

# ROTH'S THEOREM ON 3-TERM ARITHMETIC PROGRESSIONS

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# Outline

1. The Question: 3-Term APs
2. Attempts and Intuition
3. Proof Sketch (The 5 Ideas)
4. The 5 Ideas, Explained
5. Beyond Roth's Theorem

## Notation

- $\{1, 2, \dots, N\}$  and  $[N]$  are used interchangeably.
- $C$  and  $c$  refer to generic large and small constants respectively. In different places, they may refer to different constants. For example,  $(cx^2 - Cx - C)/(Cx + C) \geq cx$  for  $x \geq C$ .
- Some equations, although not explicitly stated, only apply to large  $N$ .
- Blue items are additional technical pointers.

## The Question: 3-Term APs

$$0 < \alpha \leq 1$$

$A \subseteq \{1, 2, \dots, N\}$ ,  $|A| = \alpha N$ , s.t.  $A$  has no 3-term APs. How big can  $|A|$  be?

$$\alpha_{\max}(N) := \max_{A \subseteq [N]} \{\alpha\} \tag{1}$$

When  $N \rightarrow +\infty$ , does  $\alpha_{\max} \rightarrow 0$ ?

# Intuition

Arithmetic progressions have rich structure:

- AP-indexed subsequences of APs are also APs.
- The common difference may be viewed as the "period" of occurrence.
- $(a, b, c)$  is a 3-term AP iff  $a + c = 2b$ .

$\{1, 2, \dots, N\}$  may be replaced with any  $N$ -term AP.

## Intuition

Suppose the structure of interest was instead  $(a, b, c)$  where  $a + b = c$ .  
Then  $\alpha_{\max} \approx \frac{1}{2}$  for all  $N$ . (Exercise)

This tells us that:

- The rich structure of APs and 3-term APs is **very important**.
- Techniques related to number theory might be required.

## Intuition

Considering consecutive triples,  $\alpha \leq \frac{2}{3} + \frac{2}{N}$ . This settles  $N = 1, 2, 3, 4, 5, 6$ .

For  $N = 7$ , consider the middle element.  $|A|_{\max} = 4$ .

For  $N = 8$  and  $N = 9$ ,  $|A|_{\max} = 4$  and  $|A|_{\max} = 5$  respectively.

$3 \times 3$  grid pattern for  $N = 9$  may improve the  $\frac{2}{3}$  bound (consider  $N = 27$ ); possible approach to consider a  $3 \times 3 \times \cdots \times 3$  lattice. (Hales-Jewett)

# SPOILER ALERT

If you wish to try the problem yourself, stop here.



## The Answer: Roth's Theorem

(Roth, 1953) We may upper-bound  $\alpha_{\max}$  in terms of  $N$  as follows:

$$\alpha_{\max} \leq \frac{C^{568000}}{(\log \log N)^{1/5}} \quad (2)$$

In particular,  $\alpha_{\max} \rightarrow 0$  when  $N \rightarrow +\infty$ .

## Proof Sketch

Note: this is not Roth's original proof, but it is chosen for its elegance of integrating key ideas.

Density-Increment:  $\alpha' = \alpha + c\alpha^6$

Indicator  
Functions

$$f := \mathbf{1}_A - \alpha \mathbf{1}_{[N]}$$
$$AP_3 := \mathbb{E}_{x,d} g_1(x) g_2(x + d) g_3(x + 2d)$$

Cauchy-Schwarz  
and  $U^2$ -norm

$$|AP_3(g_1, g_2, g_3)| \leq \|g_i\|_{U^2}$$
$$\|f\|_{U^2} \geq c\alpha^3$$

Discrete Fourier  
Analysis

$$\exists \theta \in [0, 1] \text{ s.t.}$$
$$|\sum f(x) e(\theta x)| \geq c\alpha^6 N$$

Hardy-Littlewood  
Circle Method

$$\exists P' \subseteq [N], |P'| \geq N^{1/3},$$
$$\text{s.t.}$$
$$\sum_{x \in P'} f(x) \geq c\alpha^6 |P'|$$



$$\alpha_{\max} \leq \frac{c}{(\log \log N)^{1/5}}$$

# Idea 1: Density Increment Strategy

**Density-Increment:**  $\alpha' = \alpha + c\alpha^6$

Indicator  
Functions

$$f := \mathbf{1}_A - \alpha \mathbf{1}_{[N]}$$

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$$\alpha_{\max} \leq \frac{c}{(\log \log N)^{1/5}}$$

## Idea 1: Density-Increment Strategy

Suppose  $A \subseteq [N]$ ,  $|A| = \alpha N$ , and  $A$  has no 3-term APs. If  $N \geq C \alpha^{-\frac{3E19}{18}}$ , then:

There is a sub-AP  $P' \subseteq [N]$  of length  $|P'| \geq N^{1/3}$  with  $|A \cap P'| = \alpha' |P'|$  s.t.

$$\alpha' \geq \alpha + c \alpha^{\frac{1}{5E5}} \quad (3)$$

## Idea 1: Density-Increment Strategy

Suppose  $A \subseteq [N]$ ,  $|A| = \alpha N$ , and  $A$  has no 3-term APs. If  $N \geq C \alpha^{-3E19 \cdot 18}$ , then:

There is a sub-AP  $P' \subseteq [N]$  of length  $|P'| \geq N^{1/3}$  with  $|A \cap P'| = \alpha' |P'|$  s.t.

$$\alpha' \geq \alpha + c \alpha^6 \quad (4)$$

$1/5E5$

## Proof: Density-Increment Strategy $\Rightarrow$ (2)

Suppose the Density-Increment statement is true.

3E19 18

Let  $N_0 = N$ ,  $P_0 = [N]$ ,  $A_0 = A$ , and  $\alpha_0 = \alpha$ . Now for each  $i \in \mathbb{N}$ , if  $N_i \geq C\alpha_i^{-C}$ :

- There is some sub-AP  $P_{i+1} \subseteq P_i$  of length  $N_{i+1} \geq N_i^{1/3}$  such that,
- Setting  $A_{i+1} = A_{i+1} \cap P_{i+1}$  and  $|A_{i+1}| = \alpha_{i+1}N_{i+1}$ ,
- One has  $\alpha_{i+1} \geq \alpha_i + c\alpha_i^{1/5E56}$ .

## Proof: Density-Increment Strategy $\Rightarrow$ (2)

The  $\alpha_i$ 's are increasing rapidly:

- Starting from  $\alpha_0 = \alpha$ , it would take  $\leq \frac{1}{c\alpha^5} = C\alpha^{-5}$  iterations for it to double to  $2\alpha$ .  
 $5E5$   
 $1/5E5$
- The next double to  $4\alpha$  would be reached in  $\leq C(2\alpha)^{-5}$  more iterations.  
 $5E5$
- The  $k$ -th double to  $2^k\alpha$  would require  $\leq C\alpha^{-5}2^{-5k}$  iterations.  
 $5E5$

However, the  $\alpha_i$ 's are also bounded above by 1.

## Proof: Density-Increment Strategy $\Rightarrow$ (2)

Therefore, the total number of iterations is at most

$$i_{\max} = C\alpha^{-5}(1 + 2^{-5} + 2^{-10} + \dots) \leq C\alpha^{-5}.$$

5E5

5.17E5

Since the iterations must stop, therefore  $N_{i_{\max}} \leq C\alpha_{i_{\max}}^{-C} \leq C\alpha^{-C}$ .

3E19

18

Since  $N_{i_{\max}} \geq N^{(1/3)^{i_{\max}}} \geq N^{(1/3)^{C\alpha^{-5}}}$ , thus  $N^{(1/3)^{C\alpha^{-5}}} \leq C\alpha^{-C}$ .

5.17E5

3E19 18

5.17E5

45 18

Taking double-log,  $\log \log N - (\log 3)C\alpha^{-5} \leq \log(C - C \log \alpha)$ .



## Proof: Density-Increment Strategy $\Rightarrow$ (2)

Note that  $\log(C - C \log \alpha) \leq C\alpha^{-5}$ .

Therefore  $\log \log N \leq C\alpha^{-5}$ .

Rearranging,  $\alpha \leq \frac{C}{(\log \log N)^{1/5}} \cdot \square$

## Idea 2: Indicator Functions

**Density-Increment:**  $\alpha' = \alpha + c\alpha^6$

**Indicator Functions**

$$f := 1_A - \alpha 1_{[N]}$$

$$AP_3 := \mathbb{E}_{x,d} g_1(x) g_2(x+d) g_3(x+2d)$$

**Cauchy-Schwarz and  $U^2$ -norm**

$$|AP_3(g_1, g_2, g_3)| \leq \|g_i\|_{U^2}$$

$$\|f\|_{U^2} \geq c\alpha^3$$

**Discrete Fourier Analysis**

$$\exists \theta \in [0, 1] \text{ s.t.}$$

$$|\sum f(x) e(\theta x)| \geq c\alpha^6 N$$

**Hardy-Littlewood Circle Method**

$$\exists P' \subseteq [N], |P'| \geq N^{1/3},$$

$$\text{s.t.}$$

$$\sum_{x \in P'} f(x) \geq c\alpha^6 |P'|$$



$$\alpha_{\max} \leq \frac{c}{(\log \log N)^{1/5}}$$

## Idea 2: Indicator Functions

We want a function to count the number of 3-term APs in  $A$ .

Define indicator  $1_A$  as  $1_A(x) := \{1 \text{ if } x \in A, 0 \text{ if } x \notin A\}$ .  $N'$  is chosen this way to prevent counting of wrap-around APs.

Consider the finite abelian group  $G = \mathbb{Z}/N'\mathbb{Z}$  where  $N' = 2N + 1$ .

$|A|$  arises from the trivial  $d = 0$  solutions.

Then  $\mathbb{E}_{x,d \in G} 1_A(x) 1_A(x+d) 1_A(x+2d) = (|A| + 2 \cdot \# \text{ 3-term APs in } A)/N'^2$ .

Each 3-term AP is double-counted by  $\pm d$ .

## Idea 2: Indicator Functions

We therefore consider the linear indicator function

$$AP_3(g_1, g_2, g_3) := \mathbb{E}_{x, d \in G} g_1(x) g_2(x + d) g_3(x + 2d).$$

This term,  $\approx \alpha/4N$ , is small.

Since  $A$  has no 3-term APs, thus  $AP_3(1_A, 1_A, 1_A) = |A|/N^2 = \alpha N/(2N + 1)^2$ .

We compare  $A$  to a random subset of  $[N]$  of size  $\alpha N$ .

$$AP_3(\alpha 1_{[N]}, \alpha 1_{[N]}, \alpha 1_{[N]}) = \alpha^3 AP_3(1_{[N]}, 1_{[N]}, 1_{[N]}).$$

## Idea 2: Indicator Functions

Note that  $[N]$  has  $\geq \lfloor N/3 \rfloor^2$  3-term APs formed by  $x, d = 1, 2, \dots, \lfloor N/3 \rfloor$ .

$$\text{Thus } AP_3(1_{[N]}, 1_{[N]}, 1_{[N]}) \geq (N + 2 \overset{\approx 1/18}{\lfloor N/3 \rfloor^2}) / (2N + 1)^2 \overset{1/19}{\geq} \overset{50}{c} \text{ for } N \geq C.$$

$$\text{So for } N \geq C, AP_3(\overset{50}{\alpha} 1_{[N]}, \overset{1/19}{\alpha} 1_{[N]}, \overset{1/19}{\alpha} 1_{[N]}) \geq c \alpha^3.$$

## Idea 2: Indicator Functions

1/19

We now compare  $AP_3(1_A, 1_A, 1_A) \leq \frac{\alpha}{4N}$  to  $AP_3(\alpha 1_{[N]}, \alpha 1_{[N]}, \alpha 1_{[N]}) \geq c\alpha^3$ .

Define function  $f := 1_A - \alpha 1_{[N]}$ .  $f$  is *balanced* because  $\mathbb{E}_{x \in G} f(x) = 0$ .

Expanding linearly,  $AP_3(1_A, 1_A, 1_A) = AP_3(\alpha 1_{[N]}, \alpha 1_{[N]}, \alpha 1_{[N]})$   
 $+ AP_3(f, \alpha 1_{[N]}, \alpha 1_{[N]}) + AP_3(\alpha 1_{[N]}, f, \alpha 1_{[N]}) + AP_3(\alpha 1_{[N]}, \alpha 1_{[N]}, f)$   
 $+ AP_3(f, f, \alpha 1_{[N]}) + AP_3(f, \alpha 1_{[N]}, f) + AP_3(f, f, \alpha 1_{[N]}) + AP_3(f, f, f)$ .

Since  $f$  is *balanced*, the 7 terms of the form  $AP_3(f, *, *)$ ,  $AP_3(*, f, *)$ , and  $AP_3(*, *, f)$  should be small; however, the equation above says otherwise.

## Idea 2: Indicator Functions

Explicitly, we have that  $|AP_3(g_1, g_2, g_3)| \geq \frac{1}{7}(c\alpha^3 - \frac{\alpha}{4N})$  for 1-bounded functions  $g_1, g_2, g_3$  among which  $g_i = f$  for some  $i \in \{1, 2, 3\}$ .

For  $N \geq C\alpha^{-\frac{2}{7}}$ ,  $\frac{1}{7}(c\alpha^3 - \frac{\alpha}{4N}) \geq \frac{1}{14}c\alpha^3$ .

To make sense of the  $AP_3$  indicator function, we will relate  $|AP_3(g_1, g_2, g_3)|$  to the "norm"s of  $g_1, g_2$ , and  $g_3$ . We hope that this will give us a lower bound for the "norm" of  $f$ , which in turn gives us information about  $A$ .

## Idea 3: Cauchy-Schwarz Inequality

Density-Increment:  $\alpha' = \alpha + c\alpha^6$

Indicator  
Functions

$$f := 1_A - \alpha 1_{[N]}$$

$$AP_3 := \mathbb{E}_{x,d} g_1(x) g_2(x + d) g_3(x + 2d)$$

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$$\exists \theta \in [0, 1] \text{ s.t.}$$

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$$\alpha_{\max} \leq \frac{c}{(\log \log N)^{1/5}}$$



## Idea 3: Cauchy-Schwarz Inequality

To bound  $AP_3$ , we will attempt to use the Cauchy-Schwarz Inequality. We first rephrase the statement with the use of expected value: for any functions  $\beta, \gamma : X \rightarrow \mathbb{C}$ ,

$$|(\mathbb{E}_{x \in X} \beta(x) \overline{\gamma(x)})|^2 \leq \mathbb{E}_{x \in X} |\beta(x)|^2 \mathbb{E}_{x' \in X} |\gamma(x')|^2 \quad (5)$$

Note that by setting  $\gamma := \overline{\gamma}$ , we may remove the conjugation on the left.

If  $\beta$  is 1-bounded, then  $|\mathbb{E}_{x \in X} \beta(x) \gamma(x)|^2 \leq \mathbb{E}_{x \in X} |\gamma(x)|^2$ .

## Idea 3: Cauchy-Schwarz Inequality

We now apply Cauchy-Schwarz to bound  $AP_3(g_1, g_2, g_3)$  in terms of  $g_1$  alone; to do this, we will use the fact that  $g_2$  and  $g_3$  are 1-bounded.

Note that  $AP_3(g_1, g_2, g_3) = \mathbb{E}_{x,y \in G} g_1(2x - y)g_2(x)g_3(y)$ . There is some freeplay in parameterising 3-term APs.

Since  $g_2$  is 1-bounded, we let  $\gamma(x) := \mathbb{E}_{y \in G} g_1(2x - y)g_3(y)$ . Then,

$$|AP_3(g_1, g_2, g_3)|^2 \leq \mathbb{E}_{x \in G} |\mathbb{E}_{y \in G} g_1(2x - y)g_3(y)|^2 \quad (6)$$

which we may rewrite as  $\mathbb{E}_{x \in G} \mathbb{E}_{y \in G} g_1(2x - y)g_3(y) \overline{\mathbb{E}_{y' \in G} g_1(2x - y')g_3(y')}$ .

## Idea 3: Cauchy-Schwarz Inequality

We apply Cauchy-Schwarz once again:  $g_3(y)\overline{g_3(y')}$  is 1-bounded, thus

$$|\mathbb{E}_{y,y',x \in G} g_1(2x-y)\overline{g_1(2x-y')}g_3(y)\overline{g_3(y')}|^2 \leq \mathbb{E}_{y,y' \in G} |\mathbb{E}_{x \in G} g_1(2x-y)\overline{g_1(2x-y')}|^2$$

which we rewrite as  $\mathbb{E}_{x,x',y,y' \in G} g_1(2x-y)\overline{g_1(2x-y')}g_1(2x'-y)\overline{g_1(2x'-y')}$ ,

and reparameterise as  $\mathbb{E}_{x,d_1,d_2 \in G} g_1(x)\overline{g_1(x+d_1)}g_1(x+d_2)\overline{g_1(x+d_1+d_2)}$ .

This is solely in terms of  $g_1$ , and is non-negative! "norm"-worthy.

## Idea 3: Cauchy-Schwarz Inequality

Define the Gowers  $U^2$ -norm on function  $F : G \rightarrow \mathbb{C}$  as

$$\|F\|_{U^2} := \sqrt[4]{\mathbb{E}_{x,d_1,d_2 \in G} F(x) \overline{F(x+d_1)} F(x+d_2) \overline{F(x+d_1+d_2)}} \quad (7)$$

As its norm-related properties will not be used, we postpone the proof that it is in fact a norm to the end of the slides.

The two applications of Cauchy-Schwarz show  $|AP_3(g_1, g_2, g_3)|^4 \leq \|g_1\|_{U^2}^4$ .

Thus  $|AP_3(g_1, g_2, g_3)| \leq \|g_1\|_{U^2}$ .

## Idea 3: Cauchy-Schwarz Inequality

Similarly, we may also show  $|AP_3(g_1, g_2, g_3)| \leq \|g_2\|_{U^2}, \|g_3\|_{U^2}$ .

We know that  $|AP_3(g_1, g_2, g_3)| \geq c\alpha^3$  for  $N \geq C\alpha^{-C}$  and 1-bounded functions  $g_1, g_2, g_3$  among which  $g_i = f$  for some  $i \in \{1, 2, 3\}$ .

Thus for  $N \geq C\alpha^{-C}$ ,  $\|f\|_{U^2} \geq c\alpha^3$ .

We will now make sense of  $\|f\|_{U^2}$  using Discrete Fourier Analysis.

# Idea 4: Discrete Fourier Analysis

Density-Increment:  $\alpha' = \alpha + c\alpha^6$

Indicator  
Functions

$$f := 1_A - \alpha 1_{[N]}$$

$$AP_3 := \mathbb{E}_{x,d} g_1(x) g_2(x + d) g_3(x + 2d)$$

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$$\exists P' \subseteq [N], |P'| \geq N^{1/3},$$

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$$\alpha_{\max} \leq \frac{c}{(\log \log N)^{1/5}}$$

## Idea 4: Discrete Fourier Analysis

A trained eye would recognise that the  $U^2$ -norm can be treated with Fourier Analysis. We first establish the basic tools required.

For function  $f : G = \mathbb{Z}/N'\mathbb{Z} \rightarrow \mathbb{C}$ , define its transform  $\hat{f} : \mathbb{Z}/N'\mathbb{Z} \rightarrow \mathbb{C}$  as

$$\hat{f}(r) := \mathbb{E}_{x \in G} f(x) \exp\left(-2\pi i \frac{rx}{N'}\right) \quad (8)$$

To simplify notation, we define  $e(x) := \exp(2\pi i x)$ , which has period 1. Then

$$\hat{f}(r) := \mathbb{E}_{x \in G} f(x) e\left(-\frac{rx}{N'}\right) \quad (9)$$

## Idea 4: Discrete Fourier Analysis

The elegance of the Fourier transform shows in the inverse transform:

$$f(x) = \sum_{r \in G} \hat{f}(r) e\left(\frac{rx}{N'}\right) \quad (10)$$

To see why this is true:

$$\sum_{r \in G} \hat{f}(r) e\left(\frac{rx}{N'}\right) = \sum_{r \in G} \mathbb{E}_{y \in G} f(y) e\left(-\frac{ry}{N'}\right) e\left(\frac{rx}{N'}\right) = \sum_{y \in G} f(y) \mathbb{E}_{r \in G} e\left(\frac{r(x-y)}{N'}\right)$$

Neat trick to swap  $\sum$  with  $\mathbb{E}$  over the variables  $y$  and  $r$ .



## Idea 4: Discrete Fourier Analysis

The critical observation is

$$\mathbb{E}_{r \in G} e\left(\frac{r(x-y)}{N'}\right) = \{1 \text{ if } x = y, 0 \text{ if } x \neq y\} = \delta_{xy}$$

which we immediately apply to achieve

$$\sum_{r \in G} \hat{f}(r) e\left(\frac{rx}{N'}\right) = \sum_{y \in G} f(y) \mathbb{E}_{r \in G} e\left(\frac{r(x-y)}{N'}\right) = \sum_{y \in G} f(y) \delta_{xy} = f(x)$$

as desired.

The fact that the inverse transform involves a summation ( $\sum$ ) instead of an expected value ( $\mathbb{E}$ ) is important.

## Idea 4: Discrete Fourier Analysis

Here are other facts required: Note that  $\mathbb{E}$  is applied to the original function while  $\sum$  is applied to the transform.

1. Define the norms  $\|f\|_2 = (\mathbb{E}_{x \in G} |f(x)|^2)^{1/2}$  and  $\|\hat{f}\|_2 = (\sum_{r \in G} |\hat{f}(r)|^2)^{1/2}$ .  
Then  $\|f\|_2 = \|\hat{f}\|_2$ . Again, we omit the proof of the norm-related properties.
2. Define the *convolution* of two functions as  $f * g(x) := \mathbb{E}_{y \in G} f(y)g(x - y)$ .  
Then this operation is commutative, and  $\widehat{(f * g)}(r) = \hat{f}(r)\hat{g}(r)$ .
3. Define  $\|\hat{f}\|_4 = (\sum_{r \in G} |\hat{f}(r)|^4)^{1/4}$ . Then  $\|f\|_{U^2} = \|\hat{f}\|_4$ .

In particular, Statement 3 is relevant to interpreting the  $U^2$ -norm.

## Idea 4: Discrete Fourier Analysis

We shall prove Statements 1, 2, and 3 in sequence. For Statement 1,

$$\|f\|_2^2 = \mathbb{E}_{x \in G} |f(x)|^2 = \mathbb{E}_{x \in G} f(x) \overline{f(x)} = \mathbb{E}_{x \in G} \sum_{r, r' \in G} \hat{f}(r) e\left(\frac{rx}{N'}\right) \overline{\hat{f}(r') e\left(\frac{r'x}{N'}\right)} = \dots$$

Reshuffling,

$$\dots = \sum_{r, r' \in G} \hat{f}(r) \overline{\hat{f}(r')} \mathbb{E}_{x \in G} e\left(\frac{(r - r')x}{N'}\right) = \sum_{r, r' \in G} \hat{f}(r) \overline{\hat{f}(r')} \delta_{rr'} = \sum_{r \in G} |\hat{f}(r)|^2 = \|\hat{f}\|_2^2$$

## Idea 4: Discrete Fourier Analysis

For Statement 2,

$$\widehat{(f * g)}(r) = \mathbb{E}_{x \in G} (f * g)(x) e(-\frac{rx}{N'}) = \mathbb{E}_{x, y \in G} f(y) g(x - y) e(-\frac{rx}{N'}) = \dots$$

Reparameterising  $x$  as  $x' = x - y$ ,

$$\dots = \mathbb{E}_{x', y \in G} f(y) g(x') e(-\frac{r(y + x')}{N'}) = \mathbb{E}_{x', y \in G} f(x') e(-\frac{rx'}{N'}) f(y) e(-\frac{ry}{N'}) = \hat{f}(r) \hat{g}(r)$$

## Idea 4: Discrete Fourier Analysis

For Statement 3, we reparameterise the  $U^2$ -norm:

$$\|f\|_{U^2}^4 = \mathbb{E}_{x \in G} \mathbb{E}_{y, y' \in G} f(y) f(x-y) \overline{f(y')} \overline{f(x-y')} = \mathbb{E}_{x \in G} (f * f)(x) \overline{(f * f)(x)} = \|f * f\|_2^2$$

We then apply Statements 1 and 2:

$$\|f * f\|_2^2 = \|\widehat{f * f}\|_2^2 = \sum_{r \in G} \widehat{(f * f)}(r) \overline{\widehat{(f * f)}(r)} = \sum_{r \in G} \hat{f}(r)^2 \overline{\hat{f}(r)^2} = \sum_{r \in G} |\hat{f}(r)|^4 = \|\hat{f}\|_4^4$$

Therefore  $\|f\|_{U^2} = \|\hat{f}\|_4$ .  $\square$

## Idea 4: Discrete Fourier Analysis

For Statement 3, we reparameterise the  $U^2$ -norm:

$$\|f\|_{U^2}^4 = \mathbb{E}_{x \in G} \mathbb{E}_{y, y' \in G} f(y) f(x-y) \overline{f(y')} \overline{f(x-y')} = \mathbb{E}_{x \in G} (f * f)(x) \overline{(f * f)(x)} = \|f * f\|_2^2$$

We then apply Statements 1 and 2:

$$\|f * f\|_2^2 = \|\widehat{f * f}\|_2^2 = \sum_{r \in G} \widehat{(f * f)}(r) \overline{\widehat{(f * f)}(r)} = \sum_{r \in G} \hat{f}(r)^2 \overline{\hat{f}(r)^2} = \sum_{r \in G} |\hat{f}(r)|^4 = \|\hat{f}\|_4^4$$

Therefore  $\|f\|_{U^2} = \|\hat{f}\|_4$ .  $\square$

## Idea 4: Discrete Fourier Analysis

Linking back to the problem at hand: for  $N \geq C\alpha^{-\frac{190}{2}}$ ,  $\|\hat{f}\|_4 \geq c\alpha^{\frac{1}{140}}$ .

Each  $\hat{f}(r)$  represents the values of  $f$  weighted in a certain periodic fashion. Note that  $f(x) = \{\alpha \text{ if } x \in A, \alpha - 1 \text{ if } x \notin [N] \setminus A, 0 \text{ otherwise}\}$ . If  $|\hat{f}(r)|$  is big for some  $r \in G$ , this could indicate an unequal density distribution related to some period.

In particular, we shall show that if function  $f : G = \mathbb{Z}/N'\mathbb{Z}$  is 1-bounded and  $\|\hat{f}\|_4 \geq \delta$ , then  $|\hat{f}(r)| \geq \delta^2$  for some  $r \in G$ .

## Idea 4: Discrete Fourier Analysis

Since  $f$  is 1-bounded, thus  $\|\hat{f}\|_2 = \|f\|_2 \leq 1$ .

Let  $\|\hat{f}\|_\infty = \max_{r \in G} |\hat{f}(r)|$ . Then

$$\delta^4 \leq \|\hat{f}\|_4^4 = \sum_{r \in G} |\hat{f}(r)|^4 \leq \|\hat{f}\|_\infty^2 \sum_{r \in G} |\hat{f}(r)|^2 = \|\hat{f}\|_\infty^2 \|\hat{f}\|_2^2 \leq \|\hat{f}\|_\infty^2$$

thus  $\|\hat{f}\|_\infty \geq \delta^2$ .



## Idea 4: Discrete Fourier Analysis

So for  $N \geq C\alpha^{-\frac{190}{2}}$ , there is some  $r \in G$  such that  $|\hat{f}(r)| \geq \frac{1}{140} (c\alpha^3)^2 = \frac{1}{19600} c\alpha^6$ .

Let  $\theta := -\frac{r}{N'} \pmod{1}$ . Then we may write

$$|\hat{f}(r)| = |\mathbb{E}_{x \in G} f(x) e(-\frac{rx}{N'})| = \frac{1}{N'} \left| \sum_{x \in [N]} f(x) e(\theta x) \right|$$

and therefore ( $N' = 2N + 1$ )

$$\left| \sum_{x \in [N]} f(x) e(\theta x) \right| \geq \frac{1}{19600} c\alpha^6 N' \geq \frac{1}{9800} c\alpha^6 N$$

# Idea 5: Hardy-Littlewood Circle Method

Density-Increment:  $\alpha' = \alpha + c\alpha^6$

Indicator  
Functions

$$f := 1_A - \alpha 1_{[N]}$$

$$AP_3 := \mathbb{E}_{x,d} g_1(x) g_2(x + d) g_3(x + 2d)$$

Cauchy-Schwarz  
and  $U^2$ -norm

$$|AP_3(g_1, g_2, g_3)| \leq \|g_i\|_{U^2}$$

$$\|f\|_{U^2} \geq c\alpha^3$$

Discrete Fourier  
Analysis

$$\exists \theta \in [0, 1] \text{ s.t.}$$

$$|\sum f(x) e(\theta x)| \geq c\alpha^6 N$$

Hardy-Littlewood  
Circle Method

$$\exists P' \subseteq [N], |P'| \geq N^{1/3},$$

$$\text{s.t.}$$

$$\sum_{x \in P'} f(x) \geq c\alpha^6 |P'|$$



$$\alpha_{\max} \leq \frac{C}{(\log \log N)^{1/5}}$$

## Idea 5: Hardy-Littlewood Circle Method

Note for any subset  $S \subseteq [N]$ , the value  $\sum_{x \in S} f(x)$  tells us about the density of  $A$  in  $S$ . In particular, if  $\sum_{x \in S} f(x) > \alpha |S|$ , then  $|A \cap S| > \alpha |S|$ , and vice versa.

We therefore wish to consider sub-APs  $P \subseteq [N]$  of which for  $x \in P$ , the values of  $e(\theta x)$  are nearly identical.

This is in fact true when the common difference  $d$  is such that  $\|\theta d\|_{\mathbb{R}/\mathbb{Z}}$  is small, and the length of the sub-AP is not too large ( $|P| \cdot \|\theta d\|_{\mathbb{R}/\mathbb{Z}}^2 \ll 1$ ).

## Idea 5: Hardy-Littlewood Circle Method

We first show that we can pick  $d$  such that  $\|\theta d\|_{\mathbb{R}/\mathbb{Z}}$  is small.

Let  $\delta > 0$ ; then we can pick  $d \in \{1, 2, \dots, \lfloor 1/\delta \rfloor\}$  such that  $\|\theta d\|_{\mathbb{R}/\mathbb{Z}} \leq \delta$ .

To show this, apply Pigeonhole Principle to  $0, \theta, 2\theta, \dots, \lfloor 1/\delta \rfloor \theta \pmod{1}$  on  $\mathbb{R}/\mathbb{Z}$ . There must be at least one pair  $a\theta, b\theta$  of distance  $\|a\theta - b\theta\|_{\mathbb{R}/\mathbb{Z}} \leq \delta$ .

Set  $d = |a - b|$ ; then  $\|\theta d\|_{\mathbb{R}/\mathbb{Z}} \leq \delta$  as desired.

## Idea 5: Hardy-Littlewood Circle Method

Now we consider dividing  $[N]$  into sub-APs of common difference  $d$ .

There would be  $d$  sub-APs, each of length  $\geq \lfloor N/d \rfloor \geq \lfloor \delta N \rfloor$ .

We wish to further subdivide these APs such that the values of  $e(\theta x)$  are nearly identical, so that

$$\left| \sum_{x \in P} f(x) e(\theta x) \right| \approx \left| \sum_{x \in P} f(x) \right|$$

to sufficient proximity.

## Idea 5: Hardy-Littlewood Circle Method

Let  $P = a_0 - \lfloor \frac{|P|}{2} \rfloor d, \dots, a_0 - d, a_0, a_0 + d, \dots, a_0 + (\lceil \frac{|P|}{2} \rceil - 1)d$ .

Then  $|\sum_{x \in P} f(x)| = |\sum_{x \in P} f(x)e(\theta a_0)|$ . So

$$\left| \sum_{x \in P} f(x)e(\theta x) \right| = \left| \sum_{x \in P} f(x)e(\theta a_0) + \sum_{x \in P} f(x)(e(\theta x) - e(\theta a_0)) \right|$$

Let  $\Delta_P = \left| \left| \sum_{x \in P} f(x)e(\theta x) \right| - \left| \sum_{x \in P} f(x) \right| \right|$ . By Triangle Inequality,

$$\Delta_P \leq \left| \sum_{x \in P} f(x)(e(\theta x) - e(\theta a_0)) \right| \leq \sum_{x \in P} |e(\theta x) - e(\theta a_0)|$$

$f$  is 1-bounded

## Idea 5: Hardy-Littlewood Circle Method

We can evaluate and bound the RHS.

Note that  $|e(x) - e(y)| = 2 \sin(\pi \|x - y\|_{\mathbb{R}/\mathbb{Z}}) \leq 2\pi \|x - y\|_{\mathbb{R}/\mathbb{Z}}$ . Therefore

$$\sum_{x \in P} |e(\theta x) - e(\theta a_0)| \leq 2\pi \sum_{x \in P} \|\theta x - \theta a_0\|_{\mathbb{R}/\mathbb{Z}} = 2\pi \sum_{y \in Q} \|\theta dy\|_{\mathbb{R}/\mathbb{Z}}$$

where  $Q = \lfloor \frac{|P|}{2} \rfloor, \dots, -2, -1, 0, 1, 2, \dots, \lceil \frac{|P|}{2} \rceil - 1$ .

## Idea 5: Hardy-Littlewood Circle Method

Since  $\|\theta d\|_{\mathbb{R}/\mathbb{Z}} \leq \delta$ , thus

$$2\pi \sum_{y \in Q} \|\theta dy\|_{\mathbb{R}/\mathbb{Z}} \leq 2\pi \sum_{y \in Q} |y| \cdot \|\theta d\|_{\mathbb{R}/\mathbb{Z}} \leq 2\pi \delta \sum_{y \in Q} |y|^2 \leq 2\pi \delta \frac{|P|^2}{4} = \frac{\pi}{2} \delta |P|^2$$

and ultimately  $\Delta_P \leq \frac{\pi}{2} \delta |P|^2$ .



## Idea 5: Hardy-Littlewood Circle Method

We now return to the problem at hand. We partition  $[N]$  into sub-APs  $P_1, P_2, \dots, P_k$ , having common difference  $d$ ; then for  $N \geq c\alpha^{-C}$ ,<sup>2</sup>

$$\frac{1}{9800} c\alpha^6 N \leq \left| \sum_{x \in [N]} f(x) e(\theta x) \right| \leq \sum_{j=1}^k \left| \sum_{x \in P_j} f(x) e(\theta x) \right| \leq \sum_{j=1}^k \left( \left| \sum_{x \in P_j} f(x) \right| + \Delta_{P_j} \right)$$

Therefore

$$\sum_{j=1}^k \left| \sum_{x \in P_j} f(x) \right| \geq \frac{1}{9800} c\alpha^6 N - \sum_{j=1}^k \Delta_{P_j} \geq \frac{1}{9800} c\alpha^6 N - \sum_{j=1}^k \frac{\pi}{2} \delta |P_j|^2$$

## Idea 5: Hardy-Littlewood Circle Method

1/1.54E6

We now carefully set our variables for success. Let  $\delta = c\alpha^6 N^{-1/3}$  and  $d$  as defined previously.

Now  $[N]$  may be partitioned into  $d$  sub-APs, each of common difference  $d$  and length  $\geq \lfloor \delta N \rfloor \geq \lfloor c\alpha^6 N^{2/3} \rfloor$ . Taking away  $\lceil N^{1/3} \rceil$  at a time, the only cases left to consider is when  $\lfloor N^{1/3} \rfloor$  or  $\lfloor N^{1/3} \rfloor - 1$  remains.

1/1.54E6

Since  $\lfloor c\alpha^6 N^{2/3} \rfloor \geq 2\lceil N^{1/3} \rceil$ , thus the APs may be chopped finer into sub-APs of length  $\geq N^{1/3}$  and  $\leq 2N^{1/3}$ . Since  $N^{1/3} \geq 2$ , we may split the remainder between two of the  $\lceil N^{1/3} \rceil$  blocks.

true if  $N \geq (3E19)\alpha^{-18}$

## Idea 5: Hardy-Littlewood Circle Method

Linking back,

$$\sum_{j=1}^k \frac{\pi}{2} \delta |P_j|^2 \leq \sum_{j=1}^k \frac{\pi}{2} (c\alpha^6 N^{-1/3}) (2N^{1/3}) |P_j| = c\alpha^6 \sum_{j=1}^k |P_j| = c\alpha^6 N$$

Therefore

$$\sum_{j=1}^k \left| \sum_{x \in P_j} f(x) \right| \geq c\alpha^6 N - c\alpha^6 N = c\alpha^6 N$$

## Idea 5: Hardy-Littlewood Circle Method

Finally, we apply the fact that  $f$  is balanced:

$$\sum_{j=1}^k \sum_{x \in P_j} f(x) = \sum_{x \in [N]} f(x) = 0$$

to get

$$\sum_{j=1}^k \left( \left| \sum_{x \in P_j} f(x) \right| + \sum_{x \in P_j} f(x) \right) \stackrel{1/1E4}{\geq} c\alpha^6 N = c\alpha^6 \sum_{j=1}^k |P_j|$$

## Idea 5: Hardy-Littlewood Circle Method

Ultimately, there is some  $j \in \{1, 2, \dots, k\}$  such that

$$\sum_{x \in P_j} f(x) \geq \frac{1}{2} c \alpha^6 |P_j| = c \alpha^6 |P_j|$$

Now  $P_j$  has length  $\geq N^{1/3}$ , and  $|A \cap P_j| \geq (\alpha + c \alpha^6) |P_j|$ .

This proves the Density-Increment under the condition  $N \geq C \alpha^{-c}$ .  $\square$

# Proof Summary

Density-Increment:  $\alpha' = \alpha + c\alpha^6$

Indicator  
Functions

$$f := 1_A - \alpha 1_{[N]}$$

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$$|\sum f(x) e(\theta x)| \geq c\alpha^6 N$$

Hardy-Littlewood  
Circle Method

$$\exists P' \subseteq [N], |P'| \geq N^{1/3},$$

$$\text{s.t.}$$

$$\sum_{x \in P'} f(x) \geq c\alpha^6 |P'|$$



$$\alpha_{\max} \leq \frac{c}{(\log \log N)^{1/5}}$$

## Alternative Approaches

Apart from the approach detailed in these slides (Probabilistic Method + Fourier Analysis), it is worth exploring others:

- (Hales-Jewett) In a  $3 \times 3 \times \cdots \times 3$  hypercube with  $\alpha 3^n$  selected cells, there is a combinatorial line of  $n$  selected cells. (Ergodic Theory)
- (Corners) In an  $N \times N$  grid of  $\alpha N^2$  selected cells, there is a "corner" triplet  $\{(x, y), (x + h, y), (x, y + h)\}$ . (Graph Theory)

## Extensions

- (Szemerédi) Consider  $k$ -term APs.
- Consider multidimensional APs.
- (Leibman, Bergelson) Consider polynomial progressions.
- (Erdős) Strengthen the density condition to large sets,  $\sum \frac{1}{a_i} \rightarrow +\infty$ .