

Grammars, Functional Programming, and Scheme

Context-Free Grammar: A context-free grammar is a quadruple (NT, T, R, S) where:

- NT is a finite set of symbols (non-terminal symbols, variables, or syntactic categories)
- T is a finite set of terminal symbols
- R is a finite set of productions (or rules), each of which is composed of an expression of the form

$$V \rightarrow w \quad (\Sigma = T \cup NT, w \in \Sigma^*)$$

where V (the head of the production) is a single non-terminal symbol and w (the body) is a string composed of zero or more terminal or non-terminal symbols:

- S is the initial symbol and is an element of NT
1. Give an ambiguous CFG for **if-then-else**. Draw a derivation tree to prove that the grammar is ambiguous. Use the non-terminal S as your initial symbol and C for the if condition. You do not need to provide the production rules for C .

2. Backus-Naur form (BNF) is another notation used to describe a grammar where:
 - Non-terminal symbols are written between angle brackets
 - Productions with the same head are grouped into a single block using “|” to separate productions

Excerpt from the Scheme lexical structure in BNF:

$\langle \text{expression} \rangle \rightarrow \langle \text{variable} \rangle$	$\langle \text{number} \rangle \rightarrow \langle \text{integer} \rangle \mid \langle \text{decimal} \rangle$
$\mid \langle \text{literal} \rangle$	$\langle \text{integer} \rangle \rightarrow \langle \text{sign} \rangle \langle \text{digit} \rangle^+$
$\mid \langle \text{procedure call} \rangle$	$\langle \text{sign} \rangle \rightarrow \langle \text{empty} \rangle \mid + \mid -$
$\mid \langle \text{lambda expression} \rangle$	$\langle \text{digit} \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$
$\mid \langle \text{conditional} \rangle$	$\langle \text{conditional} \rangle \rightarrow (\text{if } \langle \text{test} \rangle \langle \text{consequent} \rangle$
$\langle \text{literal} \rangle \rightarrow \langle \text{quotation} \rangle \mid \langle \text{self-evaluating} \rangle$	$\langle \text{alternate} \rangle)$
$\langle \text{self-evaluating} \rangle \rightarrow \langle \text{boolean} \rangle \mid \langle \text{number} \rangle$	$\langle \text{test} \rangle \rightarrow \langle \text{expression} \rangle$
$\mid \langle \text{character} \rangle \mid \langle \text{string} \rangle$	$\langle \text{consequent} \rangle \rightarrow \langle \text{expression} \rangle$
$\langle \text{boolean} \rangle \rightarrow \#t \mid \#f$	$\langle \text{alternate} \rangle \rightarrow \langle \text{expression} \rangle \mid \langle \text{empty} \rangle$

Write out the derivation tree for (if #t 3)

3. Write a scheme procedure `reverse-list` that takes in a list and outputs the reversed list.
Hint: use the builtin procedure `append`; when given all list arguments, the result is a list that contains all of the elements of the given lists in order

4. Define a **tail-recursive** fibonacci function in Scheme. *Hint:* use a helper function