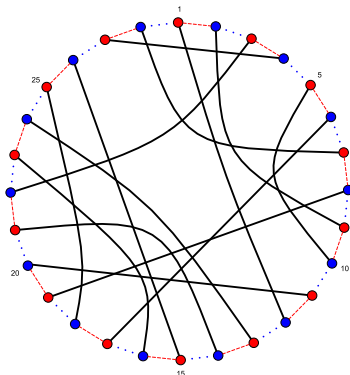


Book Embedding with Fixed Page Assignments



Daniel Hoske, 20th November 2012

Department of Informatics, Institute of Theoretical Computer Science

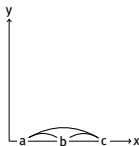
Advisors: Dipl.-Inform. Thomas Bläsius, Dr. Ignaz Rutter

Page embedding

Definition

Page embedding is planar embedding with

- vertices on a line and
- edges in half-plane above the line



Page embeddable = outerplanar

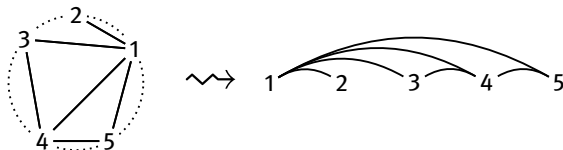
Page embedding

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Page embedding is planar embedding with

- vertices on a line and
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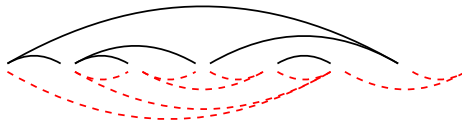
Page embeddable = outerplanar



Book embedding

Definition

Book embedding of $G_i = (V, E_i)$, $i \in \{1, \dots, k\}$ consists of page embeddings for G_i with the same vertex positions.



Book embedding

Definition

Book embedding of $G_i = (V, E_i)$, $i \in \{1, \dots, k\}$ consists of page embeddings for G_i with the same vertex positions.

Problem: BOOK-EMBEDDING

Given: Vertex set V and edge sets $E_1, \dots, E_k \subseteq \binom{V}{2}$.

Question: Is there a book embedding of (V, E_i) ?

$k = 1$: embeddable = outerplanar

$k = 2$: decidable in $\mathcal{O}(n)$ [Hong and Nagamochi, 2009]

What happens for $k = 3$?

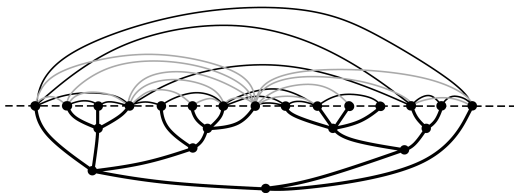
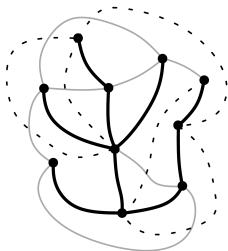
Motivation

Problem: CONNECTED-SEFE

Given: Two graphs G_1 and G_2 on V where $G_1 \cap G_2$ is connected.

Question: Are there planar embeddings of G_1 and G_2 that coincide on $G_1 \cap G_2$?

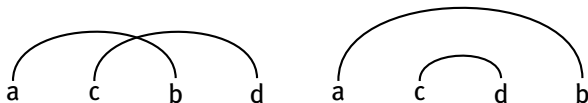
is equivalent to 2-page book embedding + a tree [Angelini et al., 2012]



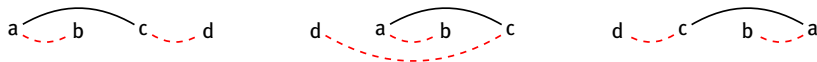
Observations

BOOK-EMBEDDING is an ordering problem:

Avoid suborder $a < c < b < d$ for $\{a, b\}, \{c, d\} \in E_i$

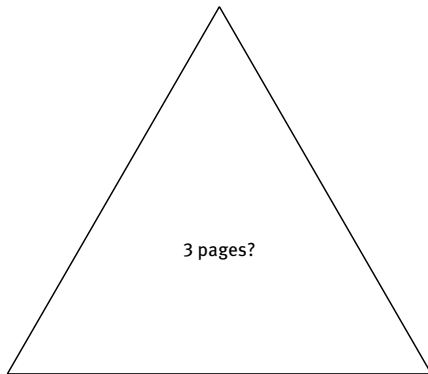


Mirror image and cyclic shifts of a valid order remain valid:



Results

2 pages: $\mathcal{O}(n)$



Contents

- 1 NP-completeness and connected pages
- 2 Disjoint perfect matchings
- 3 Tree on the vertices

Contents

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BOOK-EMBEDDING is NP-complete

Theorem

BOOK-EMBEDDING with matchings as pages is NP-complete.

Reduce from NP-complete problem BETWEENNESS [Opatrny, 1979].

Problem: BETWEENNESS

Given: Finite set $M := \{1, \dots, n\}$ and ordered triples $C \subseteq M^3$.

Question: Is there a total ordering $<$ of M such that $a < b < c$ or $a > b > c$ for all $(a, b, c) \in C$?

Reduction from BETWEENNESS

Problem: BETWEENNESS

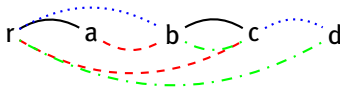
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Map triple $(a, b, c) \in C$ to two new pages (r fixed new vertex):



For example: $(a, b, c), (b, c, d)$



Reduction from BETWEENNESS

Problem: BETWEENNESS

Given: Finite set $M := \{1, \dots, n\}$ and ordered triples $C \subseteq M^3$.

Question: Is there a total ordering $<$ of M such that $a < b < c$ or $a > b > c$ for all $(a, b, c) \in C$?

Map triple $(a, b, c) \in C$ to two new pages (r fixed new vertex):



BETWEENNESS \Rightarrow BOOK-EMBEDDING:

Take r as first vertex

\rightsquigarrow Orders $r < a < b < c$ or $r < c < b < a$ are valid

Reduction from BETWEENNESS

Problem: BETWEENNESS

Given: Finite set $M := \{1, \dots, n\}$ and ordered triples $C \subseteq M^3$.

Question: Is there a total ordering $<$ of M such that $a < b < c$ or $a > b > c$ for all $(a, b, c) \in C$?

Map triple $(a, b, c) \in C$ to two new pages (r fixed new vertex):



BOOK-EMBEDDING \Rightarrow BETWEENNESS:

Rotate r to front

$\rightsquigarrow r < a < b < c$ or $r < c < b < a$ are the only valid orders

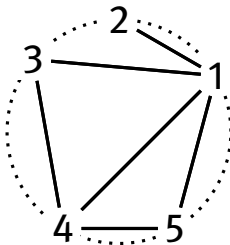
Connected pages

Theorem

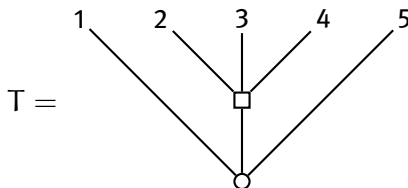
BOOK-EMBEDDING with connected pages can be solved in $\mathcal{O}(kn)$ time.

Idea:

Compute valid orders $\pi_i \subseteq \text{Sym}(n)$ for single pages and intersect them.

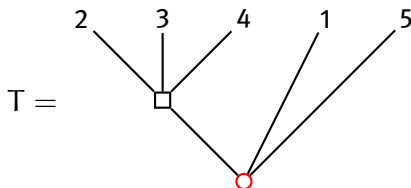


Interludium: PQ-trees



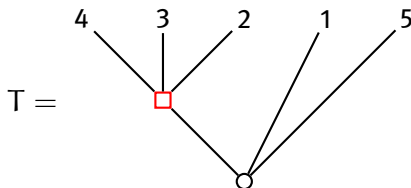
$$\pi(T) = \{\mathbf{12345}, 14325, 52341, 54321, 15234, 15432, \\ 51234, 51432, 23415, 43215, 23451, 43251\}$$

Interludium: PQ-trees



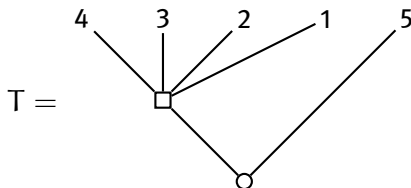
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Interludium: PQ-trees



$$\pi(T) = \{12345, 14325, 52341, 54321, 15234, 15432, \\ 51234, 51432, 23415, 43215, 23451, 43251\}$$

Permutations with the leaves in $\{1, 2\}$ adjacent

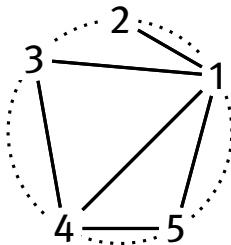
Connected pages

Theorem

BOOK-EMBEDDING with connected pages can be solved in $\mathcal{O}(kn)$ time.

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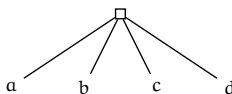
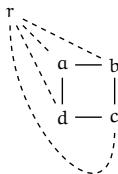


Connected pages

Theorem

BOOK-EMBEDDING with connected pages can be solved in $\mathcal{O}(kn)$ time.

- Construct PQ-trees T_i on V representing valid orders of individual pages (V, E_i) $\mathcal{O}(n)$



[Booth and Lueker et. al. 1976, Shih and Hsu 1993, Boyer and Myrvold 1999]

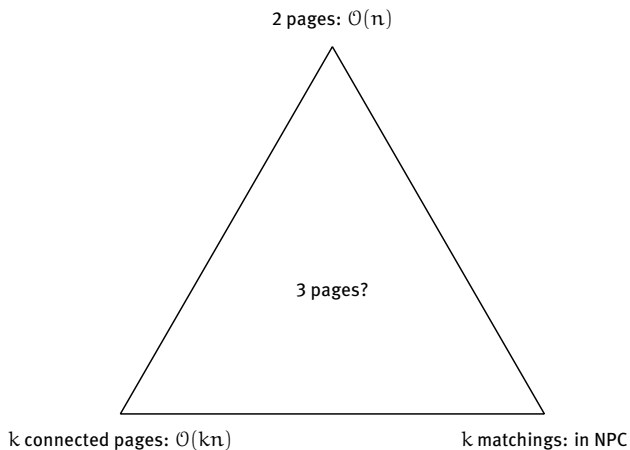
Connected pages

Theorem

BOOK-EMBEDDING with connected pages can be solved in $\mathcal{O}(kn)$ time.

- Construct PQ-trees T_i on V representing valid orders of individual pages (V, E_i) $\mathcal{O}(n)$
 - Intersect the T_i [Booth, 1975] $\mathcal{O}(n)$
- Resulting PQ-tree T represents valid book orders

Results



Contents

- 1 NP-completeness and connected pages
- 2 Disjoint perfect matchings
- 3 Tree on the vertices

Disjoint perfect matchings as pages

Problem: PERFECT-MATCHINGS-BOOK-EMBEDDING

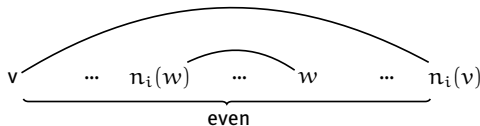
Given: Disjoint perfect matchings E_1, \dots, E_k on a vertex set V .

Question: Is there a book embedding of (V, E_i) ?

Theorem

Necessary: $G := (V, E_1 \cup \dots \cup E_k)$ is *bipartite*.

- Even number of vertices between adjacent vertices in valid order



Disjoint perfect matchings as pages

Problem: PERFECT-MATCHINGS-BOOK-EMBEDDING

Given: Disjoint perfect matchings E_1, \dots, E_k on a vertex set V .

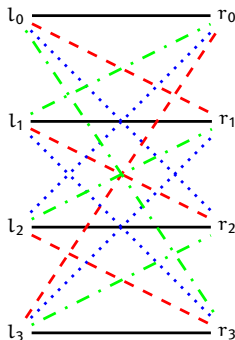
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Theorem

Necessary: $G := (V, E_1 \cup \dots \cup E_k)$ is *bipartite*.

- Even number of vertices between adjacent vertices in valid order
- ⇒ Vertices with even and odd indexes form bipartition

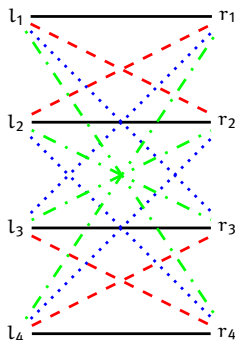
Bipartite examples



k pages: Take the partition

$$E_i := \{\{l_j, r_{(j+i) \bmod k}\} : j \in \{1, \dots, k-1\}\} \text{ of } K_{k,k}$$

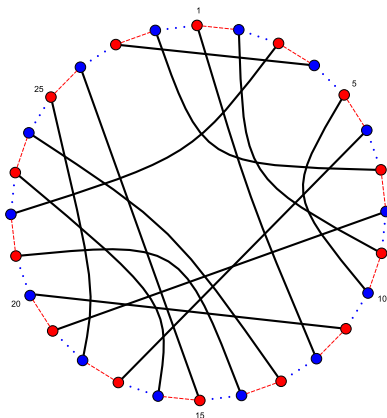
Bipartite counterexamples



$k \geq 4$ pages: Partition $K_{k,k}$ into disjoint perfect matchings that contain this counterexample

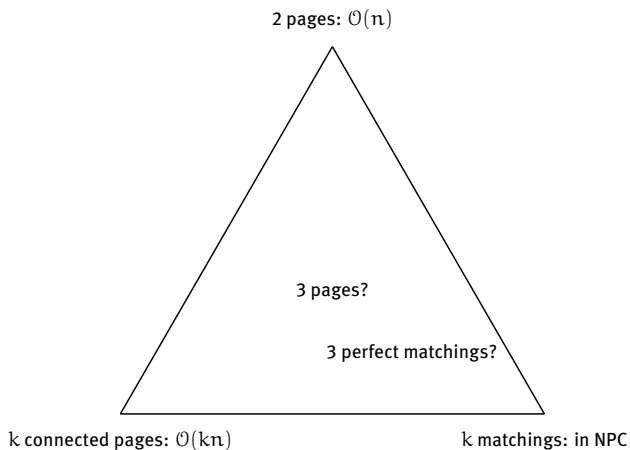
Bipartite counterexample for three pages

Smallest counterexample where two of the matchings form a cycle



Smallest unrestricted counterexample: $20 \leq n \leq 28$

Results

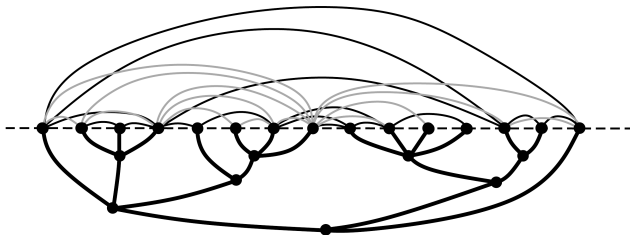


Contents

- 1 NP-completeness and connected pages
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- 3 Tree on the vertices**

Tree on the vertices

BOOK-EMBEDDING + tree:



Drawing the tree = Restricting permutations by a P-tree

Tree on the vertices

BOOK-EMBEDDING + tree:

Problem: P-TREE-BOOK-EMBEDDING

Given: BOOK-EMBEDDING instance I and a P-tree T with leaves V .

Question: Is there a total order $< \in \pi(T)$ solving I ?

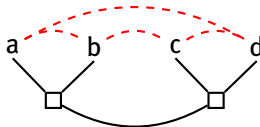
More structure (same for binary trees):

Problem: Q-TREE-BOOK-EMBEDDING

Given: BOOK-EMBEDDING instance I and a Q-tree T with leaves V .

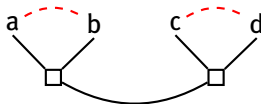
Question: Is there a total order $< \in \pi(T)$ solving I ?

Example



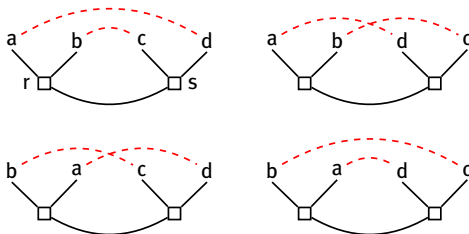
- What happens to the forbidden suborder constraint?
- ⇒ 2-CNF formula (Boolean equations) on orientation of Q-nodes

Example



- $\{a, b\}, \{c, d\}$: true

Example



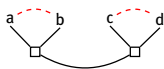
- $\{a, d\}, \{b, c\}$: $a < b \Leftrightarrow c < d$
- Fix reference orientation of inner nodes r
- Boolean variable o_r for being in reference orientation

$$\Rightarrow o_r \Leftrightarrow o_s$$

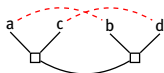
Book constraints to 2-CNF formula

Theorem

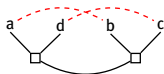
Q-TREE-BOOK-EMBEDDING is solvable in $\mathcal{O}(kn^2)$ time.



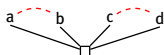
true



$\mathcal{O}_r(a,c) \Leftrightarrow \mathcal{O}_r(b,d)$ or $\mathcal{O}_r(a,c) \Leftrightarrow \neg \mathcal{O}_r(b,d)$



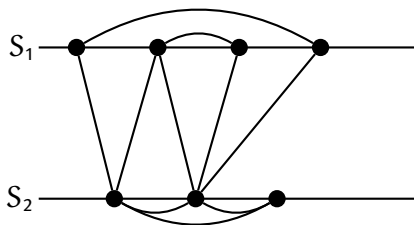
$\mathcal{O}_r(a,d) \Leftrightarrow \mathcal{O}_r(b,c)$ or $\mathcal{O}_r(a,d) \Leftrightarrow \neg \mathcal{O}_r(b,c)$



true or false

Multiple spines

Take multiple spines:



Without caps: Level planarity solvable in linear time.
[Jünger et al., 1999]

Level planarity

Level planarity is an ordering problem:

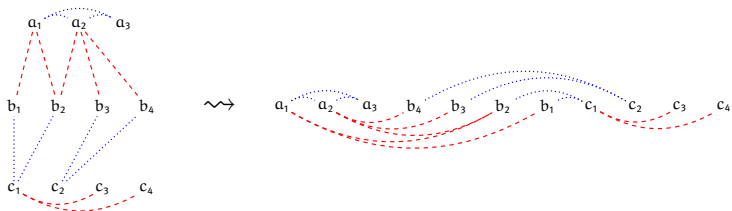


Forbidden: $a_1 <_i a_2 \wedge b_2 < b_1$ for $(a_1, b_1), (a_2, b_2) \in E_i$

Mapping to 2-page P-TREE-BOOK-EMBEDDING

Theorem

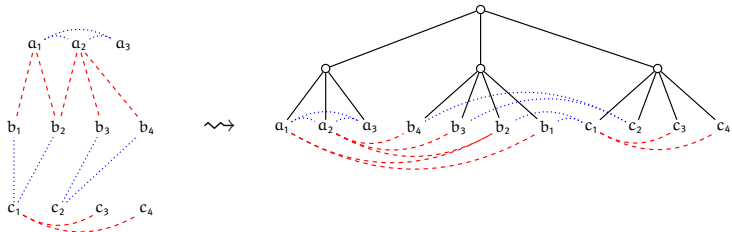
A MULTIPLE-SPINE-EMBEDDING instance is equivalent to a special 2-page P-TREE-BOOK-EMBEDDING instance.



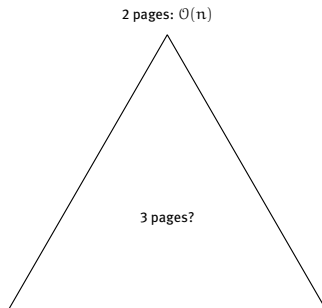
Mapping to 2-page P-TREE-BOOK-EMBEDDING

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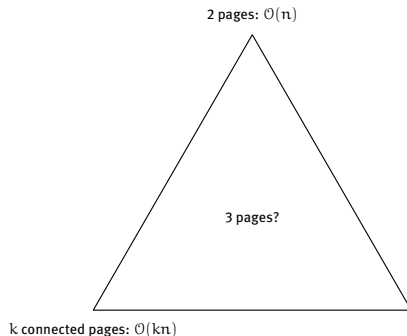
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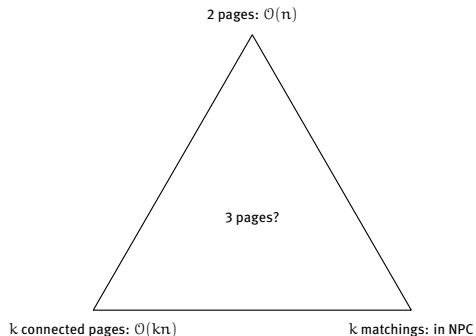
Conclusion and outlook



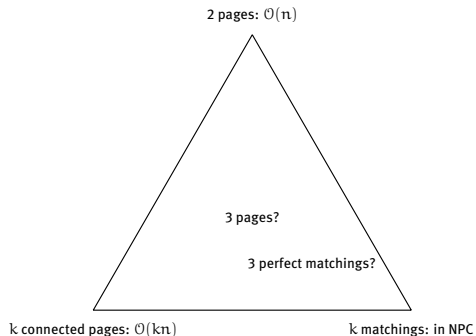
Conclusion and outlook



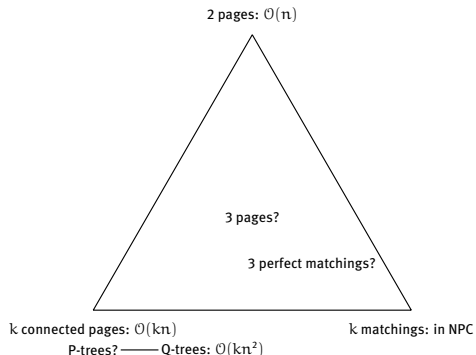
Conclusion and outlook



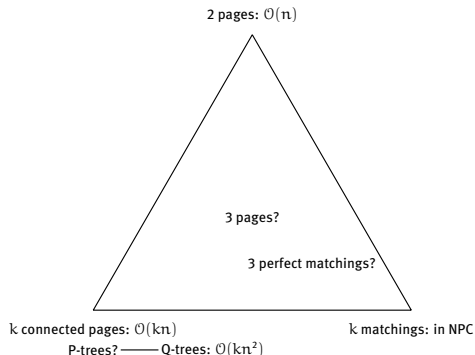
Conclusion and outlook



Conclusion and outlook



Conclusion and outlook



- Complexity for constant number of pages?
- Complexity when constrained by a P-tree?