EECS 490 – Lecture 15

Object-Oriented Programming

Announcements

- Project 3 due tomorrow at 8pm
- Midterm Tuesday 10/31 during class time
 - Will be in 1109 FXB, not in this room
 - Covers lectures 1-12
 - You are allowed one 8.5x11" note sheet, double sided
 - Review session: Sunday 10/29 2-4pm in 1690 BBB
- Mid-semester survey due Fri 11/3 at 8pm
 - http://survey2.eecs490.org

Review: Types and Type Judgments

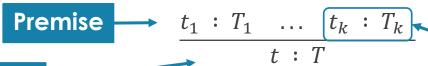
- Our language has two types: Int and Bool
- We determine the type of a **term** in the program based on its syntactic form and the types of its subterms
- A typing relation or type judgment has the form

t:T

and it specifies that term t has type T

Review: Typing Rule

Typing rules have the following familiar form:



Conclusion

Type Judgment

- This is a conditional rule that means:
 - If t_1 has type T_1 , ..., and if t_k has type T_k
 - **Then** t has type T
- This specifies a formula for computing the type of a term in a compiler
 - If the compiler sees a term of the form t, it can compute the type of t by computing the types of $t_1, ..., t_k$ that are in the premises

Review: Subsumption Rule

■ The subsumption rule allows a term to be typed as a supertype of its actual type:

$$\frac{\Gamma \vdash s : S \qquad S <: T}{\Gamma \vdash s : T}$$

The rule encodes a notion of substitutability, allowing a subtype to be used where a supertype is expected:

$$\frac{\Gamma \vdash f : Float \rightarrow Float}{\Gamma \vdash x : Int} \frac{\Gamma \vdash x : Int}{\Gamma \vdash x : Float}$$
$$\Gamma \vdash (f x) : Float$$

Joins

- We need to rewrite the arithmetic rules to work with both Ints and Floats
- The result type should be the least upper bound, or join, of the operand types
 - The join $T = T_1 \sqcup T_2$ is the minimal type T such that $T_1 <: T$ and $T_2 <: T$

$$Int = Int \sqcup Int$$

 $Float = Int \sqcup Float$
 $Float = Float \sqcup Float$

■ Rule for addition:

Require operand type to be a number

$$\frac{\Gamma \vdash t_1 : T_1 \quad \Gamma \vdash t_2 : T_2 \quad T_1 <: Float \quad T_2 <: Float \quad T = T_1 \sqcup T_2}{\Gamma \vdash (t_1 + t_2) : T}$$

The Top Type

Many languages have a Top type (also written as T), that is a supertype of every other type:

- Example: object in Python
- Adding Top to our language ensures that every pair of types has a join¹
- We can then relax the rule for conditionals:

$$\frac{\Gamma \vdash t_1 : Bool}{\Gamma \vdash (\mathbf{if} \ t_1 \ \mathbf{then} \ t_2 \ \mathbf{else} \ t_3) : T} \qquad T = T_2 \sqcup T_3$$

Contravariant Parameters

- A function that takes in a more general parameter type should be substitutable for a function that takes in a more specific parameter type
- For example, the following should be valid:

```
((\mathbf{lambda} \ f : Int \rightarrow Bool \ . \ (f \ 3)) \ (\mathbf{lambda} \ x : Float \ . \ \mathbf{true}))
```

- Thus, if $T_1 <: S_1$, then it should be that $S_1 \to U <: T_1 \to U$
- This permits a contravariant parameter type, since the direction of <: is switched between the parameter and function types</p>

Covariant Return Types

- A function that takes returns a more specific type should be substitutable for a function returns a more general type
- For example, the following should be valid:

```
((\mathbf{lambda} f : Int \rightarrow Float . (f 3)) (\mathbf{lambda} x : Int . x))
```

- Thus, if $S_2 <: T_2$, then it should be that $U \to S_2 <: U \to T_2$
- This permits a covariant return type, since the direction of <: is the same between the return and function types

Subtyping for Functions

In general, a function is substitutable for another if the parameter types are contravariant and the return types are covariant:

$$((\mathbf{lambda} f : Int \rightarrow Float . (f 3)) (\mathbf{lambda} x : Float . 0))$$

Rule for subtyping functions:

$$\frac{T_1 <: S_1 \qquad S_2 <: T_2}{S_1 \to S_2 <: T_1 \to T_2}$$

Data Abstraction

- Abstraction separates what something is from how it works
- Abstract data types (ADTs) separate the interface of a data type from its implementation
- **Encapsulation** is an important, though not universal, property of an ADT, binding the data the ADT represents along with the functions that operate on that data
- The course notes build a hierarchy of ADTs, beginning with immutable pairs all the way up to an abstraction similar to that provided by object-oriented programming
 - Required reading, like the rest of the notes!

Message Passing

- Higher-order functions can provide encapsulation in a functional ADT, allowing us to pass a *message* requesting a particular behavior
- A dispatch function takes the appropriate action given a message

```
>>> p = mutable_pair(3, 4)
>>> p('first')
3
>>> p('second')
4
>>> p('set_first', 5)
>>> p('set_second', 6)
>>> p('first')
5
>>> p('second')
6
```

```
def mutable_pair(x, y):
    def dispatch(message, value=None):
        nonlocal x, y
        if message == 'first':
            return x
        elif message == 'second':
            return y
        elif message == 'set_first':
            x = value
        elif message == 'set_second':
            y = value
        return dispatch
```

Dispatch Dictionaries

- A dispatch dictionary stores a mapping of messages to functions that perform the specified behavior
 - The dispatch function now just looks up a message in the dictionary and returns the corresponding function

```
def account(initial_balance):
    ...
    dispatch = dictionary()
    dispatch('setitem', 'balance', initial_balance)
    dispatch('setitem', 'deposit', deposit)
    dispatch('setitem', 'withdraw', withdraw)
    dispatch('setitem', 'get_balance', get_balance)

def dispatch_message(message):
    return dispatch('getitem', message)

return dispatch_message
```

Object-Oriented Programming

- Object-oriented languages provide a systematic mechanism for defining abstract data types
- Fundamental features:
 - **Encapsulation**: bundling together data of an ADT along with the functions that operate on the data
 - Information hiding: restricting access to the implementation details of an ADT
 - Inheritance: reusing code of an existing ADT when defining a new one
 - Subtype polymorphism: using an instance of a derived ADT where a base ADT is expected
 - Requires some form of dynamic binding, where the derived functionality is used at runtime

The term "encapsulation" is often used to encompass information hiding as well.

Terminology

- A class defines a pattern for the instances of an ADT
 - Specifies the data included and the functions that operate on that data
- An object is an instance of a class
- The individual data items and functions that comprise a class are its members
- Data members are also called fields or attributes
- Member functions are usually called methods

```
int x;
Foo(int x_);
int bar(int y);
Method
};
```

Static Fields

- Each object has its own set of instance fields
- Static fields are associated with a class, and there is only one copy shared by all instances of the class
 - Can generally be accessed directly through class or indirectly through an instance
- Example in Java:

```
class Foo {
   static int bar = 3;
}

class Main {
   public static void main(String[] args) {
      System.out.println(Foo.bar);
      System.out.println(new Foo().bar);
   }
}

Access through instance
```

Static Fields in C++

Example:

```
struct Foo {
   static int bar;
};

int Foo::bar = 3;

int main() {
   cout << Foo::bar << endl;
   cout << Foo().bar << endl;
}</pre>
```

Out-of-line definition required to designate storage

Access through class uses scope-resolution operator

Access through instance uses dot operator

Static Fields in Python

■ In Python, variables defined directly within the class definition are automatically static fields

```
class Foo:
    bar = 3

print(Foo.bar)
print(Foo().bar)
```

Instance fields have to be defined through self

```
class Baz:
    def __init__(self):
        self.bar = 3
```

Access Control

- Information hiding requires ability to restrict access to members of a class
- Access modifiers, in languages that have them, allow the programmer to specify what code has access

	public	private	protected		internal in	
			C++, C#	Java	C#, Java default	Python
Same instance	X	X	X	X	X	Х
Same class	X	X	X	Χ	Χ	Х
Derived classes	Х		X	Х		Х
Code in same package	X			X	Х	Х
Global access	Х					Х

Instance Methods

- Instance methods take in the instance on which to operate as a parameter
 - Often named self or this
 - Usually an implicit parameter
- Example in C++:

```
class Foo {
  int x;
public:
  Foo(int x_) : x(x_) {}
  int get_x() { return this->x; }
};
```

Object that receives method call

```
f.get_x();
Address of object implicitly passed as this
```

this-> can be elided

if x not hidden by

local variable

Methods in Python

- In Python, instance methods must take the instance as an explicit parameter
 - Named self by convention
- Example:

```
class Foo:
    def __init__(self, x):
        self.x = x

    def get_x(self):
        return self.x

f = foo(3)
f.get_x()

Cannot elide
Cannot elide
```

self.

Object that receives method call

We'll start again in five minutes.

Static Methods

- Static methods do not operate on an instance, so they do not have access to instance members
- In many languages, the static keyword denotes a static method
- In Python, the @staticmethod decorator must be used to enable access through both a class and instance

```
class Baz:
    @staticmethod
    def name():
        return 'Baz'

print(Baz.name())
print(Baz().name())
```

Class Methods in Python

Python also allows the definition of class methods, which take in the class as an argument

```
class Baz:
    @classmethod
    def name(cls):
        return cls.__name__

class Fie(Baz):
    pass

print(Baz.name())  # prints Baz
print(Baz().name())  # prints Baz
print(Fie.name())  # prints Fie
print(Fie().name())  # prints Fie
```

Property Methods

- Some languages enable property methods to be defined, which have the syntax of field access but invoke methods
 - Abstract the interface of a field from its implementation
- Example in Python:

Property Setters

- In Python, the @property decorator only specifies a getter property method
- A setter can be defined with @<method>.setter, where <method> is the name of the property method

```
@magnitude.setter
def magnitude(self, mag):
    old_angle = self.angle
    self.real = mag * math.cos(old_angle)
    self.imag = mag * math.sin(old angle)
```

```
>>> c.magnitude = math.sqrt(2)
>>> c.angle = math.pi / 4
>>> c.real
1.0000000000000002
>>> c.imag
1.0
```

Nested and Local Classes

- Many languages allow classes to be defined within another class or within a local scope
- Languages in which classes are first-class entities, such as Python, allow classes to be created dynamically
 - Generally have access to all variables in scope
- In other languages, such as C++, the primary purpose of a nested or local class is to limit the scope in which it may be used
 - In C++, a nested class has access to the private members of its enclosing class, but not vice versa
 - Local classes in C++ do not have access to local variables

Nested Classes in Java

- In Java, local classes have access to effectively final local variables
- Nested and local classes defined at non-static scope are associated with an instance of the enclosing class and have access to its members

```
class Outer {
  private int x;
  Outer(int x_) { x = x_; }
  class Inner {
    private int y;
    Inner(int y_) { y = y_; }
    int get() { return x + y; }
}

Outer out = new Outer(3);
Outer.Inner inn = out.new Inner(4);
System.out.println(inn.get());
```

OOP and Message Passing

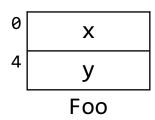
- Conceptually, object-oriented programming consists of passing messages to objects, which then respond to the message
 - Member access on an object can be thought of as sending a message to the object
- Languages differ in:
 - Whether the set of messages an object responds to (i.e. its members) is fixed at compile time
 - Whether the actual message to be sent to an object must be known at compile time

Record¹-Based Implementation

- In languages that prioritize efficiency, the members of an object are known at compile time
- Fields of an object are stored directly within the memory of the object, at offsets that can be computed at compile time
- Field access can be translated by the compiler to an offset into the object

```
class Foo {
public:
    int x, y;
    Foo(int x_, int y_);
};

Foo f(3, 4);
cout << (f.x + f.y);</pre>
```



Dictionary-Based Implementation

- In languages that allow members to be added to an object at runtime, an object's members are usually stored in a dictionary
 - Similar to our message-passing implementation from the notes
- A well-defined lookup process specifies how to lookup a member
 - In Python, check instance dictionary first, then class

```
class Foo:
    y = 2
    def __init__(self, x):
        self.x = x
```

Adds binding to instance dictionary

```
f = Foo(3)
print(f.x, f.y, Foo.y) # prints 3 2 2

→ f.y = 4
print(f.x, f.y, Foo.y) # prints 3 4 2
```

Slots in Python

class Complex(object):

 Python actually takes a hybrid approach, using a dictionary by default but allowing a record-like representation as well

Objects that are dictionary-less lose the ability to add

instance attributes at runtime.

11/2/17

Dynamic Messages

- Dictionary-based languages generally provide a means for constructing and sending a message to an object at runtime
- Example in Python:

```
>>> x = [1, 2, 3]
>>> x.__getattribute__('append')(4)
>>> x
[1, 2, 3, 4]
```

Java Reflection

- In Java, the powerful reflection API allows inspection of classes and objects at runtime
- Reflection can be used to construct and invoke a dynamic message

```
import java.lang.reflect.Method;

class Main {
   public static void main(String[] args)
       throws Exception {
       String s = "Hello World";
       Method m =
            String.class.getMethod("length", null);
       System.out.println(m.invoke(s)); // prints 11
   }
}
```