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50005 Lab 2 Report

Q2

```
(base) dylantan@Dylans-MacBook-Pro StarterCode_C % ./q2 testcases/q2_1.txt
Customer 0 requesting
0 1 0
Customer 1 requesting
2 0 0
Customer 2 requesting
3 0 2
Customer 3 requesting
2 1 1
Customer 4 requesting
0 0 2
Customer 1 requesting
1 0 2

Current state:
Available:
2 3 0
Maximum:
7 5 3
3 2 2
9 0 2
2 2 2
4 3 3

Allocation:
0 1 0
3 0 2
3 0 2
2 1 1
0 0 2

Need:
7 4 3
0 2 0
6 0 0
0 1 1
4 3 1
Customer 0 requesting
0 2 0

Current state:
Available:
2 3 0
Maximum:
7 5 3
3 2 2
9 0 2
2 2 2
4 3 3

Allocation:
0 1 0
3 0 2
3 0 2
2 1 1
0 0 2

Need:
7 4 3
0 2 0
6 0 0
0 1 1
4 3 1
```

```
(base) dylantan@Dylans-MacBook-Pro StarterCode_C % ./checkq1
For Q1: You have scored 1/1
(base) dylantan@Dylans-MacBook-Pro StarterCode_C % ./checkq2
For Q2: You have scored 1/1
```

Q3

There are 2 parts to the banker's algorithm:

- 1) Resource request algorithm $\rightarrow O(m)$
- 2) Safety algorithm $\rightarrow O(mn^2)$

where m is the number of resources and n is the number of customers.

Resource request algorithm:

The print statements, return statement, and if-statement takes $O(1)$ time if we exclude the safety algorithm.

```
printf("Customer %d requesting\n", customerIndex);
printf("\n"); // Leave a line after each customer
if (checkSafe(customerIndex, request) == 0){
    return 0;
}
return 1;
```

The 3 for-loops takes $O(m)$ time each in the worst case.

```
for (int i=0; i<numberOfResources; i++){
    printf("%d ", request[i]); // Leave a space between each request
}

for (int i=0; i<numberOfResources; i++){
    if (request[i] > need[customerIndex][i]){
        printf("Error!!!\n");
        return 0;
    }
    if (request[i] > available[i]){
        printf("Rejected!!!\n");
        return 0;
    }
}

// TODO: request is granted, update state
for (int i=0; i<numberOfResources; i++){
    available[i] -= request[i];
    need[customerIndex][i] -= request[i];
    allocation[customerIndex][i] += request[i];
}
```

Therefore, time complexity is $O(3m + 1)$, which is just $O(m)$.

Safety Algorithm:

Allocation of space:

- 1) work $\rightarrow O(m)$
- 2) tempNeed and tempAllocation $\rightarrow O(nm)$
- 3) Finish $\rightarrow O(n)$.

The declarations and return statement has $O(1)$ time complexity.

Therefore, it takes $O(m+nm+n+1)$, which is just $O(mn)$.

```
int *work = mallocIntVector(numberOfResources);
int **tempNeed = mallocIntMatrix(numberOfCustomers, numberOfResources);
int **tempAllocation = mallocIntMatrix(numberOfCustomers, numberOfResources);

int *Finish = mallocIntVector(numberOfCustomers);
int safeState = 1;
int needState;

return 1;
```

Setting of temp arrays uses a nested for-loop and has a time complexity of $O(mn)$.

```
for (int i=0; i<numberOfCustomers; i++)
{
    for (int j=0; j<numberOfResources; j++){
        tempNeed[i][j] = need[i][j];
        tempAllocation[i][j] = allocation[i][j];
    }
}
```

Temporarily granting the request uses a for-loop with time complexity of $O(m)$.

```
for (int i=0; i<numberOfResources; i++)
{
    work[i] = available[i] - request[i];
    tempNeed[customerIndex][i] = need[customerIndex][i] - request[i];
    tempAllocation[customerIndex][i] = allocation[customerIndex][i] + request[i];
}
```

Initialization and checking of all slots for the Finish array takes $O(n)$ time each.

```
for (int i=0; i<numberOfCustomers; i++)
{
    Finish[i] = 0;
}

for (int i=0; i<numberOfCustomers; i++){
    if (Finish[i] == 0){
        return 0;
    }
}
```

The worst-case scenario for the while-loop would be when only 1 slot of the Finish array is set to 1 each loop. It will thus loop for n times just to set all the slots in the Finish array to 1, taking $O(n)$ time.

Outermost for-loop takes $O(n)$ to loop through all customers.

Innermost for-loops takes $O(m)$ to loop through all resources.

Therefore, time complexity is $O(n \cdot n \cdot m)$, which is $O(mn^2)$.

```
while (safeState){
    safeState = 0;
    for (int i=0; i<numberOfCustomers; i++)
    {
        needState=1;
        if (Finish[i] == 0)
        {
            for (int j=0; j<numberOfResources; j++)
            {
                if (tempNeed[i][j] > work[j]){
                    needState = 0;
                }
            }

            if (needState == 1)
            {
                safeState = 1;
                for (int j=0; j<numberOfResources; j++)
                {
                    work[j] += tempAllocation[i][j];
                }
                Finish[i] = 1;
            }
        }
    }
}
```

Freeing:

- 1) finish and work $\rightarrow O(1)$
- 2) tempAllocation and tempNeed $\rightarrow O(n)$ time
- 3) Therefore, the whole freeing process takes $O(n)$ time.

```
// Free memory
freeIntVector(Finish);
freeIntMatrix(tempAllocation);
freeIntMatrix(tempNeed);
freeIntVector(work);
```

Therefore, the combined time complexity for the safety algorithm is $O(2mn + m + 3n + mn^2)$, which is $O(mn^2)$.

In conclusion, the whole banker's algorithm takes $O(m+mn^2)$, which is $O(mn^2)$.