

Exercise 1

Background: The question begins by defining a relation “ \sqsubseteq ” on the set $\mathbb{Z} \times \mathbb{Z}$. We should start by making sure we understand the terminology being used here. The symbol \mathbb{Z} indicates the set of all integers

$$\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}.$$

But the definition refers to something more than this: it uses $\mathbb{Z} \times \mathbb{Z}$.

Remember that in general $A \times B$ is the set of pairs of elements coming from A and B . In other words, an element of $A \times B$ is something of the form (a, b) where $a \in A$ and $b \in B$ (a is an element of A and b is an element of B). For example, if $A = \{1, 2\}$ and $B = \{x, y, z\}$, then we have that

$$A \times B = \left\{ \begin{array}{l} (1, x), (1, y), (1, z), \\ (2, x), (2, y), (2, z) \end{array} \right\}.$$

From this we can work out what is $\mathbb{Z} \times \mathbb{Z}$. Its elements are all pairs (u, v) where u and v are integers. We can visualise this as an infinite grid of elements:

$$\mathbb{Z} \times \mathbb{Z} = \left\{ \begin{array}{cccccc} \ddots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \\ \dots & (-2, -2), & (-1, -2), & (0, -2), & (1, -2), & (2, -2), & \dots \\ \dots & (-2, -1), & (-1, -1), & (0, -1), & (1, -1), & (2, -1), & \dots \\ \dots & (-2, 0), & (-1, 0), & (0, 0), & (1, 0), & (2, 0), & \dots \\ \dots & (-2, 1), & (-1, 1), & (0, 1), & (1, 1), & (2, 1), & \dots \\ \dots & (-2, 2), & (-1, 2), & (0, 2), & (1, 2), & (2, 2), & \dots \\ \ddots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{array} \right\}$$

So when we say (a, b) is an element of $\mathbb{Z} \times \mathbb{Z}$, we mean that a and b are integers. Therefore \sqsubseteq being a relation on $\mathbb{Z} \times \mathbb{Z}$ just means that \sqsubseteq is a way to compare pairs of integers with each other. The definition given in the exercise is that

$$(a, b) \sqsubseteq (c, d) \text{ if } a < c \text{ or } a = c \text{ and } b \leq d.$$

Put another way, this means

- if $a < c$ then $(a, b) \sqsubseteq (c, d)$,
- and, if $a = c$ and $b \leq d$ then $(a, b) \sqsubseteq (c, d)$.

This is the ‘lexicographical order’ or the ‘dictionary order’. If we were comparing letters instead of numbers this would be like alphabetical ordering for 2-letter words: If two words start with different letters then we compare them by looking at the first letter only. But if they start with the same first letter we have to move to the second letter to order them.

Part 1: If X is a set and \preceq is a relation on it, then we say \preceq is a partial order when the following conditions hold:

- The relation is reflexive: every element x in X is related to itself ($x \preceq x$)
- The relation is antisymmetric: the only way for $x \preceq y$ and $y \preceq x$ at the same time is if $x = y$
- The relation is transitive: if $x \preceq y$ and if $y \preceq z$ then also we must have that $x \preceq z$.

We now show that \sqsubseteq is a partial order. In the definition for \sqsubseteq we looked at earlier, there were two possible ways in which we could have $(a, b) \sqsubseteq (c, d)$:

- (i) $a < c$, or
- (ii) $a = c$ and $b \leq d$

We number these (i) and (ii) so that we can refer to the two cases individually in the proof.

Now, there are three properties we need to check to show that \sqsubseteq is a partial order. Let's prove them one-by-one:

- **Reflexivity:** Consider a pair (a, b) of integers. We have that $a = a$ and $b \leq b$. This is exactly the condition required by case (ii) in the definition of \sqsubseteq , for $(a, b) \sqsubseteq (a, b)$. So \sqsubseteq is reflexive.
- **Antisymmetry:** Assume we have two pairs of integers (a, b) and (c, d) such that $(a, b) \sqsubseteq (c, d)$ and $(c, d) \sqsubseteq (a, b)$. There are two ways for the relation $(a, b) \sqsubseteq (c, d)$ to hold. Since we have assumed two relations of this form, there are four ways for this to be true overall:
 1. $(a, b) \sqsubseteq (c, d)$ by case (i) and $(c, d) \sqsubseteq (a, b)$ by case (i)
 2. $(a, b) \sqsubseteq (c, d)$ by case (i) and $(c, d) \sqsubseteq (a, b)$ by case (ii)
 3. $(a, b) \sqsubseteq (c, d)$ by case (ii) and $(c, d) \sqsubseteq (a, b)$ by case (i)
 4. $(a, b) \sqsubseteq (c, d)$ by case (ii) and $(c, d) \sqsubseteq (a, b)$ by case (ii)

However, we can rule out some of these cases because they make no sense! We check each case individually:

1. This means that $a < c$ and $c < a$, which is impossible
2. This means that $a < c$ and that $c = a$, which is impossible
3. This means that $a = c$ and that $c < a$, which is impossible
4. This means that $a = c$ and $b \leq d$ and that $c = a$ and $d \leq b$. In particular we have that $b \leq d$ and $d \leq b$, so $d = b$.

The only case that is actually possible is the fourth one, and in that case we know $a = c$ and $d = b$, so $(a, c) = (b, d)$. Hence \sqsubseteq is antisymmetric.

- **Transitivity** Similarly to the previous case, if we have two relations $(a, b) \sqsubseteq (c, d)$ and $(c, d) \sqsubseteq (e, f)$, there are four possible cases:

1. $a < c$, and $c < e$, hence $a < e$.
2. $a < c$, and $c = e$ and $d \leq f$. Therefore $a < e$.
3. $a = c$ and $b \leq d$, and $c < e$. Therefore $a < e$.
4. $a = c$ and $b \leq d$, and $c = e$ and $d \leq f$. Therefore $a = e$ and $b \leq f$.

In the first three cases we have that $a < e$, meaning $(a, b) \sqsubseteq (e, f)$ by case (i). In the fourth case we have that $a = e$ and $b \leq f$, so $(a, b) \sqsubseteq (e, f)$. So, in any case we have that $(a, b) \sqsubseteq (e, f)$ and so \sqsubseteq is transitive.

We have shown that \sqsubseteq is reflexive, antisymmetric, and transitive, hence it is a partial order. \square

Part 2: Given a partial order \preceq , we say this is a total order if any pair of elements are related by this in some direction. That is for any x and y , we either have $x \preceq y$ or $y \preceq x$.

Therefore, to show that \sqsubseteq is a total order we must show that for any pairs of integers (a, b) and (c, d) , we either have $(a, b) \sqsubseteq (c, d)$ or $(c, d) \sqsubseteq (a, b)$. Indeed, for any a and c we will have one of three cases:

1. $a < c$
2. $a > c$
3. $a = c$

Case 1 implies that $(a, b) \sqsubseteq (c, d)$, and case 2 implies that $(c, d) \sqsubseteq (a, b)$.

The third case however, does not imply a relation in either direction. If $a = c$ then $(a, b) \sqsubseteq (c, d)$ is true exactly when $b \leq d$. Dually, if $a = c$ then $(c, d) \sqsubseteq (a, b)$ is true exactly when $d \leq b$. But we know that either $b \leq d$ or $d \leq b$ always holds (because \leq is a total order!). Hence, we have that if $a = c$ then $(a, b) \sqsubseteq (c, d)$ or $(c, d) \sqsubseteq (a, b)$, as required.

Therefore, in all three cases $(a, b) \sqsubseteq (c, d)$ or $(c, d) \sqsubseteq (a, b)$, and so \sqsubseteq is a total order. \square