Improving Software Quality using Automatic Invariant Discovery and Program Repair

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Automated program analysis techniques and tools can decrease debugging time by an average of 26% and \$41 billion annually

Program Verification





Check if a program meets a given specification

Program Repair





Fix a buggy program to satisfy a given specification

Invariant Generation

```
def intdiv(x, y):
    q = 0
    r = x
    while r \geq y:
    a = 1
    b = y
    while [??] r \geq 2b:
    a = 2a
    b = 2b
    r = r - b
    q = q + a
    [??]
    return q
```

- Discover invariant properties at certain program locations
- Answer the question "what does this program do?"

Automatic Program Repair

return q

- Localize errors and modify code to fix bugs
- A form of program synthesis

Outline

Research

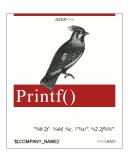
Invariant Generation

Automatic Program Repair

• Current/New Research Works

4

How We Analyze Programs



```
File Edit Options Buffers Tools Python Help

def intdiv(x, y):
    assert y != 0

# .. compute result ..

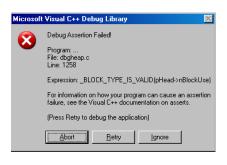
assert r >= 0
assert x >= q

return q,r
```



Software Testing course

"GCC: 9000 assertions, LLVM: 13,000 assertions [..] 1 assertion per 110 loc" "program invariants are asserted properties, such as relations among variables, at certain locations in a program"



Uses

- verify programs
- help understand programs
- reveal interesting, unexpected properties
- debug (locate errors)
- documentations
- ..

Approaches to Finding Invariants

```
int intdiv(int x, int y){
  int q=0; int r=x;
  while(r ≥ y){
    int a=1; int b=y;
    while[L](r ≥ 2*b){
        a = 2*a; b = 2*b;
    }
    r=r-b; q=q+a;
}
return q;
}
```

X	у	q	r
0	1	0	0
1	1	1	0
3	4	0	3
8	1	8	0
15	5	3	0
20	2	10	0
100	1	100	0
	:	:	

Static Analysis

- Analyze source code directly
- Pros: guaranteed results
- Cons: computationally intensive, infer simple invariants

Dynamic Analysis

- Analyze program traces
- Pros: fast, source code not required
- Cons: results depend on traces, might not hold for all runs

Example

```
int intdiv(int x, int y){
 assert(x>0 && y>0);
 int q=0; int r=x;
 while(r \ge y){
   int a=1;
   int b=y;
   while [L] (r \geq 2*b) {
     a = 2*a;
     b = 2*b;
   r=r-b;
   q=q+a;
 return q;
```

```
        x
        y
        a
        b
        q
        r

        15
        2
        1
        2
        0
        15

        15
        2
        2
        4
        0
        15

        15
        2
        1
        2
        4
        7

        4
        1
        1
        1
        0
        4

        4
        1
        2
        2
        0
        4
```

Invariants at L: b = ya, x = qy + r, $r \ge 2ya$

DIG discovers polynomial relations of the forms

Equalities
$$c_0 + c_1 x_1 + c_2 x_n + c_3 x_1 x_2 + \dots + c_m x_1^{d_1} \dots x_n^{d_n} = 0$$

Inequalities $c_0 + c_1 x_1 + c_2 x_n + c_3 x_1 x_2 + \dots + c_m x_1^{d_1} \dots x_n^{d_n} \ge 0, \quad c_i \in \mathbb{R}$

Examples

cubic
$$z-6n=6$$
, $\frac{1}{12}z^2-y-\frac{1}{2}z=-1$ extended gcd $\gcd(a,b)=ia+jb$
$$\operatorname{sqrt} \quad x+\varepsilon \geq y^2 \geq x-\varepsilon$$

Method

- Equalities: solve equations
- Inequalities: construct polyhedra

Example

```
int intdiv(int x, int y){
 assert(x>0 && y>0);
 int q=0; int r=x;
 while(r \ge y){
   int a=1;
   int b=y;
   while [L] (r \geq 2*b) {
     a = 2*a;
     b = 2*b;
   r=r-b;
   q=q+a;
 return q;
```

```
    x
    y || a
    b
    q
    r

    15
    2 || 1
    2
    0
    15

    15
    2 || 2
    4
    0
    15

    15
    2 || 1
    2
    4
    7

    4
    1 || 1
    1
    0
    4

    4
    1 || 2
    2
    0
    4
```

Invariants at L: b = ya, x = qy + r, $r \ge 2ya$

Finding Nonlinear Equations using Linear Equation Solving

Terms and degrees

$$V = \{r, y, a\}; \text{ deg} = 2$$

$$\downarrow$$

$$T = \{1, r, y, a, ry, ra, ya, r^2, y^2, a^2\}$$

$$T = \{\dots, \log(r), a^y, \sin(y), \dots\}$$

Nonlinear equation template

$$c_1 + c_2 r + c_3 y + c_4 a + c_5 r y + c_6 r a + c_7 y a + c_8 r^2 + c_9 y^2 + c_{10} a^2 = 0$$

System of linear equations

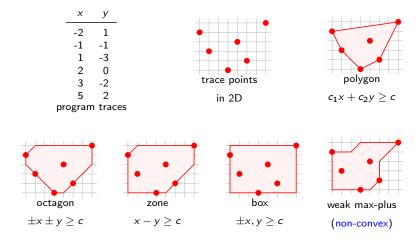
trace 1 :
$$\{r=15, y=2, a=1\}$$

eq 1 : $c_1+15c_2+2c_3+c_4+30c_5+15c_6+2c_7+225c_8+4c_9+c_{10}=0$
:

Solve for coefficients c;

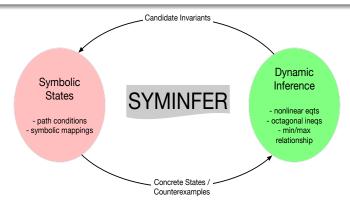
$$V = \{x, y, a, b, q, r\}; \text{ deg} = 2 \longrightarrow b = ya, x = qy + r$$

Geometric Invariant Inference



Symbolic States: logical formulae representing program semantics

- path conditions and mappings from variables to symbolic values
- obtained using a symbolic execution tool



- Dynamic Inference: learn *candidate* invariants under different templates from concrete states (execution traces)
- Symbolic States: check invariants and return counterexample states

Application: Complexity Analysis

```
void triple(int M, int N, int P){ Complexity of this program?
  assert (0 \le M):
  assert (0 \le N):
  assert (0 \le P);
  int i = 0, j = 0, k = 0;
  int t = 0;
  while(i < N){</pre>
    j = 0; t++;
    while(j < M){</pre>
      j++; k = i; t++;
      while (k < P){
       k++; t++;
     i = k;
    i++:
  [L]
```

- - Use t to count loop iterations
 - At first glance: t = O(MNP)
 - SymInfer found an interesting (and unexpected) nonlinear invariant at L:

```
P^2Mt + PM^2t - PMNt - M^2Nt -
PMt^2 + MNt^2 + PMt - PNt - 2MNt +
Pt^2 + Mt^2 + Nt^2 - t^3 - Nt + t^2 = 0
```

• Solve for t yields the most precise, unpublished bound:

```
t = 0
    when N=0,
t = P + M + 1 when N < P,
t = N - M(P - N) when N > P
```

Results and Applications

Results

- ICSE'14 Distinguish Paper award
- generates correct and sufficiently strong invariants for all 28 NLA programs in SV-COMP

Applications

- complexity and side-channel attacks (FSE'17, ASE'17, SEAD Workshop'20)
- AES analysis (ICSE'12, TOSEM'13)
- termination/liveness (OOPSLA'20)
- heap/pointer (PLDI'19)
- disjunctive/geometric invs (ICSE'14, J. Automated Reasoning'13)
- interactions in highly-configurable systems (FSE'16, ICSE'21)

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Zune Bug



Wed morning, Dec 31, 2008: Microsoft Zune music players mysteriously froze

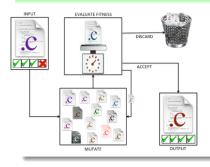
– Matt Akers (Microsoft Zune spokesman)

"By [Thursday] you should allow the battery to fully run out of power before the unit can restart successfully, then simply ensure that your device is recharged, then turn it back on"

Zune Bug

```
int zunebug(int days) {
                                   int zunebug_repair(int days) {
 int year = 1980;
                                     int year = 1980;
 while (days > 365) {
                                     while (days > 365) {
   if (isLeapYear(year)){
                                       if (isLeapYear(year)){
     if (days > 366) {
                                        if (days > 366) {
       days -= 366;
                                          // days -= 366; // repair deletes
      year += 1;
                                          year += 1;
                                        days -= 366; // repair inserts
   else {
                                       } else {
     days -= 365;
                                        days -= 365;
     year += 1;
                                        year += 1;
 return year;
                                     return year;
```

GenProg: Program Repair using Genetic Algorithm



- Isolate faults
- 2 Mutate program statements and reuse existing code
- Check repair candidates

Results

- demonstrated on bugs in real-world software (repair 16 programs over 1.25 MLocs, 2 mins avg)
- 10-year Most Influential Paper award (ICSE '19) and 10-year Most Impact Paper award (GECCO '19)

From Reachability to Synthesis

Equivalance Thm: Template-based Synthesis

Program Reachablity

```
if i: b = c_0 + c_1 \times u + c_2 \times d \text{ #syn}
template
else: b = u
if (b > d): r = 1
else: r = 0
return r
C(1, 0, 100) = 0
Q(1, 11, 110) = 1
Q(0, 100, 50) = 1
Q(1, -20, 60) = 1
Q(0, 0, 10) = 0
Q(0, 0, -10) = 1
```

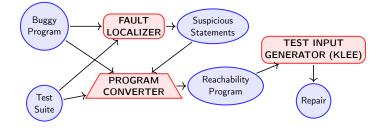
def Q(i, u, d):

Input: a synthesis instance

```
def p<sub>0</sub>(i, u, d, c<sub>0</sub>, c<sub>1</sub>, c<sub>2</sub>):
 if i:
    b = c_0 + c_1 \times u + c_2 \times d
 else: b = u
 if b > d: r = 1
  else: r = 0
  return r
def p_{main}(c_0, c_1, c_2):
   e = p_0(1, 0, 100, c_0, c_1, c_2) == 0 and
       p_0(1,11,110, c_0,c_1,c_2) == 1 and
        p_0(0.100.50, c_0.c_1.c_2) == 1 and
        p_0(1,-20,60, c_0,c_1,c_2) == 1 and
        p_0(0, 0, 10, c_0, c_1, c_2) == 0 and
       p_0(0, 0, -10, c_0, c_1, c_2) == 1
   if e:
       [L] #pass the given test suite
   return 0
```

Output: a reachability instance, solvable using a test-input generation tool

CETI: Correcting Errors using Test Inputs (TACAS'17)



Non-traditional Repairs

- Corrupted data structures (Google Summer of Code'18, FSE JPF Workshop'18)
- Fault localization in declarative models (ICSE '21)
- Repair declarative programs (ICSE '21)

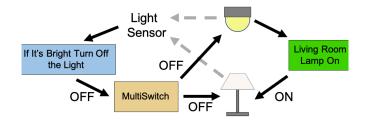
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Research

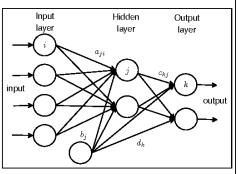
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IoT Interaction Analysis and Repair (UNL Faculty grant'21)



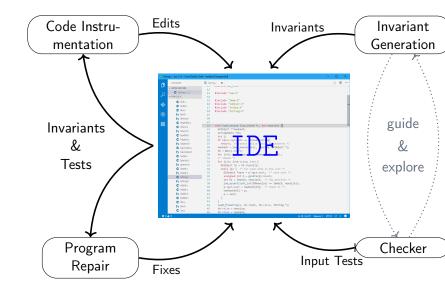
Deep Neural Networks



```
\begin{array}{l} a[1,\ldots,h][1,\ldots,n] \triangleright \text{Input weights} \\ c[1,\ldots,m][1,\ldots,h] \triangleright \text{Output weights} \\ b[1,\ldots,h] & \triangleright \text{Hidden nodes' bias} \\ d[1,\ldots,m] & \triangleright \text{Output nodes' bias} \\ \end{array} \begin{array}{l} \text{FUNCTION } \nu(x) \\ \text{for } j \leftarrow 1 \text{ to } h \text{ do} \\ r_j \leftarrow 0; \\ \text{for } i \leftarrow 1 \text{ to } n \text{ do} \\ r_j \leftarrow r_j + a_{ji} \cdot x_i + b_j \\ \text{for } k \leftarrow 1 \text{ to } m \text{ do} \\ s_k \leftarrow 0; \\ \text{for } j \leftarrow 1 \text{ to } h \text{ do} \\ s_k \leftarrow s_k + c_{kj} \cdot \sigma_h(r_j) + d_k \\ y_k \leftarrow \sigma_o(s_k) \\ \end{array}
```

- Invariants (activation patterns) discovery
- Symbolic testing

IDE Integration (https://grammatech.gitlab.io/Mnemosyne/docs/muses/)



Thank you!