Fall 2024

Section C: Algorithms Analysis

1) (10 pts) ANL (Algorithm Analysis)

Consider the task of sorting \mathbf{n}^2 integers. Using an insertion sort, this task would take $O(\mathbf{n}^4)$ time. Using a single heap sort, this task would take $O(\mathbf{n}^2 \mathbf{lg} \mathbf{n})$. Consider this hybrid approach and, with proof, determine its worst case run time, in terms of \mathbf{n} . Assume efficient implementations of each of the heap and linked list operations described. Leave your answer in Big-Oh notation.

- 1. Separate the n^2 integers into n groups of n integers each.
- 2. Create heaps out of each of the **n** groups of integers.
- 3. Call delete min on each of the \mathbf{n} heaps, storing these \mathbf{n} deleted values in a linked list, also storing which heap each value came from.
- 4. Repeat the following n^2 times:
 - a. Loop through the linked list, locating the minimum integer in it, noting which heap it was from. Name the integer \mathbf{x} and the heap \mathbf{H} .
 - b. Place **x** next in the sorted list and delete it from the linked list.
 - c. If **H** isn't empty, delete the minimum item from **H** and add it to the end of the linked list, also storing that the value came from heap **H**.

The first step takes $O(n^2)$ time as each number if processed once.

The run-time of make heap on n integers is O(n). Since we do this n times, the total run-time of step 2 is $O(n^2)$.

Each delete min call on a heap takes $O(\lg n)$ time because each of the heaps from which we delete items will have no more than n items. Ultimately, each item from each heap gets deleted, so the total run time of all the delete mins on the heap is $O(n^2 \lg n)$ time. This encompasses the total run time of steps 3 and 4c.

In step 4, the loop in part (a) takes O(n) time and we repeat it n^2 time, meaning that we spend a total of $O(n^3)$ time on step 4a across all of its iterations.

Step 4b takes O(1) time each time it's executed for a total of O(n²) time.

If we add up all of these run-times we get:

$$O(n^2) + O(n^2) + O(n^2 \lg n) + O(n^3) + O(n^2) = O(n^3)$$

An alternate way to do the analysis is to state that step 3 takes O(nlgn) time, followed by a single iteration of step 4 taking O(n) + O(1) + O(lg n) = O(n) time, then noting that since this gets repeated n^2 times, the total run time of step 4 is $O(n^3)$,

Grading: 2 pts for the correct answer

2 pts for explaining why step 2 takes either $O(n^2)$ or $O(n^2 lg \ n)$ time

3 pts for multiplying n² times the run time of steps 4a, 4b and 4c.

3 pts for the run time of steps 4a, 4b and 4c as well as step 3 with explanation