# **Moduli Spaces**

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1	Thursday January 9th
So	ome references:
	• Course Notes
	• Hilbert schemes/functors of points: Notes by Stromme

 $-\,$  Slightly more detailed: Nitsure,  $\dots$  Hilbert schemes, Fundamentals of Algebraic Geometry

- Mumford, Curves on Surfaces
- Harris-Harrison, Moduli of Curves (chatty and less rigorous)

### 1.1 Representability

Last time: Fix an S-scheme, i.e. a scheme over S.

Then there is a map

$$\operatorname{Sch}/S \longrightarrow \operatorname{Fun}(\operatorname{Sch}/S^{\operatorname{op}}, \operatorname{Set})$$
  
 $x \mapsto h_x(T) = \operatorname{hom}_{\operatorname{Sch}/S}(T, x).$ 

where  $T' \xrightarrow{f} T$  is given by

$$h_x(f): h_x(T) \longrightarrow h_x(T')$$
  
 $(T \mapsto x) \mapsto \text{triangles of the form}$ 



Theorem 1.1(Yoneda).

$$hom_{Fun}(h_x, F) = F(x).$$

Corollary 1.2.

$$hom_{Sch/S}(x,y) \cong hom_{Fun}(h_x, h_y).$$

**Definition 1.2.1** (Moduli Functor).

A moduli functor is a map

$$F: (\mathrm{Sch}/S)^{\mathrm{op}} \longrightarrow \mathrm{Set}$$
 
$$F(x) = \text{"Families of something over } x\text{"}$$
 
$$F(f) = \text{"Pullback"}.$$

#### **Definition 1.2.2** (Moduli Space).

A **moduli space** for that "something" appearing above is an  $M \in \text{Obj}(\text{Sch}/S)$  such that  $F \cong h_M$ .

Now fix S = Spec (k).

 $h_m$  is the functor of points over M.

**Remark (1)**  $h_m(\operatorname{Spec}(k)) = M(\operatorname{Spec}(k)) \cong \text{"families over Spec } k" = F(\operatorname{Spec}(k)).$ 

**Remark (2)**  $h_M(M) \cong F(M)$  are families over M, and  $\mathrm{id}_M \in \mathrm{Mor}_{\mathrm{Sch}/S}(M,M) = \xi_{Univ}$  is the universal family.

Every family is uniquely the pullback of  $\xi_{\text{Univ}}$ . This makes it much like a classifying space. For  $T \in \text{Sch}/S$ ,

$$h_M \xrightarrow{\cong} F$$

$$f \in h_M(T) \xrightarrow{\cong} F(T) \ni \xi = F(f)(\xi_{\text{Univ}}).$$

where  $T \xrightarrow{f} M$  and  $f = h_M(f)(\mathrm{id}_M)$ .

**Remark (3)** If M and M' both represent F then  $M \cong M'$  up to unique isomorphism.

$$\xi_M \qquad \qquad \xi_{M'}$$

$$M \longrightarrow f \qquad M'$$

$$M' \xrightarrow{g} M$$

$$\xi_{M'}$$
  $\xi_{M}$ 

which shows that f, g must be mutually inverse by using universal properties.

### Example 1.1.

A length 2 subscheme of  $\mathbb{A}^1_k$  (??) then

$$F(S) = \left\{ V(x^2 + bx + c) \right\} \subset \mathbb{A}_5'$$

where  $b, c \in \mathcal{O}_s(s)$ , which is functorially bijective with  $\{b, c \in \mathcal{O}_s(s)\}$  and F(f) is pullback.

Then F is representable by  $\mathbb{A}_k^2(b,c)$  and the universal object is given by

$$V(x^2 + bx + c) \subset \mathbb{A}^1(?) \times \mathbb{A}^2(b, c)$$

where  $b, c \in k[b, c]$ .

Moreover, F'(S) is the set of effective Cartier divisors in  $\mathbb{A}_5'$  which are length 2 for every geometric fiber. F''(S) is the set of subschemes of  $\mathbb{A}_5'$  which are length 2 on all geometric fibers. In both cases, F(f) is always given by pullback.

Problem: F'' is not a good moduli functor, as it is not representable. Consider Spec  $k[\varepsilon]$ .





We think of  $T_p F^{',"}$  as the tangent space at p.

If F is representable, then it is actually the Zariski tangent space.





Moreover,  $T_pM = (\mathfrak{m}_p/\mathfrak{m}_p^2)^{\vee}$ , and in particular this is a k-vector space. To see the scaling structure, take  $\lambda \in k$ .

$$\lambda: k[\varepsilon] \longrightarrow k[\varepsilon]$$

$$\varepsilon \mapsto \lambda \varepsilon$$

$$\lambda^*: \operatorname{Spec} (k[\varepsilon]) \longrightarrow \operatorname{Spec} (k[\varepsilon])$$

$$\lambda: M(\operatorname{Spec} (k[\varepsilon])) \longrightarrow M(\operatorname{Spec} (k[\varepsilon]))$$

$$\cup \qquad \cup$$

$$T_pM \longrightarrow T_pM.$$

**Conclusion**: If F is representable, for each  $p \in F(\text{Spec } k)$  there exists a unique point of  $T_pF$  that are invariant under scaling.

1. If  $F, F', G \in \text{Fun}((\text{Sch}/S)^{\text{op}}, \text{Set})$ , there exists a fiber product



where

$$(F \times_G F')(T) = F(T) \times_{G(T)} F'(T).$$

2. This works with the functor of points over a fiber product of schemes  $X \times_T Y$  for  $X, Y \longrightarrow T$ ,

where

$$h_{X \times_T Y} = h_X \times_{h_t} h_Y.$$

- 3. If F, F', G are representable, then so is the fiber product  $F \times_G F'$ .
- 4. For any functor

$$F: (\operatorname{Sch}/S)^{\operatorname{op}} \longrightarrow \operatorname{Set},$$

for any  $T \xrightarrow{f} S$  there is an induced functor

$$F_T: (\operatorname{Sch}/T) \longrightarrow \operatorname{Set}$$
  
 $x \mapsto F(x).$ 

5. F is representable by M/S implies that  $F_T$  is representable by  $M_T = M \times_S T/T$ .

#### 1.2 Projective Space

Consider  $\mathbb{P}^n_{\mathbb{Z}}$ , i.e. "rank 1 quotient of an n+1 dimensional free module".

#### Proposition 1.3.

 $\mathbb{P}^n_{\mathbb{Z}}$  represents the following functor

$$\begin{split} F: \operatorname{Sch}^{\operatorname{op}} &\longrightarrow \operatorname{Set} \\ F(S) &= \mathcal{O}_s^{n+1} &\longrightarrow L \longrightarrow 0/\sim. \end{split}$$

where  $\sim$  identifies diagrams of the following form:

and F(f) is given by pullbacks.

**Remark**  $\mathbb{P}^n_S$  represents the following functor:

$$F_S: (\mathrm{Sch}/S)^{\mathrm{op}} \longrightarrow \mathrm{Set}$$

$$T \mapsto F_S(T) = \left\{ \mathcal{O}_T^{n+1} \longrightarrow L \longrightarrow 0 \right\} / \sim .$$

This gives us a cleaner way of gluing affine data into a scheme.

Proof (of Proposition).

Note:  $\mathcal{O}^{n+1} \longrightarrow L \longrightarrow 0$  is the same as giving n+1 sections  $s_1, \dots s_n$  of L, where surjectivity ensures that they are not the zero section.

So

$$F_i(S) = \left\{ \mathcal{O}_s^{n+1} \longrightarrow L \longrightarrow 0 \right\} / \sim,$$

with the additional condition that  $s_i \neq 0$  at any point.

There is a natural transformation  $F_i \longrightarrow F$  by forgetting the latter condition, and is in fact a subfunctor.

$$F \leq G$$
 is a subfunctor iff  $F(s) \hookrightarrow G(s)$ .

**Claim:** It is enough to show that each  $F_i$  and each  $F_{ij}$  are representable, since we have natural transformations:

$$F_{i} \longrightarrow F$$

$$\uparrow \qquad \qquad \uparrow$$

$$F_{ij} \longrightarrow F_{i}$$

and each  $F_{ij} \longrightarrow F_i$  is an open embedding (on the level of their representing schemes).

#### Example.

For n = 1, we can glue along open subschemes



For n=2, we get overlaps of the following form:



This claim implies that we can glue together  $F_i$  to get a scheme M. We want to show that M represents F. F(s) (LHS) is equivalent to an open cover  $U_i$  of S and sections of  $F_i(U_i)$  satisfying the gluing (RHS).

Going from LHS to RHS isn't difficult, since for  $\mathcal{O}_s^{n+1} \longrightarrow L \longrightarrow 0$ ,  $U_i$  is the locus where  $s_i \neq 0$  and by surjectivity, this gives a cover of S.

RHS to LHS comes from gluing.

Proof (of Claim).

$$F_i(S) = \left\{ \mathcal{O}_S^{n+1} \longrightarrow L \cong \mathcal{O}_s \longrightarrow 0, s_i \neq 0 \right\},$$

but there are no conditions on the sections other than  $s_i$ .

So specifying  $F_i(S)$  is equivalent to specifying n-1 functions  $f_1 \cdots \widehat{f_i} \cdots f_n \in \mathcal{O}_S(s)$  with  $f_k \neq 0$ . We know this is representable by  $\mathbb{A}^n$ .

We also know  $F_{ij}$  is obviously the same set of sequences, where now  $s_j \neq 0$  as well, so we need to specify  $f_0 \cdots \widehat{f_i} \cdots f_j \cdots f_n$  with  $f_j \neq 0$ . This is representable by  $\mathbb{A}^{n-1} \times \mathbb{G}_m$ , i.e. Spec  $k[x_1, \cdots, \widehat{x_i}, \cdots, x_n, x_j^{-1}]$ . Moreover,  $F_{ij} \hookrightarrow F_i$  is open.

What is the compatibility we are using to glue? For any subset  $I \subset \{0, \dots, n\}$ , we can define

$$F_I = \left\{ \mathcal{O}_s^{n+1} \longrightarrow L \longrightarrow 0, s_i \neq 0 \text{ for } i \in I \right\} = \underset{i \in I}{\times} F_i,$$

and  $F_I \longrightarrow F_J$  when  $I \supset J$ .

## 2 Tuesday January 14th

Last time: Representability of functors, and specifically projective space  $\mathbb{P}^n_{\mathbb{Z}}$  constructed via a functor of points, i.e.

$$h_{\mathbb{P}^n_{\mathbb{Z}}} : \mathbb{P}^n_{\mathbb{Z}} \operatorname{Sch}^{\operatorname{op}} \longrightarrow \operatorname{Set}$$

$$s \mapsto \mathbb{P}^n_{\mathbb{Z}}(s) = \left\{ \mathcal{O}^{n+1}_s \longrightarrow L \longrightarrow 0 \right\}.$$

for L a line bundle, up to isomorphisms of diagrams:

That is, line bundles with n+1 sections that globally generate it, up to isomorphism.

The point was that for  $F_i \subset \mathbb{P}^n_{\mathbb{Z}}$  where

$$F_i(s) = \left\{ \mathcal{O}_s^{n+1} \longrightarrow L \longrightarrow 0 \mid s_i \text{ is invertible} \right\}$$

are representable and can be glued together, and projective space represents this functor.

**Remark** Because projective space represents this functor, there is a universal object:



and other functors are pullbacks of the universal one. (Moduli Space)

**Exercise** Show that  $\mathbb{P}^n_{\mathbb{Z}}$  is proper over Spec  $\mathbb{Z}$ . Use the evaluative criterion, i.e. there is a unique lift



#### **Definition 2.0.1** (Equalizer).

For a category C, we say a diagram  $X \longrightarrow Y \rightrightarrows Z$  is an equalizer iff it is universal with respect to the property:



Note that X is the universal object here.

#### Example 2.1.

For sets, 
$$X = \{y \mid f(y) = g(y)\}$$
 for  $Y \xrightarrow{f,g} Z$ .

#### Definition 2.0.2 (Coequalizer).

A coequalizer is the dual notion,



#### Example 2.2.

Take  $C = \operatorname{Sch}/S$ , X/S a scheme, and  $X_{\alpha} \subset X$  an open cover. We can take two fiber products,  $X_{\alpha\beta}, X_{\beta,\alpha}$ :





These are canonically isomorphic.

In Sch/S, we have

$$\coprod_{\alpha\beta} X_{\alpha\beta} \xrightarrow{f_{\alpha\beta}} \coprod_{\alpha} X_{\alpha} \longrightarrow X$$

where

$$f_{\alpha\beta}: X_{\alpha\beta} \longrightarrow X_{\alpha}$$
  
 $g_{\alpha\beta}: X_{\alpha\beta} \longrightarrow X_{\beta};$ 

this is a coequalizer.

Conversely, we can glue schemes. Given  $X_{\alpha} \longrightarrow X_{\alpha\beta}$  (schemes over open subschemes), we need to check triple intersections:



Then  $\varphi_{\alpha\beta}: X_{\alpha\beta} \xrightarrow{\cong} X_{\beta\alpha}$  must satisfy the **cocycle condition**:

1.

$$\varphi_{\alpha\beta}^{-1}(X_{\beta\alpha} \cap X_{\beta\gamma}) = X_{\alpha\beta} \cap X_{\alpha\gamma},$$

noting that the intersection is exactly the fiber product  $X_{\beta\alpha} \times_{X_{\beta}} X_{\beta\gamma}$ .

2. The following diagram commutes:



Then there exists a scheme X/S such that  $\coprod_{\alpha\beta} X_{\alpha\beta} \rightrightarrows \coprod_{\alpha\beta} X_{\alpha} \longrightarrow X$  is a coequalizer; this is the gluing.

Subfunctors satisfy a patching property because morphisms to schemes are locally determined. Thus representable functors (e.g. functors of points) have to be (Zariski) sheaves.

#### Definition 2.0.3 (Zariski Sheaf).

A functor  $F: (\mathrm{Sch}/S)^{\mathrm{op}} \longrightarrow \mathrm{Set}$  is a Zariski sheaf iff for any scheme T/S and any open cover  $T_{\alpha}$ , the following is an equalizer:

$$F(T) \longrightarrow \prod F(T_{\alpha}) \Longrightarrow \prod_{\alpha\beta} F(T_{\alpha\beta})$$

where the maps are given by restrictions.

#### Example 2.3.

Any representable functor is a Zariski sheaf precisely because the gluing is a coequalizer. Thus if you take the cover

$$\coprod_{\alpha\beta} T_{\alpha\beta} \longrightarrow \coprod_{\alpha} T_{\alpha} \longrightarrow T,$$

since giving a local map to X that agrees on intersections if enough to specify a map from  $T \longrightarrow X$ .

Thus any functor represented by a scheme automatically satisfies the sheaf axioms.

### **Definition 2.0.4** (Subfunctors, Open/Closed Functors).

Suppose we have a morphism  $F' \longrightarrow F$  in the category Fun(Sch/S, Set).

- This is a **subfunctor** if  $\iota(T)$  is injective for all T/S.
- $\iota$  is **open/closed/locally closed** iff for any scheme T/S and any section  $\xi \in F(T)$  over T, then there is an open/closed/locally closed set  $U \subset T$  such that for all maps of schemes  $T' \xrightarrow{f} T$ , we can take the pullback  $f^*\xi$  and  $f^*\xi \in F'(T')$  iff f factors through U.

I.e. we can test if pullbacks are contained in a subfunctors by checking factorization.

**Note** This is the same as asking if the subfunctor F', which maps to F (noting a section is the same as a map to the functor of points), and since  $T \longrightarrow F$  and  $F' \longrightarrow F$ , we can form the fiber product  $F' \times_F T$ :



and  $F' \times_F T \cong U$ .

Note: this is almost tautological!

Thus  $F' \longrightarrow F$  is open/closed/locally closed iff  $F' \times_F T$  is representable and g is open/closed/locally closed.

I.e. base change is representable, and (?).

#### **Exercise (Tautologous)**

- 1. If  $F' \longrightarrow F$  is open/closed/locally closed and F is representable, then F' is representable as an open/closed/locally closed subscheme
- 2. If F is representable, then open/etc subschemes yield open/etc subfunctors

Mantra: Treat functors as spaces. We have a definition of open, so now we'll define coverings.

#### **Definition 2.0.5** (Open Covers).

A collection of open subfunctors  $F_{\alpha} \subset F$  is an **open cover** iff for any T/S and any section  $\xi \in F(T)$ , i.e.  $\xi : T \longrightarrow F$ , the  $T_{\alpha}$  in the following diagram are an open cover of T:



#### Example 2.4.

Given

$$F(s) = \left\{ \mathcal{O}_s^{n+1} \longrightarrow L \longrightarrow 0 \right\}$$

and  $F_i(s)$  given by those where  $s_i \neq 0$  everywhere, the  $F_i \longrightarrow F$  are an open cover. Because the sections generate everything, taking the  $T_i$  yields an open cover.

### Proposition 2.1.

A Zariski sheaf  $F: (Sch/S)^{op} \longrightarrow Set$  with a representable open cover is representable.

#### Proof

Let  $F_{\alpha} \subset F$  be an open cover, say each  $F_{\alpha}$  is representable by  $x_{\alpha}$ . Form the fiber product  $F_{\alpha\beta} = F_{\alpha} \times_F F_{\beta}$ . Then  $x_{\beta}$  yields a section (plus some openness condition?), so  $F_{\alpha\beta} = x_{\alpha\beta}$ 

representable. Because  $F_{\alpha} \subset F$ , the  $F_{\alpha\beta} \longrightarrow F_{\alpha}$  have the correct gluing maps.

This follows from Yoneda (schemes embed into functors), and we get maps  $x_{\alpha\beta} \longrightarrow x_{\alpha}$  satisfying the gluing conditions. Call the gluing scheme x; we'll show that x represents F. First produce a map  $x \longrightarrow F$  from the sheaf axioms. We have a map  $\xi \in \prod_{\alpha} F(x_{\alpha})$ , and because we can pullback, we get a unique element  $\xi \in F(X)$  coming from the diagram

$$F(x) \longrightarrow \prod F(x_{\alpha}) \rightrightarrows \prod_{\alpha\beta} F(x_{\alpha\beta}).$$

#### Lemma 2.2.

If  $E \longrightarrow F$  is a map of functors and E, F are zariski sheaves, where there are open covers  $E_{\alpha} \longrightarrow E, F_{\alpha} \longrightarrow F$  with commutative diagrams

$$E \longrightarrow F$$

$$\uparrow \qquad \uparrow$$

$$E_{\alpha} \stackrel{\cong}{\longrightarrow} F_{\alpha}$$

(i.e. these are isomorphisms locally) then the map is an isomorphism.

With the following diagram, we're done by the lemma:

$$\begin{array}{ccc} X & \longrightarrow & F \\ \uparrow & & \uparrow \\ \downarrow & & \downarrow \\ X_{\alpha} & \stackrel{\cong}{\longrightarrow} & F_{\alpha} \end{array}$$

#### Example 2.5.

For S and E a locally free coherent  $\mathcal{O}_s$  module,

$$\mathbb{P}E(T) = \{f^*E \longrightarrow L \longrightarrow 0\} / \sim$$

is a generalization of projectivization, then S admits a cover  $U_i$  trivializing E.

Then the restriction  $F_i \longrightarrow \mathbb{P}E$  were  $F_i(T)$  is the above set if f factors through  $U_i$  and empty otherwise. On  $U_i$ ,  $E \cong \mathcal{O}_{U_i}^{n_i}$ , so  $F_i$  is representable by  $\mathbb{P}_{U_i}^{n_i-1}$  by the proposition. (Note that this is clearly a sheaf.)

#### Example 2.6.

For E locally free over S of rank n, take r < n and consider the functor  $Gr(k, E)(T) = \{f^*E \longrightarrow Q \longrightarrow 0\} / \sim$  (a Grassmannian) where Q is locally free of rank k.

#### **Exercise**

- a. Show that this is representable
- b. For the Plucker embedding

$$Gr(k, E) \longrightarrow \mathbb{P} \wedge^k E$$
,

a section over T is given by  $f^*E \longrightarrow Q \longrightarrow 0$  corresponding to

$$\wedge^k f^*E \longrightarrow \wedge^k Q \longrightarrow 0,$$

noting that the left-most term is  $f^* \wedge^k E$ .

Show that this is a closed subfunctor. (That it's a functor is clear, that it's closed is not.)

Take  $S = \operatorname{Spec} k$ , then E is a k-vector space V, then sections of the Grassmannian are quotients of  $V \otimes \mathcal{O}$  that are free of rank n.

Take the subfunctor  $G_w \subset Gr(k, V)$  where

$$G_w(T) = \{ \mathcal{O}_T \otimes V \longrightarrow Q \longrightarrow 0 \} \text{ with } Q \cong \mathcal{O}_t \otimes W \subset \mathcal{O}_t \otimes V.$$

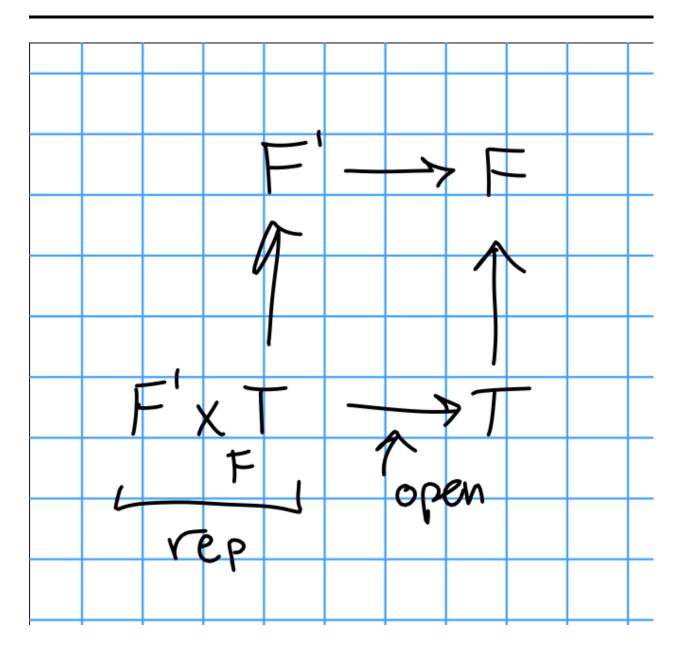
If we have a splitting  $V = W \oplus U$ , then  $G_W = \mathbb{A}(\text{hom}(U, W))$ . If you show it's closed, it follows that it's proper by the exercise at the beginning.

Thursday: Define the Hilbert functor, show it's representable. The Hilbert scheme functor gives e.g. for  $\mathbb{P}^n$  of all flat families of subschemes.

# 3 Thursday January 16th

Subfunctors:

A functor  $F' \subset F : (\operatorname{Sch}/S)^{op} \longrightarrow \operatorname{Set}$  is open iff for all  $T \xrightarrow{\xi} F$  where  $T = h_T$  and  $\xi \in F(T)$ . We can take fiber products:



So we can think of "inclusion in F" as being an open condition: for all T/S and  $\xi \in F(T)$ , there exists an open  $U \subset T$  such that for all covers  $f: T' \longrightarrow T$ , we have  $F(f)(\xi) = f^*(\xi) \in F'(T')$  iff f factors through U.

Suppose  $U \subset T$  in Sch/T, we then have

$$h_{U/T}(T') = \begin{cases} \emptyset & T' \longrightarrow T \text{ doesn't factor} \\ \{\text{pt}\} & \text{otherwise} \end{cases}.$$

which follows because the literal statement is  $h_{U/T}(T') = \text{hom}_T(T', U)$ .

By the definition of the fiber product,  $(F' \times_F T)(T') = \{(a,b) \in F'(T) \times T(T) \mid \xi(b) = \iota(a) \text{ in } F(T)\},\$  where  $F' \xrightarrow{\iota} F$  and  $T \xrightarrow{\xi} F$ .

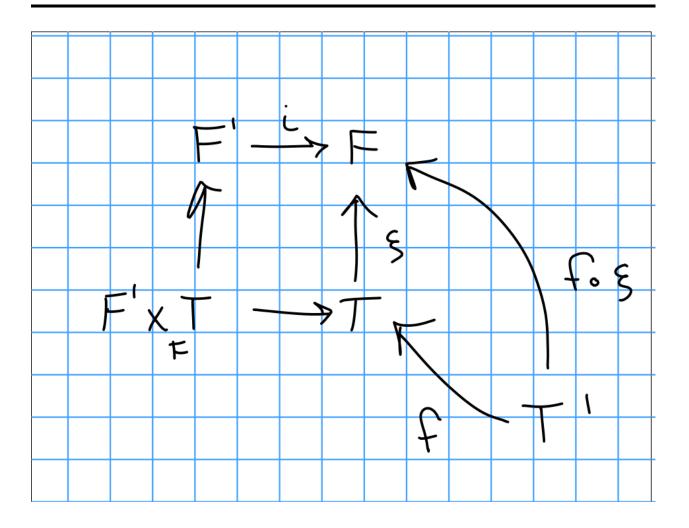
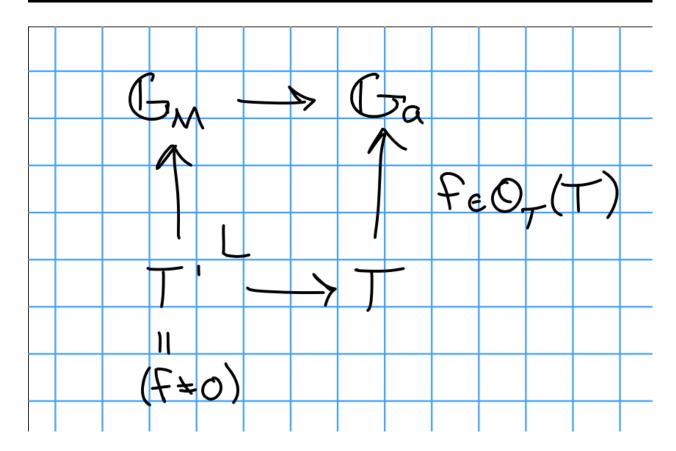


Figure 1: Image

So note that the RHS diagram here is exactly given by pullbacks, since we identify sections of F/T' as sections of F over T/T' (?).

We can thus identify  $(F' \times_F T)(T') = h_{U/S}(T')$ , and so for  $U \subset T$  in Sch/S we have  $h_{U/S} \subset h_{T/S}$  is the functor of maps that factor through U. We just identify  $h_{U/S}(T') = \text{hom}_S(T', U)$  and  $h_{T/S}(T') = \text{hom}_S(T', T)$ .

Example:  $GG_m$ ,  $\mathbb{G}_a$ .  $\mathbb{G}_a$  represents giving a global function,  $\mathbb{G}_m$  represents giving an invertible function.



Where  $\mathcal{O}_T(T)$  are global functions.

#### 3.1 Actual Geometry: Hilbert Schemes

The best moduli space!

Want to parameterize families of subschemes over a fixed object. Fix k a field, X/k a scheme; we'll parameterize subschemes of X.

Definition: The hilbert functor is given by

$$mathrmHilb_{X/S}: (Sch/S)^{op} \longrightarrow Set.$$

which sends T to closed subschemes  $Z \subset X \times_S T \longrightarrow T$  which are flat over T.

Here flatness replaces the Cartier condition.

Recall (flatness): For  $X \xrightarrow{f} Y$  and  $\mathbb{F}$  a coherent sheaf on X, f is flat over Y iff for all  $x \in X$  the stalk  $F_x$  is a flat  $\mathcal{O}_{y,f(x)}$ -module. Note that f is flat if  $\mathcal{O}_x$  is.

Flatness corresponds to varying continuously.

Warning: Unless otherwise stated, assume schemes are Noetherian.

Note that everything works out if we only path with finite covers.



Figure 2: Image

Remark: If X/k is projective, so  $X \subset \mathbb{P}^n_k$ , we have line bundles  $\mathcal{O}_x(1) = \mathcal{O}(1)$ . For any sheaf F over X, there is a hilbert polynomial  $P_F(n) = \chi(F(n)) \in \mathbb{Z}[n]$ . (i.e. we twist by  $\mathcal{O}(1)$  n times.) The cohomology of F isn't change by the pushforward into  $\mathbb{P}_n$  since it's a closed embedding, i.e.  $\chi(X,F) = \chi(\mathbb{P}^n,i_*F) = \sum (-1)^i \dim_k H^i(\mathbb{P}^n,i_*F(n))$ .

First fact: For  $n \gg 0$ ,  $\dim_k H^0 = \dim M_n$ , the *n*th graded piece of M, which is a graded module over the homogeneous coordinate ring whose  $i_*F = \tilde{M}$ .

In general, for L ample of X and F coherent on X, we can define a hilbert polynomial  $P_F(n) = \chi(F \otimes L^n)$ .

This is an invariant of a polarized projective variety, and in particular subschemes. Over irreducible bases, flatness corresponds to this invariant being constant.

Proposition: For  $f: X \longrightarrow S$  projective, i.e. there is a factorization:

If S is reduced, irreducible, locally Noetherian, then f is flat iff  $P_{\mathcal{O}_{x_s}}$  is constant for all  $s \in S$ . (To be more precise, look the base change to  $X_1$ , and the pullback of the fiber?  $\mathcal{O} \mid_{T_s}$ ?)

Note: not using the word "integral" here! S is flat iff the hilbert polynomial over the fibers are constant.

Example: The zero-dimensional subschemes  $Z \in \mathbb{P}_k^n$ , then  $P_Z$  is the length of Z, i.e.  $\dim_k(\mathcal{O}_Z)$ . And  $P_Z(n) = \chi(\mathcal{O}_Z \otimes \mathcal{O}(n)) = \chi(\mathcal{O}_Z) = \dim_k H^0(Z; \mathcal{O}_Z) = \dim_k \mathcal{O}_Z(Z)$ .

For two closed points in  $\mathbb{P}^2$ ,  $P_Z = 2$ .

Consider the affine chart  $\mathbb{A}^2 \subset \mathbb{P}^2$ , which is given by Spec  $k[x,y]/(y,x^2) \cong k[x]/(x^2)$  and  $P_Z = 2$ . I.e. in flat families, it has to record how the tangent directions come together.

Example: Consider the flat family xy = 1 (flat because it's an open embedding) over k[x], here we have points running off to infinity.

**Modified proposition:** A sheaf F is flat iff  $P_{F_S}$  is constant.

Proof of proposition: Assume S = Spec A for A a local Noetherian domain.

**Lemma:** For F a coherent sheaf on X/A is flat, we can take the cohomology via global sections  $H^0(X; F(n))$ . This is an A-module, and is a free A-module for  $n \gg 0$ .

Proof of Lemma: Assumed X was projective, so just take  $X = \mathbb{P}_A^n$  and let F be the pushforward. There is a correspondence sending F to its ring of homogeneous sections constructed by taking the sheaf associated to the graded module  $\sum_{n\gg 0} H^0(\Pi_A^m; F(n))$  This is equal to  $\bigoplus_{n\gg 0} H^0(\mathbb{P}_A^m; F(n))$  and

taking the associated sheaf  $(Y \mapsto \tilde{Y})$ , as per Hartshorne's notation) which is free, and thus F is free.

See tilde construction in Hartshorne, essentially amounts to localizing free tings.

Conversely, take an affine cover  $U_i$  of X. We can compute the cohomology using Čech cohomology, i.e. taking the Čech resolution. We can also assume  $H^i(\mathbb{P}^m; F(n)) = 0$  for  $n \gg 0$ , and the Čech complex vanishes in high enough degree. But then there is an exact sequence

$$0 \longrightarrow H^0(\mathbb{P}^m; F(n)) \longrightarrow \mathcal{C}^0(U; F(n)) \longrightarrow \cdots \longrightarrow C^m(U; F(n)) \longrightarrow 0.$$

Assuming F is flat, and using the fact that flatness is a 2 out of 3 property, the images of these maps are all flat by induction from the right.

Finally, local Noetherian + finitely generated flat implies free.

By the lemma, we want to show  $H^0(\mathbb{P}^m; F(n))$  is free for  $n \gg 0$  iff the hilbert polynomials on the fibers  $P_{F_S}$  are all constant.

Claim 1: It suffices to show that for each point  $s \in \text{Spec } A$ , we have  $H^0(X_s; F_S(n)) = H^0(X; F(n)) \otimes k(S)$  for k(S) the residue field, for  $n \gg 0$ .

Note that  $P_{F_s}$  measures the rank of the LHS.

 $\implies$ : The dimension of RHS is constant, whereas the LHS equals  $P_{F_S}(n)$ .

⇐ : If the dimension of the RHS is constant, so the LHS is free.

For a f.g. module over a local ring, testing if localization at closed point and generic point have the same rank.

For M a finitely generated module over A, find  $0 \longrightarrow A^n \longrightarrow M \longrightarrow Q$  is surjective after tensoring with Frac(A), and tensoring with k(S) for 4s\$a closed point, if dim  $A^n = \dim M$  then Q = 0.

Proof of Claim 1: By localizing, we can assume s is a closed point. Since A is Noetherian, its ideal is f.g. and we have  $A^m \longrightarrow A \longrightarrow k(S) \longrightarrow 0$ . We can tensor with F (viewed as restricting to fiber) to obtain  $F(n)^m \longrightarrow F(n) \longrightarrow F_S(n) \longrightarrow 0$ . Because F is flat, this is still exact.

We can take  $H^*(x, \cdot)$ , and for  $n \gg 0$  only  $H^0$  survives. This is the same as tensoring with  $H^0(x, F(n))$ .

Definition: Given a polynomial  $P \in \mathbb{Z}[n]$  for X/S projective, we define a subfunctor by picking only those with hilbert polynomial p fiberwise as  $mathrmHilb_{X/S}^P \subset mathrmHilb_{X/S}$ . This is given by  $Z \subset X \times_S T$  with  $P_Z = P$ .

Theorem (Grothendieck): If S is Noetherian and X/S projective, then  $\operatorname{mathrmHilb}_{X/S}^P$  is representable by a projective S-scheme.

See cycle spaces in analytic geometry.