forms). The quadratic forms  $Q_1$  and  $Q_2$  are **equivalent**, denoted  $Q_1 \sim Q_2$ , if there exists some matrix  $\gamma \in SL_2(\mathbb{Z})$  (2 × 2 matrices with determinant 1) such that  $Q_1 = Q_2 \circ \gamma$ .

Given a form Q, the **equivalence class** of Q, denoted [Q], is the set of all forms equivalent to Q:

$$[Q] = \{Q' \mid Q' \sim Q\}.$$

**Lemma 4.1** (Equivalent Forms Represent Same Numbers). If  $Q_1$  and  $Q_2$  are two equivalent quadratic forms, then

 $\{n \in \mathbb{Z} \mid \text{ there exists } x, y \in \mathbb{Z}, Q_1(x, y) = n\} = \{n \in \mathbb{Z} \mid \text{ there exists } x, y \in \mathbb{Z}, Q_2(x, y) = n\}$ 

meaning that equivalent forms represent the same numbers.

**Proof.** If 
$$n = Q_1(x, y)$$
 then  $n = Q_2(x_1, y_1)$ , where  $\gamma \begin{pmatrix} x_1 \\ y_1 \end{pmatrix} = \begin{pmatrix} x \\ y \end{pmatrix}$ , so 
$$\begin{pmatrix} x_1 \\ y_1 \end{pmatrix} = \gamma^{-1} \begin{pmatrix} x \\ y \end{pmatrix}.$$

**Remark.** This lemma tells us that if we started with Q = [1, -1, 5] instead of  $Q_0 = [7, -11, 5]$  and wrote out the values for different x and y, we would get the same values but shuffled around.

**Theorem 4.1.** If two quadratic forms  $Q_1$  and  $Q_2$  are equivalent, their discriminants are equal.

One proof follows directly from algebra involving the coefficients of an equivalent quadratic form, described previously. Our proof will be more elegant.

**Definition 4.5** (Half-Hessian Matrix (Gram Matrix)). If Q = [A, B, C], its **Gram Matrix** (or **Half-Hessian Matrix**) is

$$M = \frac{1}{2} \begin{pmatrix} 2A & B \\ B & 2C \end{pmatrix} = \begin{pmatrix} A & \frac{B}{2} \\ \frac{B}{2} & C \end{pmatrix}$$

with det  $M = -\frac{1}{4}D_Q$ . Furthermore,

$$Q(x,y) = \begin{pmatrix} x & y \end{pmatrix} M \begin{pmatrix} x \\ y \end{pmatrix}.$$

**Remark.** If  $Q_2 = Q_1 \circ \gamma$  (or  $Q_1 \sim Q_2$ ) then

$$(x \quad y) M_2 \begin{pmatrix} x \\ y \end{pmatrix} = Q_2(x, y) = Q_1 \left( \left[ \gamma \begin{pmatrix} x \\ y \end{pmatrix} \right]^T \right)$$

$$= \left[ \gamma \begin{pmatrix} x \\ y \end{pmatrix} \right]^T M_1 \left[ \gamma \begin{pmatrix} x \\ y \end{pmatrix} \right].$$

As matrices,  $M_2 = \gamma^T M_1 \gamma$ .

**Proof (of Theorem 4.1).** Suppose that  $Q_1 \sim Q_2$ . From the above remark, we know that  $M_2 = \gamma^T M_1 \gamma$ , so

$$\det M_2 = \det(\gamma^T M_1 \gamma) = \det(\gamma^T) \det(M_1) \det(\gamma) = \det(M_1)$$

as  $\gamma \in SL_2(\mathbb{Z})$ . Furthermore, since det  $M_1 = -\frac{1}{4}D_{Q_1}$  and det  $M_2 = -\frac{1}{4}D_{Q_2}$ , it follows that  $D_{Q_1} = D_{Q_2}$ , as desired.

One lingering question that remains is to understand equivalences on quadratic forms  $Q_1$  and  $Q_2$  in terms of what they do to the roots  $\alpha_1$  and  $\alpha_2$  of  $Q_1(x,1)$  and  $Q_2(x,1)$ , respectively. For the

notes below, we will focus on the definite form case  $(D_{Q_1} = D_{Q_2} < 0)$ .

**Theorem 4.2.** If  $Q_2 = \gamma \circ Q_1$  (so  $Q_2 \sim Q_1$ , with  $\gamma \in SL_2(\mathbb{Z})$ ), then

$$\alpha_2 = \gamma^{-1} \circ \alpha_1.$$

Note: the  $\gamma^{-1}$  acts as a fractional linear transformation.

**Proof.** Let  $Q_1 = [A_1, B_1, C_1]$  and  $Q_2 = [A_2, B_2, C_2]$  (where each discriminant is negative) and let  $Q_2 \sim Q_1$ , so that

$$Q_2 = Q_1 \circ \begin{pmatrix} a & c \\ b & d \end{pmatrix}.$$

The roots are

$$\alpha_0 = \frac{-B_1 + \sqrt{|D|}i}{2A_1}$$
 and  $\alpha_2 = \frac{-B_2 + \sqrt{|D|}i}{2A_2}$ .

Recall that  $A_2 = Q_1(a, c) = A_1a^2 + B_1ac + C_1c^2$  and  $B_2 = A_1(2ab) + B_2(ad + bc) + C_2(2cd)$ . Substituting these results into the above expressions, we find that

$$\alpha_2 = \gamma^{-1} \circ \alpha_1. \qquad \Box$$

## 4.2 Reduction Theory of Quadratic Forms

With these tools under our belt, we now have a geometric reduction algorithm.

Let

$$T = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \text{ and } S = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}.$$

Note that  $T \circ \alpha = \alpha + 1$  and  $S \circ \alpha = -\frac{\overline{\alpha}}{|\alpha|^2}$  (S is an inversion about the unit circle: it brings points outside in and points inside out composed with a reflection about the x-axis.)

Our reduction algorithm works as follows:

- Apply T or  $T^{-1}$  to make  $|\operatorname{Re} \alpha| < \frac{1}{2}$ .
- Use S to get outside the unit circle.
- Continue above two steps until neither of them reduces  $\alpha$ , and we reach some  $\alpha' \in \{z \in \mathbb{H} \mid |\operatorname{Re} \alpha'| < \frac{1}{2}, |\alpha'| > 1\}$ . The region described above is known as the **fundamental domain**, which we denote as  $\mathcal{F}$ .

**Proof Sketch of Termination.** As we move z by S and T, we apply S if  $z_1 = \gamma \circ z$  is in

$$\{w\in\mathbb{H}\mid |\mathrm{Re}\,w|<\frac{1}{2},\, |w|\leq 1,\, \mathrm{Im}w\geq \mathrm{Im}z\},$$

a compact region. If this process continues forever, we would get a limit point. However, the

matrix  $\gamma$  has integer entries – the action is discrete and so there cannot be any limit points.

Thus, our process halts, and does so once we are in the fundamental domain.

**Example.** Let  $\alpha = \frac{10}{3} + \frac{1}{2}i$ .

- $\alpha_1 = T^{-3} \circ \alpha = \frac{1}{3} + \frac{1}{2}i$   $\alpha_2 = S \circ \alpha_1 = \frac{-12}{13} + \frac{18}{13}i$   $\alpha_3 = T \circ \alpha_2 = \frac{1}{13} + \frac{18}{13}i$

To apply this idea to reducing quadratic forms, we will reduce a quadratic form Q by reducing

**Example.** Begin with Q = [7, -11, 5], which has  $D_Q = -19$  and  $\alpha_Q = \frac{11}{14} + \frac{\sqrt{19}i}{14}$ .

- $\alpha_1 = T^{-1} \circ \alpha = \frac{11}{14} + \frac{\sqrt{19}i}{14}$ .  $\alpha_2 = S \circ \alpha_1 = \frac{42}{19} + \frac{14\sqrt{19}i}{19}$ .
- $\alpha_3 = T^{-2} \circ \alpha_2 = \frac{4}{19} + \frac{14\sqrt{19}i}{19}$

In summary, we have that

$$\alpha_3 = (T^{-2}ST^{-1}) \circ \alpha = \begin{pmatrix} 2 & -1 \\ -1 & 1 \end{pmatrix} \circ \alpha$$

Consequently, our reduced quadratic form is

$$Q_2 = Q_1 \circ \left( \begin{pmatrix} 2 & -1 \\ -1 & 1 \end{pmatrix} \right)^{-1}$$
$$= Q_1 \circ \left( \begin{pmatrix} 1 & 1 \\ 1 & 2 \end{pmatrix} \right)^{-1}$$
$$= [1, 1, 5].$$

What makes our resulting form Q with  $\alpha_Q \in \mathcal{F}$  reduced?

First, note that for  $\alpha_Q = \frac{-B + \sqrt{|D|}i}{2A}$  to be in the fundamental domain, we must have that

$$|\operatorname{Re} \alpha| = \left| \frac{-B}{2A} \right| < \frac{1}{2}.$$

This occurs if and only if  $|B| \leq |A|$ . Assume that Q is positive definite. Since Q(1,0) = A > 0, then this condition is equivalent to the condition |B| < A.

The other condition for  $\alpha_Q$  to be in the fundamental region is for the norm to be greater than or

equal to 1:

$$1 \le |\alpha_Q|^2 = \frac{1}{4A^2} [B^2 + |D|]$$

$$= \frac{1}{4A^2} [B^2 - |D|]$$

$$= \frac{1}{4A^2} [B^2 - (B^2 - 4AC)]$$

$$= \frac{C}{A}.$$

Simplifying, we get the condition that  $A \leq C$ .

**Lemma 4.2** (Reduced Quadratic Form). A quadratic form Q = [A, B, C] is **reduced** if

$$|B| \le A \le C$$
.

**Exercise.** Note that  $Q(x,y) = x^2 + y^2 = [1,0,1]$  has discriminant -4. Are there any other reduced quadratic forms with discriminant -4?

**Answer.** No. Note that  $D = B^2 - 4AC = -4$  implies that  $4AC - B^2 = 4$ . Equivalently,  $4AC = 4 + B^2$ . Since  $A \le C$  and  $|B| \le A$ , we also have that

$$4A^2 < 4AC = 4 + B^2 < 4 + A^2$$
.

It follows that  $4A^2 \le 4 + A^2$ , so  $3A^2 \le 4$ , implying that A = 0 or 1. Furthermore, since  $4AC = 4 + B^2$ , we know  $B \equiv 0 \pmod{2}$ .

The only reduced quadratic form that satisfies these conditions is Q = [1, 0, 1]. Thus, if Q = [A, B, C] and  $D_Q = -4$ , then  $Q \sim Q_0 = [1, 0, 1]$ .

**Exercise.** What about for the norm form of the Eisenstein integers:  $Q(x,y) = x^2 - xy + y^2$ ? Are there other reduced forms that are not equivalent to Q?

**Answer.** Following a similar approach as the previous example, we find two solutions: Q = [1, -1, 1] and  $Q_1 = [1, 1, 1]$ : however, the roots of Q and  $Q_1$  are  $1 + \omega$  and  $\omega$ , respectively. These solutions are equivalent (under the transformation T), but they both lie on the *boundary* of the fundamental domain.

**Lemma 4.3.** If Q = [A, B, C] and  $D_Q = -3$ , then  $Q \sim [1, -1, 1]$ .

**Remark.** This does not hold for every quadratic form though. Recall the norm form of  $\mathbb{Z}[\sqrt{-5}]$  which is not a PID:  $Q = x^2 + 5y^2 = [1,0,5]$ . It can be shown that [2,2,3] and [1,0,5] are two valid quadratic forms (roots  $\frac{-1}{2} + \frac{\sqrt{5}i}{2}$  and  $\sqrt{-5}i$ ), each of which have discriminant -20.

**Lemma 4.4.** If Q = [A, B, C] and  $D_Q = -20$ , then  $Q \sim [1, 0, 5]$  or  $Q \sim [2, 2, 3]$ .

#### 4.3 Finiteness of Class Number

**Theorem 4.3** (Gauss). Let Q be a positive definite form of discriminant D < 0. Then there exists finitely many reduced forms of discriminant D such that Q is equivalent to one of them.

Equivalently, the number of reduced forms of a given negative discriminant is finite (finiteness of class number).

**Proof.** If Q has discriminant  $D = B^2 - 4AC = -|D|$  and is reduced,  $|B| \le A \le C$ .

Consequently,

$$4A^2 \le 4AC = B^2 + |D| \le A^2 + |D|$$

so 
$$3A^2 \le |D|$$
 and  $A \le \sqrt{\frac{|D|}{3}}$ .

Thus, the number of A's is finite, and since  $|B| \leq A$ , the number of B's is finite. A, B, and D uniquely determine C, so the number of reduced forms of discriminant D is finite.

# 5 Summary

Topics covered, in rough order, include:

- Pythagorean Triples
- Primitive integer points:

$$\begin{cases} x = u^2 - v^2 \\ y = 2uv \\ z = u^2 + v^2 \end{cases}$$

with extension to rational points on  $S^1$ .

 $\bullet$  Discussion of which z are hypotenuses (sums of two squares). This led to the following theorem

**Theorem 5.1** (Sum of Squares). A prime p is a sum of squares:  $p = u^2 + v^2$  if and only if p = 2 or  $p \equiv 1 \pmod{4}$ .

A number n is a sum of squares:  $n = u^2 + v^2$  if and only if for every prime  $p \equiv 3 \pmod{4}$ , ord<sub>p</sub>(n) is even.

which came from  $\mathbb{Z}[i]$ , a UFD, and we had an algorithm to find a given u, v for  $p \equiv 1 \pmod{4}$ .

• We expanded this discussion to other norm forms, like the norm form of the Eisenstein integers:  $u^2 - uv + v^2$ .

**Theorem 5.2** (Primes in Eisenstein Norm Form). A prime p can be expressed in Eisenstein norm form  $p = u^2 - uv + v^2$  if and only if p = 3 or  $p \equiv 1 \pmod 3$ .

Once again, the proof came from a study of  $\mathbb{Z}[\omega]$ , and we developed an algorithm to find u, v.

- More generally, we discussed quadratic forms: n = Q(x, y).
- Equivalence of Quadratic Forms
- Reduction Theory of (Definite) Quadratic Forms
- Finiteness of Class Number