EECS 490 – Lecture 10

Continuations

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Announcements

Project 2 due Fri 10/6 at 8pm

Agenda

■ Restricted Continuations

■ First-Class Continuations

Review: First-Class Entities

- We use entity to denote something that can be named in a program
 - Other terms also used: citizen, object
 - Examples: types, functions, data objects, values
- A first-class entity is an entity that supports all operations generally available to other entities
 - e.g. can be assigned to a variable, passed to or returned from a function

	C++	Java	Python	Scheme
Functions	sort of	no	yes	yes
Types	no	no	yes	no
Control	no	no	no	yes

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Continuations



- A continuation represents the control state of a program
 - The sequence of active functions
 - Code location within each active function
 - Intermediate results
- A continuation can be invoked to return control to a previous state
- Only control state is restored, not state of data

Continuation Analogy

Say you're in the kitchen in front of the refrigerator, thinking about a sandwitch [sic]. You take a continuation right there and stick it in your pocket. Then you get some turkey and bread out of the refrigerator and make yourself a sandwitch, which is now sitting on the counter. You invoke the continuation in your pocket, and you find yourself standing in front of the refrigerator again, thinking about a sandwitch. But fortunately, there's a sandwitch on the counter, and all the materials used to make it are gone. So you eat it.:-)

Luke Palmer

Types of Continuations

- A language may provide restricted forms of continuations that can only be invoked at specific times
 - Subroutines (i.e. functions)
 - Coroutines
 - Exceptions
 - Generators
- Some languages have first-class continuations that can be stored in a variable and invoked at arbitrary times

Subroutines

- A subroutine involves transfer of control between a caller and a callee
- Before control is transferred to the callee, the state of the caller, i.e. its continuation, must be saved
 - Intermediate results stored in caller's activation record
 - Information about how to return control to caller stored in callee's activation record
- Upon completion of call, caller's continuation invoked

Continuation of foo()

must be saved
before call to bar()

Abrupt Termination

- In some languages the caller's continuation is only invoked when the callee completes normally
- Other languages allow early termination of a call, also called abrupt termination, with a return statement

```
def foo(x):
    return x
    # dead code
    if x < 0:
        bar(x)
    baz(x)</pre>
Invoke caller's
    continuation

Code never
    reached
```

Control vs. Data State

 A continuation only represents control state, so invoking it does not restore the state of data

```
def outer():
\rightarrow X = 0
    def inner():
         nonlocal x
         x += 1
                        Invoke continuation
                             of outer()
   inner()
print(x) # prints 1
                        Value of x is
                        not restored
```

Coroutines

 Generalize subroutines to allow multiple routines to invoke each other's continuations

```
var q := new queue
coroutine produce
    loop
        while q is not full
            create some new items
          add the items to q ield to consume
coroutine consume
    loop
        while q is not empty
            remove some items from q
            use the items
          ield to produce 🥧
```

This is pseudocode.

Exceptions

 Allow control to be passed to a function further up in the call chain, rather than just the direct caller

```
def foo(x):
                      Save continuation of
    try: <
                     foo(), add exception
        bar(x)
                    handler to handler stack
    except:
        print('Exception')
def bar(x):
    baz(x)
def baz(x):
                                Invoke continuation
    raise Exception -
                                   of foo(), run
                                exception handler
```

Generators



- Like a subroutine, but allow execution to be paused and resumed
- Also called semicoroutine
 - Generator can be resumed by any caller
 - However, generator can only yield execution to caller that invoked it

```
def naturals():
    num = 0
    while True:
    yield num *
    num += 1
```

Pause execution and yield an item to caller

Generators and Iterators

- In Python, generators implement the same interface as an iterator
- Often simpler to write generator than a class that implements the iterator interface

```
def naturals():
    num = 0
    while True:
        yield num
    num += 1
```

```
>>> numbers = naturals()
>>> next(numbers)
0
>>> next(numbers)
1
>>> next(numbers)
2
```

Finite Generators

- A finite generator automatically raises a
 StopIteration exception when it completes
 - Used by a for loop to determine the end of an iterator

```
def range2(start, stop, step = 1):
    while start < stop:
        yield start
        start += step</pre>
```

```
>>> values = range2(0, 5, 3)
>>> next(values)
0
>>> next(values)
3
>>> next(values)
Traceback (most recent call last):
File "<stdin>", line 1, in <module>
StopIteration
```

```
>>> for i in range2(0, 4):
... print(i)
...
0
1
2
3
```

Generator Expressions

 Similar to list comprehensions, but produce a generator instead

```
def naturals():
    num = 0
    while True:
        yield num
    num += 1
```

```
>>> negatives = (-i for i in naturals() if i != 0)
>>> next(negatives)
-1
>>> next(negatives)
-2
>>> next(negatives)
-3
Generator
expression
```

■ We'll start again in five minutes.

First-Class Continuations

- Many functional languages allow the current continuation to be captured in an explicit data structure
- Continuation can be passed as a parameter, returned, saved as a variable, etc.
- Depending on the language, the continuation may be invoked only once or an arbitrary number of times

Call/cc call/cc call-with-current-continuation

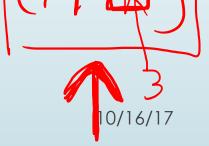
- In Scheme, the call-with-current-continuation procedure, often abbreviated call/cc, creates an object representing the current continuation
- It then calls another procedure with the continuation as the argument

(call-with-current-continuation continuation continuation

- The called procedure can invoke the continuation, return it, discard it,etc.
 - If the procedure returns normally, the call/cc call evaluates to its result

> (+ 1 (call/cc (lambda (cc) 3)))
4

Must be a one-argument procedure





Invoking a Continuation



 A continuation is invoked with a value, which en becomes the "return value" of the call/cc eall

```
> (+ 1 (call/cc (lambda (cc) (cc 5) 3)))
6
```

Call to call/cc replaced with value 5

Continuation invoked with value 5

Not evaluated



Allows a continuation to be invoked multiple times

```
> (define var (call/cc (lambda (cc) cc)))
var
> (define cont var)
cont
> (cont 3) •
                         Becomes
var
                      (define var 3)
> var
  (cont 4)
                         Becomes
var
                      (define var 4)
> var
```

Example: Factorial

Continuation that multiplies a number by the factorial of another number:
 (define cont '())

```
> (factorial 3)
6
> (cont 1)
6
> (cont 3)
18
> (factorial 5)
120
> (cont 4)
480
```

Emulating Call and Return

- Scheme does not provide abrupt termination, but we can emulate it with continuations
- We need:
 - A data structure to explicitly represent the call stack
 - A mechanism for calling a procedure while saving the caller's continuation
 - A mechanism for returning from a procedure by invoking the continuation of the caller
- For simplicity, we will only implement this for oneargument procedures

Call Stack

- A standard stack data structure using a list
- We need set! to modify the structure

```
(define call-stack '())

(define (push-call call)
   (set! call-stack (cons call call-stack)))

(define (pop-call)
   (let ((caller (car call-stack)))
      (set! call-stack (cdr call-stack))
      caller))
```

Call and Return

The return procedure just pops off the caller's continuation from the stack and invokes it

```
(define (return value)
  ((pop-call) value))
```

■ The call procedure must push the current continuation on the stack and then call the target procedure

Using Call and Return

We can now use call and return:

```
> (+ 1 (call foo 3))
4
> (+ 1 (call foo 20))
23
> (+ 2 (call bar 3))
-1
> (+ 2 (call bar 20))
-20
```

Yin-Yang Puzzle

Prints out unary representations of the natural numbers

Continuations and Goto

- First-class continuations are often criticized for the same reasons as goto, since they allow unstructured transfer of control
- As with goto, continuations should be used judiciously
 - Implementing more restricted forms of control transfer such as exceptions
 - Adhering to conventions as in continuation-passing style