Chapter 3:

Control, Conditional Branching, Loops

Topics

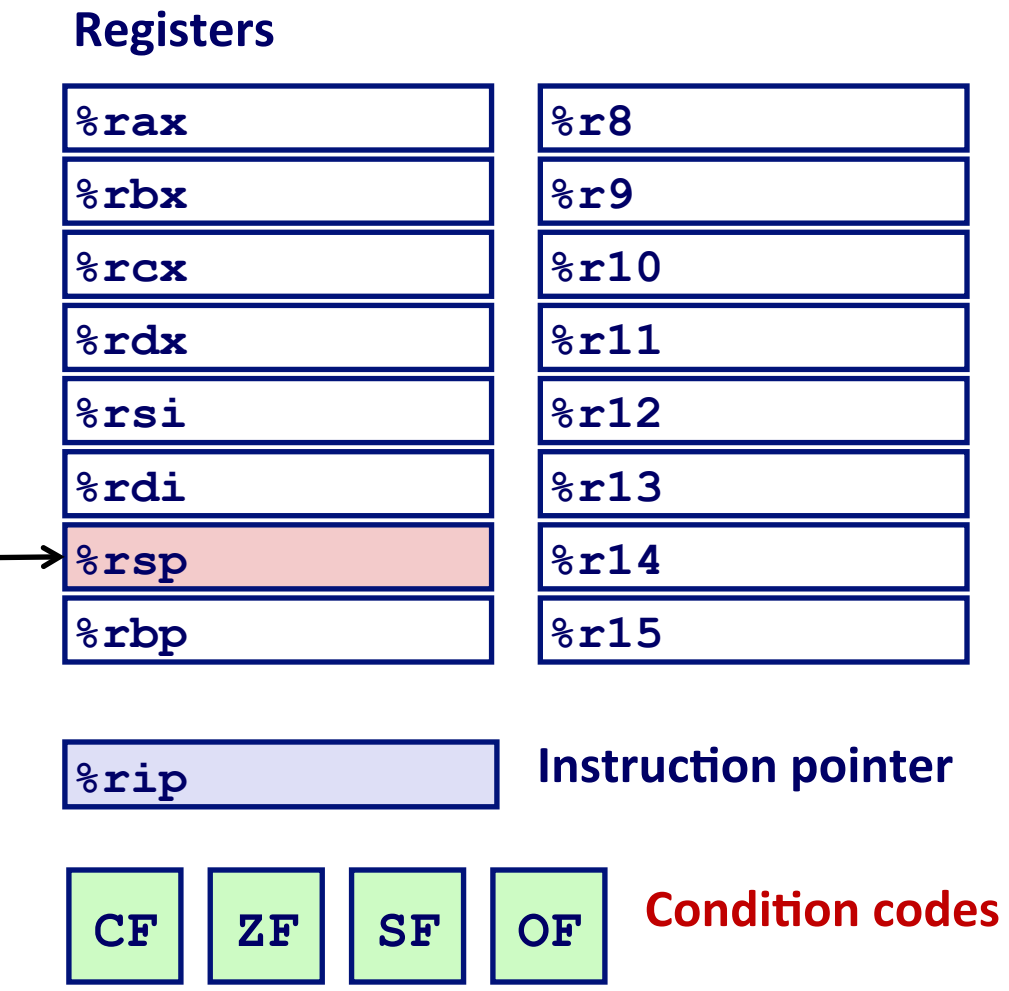
- Condition Codes
- Conditional Jumping/ Branching
- Loops

Announcements

- Data Lab grading this week
  - Sign up for 12-minute slots, come prepared and don't forget!
- Next Assembly Quiz on moodle, due Mon Feb 13 by noon
- Bomb Lab #2 released, due Friday Feb 24 by 11:55 pm
  - TAs will cover phase one this week in recitation
  - Can form teams of two
- Read Chapter 3.1-3.12 (except 3.11) and do practice problems

Processor State (x86-64, Partial)

- **Information about currently executing program**
  - **Temporary data (`%rax`, ...)**
  - **Location of runtime stack (`%rsp`)**
  - **Location of current code control point (`%rip`, ...)**
  - **Status of recent tests (CF, ZF, SF, OF)**



Condition Codes (Implicit Setting)

- Single bit registers
  - `CF` Carry Flag (for unsigned)
  - `SF` Sign Flag (for signed)
  - `ZF` Zero Flag
  - `OF` Overflow Flag (for signed)

- Implicitly set (think of it as side effect) by arithmetic operations
  - `addq` ***Src,Dest***
  - C analog: `t = a+b`
    - CF set if carry out from most significant bit
      - Used to detect unsigned overflow
    - ZF set if `t == 0`
    - SF set if `t < 0` (as signed)
    - OF set if two's complement (signed) overflow

```
(a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)
```

- Codes set differently depending on instructions, and in some cases not set at all, e.g. `lea` instruction

Condition Codes (Explicit Setting: Compare)

- **Explicit Setting by Compare Instruction**

  `cmpq` ***Src2,Src1***

  `cmpq b,a` like computing `a-b` without setting destination

  - **CF set** if carry out from most significant bit (used for unsigned comparisons)
  - **ZF set** if `a == b`
  - **SF set** if `(a-b) < 0` (as signed)
  - **OF set** if two's complement (signed) overflow

    `(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)`

Condition Codes (Explicit Setting: Test)

- **Explicit Setting by Test instruction**

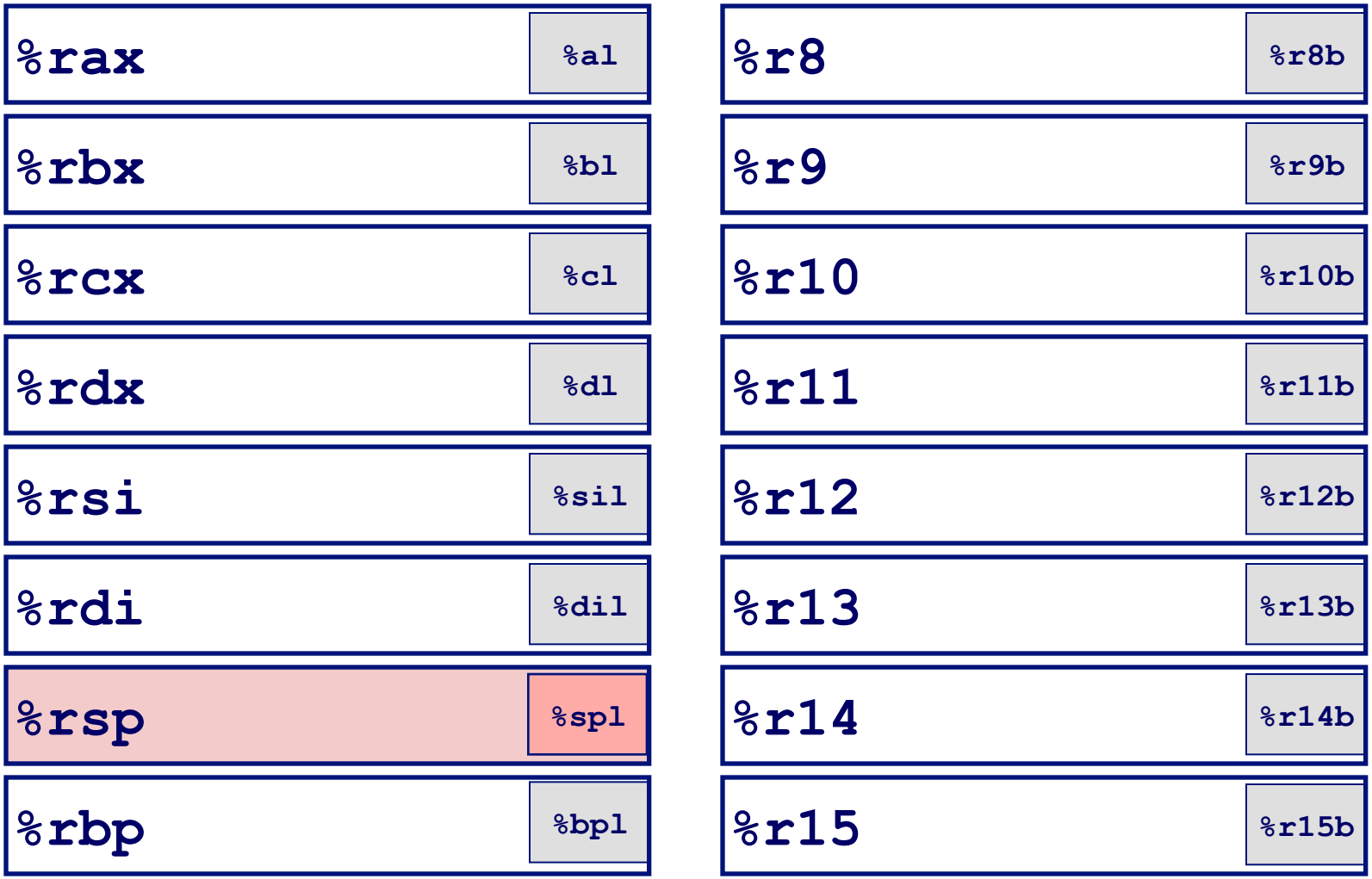
  `testq` ***Src2,Src1***

  `testq b,a` like computing `a&b` without setting destination

  - Sets condition codes based on value of *Src1* `&` *Src2*
  - Useful to have one of the operands be a mask

  - ZF set when `a&b == 0`
  - SF set when `a&b < 0`

x86-64 Integer Registers



- Can reference low-order byte

Reading Condition Codes

- **setX dest // e.g.** `setl %al`
  - **Set low-order byte of destination to 0 or 1 based on combinations of condition codes**
  - **Does not alter remaining 7 bytes**
  - **Previous instruction should be a** `cmp`

SetX	Condition	Description
`sete`	`ZF`	Equal / Zero
`setne`	`~ZF`	Not Equal / Not Zero
`sets`	`SF`	Negative
`setns`	`~SF`	Nonnegative
`setg`	`~(SF^OF)&~ZF`	Greater (Signed)
`setge`	`~(SF^OF)`	Greater or Equal (Signed)
`setl`	`(SF^OF)`	Less (Signed) – derivation in text
`setle`	`(SF^OF)|ZF`	Less or Equal (Signed)
`seta`	`~CF&~ZF`	Above (unsigned)
`setb`	`CF`	Below (unsigned)

Reading Condition Codes (Cont.)

- **SetX Instructions:**
  - Set single byte based on combination of condition codes
- **One of addressable byte registers**
  - Does not alter remaining bytes
  - Typically use `movzbl` to finish job
    - Also a `movzbq`, etc.

```
int gt (long x, long y)
{
  return x > y;
}
```

Register	Use(s)
`%rdi`	Argument `x`
`%rsi`	Argument `y`
`%rax`	Return value

```
    cmpq    %rsi, %rdi    # Compare x:y
    setg    %al           # Set when >
    movzbl  %al, %eax     # Zero rest of %rax
    ret
```

Jumping – Conditional or not

- **jX Instructions**
  - **Unconditional jump: `jmp` Label or `jmp` *%eax or `jmp` *(%eax)**
  - **Conditional jumps to different part of code depending on condition codes, e.g. `jle` Label**

jX	Condition	Description
`jmp`	`1`	Unconditional
`je`	`ZF`	Equal / Zero
`jne`	`~ZF`	Not Equal / Not Zero
`js`	`SF`	Negative
`jns`	`~SF`	Nonnegative
`jg`	`~(SF^OF)&~ZF`	Greater (Signed)
`jge`	`~(SF^OF)`	Greater or Equal (Signed)
`jl`	`(SF^OF)`	Less (Signed)
`jle`	`(SF^OF)|ZF`	Less or Equal (Signed)
`ja`	`~CF&~ZF`	Above (unsigned)
`jb`	`CF`	Below (unsigned)

Conditional Branch Example (Old Style)

- If-then-else converted to assembly form

```
long absdiff
  (long x, long y)
{
  long result;
  if (x > y)
    result = x-y;
  else
    result = y-x;
  return result;
}
```

```
absdiff:
    cmpq    %rsi, %rdi   # x:y
    jle     .L4
    movq    %rdi, %rax
    subq    %rsi, %rax
    ret
.L4:                     # x <= y
    movq    %rsi, %rax
    subq    %rdi, %rax
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

Expressing with Goto Code

- C allows `goto` statement
- Jump to position designated by label
- Generally considered bad programming practice

```
long absdiff
  (long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
long absdiff_j
  (long x, long y)
{
    long result;
    int ntest = x <= y;
    if (ntest) goto Else;
    result = x-y;
    goto Done;
 Else:
    result = y-x;
 Done:
    return result;
}
```

General Conditional Expression Translation (Using Branches)

C Code

```
val = Test ? Then_Expr : Else_Expr;
```

```
val = x>y ? x-y : y-x;
```

Goto Version

```
  ntest = !Test;
  if (ntest) goto Else;
  val = Then_Expr;
  goto Done;
Else:
  val = Else_Expr;
Done:
  . . .
```

- *Test* is expression returning integer
  - = 0 interpreted as false
  - ≠0 interpreted as true
- Create separate code regions for then & else expressions
- Execute appropriate one

Conditional Move

- **cmovX src, dest**
  - Set dest=src only if condition X holds
  - More efficient than conditional branching for highly pipelined processors – easier to guess the next instruction to execute
  - But overhead: both branches are evaluated

cmovX	Condition	Description
`cmove`	`ZF`	Equal / Zero
`cmovne`	`~ZF`	Not Equal / Not Zero
`cmovs`	`SF`	Negative
`cmovns`	`~SF`	Nonnegative
`cmovg`	`~(SF^OF)&~ZF`	Greater (Signed)
`cmovge`	`~(SF^OF)`	Greater or Equal (Signed)
`cmovl`	`(SF^OF)`	Less (Signed)
`cmovle`	`(SF^OF)|ZF`	Less or Equal (Signed)
`cmova`	`~CF&~ZF`	Above (unsigned)
`cmovb`	`CF`	Below (unsigned)

Using Conditional Moves

- **Conditional Move Instructions**
  - Instruction supports:
    if (Test) Dest ← Src
  - Supported in post-1995 x86 processors
  - GCC tries to use them
    - But, only when known to be safe
- **Why?**
  - Branches are very disruptive to instruction flow through pipelines
  - Conditional moves do not require control transfer

C Code

```
val = Test
    ? Then_Expr
    : Else_Expr;
```

Goto Version

```
result = Then_Expr;
eval = Else_Expr;
nt = !Test;
if (nt) result = eval;
return result;
```

Conditional Move Example

```
long absdiff
   (long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
absdiff:
    movq    %rdi, %rax  # x
    subq    %rsi, %rax  # result = x-y
    movq    %rsi, %rdx
    subq    %rdi, %rdx  # eval = y-x
    cmpq    %rsi, %rdi  # x:y
    cmovle  %rdx, %rax  # if <=, result = eval
    ret
```

Bad Cases for Conditional Move

Expensive Computations

```
val = Test(x) ? Hard1(x) : Hard2(x);
```

- **Both values get computed**
- **Only makes sense when computations are very simple**

Risky Computations

```
val = p ? *p : 0;
```

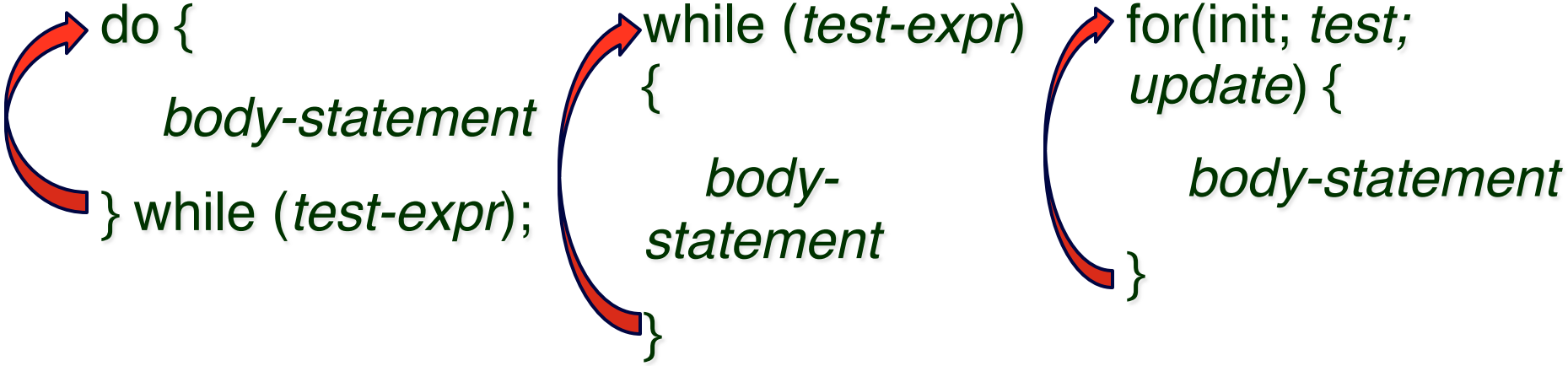
- Both values get computed
- **May have undesirable effects**

Computations with side effects

```
val = x > 0 ? x*=7 : x+=3;
```

- Both values get computed
- **Must be side-effect free**

Loops in C



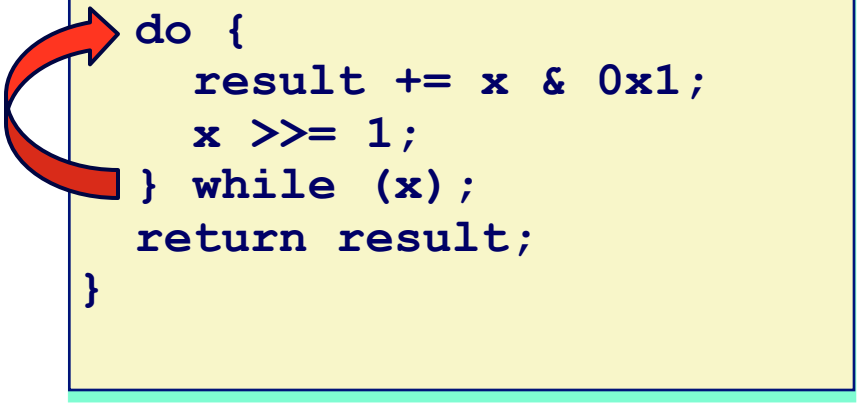
- **Executes body-statement**
- **Then tests expression**
  - **If true, loops back to 'do'**
  - **Else exit**

- **Tests expression first**
  - **If true, executes body & loops back to 'while'**
  - **Else exit**

- **Initializes first**
- **If test is true**
  - **Execute body statement**
  - **Execute update & loop back to 'for'**
- **Else exit**

"Do-While" Loop Example

C Code



```
long pcount_do
  (unsigned long x) {
  long result = 0;
  do {
    result += x & 0x1;
    x >>= 1;
  } while (x);
  return result;
}
```

Goto Version

```
long pcount_goto
  (unsigned long x) {
  long result = 0;
 loop:
  result += x & 0x1;
  x >>= 1;
  if(x) goto loop;
  return result;
}
```

- Count number of 1's in argument `x` ("popcount")
- Use conditional branch to either continue looping or to exit loop

“Do-While” Loop Compilation

Goto Version

```
long pcount_goto
   (unsigned long x) {
  long result = 0;
 loop:
  result += x & 0x1;
  x >>= 1;
  if(x) goto loop;
  return result;
}
```

Register	Use(s)
%rdi	Argument x
%rax	result

```
   movl    $0, %eax    #  result = 0
.L2:                   # loop:
   movq    %rdi, %rdx
   andl    $1, %edx    #  t = x & 0x1
   addq    %rdx, %rax  #  result += t
   shrq    %rdi        #  x >>= 1
   jne     .L2         #  if (x) goto loop
   rep; ret
```

General "Do-While" Translation

C Code

```
do
   Body
   while (Test);
```

Goto Version

```
loop:
   Body
   if (Test)
      goto loop
```

- **Body:**

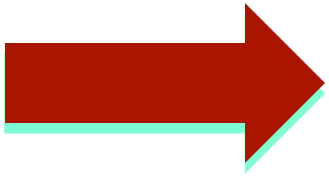
```
{
   Statement1;
   Statement2;
      …
   Statementn;
}
```

General "While" Translation #1

- "Jump-to-middle" translation
- Used with `-Og`

While version

```
while (Test)
  Body
```



Goto Version

```
  goto test;
loop:
  Body
test:
  if (Test)
    goto loop;
done:
```

While Loop Example #1

C Code

```
long pcount_while
   (unsigned long x) {
   long result = 0;
   while (x) {
     result += x & 0x1;
     x >>= 1;
   }
   return result;
}
```

Jump to Middle

```
long pcount_goto_jtm
   (unsigned long x) {
   long result = 0;
   goto test;
 loop:
   result += x & 0x1;
   x >>= 1;
 test:
   if(x) goto loop;
   return result;
}
```

- Compare to do-while version of function
- Initial goto starts loop at test

General "While" Translation #2

While version

```
while (Test)
  Body
```

Do-While Version

```
  if (!Test)
    goto done;
  do
    Body
    while(Test);
done:
```

Goto Version

```
  if (!Test)
    goto done;
loop:
  Body
  if (Test)
    goto loop;
done:
```

- "Do-while" conversion
- Used with `-O1`

While Loop Example #2

C Code

```
long pcount_while
   (unsigned long x) {
   long result = 0;
   while (x) {
     result += x & 0x1;
     x >>= 1;
   }
   return result;
}
```

Do-While Version

```
long pcount_goto_dw
   (unsigned long x) {
   long result = 0;
   if (!x) goto done;
 loop:
   result += x & 0x1;
   x >>= 1;
   if(x) goto loop;
 done:
   return result;
}
```

- Compare to do-while version of function
- Initial conditional guards entrance to loop

“For” Loop Form

General Form

```
for (Init; Test; Update)
    Body
```

```
#define WSIZE 8*sizeof(int)
long pcount_for
  (unsigned long x)
{
  size_t i;
  long result = 0;
  for (i = 0; i < WSIZE; i++)
  {
    unsigned bit =
      (x >> i) & 0x1;
    result += bit;
  }
  return result;
}
```

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

Body

```
{
  unsigned bit =
      (x >> i) & 0x1;
  result += bit;
}
```

"For" Loop → While Loop

For Version

```
for (Init; Test; Update)
        Body
```

While Version

```
Init;
while (Test) {
    Body
    Update;
}
```

For-While Conversion

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

Body

```
{
  unsigned bit =
     (x >> i) & 0x1;
  result += bit;
}
```

```
long pcount_for_while
   (unsigned long x)
{
   size_t i;
   long result = 0;
   i = 0;
   while (i < WSIZE)
   {
      unsigned bit =
         (x >> i) & 0x1;
      result += bit;
      i++;
   }
   return result;
}
```

"For" Loop Do-While Conversion

C Code

```
long pcount_for
  (unsigned long x)
{
  size_t i;
  long result = 0;
  for (i = 0; i < WSIZE; i++)
  {
    unsigned bit =
      (x >> i) & 0x1;
    result += bit;
  }
  return result;
}
```

Goto Version

```
long pcount_for_goto_dw
  (unsigned long x) {
  size_t i;
  long result = 0;
  i = 0;
  if (!(i < WSIZE))
    goto done;
 loop:
  {
    unsigned bit =
      (x >> i) & 0x1;
    result += bit;
  }
  i++;
  if (i < WSIZE)
    goto loop;
 done:
  return result;
}
```

Init

! Test

Body

Update

Test

- Initial test can be optimized away

```
typedef enum
 {ADD, MULT, MINUS, DIV, MOD, BAD}
    op_type;

char unparse_symbol(op_type op)
{
  switch (op) {
    case ADD :
      return '+';
    case MULT:
      return '*';
    case MINUS:
      return '-';
    case DIV:
      return '/';
    case MOD:
      return '%';
    case BAD:
      return '?';
  }
}
```

Switch Statements

Implementation Options

- Series of conditionals, e.g. if- else if - else if - else if…
  - Good if few cases
  - Slow if many cases, e.g. many compares and conditional jumps
- Jump Table (array of addresses)
  - Index into array and jump to branch target
  - Avoids conditionals
  - Good when cases are small integer constants
- GCC picks one based on case structure

switch() is missing default case

```
long switch_eg
   (long x, long y, long z)
{
    long w = 1;
    switch(x) {
    case 1:
        w = y*z;
        break;
    case 2:
        w = y/z;
        /* Fall Through */
    case 3:
        w += z;
        break;
    case 5:
    case 6:
        w -= z;
        break;
    default:
        w = 2;
    }
    return w;
}
```

Switch Statement Example

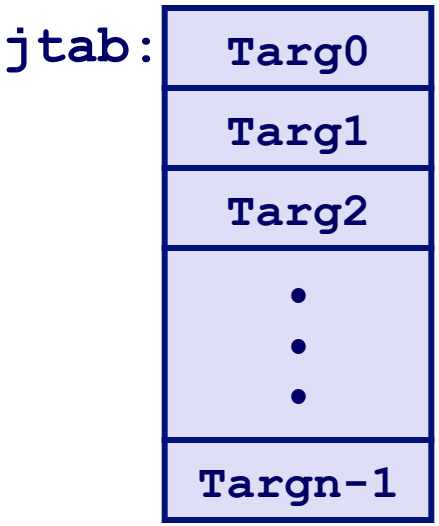
- Multiple case labels
  - Here: 5 & 6
- Fall through cases
  - Here: 2
- Missing cases
  - Here: 4

Jump Table Structure

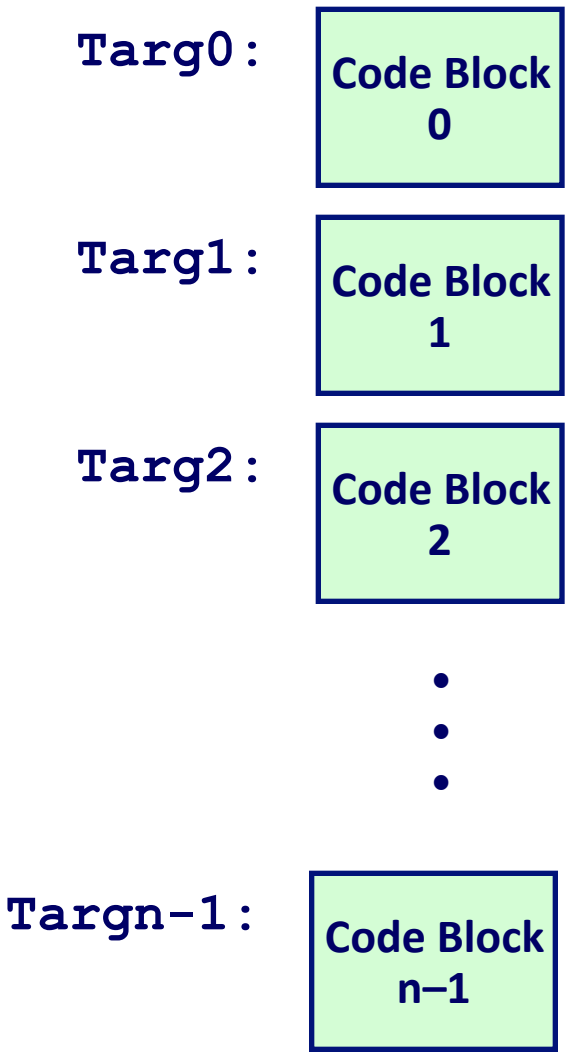
Switch Form

```
switch(x) {
  case val_0:
    Block 0
  case val_1:
    Block 1
    • • •
  case val_n-1:
    Block n–1
}
```

Jump Table



Jump Targets



Switch Statement Example

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
      . . .
    }
    return w;
}
```

Setup:

```
switch_eg:
    movq    %rdx, %rcx
    cmpq    $6, %rdi    # x:6
    ja      .L8
    jmp     *.L4(,%rdi,8)
```

What range of values takes default?

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Note that w not initialized here

Switch Statement Example

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
      . . .
    }
    return w;
}
```

Jump table

```
.section    .rodata
  .align 8
.L4:
  .quad     .L8  # x = 0
  .quad     .L3  # x = 1
  .quad     .L5  # x = 2
  .quad     .L9  # x = 3
  .quad     .L8  # x = 4
  .quad     .L7  # x = 5
  .quad     .L7  # x = 6
```

Setup:

```
switch_eg:
    movq    %rdx, %rcx
    cmpq    $6, %rdi        # x:6
    ja      .L8             # Use default
    jmp     *.L4(,%rdi,8)   # goto *JTab[x]
```

Indirect jump

Assembly Setup Explanation

- **Table Structure**
  - Each target requires 8 bytes
  - Base address at `.L4`

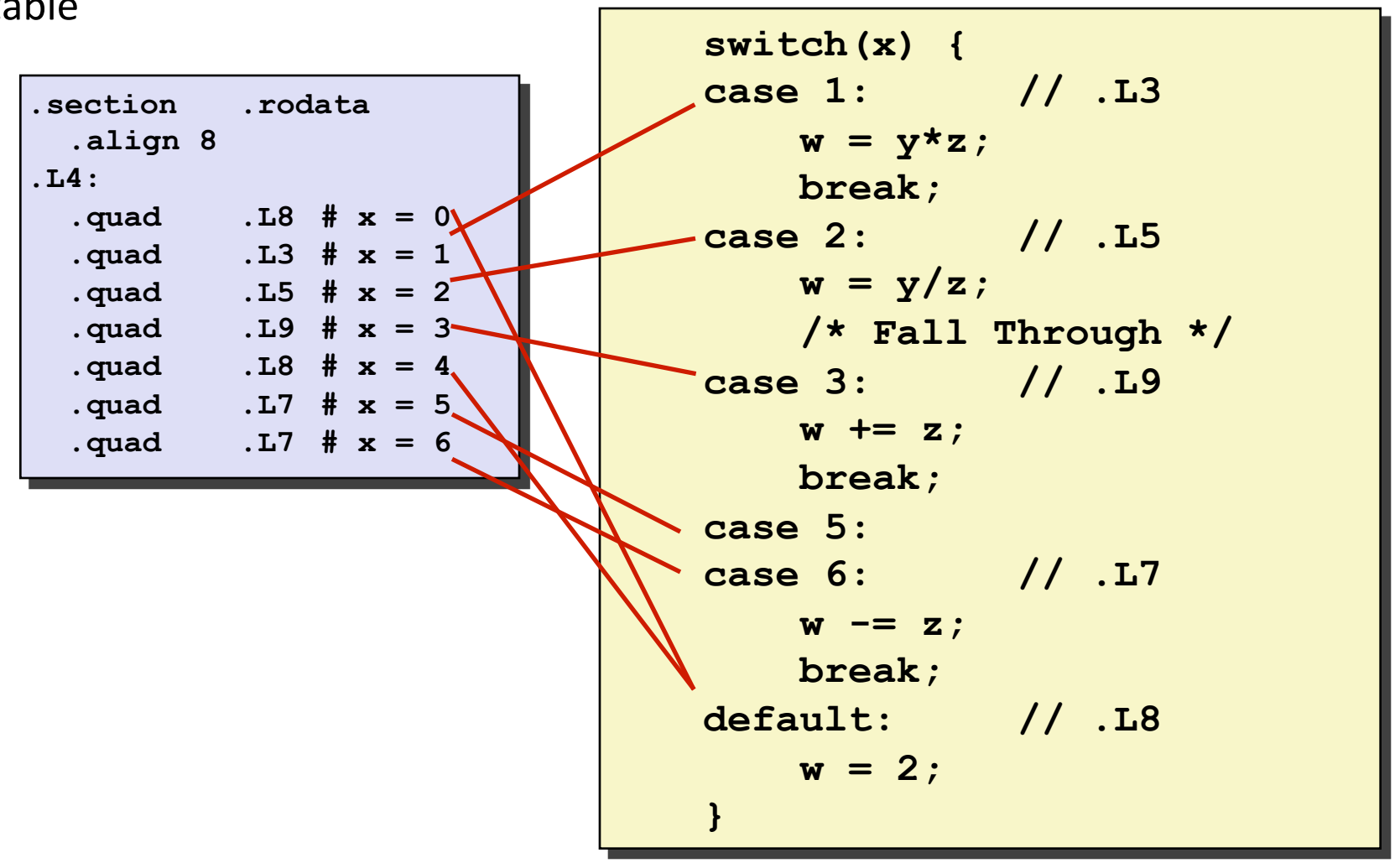
Jump table

```
.section    .rodata
  .align 8
.L4:
  .quad     .L8  # x = 0
  .quad     .L3  # x = 1
  .quad     .L5  # x = 2
  .quad     .L9  # x = 3
  .quad     .L8  # x = 4
  .quad     .L7  # x = 5
  .quad     .L7  # x = 6
```

- **Jumping**
  - Direct: `jmp .L8`
  - Jump target is denoted by label `.L8`
  - Indirect: `jmp *.L4(,%rdi,8)`
  - Start of jump table: `.L4`
  - Must scale by factor of 8 (addresses are 8 bytes)
  - Fetch target from effective Address `.L4 + x*8`
    - Only for $0 \le x \le 6$

Jump Table

Jump table



Code Blocks (x == 1)

```
  switch(x) {
  case 1:       // .L3
        w = y*z;
        break;
    . . .
}
```

```
.L3:
   movq     %rsi, %rax   # y
   imulq    %rdx, %rax   # y*z
   ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Handling Fall-Through

```
long w = 1;
   . . .
switch(x) {
   . . .
case 2:
    w = y/z;
    /* Fall Through */
case 3:
    w += z;
    break;
   . . .
}
```

```
case 2:
    w = y/z;
    goto merge;
```

```
case 3:
        w = 1;
```

```
merge:
        w += z;
```

Code Blocks (x == 2, x == 3)

```
    long w = 1;
       . . .
    switch(x) {
       . . .
    case 2:
        w = y/z;
        /* Fall Through */
    case 3:
        w += z;
        break;
       . . .
    }
```

```
.L5:                        # Case 2
   movq     %rsi, %rax
   cqto
   idivq    %rcx            #  y/z
   jmp      .L6             #  goto merge
.L9:                        # Case 3
   movl     $1, %eax        #  w = 1
.L6:                        # merge:
   addq     %rcx, %rax      #  w += z
   ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Code Blocks (x == 5, x == 6, default)

```
  switch(x) {
    . . .
    case 5:  // .L7
    case 6:  // .L7
        w -= z;
        break;
    default: // .L8
        w = 2;
}
```

```
.L7:                    # Case 5,6
  movl  $1, %eax        #  w = 1
  subq  %rdx, %rax      #  w -= z
  ret
.L8:                    # Default:
  movl  $2, %eax        #  2
  ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Supplementary Slides

Example: Left Shifting <<

- `shl` **or** `sal`
- **How are condition flags set?**
  - **shl/sal sets carry flag to last shifted out bit.**
    - **For unsigned, if CF=1, then overflow occurred.**
  - **for overflow flag OF,**
    - **if shift by one,**
      - » **if CF & MSbit identical after shift, OF = 0**
      - » **else OF = 1 (MS bit changes from 1->0 or 0->1),**

      i.e. OF = CF ^ MSbit
    - **else if shift by > 1, OF undefined.**
    - **For signed, if OF=1, then overflow occurred.**

Conditional Move Example

- Rewrite the absdiff example using conditional moves instead of conditional branching

```
int absdiff(
    int x, int y)
{
    int result;
    if (x > y) {
        result = x-y;
    } else {
        result = y-x;
    }
    return result;
}
```

```
absdiff: # x in %edi, y in %esi
  movl   %edi, %eax  # eax = x
  movl   %esi, %edx  # edx = y
  subl   %esi, %eax  # eax = x-y
  subl   %edi, %edx  # edx = y-x
  cmpl   %esi, %edi  # x<y?
  cmovl  %edx, %eax  # eax=edx if <=
  ret
```

- Note how the control flow is much easier to predict than all the jumping around with labels, e.g. `jle`, in the conditionally branched version of absdiff

Conditional Move Example

```
int absdiff(
    int x, int y)
{
    int result;
    if (x > y) {
        result = x-y;
    } else {
        result = y-x;
    }
    return result;
}
```

```
int cmovdiff(int x, int
y)
{
  int rval = x-y;
  int tval = y-x;
  int test = x < y;
  if (test) rval = tval;
  return rval;
}
```

Conditional Move Example

```
int cmovdiff(int x, int
y)
{
  int rval = x-y;
  int tval = y-x;
  int test = x < y;
  if (test) rval = tval;
  return rval;
}
```

```
absdiff: # x in %edi, y in %esi
  movl   %edi, %eax  # eax = x
  movl   %esi, %edx  # edx = y
  subl   %esi, %eax  # eax = x-y
  subl   %edi, %edx  # edx = y-x
  cmpl   %esi, %edi  # x<y?
  cmovl  %edx, %eax  # eax=edx if <=
  ret
```

Conditional Move Example

```
int cmovdiff(int x, int
y)
{
  int rval = x-y;
  int tval = y-x;
  int test = x < y;
  if (test) rval = tval;
  return rval;
}
```

```
absdiff: # x in %edi, y in %esi
  movl   %edi, %eax  # eax = x
  movl   %esi, %edx  # edx = y
  subl   %esi, %eax  # eax = x-y
  subl   %edi, %edx  # edx = y-x
  cmpl   %esi, %edi  # x<y?
  cmovl  %edx, %eax  # eax=edx if <=
  ret
```

Conditional Move Example

```
int cmovdiff(int x, int
y)
{
  int rval = x-y;
  int tval = y-x;
  int test = x < y;
  if (test) rval = tval;
  return rval;
}
```

```
absdiff: # x in %edi, y in %esi
  movl   %edi, %eax  # eax = x
  movl   %esi, %edx  # edx = y
  subl   %esi, %eax  # eax = x-y
  subl   %edi, %edx  # edx = y-x
  cmpl   %esi, %edi  # x<y?
  cmovl  %edx, %eax  # eax=edx if <=
  ret
```

Conditional Move Example

```
int cmovdiff(int x, int
y)
{
  int rval = x-y;
  int tval = y-x;
  int test = x < y;
  if (test) rval = tval;
  return rval;
}
```

```
absdiff: # x in %edi, y in %esi
  movl   %edi, %eax  # eax = x
  movl   %esi, %edx  # edx = y
  subl   %esi, %eax  # eax = x-y
  subl   %edi, %edx  # edx = y-x
  cmpl   %esi, %edi  # x<y?
  cmovl  %edx, %eax  # eax=edx if <=
  ret
```

- Control flow is more predictable, but both branches must be evaluated

General Form with Conditional Move

C Code

```
val = Test ? Then-Expr : Else-Expr;
```

Conditional Move Version

```
val1  = Then-Expr;
val2  = Else-Expr;
val1  = val2 if !Test;
```

- Both values get computed
- Overwrite then-value with else-value if condition doesn't hold
- Don't use when:
  - Then or else expression have side effects, like dereferencing a null pointer or incrementing a global variable (then & else expressions always evaluated)
  - Then and else expressions are too expensive

Jump-to-Middle "While" Loop Translation

Goto Version #2

```
int fact_while_goto2(int x)
{
  int result = 1;
  if (!(x > 1))
    goto done;
loop:
  result *= x;
  x = x-1;
  if (x > 1)
    goto loop;
done:
  return result;
}
```

Goto Version #3

```
int fact_while_goto3(int x)
{
  int result = 1;
  goto middle;
loop:
  result *= x;
  x = x-1;
middle:
  if (x > 1)
    goto loop;
  return result;
}
```

- Recent technique for GCC
  - Both IA32 & x86-64
- First iteration jumps over body computation within loop

Jump-to-Middle While Loop Translation

C Code

```
while (Test)
    Body
```


Goto Version

```
goto middle;
loop:
  Body
middle:
  if (Test)
    goto loop;
```

- Avoids duplicating test code
- Unconditional `goto` incurs no performance penalty
- `for` loops compiled in similar fashion

Goto (Previous) Version

```
  if (!Test)
    goto done;
loop:
  Body
  if (Test)
    goto loop;
done:
```

Jump-to-Middle Assembly Example

Goto Version #3

```
int
fact_while_goto3(int
x)
{
  int result = 1;
  goto middle;
loop:
  result *= x;
  x = x-1;
middle:
  if (x > 1)
    goto loop;
  return result;
}
```

```
# x in %edx, result in %eax
  jmp    .L34          #   goto Middle
.L35:                  # Loop:
  imull %edx, %eax     #   result *= x
  decl   %edx          #   x--
.L34:                  # Middle:
  cmpl   $1, %edx      #   x:1
  jg     .L35          #   if >, goto Loop
```

C operators – Assembly Equivalents?

Operators	Associativity
`() [] -> .`	left to right
`! ~ ++ -- + - * & (type) sizeof`	right to left
`* / %`	left to right
`+ -`	left to right
`<< >>`	left to right
`< <= > >=`	left to right
`== !=`	left to right
`&`	left to right
`^`	left to right
`|`	left to right
`&&`	left to right
`||`	left to right
`?:`	right to left
`= += -= *= /= %= &= ^= != <<= >>=`	right to left
`,`	left to right

Note: Unary +, –, and * have higher precedence than binary forms