

# Title

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# 1 | Saturday, November 28: Introduction to $\infty$ -categories

Dealing with size issues: take a Grothendieck Universe  $\mathcal{U}$ : sets whose subsets are closed under all of the usual set operations (small).

## Definition 1.0.1 ( $\infty$ -Category)

An  $\infty$ -category  $\mathcal{C}$  is a (large) simplicial set  $\mathcal{C}$  such that any diagram of the form

$$\begin{array}{ccc} \Lambda_i^n & \xrightarrow{\quad} & \mathcal{C} \\ \downarrow & \nearrow \exists & \\ \Delta_n & & \end{array}$$

admits the indicated lift, where  $\Lambda_i^n$  is an  $i$ -horn (a simplex missing the  $i$ th face) for  $0 < i < n$ .

**Remark 1.0.2:** This is a specialized notion of a Kan complex, and in particular all  $\infty$ -categories are Kan complexes. All inner horns are fillable, i.e. simplicial sets are *inner* Kan complexes. Different to Kan complexes, which include all  $i$ .

## Definition 1.0.3 (Functors between $\infty$ -categories)

A  $\infty$ -functor between two  $\infty$ -categories is a map between simplicial sets.

## Definition 1.0.4 (Nerve of a category)

Given an ordinary category  $\mathcal{C}$ , define the **nerve** of  $\mathcal{C}$  to be the simplicial set given by

$$N(\mathcal{C})_n := \{\text{Functors } F : [n] \rightarrow \mathcal{C}\}$$

where  $[n]$  is the poset category on  $\{1, 2, \dots, n\}$ . So an  $n$ -simplex is a diagram of objects  $X_0, \dots, X_n \in \text{Ob}(\mathcal{C})$  and a sequence of maps. This defines an  $\infty$ -category, and there is a correspondence

$$\{\text{Functors } F : \mathcal{C} \rightarrow \mathcal{D}\} \iff \{\infty\text{-Functors } \hat{F} : N(\mathcal{C}) \rightarrow N(\mathcal{D})\}.$$

Note that taking the nerve of a category preserves the usual categorical structure, since the objects are the 0-simplices and the morphisms are the 1-simplices.

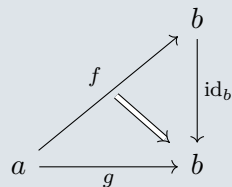
**Remark 1.0.5:** For  $\mathcal{C}$  an  $\infty$ -category, we can define  $\mathcal{C}_0$  to be the “objects” and  $\mathcal{C}_1$  to be the “morphisms”, although we don’t have a good notion of composition yet. There will be boundary map: a 1-simplex has two boundary points, i.e. two objects  $a, b \in \mathcal{C}_0$ , so we can think of this as a map  $f : a \rightarrow b$  where  $a = \partial_1 f, b = \partial_0 f^1$  are the first and second vertices respectively. We’ll also have “degeneracy” maps going up from  $\mathcal{C}_0 \rightarrow \mathcal{C}_1$ , which we should think of as assigning identity morphisms to objects, or conversely that the identity morphism is the degenerate 1-simplex at an

<sup>1</sup>This notation  $\partial_i$  denotes the boundary operator that drops the  $i$ th vertex.

object.

**Definition 1.0.6** (Equivalence of Morphisms)

Given two morphisms  $f, g : a \rightarrow b$  in an  $\infty$ -category, we say  $f \simeq g$  are **equivalent** iff there is a 2-simplex filling in the following diagram:

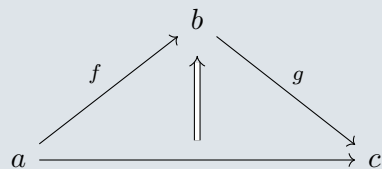


[Link to diagram](#)

**Remark 1.0.7:** This turns out to be an equivalence relation. Note that in an ordinary category, if two morphisms are equivalent then they are already equal.

**Definition 1.0.8** (Composition of morphisms)

For 1-simplices  $f : a \rightarrow b, g : b \rightarrow c$ , a **composition** of  $f$  and  $g$  is a 2-simplex



[Link to diagram](#)