

# EECS 490 – Lecture 2

Names and Environments

1

10/16/17

# Announcements

- Homework 1 due 9/15 at 8pm
- Entry survey due at 8pm tonight

# Agenda

- Environments and Name Lookup
- Static and Dynamic Scope
- Point of Declaration

# Names

- Fundamental form of abstraction
  - Allow entities of arbitrary complexity to be referenced by a single name
- A name is distinct from the entity it names
  - The same name can refer to different entities in different contexts or at different times
  - An entity may have multiple names that refer to it
- Languages define built-in names and also provide a mechanism for users to define their own names

```
void foo() {  
    int x;  
}
```

```
void bar() {  
    double x;  
}
```

# Scope and Frames

- In order to properly implement abstraction, names in general must have a restricted scope
  - Avoid conflict between internal names defined in different contexts
- The mapping of names to entities is tracked at runtime in individual *frames* or *activation* records for each region of scope
  - A name is *bound* to an entity in a frame or scope

```
void foo() {  
    int x = 3;  
}
```

```
void bar() {  
    double x = 3.1;  
}
```

```
foo  
x: 3
```

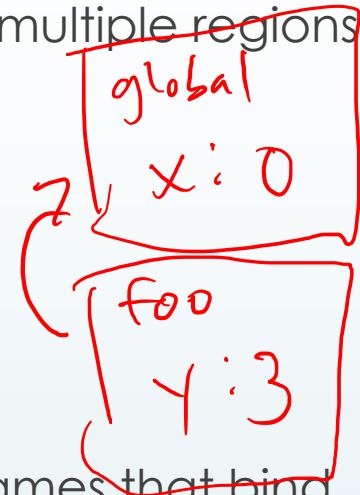
```
bar  
x: 3.1
```

# Frames and Environments

- A piece of code may be located in multiple regions of scope

*foo(3)*

```
int x = 0;
void foo(int y) {
    cout << (x + y) << endl;
}
```



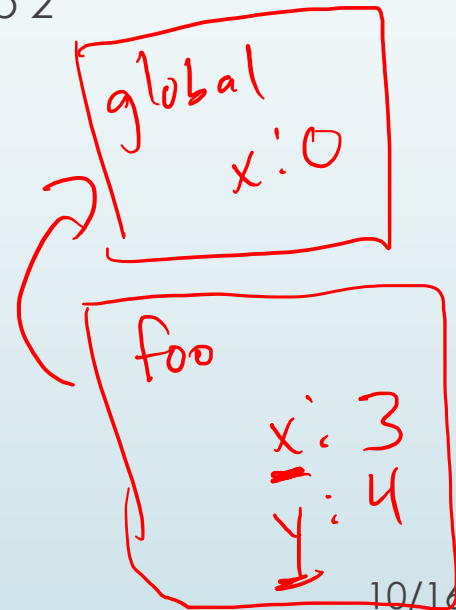
- It therefore has access to multiple frames that bind names to entities
- Frames are generally ordered by how restricted their corresponding scope regions are
- The set of frames available to a piece of code is called its *environment*

# Name Lookup

- Names have a well-defined procedure to look them up in an environment with multiple frames:
  1. Start lookup in the innermost frame
  2. If the name is bound in the current frame, then use that binding
  3. If the name is not bound in the current frame, proceed to the next frame and go to step 2

Example:  
`int x = 0;  
void foo(int x, int y) {  
 cout << (x + y);  
}`

`foo(3, 4)`



# Overloading

- ▶ A name is *overloaded* if it has multiple bindings in the same frame
- ▶ A language that allows overloading must define how overloads are resolved

```
void foo(int x);  
int foo(const string &s);  
foo(3);  
foo("hello");
```

- ▶ Some languages, such as Java, use similar rules to disambiguate names in separate frames

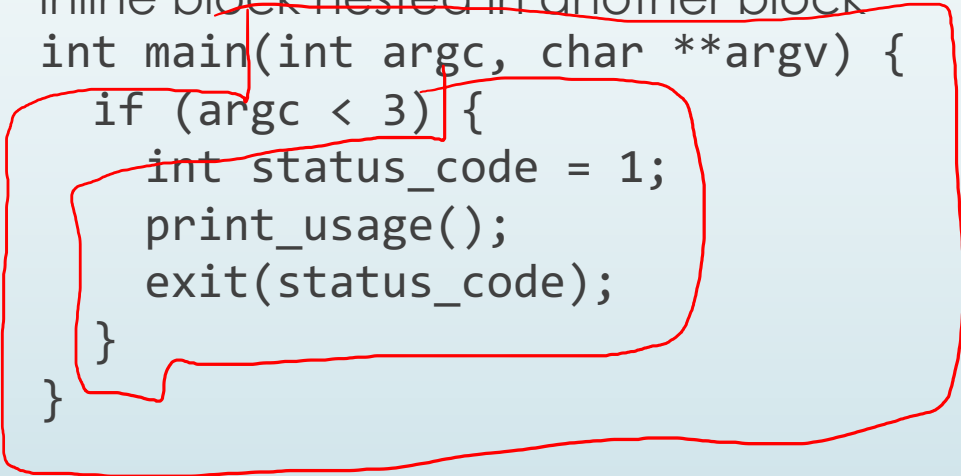
```
public static void main(String[] args) {  
    int main = 3;  
    main(null); // recursive call  
}
```



# Blocks

- A *block* is a compound statement that groups together other statements  
`{ statement1; statement2; ...; statmentN; }`
- A block usually defines a region of scope and therefore has its own frame
- Blocks can be associated with a function or be an inline block nested in another block

```
int main(int argc, char **argv) {  
    if (argc < 3) {  
        int status_code = 1;  
        print_usage();  
        exit(status_code);  
    }  
}
```



# Suites in Python

- ▶ Python does not have inline blocks
- ▶ Compound statements can be composed of a *header* followed by a *suite* of statements
- ▶ In general, a suite does not have its own frame

```
def foo(x):  
    if x < 0:  
        negative = True  
    else:  
        negative = False  
    print(negative)
```

# Blocks in Scheme

- The *let* forms in Scheme introduce a new frame

```
(let ((x 3) (y 4))  
  (display (+ x y))  
  (display (- x y)))
```

- This is commonly implemented by translating into a function definition and call:

```
((lambda (x y)  
  (display (+ x y))  
  (display (- x y)))  
 3 4)
```

Anonymous  
function; more on  
this in a few weeks

# Environments and Nested Blocks

- Nested blocks result in nested frames in the environment
- *Visibility rules* correspond to the lookup procedure

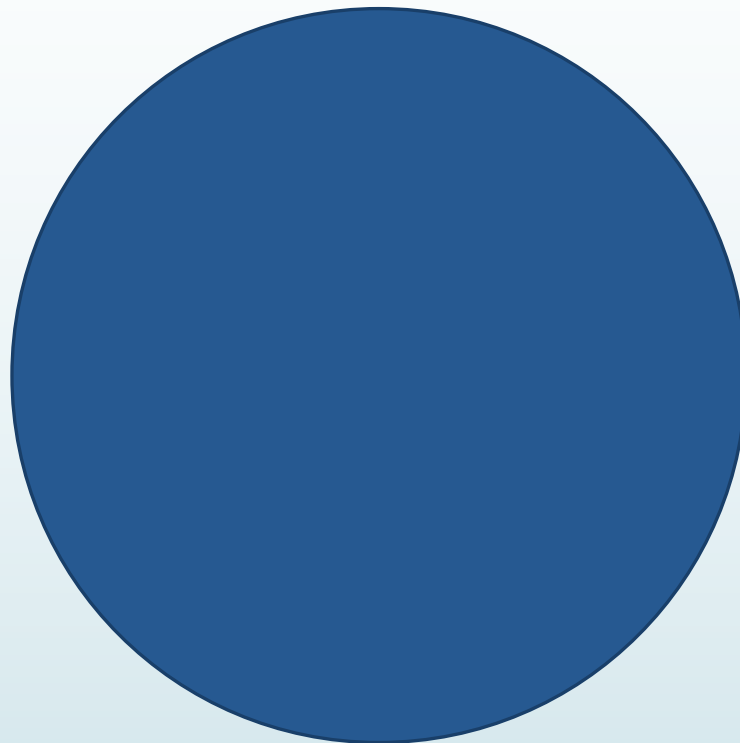
```
{  
  int x = 0;  
  int y = 1;  
  {  
    int x = 2;  
    int z = 3;  
    cout << (x + y + z);  
  }  
}
```

Inner x  
shadows  
outer x

Binding *hidden* by  
declaration of x in  
inner block

Binding *visible* in  
inner block

- ▶ We'll start again in five minutes.



# Functions and Environments

- Functions differ from inline blocks in that the context in which they are defined differs from the context in which they execute

```
int x = 0;
```

```
void foo() {  
    print(x);  
}
```

```
void bar() {  
    int x = 1;  
    foo();  
}
```

bar();

Which x is printed?  
Either is a valid choice

global

x : 0

bar

x : 1

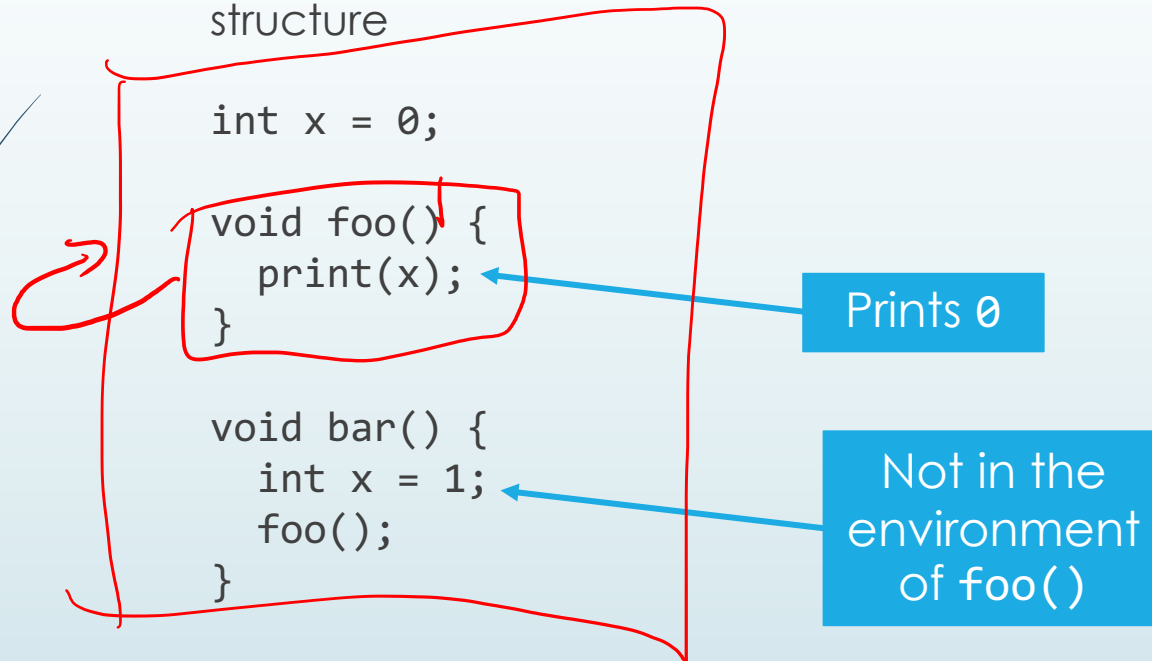
foo

# Kinds of Environments

- The environment in which a function executes is often divided into three components
  - The *local* environment is the part that is internal to the function
  - The *global* environment is the part defined at the top-level of a program, at global or module scope
  - The *non-local* environment consists of the bindings that are visible to a function but not part of the local or global environment
- The two possibilities for which *x* is printed correspond to different choices about what constitutes the non-local environment

# Static Scope

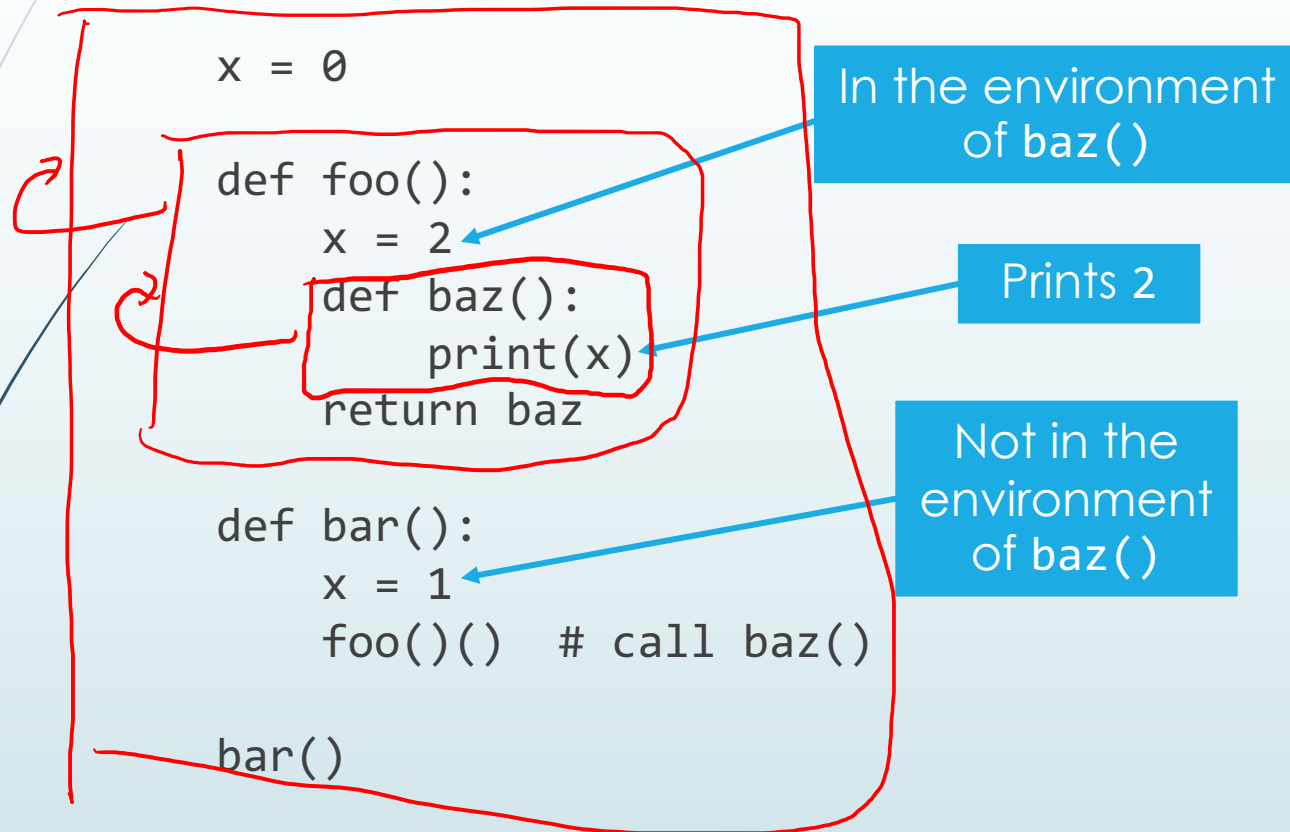
- In *static* or *lexical* scope, the non-local environment of a function is the environment in which the function is defined
  - Can be determined directly from the program's syntactic structure





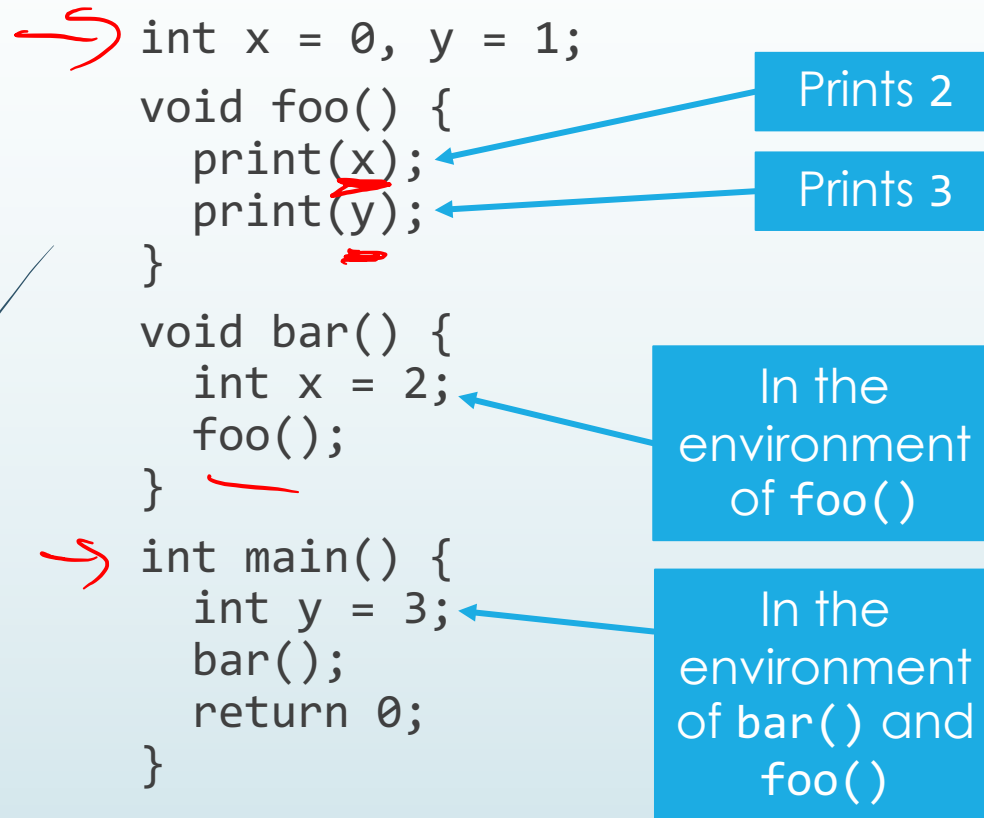
# Nested Function Definitions

- Nested function definitions result in more complex environments in static scope



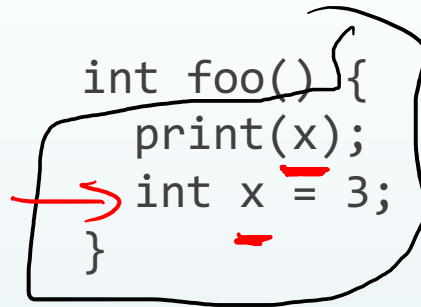
# Dynamic Scope

- In *dynamic* scope, the non-local environment of a function is the environment in which it is called



# Use Before Declaration

- An exact correspondence between blocks, frames, and scope allows code such as the following:



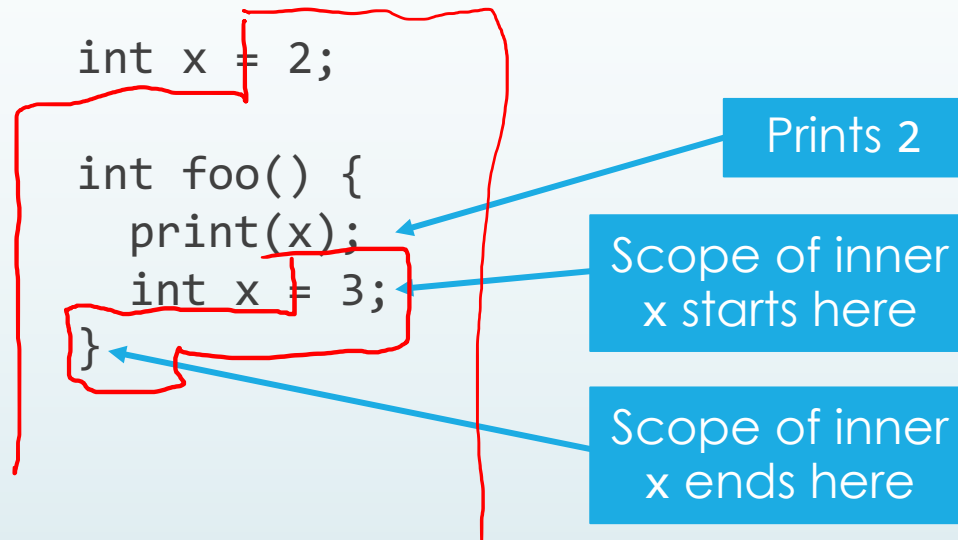
```
int foo() {  
    print(x);  
    int x = 3;  
}
```

A hand-drawn diagram of a function scope. A black bracket on the right side of the function definition groups the lines. A red arrow points from the left to the variable `x` in the `print(x);` statement, indicating its use before it has been declared. The variable `x` in the initialization `int x = 3;` is underlined in red.

- This should be invalid, since `x` is used before it is initialized

# Point of Declaration

- In some languages, including the C family, the scope of a name extends from its *point of declaration* to the end of the enclosing block



# Assignments in Python

- ▶ Python assumes that an assignment to a variable is intended to target a local variable
- ▶ Furthermore, the scope of a local variable starts at the beginning of a function
- ▶ Using a variable before it is initialized is an error

```
x = 2
```

```
def foo():  
    print(x)  
    x = 3
```

Scope of local  
x starts here

Error: use before  
initialization

Defines local  
variable x

# global and nonlocal in Python

- A programmer can specify that a name is meant to refer to a global or non-local variable using the `global` and `nonlocal` statements

```
x = 2
```

```
def foo():  
    ↪ global x  
    print(x)  
    x = 3
```

Specifies that x  
refers to the  
global variable

Prints 2

Assigns 3 to the  
global variable x

```
foo()  
print(x)
```

Prints 3

# Mutually Recursive Entities

- The C-style point of declaration rules are insufficient for defining mutually recursive entities

```
int foo(int x) {  
    return bar(x + 1);  
}
```

Error: use before  
declaration

```
int bar(int x) {  
    return foo(x - 1);  
}
```

Scope of bar  
starts here

# Incomplete Declarations

- C and C++ allow *incomplete declarations* that allow an entity to be declared without being defined

```
int foo(int x) {  
    int bar(int);  
    return bar(x + 1);  
}
```

Incomplete  
declaration

Scope of incomplete  
declaration ends here

```
int bar(int x) {  
    return foo(x - 1);  
}
```